



[54] POWER REDUCTION IN A MULTIPROCESSOR DIGITAL SIGNAL PROCESSOR BASED ON PROCESSOR LOAD

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[57] ABSTRACT

Improved operation of multi-processor chips is achieved by dynamically controlling processing load of chips and controlling, significantly greater than on/off granularity, the operating voltages of those chips so as to minimize overall power consumption. A controller in a multi-processor chip allocates tasks to the individual processors to equalize processing load among the chips, then the controller lowers the clock frequency on the chip to as low a level as possible while assuring proper operation, and finally reduces the supply voltage. Further improvement is possible by controlling the supply voltage of individual processing elements within the multi-processor chip, as well as controlling the supply voltage of other elements in the system within which the multi-processor chip operates.

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[52] U.S. Cl. .... 713/300; 713/320; 713/501; 709/100

[58] Field of Search ..... 713/300, 320, 713/321, 322, 323, 340, 501, 600; 709/100, 202; 327/291, 540; 365/227; 712/10

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46 Claims, 2 Drawing Sheets

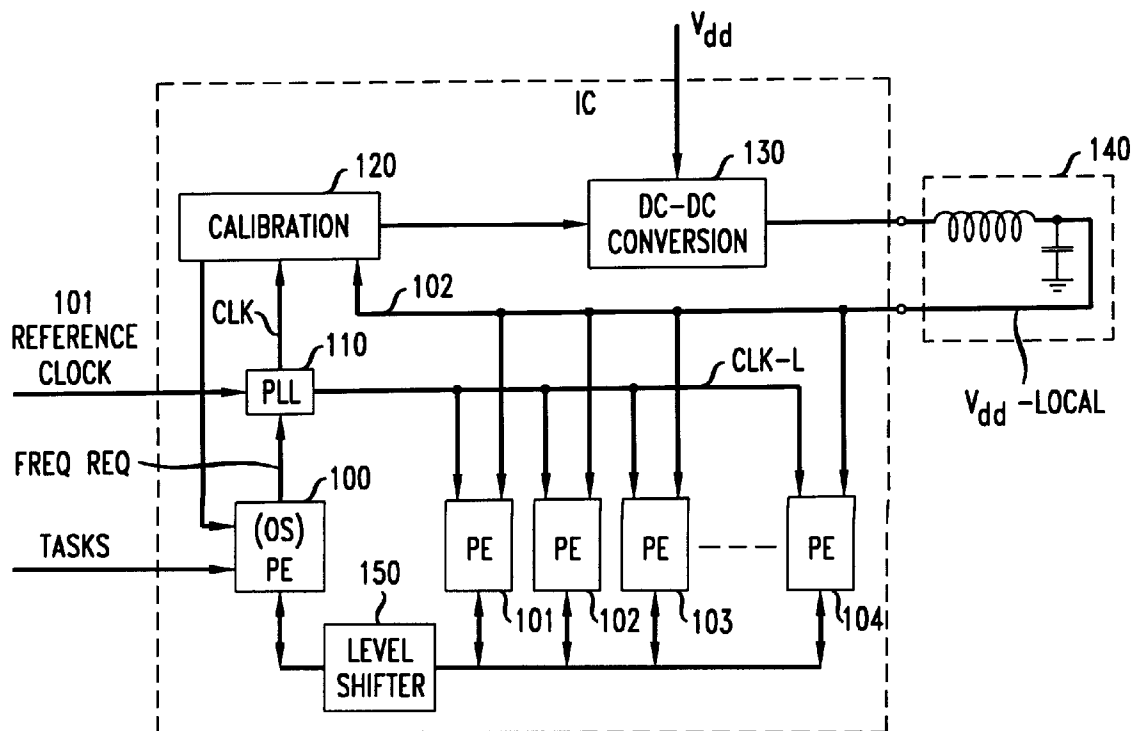


FIG. 1

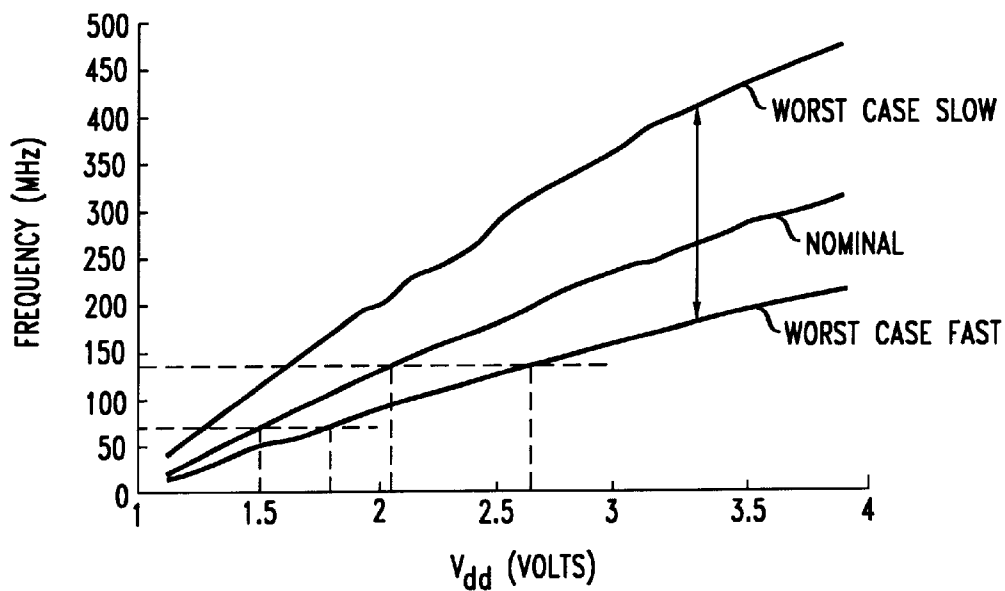


FIG. 2

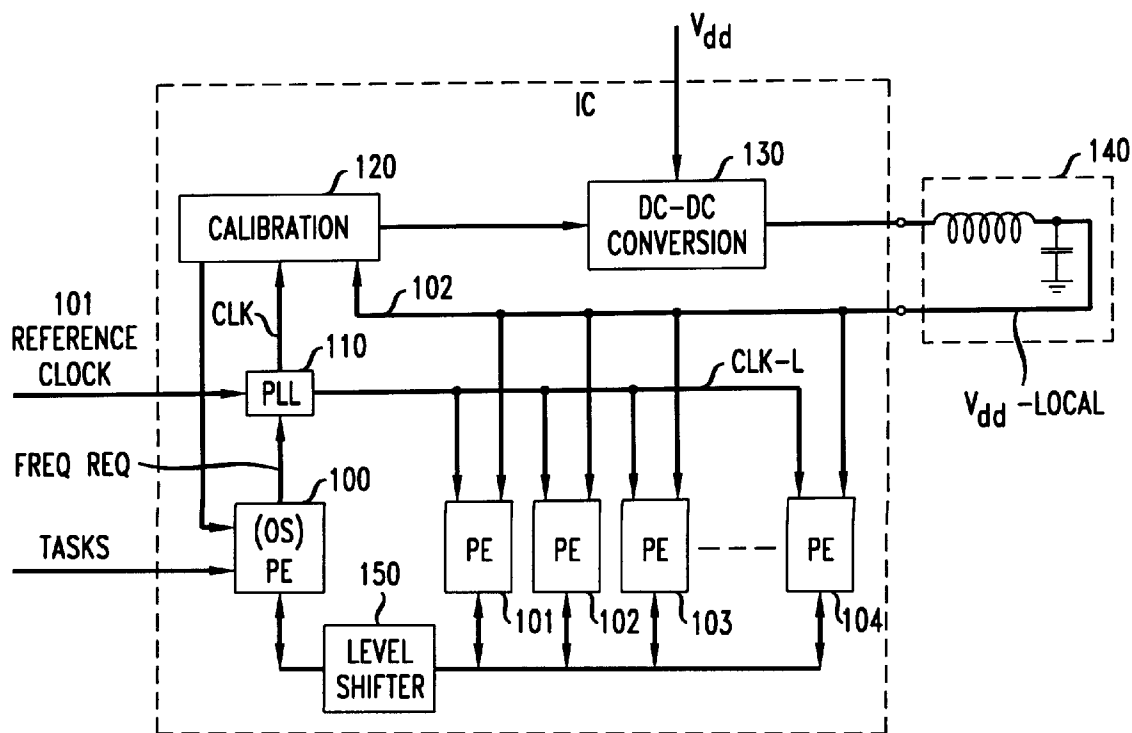


FIG. 3

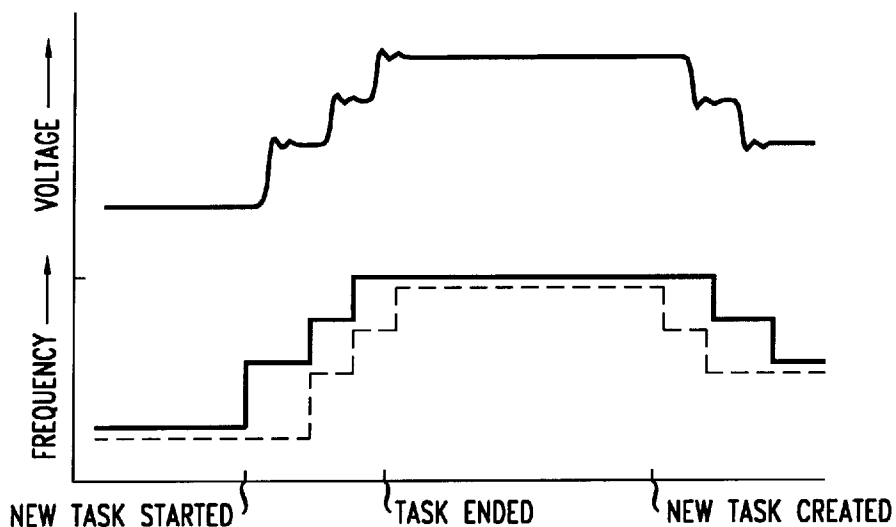
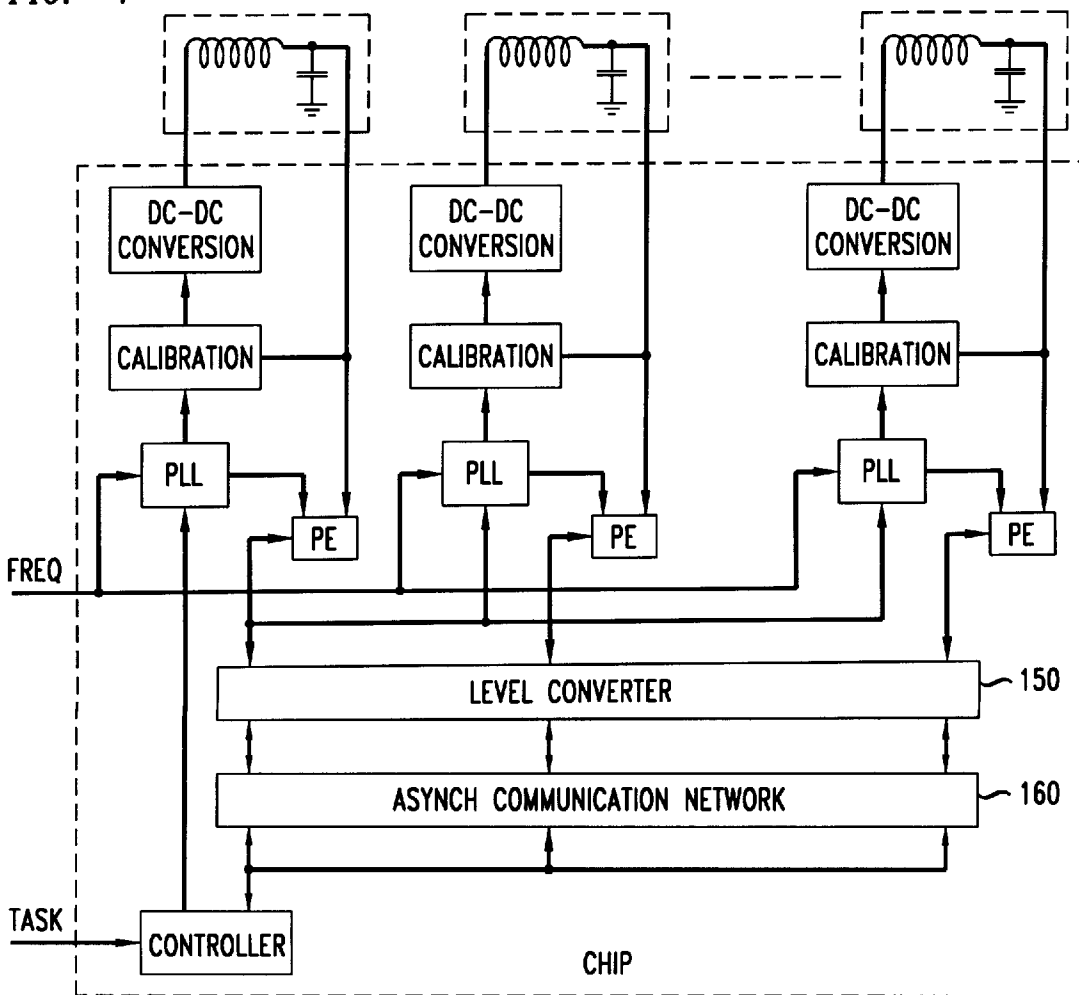


FIG. 4



**POWER REDUCTION IN A  
MULTIPROCESSOR DIGITAL SIGNAL  
PROCESSOR BASED ON PROCESSOR LOAD**

**BACKGROUND**

This invention relates to electronic circuits and, more particularly to power consumption within electronic circuits.

Integrated circuits are designed to meet speed requirements under worst-case operating conditions. In Lucent Technology's 0.35  $\mu\text{m}$  3.3V CMOS technology, the "worst-case-slow" condition is specified for a temperature of 125C. and a chip supply voltage,  $V_{dd}$ , of 2.7V. The worst-case power consumption of the chip is quoted at the maximum supply voltage of 3.6V. The difference in chip performance at the "worst-case slow", nominal, and "worst-case-fast" conditions is shown in FIG. 1, where the frequency of a 25-stage ring oscillator is shown at different supply voltages and process corners. At the nominal operating voltage of 3.3V, the speed difference between "worst case slow" (WCS) and "worst case fast" (WCF) is a factor of 2.2. From the graph it can be seen that if a chip is designed to operate at 140 MHz and at 2.1V supply even when it is "worst-case-slow", a manufactured chip whose characteristics happen to be nominal will continue to operate at 140 MHz even when the chip supply is reduced to 2.1V.

The power consumption of a CMOS circuit increases linearly with operating frequency and quadratically with supply voltage. Therefore, a reduction in supply voltage can significantly reduce power consumption. For example, by reducing the nominal operating voltage from 3.3V to 2.1 V, the nominal power consumption of a 140 MHz chip is reduced by 60% without altering the circuit. This, of course, presumes an ability to identify and measure a chip's variation from nominal characteristics, and an ability to modify the supply voltage based on this measurement.

To achieve variable power supply voltage scaling, a programmable dc—dc converter may be used. Probably, the most efficient approach in use today is the buck converter circuit. These are well known in the art.

Voltage scaling as a function of temperature has been incorporated into the Intel Pentium product family as a technique to achieve high performance at varying operating temperatures and process corners. It is described in U.S. Pat. No. 5,440,520. The approach uses an on-chip temperature sensor and associated processing circuitry which issues a code to the off-chip power supply to provide a particular supply voltage. The process variation information is hard-coded into each device as a final step of manufacturing. This approach has the disadvantage of costly testing of each chip to determine its variance from nominal processing. Several manufacturers make Pentium-compatible dc—dc converter circuits, which are highlighted in "Powering the Big Microprocessors", by B. Travis, EDN, Aug. 15, pp. 31–44, 1997.

Recently, there has been considerable interest in integrating much of the buck controller circuit onto the chip. The only off-chip components are the inductor (typically about 10  $\mu\text{H}$ ) and capacitor (typically about 30  $\mu\text{F}$ ) used in the buck converter. Efficiencies in excess of 80% are typical for a range of voltages and load currents. See, for example, "A High-Efficiency Variable Voltage CMOS Dynamic dc—dc Switching Regulator," by W. Namgoong, M. Yu, and T. Meng, Proceedings ISSCC97 pp. 380–381, February, 1997. Researchers have been also experimenting with on-chip voltage scaling techniques to counter process and temperature variations. See "Variable Supply-Voltage Scheme for

Low Power High-Speed COMS Digital Design," by T. Kuroda et al, CICC97 *Conference Proceedings, and JSSC Issue of CISS97*, May, 1998. The Kuroda et al paper demonstrates that the speed of the circuit can be maintained (or at least the speed degradation can be minimized) by tuning the threshold voltages even as the supply voltage is lowered. The tuning is achieved on-chip by varying the substrate-bias voltage. These techniques are needed to ensure that the leakage current, which increasing as the threshold voltage is reduced, does not become too large.

Thus, it is known that varying supply voltage to a chip can improve performance by eliminating unexpected variability in the supply voltage, and by accounting for process and operating temperature variations.

**SUMMARY OF THE INVENTION**

Improved performance of multi-processor chips is achieved by dynamically controlling the processing load of chips and controlling, which significantly greater than on/off granularity, the operating voltages of those chips so as to minimize overall power consumption. A controller in a multi-processor chip allocates tasks to the individual processors to equalize processing load among the chips, then the controller lowers the clock frequency on the chip to as low a level as possible while assuring proper operation, and finally reduces the supply voltage. Further improvement is possible by controlling the supply voltage of individual processing elements within the multi-processor chip, as well as controlling the supply voltage of other elements in the system within which the multi-processor chip operates.

**BRIEF DESCRIPTION OF THE DRAWINGS**

FIG. 1 illustrates the maximum operating frequency that is achievable with a 0.35  $\mu\text{m}$  technology CMOS chip as a function of supply voltage;

FIG. 2 presents a block diagram of a multi-processor chip with supply voltage control in accordance with the principles disclosed herein;

FIG. 3 shows the relationship between the voltage control clock, Clk, of FIG. 2, the clock applied to the processing elements of FIG. 2, Clk-L, and the supply voltage applied to the processing elements,  $V_{dd}$ -local; and

FIG. 4 depicts the block diagram of a multi-processor chip with supply voltage control that is individual to each of the processing elements.

**DETAILED DESCRIPTION**

FIG. 2 depicts a block diagram of a multi-processor chip. It contains processing elements (PEs) **100**, **101**, **102**, **103**, . . . **104**, and each PE contains a central processing unit (CPU) and a local cache memory (not shown). A real-time operating system resides in PE **100** and allocates tasks to the other PEs from a mix of many digital signal processing applications. The load of the FIG. 2 system is time varying and is dependent on the applications that are being executed at any given time. For example, a set-top-box for a multi-media broadband access system might need to receive an HDTV signal. It could also be transmitting data from a computer, to the Internet, and responding to button requests from a remote control handset. Over time, this dynamic mix of applications places different load requirements on the system.

For a maximally utilized system, all of the available processors ought to be operating at full speed when satisfying the maximum load encountered by the system. At such

a time, the power consumption of the multiprocessor chip is at its maximum level. However, as the load requirements are lowered, the system should, advantageously, reduce its power consumption. It may be noted that, typically, computers spend 99% of their time waiting for a user to press a key. This presents a great opportunity to drastically reduce the average power consumption. The specific approach by which the system “scales back” its performance can greatly impact the realizable power savings.

In the FIG. 2 arrangement, in accordance with the principles disclosed herein, the applications that need to be processed are mapped to the N PEs under control of real time operating system (RTOS) executed on PE 100. If the number of instructions that need to be executed for each task is known and made available to the operating system, a scheduler within the operating system can use this information to determine the best way to allocate the tasks to the available processors in order to balance the computation. The intermediate goal, of course, is to maximize the parallelism and to evenly distribute the load presented to the FIG. 2 system among all of the PE's.

When an application that is running on the FIG. 2 system is subdivided into N concurrent task streams, as suggested above, each of the PEs become lightly loaded. This allows the clock frequency of the PEs to be reduced, and if the task division can be carried out perfectly, then the clock frequency of the FIG. 2 system can be reduced by a factor of N. Reducing the frequency, as indicated above, allows reducing the necessary supply voltage, and reducing the supply voltage reduces the system's power consumption (quadratically). To illustrate, if a given application that is executed on 1 PE requires operating the PE at 140 MHz, it is known from FIG. 1 that the PE can be operated at approximately a 2.7V supply. When the application is divided into two concurrent tasks and assigned to two PEs that are designed to operate at 140 MHz from a 2.7V supply, then the PEs can be operated at 70 MHz and at a supply voltage of 1.8V. This reduction in operating voltage represents a power saving of 55%. Of course, it is unlikely that an application can be perfectly divided into two equal load task streams and, therefore, the 55% power saving is the maximum achievable power saving for two PEs.

It should be understood that in the above example, when two PEs are employed and their operating frequency can be reduced to 70 MHz, the indicated reduction presumes that it is desired to perform the given tasks as if there was a single PE that operates at 140 MHz. That is, the presumption is that there is a certain time when the tasks assigned to the chip must be finished. In fact, there might not be any particular requirement for when the tasks are to be finished. Alternatively, a requirement for when the tasks are to be finished might not be related to the highest operating frequency of the chip.

For example, the above-illustrated chip (where each of the PEs is designed to operate at 140 MHz) might be employed in a system whose basic frequency is related to 160 MHz. In such an arrangement, dividing tasks between the two PEs of the chip and operating each of the PEs at 80 MHz would be preferable because it would be easier to synchronize the chip's input and output functions to the other elements in the system. Thus, in a sense it is the expected completion time for the collection of assigned tasks that is controlling, and the reduction of frequency from the maximum that the chip can support may be controlled by the division of tasks that may be accomplished.

Hence, the operating system of PE 100 needs to ascertain the required completion time, divide the collection of tasks

as evenly as possible (in terms of needed processing time), consider the PE with the tasks that require the most time to carry out, and adjust the clock frequency to insure that the most heavily loaded PE carries out its assigned tasks within the required completion time. Once the frequency is thus determined, a minimum supply voltage can be determined. The supply voltage determination can be made by reference to a plot like the one shown in FIG. 1 or, advantageously, by evaluating the actual performance of the multiprocessor at hand.

As indicated above, the operating system can reduce the supply voltage even further by tracking temperature and process variations. For example, when the chip is nominal in its characteristics, then it can be operated along line 20 of FIG. 1, which calls for only 1.5V supply when operating at 70 MHz.

Returning the discussion to FIG. 2, the programmable-frequency clock is generated using an appropriately multiplied input reference clock (line 101) via a phase lock loop frequency synthesizer circuit 110 which has a high resolution, e.g., can be altered in increments of 5 MHz. Advantageously, two clocks are generated by PLL 110 (requiring two synthesizer circuits), a Clk clock, and a Clk-L which is 1 frequency step lower than Clk when Clk is being increased. For example, in a PLL 110 unit that provides 5 MHz resolution, when Clk is being increased from 75 MHz to 80 MHz, the value of Clk-L is set to 75 MHz.

Clk-L is applied to the PEs, while Clk is applied to calibration circuit 120, which generates a supply voltage command. The supply voltage command is applied to dc—dc converter 130 followed by L-C circuit 140 to cause the combination of converter 130 and L-C circuit 140 to create the supply voltage  $V_{dd}$ -local, which is fed back to calibration circuit 120 via line 102. The  $V_{dd}$ -local supply voltage is also applied to all of the PEs (excluding perhaps the operating system PE 100).

The reason for having the frequency Clk-L lag behind the frequency Clk is that the clock frequency applied to the PEs should not be increased prior to the supply voltage being increased to accommodate the higher frequency. Otherwise, the PEs might fail to perform properly. Circuit 120 observes the level on line 102 to determine whether it corresponds to the voltage necessary to make PEs 100–104 operate properly (described below), and it also waits till the signal on line 102 is stable (following whatever ringing occurs at the output of L-C circuit 140). The signal on line 121 provides information to PE 100 (yes/no) to inform the operating system of when the supply voltage is stable. When the voltage is stable and Clk has reached the required frequency, the operating system sets Clk-L to Clk and then changes the task allocation on the PEs to correspond to that which the PEs were set up to accommodate.

FIG. 3 demonstrates the timing associated with increasing Clk, Clk-L and  $V_{dd}$ -local when a new task is created and the load on the multiprocessor is thus increased, and the timing associated with decreasing Clk, Clk-L and  $V_{dd}$ -local when the load on the multiprocessor is decreased. Specifically, it shows the system operating at 70 MHz from a 1.8V supply when the load is increased in three steps to 140 MHz. When the 2.7V supply is stable, as shown by the supply voltage plot, the new task is enabled for execution. Some time thereafter according to FIG. 3, a task completes, which reduces the load on the multiprocessor. The reduced load permits lowering the clock frequency to 100 MHz and lowering the supply voltage to 2.1V. This, too, is accommodated in steps (two steps, this time), with Clk-L preceding

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