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PKCS #1: RSA Encryption  
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Status of this Memo

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Overview

This document describes a method for encrypting data using the RSA public-key cryptosystem.

1. Scope

This document describes a method for encrypting data using the RSA public-key cryptosystem. Its intended use is in the construction of digital signatures and digital envelopes, as described in PKCS #7:

- o For digital signatures, the content to be signed is first reduced to a message digest with a message-digest algorithm (such as MD5), and then an octet string containing the message digest is encrypted with the RSA private key of the signer of the content. The content and the encrypted message digest are represented together according to the syntax in PKCS #7 to yield a digital signature. This application is compatible with Privacy-Enhanced Mail (PEM) methods.
- o For digital envelopes, the content to be enveloped is first encrypted under a content-encryption key with a content-encryption algorithm (such as DES), and then the content-encryption key is encrypted with the RSA public keys of the recipients of the content. The encrypted content and the encrypted

content-encryption key are represented together according to the syntax in PKCS #7 to yield a digital envelope. This application is also compatible with PEM methods.

The document also describes a syntax for RSA public keys and private keys. The public-key syntax would be used in certificates; the private-key syntax would be used typically in PKCS #8 private-key information. The public-key syntax is identical to that in both X.509 and Privacy-Enhanced Mail. Thus X.509/PEM RSA keys can be used in this document.

The document also defines three signature algorithms for use in signing X.509/PEM certificates and certificate-revocation lists, PKCS #6 extended certificates, and other objects employing digital signatures such as X.401 message tokens.

Details on message-digest and content-encryption algorithms are outside the scope of this document, as are details on sources of the pseudorandom bits required by certain methods in this document.

## 2. References

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### 3. Definitions

For the purposes of this document, the following definitions apply.

**AlgorithmIdentifier:** A type that identifies an algorithm (by object identifier) and associated parameters. This type is defined in X.509.

**ASN.1:** Abstract Syntax Notation One, as defined in X.208.

**BER:** Basic Encoding Rules, as defined in X.209.

**DES:** Data Encryption Standard, as defined in FIPS PUB 46-1.

**MD2:** RSA Data Security, Inc.'s MD2 message-digest algorithm, as defined in RFC 1319.

**MD4:** RSA Data Security, Inc.'s MD4 message-digest algorithm, as defined in RFC 1320.

**MD5:** RSA Data Security, Inc.'s MD5 message-digest algorithm, as defined in RFC 1321.

**modulus:** Integer constructed as the product of two primes.

**PEM:** Internet Privacy-Enhanced Mail, as defined in RFC 1423 and related documents.

**RSA:** The RSA public-key cryptosystem, as defined in [RSA78].

**private key:** Modulus and private exponent.

**public key:** Modulus and public exponent.

### 4. Symbols and abbreviations

Upper-case symbols (e.g., BT) denote octet strings and bit strings (in the case of the signature S); lower-case symbols (e.g., c) denote integers.

ab	hexadecimal octet value	c	exponent
BT	block type	d	private exponent
D	data	e	public exponent
EB	encryption block	k	length of modulus in octets
ED	encrypted data	n	modulus
M	message	p, q	prime factors of modulus
MD	message digest	x	integer encryption block
MD'	comparative message digest	y	integer encrypted data
PS	padding string	mod n	modulo n
S	signature	X    Y	concatenation of X, Y
		X	length in octets of X

## 5. General overview

The next six sections specify key generation, key syntax, the encryption process, the decryption process, signature algorithms, and object identifiers.

Each entity shall generate a pair of keys: a public key and a private key. The encryption process shall be performed with one of the keys and the decryption process shall be performed with the other key. Thus the encryption process can be either a public-key operation or a private-key operation, and so can the decryption process. Both processes transform an octet string to another octet string. The processes are inverses of each other if one process uses an entity's public key and the other process uses the same entity's private key.

The encryption and decryption processes can implement either the classic RSA transformations, or variations with padding.

## 6. Key generation

This section describes RSA key generation.

Each entity shall select a positive integer  $e$  as its public exponent.

Each entity shall privately and randomly select two distinct odd primes  $p$  and  $q$  such that  $(p-1)$  and  $e$  have no common divisors, and  $(q-1)$  and  $e$  have no common divisors.

The public modulus  $n$  shall be the product of the private prime factors  $p$  and  $q$ :

$$n = pq .$$

The private exponent shall be a positive integer  $d$  such that  $de-1$  is divisible by both  $p-1$  and  $q-1$ .

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