UNITED STATES PATENT AND TRADEMARK OFFICE



Commissioner for Patents United States Patent and Trademark Office P.O. Box 1450 Alexandria, VA 22313-1450 www.uspto.gov

Konda Technologies, Inc 6278 GRAND OAK WAY SAN JOSE CA 95135 In re Application of : KONDA, Venkat : Application No.: 12/601,275 : PCT No.: PCT/US08/64605 : Int. Filing Date: 22 May 2008 : Priority Date: 25 May 2007 : Att. Doc. No.: V-0045US : For: VLSI LAYOUTS OF FULLY : CONNECTED GENERALIZED : NETWORKS : MAILED

AUG 192019

INTERNATIONAL PATENT LEGAL ADM.

DECISION

Applicant's Petition To Accept Late Payment of the Basic National Fee Under 37 CFR 1.137(a), filed in the above-captioned application on 08 August 2019 is **GRANTED**.

Applicant states that the entire delay in filing the required reply from the due date for the required reply until the filing of a grantable petition was unintentional. The proper response was previously filed on 22 December 2009. The petition fee has been furnished. A terminal disclaimer is not required as the application was filed on or after 08 June 1995. Accordingly, all requirements under 37 CFR 1.137(a) have been satisfied.

/Erin P. Thomson/ Erin P. Thomson Attorney Advisor International Patent Legal Administration 571-272-3292 Application Number: 12/601,275

Art Unit: 2819

Communication Dated: August 8, 2019

In The United States Patent And Trademark Office

Application Number: 12/601,275

Application Filed: 11/22/2009

Applicant(s): Venkat Konda

Patent No. 8,269,523

5 Title: VLSI Layouts of Fully Connected Generalized Networks Confirmation No: 6372

Examiner/Art Unit: Vibol Tan / 2819

Priority Date: 5/25/2007

San Jose, 2019 August 8

PETITION TO ACCPET LATE PAYMENT OF THE BASIC NATIONAL FEE UNDER 37 CFR 1.137

Attn: Richard Cole Office of PCT Legal Mail Stop PCT

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15 Commissioner for PatentsP.O. Box 1450Alexandria, Virginia, 22313-1450

Dear Sir/Madam:

20 For the above referenced application, please accept the payment of the basic national fee as being unintentionally delayed. The entire delay in filing the basic national fee from the due date for the fee until the filing of a grantable petition under 37 CFR 1.137 was unintentional. A check of \$1000, the small entity fee, for this petition is enclosed.

- 25 Very respectfully, /Venkat Konda/
 Venkat Konda
 6278 Grand Oak Way
 San Jose, CA 95135
- 30 Phone: 408-472-3273

Electronic Patent Application Fee Transmittal						
Application Number:	12	12601275				
Filing Date:	31-	31-May-2010				
Title of Invention:	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS					
First Named Inventor/Applicant Name:	Venkat Konda					
Filer:	Venkat Konda					
Attorney Docket Number:	V-(0045US				
Filed as Small Entity						
Filing Fees for U.S. National Stage under 35 USC 371						
Description		Fee Code	Quantity	Amount	Sub-Total in USD(\$)	
Basic Filing:						
Pages:						
Claims:						
Miscellaneous-Filing:						
Petition:						
PET. REVIVE ABANDON APP, DELAY PYMT-RESP		2453	1	1000	1000	
Patent-Appeals-and-Interference:						
Post-Allowance-and-Post-Issuance:						

Description	Fee Code	Quantity	Amount	Sub-Total in USD(\$)
Extension-of-Time:				
Miscellaneous:				
	Tot	al in USD	(\$)	1000

Electronic Ac	knowledgement Receipt
EFS ID:	36825423
Application Number:	12601275
International Application Number:	
Confirmation Number:	6372
Title of Invention:	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS
First Named Inventor/Applicant Name:	Venkat Konda
Customer Number:	38139
Filer:	Venkat Konda
Filer Authorized By:	
Attorney Docket Number:	V-0045US
Receipt Date:	08-AUG-2019
Filing Date:	31-MAY-2010
Time Stamp:	16:32:33
Application Type:	U.S. National Stage under 35 USC 371

Payment information:

Submitted with Payment	yes		
Payment Type	CARD		
Payment was successfully received in RAM	\$1000		
RAM confirmation Number	E201988G34075103		
Deposit Account			
Authorized User			
The Director of the USPTO is hereby authorized to charge indicated fees and credit any overnayment as follows:			

The Director of the USPTO is hereby authorized to charge indicated fees and credit any overpayment as follows:

File Listing:

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Document Number	Document Description	File Name	File Size(Bytes)/ Message Digest	Multi Part /.zip	Pages (if appl.)
1	Petition for review by the PCT legal office	PetUnintDelay-V0045US.pdf	75026 dc16c6a0f247024d21c0f54128af4f7c5e1ee 3a0	no	1
Warnings:					
Information:					
			30076		
2	Fee Worksheet (SB06)	fee-info.pdf	f6cb27709e04fa0fb586944ee8732dc11110 b06f	no	2
Warnings:					
Information:					
		Total Files Size (in bytes)	1	05102	
characterized Post Card, as <u>New Applicat</u> If a new appli 1.53(b)-(d) an Acknowledge <u>National Stac</u> If a timely sul U.S.C. 371 an national stag <u>New Internat</u> If a new inter an internatio and of the Int	ledgement Receipt evidences receip d by the applicant, and including pay described in MPEP 503. tions Under 35 U.S.C. 111 ication is being filed and the applica of MPEP 506), a Filing Receipt (37 CF ement Receipt will establish the filin ge of an International Application un bmission to enter the national stage d other applicable requirements a F ie submission under 35 U.S.C. 371 with ional Application Filed with the USP national application is being filed an nal filing date (see PCT Article 11 an ternational Filing Date (Form PCT/Re urity, and the date shown on this Ack on.	ge counts, where applicable. tion includes the necessary of R 1.54) will be issued in due g date of the application. <u>Inder 35 U.S.C. 371</u> of an international applicati orm PCT/DO/EO/903 indicati ill be issued in addition to the <u>PTO as a Receiving Office</u> and the international applicat d MPEP 1810), a Notification D/105) will be issued in due c	It serves as evidence components for a filir course and the date s ion is compliant with ing acceptance of the e Filing Receipt, in du ion includes the nece of the International ourse, subject to pres	of receipt s ing date (see shown on th the condition application e course. essary comp Application scriptions co	imilar to a 37 CFR is ons of 35 as a onents for Number oncerning

Application Number: 12/601,275

Art Unit: 2819

Communication Dated: August 8, 2019

In The United States Patent And Trademark Office

Application Number: 12/601,275

Application Filed: 11/22/2009

Applicant(s): Venkat Konda

Patent No. 8,269,523

5 Title: VLSI Layouts of Fully Connected Generalized Networks Confirmation No: 6372

Examiner/Art Unit: Vibol Tan / 2819

Priority Date: 5/25/2007

San Jose, 2019 August 8

PETITION TO ACCPET LATE PAYMENT OF THE BASIC NATIONAL FEE UNDER 37 CFR 1.137

Attn: Richard Cole Office of PCT Legal Mail Stop PCT

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15 Commissioner for PatentsP.O. Box 1450Alexandria, Virginia, 22313-1450

Dear Sir/Madam:

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- 25 Very respectfully, /Venkat Konda/
 Venkat Konda
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 San Jose, CA 95135
- 30 Phone: 408-472-3273

Electronic Patent Application Fee Transmittal						
Application Number:	12	12601275				
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Title of Invention:	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS					
First Named Inventor/Applicant Name:	Venkat Konda					
Filer:	Venkat Konda					
Attorney Docket Number:	V-(0045US				
Filed as Small Entity						
Filing Fees for U.S. National Stage under 35 USC 371						
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PET. REVIVE ABANDON APP, DELAY PYMT-RESP		2453	1	1000	1000	
Patent-Appeals-and-Interference:						
Post-Allowance-and-Post-Issuance:						

Description	Fee Code	Quantity	Amount	Sub-Total in USD(\$)
Extension-of-Time:				
Miscellaneous:				
	Tot	al in USD	(\$)	1000

Electronic Ac	knowledgement Receipt
EFS ID:	36824269
Application Number:	12601275
International Application Number:	
Confirmation Number:	6372
Title of Invention:	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS
First Named Inventor/Applicant Name:	Venkat Konda
Customer Number:	38139
Filer:	Venkat Konda
Filer Authorized By:	
Attorney Docket Number:	V-0045US
Receipt Date:	08-AUG-2019
Filing Date:	31-MAY-2010
Time Stamp:	15:45:01
Application Type:	U.S. National Stage under 35 USC 371

Payment information:

Submitted with Payment			no				
File Listing	g:						
Document Number	Document Description		File Name	File Size(Bytes)/ Message Digest	Multi Part /.zip	Pages (if appl.)	
		PetUnintDelay-V0045US.pdf		75026			
Petition for review by the PCT legal office	dc16c6a0f247024d21c0f54128af4f7c5e1ee 3a0			no	1		
Warnings:				ł	ľ		

30076	
2 Fee Worksheet (SB06) fee-info.pdf no 2	2
Warnings:	ngs:
Information:	nation:
Total Files Size (in bytes): 105102	
characterized by the applicant, and including page counts, where applicable. It serves as evidence of receipt similar t Post Card, as described in MPEP 503. <u>New Applications Under 35 U.S.C. 111</u> If a new application is being filed and the application includes the necessary components for a filing date (see 37 CFR 1.53(b)-(d) and MPEP 506), a Filing Receipt (37 CFR 1.54) will be issued in due course and the date shown on this Acknowledgement Receipt will establish the filing date of the application. <u>National Stage of an International Application under 35 U.S.C. 371</u> If a timely submission to enter the national stage of an international application is compliant with the conditions of 3 U.S.C. 371 and other applicable requirements a Form PCT/DO/EO/903 indicating acceptance of the application as a national stage submission under 35 U.S.C. 371 will be issued in addition to the Filing Receipt, in due course. <u>New International Application Filed with the USPTO as a Receiving Office</u> If a new international application is being filed and the international application includes the necessary components an international filing date (see PCT Article 11 and MPEP 1810), a Notification of the International Application Numbu and of the International Filing Date (Form PCT/RO/105) will be issued in due course, subject to prescriptions concernin national security, and the date shown on this Acknowledgement Receipt will establish the international filing date of the shown on this Acknowledgement Receipt will establish the international filing date of the shown on this Acknowledgement Receipt will establish the international filing date of the shown on this Acknowledgement Receipt will establish the international filing date of the shown on this Acknowledgement Receipt will establish the international filing date of the shown on this Acknowledgement Receipt will establish the international filing date of the shown on this Acknowledgement Receipt will establish the international filing date of	ard, as d pplication w applic b)-(d) and bwledger nal Stage nely sub- al stage nternation ernation f the Inte

505428057 04/15/2019

PATENT ASSIGNMENT COVER SHEET

Electronic Version v1.1 Stylesheet Version v1.2 EPAS ID: PAT5474857

SUBMISSION TYPE:		NE	W ASSIGNMENT		
NATURE OF CONVEYANCE:			SIGNMENT		
CONVEYING PARTY D	ΑΤΑ	I			
		Nai	me		Execution Date
KONDA TECHNOLOGIE	ES INC.				04/15/2019
RECEIVING PARTY DA					
Name:			A		
Street Address:		ARAND O			
City:	SAN JO				
State/Country:	CALIFO				
Postal Code:	95135				
PROPERTY NUMBERS	Total: 2	2			
Property Type			Number		
Patent Number: 8269		8269523			
Application Number: 16202067					
Fax Number:		(408)238	-2478		
		. ,	ail address first; if that	is unsud	cessful, it will be sent
using a fax number, if	providea	d; if that i	is unsuccessful, it will l		
Phone:		4084723			
Email:			kondatech.com		
Correspondent Name: KONDA TECHNOLOGIES INC.					
•					
Address Line 1: Address Line 4:			SE, CALIFORNIA 95135		
Address Line 1: Address Line 4:		SAN JOS			
Address Line 1: Address Line 4: NAME OF SUBMITTER:		SAN JOS	ENKAT KONDA		
Address Line 1: Address Line 4: NAME OF SUBMITTER: SIGNATURE:		SAN JOS VE /ve	ENKAT KONDA enkat Konda/		
Address Line 1: Address Line 4: NAME OF SUBMITTER: SIGNATURE:		SAN JOS VE /ve 04	NKAT KONDA enkat Konda/ /15/2019		\sim
Address Line 1: Address Line 4: NAME OF SUBMITTER: SIGNATURE: DATE SIGNED:		SAN JOS VE /ve 04	NKAT KONDA enkat Konda/ /15/2019	n Oath/De	eclaration (37 CFR 1.63).
Address Line 1: Address Line 4: NAME OF SUBMITTER: SIGNATURE:		SAN JOS VE /ve 04	NKAT KONDA enkat Konda/ /15/2019	n Oath/De	eclaration (37 CFR 1.63).

ASSIGNMENT OF PATENTS

WHEREAS, Konda Technologies, Inc., a California corporation, located at 6278 Grand Oak Way, San Jose, CA 95135, hereinafter referred to as "Assignor," is the record owner of the patents described as

- Title: VLSI Layouts of Fully Connected Generalized Networks US Patent No.: 8,269,523 Patent Issue Date: September 18, 2012;
- Title: VLSI Layouts of Fully Connected Generalized Networks US Patent Application No.: 16/202,067 Patent Filing Date: November 27, 2018;

WHEREAS, Konda, Venkat of San Jose, California, having a residence at 6278 Grand Oak Way, San Jose, CA, hereinafter referred to as "ASSIGNEE," wishes to acquire all right, title and interest to and under the patents described above owned by Assignor and in and to any and all improvements to said applications and in any Letters Patent and Registrations which may be granted on the same in the United States or any country throughout the world;

For good and valuable consideration, receipt of which is hereby acknowledged by Assignor, effective 15th day of April, 2019, has assigned, and by these presents does assign to Assignee all right, title and interest for the United States and all foreign countries, in and to any and all improvements for the applications described above and in and to said applications and to all utility divisional continuing, substitute, renewal, reissue, and all other patent applications which have been or shall be filed in the United States and all foreign counterparts (including patent, utility model and industrial designs), and in and to any Letters Patent and Registrations which may hereafter be granted on the same in the United States and all countries throughout the world, and to claim the priority from the application as provided by the Paris Convention. The right, title and interest is to be held and enjoyed by Assignee and Assignee's successors and assigns as fully and exclusively as it would have been held and enjoyed by Assignor had this Assignment not been made, for the full term of any Letters Patent and Registrations which may be granted thereon, or of any division, renewal, continuation in whole or in part, substitution, conversion, reissue, prolongation or extension thereof.

Assignor further agrees that they will, without charge to Assignee, but at Assignee's expense, (a) cooperate with Assignee in the prosecution of U.S. Patent applications and foreign counterparts on the applications and any improvements, (b) execute, verify, acknowledge and deliver all such further papers, including patent applications and instruments of transfer, and (c) perform such other acts as Assignee lawfully may request to obtain or maintain Letters Patent and Registrations for the applications and improvements in any and all countries, and to vest title thereto in Assignee, or Assignee's successors and assigns. Page 2

IN TESTIMONY WHEREOF, Assignor has signed on the date indicated.

Date: 4/15/2019 By: Vanka

- 1(r Venkat Konda (Founder/CEO)

Venkat Konda (Founder/CEC Konda Technologies Inc. 6278 Grand Oak Way San Jose, CA 95135



UNITED STATES PATENT AND TRADEMARK OFFICE

UNITED STATES DEPARTMENT OF COMMERCE United States Patent and Trademark Office
Address: COMMISSIONER FOR PATENTS
P.O. Box 1450
Alexandria, Virginia 22313-1450
www.uspto.gov

APPLICATION NO.	ISSUE DATE	PATENT NO.	ATTORNEY DOCKET NO.	CONFIRMATION NO.
12/601,275	09/18/2012	8269523	V-0045US	6372
38139 7	590 08/29/2012			

Konda Technologies, Inc 6278 GRAND OAK WAY SAN JOSE, CA 95135

ISSUE NOTIFICATION

The projected patent number and issue date are specified above.

Determination of Patent Term Adjustment under 35 U.S.C. 154 (b)

(application filed on or after May 29, 2000)

The Patent Term Adjustment is 183 day(s). Any patent to issue from the above-identified application will include an indication of the adjustment on the front page.

If a Continued Prosecution Application (CPA) was filed in the above-identified application, the filing date that determines Patent Term Adjustment is the filing date of the most recent CPA.

Applicant will be able to obtain more detailed information by accessing the Patent Application Information Retrieval (PAIR) WEB site (http://pair.uspto.gov).

Any questions regarding the Patent Term Extension or Adjustment determination should be directed to the Office of Patent Legal Administration at (571)-272-7702. Questions relating to issue and publication fee payments should be directed to the Application Assistance Unit (AAU) of the Office of Data Management (ODM) at (571)-272-4200.

APPLICANT(s) (Please see PAIR WEB site http://pair.uspto.gov for additional applicants):

Venkat Konda, San Jose, CA;

The United States represents the largest, most dynamic marketplace in the world and is an unparalleled location for business investment, innovation, and commercialization of new technologies. The USA offers tremendous resources and advantages for those who invest and manufacture goods here. Through SelectUSA, our nation works to encourage and facilitate business investment. To learn more about why the USA is the best country in the world to develop technology, manufacture products, and grow your business, visit <u>SelectUSA.gov</u>.

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Change(s) applied to document, /D.H.P./ 6/29/2012 This application is related to and incorporates by reference in its entirety the US 12/601273 Patent Application Docket No. V-0038US entitled "FULLY CONNECTED GENERALIZED BUTTERFLY FAT TREE NETWORKS" by Venkat Konda assigned to the same assignee as the current application filed concurrently, the PCT Application Serial No. PCT/US08/64603 entitled "FULLY CONNECTED GENERALIZED BUTTERFLY FAT TREE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 22, 2008, the U.S. Provisional Patent Application Serial No. 60/940, 387 entitled "FULLY CONNECTED GENERALIZED BUTTERFLY FAT TREE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007, and the U.S. Provisional Patent Application Serial No. 60/940, 390 entitled "FULLY CONNECTED GENERALIZED BUTTERFLY FAT TREE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007, and the U.S. Provisional Patent Application Serial No. 60/940, 390 entitled "FULLY CONNECTED GENERALIZED MULTI-LINK BUTTERFLY FAT TREE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007.

This application is related to and incorporates by reference in its entirety the US 12/601274

- 15 Patent Application Doeket No. V-0039US entitled "FULLY CONNECTED GENERALIZED MULTI-LINK MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application filed concurrently, the PCT Application Serial No. PCT/US08/64604 entitled "FULLY CONNECTED GENERALIZED MULTI-LINK MULTI-STAGE NETWORKS" by Venkat Konda
- 20 assigned to the same assignee as the current application, filed May 22, 2008, the U.S. Provisional Patent Application Serial No. 60/940, 389 entitled "FULLY CONNECTED GENERALIZED REARRANGEABLY NONBLOCKING MULTI-LINK MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007, the U.S. Provisional Patent Application Serial No.
- 25 60/940, 391 entitled "FULLY CONNECTED GENERALIZED FOLDED MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007 and the U.S. Provisional Patent Application Serial No. 60/940, 392 entitled "FULLY CONNECTED GENERALIZED STRICTLY NONBLOCKING MULTI-LINK MULTI-STAGE NETWORKS" by Venkat Konda
- 30 assigned to the same assignee as the current application, filed May 25, 2007.

PART B - FEE(S) TRANSMITTAL

Complete and send this form, together with applicable fee(s), to: Mail Mail Stop ISSUE FEE

Commissioner for Patents P.O. Box 1450 Alexandria, Virginia 22313-1450 or Fax (571) 273-2885

INSTRUCTIONS: This form should be used for transmitting the ISSUE FEE and PUBLICATION FEE (if required). Blocks 1 through 5 should be completed where appropriate. All further correspondence including the Patent, advance orders and notification of maintenance fees will be mailed to the current correspondence address as indicated unless corrected below or directed otherwise in Block 1, by (a) specifying a new correspondence address; and/or (b) indicating a separate "FEE ADDRESS" for maintenance fees will be mailed to the current correspondence of the second maintenance fee notifications

CURRENT CORRESPONDENCE ADDRESS (Note: Use Block 1 for any change of address)

Note: A certificate of mailing can only be used for domestic mailings of the Fee(s) Transmittal. This certificate cannot be used for any other accompanying papers. Each additional paper, such as an assignment or formal drawing, must have its own certificate of mailing or transmission.

Certificate of Mailing or Transmission

I hereby certify that this Fee(s) Transmittal is being deposited with the United States Postal Service with sufficient postage for first class mail m an envelope addressed to the Mail Stop ISSUE FEE address above, or being facsimile transmitted to the USPTO (571) 273-2885, on the date indicated below.

	VEC	¢970	¢200	¢1170	8/17/0010
APPLN. TYPE	SMALL ENTITY	ISSUE FEE	PUBLICATION FEE	TOTAL FEE(S) DUE	DATE DUE
ILE OF INVENTION:					
12/601,275	5/31/2010	Venkat Konda		V-0045US	6372
APPLICATION NO.	FILING DATE	FIRST NAM	IED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
					(Date)
					(Signature)
					(Depositor's name)

	nonprovisional	YES	\$870		\$300	\$117	0	8/17/2012	
	EXAMINER		ART UN	IT	CLASS-SUBCLASS				
	1. Change of correspondence addre	ess or indication of "Fe	ee Address" (37	2. For prir	ting on the patent front page, list				
ę	CFR 1.363). Change of correspondence a	ddress (or Change of	Correspondence	(1) the nation of agents (1)	mes of up to 3 registered patent DR, alternatively,	attorneys 1			
	Address form PTO/SB/122) atta	ached.		(2) the nor	ne of a single firm (having as a r	nembera 2			

Address form P1O/SB/122) attached. "Fee Address" indication (or "Fee Address" Indication form PTO/SB/47; Rev 03-02 or more recent) attached. Use of a Customer Number is required.	(2) the name of a single firm (having as a member a registered attorney or agent) and the names of up to 2 registered patent attorneys or agents. If no name is listed, no name will be printed.	2 3	
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3. ASSIGNEE NAME AND RESIDENCE DATA TO BE PRINTED ON THE PATENT (print or type)

PLEASE NOTE: Unless an assignee is identified below, no assignee data will appear on the patent. If an assignee is identified below, the document has been filed for recordation as set forth in 37 CFR 3.11. Completion of this form is NOT a substitute for filing an assignment.

(A) NAME OF ASSIGNEE	(B) RESIDENCE: (CITY and STATE OR COUNTRY)
Konda Technologies Inc.	San Jose, CA

Please check the appropriate assignee category or ca	ategories (will not be printed on the patent) : Individ	ual Corporation or other private group entity Government
4a. The following fee(s) are enclosed:	4b. Payment of Fee(s):	
🗹 Issue Fee	A check in the amount of the t	fee(s) is enclosed.

Publication Fee (No small entity discount permitted)	Payment by credit card. Form PTO-2038 is attached.
Advance Order - # of Copies _	The Director is hereby authorized to charge the required fee(s), or credit any overpayment, to Deposit Account Number
5. Change in Entity Status (from status indicated above)	

a. Applicant claims SMALL ENTITY status. See 37 CFR 1.27. b. Applicant is no longer claiming SMALL ENTITY status. See 37 CFR 1.27(g)(2).

The Director of the USPTO is requested to apply the Issue Fee and Publication Fee (if any) or to re-apply any previously paid issue fee to the application identified above. NOTE: The Issue Fee and Publication Fee (if required) will not be accepted from anyone other than the applicant; a registered attorney or agent; or the assignee or other party in interest as shown by the records of the United States Patent and Trademark Office.

Authorized Signature /	Venkat Konda/	Date	8/15/2012
Typed or printed name	Venkat Konda	Regist	tration No.

This collection of information is required by 37 CFR 1.311. The information is required to obtain or retain a benefit by the public which is to file (and by the USPTO to process) an application. Confidentiality is governed by 35 U.S.C. 122 and 37 CFR 1.14. This collection is estimated to take 12 minutes to complete, including gathering, preparing, and submitting the completed application form to the USPTO. Time will vary depending upon the individual case. Any comments on the amount of time you require to complete this form and/or suggestions for reducing this burden, should be sent to the Chief Information Officer, U.S. Patent and Trademark Office, U.S. Department of Commerce, P.O. Box 1450, Alexandria, Virginia 22313-1450. DO NOT SEND FEES OR COMPLETED FORMS TO THIS ADDRESS. SEND TO: Commissioner for Patents, P.O. Box 1450, Alexandria, Virginia 22313-1450.

Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number.

Mankat Kanda/

Privacy Act Statement

The Privacy Act of 1974 (P.L. 93-579) requires that you be given certain information in connection with your submission of the attached form related to a patent application or patent. Accordingly, pursuant to the requirements of the Act, please be advised that: (1) the general authority for the collection of this information is 35 U.S.C. 2(b)(2); (2) furnishing of the information solicited is voluntary; and (3) the principal purpose for which the information is used by the U.S. Patent and Trademark Office is to process and/or examine your submission related to a patent application or patent. If you do not furnish the requested information, the U.S. Patent and Trademark Office may not be able to process and/or examine your submission, which may result in termination of proceedings or abandonment of the application or expiration of the patent.

The information provided by you in this form will be subject to the following routine uses:

- 1. The information on this form will be treated confidentially to the extent allowed under the Freedom of Information Act (5 U.S.C. 552) and the Privacy Act (5 U.S.C 552a). Records from this system of records may be disclosed to the Department of Justice to determine whether disclosure of these records is required by the Freedom of Information Act.
- 2. A record from this system of records may be disclosed, as a routine use, in the course of presenting evidence to a court, magistrate, or administrative tribunal, including disclosures to opposing counsel in the course of settlement negotiations.
- 3. A record in this system of records may be disclosed, as a routine use, to a Member of Congress submitting a request involving an individual, to whom the record pertains, when the individual has requested assistance from the Member with respect to the subject matter of the record.
- 4. A record in this system of records may be disclosed, as a routine use, to a contractor of the Agency having need for the information in order to perform a contract. Recipients of information shall be required to comply with the requirements of the Privacy Act of 1974, as amended, pursuant to 5 U.S.C. 552a(m).
- 5. A record related to an International Application filed under the Patent Cooperation Treaty in this system of records may be disclosed, as a routine use, to the International Bureau of the World Intellectual Property Organization, pursuant to the Patent Cooperation Treaty.
- 6. A record in this system of records may be disclosed, as a routine use, to another federal agency for purposes of National Security review (35 U.S.C. 181) and for review pursuant to the Atomic Energy Act (42 U.S.C. 218(c)).
- 7. A record from this system of records may be disclosed, as a routine use, to the Administrator, General Services, or his/her designee, during an inspection of records conducted by GSA as part of that agency's responsibility to recommend improvements in records management practices and programs, under authority of 44 U.S.C. 2904 and 2906. Such disclosure shall be made in accordance with the GSA regulations governing inspection of records for this purpose, and any other relevant (i.e., GSA or Commerce) directive. Such disclosure shall not be used to make determinations about individuals.
- 8. A record from this system of records may be disclosed, as a routine use, to the public after either publication of the application pursuant to 35 U.S.C. 122(b) or issuance of a patent pursuant to 35 U.S.C. 151. Further, a record may be disclosed, subject to the limitations of 37 CFR 1.14, as a routine use, to the public if the record was filed in an application which became abandoned or in which the proceedings were terminated and which application is referenced by either a published application, an application open to public inspection or an issued patent.
- 9. A record from this system of records may be disclosed, as a routine use, to a Federal, State, or local law enforcement agency, if the USPTO becomes aware of a violation or potential violation of law or regulation.

Electronic Patent Application Fee Transmittal						
Application Number:	126	12601275				
Filing Date:	31-I	May-2010				
Title of Invention:	VLS	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS				
First Named Inventor/Applicant Name:	Ven	kat Konda				
Filer:	Ven	Venkar Konda				
Attorney Docket Number:	V-00	045US				
Filed as Small Entity						
U.S. National Stage under 35 USC 371 Filing F	Fees	;				
Description		Fee Code	Quantity	Amount	Sub-Total in USD(\$)	
Basic Filing:						
Pages:						
Claims:						
Miscellaneous-Filing:						
Publ. Fee- early, voluntary, or normal		1504	1	300	300	
Petition:						
Patent-Appeals-and-Interference:						
Post-Allowance-and-Post-Issuance:						
Utility Appl issue fee		2501	1	870	870	

Description	Fee Code	Quantity	Amount	Sub-Total in USD(\$)
Extension-of-Time:				
Miscellaneous:				
	Tot	al in USD	(\$)	1170

Electronic A	Electronic Acknowledgement Receipt					
EFS ID:	13505897					
Application Number:	12601275					
International Application Number:						
Confirmation Number:	6372					
Title of Invention:	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS					
First Named Inventor/Applicant Name:	Venkat Konda					
Customer Number:	38139					
Filer:	Venkar Konda					
Filer Authorized By:						
Attorney Docket Number:	V-0045US					
Receipt Date:	15-AUG-2012					
Filing Date:	31-MAY-2010					
Time Stamp:	16:44:45					
Application Type:	U.S. National Stage under 35 USC 371					

Payment information:

Document Number Page 21 of 3	Document Description	File Name	File Size(Bytes)/ Message Digest	Multi Part /.zip	Pages (if appl.)			
File Listing	:							
Authorized Use	r							
Deposit Accour	nt							
RAM confirmati	on Number	3603						
Payment was su	uccessfully received in RAM	\$1170						
Payment Type		Credit Card						
Submitted with	Payment	yes	yes					

	/ledgement Receipt evidences receip	Total Files Size (in bytes)		55462	
Information	:		1		
Warnings:					
2	Fee Worksheet (SB06)	fee-info.pdf	31848 0d3037333ad300df6fcb48c6df4c5fc1941fa ee5	no	2
Information	:	I	i		
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1	lssue Fee Payment (PTO-85B)	ptol85b-V0045US.pdf	223614	no	2

New Applications Under 35 U.S.C. 111

If a new application is being filed and the application includes the necessary components for a filing date (see 37 CFR 1.53(b)-(d) and MPEP 506), a Filing Receipt (37 CFR 1.54) will be issued in due course and the date shown on this Acknowledgement Receipt will establish the filing date of the application.

National Stage of an International Application under 35 U.S.C. 371

If a timely submission to enter the national stage of an international application is compliant with the conditions of 35 U.S.C. 371 and other applicable requirements a Form PCT/DO/EO/903 indicating acceptance of the application as a national stage submission under 35 U.S.C. 371 will be issued in addition to the Filing Receipt, in due course.

New International Application Filed with the USPTO as a Receiving Office

If a new international application is being filed and the international application includes the necessary components for an international filing date (see PCT Article 11 and MPEP 1810), a Notification of the International Application Number and of the International Filing Date (Form PCT/RO/105) will be issued in due course, subject to prescriptions concerning national security, and the date shown on this Acknowledgement Receipt will establish the international filing date of the application. UNITED STATES PATENT AND TRADEMARK OFFICE



UNITED STATES DEPARTMENT OF COMMERCE United States Patent and Trademark Office Address: COMMISSIONER FOR PATENTS P.O. Box 1450 Alexandria, Virginia 22313-1450 www.uspto.gov

NOTICE OF ALLOWANCE AND FEE(S) DUE

 38139
 7590
 05/17/2012

 Konda Technologies, Inc
 6278 GRAND OAK WAY
 5AN JOSE, CA 95135

EXAMINER

TAN, VIBOL

ART UNIT PAPER NUMBER
2819

DATE MAILED: 05/17/2012

APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
12/601,275	05/31/2010	Venkat Konda	V-0045US	6372

TITLE OF INVENTION: VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS

APPLN. TYPE	SMALL ENTITY	ISSUE FEE DUE	PUBLICATION FEE DUE	PREV. PAID ISSUE FEE	TOTAL FEE(S) DUE	DATE DUE
nonprovisional	YES	\$870	\$300	\$O	\$1170	08/17/2012

THE APPLICATION IDENTIFIED ABOVE HAS BEEN EXAMINED AND IS ALLOWED FOR ISSUANCE AS A PATENT. <u>PROSECUTION ON THE MERITS IS CLOSED</u>. THIS NOTICE OF ALLOWANCE IS NOT A GRANT OF PATENT RIGHTS. THIS APPLICATION IS SUBJECT TO WITHDRAWAL FROM ISSUE AT THE INITIATIVE OF THE OFFICE OR UPON PETITION BY THE APPLICANT. SEE 37 CFR 1.313 AND MPEP 1308.

THE ISSUE FEE AND PUBLICATION FEE (IF REQUIRED) MUST BE PAID WITHIN <u>THREE MONTHS</u> FROM THE MAILING DATE OF THIS NOTICE OR THIS APPLICATION SHALL BE REGARDED AS ABANDONED. <u>THIS STATUTORY PERIOD CANNOT BE EXTENDED</u>. SEE 35 U.S.C. 151. THE ISSUE FEE DUE INDICATED ABOVE DOES NOT REFLECT A CREDIT FOR ANY PREVIOUSLY PAID ISSUE FEE IN THIS APPLICATION. IF AN ISSUE FEE HAS PREVIOUSLY BEEN PAID IN THIS APPLICATION (AS SHOWN ABOVE), THE RETURN OF PART B OF THIS FORM WILL BE CONSIDERED A REQUEST TO REAPPLY THE PREVIOUSLY PAID ISSUE FEE TOWARD THE ISSUE FEE NOW DUE.

HOW TO REPLY TO THIS NOTICE:

I. Review the SMALL ENTITY status shown above.

If the SMALL ENTITY is shown as YES, verify your current SMALL ENTITY status:	If the SMALL ENTITY is shown as NO:
A. If the status is the same, pay the TOTAL FEE(S) DUE shown above.	A. Pay TOTAL FEE(S) DUE shown above, or
B. If the status above is to be removed, check box 5b on Part B - Fee(s) Transmittal and pay the PUBLICATION FEE (if required) and twice the amount of the ISSUE FEE shown above, or	B. If applicant claimed SMALL ENTITY status before, or is now claiming SMALL ENTITY status, check box 5a on Part B - Fee(s) Transmittal and pay the PUBLICATION FEE (if required) and 1/2 the ISSUE FEE shown above.

II. PART B - FEE(S) TRANSMITTAL, or its equivalent, must be completed and returned to the United States Patent and Trademark Office (USPTO) with your ISSUE FEE and PUBLICATION FEE (if required). If you are charging the fee(s) to your deposit account, section "4b" of Part B - Fee(s) Transmittal should be completed and an extra copy of the form should be submitted. If an equivalent of Part B is filed, a request to reapply a previously paid issue fee must be clearly made, and delays in processing may occur due to the difficulty in recognizing the paper as an equivalent of Part B.

III. All communications regarding this application must give the application number. Please direct all communications prior to issuance to Mail Stop ISSUE FEE unless advised to the contrary.

IMPORTANT REMINDER: Utility patents issuing on applications filed on or after Dec. 12, 1980 may require payment of maintenance fees. It is patentee's responsibility to ensure timely payment of maintenance fees when due.

PART B - FEE(S) TRANSMITTAL

Complete and send this form, together with applicable fee(s), to: <u>Mail</u> Mail Stop ISSUE FEE Commissioner for Patents P.O. Box 1450 Alexandria, Virginia 22313-1450

or Fax (571)-273-2885

INSTRUCTIONS: This form should be used for transmitting the ISSUE FEE and PUBLICATION FEE (if required). Blocks 1 through 5 should be completed where appropriate. All further correspondence including the Patent, advance orders and notification of maintenance fees will be mailed to the current correspondence address as indicated unless corrected below or directed otherwise in Block 1, by (a) specifying a new correspondence address; and/or (b) indicating a separate "FEE ADDRESS" for maintenance fee notifications. Note: A certificate of mailing can only be used for domestic mailings of the Fee(s) Transmittal. This certificate cannot be used for any other accompanying papers. Each additional paper, such as an assignment or formal drawing, must CURRENT CORRESPONDENCE ADDRESS (Note: Use Block 1 for any change of address) have its own certificate of mailing or transmission. 38139 7590 05/17/2012 **Certificate of Mailing or Transmission** Konda Technologies, Inc I hereby certify that this Fee(s) Transmittal is being deposited with the United States Postal Service with sufficient postage for first class mail in an envelope addressed to the Mail Stop ISSUE FEE address above, or being facsimile transmitted to the USPTO (571) 273-2885, on the date indicated below. 6278 GRAND OAK WAY SAN JOSE, CA 95135 (Depositor's name (Signature Date APPLICATION NO. FILING DATE FIRST NAMED INVENTOR ATTORNEY DOCKET NO. CONFIRMATION NO. 12/601,275 05/31/2010 Venkat Konda V-0045US 6372 TITLE OF INVENTION: VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS PUBLICATION FEE DUE PREV. PAID ISSUE FEE TOTAL FEE(S) DUE APPLN, TYPE SMALL ENTITY ISSUE FEE DUE DATE DUE YES \$870 \$300 \$0 \$1170 08/17/2012 nonprovisional EXAMINER CLASS-SUBCLASS ART UNIT

TAN, VIBOL	2819	326-041000	
 1. Change of correspondence address or indication of CFR 1.363). Change of correspondence address (or Change Address form PTO/SB/122) attached. "Fee Address" indication (or "Fee Address" I PTO/SB/47; Rev 03-02 or more recent) attached Number is required. 	e of Correspondence	 For printing on the patent front page, list (1) the names of up to 3 registered patent attorneys or agents OR, alternatively, (2) the name of a single firm (having as a member a registered attorney or agent) and the names of up to 2 registered patent attorneys or agents. If no name is listed, no name will be printed. 	1 2 3
Number is required.		nores, no mane traine er printeet	

3. ASSIGNEE NAME AND RESIDENCE DATA TO BE PRINTED ON THE PATENT (print or type)

PLEASE NOTE: Unless an assignee is identified below, no assignee data will appear on the patent. If an assignee is identified below, the document has been filed for recordation as set forth in 37 CFR 3.11. Completion of this form is NOT a substitute for filing an assignment. (A) NAME OF ASSIGNEE (B) RESIDENCE: (CITY and STATE OR COUNTRY)

Please check the appropriate assignee category or categories (will not be printed on the patent):	Individual Corporation or other private group entity Governm	ient

 4a. The following fee(s) are submitted: Issue Fee Publication Fee (No small entity discount permitted) Advance Order - # of Copies	 4b. Payment of Fee(s): (Please first reapply any previously paid issue fee shown above) A check is enclosed. Payment by credit card. Form PTO-2038 is attached. The Director is hereby authorized to charge the required fee(s), any deficiency, or credit any overpayment, to Deposit Account Number (enclose an extra copy of this form).
 5. Change in Entity Status (from status indicated above) a. Applicant claims SMALL ENTITY status. See 37 CFR 1.27. 	b. Applicant is no longer claiming SMALL ENTITY status. See 37 CFR 1.27(g)(2).
NOTE: The Issue Fee and Publication Fee (if required) will not be acce interest as shown by the records of the United States Patent and Traden	pted from anyone other than the applicant; a registered attorney or agent; or the assignee or other party in tark Office.
Authorized Signature	Date
Typed or printed name	Registration No
This collection of information is required by 37 CFR 1.311. The inform an application Confidentiality is governed by 35 U.S.C. 122 and 37 C	nation is required to obtain or retain a benefit by the public which is to file (and by the USPTO to process) FR 1.14. This collection is estimated to take 12 minutes to complete including gathering, preparing, and

This collection of information is required by 37 CFR 1.311. The information is required to obtain or retain a benefit by the public which is to file (and by the USPTO to process) an application. Confidentiality is governed by 35 U.S.C. 122 and 37 CFR 1.14. This collection is estimated to take 12 minutes to complete, including gathering, preparing, and submitting the completed application form to the USPTO. Time will vary depending upon the individual case. Any comments on the amount of time you require to complete this form and/or suggestions for reducing this burden, should be sent to the Chief Information Officer, U.S. Patent and Trademark Office, U.S. Department of Commerce, P.O. Box 1450, Alexandria, Virginia 22313-1450. DO NOT SEND FEES OR COMPLETED FORMS TO THIS ADDRESS. SEND TO: Commissioner for Patents, P.O. Box 1450, Alexandria, Virginia 22313-1450.

Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number.

	ted States Pate	NT AND TRADEMARK OFFICE	UNITED STATES DEPAR United States Patent and Address: COMMISSIONER F P.O. Box 1450 Alexandria, Virginia 223 www.uspto.gov	FOR PATENTS
APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
12/601,275	05/31/2010	Venkat Konda	V-0045US	6372
38139 75	90 05/17/2012		EXAM	IINER
Konda Technolog 6278 GRAND OAI	gies, Inc		TAN, Y	VIBOL
SAN JOSE, CA 95	135		ART UNIT	PAPER NUMBER
			2819	
			DATE MAILED: 05/17/201	2

Determination of Patent Term Adjustment under 35 U.S.C. 154 (b)

(application filed on or after May 29, 2000)

The Patent Term Adjustment to date is 183 day(s). If the issue fee is paid on the date that is three months after the mailing date of this notice and the patent issues on the Tuesday before the date that is 28 weeks (six and a half months) after the mailing date of this notice, the Patent Term Adjustment will be 183 day(s).

If a Continued Prosecution Application (CPA) was filed in the above-identified application, the filing date that determines Patent Term Adjustment is the filing date of the most recent CPA.

Applicant will be able to obtain more detailed information by accessing the Patent Application Information Retrieval (PAIR) WEB site (http://pair.uspto.gov).

Any questions regarding the Patent Term Extension or Adjustment determination should be directed to the Office of Patent Legal Administration at (571)-272-7702. Questions relating to issue and publication fee payments should be directed to the Customer Service Center of the Office of Patent Publication at 1-(888)-786-0101 or (571)-272-4200.

Privacy Act Statement

The Privacy Act of 1974 (P.L. 93-579) requires that you be given certain information in connection with your submission of the attached form related to a patent application or patent. Accordingly, pursuant to the requirements of the Act, please be advised that: (1) the general authority for the collection of this information is 35 U.S.C. 2(b)(2); (2) furnishing of the information solicited is voluntary; and (3) the principal purpose for which the information is used by the U.S. Patent and Trademark Office is to process and/or examine your submission related to a patent application or patent. If you do not furnish the requested information, the U.S. Patent and Trademark Office may not be able to process and/or examine your submission, which may result in termination of proceedings or abandonment of the application or expiration of the patent.

The information provided by you in this form will be subject to the following routine uses:

- 1. The information on this form will be treated confidentially to the extent allowed under the Freedom of Information Act (5 U.S.C. 552) and the Privacy Act (5 U.S.C 552a). Records from this system of records may be disclosed to the Department of Justice to determine whether disclosure of these records is required by the Freedom of Information Act.
- 2. A record from this system of records may be disclosed, as a routine use, in the course of presenting evidence to a court, magistrate, or administrative tribunal, including disclosures to opposing counsel in the course of settlement negotiations.
- 3. A record in this system of records may be disclosed, as a routine use, to a Member of Congress submitting a request involving an individual, to whom the record pertains, when the individual has requested assistance from the Member with respect to the subject matter of the record.
- 4. A record in this system of records may be disclosed, as a routine use, to a contractor of the Agency having need for the information in order to perform a contract. Recipients of information shall be required to comply with the requirements of the Privacy Act of 1974, as amended, pursuant to 5 U.S.C. 552a(m).
- 5. A record related to an International Application filed under the Patent Cooperation Treaty in this system of records may be disclosed, as a routine use, to the International Bureau of the World Intellectual Property Organization, pursuant to the Patent Cooperation Treaty.
- 6. A record in this system of records may be disclosed, as a routine use, to another federal agency for purposes of National Security review (35 U.S.C. 181) and for review pursuant to the Atomic Energy Act (42 U.S.C. 218(c)).
- 7. A record from this system of records may be disclosed, as a routine use, to the Administrator, General Services, or his/her designee, during an inspection of records conducted by GSA as part of that agency's responsibility to recommend improvements in records management practices and programs, under authority of 44 U.S.C. 2904 and 2906. Such disclosure shall be made in accordance with the GSA regulations governing inspection of records for this purpose, and any other relevant (i.e., GSA or Commerce) directive. Such disclosure shall not be used to make determinations about individuals.
- 8. A record from this system of records may be disclosed, as a routine use, to the public after either publication of the application pursuant to 35 U.S.C. 122(b) or issuance of a patent pursuant to 35 U.S.C. 151. Further, a record may be disclosed, subject to the limitations of 37 CFR 1.14, as a routine use, to the public if the record was filed in an application which became abandoned or in which the proceedings were terminated and which application is referenced by either a published application, an application open to public inspection or an issued patent.
- 9. A record from this system of records may be disclosed, as a routine use, to a Federal, State, or local law enforcement agency, if the USPTO becomes aware of a violation or potential violation of law or regulation.

	Application No.	Applicant(s)
	12/601,275	KONDA, VENKAT
Notice of Allowability	Examiner	Art Unit
	Vibol Tan	2819
The MAILING DATE of this communication app All claims being allowable, PROSECUTION ON THE MERITS IS herewith (or previously mailed), a Notice of Allowance (PTOL-85 NOTICE OF ALLOWABILITY IS NOT A GRANT OF PATENT F of the Office or upon petition by the applicant. See 37 CFR 1.31 1. X This communication is responsive to <u>amendment filed 5/8/</u>	6 (OR REMAINS) CLOSED ir i) or other appropriate community RIGHTS. This application is a 3 and MPEP 1308.	n this application. If not included unication will be mailed in due course. THIS
 An election was made by the applicant in response to a restriction requirement and election have been incorporat 		during the interview on;
3. ⊠ The allowed claim(s) is/are <u>1 and 3-49</u> .		
 4. Acknowledgment is made of a claim for foreign priority undated a) All b) Some* c) None of the: Certified copies of the priority documents have Certified copies of the priority documents have Copies of the certified copies of the priority documents have Copies of the certified copies of the priority documents have Copies of the certified copies of the priority documents have Copies of the certified copies of the priority documents have Copies of the certified copies of the priority documents have Copies of the certified copies of the priority documents have Copies of the certified copies of the priority documents have Certified copies not received: Applicant has THREE MONTHS FROM THE "MAILING DATE" noted below. Failure to timely comply will result in ABANDONI THIS THREE-MONTH PERIOD IS NOT EXTENDABLE. A SUBSTITUTE OATH OR DECLARATION must be submininformation (PTO-152) which give CORRECTED DRAWINGS (as "replacement sheets") must (a) including changes required by the Notice of Draftspering including changes required by the Notice of Draftspering (b) including changes required by the attached Examiner Paper No./Mail Date	e been received. The been received in Application ocuments have been received i' of this communication to file MENT of this application. witted. Note the attached EXA res reason(s) why the oath or st be submitted. reson's Patent Drawing Review ''''''''''''''''''''''''''''''''''	on No d in this national stage application from the e a reply complying with the requirements MINER'S AMENDMENT or NOTICE OF r declaration is deficient. v (PTO-948) attached r in the Office action of the drawings in the front (not the back) of rR 1.121(d). ust be submitted. Note the
Attachment(s) 1. □ Notice of References Cited (PTO-892) 2. □ Notice of Draftperson's Patent Drawing Review (PTO-948) 3. □ Information Disclosure Statements (PTO/SB/08), Paper No./Mail Date	6. ☐ Interview Su Paper No./ 7. ☐ Examiner's	formal Patent Application ummary (PTO-413), 'Mail Date Amendment/Comment Statement of Reasons for Allowance

U.S. Patent and Trademark Office PTOL-37 (Rev. 03-11)

Allowable Subject Matter

1. The following is an examiner's statement of reasons for allowance: in combination with other limitations of the claims, the cited prior arts fail to teach said all straight links are connecting from switches in each said sub-integrated circuit block are connecting to switches in the same said sub-integrated circuit block; and said all cross links are connecting as either vertical or horizontal links between switches in two different said sub-integrated circuit blocks which are either placed vertically above or below, or placed horizontally to the left or to the right, each said plurality of sub-integrated circuit blocks comprising same number of said stares and said switches in each said stage, regardless of the size of said two-dimensional grid so that each said plurality of sub-integrated circuit block with its corresponding said stages and said switches in each stage is replicable in both vertical direction or horizontal direction of said two-dimensional grid, as required by amended claim 1.

2. Claims 1 and 3-49 are now in condition for allowance.

Any comments considered necessary by applicant must be submitted no later than the payment of the issue fee and, to avoid processing delays, should preferably accompany the issue fee. Such submissions should be clearly labeled "Comments on Statement of Reasons for Allowance."

Any inquiry concerning this communication or earlier communications from the examiner should be directed to Vibol Tan whose telephone number is (571)272-1811. The examiner can normally be reached on Monday-Friday (7:00 AM-4:30 PM).

Application/Control Number: 12/601,275 Art Unit: 2819

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Shawki Ismail can be reached on (571) 272-3985. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see http://pair-direct.uspto.gov. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

/Vibol Tan/ Primary Examiner, Art Unit 2819



UNITED STATES PATENT AND TRADEMARK OFFICE

UNITED STATES DEPARTMENT OF COMMERCE United States Patent and Trademark Office Address: COMMISSIONER FOR PATENTS P.O. Box 1450 Alexandria, Virginia 22313-1450 www.uspto.gov

BIB DATA SHEET

CONFIRMATION NO. 6372

SERIAL NUM	BER	FILING or	371(c)		CLASS	GR	OUP ART UNIT ATTORNEY DOC			
12/601,275	275 DAT 275 05/31/2				326		2819		NO. V-0045US	
APPLICANTS	3								1	
Venkat Konda, San Jose, CA;										
** CONTINUING DATA ***********************************										
** FOREIGN AF	PLICA	TIONS ******	*******	******	*					
** IF REQUIRE 04/08/201		EIGN FILING	LICENS	E GRA	NTED ** ** SMA		NTITY **			
Foreign Priority claimed		Yes Vo			STATE OR		HEETS	тот		INDEPENDENT
35 USC 119(a-d) cond Verified and	litions met VIBOL TAN		Met after Allowance VTan		COUNTRY				MS	CLAIMS
	Examiner's S		Initials		CA		39	49)	1
ADDRESS										
Konda Te 6278 GRA SAN JOSI UNITED S	AND OÀ E, CA 9	AK WAY 95135								
TITLE										
VLSI LAY	OUTS	OF FULLY CO		ED GE	NERALIZED NE	TWO	RKS			
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I FILING FEE I Processing Ext of tir						ES: Authority has been given in Paper			ing Ext. of time)	
	Noto charge/credit DEPOSIT ACCOUNT Nofor following:						<u> </u>			
							Credit			
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BIB (Rev. 05/07).

EAST Search History

EAST Search History (Prior Art)

Ref #	Hits	Search Query	DBs	Default Operator	Plurals	Time Stamp
S25	35	(US-20080267182-\$ or US- 20040114586-\$ or US-20030021267-\$ or US-20030112797-\$ or US- 20030053456-\$ or US-20020051447-\$ or US-20020031118-\$ or US- 20020031117-\$).did. or (US-8098081-\$ or US-7924052-\$ or US-7468974-\$ or US-7424011-\$ or US-7424010-\$ or US- 7397796-\$ or US-7349387-\$ or US- 7346049-\$ or US-7349387-\$ or US- 6567858-\$ or US-7130920-\$ or US- 6456838-\$ or US-6157643-\$ or US- 6456838-\$ or US-6157643-\$ or US- 5406556-\$ or US-856977-\$ or US- 5406556-\$ or US-7103059-\$ or US- 7139266-\$ or US-7065074-\$ or US- 7050429-\$ or US-7042873-\$ or US- 7051303-\$ or US-7016345-\$ or US- 6952418-\$).did. or (US-6201808-\$).did.	US-PGPUB; USPAT	OR	OFF	2012/05/08 15:44
S26	42	konda-venkat.in.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/05/08 15:44
S27	84	(routing adj network) and (multicast adj connections)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/05/08 15:58
S28	1365	(routing adj network) and (inputs) and (outputs)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/05/08 15:59
S29	229142	(("326") or ("370")).CLAS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/05/08 15:59
S30	693	S28 and S29	US-PGPUB; USPAT; USOCR; FPRS;	OR	OFF	2012/05/08 15:59

Page 31 of 374 file:///Cl/Users/vtan/Documents/e-Red%20Folder/12601275/EASTSearchHistory.12601275_AccessibleVersion.htm[5/10/2012 7:09:33 AM]

			EPO; JPO; DERWENT; IBM_TDB			
S31	380	S30 and links	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/05/08 16:00
S 32	1	S30 and (straight adj links) and (vertical adj links) and (horizontal adj links)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/05/08 16:01
S33	1	((straight adj links) and (vertical adj links) and (horizontal adj links)).clm.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/05/08 16:02

EAST Search History (Interference)

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	Application/Control No.	Applicant(s)/Patent Under Reexamination
Search Notes	12601275	KONDA, VENKAT
	Examiner	Art Unit
	VIBOL TAN	2819

	SEARCHED											
Class	Subclass	Date	Examiner									
326	38-41	2/2/12	VTan									
370	390, 312, 360, 388, 412	2/2/12	VTan									

SEARCH NOTES										
Search Notes	Date	Examiner								
Inventor search	2/2/12	VTan								
search report pct/us2008/064605	2/1/12	VTan								
EAST text search	2/2/12	VTan								
EAST text update	5/8/12	VTan								

	INTERFERENCE SEARCH		
Class	Subclass	Date	Examiner
uspgpub, uspat & upad	text search	5/8/12	VTan

	Application/Control No.	Applicant(s)/Patent Under Reexamination
Issue Classification	12601275	KONDA, VENKAT
	Examiner	Art Unit
	VIBOL TAN	2819

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NONE		Total Clain	ns Allowed:
(Assistant Examiner)	(Date)	4	8
/VIBOL TAN/ Primary Examiner.Art Unit 2819	05/10/2012	O.G. Print Claim(s)	O.G. Print Figure
(Primary Examiner)	(Date)	1	1c

U.S. Patent and Trademark Office

Index of Claims						Application/Control No. 12601275 Examiner VIBOL TAN Cancelled N Non-Ele							Applicant(s)/Patent Under Reexamination KONDA, VENKAT Art Unit 2819 Acted Appeal					
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U.S. Patent and Trademark Office

Part of Paper No. : 20120508

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In The United States Patent And Trademark Office

Application Number: 12/601,275

Application Filed: 11/22/2009

5 Applicant(s): Venkat Konda
Title: VLSI Layouts of Fully Connected Generalized Networks
Examiner/Art Unit: Vibol Tan / 2819
International Application Number: PCT/US08/64605
International Application Filed: 5/22/2008
10 Priority Date: 5/25/2007

San Jose, 2012 May 8, Tue

SUPPLEMENTAL AMENDMENT

(and the response to office letter dated 2/7/2012)

15

Mail Stop Amendment

Commissioner for Patents

P.O. Box 1450

Alexandria, Virginia, 22313-1450

20

Dear Sir/Madam:

In response to the office action mailed 2012 February 7, Applicant has submitted Amendment A on April 30, 2012. In the Amendment A filed on April 30, 2012,

Applicant has made mistakes with the format of the claims and not identifying the cancelled claim. To correct those mistakes, please consider the following Supplemental Amendment.

5 AMENDMENTS TO THE CLAIMS

This listing of claims will replace all prior versions and listings of claims in the application.

CLAIMS

What is claimed is:

1. (currently amended): An integrated circuit device comprising a plurality of subintegrated circuit blocks and a routing network, and

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Said each plurality of sub-integrated circuit blocks comprising a plurality of inlet links and a plurality of outlet links; and

Said routing network interconnects any one of said outlet link of one of said subintegrated circuit block to one or more said inlet links of one or more of said subintegrated circuit blocks; and

10 Said routing network comprising of a plurality of stages y, in each said subintegrated circuit block, starting from the lowest stage of 1 to the highest stage of y, where $y \ge 1$; and

Said routing network comprising a plurality of switches of size $d \times d$, where $d \ge 2$, in each said stage and each said switch of size $d \times d$ having d inlet links and d 15 outlet links; and

Said each sub-integrated circuit block comprising a plurality of said switches corresponding to each said stage; and

Said plurality of outlet links of said each sub-integrated circuit block are directly connected to said inlet links of said switches of its corresponding said lowest stage of 1,

20 <u>and said plurality of inlet links of said each sub-integrated circuit block are directly</u> connected from said outlet links of said switches of its corresponding said lowest stage of <u>1; and</u>

Said each sub-integrated circuit block comprising a plurality of forward connecting links connecting from switches in \underline{a} lower stage to switches in the its

25 immediate succeeding higher stage, and also comprising a plurality of backward connecting links connecting from switches in <u>a</u> higher stage to switches in <u>the its</u> immediate preceding lower stage; and

Said each sub-integrated circuit block comprising a plurality straight links in said forward connecting links from switches in <u>said each</u> lower stage to switches in <u>the its</u> immediate succeeding higher stage and a plurality cross links in said forward connecting links from switches in <u>said each</u> lower stage to switches in <u>the its</u> immediate succeeding

5 higher stage, and further comprising a plurality of straight links in said backward connecting links from switches in <u>said each</u> higher stage to switches in <u>the its</u> immediate preceding lower stage and a plurality of cross links in said backward connecting links from switches in <u>said each</u> higher stage to switches in <u>the its</u> immediate preceding lower stage.

10 <u>said plurality of sub-integrated circuit blocks arranged in a two-dimensional grid</u> of rows and columns, and

said all straight links are connecting from switches in each said sub-integrated circuit block are connecting to switches in the same said sub-integrated circuit block; and said all cross links are connecting as either vertical or horizontal links between switches

15 in two different said sub-integrated circuit blocks which are either placed vertically above or below, or placed horizontally to the left or to the right,

each said plurality of sub-integrated circuit blocks comprising same number of said stages and said switches in each said stage, regardless of the size of said twodimensional grid so that each said plurality of sub-integrated circuit block with its

20 <u>corresponding said stages and said switches in each stage is replicable in both vertical</u> <u>direction or horizontal direction of said two-dimensional grid</u>.

 (cancelled): The integrated circuit device of claim 1, wherein said all straight links are connecting from switches in each said sub-integrated circuit block are
 connecting to switches in the same said sub-integrated circuit block; and said all cross links are connecting as either vertical or horizontal links between switches in two different said sub-integrated circuit blocks.

3. (currently amended): The integrated circuit device of claim 1, wherein said plurality of sub-integrated circuit blocks arranged in a two-dimensional grid

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said two-dimensional grid of said sub-integrated circuit blocks with their corresponding said stages and said switches in each stage is scalable by any power of 2, and,

for each multiplication of 2 of the size of total said sub-integrated circuit blocks.
 by adding one more stage of switches and the layout is placed in hypercube format and also the cross links between said one more stage of switches are connected in hypercube format.

4. (Original): The integrated circuit device of claim 3, wherein said cross links in succeeding stages are connecting as alternative vertical and horizontal links
10 between switches in said sub-integrated circuit blocks.

5. (Original): The integrated circuit device of claim 4, wherein said cross links from switches in a stage in one of said sub-integrated circuit blocks are connecting to switches in the succeeding stage in another of said sub-integrated circuit blocks so that said cross links are either vertical links or horizontal and vice versa, and hereinafter such cross links are "shuffle exchange links").

6. (Original): The integrated circuit device of claim 5, wherein said all
horizontal shuffle exchange links between switches in any two corresponding said
succeeding stages are substantially of equal length and said vertical shuffle exchange
links between switches in any two corresponding said succeeding stages are substantially
of equal length in the entire said integrated circuit device.

7. (Original): The integrated circuit device of claim 6, wherein the shortest horizontal shuffle exchange links are connecting at the lowest stage and between switches in two nearest neighboring said sub-integrated circuit blocks, and length of the horizontal shuffle exchange links is doubled in each succeeding stage; and the shortest vertical

25 shuffle exchange links are connecting at the lowest stage and between switches in two nearest neighboring said sub-integrated circuit blocks, and length of the vertical shuffle exchange links is doubled in each succeeding stage.

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8. (currently amended): The integrated circuit device of claim 7, wherein $y \ge (\log_2 N)$, where $N > 1_{a}$ so that the length of the horizontal shuffle exchange links in the highest stage is equal to half the size of the horizontal size of said two dimensional grid of sub-integrated circuit blocks and the length of the vertical shuffle exchange links in the highest stage is equal to half the size of the vertical size of said two dimensional

grid of sub-integrated circuit blocks.

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9. (Original): The integrated circuit device of claim 8, wherein d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage

10 in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast Benes network with full bandwidth.

10. (Original): The integrated circuit device of claim 8, wherein d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block
15 connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast Benes network and rearrangeably nonblocking for arbitrary fan-out multicast Benes network with full bandwidth.

- 20 11. (Original): The integrated circuit device of claim 8, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast Benes 25 network with full bandwidth.
 - 12. (currently amended): The integrated circuit device of claim 7, wherein $y \ge (\log_2 N)$, where N > 1, so that the length of the horizontal shuffle exchange links in the highest stage is equal to half the size of the horizontal size of said two dimensional grid of sub-integrated circuit blocks and the length of the vertical shuffle exchange links

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in the highest stage is equal to half the size of the vertical size of said two dimensional grid of sub-integrated circuit blocks, and

said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks.

- 5 13. (Original): The integrated circuit device of claim 12, wherein d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast butterfly fat tree network
- 10 with full bandwidth.

14. (Original): The integrated circuit device of claim 12, wherein d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links
15 and said routing network is strictly nonblocking for unicast butterfly fat tree network and rearrangeably nonblocking for arbitrary fan-out multicast butterfly fat tree network with full bandwidth.

15. (Original): The integrated circuit device of claim 12, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block
20 connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast butterfly fat tree network with full bandwidth.

16. (Original): The integrated circuit device of claim 1, wherein said25 horizontal and vertical links are implemented on two or more metal layers.

17. (Original): The integrated circuit device of claim 1, wherein said switches comprising active and reprogrammable cross points and said each cross point is programmable by an SRAM cell or a Flash Cell.

18. (Original): The integrated circuit device of claim 1, wherein said sub-5 integrated circuit blocks are of equal die size.

19. (Original): The integrated circuit device of claim 16, wherein said subintegrated circuit blocks are Lookup Tables (hereinafter "LUTs") and said integrated circuit device is a field programmable gate array (FPGA) device or field programmable gate array (FPGA) block embedded in another integrated circuit device.

10 20. (Original): The integrated circuit device of claim 16, wherein said subintegrated circuit blocks are AND or OR gates and said integrated circuit device is a programmable logic device (PLD).

21. (Original): The integrated circuit device of claim 1, wherein said subintegrated circuit blocks comprising any arbitrary hardware logic or memory circuits.

15 22. (Original): The integrated circuit device of claim 1, wherein said switches comprising active one-time programmable cross points and said integrated circuit device is a mask programmable gate array (MPGA) device or a structured ASIC device.

23. (Original): The integrated circuit device of claim 1, wherein said switches comprising passive cross points or just connection of two links or not and said integrated
20 circuit device is a Application Specific Integrated Circuit (ASIC) device.

24. (Original): The integrated circuit device of claim 1, wherein said subintegrated circuit blocks further recursively comprise one or more super-sub-integrated circuit blocks and a sub-routing network.

25. (currently amended): The integrated circuit device of claim 5, wherein25 said all horizontal shuffle exchange links between switches in any two corresponding said

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succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and $y \ge (\log_2 N)$, where N > 1.

- 26. (Original): The integrated circuit device of claim 25, wherein d = 2 and
 5 there is only one switch in each said stage in each said sub-integrated circuit block
 connecting said forward connecting links and there is only one switch in each said stage
 in each said sub-integrated circuit block connecting said backward connecting links and
 said routing network is rearrangeably nonblocking for unicast generalized multi-stage
 network with full bandwidth.
- 10 27. (Original): The integrated circuit device of claim 25, wherein d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast generalized multi-stage
 15 network and rearrangeably nonblocking for arbitrary fan-out multicast generalized multi-stage network with full bandwidth.

28. (Original): The integrated circuit device of claim 25, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said
20 stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast generalized multi-stage network with full bandwidth.

29. (currently amended): The integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said
25 succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and y ≥ (log₂ N), where N > 1, and

said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks.

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30. (Original): The integrated circuit device of claim 29, wherein d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized butterfly fat tree network with full bandwidth.

31. (Original): The integrated circuit device of claim 29, wherein d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said
stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast generalized butterfly fat tree Network and rearrangeably nonblocking for arbitrary fan-out multicast generalized butterfly fat tree network with full bandwidth.

32. (Original): The integrated circuit device of claim 29, wherein d = 2 and 15 there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast generalized butterfly fat tree network with full bandwidth.

20 33. (Original): The integrated circuit device of claim 1, wherein said straight links connecting from switches in each said sub-integrated circuit block are connecting to switches in the same said sub-integrated circuit block; and

said cross links are connecting as vertical or horizontal or diagonal links between two different said sub-integrated circuit blocks.

25 34. (Original): The integrated circuit device of claim 8, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and

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said routing network is rearrangeably nonblocking for unicast multi-link Benes network with full bandwidth.

35. (Original): The integrated circuit device of claim 8, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block
5 connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast multi-link Benes network and rearrangeably nonblocking for arbitrary fan-out multicast multi-link Benes network with full bandwidth.

36. (Original): The integrated circuit device of claim 8, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast multi-link
 Benes network with full bandwidth.

37. (Original): The integrated circuit device of claim 12, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast multi-link butterfly fat tree

network with full bandwidth.

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38. (Original): The integrated circuit device of claim 12, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said
25 stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast multi-link butterfly fat tree network and rearrangeably nonblocking for arbitrary fan-out multicast multi-link butterfly fat tree network with full bandwidth.

39. (Original): The integrated circuit device of claim 12, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast multi-link butterfly fat tree network with full bandwidth.

40. (currently amended): The integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between

10 switches in any two corresponding said succeeding stages are of different length and $y \ge (\log_2 N)$, where N > 1.

41. (Original): The integrated circuit device of claim 40, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage
15 in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized multi-link multi-stage network with full bandwidth.

42. (Original): The integrated circuit device of claim 40, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block
20 connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast generalized multi-link multi-stage network and rearrangeably nonblocking for arbitrary fan-out multicast generalized multi-link multi-stage network with full bandwidth.

25 43. (Original): The integrated circuit device of claim 40, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links

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and said routing network is strictly nonblocking for arbitrary fan-out multicast generalized multi-link multi-stage network with full bandwidth.

44. (currently amended): The integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said
5 succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and y ≥ (log₂ N), where N > 1, and

said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks.

10 45. (Original): The integrated circuit device of claim 44, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized multi-link
15 butterfly fat tree network with full bandwidth.

46. (Original): The integrated circuit device of claim 44, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links
and said routing network is strictly nonblocking for unicast generalized multi-link butterfly fat tree Network and rearrangeably nonblocking for arbitrary fan-out multicast generalized multi-link butterfly fat tree network with full bandwidth.

47. (Original): The integrated circuit device of claim 44, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block
25 connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast generalized multi-link butterfly fat tree network with full bandwidth.

48. (currently amended): The integrated circuit device of claim 1, wherein said plurality of forward connecting links use a plurality of buffers to amplify signals driven through them and said plurality of backward connecting links use a plurality of buffers to amplify signals driven through them; and said buffers can be inverting or non-inverting buffers

5 inverting buffers.

49. (currently amended): The integrated circuit device of claim 1, wherein said wherein said all switches of size $d \times d$ are either fully populated or partially populated.

CONCLUSION

For all of the above reasons, applicant submits that the Claims are now in proper form, and that the Claims all define patentably over the prior art. Therefore applicant submits that this application is now in condition for allowance, which action he respectfully solicits.

Conditional request for Constructive Assistance

Applicant has amended the claims of this application so that they are proper, definite, and define novel structure which is also unobvious. If, for any reason this application is not believed to be in full condition for allowance, applicant respectfully request the

constructive assistance and suggestions of the Examiner pursuant to M.P.E.P §
 2173.02 and § 707.07(j) in order that the undersigned can place this application in allowable condition as soon as possible and without the need for further proceedings.

Very respectfully,

15 /Venkat Konda/

5

Venkat Konda

Konda Technologies, Inc (USPTO Customer Number: 38139)

6278 Grand Oak Way

San Jose, CA 95135

20 Phone: 408-472-3273

Fax: 408-238-2478

Electronic Acknowledgement Receipt					
EFS ID:	12733551				
Application Number:	12601275				
International Application Number:					
Confirmation Number:	6372				
Title of Invention:	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS				
First Named Inventor/Applicant Name:	Venkat Konda				
Customer Number:	38139				
Filer:	Venkar Konda				
Filer Authorized By:					
Attorney Docket Number:	V-0045US				
Receipt Date:	08-MAY-2012				
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Time Stamp:	22:22:15				
Application Type:	U.S. National Stage under 35 USC 371				

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File Listing:							
Document Number	Document Description		File Name	File Size(Bytes)/ Message Digest	Multi Part /.zip	Pages (if appl.)	
1	Supplemental Response or		uppleAmend-V0045US.pdf	144109	no	15	
	Supplemental Amendment			3e7bd4b5740546aeb12e5e0cfd93e5979ed c5277			
Warnings:				· ·			
Information:							

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New Applications Under 35 U.S.C. 111

If a new application is being filed and the application includes the necessary components for a filing date (see 37 CFR 1.53(b)-(d) and MPEP 506), a Filing Receipt (37 CFR 1.54) will be issued in due course and the date shown on this Acknowledgement Receipt will establish the filing date of the application.

National Stage of an International Application under 35 U.S.C. 371

If a timely submission to enter the national stage of an international application is compliant with the conditions of 35 U.S.C. 371 and other applicable requirements a Form PCT/DO/EO/903 indicating acceptance of the application as a national stage submission under 35 U.S.C. 371 will be issued in addition to the Filing Receipt, in due course.

New International Application Filed with the USPTO as a Receiving Office

If a new international application is being filed and the international application includes the necessary components for an international filing date (see PCT Article 11 and MPEP 1810), a Notification of the International Application Number and of the International Filing Date (Form PCT/RO/105) will be issued in due course, subject to prescriptions concerning national security, and the date shown on this Acknowledgement Receipt will establish the international filing date of the application.

PTO/SB/06 (07-06)

Approved for use through 1/31/2007. OMB 0651-0032 U.S. Patent and Trademark Office; U.S. DEPARTMENT OF COMMERCE

-	Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number.										
P/	ATENT APPL	Substitute fo			N RECORD	A		Docket Number 1,275		ing Date 31/2010	To be Mailed
APPLICATION AS FILED – PART I (Column 1) (Column 2)						SMALL	entity 🛛	OR		HER THAN	
	FOR	N	UMBER FIL	.ED NU	MBER EXTRA		RATE (\$)	FEE (\$)		RATE (\$)	FEE (\$)
	BASIC FEE (37 CFR 1.16(a), (b), c	or (c))	N/A		N/A		N/A			N/A	
	SEARCH FEE (37 CFR 1.16(k), (i), c	or (m))	N/A		N/A		N/A			N/A	
	EXAMINATION FE (37 CFR 1.16(o), (p), o		N/A		N/A		N/A			N/A	
	AL CLAIMS CFR 1.16(i))		min	us 20 = *			X \$ =		OR	X \$ =	
	EPENDENT CLAIM CFR 1.16(h))	S	mi	nus 3 = *			X \$ =			X \$ =	
	APPLICATION SIZE 37 CFR 1.16(s))	FEE shee is \$2 addi	ts of pape 50 (\$125 ional 50 s	ation and drawin er, the application for small entity) sheets or fraction a)(1)(G) and 37	on size fee due for each n thereof. See						
	MULTIPLE DEPEN		,	477							
* If t	he difference in colu						TOTAL			TOTAL	
	APPI	LICATION AS	AMEND	ED – PART II						OTU	
		(Column 1)		(Column 2)	(Column 3)		OTHER THAN SMALL ENTITY OR SMALL ENTITY				
NT	05/08/2012	CLAIMS REMAINING AFTER AMENDMENT		HIGHEST NUMBER PREVIOUSLY PAID FOR	PRESENT EXTRA		RATE (\$)	ADDITIONAL FEE (\$)		RATE (\$)	ADDITIONAL FEE (\$)
ΜE	Total (37 CFR 1.16(i))	* 48	Minus	** 49	= 0		X \$30 =	0	OR	X \$ =	
ND	Independent (37 CFR 1.16(h))	* 1	Minus	***3	= 0		X \$125 =	0	OR	X \$ =	
AMENDMENT		ze Fee (37 CFR 1	.16(s))		•						
A		ITATION OF MULTI	PLE DEPEN	DENT CLAIM (37 CF	R 1.16(j))				OR		
							TOTAL ADD'L FEE	0	OR	TOTAL ADD'L FEE	
		(Column 1)		(Column 2)	(Column 3)					-	
_		CLAIMS REMAINING AFTER AMENDMENT		HIGHEST NUMBER PREVIOUSLY PAID FOR	PRESENT EXTRA		RATE (\$)	ADDITIONAL FEE (\$)		RATE (\$)	ADDITIONAL FEE (\$)
ENT	Total (37 CFR 1.16(i))	×	Minus	**	=		X \$ =		OR	X \$ =	
ENDM	Independent (37 CFR 1.16(h))	*	Minus	***	=		X \$ =		OR	X \$ =	
ШN	Application Si	ze Fee (37 CFR 1	.16(s))								
FIRST PRESENTATION OF MULTIPLE DEPENDENT CLAIM (37 CFR 1.16(j))								OR			
	the entry in column ⁻ the "Highest Numbe							nstrument Ex	OR amin	TOTAL ADD'L FEE er:	
***	f the "Highest Numb	er Previously Pai	d For" IN T	HIS SPACE is les	s than 3, enter "3".						
	The "Highest Number Previously Paid For" (Total or Independent) is the highest number found in the appropriate box in column 1. This collection of information is required by 37 CFR 1.16. The information is required to obtain or retain a benefit by the public which is to file (and by the USPTO to										

This collection of information is required by 37 CFR 1.16. The information is required to obtain or retain a benefit by the public which is to file (and by the USPTO to process) an application. Confidentiality is governed by 35 U.S.C. 122 and 37 CFR 1.14. This collection is estimated to take 12 minutes to complete, including gathering, preparing, and submitting the completed application form to the USPTO. Time will vary depending upon the individual case. Any comments on the amount of time you require to complete this form and/or suggestions for reducing this burden, should be sent to the Chief Information Officer, U.S. Patent and Trademark Office, U.S. Department of Commerce, P.O. Box 1450, Alexandria, VA 22313-1450. DO NOT SEND FEES OR COMPLETED FORMS TO THIS ADDRESS. SEND TO: Commissioner for Patents, P.O. Box 1450, Alexandria, VA 22313-1450.

If you need assistance in completing the form, call 1-800-PTO-9199 and select option 2.

PTO/SB/06 (07-06)

Approved for use through 1/31/2007. OMB 0651-0032 U.S. Patent and Trademark Office; U.S. DEPARTMENT OF COMMERCE

	Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it displays a valid OMB control number.										
PATENT APPLICATION FEE DETERMINATION RECORD Substitute for Form PTO-875						A		Docket Number 1,275		ing Date 31/2010	To be Mailed
APPLICATION AS FILED – PART I (Column 1) (Column 2)							SMALL	entity 🛛	OR		HER THAN
	FOR	Ν	UMBER FIL	.ED NUI	MBER EXTRA		RATE (\$)	FEE (\$)		RATE (\$)	FEE (\$)
	BASIC FEE (37 CFR 1.16(a), (b), c	or (c))	N/A		N/A		N/A			N/A	
	SEARCH FEE (37 CFR 1.16(k), (i), c	or (m))	N/A		N/A		N/A			N/A	
	EXAMINATION FE (37 CFR 1.16(o), (p), c		N/A		N/A		N/A			N/A	
(37	AL CLAIMS CFR 1.16(i))		mir	us 20 = *			X \$ =		OR	X \$ =	
	EPENDENT CLAIM CFR 1.16(h))	S	m	nus 3 = *			X \$ =			X \$ =	
	APPLICATION SIZE 37 CFR 1.16(s))	FEE shee is \$2 addi 35 L	ts of pap 50 (\$125 tional 50 s I.S.C. 41(ation and drawing er, the application for small entity) sheets or fraction a)(1)(G) and 37	n size fee due for each n thereof. See						
	MULTIPLE DEPEN		,				TOTAL			TOTAL	
° If t	he difference in colu		,				TOTAL			TOTAL	
	APPL	LICATION AS	AMEND)ED – PART II						OTU	
		(Column 1)		(Column 2)	(Column 3)		OTHER THAN SMALL ENTITY OR SMALL ENTITY				
١T	05/08/2012	CLAIMS REMAINING AFTER AMENDMENT		HIGHEST NUMBER PREVIOUSLY PAID FOR	PRESENT EXTRA		RATE (\$)	ADDITIONAL FEE (\$)		RATE (\$)	ADDITIONAL FEE (\$)
ME	Total (37 CFR 1.16(i))	* 48	Minus	** 49	= 0		X \$30 =	0	OR	X \$ =	
ND	Independent (37 CFR 1.16(h))	* 1	Minus	***3	= 0		X \$125 =	0	OR	X \$ =	
AMENDMENT		ze Fee (37 CFR [.]	.16(s))								
A	FIRST PRESEN	ITATION OF MULTI	PLE DEPEN	DENT CLAIM (37 CFI	R 1.16(j))				OR		
							TOTAL ADD'L FEE	0	OR	TOTAL ADD'L FEE	
		(Column 1)		(Column 2)	(Column 3)						
		CLAIMS REMAINING AFTER AMENDMENT		HIGHEST NUMBER PREVIOUSLY PAID FOR	PRESENT EXTRA		RATE (\$)	ADDITIONAL FEE (\$)		RATE (\$)	ADDITIONAL FEE (\$)
	Total (37 CFR 1.16(i))	*	Minus	**	=		X \$ =		OR	X \$ =	
ENDMENT	Independent (37 CFR 1.16(h))	*	Minus	***	=		X \$ =		OR	X \$ =	
ЫN	Application Si	ze Fee (37 CFR [.]	.16(s))								
AM	FIRST PRESENTATION OF MULTIPLE DEPENDENT CLAIM (37 CFR 1.16(j))								OR		
** lf ***	* If the entry in column 1 is less than the entry in column 2, write "0" in column 3. ** If the "Highest Number Previously Paid For" IN THIS SPACE is less than 20, enter "20". *** If the "Highest Number Previously Paid For" IN THIS SPACE is less than 3, enter "3". The "Highest Number Previously Paid For" (Total or Independent) is the highest number found in the appropriate box in column 1.										
	Figure realized realized realized in the provided of the information is required to obtain or retain a benefit by the public which is to file (and by the USPTO to										

This collection of information is required by 37 CFR 1.16. The information is required to obtain or retain a benefit by the public which is to file (and by the USP10 to process) an application. Confidentiality is governed by 35 U.S.C. 122 and 37 CFR 1.14. This collection is estimated to take 12 minutes to complete, including gathering, preparing, and submitting the completed application form to the USP10. Time will vary depending upon the individual case. Any comments on the amount of time you require to complete this form and/or suggestions for reducing this burden, should be sent to the Chief Information Officer, U.S. Patent and Trademark Office, U.S. Department of Commerce, P.O. Box 1450, Alexandria, VA 22313-1450. DO NOT SEND FEES OR COMPLETED FORMS TO THIS ADDRESS. **SEND TO: Commissioner for Patents, P.O. Box 1450, Alexandria, VA 22313-1450.**

If you need assistance in completing the form, call 1-800-PTO-9199 and select option 2.

	ted States Patent a	ND TRADEMARK OFFICE	UNITED STATES DEPAR United States Patent and Address: COMMISSIONER I P.O. Box 1450 Alexandria, Virginia 22 www.uspto.gov	FOR PATENTS			
APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.			
12/601,275	05/31/2010	Venkat Konda	V-0045US	6372			
38139 Konda Techno	7590 05/01/2012 logies Inc		EXAN	EXAMINER			
6278 GRAND	OAK WAY		TAN,	VIBOL			
SAN JOSE, CA	A 95135		ART UNIT	PAPER NUMBER			
			2819				
			MAIL DATE	DELIVERY MODE			
			05/01/2012	PAPER			

Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

	Application No. Applicant(s)								
Applicant-Initiated Interview Summary	12/601,275	KONDA, VENKA	т						
	Examiner	Art Unit							
	Vibol Tan	2819							
All participants (applicant, applicant's representative, PTO personnel):									
(1) <u>Vibol Tan</u> .									
(2) <u>Venkat Konda</u> .	(4)								
Date of Interview: <u>30 April 2012</u> .									
Type: 🛛 Telephonic 🔲 Video Conference 🗌 Personal [copy given to: 🗌 applicant 🛛 [applicant's representative]								
Exhibit shown or demonstration conducted: Yes [If Yes, brief description:	☐ No.								
Issues Discussed \Box 101 \Box 112 \boxtimes 102 \Box 103 \Box Othe (For each of the checked box(es) above, please describe below the issue and detail									
Claim(s) discussed: <u>1</u> .									
Identification of prior art discussed: <u>6,940,308 (Wong)</u> .									
Substance of Interview (For each issue discussed, provide a detailed description and indicate if agreement was reached. Some topics may include: identification or clarification of a reference or a portion thereof, claim interpretation, proposed amendments, arguments of any applied references etc)									
Agreement was reached for the proposed amendment of claim 1 to read as an integrated circuit device comprising a plurality of sub-integrated circuit blocks and a noting network, and Said each plurality of sub-integrated circuit blocks comprising a plurality of inlet links and a plurality of outlet links: andSaid routing network comprising of a plurality of stages , in each said sub-integrated circuit block, starting from the lowest stage of 1 to the highest stage of, where; and Said routing network comprising a plurality of switches of size, where, in each said stage and each said switch of size having inlet links and outlet links; andSaid plurality of outlet links of said each sub-integrated circuit block are directly connected to said inlet links; andSaid plurality of outlet links of said each sub-integrated circuit block are directly connected to said inlet links; and said switches of its corresponding said lowest stage of 1, and said plurality of inlet links of said each sub-integrated circuit block are directly connected from said outlet links of said switches of its corresponding said lowest stage of 1: and Said each sub-integrated circuit block comprising a plurality of forward connecting links connecting from switches in a lower stage to switches in the its immediate succeeding higher stage, and also comprising a plurality of backward connecting links connecting from switches in said each lower stage to switches in the its immediate succeeding higher stage. In said each lower stage to switches in said each lower stage to switches in said each lower stage to switches in the its immediate succeeding higher stage, and further comprising a plurality of sraight links in said forward connecting links from switches in said each higher stage to switches in the its immediate preceding lower stage said plurality of sub-integrated circuit block: and said all cross links in said all cross links and cannecting links from switches in said each higher stage to switches in the its immediate preceding lowe									

Applicant recordation instructions: The formal written reply to the last Office action must include the substance of the interview. (See MPEP section 713.04). If a reply to the last Office action has already been filed, applicant is given a non-extendable period of the longer of one month or thirty days from this interview date, or the mailing date of this interview summary form, whichever is later, to file a statement of the substance of the interview

Examiner recordation instructions: Examiners must summarize the substance of any interview of record. A complete and proper recordation of the substance of an interview should include the items listed in MPEP 713.04 for complete and proper recordation including the identification of the general thrust of each argument or issue discussed, a general indication of any other pertinent matters discussed regarding patentability and the general results or outcome of the interview, to include an indication as to whether or not agreement was reached on the issues raised.

Attachment

/Vibol Tan/ Primary Examiner, Art Unit 2819

U.S. Patent and Trademark Office PTOL-413 (Rev. 8/11/2010)

Interview Summary

Paper No. 20120430

Application No. 12601275 Summary of Record of Interview Requirements

Manual of Patent Examining Procedure (MPEP), Section 713.04, Substance of Interview Must be Made of Record

A complete written statement as to the substance of any face-to-face, video conference, or telephone interview with regard to an application must be made of record in the application whether or not an agreement with the examiner was reached at the interview.

Title 37 Code of Federal Regulations (CFR) § 1.133 Interviews

Paragraph (b)

In every instance where reconsideration is requested in view of an interview with an examiner, a complete written statement of the reasons presented at the interview as warranting favorable action must be filed by the applicant. An interview does not remove the necessity for reply to Office action as specified in §§ 1.111, 1.135. (35 U.S.C. 132)

37 CFR §1.2 Business to be transacted in writing.

All business with the Patent or Trademark Office should be transacted in writing. The personal attendance of applicants or their attorneys or agents at the Patent and Trademark Office is unnecessary. The action of the Patent and Trademark Office will be based exclusively on the written record in the Office. No attention will be paid to any alleged oral promise, stipulation, or understanding in relation to which there is disagreement or doubt.

The action of the Patent and Trademark Office cannot be based exclusively on the written record in the Office if that record is itself incomplete through the failure to record the substance of interviews.

It is the responsibility of the applicant or the attorney or agent to make the substance of an interview of record in the application file, unless the examiner indicates he or she will do so. It is the examiner's responsibility to see that such a record is made and to correct material inaccuracies which bear directly on the question of patentability.

Examiners must complete an Interview Summary Form for each interview held where a matter of substance has been discussed during the interview by checking the appropriate boxes and filling in the blanks. Discussions regarding only procedural matters, directed solely to restriction requirements for which interview recordation is otherwise provided for in Section 812.01 of the Manual of Patent Examining Procedure, or pointing out typographical errors or unreadable script in Office actions or the like, are excluded from the interview recordation procedures below. Where the substance of an interview is completely recorded in an Examiners Amendment, no separate Interview Summary Record is required.

The Interview Summary Form shall be given an appropriate Paper No., placed in the right hand portion of the file, and listed on the "Contents" section of the file wrapper. In a personal interview, a duplicate of the Form is given to the applicant (or attorney or agent) at the conclusion of the interview. In the case of a telephone or video-conference interview, the copy is mailed to the applicant's correspondence address either with or prior to the next official communication. If additional correspondence from the examiner is not likely before an allowance or if other circumstances dictate, the Form should be mailed promptly after the interview rather than with the next official communication.

The Form provides for recordation of the following information:

- Application Number (Series Code and Serial Number)
- Name of applicant
- Name of examiner
- Date of interview
- Type of interview (telephonic, video-conference, or personal)
- Name of participant(s) (applicant, attorney or agent, examiner, other PTO personnel, etc.)
- An indication whether or not an exhibit was shown or a demonstration conducted
- An identification of the specific prior art discussed
- An indication whether an agreement was reached and if so, a description of the general nature of the agreement (may be by attachment of a copy of amendments or claims agreed as being allowable). Note: Agreement as to allowability is tentative and does not restrict further action by the examiner to the contrary.
- The signature of the examiner who conducted the interview (if Form is not an attachment to a signed Office action)

It is desirable that the examiner orally remind the applicant of his or her obligation to record the substance of the interview of each case. It should be noted, however, that the Interview Summary Form will not normally be considered a complete and proper recordation of the interview unless it includes, or is supplemented by the applicant or the examiner to include, all of the applicable items required below concerning the substance of the interview.

A complete and proper recordation of the substance of any interview should include at least the following applicable items:

- 1) A brief description of the nature of any exhibit shown or any demonstration conducted,
- 2) an identification of the claims discussed,
- 3) an identification of the specific prior art discussed,
- 4) an identification of the principal proposed amendments of a substantive nature discussed, unless these are already described on the Interview Summary Form completed by the Examiner,
- 5) a brief identification of the general thrust of the principal arguments presented to the examiner,
 - (The identification of arguments need not be lengthy or elaborate. A verbatim or highly detailed description of the arguments is not required. The identification of the arguments is sufficient if the general nature or thrust of the principal arguments made to the examiner can be understood in the context of the application file. Of course, the applicant may desire to emphasize and fully describe those arguments which he or she feels were or might be persuasive to the examiner.)
- 6) a general indication of any other pertinent matters discussed, and
- 7) if appropriate, the general results or outcome of the interview unless already described in the Interview Summary Form completed by the examiner.

Examiners are expected to carefully review the applicant's record of the substance of an interview. If the record is not complete and accurate, the examiner will give the applicant an extendable one month time period to correct the record.

Examiner to Check for Accuracy

If the claims are allowable for other reasons of record, the examiner should send a letter setting forth the examiner's version of the statement attributed to him or her. If the record is complete and accurate, the examiner should place the indication, "Interview Record OK" on the paper recording the substance of the interview along with the date and the examiner's initials.

In The United States Patent And Trademark Office

Application Number: 12/601,275

Application Filed: 11/22/2009

5 Applicant(s): Venkat Konda
Title: VLSI Layouts of Fully Connected Generalized Networks
Examiner/Art Unit: Vibol Tan / 2819
International Application Number: PCT/US08/64605
International Application Filed: 5/22/2008
10 Priority Date: 5/25/2007

San Jose, 2012 Apr 30, Mon

AMENDMENT A

(and the response to office letter dated 2/7/2012)

15

Mail Stop Amendment

Commissioner for Patents

P.O. Box 1450

Alexandria, Virginia, 22313-1450

20

Dear Sir/Madam:

In response to the office action mailed 2012 February 7, please consider the following Amendment A.

First Applicant addresses the novelty and unobviousness of the current invention over the prior art, including the U.S. Patent 6,940,308 by Wong and the U.S Patent 7,154,887 by Wu et. al. Applicant also submits that he has reviewed all the other cited references and they do not show the current invention or render it obvious.

5

I. IN RESPONSE TO CLAIM REJECTIONS – USC 102(b):

First Applicant submits that the current patent application discloses the VLSI layouts or floor plans of the Benes Network, Butterfly Fat Tree Networks including more

10 generalizations of these networks. The layout is critical to implement these networks in a semiconductor chip which is typically a 2D-plane. As such topologies of the Benes Network and Butterfly Fat Tree networks are known in the prior art and the topologies of the Benes Network and Butterfly Fat Tree networks disclosed in the current invention are the same as in the prior art. However the layouts of these networks are novel.

- 15 In addition the layouts disclosed in the current invention include generalized multi-stage networks $V(N_1, N_2, d, s)$, generalized folded multi-stage networks $V_{fold}(N_1, N_2, d, s)$, generalized butterfly fat tree networks $V_{bft}(N_1, N_2, d, s)$, generalized multi-link multi-stage networks $V_{mlink}(N_1, N_2, d, s)$, generalized folded multi-link multistage networks $V_{fold-mlink}(N_1, N_2, d, s)$, generalized multi-link butterfly fat tree networks
- 20 $V_{mlink-bft}(N_1, N_2, d, s)$, and generalized hypercube networks $V_{hcube}(N_1, N_2, d, s)$ for s = 1,2,3 or any number in general.

AMENDMENT A, Contd.

<u>The following table addresses the list of items where the current</u> <u>invention is different/superior over column-based layout (or floor plan)</u> <u>disclosed in U.S. Patent 6,940,308 by Wong:</u>

Item addressed	Solutions in US patent	Solutions in Current
	6,940,308 by Wong	Application
Layout format	Column-based layout	2D-grid in row-column layout (Key insight in the disclosure is to arrange the sub-integrated circuits in a hypercube topology)
Fundamental sub-	As disclosed in Fig. 13A, the	As disclosed in Figs. 1C –
integrated circuit	 basic building block of the column-based layout is based on (as mentioned in Col. 13, lines 22-24) "There are two logic cells 81 per switch cell 82, connected in the so called "butterfly pattern", consistent with the Benes Network topology:". It is clear that column-based layout has the following properties: 1) Basic building block (or sub-integrated circuit) has cross links connected 	1G, the 2D-grid in row- column layout disclosed in the current patent application has no butterfly pattern in the layout. In other words the cross links between switches of two succeeding stages are always connected between two sub- integrated circuits which are placed either horizontally in the same column or vertically in the

AMENDMENT A, Contd.

To scale the layout for	succeeding stages with in the same sub-integrated circuit. Otherwise there will not be a butterfly pattern. As disclosed in Fig. 13B,	links between two succeeding stages are either vertical links (as shown in Figs. 1D & 1F) or horizontal links (as shown in Figs. 1E & 1G). As disclosed in Fig. 1H,
bigger size of network	columns are doubled by adding	the layout doubles in both
(or sub-integrated	one more stage of switches in	columns and rows
circuits)	both the columns (for example	alternatively, as the size
ch cultoj	Fig. 13A is replicated to get Fig.	of the network is doubled
	13B and a new stage of switches	repeatedly as needed, in
	83 are added and the cross links	hypercube layout format.
	between the new stage of	hypereube layout format.
	switches in the two columns	
	become horizontal links).	
	 It is clear as the network is scaled up, only columns are increasing. And the basic building block is the same as that shown in Fig. 13A. 	
Industry Applicability	The applicant believes column-	Same basic block can be
or practical use of the	based layout is impractical for an	easily used to
layout	FPGA vendor to manufacture	manufacture FPGAs of
(FPGA vendors typically take the same basic	FPGA devices with different size of basic blocks, as the number of	different sizes. (The layout disclosed in the
block (consisting of a	basic blocks in each device determine the size of each	current invention is already adopted by an

	I	
certain set of LUTs and	column (as that shown in FIG.	FPGA vendor namely
registers) and	13A) and each device will	Quicklogic corporation,
manufacture devices of	require a different size of each	located in Sunnyvale,
different sizes of basic	column to manufacture the	CA).
blocks (typically called	devices close to square in size.	
slice, CLB or PLC etc.)	And hence it is impractical for an	
	FPGA vendor to employ	
	column-based layout.	
Locality between two	From Fig. 13A, it can be seen	For most of the logic cells
neighboring basic	that for a signal from a logic cell	in this layout, a signal
blocks either	to be routed horizontally to the	from a logic cell can be
horizontally or	nearest neighboring column, the	routed vertically or
vertically	signal needs to go through all the	horizontally to the nearest
	stages 82 (whose cross links are	neighboring logic cell by
	connected with in the same	passing through one or
	column) before connecting out to	two stages of switches.
	the logic cell in the basic block	
	of the nearest neighboring	
	column.	

AMENDMENT A, Contd.

2. IMPRACTCIAL TREE-BASED LAYOUT DISCLOSED IN PATENT 6,940,308 BY WONG:

1) The tree-based layout disclosed in patent 6,940,308 by Wong is also not practical to implement by an FPGA vendor because:

- 5 a) From Fig. 14C it is evident that the switches corresponding to the higher stages get pulled into the middle of the layout. For example in Fig. 14C, the switches A4 P4 are placed in the center or middle of the layout. As the size of the logic cells increase from 16 to bigger sizes, this problem will get even worse as the highest stage switches get pulled into the middle of the layout. The next highest stage switches get pulled into close
- 10 to the middle of the layout. Similarly the higher the stage of switches they are placed close to the middle of the layout and the lower the stage of switches they are placed close to the logic cells.

b) Evidently the logic cells and the switches corresponding to the all the stages for that logic cell cannot be designed as a repeatable basic block and hence this tree-based
layout is not practical to implement by an FPGA vendor.

II. IN RESPONSE TO CLAIM REJECTIONS – USC 103(a):

1. LAYOUT FOR REARRANGEBALY AND STRICLTY NON-BOCKING NETWORKS FOR ARBITRARY FAN-OUT MULTICAST CONNECTIONS:

- 5 1) US patent 7,154,887 by Wu et. al. disclosed a grooming switch based on a rearrangeably non-blocking clos network for arbitrary fan-out multicast traffic. To make 3-stage clos network which is rearrangeably non-blocking for unicast connections, US patent 7,154,887 by Wu et. al. disclosed a rearrangeably non-blocking clos network with 5 stages for arbitrary fan-out multicast traffic by adding 2 more stages.
- 2) As disclosed in the "CROSS REFERENCE TO RELATED APPLICATIONS" of the current patent application, by the same applicant, the current patent application also discloses VLSI layouts for:

a) Strictly and rearrangeably nonblocking for arbitrary fan-out multicast and unicast for generalized multi-stage networks disclosed in the US Application Serial No.
15 12/530,207 entitled "FULLY CONNECTED GENERALIZED MULTI-STAGE NETWORKS" by the same applicant.

b) Strictly and rearrangeably nonblocking for arbitrary fan-out multicast and unicast for generalized multi-stage networks disclosed in the US Application Serial No. 12/601,273 entitled "FULLY CONNECTED GENERALIZED BUTTERFLY FAT TREE NETWORKS" by the same applicant.

c) Strictly and rearrangeably nonblocking for arbitrary fan-out multicast and unicast for generalized multi-stage networks disclosed in the US Application Serial No.
 12/601,274 entitled "FULLY CONNECTED GENERALIZED MULTI-LINK MULTI-STAGE NETWORKS" by the same applicant.

25 Claims: Claims 1-3, 8, 12, 25, 29, 40, 44, 48-49 of record are amended and substitute new claims as follows.

20

<u>CLAIMS</u>

What is claimed is:

1. (currently amended): An integrated circuit device comprising a plurality of subintegrated circuit blocks and a routing network, and

5

Said each plurality of sub-integrated circuit blocks comprising a plurality of inlet links and a plurality of outlet links; and

Said routing network interconnects any one of said outlet link of one of said subintegrated circuit block to one or more said inlet links of one or more of said subintegrated circuit blocks; and

10 Said routing network comprising of a plurality of stages y, in each said subintegrated circuit block, starting from the lowest stage of 1 to the highest stage of y, where $y \ge 1$; and

Said routing network comprising a plurality of switches of size $d \times d$, where $d \ge 2$, in each said stage and each said switch of size $d \times d$ having d inlet links and d 15 outlet links; and

Said each sub-integrated circuit block comprising a plurality of said switches corresponding to each said stage; and

Said plurality of outlet links of said each sub-integrated circuit block are directly connected to said inlet links of said switches of its corresponding said lowest stage of 1,

20 and said plurality of inlet links of said each sub-integrated circuit block are directly connected from said outlet links of said switches of its corresponding said lowest stage of <u>1</u>; and

Said each sub-integrated circuit block comprising a plurality of forward connecting links connecting from switches in <u>a</u> lower stage to switches in <u>the its</u>

25 immediate succeeding higher stage, and also comprising a plurality of backward connecting links connecting from switches in <u>a</u> higher stage to switches in <u>the its</u> immediate preceding lower stage; and

Said each sub-integrated circuit block comprising a plurality straight links in said forward connecting links from switches in <u>said each</u> lower stage to switches in <u>the its</u> immediate succeeding higher stage and a plurality cross links in said forward connecting links from switches in <u>said each</u> lower stage to switches in <u>the its</u> immediate succeeding

5 higher stage, and further comprising a plurality of straight links in said backward connecting links from switches in <u>said each</u> higher stage to switches in <u>the its</u> immediate preceding lower stage and a plurality of cross links in said backward connecting links from switches in <u>said each</u> higher stage to switches in <u>the its</u> immediate preceding lower stage.

10 <u>said plurality of sub-integrated circuit blocks arranged in a two-dimensional grid</u> of rows and columns, and

said all straight links are connecting from switches in each said sub-integrated circuit block are connecting to switches in the same said sub-integrated circuit block; and said all cross links are connecting as either vertical or horizontal links between switches

15 in two different said sub-integrated circuit blocks which are either placed vertically above or below, or placed horizontally to the left or to the right,

each said plurality of sub-integrated circuit blocks comprising same number of said stages and said switches in each said stage, regardless of the size of said twodimensional grid so that each said plurality of sub-integrated circuit block with its

20 <u>corresponding said stages and said switches in each stage is replicable in both vertical</u> <u>direction or horizontal direction of said two-dimensional grid</u>.

(currently amended): The integrated circuit device of claim 1, wherein said plurality of sub-integrated circuit blocks arranged in a two-dimensional grid
 <u>said two-dimensional grid of said sub-integrated circuit blocks with their</u>
 <u>corresponding said stages and said switches in each stage is scalable by any power of 2,</u>
 and,

for each multiplication of 2 of the size of total said sub-integrated circuit blocks, by adding one more stage of switches and the layout is placed in hypercube format and

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also the cross links between said one more stage of switches are connected in hypercube format.

3. (Original): The integrated circuit device of claim 2, wherein said cross links in succeeding stages are connecting as alternative vertical and horizontal links
 5 between switches in said sub-integrated circuit blocks.

4. (Original): The integrated circuit device of claim 3, wherein said cross links from switches in a stage in one of said sub-integrated circuit blocks are connecting to switches in the succeeding stage in another of said sub-integrated circuit blocks so that said cross links are either vertical links or horizontal and vice versa, and hereinafter such cross links are "shuffle exchange links").

5. (Original): The integrated circuit device of claim 4, wherein said all
horizontal shuffle exchange links between switches in any two corresponding said
succeeding stages are substantially of equal length and said vertical shuffle exchange
links between switches in any two corresponding said succeeding stages are substantially
of equal length in the entire said integrated circuit device.

6. (Original): The integrated circuit device of claim 5, wherein the shortest horizontal shuffle exchange links are connecting at the lowest stage and between switches in two nearest neighboring said sub-integrated circuit blocks, and length of the horizontal shuffle exchange links is doubled in each succeeding stage; and the shortest vertical
20 shuffle exchange links are connecting at the lowest stage and between switches in two nearest neighboring said sub-integrated circuit blocks, and length of the vertical
20 shuffle exchange links are connecting at the lowest stage and between switches in two nearest neighboring said sub-integrated circuit blocks, and length of the vertical shuffle exchange links is doubled in each succeeding stage.

7. (currently amended): The integrated circuit device of claim 6, wherein y ≥ (log₂ N), where N > 1, so that the length of the horizontal shuffle exchange links in
25 the highest stage is equal to half the size of the horizontal size of said two dimensional grid of sub-integrated circuit blocks and the length of the vertical shuffle exchange links

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in the highest stage is equal to half the size of the vertical size of said two dimensional grid of sub-integrated circuit blocks.

8. (Original): The integrated circuit device of claim 7, wherein d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block
5 connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast Benes network with full bandwidth.

9. (Original): The integrated circuit device of claim 7, wherein d = 2 and
10 there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast Benes network and rearrangeably nonblocking for arbitrary fan-out multicast Benes network with full
15 bandwidth.

10. (Original): The integrated circuit device of claim 7, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links
20 and said routing network is strictly nonblocking for arbitrary fan-out multicast Benes network with full bandwidth.

11. (currently amended): The integrated circuit device of claim 6, wherein $y \ge (\log_2 N)$, where $N > 1_x$ so that the length of the horizontal shuffle exchange links in the highest stage is equal to half the size of the horizontal size of said two dimensional grid of sub-integrated circuit blocks and the length of the vertical shuffle exchange links in the highest stage is equal to half the size of the vertical size of said two dimensional grid of sub-integrated circuit blocks, and

said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks.

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12. (Original): The integrated circuit device of claim 11, wherein d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast butterfly fat tree network

with full bandwidth.

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13. (Original): The integrated circuit device of claim 11, wherein d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said
stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast butterfly fat tree network and rearrangeably nonblocking for arbitrary fan-out multicast butterfly fat tree network with

full bandwidth.

14. (Original): The integrated circuit device of claim 11, wherein d = 2 and 15 there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast butterfly fat tree network with full bandwidth.

20 15. (Original): The integrated circuit device of claim 1, wherein said horizontal and vertical links are implemented on two or more metal layers.

16. (Original): The integrated circuit device of claim 1, wherein said switches comprising active and reprogrammable cross points and said each cross point is programmable by an SRAM cell or a Flash Cell.

25 17. (Original): The integrated circuit device of claim 1, wherein said subintegrated circuit blocks are of equal die size.

18. (Original): The integrated circuit device of claim 15, wherein said subintegrated circuit blocks are Lookup Tables (hereinafter "LUTs") and said integrated circuit device is a field programmable gate array (FPGA) device or field programmable gate array (FPGA) block embedded in another integrated circuit device.

5 19. (Original): The integrated circuit device of claim 15, wherein said subintegrated circuit blocks are AND or OR gates and said integrated circuit device is a programmable logic device (PLD).

20. (Original): The integrated circuit device of claim 1, wherein said subintegrated circuit blocks comprising any arbitrary hardware logic or memory circuits.

10 21. (Original): The integrated circuit device of claim 1, wherein said switches comprising active one-time programmable cross points and said integrated circuit device is a mask programmable gate array (MPGA) device or a structured ASIC device.

22. (Original): The integrated circuit device of claim 1, wherein said switches comprising passive cross points or just connection of two links or not and said integrated
 15 circuit device is a Application Specific Integrated Circuit (ASIC) device.

23. (Original): The integrated circuit device of claim 1, wherein said subintegrated circuit blocks further recursively comprise one or more super-sub-integrated circuit blocks and a sub-routing network.

24. (currently amended): The integrated circuit device of claim 4, wherein
20 said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and y ≥ (log₂ N), where N > 1.

25. (Original): The integrated circuit device of claim 24, wherein d = 2 and
there is only one switch in each said stage in each said sub-integrated circuit block
connecting said forward connecting links and there is only one switch in each said stage

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in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized multi-stage network with full bandwidth.

- 26. (Original): The integrated circuit device of claim 24, wherein d = 2 and
 5 there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast generalized multi-stage network and rearrangeably nonblocking for arbitrary fan-out multicast generalized multi-
- 10 stage network with full bandwidth.

27. (Original): The integrated circuit device of claim 24, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast generalized multi-stage network with full bandwidth.

28. (currently amended): The integrated circuit device of claim 4, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and $y \ge (\log_2 N)$, where $N > 1_*$ and

said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks.

29. (Original): The integrated circuit device of claim 28, wherein d = 2 and
25 there is only one switch in each said stage in each said sub-integrated circuit block
connecting said forward connecting links and there is only one switch in each said stage
in each said sub-integrated circuit block connecting said backward connecting links and

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said routing network is rearrangeably nonblocking for unicast generalized butterfly fat tree network with full bandwidth.

30. (Original): The integrated circuit device of claim 28, wherein d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block
5 connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast generalized butterfly fat tree Network and rearrangeably nonblocking for arbitrary fan-out multicast generalized butterfly fat tree butterfly fat tree network with full bandwidth.

31. (Original): The integrated circuit device of claim 28, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast
 generalized butterfly fat tree network with full bandwidth.

32. (Original): The integrated circuit device of claim 1, wherein said straight links connecting from switches in each said sub-integrated circuit block are connecting to switches in the same said sub-integrated circuit block; and

said cross links are connecting as vertical or horizontal or diagonal links betweentwo different said sub-integrated circuit blocks.

33. (Original): The integrated circuit device of claim 7, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and

25 said routing network is rearrangeably nonblocking for unicast multi-link Benes network with full bandwidth.

Application Number: 12/601,275 (Venkat Konda) Art Unit: 2819 AMENDMENT A, Contd.

34. (Original): The integrated circuit device of claim 7, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast multi-link Benes network and rearrangeably nonblocking for arbitrary fan-out multicast multi-link Benes network with full bandwidth.

35. (Original): The integrated circuit device of claim 7, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block
connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast multi-link Benes network with full bandwidth.

36. (Original): The integrated circuit device of claim 11, wherein d = 4 and 15 there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast multi-link butterfly fat tree network with full bandwidth.

20 37. (Original): The integrated circuit device of claim 11, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast multi-link butterfly fat tree

25 network and rearrangeably nonblocking for arbitrary fan-out multicast multi-link butterfly fat tree network with full bandwidth.

38. (Original): The integrated circuit device of claim 11, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said

stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast multi-link butterfly fat tree network with full bandwidth.

- 39. (currently amended): The integrated circuit device of claim 4, wherein
 5 said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and y ≥ (log₂ N), where N > 1.
- 40. (Original): The integrated circuit device of claim 39, wherein d = 4 and 10 there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized multi-link multi-stage network with full bandwidth.
- 41. (Original): The integrated circuit device of claim 39, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast generalized multi-link multistage network and rearrangeably nonblocking for arbitrary fan-out multicast generalized multi-link multi-stage network with full bandwidth.

42. (Original): The integrated circuit device of claim 39, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast

generalized multi-link multi-stage network with full bandwidth.

Application Number: 12/601,275 (Venkat Konda) Art Unit: 2819 AMENDMENT A, Contd.

43. (currently amended): The integrated circuit device of claim 4, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and

5 $y \ge (\log_2 N)$, where N > 1, and

said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks.

44. (Original): The integrated circuit device of claim 43, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block
connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized multi-link butterfly fat tree network with full bandwidth.

45. (Original): The integrated circuit device of claim 43, wherein d = 4 and
15 there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast generalized multi-link butterfly fat tree Network and rearrangeably nonblocking for arbitrary fan-out multicast
20 generalized multi-link butterfly fat tree network with full bandwidth.

46. (Original): The integrated circuit device of claim 43, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast generalized multi-link butterfly fat tree network with full bandwidth.

47. (currently amended): The integrated circuit device of claim 1, wherein said plurality of forward connecting links use a plurality of buffers to amplify signals

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driven through them and said plurality of backward connecting links use a plurality of buffers to amplify signals driven through them; and said buffers can be inverting or noninverting buffers.

48. (currently amended): The integrated circuit device of claim 1, wherein
5 said wherein said all switches of size d×d are either fully populated or partially populated.

The rejection of Claims 1-8, 12, 16-25, 29, 33, 40, 44, and 48-49 under 35 USC 102(b)

Accordingly applicant submit that the claims do comply with § 102(b) and therefore request withdrawal of this rejection.

2) The rejection of Claims 9-11, 13-15, 26-28, 30-32, 34-39, 41-43, and 45-47 under 35 USC 103(a)

Accordingly applicant submit that the claims do comply with § 103(a) and therefore request withdrawal of this rejection.

CONCLUSION

For all of the above reasons, applicant submits that the Claims are now in proper form, and that the Claims all define patentably over the prior art. Therefore applicant submits that this application is now in condition for allowance, which action he respectfully solicits.

Conditional request for Constructive Assistance

Applicant has amended the claims of this application so that they are proper, definite, and define novel structure which is also unobvious. If, for any reason this application is not believed to be in full condition for allowance, applicant respectfully request the constructive assistance and suggestions of the Examiner pursuant to M.P.E.P §

20 2173.02 and § 707.07(j) in order that the undersigned can place this application in allowable condition as soon as possible and without the need for further proceedings.

Very respectfully,

/Venkat Konda/

25 Venkat Konda

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Application Number: 12/601,275 (Venkat Konda) AMENDMENT A, Contd.

Konda Technologies, Inc (USPTO Customer Number: 38139)

6278 Grand Oak Way

San Jose, CA 95135

Phone: 408-472-3273

5 Fax: 408-238-2478

Electronic A	Electronic Acknowledgement Receipt						
EFS ID:	12667779						
Application Number:	12601275						
International Application Number:							
Confirmation Number:	6372						
Title of Invention:	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS						
First Named Inventor/Applicant Name:	Venkat Konda						
Customer Number:	38139						
Filer:	Venkar Konda						
Filer Authorized By:							
Attorney Docket Number:	V-0045US						
Receipt Date:	30-APR-2012						
Filing Date:	31-MAY-2010						
Time Stamp:	18:15:10						
Application Type:	U.S. National Stage under 35 USC 371						

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Submitted wi	th Payment	no					
File Listin	g:						
Document Number	Document Description		File Name	File Size(Bytes)/ Message Digest	Multi Part /.zip	Pages (if appl.)	
1	Amendment/Req. Reconsideration-After 1 Non-Final Reject		mendmentA-V0045US.pdf	188498	no	21	
			anenamenta voo4505.pai	1ef528eeaf46ce4d077d8e28f40c4a871729 bff2	110	21	
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If a new application is being filed and the application includes the necessary components for a filing date (see 37 CFR 1.53(b)-(d) and MPEP 506), a Filing Receipt (37 CFR 1.54) will be issued in due course and the date shown on this Acknowledgement Receipt will establish the filing date of the application.

National Stage of an International Application under 35 U.S.C. 371

If a timely submission to enter the national stage of an international application is compliant with the conditions of 35 U.S.C. 371 and other applicable requirements a Form PCT/DO/EO/903 indicating acceptance of the application as a national stage submission under 35 U.S.C. 371 will be issued in addition to the Filing Receipt, in due course.

New International Application Filed with the USPTO as a Receiving Office

If a new international application is being filed and the international application includes the necessary components for an international filing date (see PCT Article 11 and MPEP 1810), a Notification of the International Application Number and of the International Filing Date (Form PCT/RO/105) will be issued in due course, subject to prescriptions concerning national security, and the date shown on this Acknowledgement Receipt will establish the international filing date of the application.

PTO/SB/06 (07-06)

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						nd to	a collection of	of information unle			OMB control number.
P	PATENT APPLICATION FEE DETERMINATION RECORD Substitute for Form PTO-875					A		Docket Number 1,275		ing Date 31/2010	To be Mailed
APPLICATION AS FILED – PART I (Column 1) (Column 2)							SMALL	entity 🛛	OR		HER THAN
	FOR	N	JMBER FIL	.ED NU	MBER EXTRA		RATE (\$)	FEE (\$)		RATE (\$)	FEE (\$)
	BASIC FEE (37 CFR 1.16(a), (b), o	or (c))	N/A		N/A		N/A			N/A	
	SEARCH FEE (37 CFR 1.16(k), (i), c	or (m))	N/A		N/A		N/A			N/A	
	EXAMINATION FE (37 CFR 1.16(o), (p), o		N/A		N/A		N/A			N/A	
	TAL CLAIMS CFR 1.16(i))		min	us 20 = *			X \$ =		OR	X \$ =	
	EPENDENT CLAIM CFR 1.16(h))	S	mi	nus 3 = *			X \$ =			X \$ =	
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* If t	he difference in colu	umn 1 is less than	zero, ente	r "0" in column 2.			TOTAL			TOTAL	
	APPI	LICATION AS	AMEND	DED – PART II						OTU	
		(Column 1)		(Column 2)	(Column 3)		SMAL	L ENTITY	OR		ER THAN ALL ENTITY
		CLAIMS		HIGHEST		1				0	
AMENDMENT	04/30/2012	REMAINING AFTER AMENDMENT		NUMBER PREVIOUSLY PAID FOR	PRESENT EXTRA		RATE (\$)	ADDITIONAL FEE (\$)		RATE (\$)	ADDITIONAL FEE (\$)
OME	Total (37 CFR 1.16(i))	* 48	Minus	** 49	= 0		X \$30 =	0	OR	X \$ =	
IN I	Independent (37 CFR 1.16(h))	* 1	Minus	***3	= 0		X \$125 =	0	OR	X \$ =	
AME	Application Si	ze Fee (37 CFR 1	.16(s))								
	FIRST PRESEN	TATION OF MULTIF	LE DEPEN	DENT CLAIM (37 CFI	R 1.16(j))				OR		
							TOTAL ADD'L FEE	0	OR	TOTAL ADD'L FEE	
		(Column 1)		(Column 2)	(Column 3)				-		
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AM		TATION OF MULTIF	LE DEPEN	DENT CLAIM (37 CFI	R 1.16(j))				OR		
							TOTAL ADD'L FEE		OR	TOTAL ADD'L FEE	
** lf ***	 * If the entry in column 1 is less than the entry in column 2, write "0" in column 3. ** If the "Highest Number Previously Paid For" IN THIS SPACE is less than 20, enter "20". *** If the "Highest Number Previously Paid For" IN THIS SPACE is less than 3, enter "3". The "Highest Number Previously Paid For" (Total or Independent) is the highest number found in the appropriate box in column 1. 										
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Application Da	ta Shoot 37 CED 1 76	Attorney Docket Number	V-0045US				
	Application Data Sheet 37 CFR 1.76						
Title of Invention	VLSI Layouts of Fully Connec	ted Generalized Networks					
bibliographic data arrar	The application data sheet is part of the provisional or nonprovisional application for which it is being submitted. The following form contains the bibliographic data arranged in a format specified by the United States Patent and Trademark Office as outlined in 37 CFR 1.76. This document may be completed electronically and submitted to the Office in electronic format using the Electronic Filing System (EFS) or the						

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Portions or all of the application associated with this Application Data Sheet may fall under a Secrecy Order pursuant to 37 CFR 5.2 (Paper filers only. Applications that fall under Secrecy Order may not be filed electronically.)

Applicant Information:

Applic	Applicant 1 Remove											
Applic	cant A	Authority 🖲	Inventor	OLega	al Representativ	Il Representative under 35 U.S.C. 117				OParty of Interest under 35 U.S.C. 118		
Prefix	Giv	en Name			Middle Na	me			Farr	nily Name		Suffix
	Ven	kat							Kono	da		
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Correspondence Information:

Enter either Customer Number or complete the Correspondence Information section below. For further information see 37 CFR 1.33(a).								
An Address is being provided for the correspondence Information of this application.								
Customer Number 38139								
Email Address venkat@kondatech.com Add Email Remove Email								

Application Information:

Title of the Invention	VLSI Layouts of Fu	/LSI Layouts of Fully Connected Generalized Networks					
Attorney Docket Number	V-0045US	-0045US Small Entity Status Claimed X					
Application Type	Nonprovisional	onprovisional					
Subject Matter	Utility	Utility					
Suggested Class (if any)			Sub Class (if any)				
Suggested Technology C	enter (if any)						
Total Number of Drawing	Sheets (if any) 39 Suggested Figure for Publication (if any) 1G						

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Application Da	ta Shoot 37 CEP 1 76	Attorney Docket Number	V-0045US
Application Data Sheet 37 CFR 1.76		Application Number	
Title of Invention	VLSI Layouts of Fully Connec	ted Generalized Networks	

Publication Information:

Request Early Publication (Fee required at time of Request 37 CFR 1.219)
 Request Not to Publish. I hereby request that the attached application not be published under 35 U.S.
 C. 122(b) and certify that the invention disclosed in the attached application has not and will not be the subject of an application filed in another country, or under a multilateral international agreement, that requires publication at eighteen months after filing.

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Repres	entative	information	should b	e provi	ided for all	practi	itioners having a	power of	of attorney	in the	appli	cation.	Providing
this info	this information in the Application Data Sheet does not constitute a power of attorney in the application (see 37 CFR 1.32).												
Enter	either	Customer	Number	or	complete	the	Representative	Name	section	below.	lf	both	sections
are con	Enter either Customer Number or complete the Representative Name section below. If both sections are completed the Customer Number will be used for the Representative Information during processing.												

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Customer Number	38139		

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This section allows for the applicant to either claim benefit under 35 U.S.C. 119(e), 120, 121, or 365(c) or indicate National Stage entry from a PCT application. Providing this information in the application data sheet constitutes the specific reference required by 35 U.S.C. 119(e) or 120, and 37 CFR 1.78(a)(2) or CFR 1.78(a)(4), and need not otherwise be made part of the specification.

Prior Application Status	Expired		Remove				
Application Number	Continuity Type	Prior Application Number	Filing Date (YYYY-MM-DD)				
Additional Domestic Benefit/National Stage Data may be generated within this form by selecting the Add button.							

Foreign Priority Information:

This section allows for the applicant to claim benefit of foreign priority and to identify any prior foreign application for which priority is not claimed. Providing this information in the application data sheet constitutes the claim for priority as required by 35 U.S.C. 119(b) and 37 CFR 1.55(a).							
Remove							
Application Number	Country ⁱ	Parent Filing Date (YYYY-MM-DD)	Priority Claimed				
PCT/US08/64605 US 2008-05-22 • Yes No							
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Assignee Information:

Providing this information in the application data sheet does not substitute for compliance with any requirement of part 3 of Title 37 of the CFR to have an assignment recorded in the Office.

Assignee 1

Remove

PTO/SB/14 (11-08) Approved for use through 01/31/2014. OMB 0651-0032 U.S. Patent and Trademark Office; U.S. DEPARTMENT OF COMMERCE

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Application Data Sheet 37 CFR 1.76		Attorney Doo	ket Number	V-0045	US	
		Application Number				
Title of Invention	VLSI L	ayouts of Fully Connected Generalized Networks				
If the Assignee is an Organization check here.						
Organization Name Konda Technologies Inc.						
Mailing Address I	nforma	tion:				
Address 1 6278 Grand Oak Wa		6278 Grand Oak Way	y			
Address 2						
City San Jos		San Jose	e		nce	СА
Country i US			Postal Code		95135	
Phone Number 408-472		408-472-3273		Fax Number		408-238-2478
Email Address venkat@kondatech.c		om				
Additional Assignee Data may be generated within this form by selecting the Add Add button.						

Signature:

A signature of the applicant or representative is required in accordance with 37 CFR 1.33 and 10.18. Please see 37 CFR 1.4(d) for the form of the signature.					
Signature	/Venkat Konda/			Date (YYYY-MM-DD)	2012-04-18
First Name	Venkat	Last Name	Konda	Registration Number	

This collection of information is required by 37 CFR 1.76. The information is required to obtain or retain a benefit by the public which is to file (and by the USPTO to process) an application. Confidentiality is governed by 35 U.S.C. 122 and 37 CFR 1.14. This collection is estimated to take 23 minutes to complete, including gathering, preparing, and submitting the completed application data sheet form to the USPTO. Time will vary depending upon the individual case. Any comments on the amount of time you require to complete this form and/or suggestions for reducing this burden, should be sent to the Chief Information Officer, U.S. Patent and Trademark Office, U.S. Department of Commerce, P.O. Box 1450, Alexandria, VA 22313-1450. DO NOT SEND FEES OR COMPLETED FORMS TO THIS ADDRESS. SEND TO: Commissioner for Patents, P.O. Box 1450, Alexandria, VA 22313-1450.

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Electronic Acknowledgement Receipt				
EFS ID:	12576502			
Application Number:	12601275			
International Application Number:				
Confirmation Number:	6372			
Title of Invention:	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS			
First Named Inventor/Applicant Name:	Venkat Konda			
Customer Number:	38139			
Filer:	Venkar Konda			
Filer Authorized By:				
Attorney Docket Number:	V-0045US			
Receipt Date:	19-APR-2012			
Filing Date:	31-MAY-2010			
Time Stamp:	04:53:51			
Application Type:	U.S. National Stage under 35 USC 371			

Payment information:

Submitted with I	Payment	no	no			
File Listing:						
Document Number	Document Description	File Name	File Size(Bytes)/ Message Digest	Multi Part /.zip	Pages (if appl.)	
1	Application Data Sheet	sb0014-V0045US.pdf	1420688	no	4	
		550014 V004505.par	1b1fcb4236336b2fb79c4394e98f66b92e83 9166	10		
Warnings:			· · ·			
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This Acknowledgement Receipt evidences receipt on the noted date by the USPTO of the indicated documents, characterized by the applicant, and including page counts, where applicable. It serves as evidence of receipt similar to a Post Card, as described in MPEP 503.

New Applications Under 35 U.S.C. 111

If a new application is being filed and the application includes the necessary components for a filing date (see 37 CFR 1.53(b)-(d) and MPEP 506), a Filing Receipt (37 CFR 1.54) will be issued in due course and the date shown on this Acknowledgement Receipt will establish the filing date of the application.

National Stage of an International Application under 35 U.S.C. 371

If a timely submission to enter the national stage of an international application is compliant with the conditions of 35 U.S.C. 371 and other applicable requirements a Form PCT/DO/EO/903 indicating acceptance of the application as a national stage submission under 35 U.S.C. 371 will be issued in addition to the Filing Receipt, in due course.

New International Application Filed with the USPTO as a Receiving Office

If a new international application is being filed and the international application includes the necessary components for an international filing date (see PCT Article 11 and MPEP 1810), a Notification of the International Application Number and of the International Filing Date (Form PCT/RO/105) will be issued in due course, subject to prescriptions concerning national security, and the date shown on this Acknowledgement Receipt will establish the international filing date of the application.

	ed States Patent a	ND TRADEMARK OFFICE	UNITED STATES DEPAR United States Patent and Address: COMMISSIONER I P.O. Box 1450 Alexandria, Virginia 22 www.uspto.gov	FOR PATENTS		
APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.		
12/601,275	05/31/2010	Venkat Konda	V-0045US	6372		
³⁸¹³⁹ Konda Technol	7590 02/07/2012	EXAMINER				
6278 GRAND OAK WAY			TAN, VIBOL			
SAN JOSE, CA 95135		ART UNIT	PAPER NUMBER			
			2819			
			MAIL DATE	DELIVERY MODE		
			02/07/2012	PAPER		

Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

	Application No.	Applicant(s)		
	12/601,275	KONDA, VENKAT		
Office Action Summary	Examiner	Art Unit		
	Vibol Tan	2819		
The MAILING DATE of this communication Period for Reply	appears on the cover sheet wit	h the correspondence address		
 A SHORTENED STATUTORY PERIOD FOR REWHICHEVER IS LONGER, FROM THE MAILING Extensions of time may be available under the provisions of 37 CFI after SIX (6) MONTHS from the mailing date of this communication If NO period for reply is specified above, the maximum statutory pe Failure to reply within the set or extended period for reply will, by st Any reply received by the Office later than three months after the meanned patent term adjustment. See 37 CFR 1.704(b). 	DATE OF THIS COMMUNIC 1.136(a). In no event, however, may a re riod will apply and will expire SIX (6) MONT atute, cause the application to become AB/	ATION. ply be timely filed THS from the mailing date of this communication. ANDONED (35 U.S.C. § 133).		
Status				
1)∑ Responsive to communication(s) filed on <u>2</u>	2 November 2009.			
	This action is non-final.			
3) An election was made by the applicant in re		ement set forth during the interview on		
; the restriction requirement and elec	tion have been incorporated ir	nto this action.		
4) Since this application is in condition for allo	wance except for formal matte	ers, prosecution as to the merits is		
closed in accordance with the practice und	er <i>Ex parte Quayle</i> , 1935 C.D.	11, 453 O.G. 213.		
Disposition of Claims				
5) Claim(s) <u>1-49</u> is/are pending in the applicat	tion.			
5a) Of the above claim(s) is/are with				
6) Claim(s) is/are allowed.				
7)⊠ Claim(s) <u>1-49</u> is/are rejected.				
8) Claim(s) is/are objected to.				
9) Claim(s) are subject to restriction ar	d/or election requirement.			
Application Papers				
10) The specification is objected to by the Exan	niner.			
11) The drawing(s) filed on is/are: a)	accepted or b) 🗌 objected to b	by the Examiner.		
Applicant may not request that any objection to	the drawing(s) be held in abeyand	ce. See 37 CFR 1.85(a).		
Replacement drawing sheet(s) including the cor	rection is required if the drawing(s) is objected to. See 37 CFR 1.121(d).		
12) The oath or declaration is objected to by the	Examiner. Note the attached	Office Action or form PTO-152.		
Priority under 35 U.S.C. § 119				
13) Acknowledgment is made of a claim for fore	eign priority under 35 U.S.C. §	119(a)-(d) or (f).		
a) All b) Some * c) None of:				
1. Certified copies of the priority documents have been received.				
 2. Certified copies of the priority documents have been received in Application No 3. Copies of the certified copies of the priority documents have been received in this National Stage 				
application from the International Bu	•	eceived in this National Stage		
* See the attached detailed Office action for a		received		
Attachment/a)				
Attachment(s) 1) X Notice of References Cited (PTO-892)		ummary (PTO-413)		
 2) Notice of Draftsperson's Patent Drawing Review (PTO-948) 	Paper No(s)/Mail Date		
3) Information Disclosure Statement(s) (PTO/SB/08)	5) 📃 Notice of In 6) 🗌 Other:	formal Patent Application		
Paper No(s)/Mail Date U.S. Patent and Trademark Office		_·		
	e Action Summary	Part of Paper No./Mail Date 20120201		

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DETAILED ACTION

Claim Rejections - 35 USC § 102

1. The following is a quotation of the appropriate paragraphs of 35 U.S.C. 102 that

form the basis for the rejections under this section made in this Office action:

A person shall be entitled to a patent unless -

(b) the invention was patented or described in a printed publication in this or a foreign country or in public use or on sale in this country, more than one year prior to the date of application for patent in the United States.

2. Claims 1-8, 12, 16-25, 29, 33, 40, 44, and 48-49 are rejected under 35

U.S.C. 102(b) as being anticipated by Wong (U.S. PAT. 6,940,308).

In claim 1, Wong teaches an integrated circuit device comprising a plurality of

sub-integrated circuit blocks and a routing network (ASIC; col. 13, lines 4-5), and

Said each plurality of sub-integrated circuit blocks comprising a plurality of inlet

links and a plurality of outlet links (input-output pin-pair; Col. 13, Line 25); and

Said routing network interconnects any one of said outlet link of one of said sub-

integrated circuit block to one or more said inlet links of one or more of said sub-

integrated circuit blocks (peripheral blocks of an integrated circuit; col. 14, line 48); and

Said routing network comprising of a plurality of stages y, starting from the lowest stage to the highest stage (hierarchical levels; col. 2, lines 7-12); and

Said routing network comprising a plurality of switches of size d x d, where d \ge 2, in each said stage and each said switch of size d x d having d inlet links and d outlet links (2x2 switch; col. 2, line 27); and

Said each sub-integrated circuit block comprising a plurality of said switches corresponding to each said stage (switches at first rank of hierarchy; col. 2, lines 34-35; Fig. 3C); and

Said each sub-integrated circuit block comprising a plurality of forward connecting links connecting from switches in lower stage to switches in the immediate succeeding higher stage (col. 2; lines 16-22; see input and output links in Fig. 2B; Fig. 3A), and also comprising a plurality of backward connecting links connecting from switches in higher stage to switches in the immediate preceding lower stage (see upper network and lower network in Fig 3B; see upper and lower in Fig. 3C; inputs and outputs, Fig. 4B); and

Said each sub-integrated circuit block comprising a plurality straight links in said forward connecting links from switches in lower stage to switches in the immediate succeeding higher stage and a plurality cross links in said forward connecting links from switches in lower stage to switches in the immediate succeeding higher stage (pass (straight) mode or cross mode; col. 5, lines 4-13; cross, Fig. 2B; see upper network and lower network in Fig. 3B; see upper and lower in Fig. 3C; inputs and outputs Fig. 4B), and further comprising a plurality of straight links in said backward connecting links from switches in higher stage to switches in the immediate preceding lower stage and a plurality of cross links in said backward connecting links from switches in the immediate preceding lower stage to switches in the immediate preceding lower stage and a plurality of cross links in said backward connecting links from switches in the immediate preceding lower stage to switches in the immediate preceding lower stage to switches in the immediate preceding lower stage to switches in higher stage to switches in the immediate preceding lower stage and a plurality of cross links in said backward connecting links from switches in higher stage to switches in the immediate preceding lower stage (pass (straight) mode or cross mode; col. 5, lines 4-13; cross, Fig. 2B; see upper network and lower network in Fig 3B; see upper and lower in Fig. 3C; inputs and outputs, Fig. 4B).

In claim 2, Wong further teaches the integrated circuit device of claim 1, wherein said all straight links are connecting from switches in each said sub-integrated circuit block are connecting to switches in the same said sub-integrated circuit block (Fig.

13A); and

said all cross links are connecting as either vertical or horizontal links between switches in two different said sub-integrated circuit blocks (Fig. 13A).

In claim 3, Wong further teaches the integrated circuit device of claim 2, wherein said plurality of sub-integrated circuit blocks arranged in a two-dimensional grid (width and height dimensions; col. 12, lines 42-51).

In claim 4, Wong further teaches the integrated circuit device of claim 3, wherein said cross links in succeeding stages are connecting as alternative vertical and horizontal links between switches in said sub-integrated circuit blocks (col. 2, lines 7-12; Fig. 10, Fig. 14B).

In claim 5, Wong further teaches the integrated circuit device of claim 4, wherein said cross links from switches in a stage in one of said sub-integrated circuit blocks are connecting to switches in the succeeding stage in another of said sub-integrated circuit blocks so that said cross links are either vertical links or horizontal and vice versa, and hereinafter such cross links are "shuffle exchange links" (rearrangeable interconnection network; col. 1, line 61 - col. 2, Line 6; rearrangeable; col. 4, lines 12-17).

In claim 6, Wong further teaches the integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are substantially of equal length and said vertical shuffle

exchange links between switches in any two corresponding said succeeding stages are substantially of equal length in the entire said integrated circuit device (original Benes network and every route must travel through all levels, col. 7, lines 26-27; col. 11, lines 23-25).

In claim 7, Wong further teaches the integrated circuit device of claim 6, wherein the shortest horizontal shuffle exchange links are connecting at the lowest stage and between switches in two nearest neighboring said sub-integrated circuit blocks, and length of the horizontal shuffle exchange links is doubled in each succeeding stage (shorter routes; col. 7, lines 26-27; col. 11, lines 23-25); and

the shortest vertical shuffle exchange links are connecting at the lowest stage and between switches in two nearest neighboring said sub-integrated circuit blocks, and length of the vertical shuffle exchange links is doubled in each succeeding stage (length = $2^{*}(\log 2 N)$; col. 11, lines 23-25).

In claim 8, Wong further teaches the integrated circuit device of claim 7, wherein $y \ge (\log 2 N)$ so that the length of the horizontal shuffle exchange links in the highest stage is equal to half the size of the horizontal size of said two dimensional grid of sub-integrated circuit blocks and the length of the vertical shuffle exchange links in the highest stage is equal to half the size of the vertical size of said two dimensional grid of sub-integrated circuit blocks (length = 2*(log2 N); col. 11, lines 23-25).

In claim 12, Wong further teaches the integrated circuit device of claim 7, wherein $y \ge (\log 2 N)$ so that the length of the horizontal shuffle exchange links in the highest stage is equal to half the size of the horizontal size of said two dimensional grid

of sub-integrated circuit blocks and the length of the vertical shuffle exchange links in the highest stage is equal to half the size of the vertical size of said two dimensional grid of sub-integrated circuit blocks (length = $2^*(\log 2 N)$; col. 11, lines 23-25), and said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks (see u-turn links; Figs. 5 and 6B).

In claim 16, Wong further teaches the integrated circuit device of claim 1, wherein said horizontal and vertical links are implemented on two or more metal layers (CMOS integrated circuit; col. 5, lines 22-25).

In claim 17, Wong further the integrated circuit device of claim 1, wherein said switches comprising active and reprogrammable cross points and said each cross point is programmable by an SRAM cell or a Flash Cell (SRAM based FPGA; col. 2, lines 25-26; FPGA, col. 1, lines 14-17).

In claim 18, Wong further teaches the integrated circuit device of claim 1, wherein said sub-integrated circuit blocks are of equal die size (sub-network of equal size; col. 6, lines 10-11).

In claim 19, Wong further teaches the integrated circuit device of claim 16, wherein said sub-integrated circuit blocks are Lookup Tables (hereinafter "LUTs") and said integrated circuit device is a field programmable gate array (FPGA) device or field programmable gate array (FPGA) block embedded in another integrated circuit device (look-up tables; col. 8, line 51; FPGA (Field Programmable Gate Array); col. 1, lines 14-17).

In claim 20, Wong further teaches the integrated circuit device of claim 16, wherein said sub-integrated circuit blocks are AND or OR gates and said integrated circuit device is a programmable logic device (PLD) ('logic gates'; col. 7, line 36; 'programmable gate arrays'; col. 14, lines 29-31).

In claim 21, Wong further teaches the integrated circuit device of claim 1, wherein said sub-integrated circuit blocks comprising any arbitrary hardware logic or memory circuits (col. 14, lines 44-51).

In claim 22, Wong further teaches the integrated circuit device of claim 1, wherein said switches comprising active one-time programmable cross points and said integrated circuit device is a mask programmable gate array (MPGA) device or a structured ASIC device (MPGA, col. 14, lines 29-31; ASIC, col. 13, lines 4-5).

In claim 23, Wong further teaches the integrated circuit device of claim 1, wherein said switches comprising passive cross points or just connection of two links or not and said integrated circuit device is a Application Specific Integrated Circuit (ASIC) device (logic cells, col. 1, lines 29-31; ASIC, col. 13, lines 4-5).

In claim 24, Wong the integrated circuit device of claim 1, wherein said subintegrated circuit blocks further recursively comprise one or more super-sub-integrated circuit blocks and a sub-routing network (one or more sub-networks; Col. 6, lines 7-9).

In claim 25, Wong further teaches the integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle

exchange links between switches in any two corresponding said succeeding stages are of different length and $y \ge (\log 2 N)$ (length = 2*(log2 N); col. 11, lines 23-25).

In claim 29, Wong further teaches the integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and $y \ge (\log 2 N)$ (length = 2*(log2 N); col. 11, lines 23-25), and said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks (see u-turn links; Figs. 5 and 6B).

In claim 33, Wong further teaches the integrated circuit device of claim 1, wherein said straight links connecting from switches in each said sub-integrated circuit block are connecting to switches in the same said sub-integrated circuit block; and said cross links are connecting as vertical or horizontal or diagonal links between two different said sub-integrated circuit blocks (see link arrangement in Fig. 3B).

In claim 40, Wong teaches the integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and $y \ge (\log 2 N)$ (length = 2*(log2 N); col. 11, lines 23-25).

In claim 44, Wong further teaches the integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two

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corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and $y \ge (\log 2 N)$ (length = 2*(log2 N); col. 11, lines 23-25), and said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks (see u-turn links; Figs. 5 and 6B).

In claim 48, Wong further teaches the integrated circuit device of claim 1, wherein said plurality of forward connecting links use a plurality of buffers to amplify signals driven through them and said plurality of backward connecting links use a plurality of buffers to amplify signals driven through them; and said buffers can be inverting or non-inverting buffers (buffers; col. 10, lines 59-64).

In claim 49, Wong further teaches the integrated circuit device of claim 1, wherein said wherein said all switches of size d x d are either fully populated or partially populated (populated and depopulated cells; col. 3, lines 34-41).

Claim Rejections - 35 USC § 103

3. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negatived by the manner in which the invention was made.

4. Claims 9-11, 13-15, 26-28, 30-32, 34-39, 41-43, and 45-47 are rejected under 35 U.S.C. 103(a) as being unpatentable over Wong in view of Wu et al. (U. S. PAT. 7,154,887), hereinafter Wu.

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In claim 9, Wong discloses a Benes network wherein d =2 and there is only one switch in each said stags in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links (col. 2, lines 27-30; 2x2 Benes network; col. 5, lines 26-31); but does not disclose said routing network is rearrangeably nonblocking for unicast Benes network with full bandwidth. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being rearrangeably nonblocking for unicast Benes network with full bandwidth (rearrangeably nonblocking for unicast traffic; col. 2, lines 52-54).

Therefore, it would have been obvious to one ordinarily skilled in the art at the time the invention was made to supplement the teachings of Wong and have routing network being rearrangeably nonblocking for unicast Benes network with full bandwidth, in order to satisfy specific routing requirements.

In claim 10, Wong discloses d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links (8x8 Benes network (each switch has two logic cells that are 2x2, hence 8x8 is 4 switches); col. 2, lines 31-33), but does not disclose said routing network is strictly nonblocking for unicast Benes network and rearrangeably nonblocking for arbitrary fan-out multicast Benes network with full bandwidth. However, in the same field of endeavor, Wu discloses a rearrangeable non-

blocking switch (Abstract) and a routing network being strictly nonblocking for unicast Benes network (col. 2, lines 9-11; see Fig. 9) and rearrangeably nonblocking for arbitrary fan-out multicast Benes network with full bandwidth (col. 4, lines 6-7; see fanouts, Fig. 7B).

Therefore, it would have been obvious to one ordinarily skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be strictly nonblocking for unicast Benes and rearrangeably nonblocking for arbitrary fan-out multicast Benes network with full bandwidth, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 11, Wong discloses d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links ('8x8 Benes network'; col. 2, lines 31-33), but does not disclose said routing network is strictly nonblocking for arbitrary fan-out multicast Benes network with full bandwidth. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network strictly nonblocking for arbitrary fan-out multicast Benes network fan-out multicast Benes network with full bandwidth. (col. 2, lines 8-11; see fanouts, Fig. 7B).

Therefore it would have been obvious to one ordinarily skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be strictly nonblocking for arbitrary fan-out multicast Benes network with full bandwidth, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 13, Wong discloses d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub integrated circuit block connecting said backward connecting links ('8x8 Benes network'; col. 2, lines 31-33) and a muting network being a butterfly fat tree network 7, but does not disclose said routing network is rearrangeably nonblocking for unicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being rearrangeably nonblocking for unicast traffic ('rearrangeably nonblocking for unicast traffic'; col. 2, lines 52-54).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a muting network be rearrangeably nonblocking for unicast traffic, as disclosed by Wu, in order to satisfy specific muting requirements.

In claim 14, Wong discloses wherein d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said subintegrated circuit block connecting said backward connecting links ('8x8 Benes network'; col. 2, lines 31-33) and a routing network being a butterfly fat tree network (col. 13, lines 23-24), but does not disclose said routing network is strictly nonblocking for unicast traffic and rearrangeably nonblocking for arbitrary fan-out multicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being strictly nonblocking for unicast traffic (col. 2, lines

9-11, see Fig. 9) and rearrangeably nonblocking for arbitrary fan-out multicast Benes network with full bandwidth (col. 4, lines 6-7; see 'fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinarily skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be strictly nonblocking for unicast Benes and rearrangeably nonblocking for arbitrary fan-out multicast Benes network with full bandwidth, as disclosed by Wu, in order to satisfy specific muting requirements.

In claim 15, Wong discloses wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links ('8x8 Benes network'; col. 2, Lines 31-33) and a routing network being a butterfly fat tree network with full-bandwidth (col. 13, lines 23-24), but does not disclose said routing network is strictly nonblocking for arbitrary fan-out multicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) wherein a routing network is strictly nonblocking for arbitrary fan-out multicast fan-out (col. 2, lines 8-11; see "fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinarily skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be strictly nonblocking for arbitrary fan-out multicast Benes network with full bandwidth, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 26, Wong discloses d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub integrated circuit block connecting said backward connecting links and said routing network is generalized with full bandwidth (col. 2, lines 27-30; '2x2 Benes network'; col. 5, lines 26-31; "networks can be generalized'; col. 6, lines 7-9), but does net disclose said routing network is rearrangeably nonblocking for unicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being rearrangeably nonblocking for unicast traffic ('rearrangeably nonblocking for unicast traffic ('rearrange

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be rearrangeably nonblocking for unicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 27, Wong discloses d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links ('8x8 Benes network' (each switch has two logic cells that are 2x2, hence 8x8 is 4 switches); col. 2, lines 31-33), and said routing network being a generalized network ('networks can be generalized'; col. 6, lines 7-9), but does not disclose said routing network is strictly nonblocking for unicast generalized multi-stage network and rearrangeably nonblocking for arbitrary fan-

out multicast generalized multi-stage network with full bandwidth. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being rearrangeably nonblocking for arbitrary fan-out multicast multi-stage network with full bandwidth (col. 4, lines 6-7; see 'fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be rearrangeably nonblocking for arbitrary fan-out multicast generalized multistage network with full bandwidth, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 28, Wong discloses d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is a generalized routing network ('8x8 Benes network' (each switch has two logic cells that are 2x2, hence 8x8 is 4 switches); col. 2, lines 31-33; 'networks can be generalized'; col. 6, lines 7-9), but does not disclose said routing network is strictly nonblocking for arbitrary fan-out multicast multistage. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network strictly nonblocking for arbitrary fan-out multicast multicast multistage (col. 2, lines 8-11; see 'fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing

network be strictly nonblocking for arbitrary fan-out multicast multistage network with full bandwidth, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 30, Wong discloses d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub integrated circuit block connecting said backward connecting links and said routing network is a generalized butterfly fat tree network with full bandwidth (col. 2, lines 27-30; '2x2 Benes network'; col. 5, lines 26-31; 'networks can be generalized'; col. 6, lines 7-9; 'butterfly pattern'; col. 13, lines 23-24), but does not disclose said routing network is rearrangeably nonblocking for unicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being rearrangeably nonblocking for unicast traffic ('rearrangeably nonblocking for unicast traffic'; col. 2, lines 52-54).

Therefore it would have been obvious to one ordinarily skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be rearrangeably nonblocking for unicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 31, Wong discloses d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network being a generalized butterfly fat tree network with full-bandwidth ('8x8 Benes network' (each

switch has two logic cells that are 2x2, hence 8x8 is 4 switches); col. 2, lines 31-33; 'networks can be generalized'; col. 6, lines 7-9; 'butterfly pattern'; col. 13, lines 23-24), but does not disclose said routing network is strictly nonblocking for unicast traffic and rearrangeably nonblocking for arbitrary fan-out multicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being strictly nonblocking for unicast traffic (col. 2, lines 9-11, see Fig. 9) and rearrangeably nonblocking for arbitrary fan-out multicast traffic (col. 4, lines 6-7; see 'fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a muting network be strictly nonblocking for unicast traffic and rearrangeably nonblocking for arbitrary fan-out multicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 32, Wong discloses wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network generalized butterfly fat tree network with full bandwidth ('8x8 Benes network' (each switch has two logic calls that are 2x2, hence 8x8 is 4 switches); col. 2, lines 31-33; 'networks can be generalized'; col. 6, Lines 7-9), but does not disclose said routing network is strictly nonblocking for arbitrary fan-out multicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract)

and a routing network strictly nonblocking for arbitrary fan-out multicast traffic (col. 2, lines 8-11; see 'fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be strictly nonblocking for arbitrary fan-out multicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 34, Wong discloses d = 4 ('plurality of input terminals and a number of output terminals', claim 1) and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network multi-link Benes network with full bandwidth (col. 2, lines 27-30; '2x2 Benes network'; col. 5, lines 26-31), but does not disclose said routing network is rearrangeably nonblocking for unicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being rearrangeably nonblocking for unicast traffic ('rearrangeably nonblocking for unicast traffic'; col. 2, lines 52-54).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be rearrangeably nonblocking for unicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 35, Wong discloses d = 4 ('plurality of input terminals and a number of output terminals', claim 1) and there are at least two switches in each said stage in each

said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is a multi-link Benes network ('8x8 Benes network' (each switch has two logic cells that are 2x2, hence 8x8 is 4 switches); col. 2, lines 31-33; col. 6, lines 7-9), but does not disclose said routing network is strictly nonblocking for unicast traffic and rearrangeably nonblocking for arbitrary fan-out multicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being strictly nonblocking for unicast traffic (col. 2, lines 9-11, see Fig. 9) and rearrangeably nonblocking for arbitrary fan-out multicast traffic (col. 4, lines 6-7; see 'fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be strictly nonblocking for unicast traffic and rearrangeably nonblocking for arbitrary fan-out multicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 36, Wong discloses d = 4 ('plurality of input terminals and a number of output terminals', claim 1) and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting links and said routing network multi-link Benes network with full bandwidth ('6x8 Benes network' (each switch has two logic cells that are 2x2, hence 8x8 is 4 switches); col. 2, lines 31-33; col. 6, lines 7-9), but does not

disclose said routing network is strictly nonblocking for arbitrary fan-out multicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being rearrangeably nonblocking for arbitrary fan-out multicast traffic (col. 4, lines 6-7; see 'fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be rearrangeably nonblocking for arbitrary fan-out multicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 37, Wong discloses d = 4 ('plurality of input terminals and a number of output terminals', claim 1) and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is a butterfly fat tree network with full bandwidth (col. 2, lines 27-30; col. 2. lines 31-33; 'butterfly pattern'; col. 13, lines 23-24); but does not disclose said routing network is rearrangeably nonblocking for unicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being rearrangeably nonblocking for unicast traffic ('rearrangeably nonblocking for unicast traffic'; col. 2, lines 52-54).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be rearrangeably nonblocking for unicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 38, Wong discloses wherein d = 4 ('plurality of input terminals and a number of output terminals', claim 1) and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links ('8x8 Benes network' (each switch has two logic cells that are 2x2. hence 8x8 is 4 switches); col. 2, lines 31-33; col. 6, lines 7-9) and said routing network is a multi-link butterfly fat tree network (col. 2, lines 27-30; col. 2, lines 31-33; 'butterfly pattern'; col. 13, lines 23-24), but does not disclose said routing network is strictly nonblocking for unicast and rearrangeably nonblocking for arbitrary fan-out multicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being strictly nonblocking for unicast traffic (col. 2, lines 9-11, see Fig. 9) and rearrangeably nonblocking for arbitrary fan-out multicast traffic (col. 2, lines 9-11, see Fig. 9) and rearrangeably nonblocking for 3.

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be strictly nonblocking for unicast traffic and rearrangeably nonblocking for arbitrary fan-out multicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 39, Wong discloses d = 4 ('plurality of input terminals and a number of output terminals', claim 1) and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit

block connecting said backward connecting links and said routing network is a multi-link butterfly fat tree network with full bandwidth ('8x8 Benes network' (each switch has two logic cells that are 2x2. hence 8x8 is 4 switches; 'butterfly pattern'; col. 13, lines 23-24); col 2, lines 31-33; col. 6, lines 7-9), but does not discloses said routing network is strictly nonblocking for arbitrary fan-out multicast. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network strictly nonblocking for arbitrary fan-out multicast traffic (col. 2, lines 8-11; see 'fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be strictly nonblocking for arbitrary fan-out multicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 41, Wong discloses d = 4 ('plurality of input terminals and a number of output terminals', claim 1) and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links (col. 2, lines 27-30; col. 2, lines 31-33) and said routing network is a generalized multi-link multi-stage network with full bandwidth ('networks can be generalized'; col. 6, lines 7-9), but does not disclose said routing network is rearrangeably nonblocking for unicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network

being rearrangeably nonblocking for unicast traffic ('rearrangeably nonblocking for unicast traffic'; col. 2, lines 52-54).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be rearrangeably nonblocking for unicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 42, Wong discloses d = 4 ('plurality of input terminals and a number of output terminals', claim 1) and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links ('8x8 Benes network' (each switch has two logic cells that are 2x2, hence 8x8 is 4 switches); col. 2, lines 31-33; col. 6, lines 7-9) and said routing network is a generalized multi-link multi-stage network ('networks can be generalized'; col. 6, lines 7-9), but does not disclose said routing network is strictly nonblocking for unicast and rearrangeably nonblocking for arbitrary fan-out multicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network being strictly nonblocking for unicast traffic (col. 2, lines 9-11, see Fig. 9) and rearrangeably nonblocking for arbitrary fan-out multicast traffic (col. 4, lines 6-7; see "fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be strictly nonblocking for unicast traffic and rearrangeably nonblocking for

arbitrary fan-out multicast traffic, as disclosed by Wu, In order to satisfy specific routing requirements.

In claim 43, Wong discloses d = 4 ('plurality of input terminals and a number of output terminals', claim 1) and there are at least three switches in each said stage in each said sub-integrated circuit connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is a generalized multi-link multistage network with full bandwidth ('8x8 Benes network' (each switch has two logic cells that are 2x2, hence 8x8 is 4 switches); col. 2, lines 31-33; col. 6, lines 7-9; 'networks can be generalized'; col. 6, lines 7-9), but does not disclose said routing network is strictly nonblocking for arbitrary fan-out multicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network strictly nonblocking for arbitrary fan-out multicast traffic (col. 2, lines 8-11; see 'fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be strictly nonblocking for arbitrary fan-out multicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 45, Wong discloses d = 4 ('plurality of input terminals and a number of output terminals', claim 1} and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said

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backward connecting links and said routing network is a generalized multi-link butterfly fat tree network with full bandwidth (col. 2, lines 27-30; col. 2, lines 31-33; 'networks can be generalized'; col. 6, lines 7-9; 'butterfly pattern'; col. 13, lines 23-24), but does not disclose said routing network is rearrangeably nonblocking for unicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeeble non-blocking switch (Abstract) and a routing network being rearrangeably nonblocking for unicast traffic ('rearrangeably nonblocking for unicast traffic'; col. 2, lines 52-54).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be rearrangeably nonblocking for unicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 46, Wong discloses d = 4 ('plurality of input terminals and a number of output terminals', claim 1) and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting links and said routing network generalized multi-link butterfly fat tree network with full bandwidth ('8x8 Benes network' (each switch has two logic cells that are 2x2, hence 8x8 is 4 switches); col. 2, lines 31-33; col. 6, lines 7-9; networks can be generalized'; col. 6, lines 7-9; 'butterfly pattern'; col. 13, lines 23-24), but does not disclose said routing network is strictly nonblocking for unicast traffic and rearrangeably nonblocking for arbitrary fan-out multicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a

routing network being strictly nonblocking for unicast traffic (col. 2, lines 9-11, see Fig. 9) and rearrangeably nonblocking for arbitrary fan-out multicast traffic (col. 4, lines 6-7; see 'fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be strictly nonblocking for unicast traffic and rearrangeably nonblocking for arbitrary fan-out multicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

In claim 47, Wong discloses d = 4 ('plurality of input terminals and a number of output terminals', claim 1) and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting links and said routing network is a generalized multilink butterfly fat tree network with full bandwidth ('8x8 Benes network' (each switch has two logic cells that are 2x2, hence 8x8 is 4 switches); col. 2, lines 31-33; col. 6, lines 7-9; networks can be generalized'; col. 6, lines 7-9; 'butterfly pattern'; Col. 13, Lines 23-24), but does not disclose said routing network is strictly nonblocking for arbitrary fan-out multicast traffic. However, in the same field of endeavor, Wu discloses a rearrangeable non-blocking switch (Abstract) and a routing network strictly nonblocking for arbitrary fan-out multicast traffic (col. 2, lines 8-11; see 'fanouts', Fig. 7B).

Therefore it would have been obvious to one ordinary skilled in the art at the time of the invention was made to supplement the teachings of Wong and have a routing network be strictly nonblocking for arbitrary fan-out multicast traffic, as disclosed by Wu, in order to satisfy specific routing requirements.

Any inquiry concerning this communication or earlier communications from the examiner should be directed to Vibol Tan whose telephone number is (571)272-1811. The examiner can normally be reached on Monday-Friday (7:00 AM-4:30 PM).

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Shawki Ismail can be reached on (571) 272-3985. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see http://pair-direct.uspto.gov. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

/Vibol Tan/ Primary Examiner, Art Unit 2819

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EAST Search History

EAST Search History (Prior Art)

Ref #	Hits	Search Query	DBs	Default Operator	Plurals	Time Stamp
L1	35	(US-20080267182-\$ or US- 20040114586-\$ or US-20030021267-\$ or US-20030112797-\$ or US- 20030053456-\$ or US-20020051447-\$ or US-20020031118-\$ or US- 20020031117-\$).did. or (US-8098081-\$ or US-7924052-\$ or US-7468974-\$ or US-7424011-\$ or US-7424010-\$ or US- 7397796-\$ or US-7349387-\$ or US- 7346049-\$ or US-7130920-\$ or US- 6567858-\$ or US-7136380-\$ or US- 6456838-\$ or US-6157643-\$ or US- 6456838-\$ or US-6157643-\$ or US- 64456556-\$ or US-856977-\$ or US- 5406556-\$ or US-7103059-\$ or US- 7139266-\$ or US-7065074-\$ or US- 7099314-\$ or US-7065074-\$ or US- 7050429-\$ or US-7016345-\$ or US- 7031303-\$ or US-7016345-\$ or US- 6952418-\$).did. or (US-6201808-\$).did.	US-PGPUB; USPAT	OR	OFF	2012/02/02 10:27
S1	2	(("6940308") or ("7154887")).PN.	USPAT; USOCR	OR	OFF	2012/02/01 07:26
S2	3	(("6940308") or ("7154887") or ("20110037498")).PN.	US-PGPUB; USPAT; USOCR	OR	OFF	2012/02/01 07:30
S3	81	nonblocking adj multicast	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 08:51
S4	46	non-blocking adj multicast	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 08:51
S5	8	non-blocking adj multi-cast	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 08:52
S6	111	S3 S4 S5	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT;		OFF	2012/02/02 08:52

Page 121 of 374 file:///Cl/Users/vtan/Documents/e-Red%20Folder/12601275/EASTSearchHistory.12601275_AccessibleVersion.htm[2/2/2012 10:33:45 AM]

~7		<u> </u>	IBM_TDB			
S7	35	S6 and rearrangeably	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 08:53
S8	21	S6 and rearrangeable	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 08:53
S9	39	S7 S8	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 08:53
S10	2	(routing adj network) same (inlet adj links) same (outlet adj links)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 09:18
S11	173	(routing adj network) same (inputs) same (outputs)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 09:19
S12	223328	(("326") or ("370")).CLAS.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 09:22
S13	108	S11 and S12	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 09:22
S14	0	(routing adj network) same (inputs) same (outputs) same (higher adj stage) same (lower adj stage)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 09:31
S15	1	(routing adj network) and (inputs) and (outputs) and (higher adj stage) and (lower adj stage)	US-PGPUB; USPAT; USOCR; FPRS;	OR	OFF	2012/02/02 09:31

			EPO; JPO; DERWENT; IBM_TDB			
S16	2	(routing adj network) and (higher adj stage) and (lower adj stage)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 09:31
S17	24	(routing adj network) and (cross adj links)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 09:32
S18	5	(routing adj network) and (cross adj switches)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 09:32
S19	26	S17 S18	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 09:33
S20	882	(routing adj network) and stages	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 09:37
S21	20	(US-20080267182-\$ or US- 20040114586-\$ or US-20030021267- \$).did. or (US-8098081-\$ or US- 7924052-\$ or US-7468974-\$ or US- 7424011-\$ or US-7424010-\$ or US- 7397796-\$ or US-7349387-\$ or US- 7346049-\$ or US-7130920-\$ or US- 6567858-\$ or US-7136380-\$ or US- 6456838-\$ or US-6157643-\$ or US- 64456838-\$ or US-6157643-\$ or US- 6049542-\$ or US-5856977-\$ or US- 5406556-\$ or US-4813038-\$).did.	US-PGPUB; USPAT	OR	OFF	2012/02/02 09:39
S22	41	konda-venkat.in.	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO; DERWENT; IBM_TDB	OR	OFF	2012/02/02 09:40
S23	84	(routing adj network) and (multicast adj connections)	US-PGPUB; USPAT; USOCR; FPRS; EPO; JPO;	OR	OFF	2012/02/02 09:45



EAST Search History (Interference)

< This search history is empty>

2/ 2/ 2012 10:33:37 AM C:\ Users\ vtan\ Documents\ EAST\ Workspaces\ 12601275.wsp

	Application/Control No.	Applicant(s)/Patent Under Reexamination
Search Notes	12601275	KONDA, VENKAT
	Examiner	Art Unit
	VIBOL TAN	2819

SEARCHED							
Class	Subclass	Date	Examiner				
326	38-41	2/2/12	VTan				
370	390, 312, 360, 388, 412	2/2/12	VTan				

SEARCH NOTES							
Search Notes	Date	Examiner					
Inventor search	2/2/12	VTan					
search report pct/us2008/064605	2/1/12	VTan					
EAST text search	2/2/12	VTan					

	INTERFERENCE SEARCH		
Class	Subclass	Date	Examiner

UNITED ST	ates Patent and Trademan	UNITED STA' United States Address: COMMI P.O. Box I	a, Virginia 22313-1450
APPLICATION NUMBER	FILING OR 371(C) DATE	FIRST NAMED APPLICANT	ATTY. DOCKET NO./TITLE
12/601,275	05/31/2010	Venkat Konda	V-0045US
			CONFIRMATION NO. 6372
38139		PUBLICAT	TION NOTICE
Konda Technologies, Inc 6278 GRAND OAK WAY SAN JOSE, CA 95135			C000000046063228*

Title:VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS

Publication No.US-2011-0037498-A1 Publication Date:02/17/2011

NOTICE OF PUBLICATION OF APPLICATION

The above-identified application will be electronically published as a patent application publication pursuant to 37 CFR 1.211, et seq. The patent application publication number and publication date are set forth above.

The publication may be accessed through the USPTO's publically available Searchable Databases via the Internet at www.uspto.gov. The direct link to access the publication is currently http://www.uspto.gov/patft/.

The publication process established by the Office does not provide for mailing a copy of the publication to applicant. A copy of the publication may be obtained from the Office upon payment of the appropriate fee set forth in 37 CFR 1.19(a)(1). Orders for copies of patent application publications are handled by the USPTO's Office of Public Records. The Office of Public Records can be reached by telephone at (703) 308-9726 or (800) 972-6382, by facsimile at (703) 305-8759, by mail addressed to the United States Patent and Trademark Office, Office of Public Records, Alexandria, VA 22313-1450 or via the Internet.

In addition, information on the status of the application, including the mailing date of Office actions and the dates of receipt of correspondence filed in the Office, may also be accessed via the Internet through the Patent Electronic Business Center at www.uspto.gov using the public side of the Patent Application Information and Retrieval (PAIR) system. The direct link to access this status information is currently http://pair.uspto.gov/. Prior to publication, such status information is confidential and may only be obtained by applicant using the private side of PAIR.

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Office of Data Managment, Application Assistance Unit (571) 272-4000, or (571) 272-4200, or 1-888-786-0101

UNITED STATES PATE
U.S. APPLICATION NUMBER NO.
10/(01.075

ent and Trademark Office

	Address: COMMIS P.O. Box 1	SSIONER FOR P. 1450 a, Virginia 22313-145		
U.S. APPLICATION NUMBER NO.	FIRST NAMED APPLICANT		ATT	Y. DOCKET NO.
12/601,275	Venkat Konda	V-0045US		
38139		INTER	NATIONAL AP	PLICATION NO.
Konda Technologies, Inc	PCT/US08/64605			
6278 GRAND OAK WAY			NG DATE	PRIORITY DATE
SAN JOSE, CA 95135		05/22	2/2008	05/25/2007
		37		IATION NO. 6372 TANCE LETTER

Date Mailed: 11/09/2010

NOTICE OF ACCEPTANCE OF APPLICATION UNDER 35 U.S.C 371 AND 37 CFR 1.495

The applicant is hereby advised that the United States Patent and Trademark Office in its capacity as a Designated / Elected Office (37 CFR 1.495), has determined that the above identified international application has met the requirements of 35 U.S.C. 371, and is ACCEPTED for national patentability examination in the United States Patent and Trademark Office.

The United States Application Number assigned to the application is shown above and the relevant dates are:

05/31/2010 DATE OF RECEIPT OF 35 U.S.C. 371(c)(1), (c)(2) and (c)(4) REQUIREMENTS

05/31/2010 DATE OF COMPLETION OF ALL 35 U.S.C. 371 REQUIREMENTS

UNITED STATES DEPARTMENT OF COMMERCE

A Filing Receipt (PTO-103X) will be issued for the present application in due course. THE DATE APPEARING ON THE FILING RECEIPT AS THE "FILING DATE" IS THE DATE ON WHICH THE LAST OF THE 35 U.S.C. 371 (c)(1), (c)(2) and (c)(4) REQUIREMENTS HAS BEEN RECEIVED IN THE OFFICE. THIS DATE IS SHOWN ABOVE. The filing date of the above identified application is the international filing date of the international application (Article 11(3) and 35 U.S.C. 363). Once the Filing Receipt has been received, send all correspondence to the Group Art Unit designated thereon.

The following items have been received:

- Indication of Small Entity Status
- Copy of the International Application filed on 11/22/2009
- Copy of the International Search Report filed on 11/22/2009
- Oath or Declaration filed on 05/31/2010
- U.S. Basic National Fees filed on 11/22/2009
- Specification filed on 11/22/2009
- Claims filed on 11/22/2009
- Abstracts filed on 11/22/2009
- Drawings filed on 11/22/2009

Applicant is reminded that any communications to the United States Patent and Trademark Office must be mailed to the address given in the heading and include the U.S. application no. shown above (37 CFR 1.5)

NISA F GILCHRIST

Telephone: (703) 756-1418

UNITED STATES PATENT AND TRADEMARK OFFICE United States Patent and Trademark Office Address: COMMISSIONER FOR PATENTS PO Box 1450 Advandra, Virginia 22313-1450 www.uspto.gov								
APPLICATION NUMBER	FILING or 371(c) DATE	GRP ART UNIT	FIL FEE REC'D	ATTY.DOCKET.NO	TOT CLAIMS IND CLAIMS			
12/601,275	05/31/2010	2819	1279	V-0045US	49 1			
				CON	FIRMATION NO. 6372			
38139				FILING RECEI	PT			
Konda Technologies, Inc								
SAN JOSE, C/	A 95135				000043036610			

Date Mailed: 11/09/2010

Receipt is acknowledged of this non-provisional patent application. The application will be taken up for examination in due course. Applicant will be notified as to the results of the examination. Any correspondence concerning the application must include the following identification information: the U.S. APPLICATION NUMBER, FILING DATE, NAME OF APPLICANT, and TITLE OF INVENTION. Fees transmitted by check or draft are subject to collection. Please verify the accuracy of the data presented on this receipt. If an error is noted on this Filing Receipt, please submit a written request for a Filing Receipt Correction. Please provide a copy of this Filing Receipt with the changes noted thereon. If you received a "Notice to File Missing Parts" for this application, please submit any corrections to this Filing Receipt with your reply to the Notice. When the USPTO processes the reply to the Notice, the USPTO will generate another Filing Receipt incorporating the requested corrections

Applicant(s)

Venkat Konda, San Jose, CA;

Power of Attorney: None

Domestic Priority data as claimed by applicant

This application is a 371 of PCT/US08/64605 05/22/2008 which claims benefit of 60/940,394 05/25/2007

Foreign Applications

If Required, Foreign Filing License Granted: 04/08/2010

The country code and number of your priority application, to be used for filing abroad under the Paris Convention, is **US 12/601,275**

Projected Publication Date: 02/17/2011

Non-Publication Request: No

Early Publication Request: No ** SMALL ENTITY **

VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS

Preliminary Class

326

PROTECTING YOUR INVENTION OUTSIDE THE UNITED STATES

Since the rights granted by a U.S. patent extend only throughout the territory of the United States and have no effect in a foreign country, an inventor who wishes patent protection in another country must apply for a patent in a specific country or in regional patent offices. Applicants may wish to consider the filing of an international application under the Patent Cooperation Treaty (PCT). An international (PCT) application generally has the same effect as a regular national patent application in each PCT-member country. The PCT process **simplifies** the filing of patent applications on the same invention in member countries, but **does not result** in a grant of "an international patent" and does not eliminate the need of applicants to file additional documents and fees in countries where patent protection is desired.

Almost every country has its own patent law, and a person desiring a patent in a particular country must make an application for patent in that country in accordance with its particular laws. Since the laws of many countries differ in various respects from the patent law of the United States, applicants are advised to seek guidance from specific foreign countries to ensure that patent rights are not lost prematurely.

Applicants also are advised that in the case of inventions made in the United States, the Director of the USPTO must issue a license before applicants can apply for a patent in a foreign country. The filing of a U.S. patent application serves as a request for a foreign filing license. The application's filing receipt contains further information and guidance as to the status of applicant's license for foreign filing.

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For information on preventing theft of your intellectual property (patents, trademarks and copyrights), you may wish to consult the U.S. Government website, http://www.stopfakes.gov. Part of a Department of Commerce initiative, this website includes self-help "toolkits" giving innovators guidance on how to protect intellectual property in specific countries such as China, Korea and Mexico. For questions regarding patent enforcement issues, applicants may call the U.S. Government hotline at 1-866-999-HALT (1-866-999-4158).

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Title 37, Code of Federal Regulations, 5.11 & 5.15

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The grant of a license does not in any way lessen the responsibility of a licensee for the security of the subject matter as imposed by any Government contract or the provisions of existing laws relating to espionage and the national security or the export of technical data. Licensees should apprise themselves of current regulations especially with respect to certain countries, of other agencies, particularly the Office of Defense Trade Controls, Department of State (with respect to Arms, Munitions and Implements of War (22 CFR 121-128)); the Bureau of Industry and Security, Department of Commerce (15 CFR parts 730-774); the Office of Foreign AssetsControl, Department of Treasury (31 CFR Parts 500+) and the Department of Energy.

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UNITED STATES PATENT AND TRADEMARK OFFICE United States Patent and Trademark Office Address: COMMISSIONER FOR PATENTS P.D. Box 1450 Alexandria, Virginia 22313-1450 www.uspto.gov							
APPLICATION	FILING or	GRP ART					
NUMBER	371(c) DATE	UNIT	FIL FEE REC'D	ATTY.DOCKET.NO	TOT CLAIMS IND CLAIMS		
12/601,275	05/31/2010		1279	V-0045US	49 1		
				CON	IFIRMATION NO. 6372		
38139 Kanda Taabaa	la sia a lua a			FILING RECE	IPT		
Konda Technologies, Inc 6278 GRAND OAK WAY SAN JOSE, CA 95135							

Date Mailed: 08/18/2010

Receipt is acknowledged of this non-provisional patent application. The application will be taken up for examination in due course. Applicant will be notified as to the results of the examination. Any correspondence concerning the application must include the following identification information: the U.S. APPLICATION NUMBER, FILING DATE, NAME OF APPLICANT, and TITLE OF INVENTION. Fees transmitted by check or draft are subject to collection. Please verify the accuracy of the data presented on this receipt. If an error is noted on this Filing Receipt, please submit a written request for a Filing Receipt Correction. Please provide a copy of this Filing Receipt with the changes noted thereon. If you received a "Notice to File Missing Parts" for this application, please submit any corrections to this Filing Receipt with your reply to the Notice. When the USPTO processes the reply to the Notice, the USPTO will generate another Filing Receipt incorporating the requested corrections

Applicant(s)

Venkat Konda, San Jose, CA;

Power of Attorney: None

Domestic Priority data as claimed by applicant

This application is a 371 of PCT/US08/64605 05/22/2008 which claims benefit of 60/940,394 05/25/2007

Foreign Applications

If Required, Foreign Filing License Granted: 04/08/2010

The country code and number of your priority application, to be used for filing abroad under the Paris Convention, is **US 12/601,275**

Projected Publication Date: 371 Perfected

Non-Publication Request: No

Early Publication Request: No ** SMALL ENTITY **

VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS

Preliminary Class

PROTECTING YOUR INVENTION OUTSIDE THE UNITED STATES

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NOT GRANTED

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UNITED STATES PATENT AND TRADEMARK OFFICE

	UNITED 51 AT ES DEFARTMENT OF COMMERCE United States Patent and Trademark Office Address: COMMISSIONER FOR PATENTS PO. Box 1450 Alexandria, Virginia 22313-1450 www.uspto.gov			
U.S. APPLICATION NUMBER NO.	FIRST NAMED APPLICANT		ATT	Y. DOCKET NO.
12/601,275	Venkat Konda		I	7-0045US
38139		INTERI	NATIONAL AP	PLICATION NO.
Konda Technologies, Inc			PCT/US08	/64605
6278 GRAND OAK WAY		I.A. FILII	NG DATE	PRIORITY DATE
SAN JOSE, CA 95135		05/22	2/2008	05/25/2007
				IATION NO. 6372 TANCE LETTER

Date Mailed: 08/18/2010

NOTICE OF ACCEPTANCE OF APPLICATION UNDER 35 U.S.C 371 AND 37 CFR 1.495

The applicant is hereby advised that the United States Patent and Trademark Office in its capacity as a Designated / Elected Office (37 CFR 1.495), has determined that the above identified international application has met the requirements of 35 U.S.C. 371, and is ACCEPTED for national patentability examination in the United States Patent and Trademark Office.

The United States Application Number assigned to the application is shown above and the relevant dates are:

<u>05/31/2010</u> DATE OF RECEIPT OF 35 U.S.C. 371(c)(1), (c)(2) and (c)(4) REQUIREMENTS 05/31/2010 DATE OF COMPLETION OF ALL 35 U.S.C. 371 REQUIREMENTS

A Filing Receipt (PTO-103X) will be issued for the present application in due course. **THE DATE APPEARING ON THE FILING RECEIPT AS THE "FILING DATE" IS THE DATE ON WHICH THE LAST OF THE 35 U.S.C. 371 (c)(1), (c)(2) and (c)(4) REQUIREMENTS HAS BEEN RECEIVED IN THE OFFICE. THIS DATE IS SHOWN ABOVE.** *The filing date of the above identified application is the international filing date of the international application (Article 11(3) and 35 U.S.C. 363).* Once the Filing Receipt has been received, send all correspondence to the Group Art Unit designated thereon.

The following items have been received:

- · Indication of Small Entity Status
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- Copy of the International Search Report filed on 11/22/2009
- Oath or Declaration filed on 05/31/2010
- U.S. Basic National Fees filed on 11/22/2009
- Specification filed on 11/22/2009
- Claims filed on 11/22/2009
- Abstracts filed on 11/22/2009
- Drawings filed on 11/22/2009

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NISA F GILCHRIST

Telephone: (703) 756-1418

Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it contains a valid OMB control number

\bigcap	DECLA	_	-	UTILITY OR	Attorney Docket Number	
	D/				First Named Inventor	Venkat Konda
PATENT APPLICATION (37 CFR 1.63)					COMF	PLETE IF KNOWN
		\		Declaration	Application Number	12/601,275
	Declaration Submitted	OR	X	Submitted After Initial	Filing Date	11/22/2009
	With Initial Filing	OR		Filing (surcharge (37 CFR 1.16(f))	Art Unit	
				required)	Examiner Name	
and (2) for whi) I believe the in ch a patent is so	ventor(s) n ought on th	amed be e inventio	elow to be the original a	and first inventor(s) of the	as stated below next to their name; subject matter which is claimed and
				(Title of the	Invention)	
the ap	plication of which	ı		(The of the	mventiony	
	is attached he	ereto				
OR						
X	was filed on (MM/DD/YY	(YY)	a	s United States Application	n Number or PCT International
	Application Nu	ımber	12/601	,275 _and was am	ended on (MM/DD/YYYY)	(if applicable).
				nderstand the contents referred to above.	of the above identified ap	plication, including the claims, as
continu	uation-in-part ap	plications,	material		ame available between the	ed in 37 CFR 1.56, including for filing date of the prior application
Autho	prization To P	ermit Aco	cess To	Application by Pa	rticipating Offices	
Japan any oth filed ac applica	If checked, the Patent Office (JI ner intellectual p ccess to the abor ant does not wish	e undersig PO), the Ko roperty offi ve-identifie n the EPO,	ned here brean Int ces in wh d patent JPO, Kli	by grants the USPTO a ellectual Property Offic nich a foreign applicatio application. See 37 C PO, WIPO, or other int	authority to provide the Eu e (KIPO), the World Intelle on claiming priority to the a FR 1.14(c) and (h). This b	ropean Patent Office (EPO), the ectual Property Office (WIPO), and bove-identified patent application is ox should not be checked if the which a foreign application claiming ied patent application.
to: 1) f claims 37 CFF	the above-identi priority under 35	fied patent 5 U.S.C. 11 filed in the	applicati 9(a)-(d) above-i	on-as-filed; 2) any fore if a copy of the foreign dentified patent applica	ign application to which the application that satisfies the set of	ified patent application with respect e above-identified patent application ne certified copy requirement of ication-as-filed from which benefit is
	ordance with 37 Access to Appli				nformation concerning the	date of filing the Authorization to
					nformation is required to obtain or	retain a benefit by the public which is to file (and This collection is estimated to take 21 minutes to

by the USPTO to process) an application. Confidentiality is governed by 35 U.S.C. 122 and 37 CFR 1.11 and 1.14. This collection is estimated to take 21 minutes to complete, including gathering, preparing, and submitting the completed application form to the USPTO. Time will vary depending upon the individual case. Any comments on the amount of time you require to complete this form and/or suggestions for reducing this burden, should be sent to the Chief Information Officer, U.S. Patent and Trademark Office, U.S. Department of Commerce, P.O. Box 1450, Alexandria, VA 22313-1450. DO NOT SEND FEES OR COMPLETED FORMS TO THIS ADDRESS. **SEND TO: Commissioner for Patents, P.O. Box 1450, Alexandria, VA 22313-1450.**

DECLARATION — Utility or Design Patent Application

Claim of Foreign Priority Benefits

I hereby claim foreign priority benefits under 35 U.S.C. 119(a)-(d) or (f), or 365(b) of any foreign application(s) for patent, inventor's or plant breeder's rights certificate(s), or 365(a) of any PCT international application which designated at least one country other than the United States of America, listed below and have also identified below, by checking the box, any foreign application for patent, inventor's or plant breeder's rights certificate(s), or any PCT international application having a filing date before that of the application on which priority is claimed.

Prior Foreign Application Number(s)	Country	Foreign Filing Date (MM/DD/YYYY)	Priority Not Claimed	Certified Co YES	py Attached? NO	
PCT/US08/64605	USA	5/22/2008			X	
Additional foreign application number(s) are listed on a supplemental priority data sheet PTO/SB/02B attached hereto.						

[Page 2 of 3]

PTO/SB/01 (04-09) Approved for use through 06/30/2010. OMB 0651-0032 U.S. Patent and Trademark Office; U.S. DEPARTMENT OF COMMERCE Under the Paperwork Reduction Act of 1995, no persons are required to respond to a collection of information unless it contains a valid OMB control number.

DECLAI	RATION — Utili	ity or Desig	n Patent Appl	lication			
correspondence to: X assoc	address ciated with 381 omer Number:	139	OR		Correspondence address below		
Name							
Address							
City		State		Zip			
Country	Telephone	I	Email				
Petitioner/applicant is cautioned to		WARNING:	- '- de sumente files	l'n e nofe	-t-selication that may		
contribute to identity theft. Persona (other than a check or credit card a USPTO to support a petition or an a USPTO, petitioners/applicants shou to the USPTO. Petitioner/applicant the application (unless a non-public a patent. Furthermore, the record f referenced in a published application PTO-2038 submitted for payment p Petitioner/applicant is advised that of into the Privacy Act system of record <i>Files</i> . Documents not retained in an COMMERCE/PAT-TM-10, System I hereby declare that all statements belief are believed to be true; and fur the like so made are punishable by may jeopardize the validity of the application NAME OF SOLE OR FIRST IN	uthorization form PT application. If this ty uld consider redacting it is advised that the re- cation request in com from an abandoned a on or an issued pater purposes are not reta documents which for rds DEPARTMENT C n application file (suc name: Deposit Acco s made herein of my of urther that these stat fine or imprisonment pplication or any pate	O-2038 submitted pe of personal in g such personal record of a pater appliance with 37 application may int (see 37 CFR ined in the appli rm the record of OF COMMERCE ch as the PTO-2 bunts and Electro own knowledge tements were may it, or both, under ent issued there	ed for payment purp formation is include information from th t application is avai CFR 1.213(a) is ma also be available to .14). Checks and o cation file and there a patent application c. COMMERCE-PAT 038) are placed into mic Funds Transfer are true and that all ade with the knowle 18 U.S.C. 1001 and on.	oses) is r ed in docu e docume lable to the the public credit card fore are r (such as F-7, Syste b the Priva <i>Profiles.</i> statemen dge that we	never required by the iments submitted to the ents before submitting them he public after publication of application) or issuance of c if the application is d authorization forms not publicly available. the PTO/SB/01) are placed em name: <i>Patent Application</i> acy Act system of the statements and willful false statements and th willful false statements		
Given Name (first and middle [if any		A petition has been filed for this unsigned inventor Family Name or Surname					
Venkat		Konda					
Inventor's Signature			Date				
/Venkat Konda/			5/30/2010				
Residence: City Sta	ate	Count	ry	Ci	itizenship		
	CA	USA		USA			
Mailing Address				I			
6278 Grand Oak Way							
City Sta	ate	Zip		C	ountry		
San Jose C	A	951	35		USA		
Additional inventors or a legal rep	resentative are being nam	ied on the	supplemental sheet(s) PT	O/SB/02A or	· 02LR attached hereto		

Electronic Patent A	Application Fe	e Transmi	ttal				
Application Number:	12601275						
Filing Date:							
Title of Invention:	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS						
First Named Inventor/Applicant Name:	Venkat Konda						
Filer:	Venkar Konda						
Attorney Docket Number:	V-0045US						
Filed as Small Entity							
U.S. National Stage under 35 USC 371 Filing	Fees						
Description	Fee Code	Quantity	Amount	Sub-Total in USD(\$)			
Basic Filing:							
Natl Stage Search Fee - U.S. was the ISA	2641	1	50	50			
Natl Stage Exam Fee - all other cases	2633	1	110	110			
Pages:							
Claims:							
Miscellaneous-Filing:							
Oath/decl > 30 mo. from priority date	2617	1	65	65			
Petition:							
Patent-Appeals-and-Interference:							

Description	Fee Code	Quantity	Amount	Sub-Total in USD(\$)
Post-Allowance-and-Post-Issuance:				
Extension-of-Time:				
Miscellaneous:				
	Tot	al in USD	(\$)	225

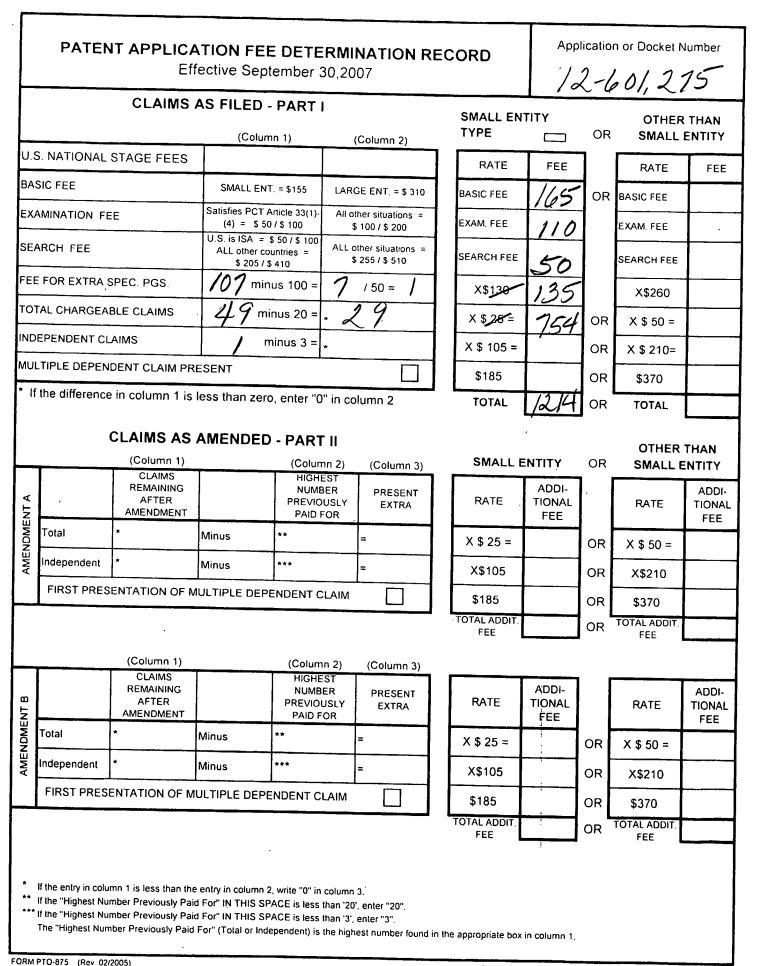
Electronic Ac	Electronic Acknowledgement Receipt							
EFS ID:	7714705							
Application Number:	12601275							
International Application Number:								
Confirmation Number:	6372							
Title of Invention:	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS							
First Named Inventor/Applicant Name:	Venkat Konda							
Customer Number:	38139							
Filer:	Venkar Konda							
Filer Authorized By:								
Attorney Docket Number:	V-0045US							
Receipt Date:	31-MAY-2010							
Filing Date:								
Time Stamp:	02:44:41							
Application Type:	U.S. National Stage under 35 USC 371							

Payment information:

Document Number Page 142 of	Document Description	File Name	File Size(Bytes)/ Message Digest	Multi Part /.zip	Pages (if appl.)					
File Listing	:									
Authorized Use	r									
Deposit Accour	nt									
RAM confirmati	on Number	6077								
Payment was su	uccessfully received in RAM	\$225								
Payment Type		Credit Card	Credit Card							
Submitted with	Payment	yes	yes							

	Application Data Sheet	sb0001.pdf	813566 367e6d42d0fb354655cf1c0bddd5d1c98b7 ef69f	no	3
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Post Card, as <u>New Applicat</u> If a new appli 1.53(b)-(d) an Acknowledge <u>National Stag</u> If a timely sub U.S.C. 371 and	described in MPEP 503. ions Under 35 U.S.C. 111 cation is being filed and the applicat d MPEP 506), a Filing Receipt (37 CFF	ion includes the necessa ? 1.54) will be issued in d date of the application. der 35 U.S.C. 371 of an international appli orm PCT/DO/EO/903 indi	ary components for a filing due course and the date sh ication is compliant with th icating acceptance of the a	date (see own on thi ne conditio pplication	37 CFR is ons of 35

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Page 145 of 374

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UNITED STATES PATENT AND TRADEMARK OFFICE

		United States Patent and Trademark Office Address: COMMISSIONER FOR PATENTS PO. Box 1450 Alexandria, Virginia 22313-1450 www.uspto.gov			
U.S. APPLICATION NUMBER NO.	FIRST NAMED APPLICANT		ATT	Y. DOCKET NO.	
12/601,275	Venkat Konda		I I	7-0045US	
38139		INTER	NATIONAL AP	PLICATION NO.	
Konda Technologies, Inc			PCT/US08	/64605	
6278 GRAND OAK WAY		I.A. FILI	NG DATE	PRIORITY DATE	
SAN JOSE, CA 95135		05/22	2/2008	05/25/2007	
		3		IATION NO. 6372 ALITIES LETTER	

UNITED STATES DEPARTMENT OF COMMERCE

Date Mailed: 04/08/2010

NOTIFICATION OF MISSING REQUIREMENTS UNDER 35 U.S.C. 371 IN THE UNITED STATES DESIGNATED/ELECTED OFFICE (DO/EO/US)

The following items have been submitted by the applicant or the IB to the United States Patent and Trademark Office as a Designated Office (37 CFR 1.494):

- Indication of Small Entity Status
- Priority Document
- Copy of the International Application filed on 11/22/2009
- Copy of the International Search Report filed on 11/22/2009
- Oath or Declaration filed on 11/22/2009
- U.S. Basic National Fees filed on 11/22/2009
- Specification filed on 11/22/2009
- Claims filed on 11/22/2009
- Abstracts filed on 11/22/2009
- Drawings filed on 11/22/2009

The applicant needs to satisfy supplemental fees problems indicated below.

The following items **MUST** be furnished within the period set forth below in order to complete the requirements for acceptance under 35 U.S.C. 371:

- Oath or declaration of the inventors, in compliance with 37 CFR 1.497(a) and (b), identifying the application by the International application number and international filing date.
- Oath or declaration of the inventors, in compliance with 37 CFR 1.497(a) and (b), identifying the application by the International application number and international filing date. The current oath or declaration does not comply with 37 CFR 1.497(a) and (b) in that it:

• THE PTO/SB/01A OATH REQUIRES AN APPLICATION DATA SHEET.(37 CFR 1.76)

• To avoid abandonment, a surcharge (for late submission of filing fee, search fee, examination fee or oath or declaration) as set forth in 37 CFR 1.492(h) of \$65 for a small entity in compliance with 37 CFR 1.27, must be submitted with the missing items identified in this letter.

SUMMARY OF FEES DUE:

Total additional fees required for this application is **\$225** for a Small Entity:

• \$65 Surcharge.

• The application search fee has not been paid. Applicant must submit \$50 to complete the search fee. Note a surcharge will be required if submitted later than commencement of the national stage (37 CFR 1.492(h)) and the basic national fee was not paid before July 1, 2005.

• The application examination fee has not been paid. Applicant must submit \$110 to complete the examination fee for a small entity in compliance with 37 CFR 1.27. Note a surcharge will be required if submitted later than commencement of the national stage (37 CFR 1.492(h)) and the basic national fee was not paid before July 1, 2005.

ALL OF THE ITEMS SET FORTH ABOVE MUST BE SUBMITTED WITHIN TWO (2) MONTHS FROM THE DATE OF THIS NOTICE OR BY 32 MONTHS FROM THE PRIORITY DATE FOR THE APPLICATION, WHICHEVER IS LATER. FAILURE TO PROPERLY RESPOND WILL RESULT IN ABANDONMENT.

The time period set above may be extended by filing a petition and fee for extension of time under the provisions of 37 CFR 1.136(a).

Applicant is reminded that any communications to the United States Patent and Trademark Office must be mailed to the address given in the heading and include the U.S. application no. shown above (37 CFR 1.5)

Registered users of EFS-Web may alternatively submit their reply to this notice via EFS-Web. <u>https://sportal.uspto.gov/authenticate/AuthenticateUserLocalEPF.html</u>

For more information about EFS-Web please call the USPTO Electronic Business Center at **1-866-217-9197** or visit our website at <u>http://www.uspto.gov/ebc.</u>

If you are not using EFS-Web to submit your reply, you must include a copy of this notice.

ANITA D JOHNSON

Telephone: (571) 272-0386

Electronic Patent Application Fee Transmittal						
Application Number:	126	01275				
Filing Date:						
Title of Invention:	VLS	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS				
First Named Inventor/Applicant Name:	Ven	kat Konda				
Filer:	Ven	kar Konda				
Attorney Docket Number:	V-00	045US				
Filed as Small Entity						
U.S. National Stage under 35 USC 371 Filing	Fees	;				
Description		Fee Code	Quantity	Amount	Sub-Total in USD(\$)	
Basic Filing:						
Basic National Stage Fee		2631	1	165	165	
Pages:						
Natl Stage Appl Sz fee per 50 pgs >100		2681	1	135	135	
Claims:						
Claims in excess of 20		2615	29	26	754	
Miscellaneous-Filing:						
Petition:						
Patent-Appeals-and-Interference:						

Description	Fee Code	Quantity	Amount	Sub-Total in USD(\$)
Post-Allowance-and-Post-Issuance:				
Extension-of-Time:				
Miscellaneous:				
	Tot	al in USD	(\$)	1054

Electronic A	Electronic Acknowledgement Receipt						
EFS ID:	6694837						
Application Number:	12601275						
International Application Number:							
Confirmation Number:	6372						
Title of Invention:	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS						
First Named Inventor/Applicant Name:	Venkat Konda						
Customer Number:	38139						
Filer:	Venkar Konda						
Filer Authorized By:							
Attorney Docket Number:	V-0045US						
Receipt Date:	22-DEC-2009						
Filing Date:							
Time Stamp:	19:05:08						
Application Type:	U.S. National Stage under 35 USC 371						

Payment information:

Document Number Page 150 of	Document Description	File Name	File Size(Bytes)/ Message Digest	Multi Part /.zip	Pages (if appl.)				
File Listing	:								
Authorized Use	r								
Deposit Accour	ıt								
RAM confirmati	on Number	6404							
Payment was su	uccessfully received in RAM	\$1054							
Payment Type		Credit Card	Credit Card						
Submitted with	Payment	yes	yes						

1	Fee Worksheet (PTO-875) fee-info.pdf		33527 86105945f3d107d6a07e8ca63ed220c3f093 31a7	no	2
Warnings:					
Information:					
		Total Files Size (in bytes)): 335	527	
New Applicat	described in MPEP 503. <u>ions Under 35 U.S.C. 111</u> ication is being filed and the applica	tion includes the necessary o	components for a filing	date (see	37 CFR
<u>New Applicat</u> If a new appli 1.53(b)-(d) ar Acknowledge <u>National Stac</u> If a timely sul U.S.C. 371 an		R 1.54) will be issued in due g date of the application. <u>ader 35 U.S.C. 371</u> of an international applicat orm PCT/DO/EO/903 indicat	course and the date sh ion is compliant with th ing acceptance of the a	own on th ne conditio pplication	is ons of 35

PTO/SB/05 (08-08) Approved for use through 06/30/2010. OMB 0651-0032 LLC Detent and

	Attorney Docket No	V-0045 US
Under the Paperwork Reduction Act of 1995, no persons are required to re-	spond to a collection of informa	ation unless it displays a valid OMB control number
	U.S. Patent and Trade	Mark Office. U.S. DEPARTMENT OF COMMERCE

(UTILITY		Attorney Docket No.	V-0045 US		
PA	TENT APPLICATION		First Inventor	Venkat Konda	1	
	TRANSMITTAL		Title			
(Only for new	nonprovisional applications under 37 CFR	1.53(b))	Express Mail Label No.			
	PPLICATION ELEMENTS	n contents.	ADDRESS TO:	P.O. Box 14	ner for Patents 50 /A 22313-1450	
1. 🖌 Fee Trans	mittal Form (e.g., PTO/SB/17)		ACCOMPAN	IYING APPL	ICATION PARTS	
2. Applicant See 37 CF	claims small entity status. FR 1 27		9. 🛛 Assignment F	apers (cover sl	neet & document(s))	
3. Specificat			Name of Ass	ignee <u>Konda T</u>	echnologies Inc.	
(For informati	ion on the preferred arrangement, see MPEP 608. s) (35 U.S.C. 113) [Total Sheets	^{01(a))} 39]				
	executed (original or copy)		10. 37 CFR 3.73(b (when there) Statement is an assignee)	Power of Attorney	
(for col	from a prior application (37 CFR 1.63 ntinuation/divisional with Box 18 comp		11. 🔲 English Trans	lation Docume	nt (if applicable)	
 DELETION OF INVENTOR(S) Signed statement attached deleting inventor(s) name in the prior application, see 37 CFR 1.63(d)(2) and 1.33(b). 				isclosure State s of citations atta	e ment (PTO/SB/08 or PTO-1449) ached	
6. 🖌 Applicatio	on Data Sheet. See 37 CFR 1.76		13. Preliminary Amendment			
<u>Compute</u>	or CD-R in duplicate, large table or r Program <i>(Appendix)</i> scape Table on CD		14. Return Receipt Postcard (MPEP 503) (Should be specifically itemized)			
8. Nucleotide an	nd/or Amino Acid Sequence Submis items a. – c. are required)	ssion	15. Certified Copy of Priority Document(s) (if foreign priority is claimed)			
a. 🗌 Con	nputer Readable Form (CRF) ecification Sequence Listing on:		16. Nonpublication Request under 35 U.S.C. 122(b)(2)(B)(i). Applicant must attach form PTO/SB/35 or equivalent.			
i. 🗌 ii. 🗌	CD-ROM or CD-R (2 copies); or Paper		17. Other:			
c. 🔲 Sta	tements verifying identity of above co	pies				
	ING APPLICATION, check appropriate ing the title, or in an Application Data			on below and in	the first sentence of the	
Continuati	ion Divisional	Continua	tion-in-part (CIP) of p	rior application No	.:	
Prior application infor	mation: Examiner		Art U	nit:		
	19. C	ORRESPON	DENCE ADDRESS			
The address as	ssociated with Customer Number:	381	39	OR Corr	espondence address below	
Name Ver	nkat Konda					
Address 627	78 Grand Oak Way					
	n Jose		CA	Zip Code	95135	
Country US	A	Telephone	408-472-3273	Email	venkat@kondatech.com	
Signature	/Venkat Konda/		Da	11/22/2003		
Name (Print/Type)	Venkat Konda			Registratio		

This collection of information is required by 37 CFR 1.53(b). The information is required to obtain or retain a benefit by the public which is to file (and by the USPTO to process) an application. Confidentiality is governed by 35 U.S.C. 122 and 37 CFR 1.11 and 1.14. This collection is estimated to take 12 minutes to complete, including gathering, preparing, and submitting the completed application form to the USPTO. Time will vary depending upon the individual case. Any comments on the amount of time you require to complete this form and/or suggestions for reducing this burden, should be sent to the Chief Information Officer, U.S. Patent and Trademark Office, U.S. Department of Commerce, P.O. Box 1450, Alexandria, VA 22313-1450. DO NOT SEND FEES OR COMPLETED FORMS TO THIS ADDRESS. SEND TO: Commissioner for Patents, P.O. Box 1450, Alexandria, VA 22313-1450.

If you need assistance in completing the form, call 1-800-PTO-9199 and select option 2.

Privacy Act Statement

The **Privacy Act of 1974 (P.L. 93-579)** requires that you be given certain information in connection with your submission of the attached form related to a patent application or patent. Accordingly, pursuant to the requirements of the Act, please be advised that: (1) the general authority for the collection of this information is 35 U.S.C. 2(b)(2); (2) furnishing of the information solicited is voluntary; and (3) the principal purpose for which the information is used by the U.S. Patent and Trademark Office is to process and/or examine your submission related to a patent application or patent. If you do not furnish the requested information, the U.S. Patent and Trademark Office may not be able to process and/or examine your submission, which may result in termination of proceedings or abandonment of the application or expiration of the patent.

The information provided by you in this form will be subject to the following routine uses:

- 1. The information on this form will be treated confidentially to the extent allowed under the Freedom of Information Act (5 U.S.C. 552) and the Privacy Act (5 U.S.C 552a). Records from this system of records may be disclosed to the Department of Justice to determine whether disclosure of these records is required by the Freedom of Information Act.
- 2. A record from this system of records may be disclosed, as a routine use, in the course of presenting evidence to a court, magistrate, or administrative tribunal, including disclosures to opposing counsel in the course of settlement negotiations.
- 3. A record in this system of records may be disclosed, as a routine use, to a Member of Congress submitting a request involving an individual, to whom the record pertains, when the individual has requested assistance from the Member with respect to the subject matter of the record.
- 4. A record in this system of records may be disclosed, as a routine use, to a contractor of the Agency having need for the information in order to perform a contract. Recipients of information shall be required to comply with the requirements of the Privacy Act of 1974, as amended, pursuant to 5 U.S.C. 552a(m).
- 5. A record related to an International Application filed under the Patent Cooperation Treaty in this system of records may be disclosed, as a routine use, to the International Bureau of the World Intellectual Property Organization, pursuant to the Patent Cooperation Treaty.
- 6. A record in this system of records may be disclosed, as a routine use, to another federal agency for purposes of National Security review (35 U.S.C. 181) and for review pursuant to the Atomic Energy Act (42 U.S.C. 218(c)).
- 7. A record from this system of records may be disclosed, as a routine use, to the Administrator, General Services, or his/her designee, during an inspection of records conducted by GSA as part of that agency's responsibility to recommend improvements in records management practices and programs, under authority of 44 U.S.C. 2904 and 2906. Such disclosure shall be made in accordance with the GSA regulations governing inspection of records for this purpose, and any other relevant (*i.e.*, GSA or Commerce) directive. Such disclosure shall not be used to make determinations about individuals.
- 8. A record from this system of records may be disclosed, as a routine use, to the public after either publication of the application pursuant to 35 U.S.C. 122(b) or issuance of a patent pursuant to 35 U.S.C. 151. Further, a record may be disclosed, subject to the limitations of 37 CFR 1.14, as a routine use, to the public if the record was filed in an application which became abandoned or in which the proceedings were terminated and which application is referenced by either a published application, an application open to public inspection or an issued patent.
- 9. A record from this system of records may be disclosed, as a routine use, to a Federal, State, or local law enforcement agency, if the USPTO becomes aware of a violation or potential violation of law or regulation.

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DEC	CLARATION (37 CFR 1.63) FOR UTILITY OR DESIGN APPLICATION DATA SHEET (37 CFR				
Title of Invention	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZ	ZED NETWORKS			
As the belo	low named inventor(s), I/we declare that:				
This declara	aration is directed to:				
	The attached application, or Application No	(if applicable);			
l/we believe sought;	eve that I/we am/are the original and first inventor(s) of the subject matter v	which is claimed and for which a patent is			
	e reviewed and understand the contents of the above-identified application, ent specifically referred to above;	including the claims, as amended by any			
I/we acknowledge the duty to disclose to the United States Patent and Trademark Office all information known to me/us to be material to patentability as defined in 37 CFR 1.56, including for continuation-in-part applications, material information which became available between the filing date of the prior application and the national or PCT International filing date of the continuation-in-part application.					
contribute in numbers (of the USPTO, per to the USP of the appli of a patent referenced PTO-2038 is All statement believed to are punish	Applicant is cautioned to avoid submitting personal information in docume to identity theft. Personal information such as social security numbers (other than a check or credit card authorization form PTO-2038 submitted for TO to support a petition or an application. If this type of personal information petitioners/applicants should consider redacting such personal information for SPTO. Petitioner/applicant is advised that the record of a patent application polication (unless a non-publication request in compliance with 37 CFR 1.2130 nt. Furthermore, the record from an abandoned application may also be a ed in a published application or an issued patent (see 37 CFR 1.14). Che 8 submitted for payment purposes are not retained in the application file and nents made herein of my/our own knowledge are true, all statements mate to be true, and further that these statements were made with the knowledg shable by fine or imprisonment, or both, under 18 U.S.C. 1001, and may jeop suing thereon.	s, bank account numbers, or credit card or payment purposes) is never required by is included in documents submitted to the om the documents before submitting them is available to the public after publication (a) is made in the application) or issuance ivailable to the public if the application is ecks and credit card authorization forms therefore are not publicly available.			
FULL NAM	ME OF INVENTOR(S)				
Inventor on	one: Venkat KondaDate:	11/22/2009			
Signature:	: /Venkat Konda/ Citize	n of: USA			
Inventor two	wo:Date:				
Signature:	:Citize	n of:			
Additio	itional inventors or a legal representative are being named on	additional form(s) attached hereto.			

This collection of information is required by 35 U.S.C. 115 and 37 CFR 1.63. The information is required to obtain or retain a benefit by the public which is to file (and by the USPTO to process) an application. Confidentiality is governed by 35 U.S.C. 122 and 37 CFR 1.11 and 1.14. This collection is estimated to take 1 minute to complete, including gathering, preparing, and submitting the completed application form to the USPTO. Time will vary depending upon the individual case. Any comments on the amount of time you require to complete this form and/or suggestions for reducing this burden, should be sent to the Chief Information Officer, U.S. Patent and Trademark Office, U.S. Department of Commerce, P.O. Box 1450, Alexandria, VA 22313-1450. DO NOT SEND FEES OR COMPLETED FORMS TO THIS ADDRESS. **SEND TO: Commissioner for Patents, P.O. Box 1450, Alexandria, VA 22313-1450.**

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- 8. A record from this system of records may be disclosed, as a routine use, to the public after either publication of the application pursuant to 35 U.S.C. 122(b) or issuance of a patent pursuant to 35 U.S.C. 151. Further, a record may be disclosed, subject to the limitations of 37 CFR 1.14, as a routine use, to the public if the record was filed in an application which became abandoned or in which the proceedings were terminated and which application is referenced by either a published application, an application open to public inspection or an issued patent.
- 9. A record from this system of records may be disclosed, as a routine use, to a Federal, State, or local law enforcement agency, if the USPTO becomes aware of a violation or potential violation of law or regulation.

Electronic Patent Application Fee Transmittal						
Application Number:						
Filing Date:						
Title of Invention:	VLSI	VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS				
First Named Inventor/Applicant Name:	Venl	kat Konda				
Filer:	Venl	kar Konda				
Attorney Docket Number:	V-00)45US				
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Claims:						
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First Named Inventor/Applicant Name:	Venkat Konda					
Customer Number:	38139					
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1	Specification		V-0045US.pdf	299380	no	69	
				debb4d8182a1df928e5922e1c44142090bb a7555			
Warnings:				· · · ·			
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Drawings-only black and white line drawings	V-0045US-FIGs.pdf		no	39
		9345637693e597fd4d92e2aa679baa691e9 5ca78		
Transmittal of New Application	sb0005_fill.pdf	468829	no	2
		959ba37f3439c0e68be3c2ba5915c1a196e 60d2f		
Application Data Sheet	sb0001a.pdf	545678	no	2
		6fdc91490275bb543e412eafc4cd318937bc 0281		
PTO supplied ADS fillable form				
Fee Worksheet (PTO-875)	fee-info.pdf	36986	no	2
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by the applicant, and including pag lescribed in MPEP 503. <u>ons Under 35 U.S.C. 111</u> ation is being filed and the applicat MPEP 506), a Filing Receipt (37 CF ment Receipt will establish the filing <u>e of an International Application un</u> mission to enter the national stage other applicable requirements a Fo submission under 35 U.S.C. 371 will <u>onal Application Filed with the USP</u> ational application is being filed an al filing date (see PCT Article 11 and ernational Filing Date (Form PCT/RC	tion includes the necessary of R 1.54) will be issued in due of g date of the application. <u>der 35 U.S.C. 371</u> of an international application orm PCT/DO/EO/903 indication be issued in addition to the <u>TO as a Receiving Office</u> and the international application of MPEP 1810), a Notification D/105) will be issued in due co	It serves as evidence omponents for a filin course and the date s on is compliant with ng acceptance of the Filing Receipt, in du ion includes the nece of the International <i>J</i> ourse, subject to pres	of receipt s of ate (see hown on th the condition application e course. ssary comp Application scriptions co	similar to a 37 CFR is ons of 35 n as a onents for Number oncerning
	Transmittal of New Application Application Data Sheet To supplied ADS fillable form Fee Worksheet (PTO-875) dgement Receipt evidences receipt by the applicant, and including pag escribed in MPEP 503. Dons Under 35 U.S.C. 111 ation is being filed and the applicat IMPEP 506), a Filing Receipt (37 CF nent Receipt will establish the filing of an International Application un nission to enter the national stage other applicable requirements a Fo submission under 35 U.S.C. 371 wil mal Application Filed with the USP ational application is being filed an al filing date (see PCT Article 11 and rnational Filing Date (Form PCT/RC ity, and the date shown on this Ack	Transmittal of New Application sb0005_fill.pdf Application Data Sheet sb0001a.pdf TO supplied ADS fillable form	Transmittal of New Application sb0005_fill.pdf 468829 Transmittal of New Application sb0005_fill.pdf 468829 Application Data Sheet sb0001a.pdf 545678 Georemotive State	Transmittal of New Application sb0005_fill.pdf 468829 no Application Data Sheet sb0001a.pdf 468829 no TO supplied ADS fillable form 669700000000000000000000000000000000000

VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS

Venkat Konda

5 CROSS REFERENCE TO RELATED APPLICATIONS

This application is related to and claims priority of the PCT Application Serial No. PCT/US08/64605 entitled "VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 22, 2008, and the U.S. Provisional Patent Application

10 Serial No. 60/940, 394 entitled "VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007.

This application is related to and incorporates by reference in its entirety the US Application Serial No. 12/530,207 entitled "FULLY CONNECTED GENERALIZED

- 15 MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed September 6, 2009, the PCT Application Serial No. PCT/US08/56064 entitled "FULLY CONNECTED GENERALIZED MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed March 6, 2008, the U.S. Provisional Patent Application Serial No. 60/905,526
- 20 entitled "LARGE SCALE CROSSPOINT REDUCTION WITH NONBLOCKING UNICAST & MULTICAST IN ARBITRARILY LARGE MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed March 6, 2007, and the U.S. Provisional Patent Application Serial No. 60/940, 383 entitled "FULLY CONNECTED GENERALIZED MULTI-STAGE NETWORKS" by
- 25 Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007.

This application is related to and incorporates by reference in its entirety the US Patent Application Docket No. V-0038US entitled "FULLY CONNECTED GENERALIZED BUTTERFLY FAT TREE NETWORKS" by Venkat Konda assigned to the same assignee as the current application filed concurrently, the PCT Application

- 5 Serial No. PCT/US08/64603 entitled "FULLY CONNECTED GENERALIZED BUTTERFLY FAT TREE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 22, 2008, the U.S. Provisional Patent Application Serial No. 60/940, 387 entitled "FULLY CONNECTED GENERALIZED BUTTERFLY FAT TREE NETWORKS" by Venkat Konda assigned to the same
- 10 assignee as the current application, filed May 25, 2007, and the U.S. Provisional Patent Application Serial No. 60/940, 390 entitled "FULLY CONNECTED GENERALIZED MULTI-LINK BUTTERFLY FAT TREE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007.

This application is related to and incorporates by reference in its entirety the US Patent Application Docket No. V-0039US entitled "FULLY CONNECTED GENERALIZED MULTI-LINK MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application filed concurrently, the PCT Application Serial No. PCT/US08/64604 entitled "FULLY CONNECTED GENERALIZED MULTI-LINK MULTI-STAGE NETWORKS" by Venkat Konda

- 20 assigned to the same assignee as the current application, filed May 22, 2008, the U.S. Provisional Patent Application Serial No. 60/940, 389 entitled "FULLY CONNECTED GENERALIZED REARRANGEABLY NONBLOCKING MULTI-LINK MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007, the U.S. Provisional Patent Application Serial No.
- 25 60/940, 391 entitled "FULLY CONNECTED GENERALIZED FOLDED MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007 and the U.S. Provisional Patent Application Serial No. 60/940, 392 entitled "FULLY CONNECTED GENERALIZED STRICTLY NONBLOCKING MULTI-LINK MULTI-STAGE NETWORKS" by Venkat Konda
- 30 assigned to the same assignee as the current application, filed May 25, 2007.

This application is related to and incorporates by reference in its entirety the U.S. Provisional Patent Application Serial No. 61/252, 603 entitled "VLSI LAYOUTS OF FULLY CONNECTED NETWORKS WITH LOCALITY EXPLOITATION" by Venkat Konda assigned to the same assignee as the current application, filed October 16, 2009.

5 This application is related to and incorporates by reference in its entirety the U.S. Provisional Patent Application Serial No. 61/252, 609 entitled "VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED AND PYRAMID NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed October 16, 2009.

10 BACKGROUND OF INVENTION

Multi-stage interconnection networks such as Benes networks and butterfly fat tree networks are widely useful in telecommunications, parallel and distributed computing. However VLSI layouts, known in the prior art, of these interconnection networks in an integrated circuit are inefficient and complicated.

15 Other multi-stage interconnection networks including butterfly fat tree networks, Banyan networks, Batcher-Banyan networks, Baseline networks, Delta networks, Omega networks and Flip networks have been widely studied particularly for self routing packet switching applications. Also Benes Networks with radix of two have been widely studied and it is known that Benes Networks of radix two are shown to be built with back to back 20 baseline networks which are rearrangeably nonblocking for unicast connections.

The most commonly used VLSI layout in an integrated circuit is based on a twodimensional grid model comprising only horizontal and vertical tracks. An intuitive interconnection network that utilizes two-dimensional grid model is 2D Mesh Network and its variations such as segmented mesh networks. Hence routing networks used in

25 VLSI layouts are typically 2D mesh networks and its variations. However Mesh Networks require large scale cross points typically with a growth rate of $O(N^2)$ where N is the number of computing elements, ports, or logic elements depending on the application. V-0045 US

Multi-stage interconnection with a growth rate of $O(N \times \log N)$ requires significantly small number of cross points. U.S. Patent 6,185,220 entitled "Grid Layouts of Switching and Sorting Networks" granted to Muthukrishnan et al. describes a VLSI layout using existing VLSI grid model for Benes and Butterfly networks. U.S. Patent

- 5 6,940,308 entitled "Interconnection Network for a Field Programmable Gate Array" granted to Wong describes a VLSI layout where switches belonging to lower stage of Benes Network are layed out close to the logic cells and switches belonging to higher stages are layed out towards the center of the layout.
- Due to the inefficient and in some cases impractical VLSI layout of Benes and butterfly fat tree networks on a semiconductor chip, today mesh networks and segmented mesh networks are widely used in the practical applications such as field programmable gate arrays (FPGAs), programmable logic devices (PLDs), and parallel computing interconnects. The prior art VLSI layouts of Benes and butterfly fat tree networks and VLSI layouts of mesh networks and segmented mesh networks require large area to
- 15 implement the switches on the chip, large number of wires, longer wires, with increased power consumption, increased latency of the signals which effect the maximum clock speed of operation. Some networks may not even be implemented practically on a chip due to the lack of efficient layouts.

20 SUMMARY OF INVENTION

When large scale sub-integrated circuit blocks with inlet and outlet links are layed out in an integrated circuit device in a two-dimensional grid arrangement, (for example in an FPGA where the sub-integrated circuit blocks are Lookup Tables) the most intuitive routing network is a network that uses horizontal and vertical links only (the most often used such a network is one of the variations of a 2D Mesh network). A direct embedding of a generalized multi-stage network on to a 2D Mesh network is neither simple nor

efficient.

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In accordance with the invention, VLSI layouts of generalized multi-stage networks for broadcast, unicast and multicast connections are presented using only

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horizontal and vertical links. The VLSI layouts employ shuffle exchange links where outlet links of cross links from switches in a stage in one sub-integrated circuit block are connected to inlet links of switches in the succeeding stage in another sub-integrated circuit block so that said cross links are either vertical links or horizontal and vice versa.

5 In one embodiment the sub-integrated circuit blocks are arranged in a hypercube arrangement in a two-dimensional plane. The VLSI layouts exploit the benefits of significantly lower cross points, lower signal latency, lower power and full connectivity with significantly fast compilation.

The VLSI layouts presented are applicable to generalized multi-stage networks 10 $V(N_1, N_2, d, s)$, generalized folded multi-stage networks $V_{fold}(N_1, N_2, d, s)$, generalized butterfly fat tree networks $V_{bft}(N_1, N_2, d, s)$, generalized multi-link multi-stage networks $V_{mlink}(N_1, N_2, d, s)$, generalized folded multi-link multi-stage networks $V_{fold-mlink}(N_1, N_2, d, s)$, generalized multi-link butterfly fat tree networks $V_{mlink-bft}(N_1, N_2, d, s)$, and generalized hypercube networks $V_{hcube}(N_1, N_2, d, s)$ for s =

15 1,2,3 or any number in general. The embodiments of VLSI layouts are useful in wide target applications such as FPGAs, CPLDs, pSoCs, ASIC placement and route tools, networking applications, parallel & distributed computing, and reconfigurable computing.

BRIEF DESCRIPTION OF DRAWINGS

- FIG. 1A is a diagram 100A of an exemplary symmetrical multi-link multi-stage network $V_{fold-mlink}(N,d,s)$ having inverse Benes connection topology of nine stages with N = 32, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention.
- FIG. 1B is a diagram 100B of the equivalent symmetrical folded multi-link multistage network $V_{fold-mlink}(N,d,s)$ of the network 100A shown in FIG. 1A, having inverse Benes connection topology of five stages with N = 32, d = 2 and s=2, strictly nonblocking

network for unicast connections and rearrangeably nonblocking network for arbitrary fanout multicast connections, in accordance with the invention.

FIG. 1C is a diagram 100C layout of the network V_{fold-mlink} (N, d, s) shown in FIG.
1B, in one embodiment, illustrating the connection links belonging with in each block
5 only.

FIG. 1D is a diagram 100D layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 1B, in one embodiment, illustrating the connection links ML(1,i) for i = [1, 64] and ML(8,i) for i = [1, 64].

FIG. 1E is a diagram 100E layout of the network V_{fold-mlink}(N,d,s) shown in FIG.
10 1B, in one embodiment, illustrating the connection links ML(2,i) for i = [1, 64] and ML(7,i) for i = [1,64].

FIG. 1F is a diagram 100F layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 1B, in one embodiment, illustrating the connection links ML(3,i) for i = [1, 64] and ML(6,i) for i = [1,64].

15 FIG. 1G is a diagram 100G layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 1B, in one embodiment, illustrating the connection links ML(4,i) for i = [1, 64] and ML(5,i) for i = [1,64].

FIG. 1H is a diagram 100H layout of a network $V_{fold-mlink}(N,d,s)$ where N = 128, d = 2, and s = 2, in one embodiment, illustrating the connection links belonging with in 20 each block only.

FIG. 1I is a diagram 100I detailed connections of BLOCK 1_2 in the network layout 100C in one embodiment, illustrating the connection links going in and coming out when the layout 100C is implementing V(N,d,s) or $V_{fold}(N,d,s)$.

FIG. 1J is a diagram 100J detailed connections of BLOCK 1_2 in the network layout 100C in one embodiment, illustrating the connection links going in and coming out when the layout 100C is implementing V(N,d,s) or $V_{fold}(N,d,s)$.

FIG. 1K is a diagram 100K detailed connections of BLOCK 1_2 in the network 5 layout 100C in one embodiment, illustrating the connection links going in and coming out when the layout 100C is implementing V(N,d,s) or $V_{fold}(N,d,s)$.

FIG. 1K1 is a diagram 100M1 detailed connections of BLOCK 1_2 in the network layout 100C in one embodiment, illustrating the connection links going in and coming out when the layout 100C is implementing V(N,d,s) or $V_{fold}(N,d,s)$ for s = 1.

10 FIG. 1L is a diagram 100L detailed connections of BLOCK 1_2 in the network layout 100C in one embodiment, illustrating the connection links going in and coming out when the layout 100C is implementing V(N,d,s) or $V_{fold}(N,d,s)$.

FIG. 1L1 is a diagram 100L1 detailed connections of BLOCK 1_2 in the network layout 100C in one embodiment, illustrating the connection links going in and coming
out when the layout 100C is implementing V(N,d,s) or V_{fold}(N,d,s) for s = 1.

FIG. 2A1 is a diagram 200A1 of an exemplary symmetrical multi-link multi-stage network $V_{fold-mlink}(N,d,s)$ having inverse Benes connection topology of one stage with N = 2, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention. FIG. 2A2 is a diagram 200A2 of the equivalent symmetrical folded multi-link multi-stage network $V_{fold-mlink}(N,d,s)$ of the network 200A1 shown in FIG. 2A1, having inverse Benes connection topology of one stage with N = 2, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in 25 accordance with the invention. FIG. 2A3 is a diagram 200A3 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2A2, in one embodiment, illustrating all the connection links,

FIG. 2B1 is a diagram 200B1 of an exemplary symmetrical multi-link multi-stage network $V_{fold-mlink}(N,d,s)$ having inverse Benes connection topology of one stage with N = 4, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in 5 accordance with the invention. FIG. 2B2 is a diagram 200B2 of the equivalent symmetrical folded multi-link multi-stage network $V_{fold-mlink}(N,d,s)$ of the network 200B1 shown in FIG. 2B1, having inverse Benes connection topology of one stage with N = 4, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in 10 accordance with the invention. FIG. 2B3 is a diagram 200B3 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2B2, in one embodiment, illustrating the connection links belonging with in each block only. FIG. 2B4 is a diagram 200B4 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2B2, in one embodiment, illustrating the connection links ML(1,i) for i = [1, 8] and ML(2,i) for i = [1,8].

- 15 FIG. 2C11 is a diagram 200C11 of an exemplary symmetrical multi-link multistage network $V_{fold-mlink}(N,d,s)$ having inverse Benes connection topology of one stage with N = 8, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention. FIG. 2C12 is a diagram 200C12 of the equivalent 20 symmetrical folded multi-link multi-stage network $V_{fold-mlink}(N,d,s)$ of the network 200C11 shown in FIG. 2C11, having inverse Benes connection topology of one stage with N = 8, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention.
- FIG. 2C21 is a diagram 200C21 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2C12, in one embodiment, illustrating the connection links belonging with in each block only. FIG. 2C22 is a diagram 200C22 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2C12, in one embodiment, illustrating the connection links ML(1,i) for i =

[1, 16] and ML(4,i) for i = [1,16]. FIG. 2C23 is a diagram 200C23 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2C12, in one embodiment, illustrating the connection links ML(2,i) for i = [1, 16] and ML(3,i) for i = [1,16].

FIG. 2D1 is a diagram 200D1 of an exemplary symmetrical multi-link multi-stage 5 network $V_{fold-mlink}(N,d,s)$ having inverse Benes connection topology of one stage with N = 16, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention.

FIG. 2D2 is a diagram 200D2 of the equivalent symmetrical folded multi-link 10 multi-stage network $V_{fold-mlink}(N,d,s)$ of the network 200D1 shown in FIG. 2D1, having inverse Benes connection topology of one stage with N = 16, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention.

FIG. 2D3 is a diagram 200D3 layout of the network $V_{fold-mlink}(N,d,s)$ shown in

15 FIG. 2D2, in one embodiment, illustrating the connection links belonging with in each block only.

FIG. 2D4 is a diagram 200D4 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2D2, in one embodiment, illustrating the connection links ML(1,i) for i = [1, 32] and ML(6,i) for i = [1,32].

20 FIG. 2D5 is a diagram 200D5 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2D2, in one embodiment, illustrating the connection links ML(2,i) for i = [1, 32] and ML(5,i) for i = [1,32].

FIG. 2D6 is a diagram 200D6 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2D2, in one embodiment, illustrating the connection links ML(3,i) for i = [1, 32] and ML(4,i) for i = [1,32].

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FIG. 3A is a diagram 300A of an exemplary symmetrical multi-link multi-stage network $V_{hcube}(N,d,s)$ having inverse Benes connection topology of nine stages with N = 32, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention.

FIG. 3B is a diagram 300B of the equivalent symmetrical folded multi-link multistage network $V_{hcube}(N,d,s)$ of the network 300A shown in FIG. 3A, having inverse Benes connection topology of five stages with N = 32, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fanout multicast connections, in accordance with the invention.

FIG. 3C is a diagram 300C layout of the network $V_{hcube}(N,d,s)$ shown in FIG. 3B, in one embodiment, illustrating the connection links belonging with in each block only.

FIG. 3D is a diagram 100D layout of the network V_{hcube}(N,d,s) shown in FIG.
3B, in one embodiment, illustrating the connection links ML(1,i) for i = [1, 64] and ML(8,i) for i = [1,64].

FIG. 3E is a diagram 300E layout of the network $V_{hcube}(N,d,s)$ shown in FIG. 3B, in one embodiment, illustrating the connection links ML(2,i) for i = [1, 64] and ML(7,i) for i = [1,64].

FIG. 3F is a diagram 300F layout of the network $V_{hcube}(N,d,s)$ shown in FIG. 3B, in one embodiment, illustrating the connection links ML(3,i) for i = [1, 64] and ML(6,i) for i = [1,64].

FIG. 3G is a diagram 300G layout of the network $V_{hcube}(N,d,s)$ shown in FIG. 3B, in one embodiment, illustrating the connection links ML(4,i) for i = [1, 64] and 25 ML(5,i) for i = [1,64]. FIG. 3H is a diagram 300H layout of a network $V_{hcube}(N, d, s)$ where N = 128, d = 2, and s = 2, in one embodiment, illustrating the connection links belonging with in each block only.

FIG. 4A is a diagram 400A layout of the network V_{fold-mlink}(N,d,s) shown in
5 FIG. 1B, in one embodiment, illustrating the connection links belonging with in each block only.

FIG. 4B is a diagram 400B layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 1B, in one embodiment, illustrating the connection links ML(1,i) for i = [1, 64] and ML(8,i) for i = [1,64].

10 FIG. 4C is a diagram 400C layout of the network $V_{fold-mlink}(N, d, s)$ shown in FIG. 4C, in one embodiment, illustrating the connection links ML(2,i) for i = [1, 64] and ML(7,i) for i = [1,64].

FIG. 4D is a diagram 400D layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 4D, in one embodiment, illustrating the connection links ML(3,i) for i = [1, 64] and ML(6,i) for i = [1,64].

FIG. 4E is a diagram 400E layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 4E, in one embodiment, illustrating the connection links ML(4,i) for i = [1, 64] and ML(5,i) for i = [1,64].

FIG. 4C1 is a diagram 400C1 layout of the network V_{fold-mlink}(N,d,s) shown in
FIG. 1B, in one embodiment, illustrating the connection links belonging with in each block only.

FIG. 5A1 is a diagram 500A1 of an exemplary prior art implementation of a two by two switch; FIG. 5A2 is a diagram 500A2 for programmable integrated circuit prior art implementation of the diagram 500A1 of FIG. 5A1; FIG. 5A3 is a diagram 500A3 for one-time programmable integrated circuit prior art implementation of the diagram 500A1

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of FIG. 5A1; FIG. 5A4 is a diagram 500A4 for integrated circuit placement and route implementation of the diagram 500A1 of FIG. 5A1.

DETAILED DESCRIPTION OF THE INVENTION

- 5 The present invention is concerned with the VLSI layouts of arbitrarily large switching networks for broadcast, unicast and multicast connections. Particularly switching networks considered in the current invention include: generalized multi-stage networks $V(N_1, N_2, d, s)$, generalized folded multi-stage networks $V_{fold}(N_1, N_2, d, s)$, generalized butterfly fat tree networks $V_{bft}(N_1, N_2, d, s)$, generalized multi-link multi-
- 10 stage networks $V_{mlink}(N_1, N_2, d, s)$, generalized folded multi-link multi-stage networks $V_{fold-mlink}(N_1, N_2, d, s)$, generalized multi-link butterfly fat tree networks $V_{mlink-bft}(N_1, N_2, d, s)$, and generalized hypercube networks $V_{hcube}(N_1, N_2, d, s)$ for s = 1,2,3 or any number in general.
- Efficient VLSI layout of networks on a semiconductor chip are very important and greatly influence many important design parameters such as the area taken up by the network on the chip, total number of wires, length of the wires, latency of the signals, capacitance and hence the maximum clock speed of operation. Some networks may not even be implemented practically on a chip due to the lack of efficient layouts. The different varieties of multi-stage networks described above have not been implemented
- 20 previously on the semiconductor chips efficiently. For example in Field Programmable Gate Array (FPGA) designs, multi-stage networks described in the current invention have not been successfully implemented primarily due to the lack of efficient VLSI layouts. Current commercial FPGA products such as Xilinx Vertex, Altera's Stratix implement island-style architecture using mesh and segmented mesh routing interconnects using
- 25 either full crossbars or sparse crossbars. These routing interconnects consume large silicon area for crosspoints, long wires, large signal propagation delay and hence consume lot of power.

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The current invention discloses the VLSI layouts of numerous types of multistage networks which are very efficient. Moreover they can be embedded on to mesh and segmented mesh routing interconnects of current commercial FPGA products. The VLSI layouts disclosed in the current invention are applicable to including the numerous generalized multi-stage networks disclosed in the following patent applications, filed concurrently:

1) Strictly and rearrangeably nonblocking for arbitrary fan-out multicast and unicast for generalized multi-stage networks $V(N_1, N_2, d, s)$ with numerous connection topologies and the scheduling methods are described in detail in the PCT Application Serial No. PCT/US08/56064 that is incorporated by reference above.

2) Strictly and rearrangeably nonblocking for arbitrary fan-out multicast and unicast for generalized butterfly fat tree networks $V_{bft}(N_1, N_2, d, s)$ with numerous connection topologies and the scheduling methods are described in detail in PCT Application Serial No. PCT/US08/64603 that is incorporated by reference above.

15 3) Rearrangeably nonblocking for arbitrary fan-out multicast and unicast, and strictly nonblocking for unicast for generalized multi-link multi-stage networks $V_{mlink}(N_1, N_2, d, s)$ and generalized folded multi-link multi-stage networks $V_{fold-mlink}(N_1, N_2, d, s)$ with numerous connection topologies and the scheduling methods are described in detail in PCT Application Serial No. PCT/US08/64604 that is 20 incorporated by reference above.

4) Strictly and rearrangeably nonblocking for arbitrary fan-out multicast and unicast for generalized multi-link butterfly fat tree networks $V_{mlink-bft}(N_1, N_2, d, s)$ with numerous connection topologies and the scheduling methods are described in detail in PCT Application Serial No. PCT/US08/64603 that is incorporated by reference above.

25 5) Strictly and rearrangeably nonblocking for arbitrary fan-out multicast and unicast for generalized folded multi-stage networks $V_{fold}(N_1, N_2, d, s)$ with numerous

connection topologies and the scheduling methods are described in detail in PCT Application Serial No. PCT/US08/64604 that is incorporated by reference above.

6) Strictly nonblocking for arbitrary fan-out multicast for generalized multi-link multi-stage networks $V_{mlink}(N_1, N_2, d, s)$ and generalized folded multi-link multi-stage

5 networks $V_{fold-mlink}(N_1, N_2, d, s)$ with numerous connection topologies and the scheduling methods are described in detail in PCT Application Serial No. PCT/US08/64604 that is incorporated by reference above.

7) VLSI layouts of numerous types of multi-stage networks with locality exploitation are described in U.S. Provisional Patent Application Serial No. 61/252, 603
10 that is incorporated by reference above.

8) VLSI layouts of numerous types of multistage pyramid networks are described in U.S. Provisional Patent Application Serial No. 61/252, 609 that is incorporated by reference above.

- In addition the layouts of the current invention are also applicable to generalized 15 multi-stage pyramid networks $V_p(N_1, N_2, d, s)$, generalized folded multi-stage pyramid networks $V_{fold-p}(N_1, N_2, d, s)$, generalized butterfly fat pyramid networks $V_{bfp}(N_1, N_2, d, s)$, generalized multi-link multi-stage pyramid networks $V_{mlink-p}(N_1, N_2, d, s)$, generalized folded multi-link multi-stage pyramid networks $V_{fold-mlink-p}(N_1, N_2, d, s)$, generalized multi-link butterfly fat pyramid networks
- 20 $V_{mlink-bfp}(N_1, N_2, d, s)$, and generalized hypercube networks $V_{hcube}(N_1, N_2, d, s)$ for s = 1,2,3 or any number in general.

Symmetric RNB generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$:

Referring to diagram 100A in FIG. 1A, in one embodiment, an exemplary

25 generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d =

2; and s = 2 with nine stages of one hundred and forty four switches for satisfying communication requests, such as setting up a telephone call or a data call, or a connection between configurable logic blocks, between an input stage 110 and output stage 120 via middle stages 130, 140, 150, 160, 170, 180 and 190 is shown where input stage 110

- 5 consists of sixteen, two by four switches IS1-IS16 and output stage 120 consists of sixteen, four by two switches OS1-OS16. And all the middle stages namely the middle stage 130 consists of sixteen, four by four switches MS(1,1) MS(1,16), middle stage 140 consists of sixteen, four by four switches MS(2,1) MS(2,16), middle stage 150 consists of sixteen, four by four switches MS(3,1) MS(3,16), middle stage 160 consists
- of sixteen, four by four switches MS(4,1) MS(4,16), middle stage 170 consists of sixteen, four by four switches MS(5,1) MS(5,16), middle stage 180 consists of sixteen, four by four switches MS(6,1) MS(6,16), and middle stage 190 consists of sixteen, four by four switches MS(7,1) MS(7,16).

As disclosed in PCT Application Serial No. PCT/US08/64604 that is

15 incorporated by reference above, such a network can be operated in rearrangeably nonblocking manner for arbitrary fan-out multicast connections and also can be operated in strictly non-blocking manner for unicast connections.

In one embodiment of this network each of the input switches IS1-IS4 and output switches OS1-OS4 are crossbar switches. The number of switches of input stage 110 and of output stage 120 can be denoted in general with the variable $\frac{N}{d}$, where N is the total number of inlet links or outlet links. The number of middle switches in each middle stage is denoted by $\frac{N}{d}$. The size of each input switch IS1-IS4 can be denoted in general with the notation d * 2d and each output switch OS1-OS4 can be denoted in general with the notation 2d * d. Likewise, the size of each switch in any of the middle stages can be

25 denoted as 2d * 2d. A switch as used herein can be either a crossbar switch, or a network of switches each of which in turn may be a crossbar switch or a network of switches. A symmetric multi-stage network can be represented with the notation $V_{mlink}(N, d, s)$, where N represents the total number of inlet links of all input switches

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(for example the links IL1-IL32), d represents the inlet links of each input switch or outlet links of each output switch, and s is the ratio of number of outgoing links from each input switch to the inlet links of each input switch.

Each of the $\frac{N}{d}$ input switches IS1 – IS16 are connected to exactly d switches in

- 5 middle stage 130 through two links each for a total of $2 \times d$ links (for example input switch IS1 is connected to middle switch MS(1,1) through the links ML(1,1), ML(1,2), and also connected to middle switch MS(1,2) through the links ML(1,3) and ML(1,4)). The middle links which connect switches in the same row in two successive middle stages are called hereinafter straight middle links; and the middle links which connect
- 10 switches in different rows in two successive middle stages are called hereinafter cross middle links. For example, the middle links ML(1,1) and ML(1,2) connect input switch IS1 and middle switch MS(1,1), so middle links ML(1,1) and ML(1,2) are straight middle links; where as the middle links ML(1,3) and ML(1,4) connect input switch IS1 and middle switch MS(1,2), since input switch IS1 and middle switch MS(1,2) belong to two
- 15 different rows in diagram 100A of FIG. 1A, middle links ML(1,3) and ML(1,4) are cross middle links.

Each of the
$$\frac{N}{d}$$
 middle switches MS(1,1) – MS(1,16) in the middle stage 130 are

connected from exactly d input switches through two links each for a total of $2 \times d$ links (for example the links ML(1,1) and ML(1,2) are connected to the middle switch MS(1,1)

- from input switch IS1, and the links ML(1,7) and ML(1,8) are connected to the middle switch MS(1,1) from input switch IS2) and also are connected to exactly d switches in middle stage 140 through two links each for a total of $2 \times d$ links (for example the links ML(2,1) and ML(2,2) are connected from middle switch MS(1,1) to middle switch MS(2,1), and the links ML(2,3) and ML(2,4) are connected from middle switch MS(1,1)
- 25 to middle switch MS(2,3)).

Each of the
$$\frac{N}{d}$$
 middle switches MS(2,1) – MS(2,16) in the middle stage 140 are connected from exactly d input switches through two links each for a total of $2 \times d$ links

(for example the links ML(2,1) and ML(2,2) are connected to the middle switch MS(2,1) from input switch MS(1,1), and the links ML(1,11) and ML(1,12) are connected to the middle switch MS(2,1) from input switch MS(1,3)) and also are connected to exactly d switches in middle stage 150 through two links each for a total of $2 \times d$ links (for

5 example the links ML(3,1) and ML(3,2) are connected from middle switch MS(2,1) to middle switch MS(3,1), and the links ML(3,3) and ML(3,4) are connected from middle switch MS(2,1) to middle switch MS(3,5)).

Each of the
$$\frac{N}{d}$$
 middle switches MS(3,1) – MS(3,16) in the middle stage 150 are

connected from exactly d input switches through two links each for a total of $2 \times d$ links

- 10 (for example the links ML(3,1) and ML(3,2) are connected to the middle switch MS(3,1) from input switch MS(2,1), and the links ML(2,19) and ML(2,20) are connected to the middle switch MS(3,1) from input switch MS(2,5)) and also are connected to exactly d switches in middle stage 160 through two links each for a total of $2 \times d$ links (for example the links ML(4,1) and ML(4,2) are connected from middle switch MS(3,1) to
- 15 middle switch MS(4,1), and the links ML(4,3) and ML(4,4) are connected from middle switch MS(3,1) to middle switch MS(4,9)).

Each of the
$$\frac{N}{d}$$
 middle switches MS(4,1) – MS(4,16) in the middle stage 160 are

connected from exactly d input switches through two links each for a total of $2 \times d$ links (for example the links ML(4,1) and ML(4,2) are connected to the middle switch MS(4,1)

from input switch MS(3,1), and the links ML(4,35) and ML(4,36) are connected to the middle switch MS(4,1) from input switch MS(3,9)) and also are connected to exactly *d* switches in middle stage 170 through two links each for a total of 2×*d* links (for example the links ML(5,1) and ML(5,2) are connected from middle switch MS(4,1) to middle switch MS(5,1), and the links ML(5,3) and ML(5,4) are connected from middle switch MS(4,1) to middle switch MS(4,1) to middle switch MS(5,9)).

Each of the
$$\frac{N}{d}$$
 middle switches MS(5,1) – MS(5,16) in the middle stage 170 are

connected from exactly d input switches through two links each for a total of $2 \times d$ links

(for example the links ML(5,1) and ML(5,2) are connected to the middle switch MS(5,1) from input switch MS(4,1), and the links ML(5,35) and ML(5,36) are connected to the middle switch MS(5,1) from input switch MS(4,9)) and also are connected to exactly d switches in middle stage 180 through two links each for a total of $2 \times d$ links (for

5 example the links ML(6,1) and ML(6,2) are connected from middle switch MS(5,1) to middle switch MS(6,1), and the links ML(6,3) and ML(6,4) are connected from middle switch MS(5,1) to middle switch MS(6,5)).

Each of the
$$\frac{N}{d}$$
 middle switches MS(6,1) – MS(6,16) in the middle stage 180 are

connected from exactly d input switches through two links each for a total of $2 \times d$ links

- 10 (for example the links ML(6,1) and ML(6,2) are connected to the middle switch MS(6,1) from input switch MS(5,1), and the links ML(6,19) and ML(6,20) are connected to the middle switch MS(6,1) from input switch MS(5,5)) and also are connected to exactly d switches in middle stage 190 through two links each for a total of $2 \times d$ links (for example the links ML(7,1) and ML(7,2) are connected from middle switch MS(6,1) to
- 15 middle switch MS(7,1), and the links ML(7,3) and ML(7,4) are connected from middle switch MS(6,1) to middle switch MS(7,3)).

Each of the
$$\frac{N}{d}$$
 middle switches MS(7,1) – MS(7,16) in the middle stage 190 are

connected from exactly d input switches through two links each for a total of $2 \times d$ links (for example the links ML(7,1) and ML(7,2) are connected to the middle switch MS(7,1)

- 20 from input switch MS(6,1), and the links ML(7,11) and ML(7,12) are connected to the middle switch MS(7,1) from input switch MS(6,3)) and also are connected to exactly d switches in middle stage 120 through two links each for a total of 2×d links (for example the links ML(8,1) and ML(8,2) are connected from middle switch MS(7,1) to middle switch MS(8,1), and the links ML(8,3) and ML(8,4) are connected from middle
- switch MS(7,1) to middle switch OS2).

Each of the $\frac{N}{d}$ middle switches OS1 – OS16 in the middle stage 120 are connected from exactly d input switches through two links each for a total of $2 \times d$ links

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(for example the links ML(8,1) and ML(8,2) are connected to the output switch OS1 from input switch MS(7,1), and the links ML(8,7) and ML(7,8) are connected to the output switch OS1 from input switch MS(7,2)).

Finally the connection topology of the network 100A shown in FIG. 1A is knownto be back to back inverse Benes connection topology.

Referring to diagram 100B in FIG. 1B, is a folded version of the multi-link multistage network 100A shown in FIG. 1A. The network 100B in FIG. 1B shows input stage 110 and output stage 120 are placed together. That is input switch IS1 and output switch OS1 are placed together, input switch IS2 and output switch OS2 are placed together, and similarly input switch IS16 and output switch OS16 are placed together. All the right going middle links (hereinafter "forward connecting links") {i.e., inlet links IL1 – IL32 and middle links ML(1,1) - ML(1,64)} correspond to input switches IS1 - IS16, and all the left going middle links (hereinafter "backward connecting links") {i.e., middle links ML(8,1) - ML(8,64) and outlet links OL1-OL32} correspond to output switches OS1 -

15 OS16.

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Middle stage 130 and middle stage 190 are placed together. That is middle switches MS(1,1) and MS(7,1) are placed together, middle switches MS(1,2) and MS(7,2) are placed together, and similarly middle switches MS(1,16) and MS(7,16) are placed together. All the right going middle links {i.e., middle links ML(1,1) - ML(1,64)

and middle links ML(2,1) – ML(2,64)} correspond to middle switches MS(1,1) –
 MS(1,16), and all the left going middle links {i.e., middle links ML(7,1) - ML(7,64) and middle links ML(8,1) and ML(8,64)} correspond to middle switches MS(7,1) –
 MS(7,16).

Middle stage 140 and middle stage 180 are placed together. That is middle switches MS(2,1) and MS(6,1) are placed together, middle switches MS(2,2) and MS(6,2) are placed together, and similarly middle switches MS(2,16) and MS(6,16) are placed together. All the right going middle links {i.e., middle links ML(2,1) - ML(2,64) and middle links ML(3,1) - ML(3,64)} correspond to middle switches MS(2,1) -MS(2,16), and all the left going middle links {i.e., middle links ML(6,1) - ML(6,64) and middle links ML(7,1) and ML(7,64)} correspond to middle switches MS(6,1) - MS(6,16).

Middle stage 150 and middle stage 170 are placed together. That is middle switches MS(3,1) and MS(5,1) are placed together, middle switches MS(3,2) and
5 MS(5,2) are placed together, and similarly middle switches MS(3,16) and MS(5,16) are placed together. All the right going middle links {i.e., middle links ML(3,1) - ML(3,64) and middle links ML(4,1) - ML(4,64)} correspond to middle switches MS(3,1) - MS(3,16), and all the left going middle links {i.e., middle links ML(5,1) - ML(5,64) and middle links ML(6,1) and ML(6,64)} correspond to middle switches MS(5,1) -

10 MS(5,16).

Middle stage 160 is placed alone. All the right going middle links are the middle links ML(4,1) - ML(4,64) and all the left going middle links are middle links ML(5,1) - ML(5,64).

- In one embodiment, in the network 100B of FIG. 1B, the switches that are placed 15 together are implemented as separate switches then the network 100B is the generalized folded multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with nine stages as disclosed in PCT Application Serial No. PCT/US08/64604 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a two by four switch and a four
- 20 by two switch. For example the switch input switch IS1 and output switch OS1 are placed together; so input switch IS1 is implemented as two by four switch with the inlet links IL1 and IL2 being the inputs of the input switch IS1 and middle links ML(1,1) ML(1,4) being the outputs of the input switch IS1; and output switch OS1 is implemented as four by two switch with the middle links ML(8,1), ML(8,2), ML(8,7) and ML(8,8) being the
- 25 inputs of the output switch OS1 and outlet links OL1 OL2 being the outputs of the output switch OS1. Similarly in this embodiment of network 100B all the switches that are placed together in each middle stage are implemented as separate switches.

Hypercube Topology layout schemes:

Referring to layout 100C of FIG. 1C, in one embodiment, there are sixteen blocks namely Block 1_2, Block 3_4, Block 5_6, Block 7_8, Block 9_10, Block 11_12, Block 13_14, Block 15_16, Block 17_18, Block 19_20, Block 21_22, Block 23_24, Block

- 5 25_26, Block 27_28, Block 29_30, and Block 31_32. Each block implements all the switches in one row of the network 100B of FIG. 1B, one of the key aspects of the current invention. For example Block 1_2 implements the input switch IS1, output Switch OS1, middle switch MS(1,1), middle switch MS(7,1), middle switch MS(2,1), middle switch MS(4,1).
- 10 For the simplification of illustration, Input switch IS1 and output switch OS1 together are denoted as switch 1; Middle switch MS(1,1) and middle switch MS(7,1) together are denoted by switch 2; Middle switch MS(2,1) and middle switch MS(6,1) together are denoted by switch 3; Middle switch MS(3,1) and middle switch MS(5,1) together are denoted by switch 4; Middle switch MS(4,1) is denoted by switch 5.
- All the straight middle links are illustrated in layout 100C of FIG. 1C. For example in Block 1_2, inlet links IL1 IL2, outlet links OL1 OL2, middle link ML(1,1), middle link ML(1,2), middle link ML(8,1), middle link ML(8,2), middle link ML(2,1), middle link ML(2,2), middle link ML(7,1), middle link ML(7,2), middle link ML(3,1), middle link ML(3,2), middle link ML(6,1), middle link ML(6,2), middle link
- 20 ML(4,1), middle link ML(4,2), middle link ML(5,1) and middle link ML(5,2) are illustrated in layout 100C of FIG. 1C.

Even though it is not illustrated in layout 100C of FIG. 1C, in each block, in addition to the switches there may be Configurable Logic Blocks (CLB) or any arbitrary digital circuit (hereinafter "sub-integrated circuit block") depending on the applications in different embodiments. There are four quadrants in the layout 100C of FIG. 1C namely top-left, bottom-left, top-right and bottom-right quadrants. Top-left quadrant implements Block 1_2, Block 3_4, Block 5_6, and Block 7_8. Bottom-left quadrant implements Block 9_10, Block 11_12, Block 13_14, and Block 15_16. Top-right quadrant implements Block 17_18, Block 19_20, Block 21_22, and Block 23_24. Bottom-right

30 quadrant implements Block 25_26, Block 27_28, Block 29_30, and Block 31_32. There

are two halves in layout 100C of FIG. 1C namely left-half and right-half. Left-half consists of top-left and bottom-left quadrants. Right-half consists of top-right and bottom-right quadrants.

Recursively in each quadrant there are four sub-quadrants. For example in top-left quadrant there are four sub-quadrants namely top-left sub-quadrant, bottom-left subquadrant, top-right sub-quadrant and bottom-right sub-quadrant. Top-left sub-quadrant of top-left quadrant implements Block 1_2. Bottom-left sub-quadrant of top-left quadrant implements Block 3_4. Top-right sub-quadrant of top-left quadrant implements Block 5_6. Finally bottom-right sub-quadrant of top-left quadrant implements Block 7_8.

- 10 Similarly there are two sub-halves in each quadrant. For example in top-left quadrant there are two sub-halves namely left-sub-half and right-sub-half. Left-sub-half of top-left quadrant implements Block 1_2 and Block 3_4. Right-sub-half of top-left quadrant implements Block 5_6 and Block 7_8. Finally applicant notes that in each quadrant or half the blocks are arranged as a general binary hypercube. Recursively in larger multi-
- 15 stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 > 32$, the layout in this embodiment in accordance with the current invention, will be such that the super-quadrants will also be arranged in d-ary hypercube manner. (In the embodiment of the layout 100C of FIG. 1C, it is binary hypercube manner since d = 2, in the network $V_{fold-mlink}(N_1, N_2, d, s)$ 100B of FIG. 1B).
- Layout 100D of FIG. 1D illustrates the inter-block links between switches 1 and 2 of each block. For example middle links ML(1,3), ML(1,4), ML(8,7), and ML(8,8) are connected between switch 1 of Block 1_2 and switch 2 of Block 3_4. Similarly middle links ML(1,7), ML(1,8), ML(8,3), and ML(8,4) are connected between switch 2 of Block 1_2 and switch 1 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 100D of FIG. 1D can be implemented as vertical tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(1,4) and ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(8,8) are implemented as two different tracks (for example middle links ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(8,8) are implemented as a time
- 30 division multiplexed single track).

Layout 100E of FIG. 1E illustrates the inter-block links between switches 2 and 3 of each block. For example middle links ML(2,3), ML(2,4), ML(7,11), and ML(7,12) are connected between switch 2 of Block 1_2 and switch 3 of Block 3_4. Similarly middle links ML(2,11), ML(2,12), ML(7,3), and ML(7,4) are connected between switch 3 of

- 5 Block 1_2 and switch 2 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 100E of FIG. 1E can be implemented as horizontal tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(2,12) and ML(7,4) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division
- 10 multiplexed single track (for example middle links ML(2,12) and ML(7,4) are implemented as a time division multiplexed single track).

Layout 100F of FIG. 1F illustrates the inter-block links between switches 3 and 4 of each block. For example middle links ML(3,3), ML(3,4), ML(6,19), and ML(6,20) are connected between switch 3 of Block 1_2 and switch 4 of Block 3_4. Similarly middle

- 15 links ML(3,19), ML(3,20), ML(6,3), and ML(6,4) are connected between switch 4 of Block 1_2 and switch 3 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 100F of FIG. 1F can be implemented as vertical tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(3,4) and ML(6,20) are implemented as two different tracks); or in an
- 20 alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track).

Layout 100G of FIG. 1G illustrates the inter-block links between switches 4 and 5 of each block. For example middle links ML(4,3), ML(4,4), ML(5,35), and ML(5,36) are
connected between switch 4 of Block 1_2 and switch 5 of Block 3_4. Similarly middle links ML(4,35), ML(4,36), ML(5,3), and ML(5,4) are connected between switch 5 of Block 1_2 and switch 4 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 100G of FIG. 1G can be implemented as horizontal tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for

30 example middle links ML(4,4) and ML(5,36) are implemented as two different tracks); or

in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(4,4) and ML(5,36) are implemented as a time division multiplexed single track).

The complete layout for the network 100B of FIG. 1B is given by combining the links in layout diagrams of 100C, 100D, 100E, 100F, and 100G. Applicant notes that in the layout 100C of FIG. 1C, the inter-block links between switch 1 and switch 2 of corresponding blocks are vertical tracks as shown in layout 100D of FIG. 1D; the interblock links between switch 2 and switch 3 of corresponding blocks are horizontal tracks as shown in layout 100E of FIG. 1E; the inter-block links between switch 3 and switch 4

- 10 of corresponding blocks are vertical tracks as shown in layout 100F of FIG. 1F; and finally the inter-block links between switch 4 and switch 5 of corresponding blocks are horizontal tracks as shown in layout 100G of FIG. 1G. The pattern is alternate vertical tracks and horizontal tracks. It continues recursively for larger networks of N > 32 as will be illustrated later.
- 15 Some of the key aspects of the current invention are discussed. 1) All the switches in one row of the multi-stage network 100B are implemented in a single block. 2) The blocks are placed in such a way that all the inter-block links are either horizontal tracks or vertical tracks; 3) Since all the inter-block links are either horizontal or vertical tracks, all the inter-block links can be mapped on to island-style architectures in current commercial
- 20 FPGA's; 4) The length of the longest wire is about half of the width (or length) of the complete layout (For example middle link ML(4,4) is about half the width of the complete layout).

In accordance with the current invention, the layout 100C in FIG. 1C can be recursively extended for any arbitrarily large generalized folded multi-link multi-stage 25 network $V_{fold-mlink}(N_1, N_2, d, s)$ the sub-quadrants, quadrants, and super-quadrants are arranged in d-ary hypercube manner and also the inter-blocks are accordingly connected in d-ary hypercube topology. Even though all the embodiments in the current invention are illustrated for $N_1 = N_2$, the embodiments can be extended for $N_1 \neq N_2$. Referring to layout 100H of FIG. 1H, illustrates the extension of layout 100C for the network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 128$; d = 2; and s = 2. There are four super-quadrants in layout 100H namely top-left super-quadrant, bottom-left super-quadrant, top-right super-quadrant, bottom-right super-quadrant. Total number of blocks

- 5 in the layout 100H is sixty four. Top-left super-quadrant implements the blocks from block 1_2 to block 31_32. Each block in all the super-quadrants has two more switches namely switch 6 and switch 7 in addition to the switches [1-5] illustrated in layout 100C of FIG. 1C. The inter-block link connection topology is the exactly the same between the switches 1 and 2; switches 2 and 3; switches 3 and 4; switches 4 and 5 as it is shown in
- 10 the layouts of FIG. 1D, FIG. 1E, FIG. 1F, and FIG. 1G respectively.

Bottom-left super-quadrant implements the blocks from block 33_34 to block 63_64. Top-right super-quadrant implements the blocks from block 65_66 to block 95_96. And bottom-right super-quadrant implements the blocks from block 97_98 to block 127_128. In all these three super-quadrants also, the inter-block link connection

topology is the exactly the same between the switches 1 and 2; switches 2 and 3; switches3 and 4; switches 4 and 5 as that of the top-left super-quadrant.

Recursively in accordance with the current invention, the inter-block links connecting the switch 5 and switch 6 will be vertical tracks between the corresponding switches of top-left super-quadrant and bottom-left super-quadrant. And similarly the

- 20 inter-block links connecting the switch 5 and switch 6 will be vertical tracks between the corresponding switches of top-right super-quadrant and bottom-right super-quadrant. The inter-block links connecting the switch 6 and switch 7 will be horizontal tracks between the corresponding switches of top-left super-quadrant and top-right super-quadrant. And similarly the inter-block links connecting the switch 6 and switch 7 will be horizontal
- 25 tracks between the corresponding switches of bottom-left super-quadrant and bottomright super-quadrant.

Referring to diagram 100I of FIG. 1I illustrates a high-level implementation of Block 1_2 (Each of the other blocks have similar implementation) of the layout 100C of FIG. 1C which represents a generalized folded multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2. Block 1_2 in 100I illustrates both the intra-block and inter-block links connected to Block 1_2. The layout diagram 100I corresponds to the embodiment where the switches that are placed together are implemented as separate switches in the network 100B of FIG. 1B. As noted before then the network 100B is the generalized folded multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with nine stages as disclosed in PCT Application Serial No. PCT/US08/64604 that is incorporated by reference

above.

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That is the switches that are placed together in Block 1_2 as shown in FIG. 1I are
namely input switch IS1 and output switch OS1 belonging to switch 1, illustrated by
dotted lines, (as noted before switch 1 is for illustration purposes only, in practice the
switches implemented are input switch IS1 and output switch OS1); middle switch
MS(1,1) and middle switch MS(7,1) belonging to switch 2; middle switch MS(2,1) and
middle switch MS(6,1) belonging to switch 3; middle switch MS(3,1) and middle switch
MS(5,1) belonging to switch 4; And middle switch MS(4,1) belonging to switch 5.

Input switch IS1 is implemented as two by four switch with the inlet links IL1 and IL2 being the inputs of the input switch IS1 and middle links ML(1,1) - ML(1,4) being the outputs of the input switch IS1; and output switch OS1 is implemented as four by two switch with the middle links ML(8,1) - ML(8,4) being the inputs of the output switch OS1 and outlet links OL1 - OL2 being the outputs of the output switch OS1.

Middle switch MS(1,1) is implemented as four by four switch with the middle links ML(1,1), ML(1,2), ML(1,7) and ML(1,8) being the inputs and middle links ML(2,1)– ML(2,4) being the outputs; and middle switch MS(7,1) is implemented as four by four switch with the middle links ML(7,1), ML(7,2), ML(7,11) and ML(7,12) being the inputs and middle links ML(8,1) - ML(8,4) being the outputs. Similarly all the other middle switches are also implemented as four by four switches as illustrated in 100I of FIG. 1I.

Now the VLSI layouts of generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$ where N₁ = N₂ < 32; d = 2; s = 2 and its corresponding version of folded generalized multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 < 32$; d = 2; s = 2 are discussed. Referring to diagram 200A1 of FIG. 2A1 is generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 2$; d = 2. Diagram 200A2 of FIG. 2A2 illustrates the corresponding folded generalized multi-link multi-

- 5 stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 2$; d = 2, version of the diagram 200A1 of FIG. 2A1. Layout 200A3 of FIG. 2A3 illustrates the VLSI layout of the network 200A2 of FIG. 2A2. There is only one block i.e., Block 1_2 comprising switch 1. Just like in the layout 100C of FIG. 1C, switch 1 consists of input switch IS1 and output switch OS1.
- 10 Referring to diagram 200B1 of FIG. 2B1 is generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 4$; d = 2; s = 2. Diagram 200B2 of FIG. 2B2 illustrates the corresponding folded generalized multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 4$; d = 2; s = 2, version of the diagram 200B1 of FIG. 2B1. Layout 200B3 of FIG. 2B3 illustrates the VLSI layout of the network 200B2
- 15 of FIG. 2B2. There are two blocks i.e., Block 1_2 and Block 3_4 each comprising switch 1 and switch 2. Switch 1 in each block consists of the corresponding input switch and output switch. For example switch 1 in Block 1_2 consists of input switch IS1 and output switch OS1. Similarly switch 2 in Block 1_2 consists of middle switch (1,1). Layout 200B4 of FIG. 2B4 illustrates the inter-block links of the VLSI layout diagram 200B3 of
- 20 FIG. 2B3. For example middle links ML(1,4) and ML(2,8). It must be noted that all the inter-block links are vertical tracks in this layout. (Alternatively all the inter-blocks can also be implemented as horizontal tracks).

Referring to diagram 200C11 of FIG. 2C11 is generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 8$; d = 2; s = 2. Diagram 200C12 of FIG.

25 2C12 illustrates the corresponding folded generalized multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 8$; d = 2; s = 2, version of the diagram 200C11 of FIG. 2C11. Layout 200C21 of FIG. 2C21 illustrates the VLSI layout of the network 200C12 of FIG. 2C12. There are four blocks i.e., Block 1_2, Block 3_4, Block 5_6, and Block 7_8 each comprising switch 1, switch 2 and switch 3. For example switch 1 in

Block 1_2 consists of input switch IS1 and output switch OS1; Switch 2 in Block 1_2 consists of MS(1,1) and MS(3,1). Switch 3 in Block 1_2 consists of MS(2,1).

Layout 200C22 of FIG. 2C22 illustrates the inter-block links between the switch 1 and switch 2 of the VLSI layout diagram 200C21 of FIG. 2C21. For example middle
links ML(1,4) and ML(4,8) are connected between Block 1_2 and Block 3_4. It must be noted that all the inter-block links between switch 1 and switch 2 of all blocks are vertical tracks in this layout. Layout 200C23 of FIG. 2C23 illustrates the inter-block links between the switch 2 and switch 3 of the VLSI layout diagram 200C21 of FIG. 2C21. For example middle links ML(2,12) and ML(3,4) are connected between Block 1_2 and

10 Block 5_6. It must be noted that all the inter-block links between switch 2 and switch 3 of all blocks are horizontal tracks in this layout

Referring to diagram 200D1 of FIG. 2D1 is generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$ where N₁ = N₂ = 16; d = 2; s = 2. Diagram 200D2 of FIG. 2D2 illustrates the corresponding folded generalized multi-link multi-stage network

- 15 $V_{fold-mlink}(N_1, N_2, d, s)$ where N₁ = N₂ = 16; d = 2; s = 2, version of the diagram 200D1 of FIG. 2D1. Layout 200D3 of FIG. 2D3 illustrates the VLSI layout of the network 200D2 of FIG. 2D2. There are eight blocks i.e., Block 1_2, Block 3_4, Block 5_6, Block 7_8, Block 9_10, Block 11_12, Block 13_14 and Block 15_16 each comprising switch 1, switch 2, switch 3 and switch 4. For example switch 1 in Block 1_2 consists of input
- 20 switch IS1 and output switch OS1; Switch 2 in Block 1_2 consists of MS(1,1) and MS(5,1). Switch 3 in Block 1_2 consists of MS(2,1) and MS(4,1), and switch 4 in Block 1_2 consists of MS(3,1).

Layout 200D4 of FIG. 2D4 illustrates the inter-block links between the switch 1 and switch 2 of the VLSI layout diagram 200D3 of FIG. 2D3. For example middle links
25 ML(1,4) and ML(6,8) are connected between Block 1_2 and Block 3_4. It must be noted that all the inter-block links between switch 1 and switch 2 of all blocks are vertical tracks in this layout. Layout 200D5 of FIG. 2D5 illustrates the inter-block links between the switch 2 and switch 3 of the VLSI layout diagram 200D3 of FIG. 2D3. For example middle links ML(2,12) and ML(5,4) are connected between Block 1_2 and Block 5_6. It

must be noted that all the inter-block links between switch 2 and switch 3 of all blocks are horizontal tracks in this layout. Layout 200D6 of FIG. 2D6 illustrates the inter-block links between the switch 3 and switch 4 of the VLSI layout diagram 200D3 of FIG. 2D3. For example middle links ML(3,4) and ML(4,20) are connected between Block 1_2 and Block 9_10. It must be noted that all the inter-block links between switch 3 and switch 4

5 Block 9_10. It must be noted that all the inter-block links between switch 3 and switch 4 of all blocks are vertical tracks in this layout.

Generalized Multi-link Butterfly Fat Tree Network Embodiment:

In another embodiment in the network 100B of FIG. 1B, the switches that are placed together are implemented as combined switch then the network 100B is the 10 generalized multi-link butterfly fat tree network $V_{mlink-bft}(N_1, N_2, d, s)$ where $N_1 = N_2 =$ 32; d = 2; and s = 2 with five stages as disclosed in PCT Application Serial No. PCT/US08/64603 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a six by six switch. For example the input switch IS1 and output switch OS1 are placed together; so

- 15 input switch IS1 and output OS1 are implemented as a six by six switch with the inlet links IL1, IL2, ML(8,1), ML(8,2), ML(8,7) and ML(8,8) being the inputs of the combined switch (denoted as IS1&OS1) and middle links ML(1,1), ML(1,2), ML(1,3), ML(1,4), OL1 and OL2 being the outputs of the combined switch IS1&OS1. Similarly in this embodiment of network 100B all the switches that are placed together are
- 20 implemented as a combined switch.

Layout diagrams 100C in FIG. 1C, 100D in FIG. 1D, 100E in FIG. 1E, 100F in FIG. 1G are also applicable to generalized multi-link butterfly fat tree network $V_{mlink-bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with five stages. The layout 100C in FIG. 1C can be recursively extended for any arbitrarily large generalized multi-link butterfly fat tree network $V_{mlink-bft}(N_1, N_2, d, s)$. Accordingly layout 100H of FIG. 1H is also applicable to generalized multi-link butterfly fat tree network $V_{mlink-bft}(N_1, N_2, d, s)$.

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Referring to diagram 100J of FIG. 1J illustrates a high-level implementation of Block 1_2 (Each of the other blocks have similar implementation) of the layout 100C of FIG. 1C which represents a generalized multi-link butterfly fat tree network $V_{mlink-bft}(N_1, N_2, d, s)$ where N₁ = N₂ = 32; d = 2; and s = 2. Block 1_2 in 100J illustrates

- 5 both the intra-block and inter-block links. The layout diagram 100J corresponds to the embodiment where the switches that are placed together are implemented as combined switch in the network 100B of FIG. 1B. As noted before then the network 100B is the generalized multi-link butterfly fat tree network $V_{mlink-bft}(N_1, N_2, d, s)$ where $N_1 = N_2 =$ 32; d = 2; and s = 2 with five stages as disclosed in PCT Application Serial No.
- 10 PCT/US08/64603 that is incorporated by reference above.

switch and each middle switch provides U-turn links.

That is the switches that are placed together in Block 1_2 as shown in FIG. 1J are namely the combined input and output switch IS1&OS1 belonging to switch 1, illustrated by dotted lines, (as noted before switch 1 is for illustration purposes only, in practice the switch implemented is combined input and output switch IS1&OS1); middle switch

MS(1,1) belonging to switch 2; middle switch MS(2,1) belonging to switch 3; middle switch MS(3,1) belonging to switch 4; And middle switch MS(4,1) belonging to switch 5.

Combined input and output switch IS1&OS1 is implemented as six by six switch with the inlet links IL1, IL2 and ML(8,1) - ML(8,4) being the inputs and middle links ML(1,1) - ML(1,4), and outlet links OL1 - OL2 being the outputs.

Middle switch MS(1,1) is implemented as eight by eight switch with the middle links ML(1,1), ML(1,2), ML(1,7), ML(1,8), ML(7,1), ML(7,2), ML(7,11) and ML(7,12) being the inputs and middle links ML(2,1) – ML(2,4) and middle links ML(8,1) – ML(8,4) being the outputs. Similarly all the other middle switches are also implemented as eight by eight switches as illustrated in 100J of FIG. 1J. Applicant observes that in middle switch MS(1,1) any one of the right going middle links can be switched to any one of the left going middle links and hereinafter middle switch MS(1,1) provides U-turn links. In general, in the network V_{mlink-bft} (N₁, N₂, d, s) each input switch, each output

In another embodiment, middle switch MS(1,1) (or the middle switches in any of the middle stage excepting the root middle stage) of Block 1_2 of

 $V_{mlink-bft}(N_1, N_2, d, s)$ can be implemented as a four by eight switch and a four by four switch to save cross points. This is because the left going middle links of these middle

- 5 switches are never setup to the right going middle links. For example, in middle switch MS(1,1) of Block 1_2 as shown FIG. 1J, the left going middle links namely ML(7,1), ML(7,2), ML(7,11), and ML(7,12) are never switched to the right going middle links ML(2,1), ML(2,2), ML(2,3), and ML(2,4). And hence to implement MS(1,1) two switches namely: 1) a four by eight switch with the middle links ML(1,1), ML(1,2),
- ML(1,7), and ML(1,8) as inputs and the middle links ML(2,1), ML(2,2), ML(2,3),
 ML(2,4), ML(8,1), ML(8,2), ML(8,3), and ML(8,4) as outputs and 2) a four by four switch with the middle links ML(7,1), ML(7,2), ML(7,11), and ML(7,12) as inputs and the middle links ML(8,1), ML(8,2), ML(8,3), and ML(8,4) as outputs are sufficient without loosing any connectivity of the embodiment of MS(1,1) being implemented as an
- 15 eight by eight switch as described before.)

Generalized multi-stage network Embodiment:

In one embodiment, in the network 100B of FIG. 1B, the switches that are placed together are implemented as two separate switches in input stage 110 and output stage 120; and as four separate switches in all the middle stages, then the network 100B is the generalized folded multi-stage network $V_{fold}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with nine stages as disclosed in PCT Application Serial No. PCT/US08/64604 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a two by four switch and a four by two switch respectively. For example the switch input switch IS1 and output switch 25 OS1 are placed together; so input switch IS1 is implemented as two by four switch with the inlet links IL1 and IL2 being the inputs and middle links ML(11) – ML(14) being

the inlet links IL1 and IL2 being the inputs and middle links ML(1,1) - ML(1,4) being the outputs; and output switch OS1 is implemented as four by two switch with the middle links ML(8,1), ML(8,4), ML(8,7) and ML(8,8) being the inputs and outlet links OL1 – OL2 being the outputs. The switches, corresponding to the middle stages that are placed together are implemented as four two by two switches. For example middle switches MS(1,1), MS(1,17), MS(7,1), and MS(7,17) are placed together; so middle switch MS(1,1) is implemented as two by two switch with middle links ML(1,1) and ML(1,7) being the

- 5 inputs and middle links ML(2,1) and ML(2,3) being the outputs; middle switch MS(1,17) is implemented as two by two switch with the middle links ML(1,2) and ML(1,8) being the inputs and middle links ML(2,2) and ML(2,4) being the outputs; middle switch MS(7,1) is implemented as two by two switch with middle links ML(7,1) and ML(7,11) being the inputs and middle links ML(8,1) and ML(8,3) being the outputs; And middle
- 10 switch MS(7,17) is implemented as two by two switch with the middle links ML(7,2) and ML(7,12) being the inputs and middle links ML(8,2) and ML(8,4) being the outputs; Similarly in this embodiment of network 100B all the switches that are placed together are implemented as separate switches.
- Layout diagrams 100C in FIG. 1C, 100D in FIG. 1D, 100E in FIG. 1E, 100F in 15 FIG. 1G are also applicable to generalized folded multi-stage network $V_{fold}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with nine stages. The layout 100C in FIG. 1C can be recursively extended for any arbitrarily large generalized folded multistage network $V_{fold}(N_1, N_2, d, s)$. Accordingly layout 100H of FIG. 1H is also applicable to generalized folded multi-stage network $V_{fold}(N_1, N_2, d, s)$.
- 20 Referring to diagram 100K of FIG. 1K illustrates a high-level implementation of Block 1_2 (Each of the other blocks have similar implementation) of the layout 100C of FIG. 1C which represents a generalized folded multi-stage network $V_{fold}(N_1, N_2, d, s)$ where N₁ = N₂ = 32; d = 2; and s = 2. Block 1_2 in 100K illustrates both the intra-block and inter-block links. The layout diagram 100K corresponds to the embodiment where the switches that are placed together are implemented as separate switches in the network 100B of FIG. 1B. As noted before then the network 100B is the generalized folded multi-stage network $V_{fold}(N_1, N_2, d, s)$ where N₁ = N₂ = 32; d = 2; and s = 2 with nine stages as disclosed in PCT Application Serial No. PCT/US08/64604 that is incorporated by reference above.

That is the switches that are placed together in Block 1_2 as shown in FIG. 1K are namely the input switch IS1 and output switch OS1 belonging to switch 1, illustrated by dotted lines, (as noted before switch 1 is for illustration purposes only, in practice the switches implemented are input switch IS1 and output switch OS1); middle switches

- MS(1,1), MS(1,17), MS(7,1) and MS(7,17) belonging to switch 2; middle switches
 MS(2,1), MS(2,17), MS(6,1) and MS(6,17) belonging to switch 3; middle switches
 MS(3,1), MS(3,17), MS(5,1) and MS(5,17) belonging to switch 4; And middle switches
 MS(4,1), and MS(4,17) belonging to switch 5.
- Input switch IS1 and output switch OS1 are placed together; so input switch IS1 is
 implemented as two by four switch with the inlet links IL1 and IL2 being the inputs and middle links ML(1,1) ML(1,4) being the outputs; and output switch OS1 is implemented as four by two switch with the middle links ML(8,1), ML(8,4), ML(8,7) and ML(8,8) being the inputs and outlet links OL1 OL2 being the outputs.

Middle switches MS(1,1), MS(1,17), MS(7,1), and MS(7,17) are placed together; so middle switch MS(1,1) is implemented as two by two switch with middle links

- 15 so middle switch MS(1,1) is implemented as two by two switch with middle links ML(1,1) and ML(1,7) being the inputs and middle links ML(2,1) and ML(2,3) being the outputs; middle switch MS(1,17) is implemented as two by two switch with the middle links ML(1,2) and ML(1,8) being the inputs and middle links ML(2,2) and ML(2,4) being the outputs; middle switch MS(7,1) is implemented as two by two switch with middle
- 20 links ML(7,1) and ML(7,11) being the inputs and middle links ML(8,1) and ML(8,3) being the outputs; And middle switch MS(7,17) is implemented as two by two switch with the middle links ML(7,2) and ML(7,12) being the inputs and middle links ML(8,2) and ML(8,4) being the outputs. Similarly all the other middle switches are also implemented as two by two switches as illustrated in 100K of FIG. 1K.

25 Generalized multi-stage network Embodiment with S = 1:

In one embodiment, in the network 100B of FIG. 1B (where it is implemented with s = 1), the switches that are placed together are implemented as two separate switches in input stage 110 and output stage 120; and as two separate switches in all the middle stages, then the network 100B is the generalized folded multi-stage network

 $V_{fold}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 1 with nine stages as disclosed in PCT Application Serial No. PCT/US08/64604 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a two by two switch and a two by two switch. For example the switch

- 5 input switch IS1 and output switch OS1 are placed together; so input switch IS1 is implemented as two by two switch with the inlet links IL1 and IL2 being the inputs and middle links ML(1,1) ML(1,2) being the outputs; and output switch OS1 is implemented as two by two switch with the middle links ML(8,1) and ML(8,3) being the inputs and outlet links OL1 OL2 being the outputs.
- 10 The switches, corresponding to the middle stages that are placed together are implemented as two, two by two switches. For example middle switches MS(1,1) and MS(7,1) are placed together; so middle switch MS(1,1) is implemented as two by two switch with middle links ML(1,1) and ML(1,3) being the inputs and middle links ML(2,1) and ML(2,2) being the outputs; middle switch MS(7,1) is implemented as two
- 15 by two switch with middle links ML(7,1) and ML(7,5) being the inputs and middle links ML(8,1) and ML(8,2) being the outputs; Similarly in this embodiment of network 100B all the switches that are placed together are implemented as two separate switches.

Layout diagrams 100C in FIG. 1C, 100D in FIG. 1D, 100E in FIG. 1E, 100F in FIG. 1G are also applicable to generalized folded multi-stage network

20 $V_{fold}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 1 with nine stages. The layout 100C in FIG. 1C can be recursively extended for any arbitrarily large generalized folded multi-stage network $V_{fold}(N_1, N_2, d, s)$. Accordingly layout 100H of FIG. 1H is also applicable to generalized folded multi-stage network $V_{fold}(N_1, N_2, d, s)$.

Referring to diagram 100K1 of FIG. 1K1 illustrates a high-level implementation 25 of Block 1_2 (Each of the other blocks have similar implementation) for the layout 100C of FIG. 1C when s = 1 which represents a generalized folded multi-stage network $V_{fold} (N_1, N_2, d, s)$ where N₁ = N₂ = 32; d = 2; and s = 1 (All the double links are replaced by single links when s = 1). Block 1_2 in 100K1 illustrates both the intra-block and inter-

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block links. The layout diagram 100K1 corresponds to the embodiment where the switches that are placed together are implemented as separate switches in the network 100B of FIG. 1B when s = 1. As noted before then the network 100B is the generalized folded multi-stage network V_{fold} (N_1, N_2, d, s) where $N_1 = N_2 = 32$; d = 2; and s = 1 with nine stages as disclosed in PCT Application Serial No. PCT/US08/64604 that is

incorporated by reference above.

That is the switches that are placed together in Block 1_2 as shown in FIG. 1K1 are namely the input switch IS1 and output switch OS1 belonging to switch 1, illustrated by dotted lines, (as noted before switch 1 is for illustration purposes only, in practice the switches implemented are input switch IS1 and output switch OS1); middle switches MS(1,1) and MS(7,1) belonging to switch 2; middle switches MS(2,1) and MS(6,1) belonging to switch 3; middle switches MS(3,1) and MS(5,1) belonging to switch 4; And middle switch MS(4,1) belonging to switch 5.

Input switch IS1 and output switch OS1 are placed together; so input switch IS1 is implemented as two by two switch with the inlet links IL1 and IL2 being the inputs and middle links ML(1,1) – ML(1,2) being the outputs; and output switch OS1 is implemented as two by two switch with the middle links ML(8,1) and ML(8,3) being the inputs and outlet links OL1 – OL2 being the outputs.

Middle switches MS(1,1) and MS(7,1) are placed together; so middle switch
20 MS(1,1) is implemented as two by two switch with middle links ML(1,1) and ML(1,3) being the inputs and middle links ML(2,1) and ML(2,2) being the outputs; And middle switch MS(7,1) is implemented as two by two switch with middle links ML(7,1) and ML(7,5) being the inputs and middle links ML(8,1) and ML(8,2) being the outputs. Similarly all the other middle switches are also implemented as two by two switches as
25 illustrated in 100K1 of FIG. 1K1.

Generalized Butterfly Fat Tree Network Embodiment:

In another embodiment in the network 100B of FIG. 1B, the switches that are placed together are implemented as two combined switches then the network 100B is the

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generalized butterfly fat tree network V_{bf} (N_1 , N_2 , d, s) where $N_1 = N_2 = 32$; d = 2; and s = 2 with five stages as disclosed in PCT Application Serial No. PCT/US08/64603 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a six by six switch. For example the input switch IS1 and output switch OS1 are placed together; so input output switch IS1&OS1 are implemented as a six by six switch with the inlet links IL1, IL2, ML(8,1), ML(8,2), ML(8,7) and ML(8,8) being the inputs of the combined switch (denoted as IS1&OS1) and middle links ML(1,1), ML(1,2), ML(1,3), ML(1,4), OL1 and OL2 being the outputs of the combined switch IS1&OS1.

10 The switches, corresponding to the middle stages that are placed together are implemented as two four by four switches. For example middle switches MS(1,1) and MS(1,17) are placed together; so middle switch MS(1,1) is implemented as four by four switch with middle links ML(1,1), ML(1,7), ML(7,1) and ML(7,11) being the inputs and middle links ML(2,1), ML(2,3), ML(8,1) and ML(8,3) being the outputs; middle switch
15 MS(1,17) is implemented as four by four switch with the middle links ML(1,2), ML(1,8), ML(7,2) and ML(7,12) being the inputs and middle links ML(2,2), ML(2,4), ML(8,2) and ML(8,4) being the outputs. Similarly in this embodiment of network 100B all the switches that are placed together are implemented as a two combined switches.

Layout diagrams 100C in FIG. 1C, 100D in FIG. 1D, 100E in FIG. 1E, 100F in 20 FIG. 1G are also applicable to generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with five stages. The layout 100C in FIG. 1C can be recursively extended for any arbitrarily large generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$. Accordingly layout 100H of FIG. 1H is also applicable to generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$.

25 Referring to diagram 100L of FIG. 1L illustrates a high-level implementation of Block 1_2 (Each of the other blocks have similar implementation) of the layout 100C of FIG. 1C which represents a generalized butterfly fat tree network $V_{bfi}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2. Block 1_2 in 100L illustrates both the intra-block and

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inter-block links. The layout diagram 100L corresponds to the embodiment where the switches that are placed together are implemented as two combined switches in the network 100B of FIG. 1B. As noted before then the network 100B is the generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with five stages as disclosed in PCT Application Serial No. PCT/US08/64603 that is incorporated by reference above.

- That is the switches that are placed together in Block 1_2 as shown in FIG. 1L are namely the combined input and output switch IS1&OS1 belonging to switch 1, illustrated by dotted lines, (as noted before switch 1 is for illustration purposes only, in practice the switch implemented is combined input and output switch IS1&OS1); middle switch MS(1,1) and MS(1,17) belonging to switch 2; middle switch MS(2,1) and MS(2,17)
 - belonging to switch 3; middle switch MS(3,1) and MS(3,17) belonging to switch 4; And middle switch MS(4,1) belonging to switch 5.

Combined input and output switch IS1&OS1 is implemented as six by six switch with the inlet links IL1, IL2, ML(8,1), ML(8,2), ML(8,7) and ML(8,8) being the inputs and middle links ML(1,1) – ML(1,4) and outlet links OL1 – OL2 being the outputs.

Middle switch MS(1,1) is implemented as four by four switch with middle links ML(1,1), ML(1,7), ML(7,1) and ML(7,11) being the inputs and middle links ML(2,1), ML(2,3), ML(8,1) and ML(8,3) being the outputs; And middle switch MS(1,17) is

- 20 implemented as four by four switch with the middle links ML(1,2), ML(1,8), ML(7,2) and ML(7,12) being the inputs and middle links ML(2,2), ML(2,4), ML(8,2) and ML(8,4) being the outputs. Similarly all the other middle switches are also implemented as two four by four switches as illustrated in 100L of FIG. 1L. Applicant observes that in middle switch MS(1,1) any one of the right going middle links can be switched to any one of the
- 25 left going middle links and hereinafter middle switch MS(1,1) provides U-turn links. In general, in the network $V_{bft}(N_1, N_2, d, s)$ each input switch, each output switch and each middle switch provides U-turn links.

In another embodiment, middle switch MS(1,1) (or the middle switches in any of the middle stage excepting the root middle stage) of Block 1_2 of $V_{bft}(N_1, N_2, d, s)$ can be implemented as a two by four switch and a two by two switch to save cross points. This is because the left going middle links of these middle switches are never setup to the

- 5 right going middle links. For example, in middle switch MS(1,1) of Block 1_2 as shown FIG. 1L, the left going middle links namely ML(7,1) and ML(7,11) are never switched to the right going middle links ML(2,1) and ML(2,3). And hence to implement MS(1,1) two switches namely: 1) a two by four switch with the middle links ML(1,1) and ML(1,7) as inputs and the middle links ML(2,1), ML(2,3), ML(8,1), and ML(8,3) as outputs and
- 2) a two by two switch with the middle links ML(7,1) and ML(7,11) as inputs and the middle links ML(8,1) and ML(8,3) as outputs are sufficient without loosing any connectivity of the embodiment of MS(1,1) being implemented as an eight by eight switch as described before.)

15 Generalized Butterfly Fat Tree Network Embodiment with S = 1:

In one embodiment, in the network 100B of FIG. 1B (where it is implemented with s = 1), the switches that are placed together are implemented as a combined switch in input stage 110 and output stage 120; and as a combined switch in all the middle stages, then the network 100B is the generalized butterfly fat tree network

- 20 $V_{bft}(N_1, N_2, d, s)$ where N₁ = N₂ = 32; d = 2; and s = 1 with five stages as disclosed in PCT Application Serial No. PCT/US08/64603 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a four by four switch. For example the switch input switch IS1 and output switch OS1 are placed together; so input and output switch IS1&OS1 is
- 25 implemented as four by four switch with the inlet links IL1, IL2, ML(8,1) and ML(8,3) being the inputs and middle links ML(1,1) ML(1,2) and outlet links OL1 OL2 being the outputs

The switches, corresponding to the middle stages that are placed together are implemented as a four by four switch. For example middle switches MS(1,1) is

implemented as four by four switch with middle links ML(1,1), ML(1,3), ML(7,1) and ML(7,5) being the inputs and middle links ML(2,1), ML(2,2), ML(8,1) and ML(8,2) being the outputs..

Layout diagrams 100C in FIG. 1C, 100D in FIG. 1D, 100E in FIG. 1E, 100F in 5 FIG. 1G are also applicable to generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 1 with five stages. The layout 100C in FIG. 1C can be recursively extended for any arbitrarily large generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$. Accordingly layout 100H of FIG. 1H is also applicable to generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$.

10 Referring to diagram 100L1 of FIG. 1L1 illustrates a high-level implementation of Block 1_2 (Each of the other blocks have similar implementation) for the layout 100C of FIG. 1C when s = 1 which represents a generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$ where N₁ = N₂ = 32; d = 2; and s = 1 (All the double links are replaced by single links when s = 1). Block 1_2 in 100K1 illustrates both the intra-block and inter-

- 15 block links. The layout diagram 100L1 corresponds to the embodiment where the switches that are placed together are implemented as a combined switch in the network 100B of FIG. 1B when s = 1. As noted before then the network 100B is the generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 1 with nine stages as disclosed in PCT Application Serial No. PCT/US08/64603 that is
- 20 incorporated by reference above.

That is the switches that are placed together in Block 1_2 as shown in FIG. 1L1 are namely the input and output switch IS1&OS1 belonging to switch 1, illustrated by dotted lines, (as noted before switch 1 is for illustration purposes only, in practice the switches implemented are input switch IS1 and output switch OS1); middle switch

25 MS(1,1) belonging to switch 2; middle switch MS(2,1) belonging to switch 3; middle switch MS(3,1) belonging to switch 4; And middle switch MS(4,1) belonging to switch 5.

Input and output switch IS1&OS1 are placed together; so input and output switch IS1&OS1 is implemented as four by four switch with the inlet links IL1, IL2, ML(8,1)

and ML(8,3) being the inputs and middle links ML(1,1) - ML(1,2) and outlet links OL1 - OL2 being the outputs.

Middle switch MS(1,1) is implemented as four by four switch with middle links ML(1,1), ML(1,3), ML(7,1) and ML(7,5) being the inputs and middle links ML(2,1),
ML(2,2), ML(8,1) and ML(8,2) being the outputs. Similarly all the other middle switches are also implemented as four by four switches as illustrated in 100L1 of FIG. 1L1.

In another embodiment, middle switch MS(1,1) (or the middle switches in any of the middle stage excepting the root middle stage) of Block 1_2 of $V_{mlink-bft}(N_1, N_2, d, s)$ can be implemented as a two by four switch and a two by two

- 10 switch to save cross points. This is because the left going middle links of these middle switches are never setup to the right going middle links. For example, in middle switch MS(1,1) of Block 1_2 as shown FIG. 1L1, the left going middle links namely ML(7,1) and ML(7,5) are never switched to the right going middle links ML(2,1) and ML(2,2). And hence to implement MS(1,1) two switches namely: 1) a two by four switch with the
- 15 middle links ML(1,1) and ML(1,3) as inputs and the middle links ML(2,1), ML(2,2), ML(8,1), and ML(8,2) as outputs and 2) a two by two switch with the middle links ML(7,1) and ML(7,5) as inputs and the middle links ML(8,1) and ML(8,2) as outputs are sufficient without loosing any connectivity of the embodiment of MS(1,1) being implemented as an eight by eight switch as described before.)

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Hypercube-like Topology layout schemes:

Referring to diagram 300A in FIG. 3A, in one embodiment, an exemplary generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with nine stages of one hundred and forty four switches for satisfying

25 communication requests, such as setting up a telephone call or a data call, or a connection between configurable logic blocks, between an input stage 110 and output stage 120 via middle stages 130, 140, 150, 170, 170, 180 and 190 is shown where input stage 110

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consists of sixteen, two by four switches IS1-IS16 and output stage 120 consists of sixteen, four by two switches OS1-OS16.

As disclosed in PCT Application Serial No. PCT/US08/64604 that is incorporated by reference above, such a network can be operated in rearrangeably nonblocking manner for arbitrary fan-out multicast connections and also can be operated in strictly non-blocking manner for unicast connections.

The diagram 300A in FIG. 3A is exactly the same as the diagram 100A in FIG. 1A excepting the connection links between middle stage 150 and middle stage 160 as well as between middle stage 160 and middle stage 170.

10 Each of the $\frac{N}{d}$ middle switches are connected to exactly *d* switches in middle stage 160 through two links each for a total of $2 \times d$ links (for example the links ML(4,1) and ML(4,2) are connected from middle switch MS(3,1) to middle switch MS(4,1), and the links ML(4,3) and ML(4,4) are connected from middle switch MS(3,1) to middle switch MS(4,15)).

- Each of the N/d middle switches MS(4,1) MS(4,16) in the middle stage 160 are connected from exactly d input switches through two links each for a total of 2×d links (for example the links ML(4,1) and ML(4,2) are connected to the middle switch MS(4,1) from input switch MS(3,1), and the links ML(4,59) and ML(4,60) are connected to the middle switch MS(4,1) from input switch MS(3,15)) and also are connected to exactly d switches in middle stage 170 through two links each for a total of 2×d links (for example the links ML(5,1) and ML(5,2) are connected from middle switch MS(4,1) to
 - middle switch MS(5,1), and the links ML(5,3) and ML(5,4) are connected from middle switch MS(4,1) to middle switch MS(5,15)).

Each of the
$$\frac{N}{d}$$
 middle switches MS(5,1) – MS(5,16) in the middle stage 170 are

connected from exactly *d* input switches through two links each for a total of $2 \times d$ links (for example the links ML(5,1) and ML(5,2) are connected to the middle switch MS(5,1)

from input switch MS(4,1), and the links ML(5,59) and ML(5,60) are connected to the middle switch MS(5,1) from input switch MS(4,15)).

Finally the connection topology of the network 100A shown in FIG. 1A is also basically back to back inverse Benes connection topology but with a slight variation. All
the cross middle links from middle switches MS(3,1) - MS(3,8) connect to middle switches MS(4,9) – MS(4,16) and all the cross middle links from middle switches MS(3,9) - MS(3,16) connect to middle switches MS(4,1) – MS(4,8). Applicant makes a key observation that there are many combinations of connections possible using this property. The difference in the connection topology between diagram 100A of FIG. 1A

- 10 and diagram 300A of FIG. 3A is that the connections formed by cross middle links between middle stage 150 and middle stage 160 are made of two different combinations otherwise both the diagrams 100A and 300A implement back to back inverse Benes connection topology. Since these networks implement back to back inverse Benes topologies since there is difference in the connections of cross middle links between
- 15 middle stage 150 and middle stage 160, the same difference in the connections of cross middle links between 160 and middle stage 170 occurs.

Referring to diagram 300B in FIG. 3B, is a folded version of the multi-link multistage network 300A shown in FIG. 3A. The network 300B in FIG. 3B shows input stage 110 and output stage 120 are placed together. That is input switch IS1 and output switch

OS1 are placed together, input switch IS2 and output switch OS2 are placed together, and similarly input switch IS16 and output switch OS16 are placed together. All the right going middle links {i.e., inlet links IL1 – IL32 and middle links ML(1,1) - ML(1,64)} correspond to input switches IS1 - IS16, and all the left going middle links {i.e., middle links ML(7,1) - ML(7,64) and outlet links OL1-OL32} correspond to output switches
OS1 - OS16.

Just the same way there is difference in the connection topology between diagram 100A of FIG. 1A and diagram 300A of FIG. 3A in the way the connections are formed by cross middle links between middle stage 150 and middle stage 160 and also between middle stage 160 and middle stage 170, the exact similar difference is there between the

30 diagram 100B of FIG. 1B and the diagram 300B of FIG. 3B, i.e., in the way the

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connections are formed by cross middle links between middle stage 150 and middle stage 160 and also between middle stage 160 and middle stage 170.

In one embodiment, in the network 300B of FIG. 3B, the switches that are placed together are implemented as separate switches then the network 300B is the generalized folded multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with nine stages as disclosed in PCT Application Serial No. PCT/US08/64604 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a two by four switch and a four by two switch. For example the switch input switch IS1 and output switch OS1 are placed together; so input switch IS1 is implemented as two by four switch with the inlet links IL1 and IL2 being the inputs of the input switch IS1; and output switch OS1 is implemented as four by two switch with the middle links ML(8,1), ML(8,2), ML(8,7) and ML(8,8) being the inputs of the output switch OS1 and outlet links OL1 – OL2 being the outputs of the

15 output switch OS1. Similarly in this embodiment of network 300B all the switches that are placed together are implemented as separate switches.

Referring to layout 300C of FIG. 3C, in one embodiment, there are sixteen blocks namely Block 1_2, Block 3_4, Block 5_6, Block 7_8, Block 9_10, Block 11_12, Block 13_14, Block 15_16, Block 17_18, Block 19_20, Block 21_22, Block 23_24, Block

- 20 25_26, Block 27_28, Block 29_30, and Block 31_32. Each block implements all the switches in one row of the network 300B of FIG. 3B, one of the key aspects of the current invention. For example Block 1_2 implements the input switch IS1, output Switch OS1, middle switch MS(1,1), middle switch MS(7,1), middle switch MS(2,1), middle switch MS(4,1).
- 25 For the simplification of illustration, Input switch IS1 and output switch OS1 together are denoted as switch 1; Middle switch MS(1,1) and middle switch MS(7,1) together are denoted by switch 2; Middle switch MS(2,1) and middle switch MS(6,1) together are denoted by switch 3; Middle switch MS(3,1) and middle switch MS(5,1) together are denoted by switch 4; And middle switch MS(4,1) is denoted by switch 5.

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All the straight middle links are illustrated in layout 300C of FIG. 3C. For example in Block 1_2, inlet links IL1 – IL2, outlet links OL1 – OL2, middle link ML(1,1), middle link ML(1,2), middle link ML(8,1), middle link ML(8,2), middle link ML(2,1), middle link ML(2,2), middle link ML(7,1), middle link ML(7,2), middle link ML(3,1), middle link ML(3,2), middle link ML(6,1), middle link ML(6,2), middle link ML(4,1), middle link ML(4,2), middle link ML(5,1) and middle link ML(5,2) are illustrated in layout 300C of FIG. 3C.

Even though it is not illustrated in layout 300C of FIG. 3C, in each block, in addition to the switches there may be Configurable Logic Blocks (CLB) or any arbitrary
digital circuit or sub-integrated circuit block depending on the applications in different embodiments. There are four quadrants in the layout 300C of FIG. 3C namely top-left, bottom-left, top-right and bottom-right quadrants. Top-left quadrant implements Block 1_2, Block 3_4, Block 5_6, and Block 7_8. Bottom-left quadrant implements Block 9_10, Block 11_12, Block 13_14, and Block 15_16. Top-right quadrant implements
Block 25_26, Block 27_28, Block 29_30, and Block 31_32. Bottom-right quadrant implements Block 17_18, Block 19_20, Block 21_22, and Block 23_24. There are two halves in layout 300C of FIG. 3C namely left-half and right-half. Left-half consists of

top-left and bottom-left quadrants. Right-half consists of top-right and bottom-right quadrants.

- 20 Recursively in each quadrant there are four sub-quadrants. For example in top-left quadrant there are four sub-quadrants namely top-left sub-quadrant, bottom-left sub-quadrant, top-right sub-quadrant and bottom-right sub-quadrant. Top-left sub-quadrant of top-left quadrant implements Block 1_2. Bottom-left sub-quadrant of top-left quadrant implements Block 3_4. Top-right sub-quadrant of top-left quadrant implements Block
- 25 7_8. Finally bottom-right sub-quadrant of top-left quadrant implements Block 5_6. Similarly there are two sub-halves in each quadrant. For example in top-left quadrant there are two sub-halves namely left-sub-half and right-sub-half. Left-sub-half of top-left quadrant implements Block 1_2 and Block 3_4. Right-sub-half of top-left quadrant implements Block 7_8 and Block 5_6. Recursively in larger multi-stage network
- 30 $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 > 32$, the layout in this embodiment in accordance

with the current invention, will be such that the super-quadrants will also be arranged in a similar manner.

Layout 300D of FIG. 3D illustrates the inter-block links (in the layout 300C of FIG. 3C all the cross middle links are inter-block links) between switches 1 and 2 of each block. For example middle links ML(1,3), ML(1,4), ML(8,7), and ML(8,8) are connected between switch 1 of Block 1_2 and switch 2 of Block 3_4. Similarly middle links ML(1,7), ML(1,8), ML(8,3), and ML(8,4) are connected between switch 2 of Block 1_2 and switch 1 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 100D of FIG. 1D can be implemented as vertical tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for example

middle links ML(1,4) and ML(8,8) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track).

15 Layout 300E of FIG. 3E illustrates the inter-block links between switches 2 and 3 of each block. For example middle links ML(2,3), ML(2,4), ML(7,11), and ML(7,12) are connected between switch 2 of Block 1_2 and switch 3 of Block 3_4. Similarly middle links ML(2,11), ML(2,12), ML(7,3), and ML(7,4) are connected between switch 3 of Block 1_2 and switch 2 of Block 3_4. Applicant notes that the inter-block links illustrated

- in layout 300E of FIG. 3E can be implemented as diagonal tracks in one embodiment.
 Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(2,12) and ML(7,4) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(2,12) and ML(7,4) are
- 25 implemented as a time division multiplexed single track).

Layout 300F of FIG. 3F illustrates the inter-block links between switches 3 and 4 of each block. For example middle links ML(3,3), ML(3,4), ML(6,19), and ML(6,20) are connected between switch 3 of Block 1_2 and switch 4 of Block 3_4. Similarly middle links ML(3,19), ML(3,20), ML(6,3), and ML(6,4) are connected between switch 4 of

30 Block 1_2 and switch 3 of Block 3_4. Applicant notes that the inter-block links illustrated

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in layout 300F of FIG. 3F can be implemented as vertical tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(3,4) and ML(6,20) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track).

Layout 300G of FIG. 3G illustrates the inter-block links between switches 4 and 5 of each block. For example middle links ML(4,3), ML(4,4), ML(5,35), and ML(5,36) are connected between switch 4 of Block 1_2 and switch 5 of Block 3_4. Similarly middle

10 links ML(4,35), ML(4,36), ML(5,3), and ML(5,4) are connected between switch 5 of Block 1_2 and switch 4 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 300G of FIG. 3G can be implemented as horizontal tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(4,4) and ML(5,36) are implemented as two different tracks); or

15 in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(4,4) and ML(5,36) are implemented as a time division multiplexed single track).

The complete layout for the network 300B of FIG. 3B is given by combining the links in layout diagrams of 300C, 300D, 300E, 300F, and 300G. Applicant notes that in 20 the layout 300C of FIG. 3C, the inter-block links between switch 1 and switch 2 are vertical tracks as shown in layout 300D of FIG. 3D; the inter-block links between switch 2 and switch 3 are horizontal tracks as shown in layout 300E of FIG. 3E; the inter-block links between switch 3 and switch 4 are vertical tracks as shown in layout 300F of FIG. 3F; and finally the inter-block links between switch 4 and switch 5 are horizontal tracks 25 as shown in layout 300G of FIG. 3G. The pattern is either vertical tracks, horizontal

tracks or diagonal tracks. It continues recursively for larger networks of N > 32 as will be illustrated later.

Some of the key aspects of the current invention related to layout diagram 300C of IFG. 3C are noted. 1) All the switches in one row of the multi-stage network 300B are implemented in a single block. 2) The blocks are placed in such a way that all the inter-

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block links are either horizontal tracks, vertical tracks or diagonal tracks; 3) The length of the longest wire is about half of the width (or length) of the complete layout (For example middle link ML(4,4) is about half the width of the complete layout.);

The layout 300C in FIG. 3C can be recursively extended for any arbitrarily large 5 generalized folded multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$. Referring to layout 300H of FIG. 3H, illustrates the extension of layout 300C for the network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 128$; d = 2; and s = 2. There are four superquadrants in layout 300H namely top-left super-quadrant, bottom-left super-quadrant, top-right super-quadrant, bottom-right super-quadrant. Total number of blocks in the 10 layout 300H is sixty four. Top-left super-quadrant implements the blocks from block 1_2 to block 31_32. Each block in all the super-quadrants has two more switches namely switch 6 and switch 7 in addition to the switches [1-5] illustrated in layout 300C of FIG. 3C. The inter-block link connection topology is the exactly the same between the switches 1 and 2; switches 2 and 3; switches 3 and 4; switches 4 and 5 as it is shown in

15 the layouts of FIG. 3D, FIG. 3E, FIG. 3F, and FIG. 3G respectively.

Bottom-left super-quadrant implements the blocks from block 33_34 to block 63_64. Top-right super-quadrant implements the blocks from block 65_66 to block 95_96. And bottom-right super-quadrant implements the blocks from block 97_98 to block 127_128. In all these three super-quadrants also, the inter-block link connection

topology is the exactly the same between the switches 1 and 2; switches 2 and 3; switches3 and 4; switches 4 and 5 as that of the top-left super-quadrant.

Recursively in accordance with the current invention, the inter-block links connecting the switch 5 and switch 6 will be vertical tracks between the corresponding switches of top-left super-quadrant and bottom-left super-quadrant. And similarly the

25 inter-block links connecting the switch 5 and switch 6 will be vertical tracks between the corresponding switches of top-right super-quadrant and bottom-right super-quadrant. The inter-block links connecting the switch 6 and switch 7 will be horizontal tracks between the corresponding switches of top-left super-quadrant and top-right super-quadrant. And similarly the inter-block links connecting the switch 6 and switch 7 will be horizontal

tracks between the corresponding switches of bottom-left super-quadrant and bottomright super-quadrant.

Ring Topology layout schemes:

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Layout diagram 400C of FIG. 4C is another embodiment for the generalized folded multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ diagram 100B in FIG. 1B.

Referring to layout 400C of FIG. 4C, there are sixteen blocks namely Block 1_2, Block 3_4, Block 5_6, Block 7_8, Block 9_10, Block 11_12, Block 13_14, Block 15_16, Block 17_18, Block 19_20, Block 21_22, Block 23_24, Block 25_26, Block 27_28,

- 10 Block 29_30, and Block 31_32. Each block implements all the switches in one row of the network 100B of FIG. 1B, one of the key aspects of the current invention. For example Block 1_2 implements the input switch IS1, output Switch OS1, middle switch MS(1,1), middle switch MS(7,1), middle switch MS(2,1), middle switch MS(6,1), middle switch MS(3,1), middle switch MS(5,1), and middle switch MS(4,1). For the simplification of
- 15 illustration, Input switch IS1 and output switch OS1 together are denoted as switch 1;
 Middle switch MS(1,1) and middle switch MS(7,1) together are denoted by switch 2;
 Middle switch MS(2,1) and middle switch MS(6,1) together are denoted by switch 3;
 Middle switch MS(3,1) and middle switch MS(5,1) together are denoted by switch 4;
 And middle switch MS(4,1) is denoted by switch 5.
- All the straight middle links are illustrated in layout 400C of FIG. 4C. For example in Block 1_2, inlet links IL1 IL2, outlet links OL1 OL2, middle link ML(1,1), middle link ML(1,2), middle link ML(8,1), middle link ML(8,2), middle link ML(2,1), middle link ML(2,2), middle link ML(7,1), middle link ML(7,2), middle link ML(3,1), middle link ML(3,2), middle link ML(6,1), middle link ML(6,2), middle link
- 25 ML(4,1), middle link ML(4,2), middle link ML(5,1) and middle link ML(5,2) are illustrated in layout 400C of FIG. 4C.

Even though it is not illustrated in layout 400C of FIG. 4C, in each block, in addition to the switches there may be Configurable Logic Blocks (CLB) or any arbitrary

digital circuit or sub-integrated circuit block depending on the applications in different embodiments. The topology of the layout 400C in FIG. 4C is a ring. For each of the neighboring rows in diagram 100B of FIG. 1B the corresponding blocks are also physically neighbors in layout diagram 400C of FIG. 4C. In addition the topmost row is

- also logically considered as neighbor to the bottommost row. For example Block 1_2 (implementing the switches belonging to a row in diagram 100B of FIG. 1B) has Block 3_4 as neighbor since Block 3_4 implements the switches in its neighboring row.
 Similarly Block 1_2 also has Block 31_32 as neighbor since Block 1_2 implements topmost row of switches and Block 31_32 implements bottommost row of switches in
- 10 diagram 100B of FIG. 1B. The ring layout scheme illustrated in 400C of FIG. 4C can be generalized for a large multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 > 32$, in accordance with the current invention.

Layout 400B of FIG. 4B illustrates the inter-block links (in the layout 400A of FIG. 4A all the cross middle links are inter-block links) between switches 1 and 2 of each

- block. For example middle links ML(1,3), ML(1,4), ML(8,7), and ML(8,8) are connected between switch 1 of Block 1_2 and switch 2 of Block 3_4. Similarly middle links ML(1,7), ML(1,8), ML(8,3), and ML(8,4) are connected between switch 2 of Block 1_2 and switch 1 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 400B of FIG. 4B are implemented as vertical tracks or horizontal tracks or diagonal
- 20 tracks. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(1,4) and ML(8,8) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track).
- 25 Layout 400C of FIG. 4C illustrates the inter-block links between switches 2 and 3 of each block. For example middle links ML(2,3), ML(2,4), ML(7,11), and ML(7,12) are connected between switch 2 of Block 1_2 and switch 3 of Block 3_4. Similarly middle links ML(2,11), ML(2,12), ML(7,3), and ML(7,4) are connected between switch 3 of Block 1_2 and switch 2 of Block 3_4. Applicant notes that the inter-block links illustrated
- 30 in layout 400C of FIG. 4C are implemented as vertical tracks or horizontal tracks or

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diagonal tracks. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(2,12) and ML(7,4) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(2,12) and ML(7,4) are implemented as a time division multiplexed single track).

Layout 400D of FIG. 4D illustrates the inter-block links between switches 3 and 4 of each block. For example middle links ML(3,3), ML(3,4), ML(6,19), and ML(6,20) are connected between switch 3 of Block 1_2 and switch 4 of Block 3_4. Similarly middle links ML(3,19), ML(3,20), ML(6,3), and ML(6,4) are connected between switch 4 of

- 10 Block 1_2 and switch 3 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 400D of FIG. 4D are implemented as vertical tracks or horizontal tracks or diagonal tracks. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(3,4) and ML(6,20) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a
- 15 time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track).

Layout 400E of FIG. 4E illustrates the inter-block links between switches 4 and 5 of each block. For example middle links ML(4,3), ML(4,4), ML(5,35), and ML(5,36) are connected between switch 4 of Block 1_2 and switch 5 of Block 3_4. Similarly middle

- 20 links ML(4,35), ML(4,36), ML(5,3), and ML(5,4) are connected between switch 5 of Block 1_2 and switch 4 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 400E of FIG. 4E are implemented as vertical tracks or horizontal tracks or diagonal tracks. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(4,4) and ML(5,36) are implemented as two
- 25 different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(4,4) and ML(5,36) are implemented as a time division multiplexed single track).

The complete layout for the network 100B of FIG. 1B is given by combining the links in layout diagrams of 400A, 400B, 400C, 400D, and 400E.

Some of the key aspects of the current invention related to layout diagram 400A of FIG. 4A are noted. 1) All the switches in one row of the multi-stage network 100B are implemented in a single block. 2) The blocks are placed in such a way that all the interblock links are either horizontal tracks, vertical tracks or diagonal tracks; 3) Length of the

5 different wires between the same two middle stages is not the same. However it gives an opportunity to implement the most connected circuits to place and route through the blocks which have shorter wires.

Layout diagram 400C1 of FIG. 4C1 is another embodiment for the generalized folded multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ diagram 100B in FIG. 1B.

- Referring to layout 400C1 of FIG. 4C1, there are sixteen blocks namely Block 1_2, Block 3_4, Block 5_6, Block 7_8, Block 9_10, Block 11_12, Block 13_14, Block 15_16, Block 17_18, Block 19_20, Block 21_22, Block 23_24, Block 25_26, Block 27_28, Block 29_30, and Block 31_32. Each block implements all the switches in one row of the network 100B of FIG. 1B, one of the key aspects of the current invention. For example
- Block 1_2 implements the input switch IS1, output Switch OS1, middle switch MS(1,1), middle switch MS(7,1), middle switch MS(2,1), middle switch MS(6,1), middle switch MS(3,1), middle switch MS(5,1), and middle switch MS(4,1). For the simplification of illustration, Input switch IS1 and output switch OS1 together are denoted as switch 1; Middle switch MS(1,1) and middle switch MS(7,1) together are denoted by switch 2;
- Middle switch MS(2,1) and middle switch MS(6,1) together are denoted by switch 3;
 Middle switch MS(3,1) and middle switch MS(5,1) together are denoted by switch 4;
 And middle switch MS(4,1) is denoted by switch 5.

All the straight middle links are illustrated in layout 400C1 of FIG. 4C1. For example in Block 1_2, inlet links IL1 – IL2, outlet links OL1 – OL2, middle link

ML(1,1), middle link ML(1,2), middle link ML(8,1), middle link ML(8,2), middle link ML(2,1), middle link ML(2,2), middle link ML(7,1), middle link ML(7,2), middle link ML(3,1), middle link ML(3,2), middle link ML(6,1), middle link ML(6,2), middle link ML(4,1), middle link ML(4,2), middle link ML(5,1) and middle link ML(5,2) are illustrated in layout 400C1 of FIG. 4C1.

Even though it is not illustrated in layout 400C1 of FIG. 4C1, in each block, in addition to the switches there may be Configurable Logic Blocks (CLB) or any arbitrary digital circuit or sub-integrated circuit block depending on the applications in different embodiments. The topology of the layout 400C1 in FIG. 4C1 is another embodiment of

5 ring layout topology. For each of the neighboring rows in diagram 100B of FIG. 1B the corresponding blocks are also physically neighbors in layout diagram 400C of FIG. 4C. In addition the topmost row is also logically considered as neighbor to the bottommost row. For example Block 1_2 (implementing the switches belonging to a row in diagram 100B of FIG. 1B) has Block 3_4 as neighbor since Block 3_4 implements the switches in

10 its neighboring row. Similarly Block 1_2 also has Block 31_32 as neighbor since Block 1_2 implements topmost row of switches and Block 31_32 implements bottommost row of switches in diagram 100B of FIG. 1B. The ring layout scheme illustrated in 400C of FIG. 4C can be generalized for a large multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 > 32$, in accordance with the current invention.

15 All the layout embodiments disclosed in the current invention are applicable to generalized multi-stage networks $V(N_1, N_2, d, s)$, generalized folded multi-stage networks $V_{fold}(N_1, N_2, d, s)$, generalized butterfly fat tree networks $V_{bft}(N_1, N_2, d, s)$, generalized multi-link multi-stage networks $V_{mlink}(N_1, N_2, d, s)$, generalized folded multilink multi-stage networks $V_{fold-mlink}(N_1, N_2, d, s)$, generalized multi-link butterfly fat tree

20 networks $V_{mlink-bft}(N_1, N_2, d, s)$, and generalized hypercube networks $V_{hcube}(N_1, N_2, d, s)$ for s = 1,2,3 or any number in general, and for both $N_1 = N_2 = N$ and $N_1 \neq N_2$, and d is any integer.

Conversely applicant makes another important observation that generalized hypercube networks $V_{hcube}(N_1, N_2, d, s)$ are implemented with the layout topology being 25 the hypercube topology shown in layout 100C of FIG. 1C with large scale cross point reduction as any one of the networks described in the current invention namely: generalized multi-stage networks $V(N_1, N_2, d, s)$, generalized folded multi-stage networks $V_{fold}(N_1, N_2, d, s)$, generalized butterfly fat tree networks $V_{bft}(N_1, N_2, d, s)$, generalized multi-link multi-stage networks $V_{mlink}(N_1, N_2, d, s)$, generalized folded multilink multi-stage networks $V_{fold-mlink}(N_1, N_2, d, s)$, generalized multi-link butterfly fat tree networks $V_{mlink-bft}(N_1, N_2, d, s)$ for s = 1,2,3 or any number in general, and for both $N_1 = N_2 = N$ and $N_1 \neq N_2$, and d is any integer.

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Applications Embodiments:

All the embodiments disclosed in the current invention are useful in many varieties of applications. FIG. 5A1 illustrates the diagram of 500A1 which is a typical two by two switch with two inlet links namely IL1 and IL2, and two outlet links namely

- OL1 and OL2. The two by two switch also implements four crosspoints namely CP(1,1), CP(1,2), CP(2,1) and CP(2,2) as illustrated in FIG. 5A1. For example the diagram of 500A1 may the implementation of middle switch MS(1,1) of the diagram 100K of FIG. 1K where inlet link IL1 of diagram 500A1 corresponds to middle link ML(1,1) of diagram 100K, inlet link IL2 of diagram 500A1 corresponds to middle link ML(1,7) of
- 15 diagram 100K, outlet link OL1 of diagram 500A1 corresponds to middle link ML(2,1) of diagram 100K, outlet link OL2 of diagram 500A1 corresponds to middle link ML(2,3) of diagram 100K.

1) Programmable Integrated Circuit Embodiments:

- 20 All the embodiments disclosed in the current invention are useful in programmable integrated circuit applications. FIG. 5A2 illustrates the detailed diagram 500A2 for the implementation of the diagram 500A1 in programmable integrated circuit embodiments. Each crosspoint is implemented by a transistor coupled between the corresponding inlet link and outlet link, and a programmable cell in programmable
- 25 integrated circuit embodiments. Specifically crosspoint CP(1,1) is implemented by transistor C(1,1) coupled between inlet link IL1 and outlet link OL1, and programmable cell P(1,1); crosspoint CP(1,2) is implemented by transistor C(1,2) coupled between inlet

link IL1 and outlet link OL2, and programmable cell P(1,2); crosspoint CP(2,1) is implemented by transistor C(2,1) coupled between inlet link IL2 and outlet link OL1, and programmable cell P(2,1); and crosspoint CP(2,2) is implemented by transistor C(2,2) coupled between inlet link IL2 and outlet link OL2, and programmable cell P(2,2).

5 If the programmable cell is programmed ON, the corresponding transistor couples the corresponding inlet link and outlet link. If the programmable cell is programmed OFF, the corresponding inlet link and outlet link are not connected. For example if the programmable cell P(1,1) is programmed ON, the corresponding transistor C(1,1) couples the corresponding inlet link IL1 and outlet link OL1. If the programmable cell P(1,1) is

- 10 programmed OFF, the corresponding inlet link IL1 and outlet link OL1 are not connected. In volatile programmable integrated circuit embodiments the programmable cell may be an SRAM (Static Random Address Memory) cell. In non-volatile programmable integrated circuit embodiments the programmable cell may be a Flash memory cell. Also the programmable integrated circuit embodiments may implement
- 15 field programmable logic arrays (FPGA) devices, or programmable Logic devices (PLD), or Application Specific Integrated Circuits (ASIC) embedded with programmable logic circuits or 3D-FPGAs.

FIG. 5A2 also illustrates a buffer B1 on inlet link IL2. The signals driven along inlet link IL2 are amplified by buffer B1. Buffer B1 can be inverting or non-inverting
buffer. Buffers such as B1 are used to amplify the signal in links which are usually long.

2) One-time Programmable Integrated Circuit Embodiments:

All the embodiments disclosed in the current invention are useful in one-time programmable integrated circuit applications. FIG. 5A3 illustrates the detailed diagram 500A3 for the implementation of the diagram 500A1 in one-time programmable

25 integrated circuit embodiments. Each crosspoint is implemented by a via coupled between the corresponding inlet link and outlet link in one-time programmable integrated circuit embodiments. Specifically crosspoint CP(1,1) is implemented by via V(1,1) coupled between inlet link IL1 and outlet link OL1; crosspoint CP(1,2) is implemented by via V(1,2) coupled between inlet link IL1 and outlet link OL2; crosspoint CP(2,1) is implemented by via V(2,1) coupled between inlet link IL2 and outlet link OL1; and crosspoint CP(2,2) is implemented by via V(2,2) coupled between inlet link IL2 and outlet link OL2.

If the via is programmed ON, the corresponding inlet link and outlet link are permanently connected which is denoted by thick circle at the intersection of inlet link and outlet link. If the via is programmed OFF, the corresponding inlet link and outlet link are not connected which is denoted by the absence of thick circle at the intersection of inlet link and outlet link. For example in the diagram 500A3 the via V(1,1) is programmed ON, and the corresponding inlet link IL1 and outlet link OL1 are connected

- 10 as denoted by thick circle at the intersection of inlet link IL1 and outlet link OL1; the via V(2,2) is programmed ON, and the corresponding inlet link IL2 and outlet link OL2 are connected as denoted by thick circle at the intersection of inlet link IL2 and outlet link OL2; the via V(1,2) is programmed OFF, and the corresponding inlet link IL1 and outlet link OL2 are not connected as denoted by the absence of thick circle at the intersection of
- 15 inlet link IL1 and outlet link OL2; the via V(2,1) is programmed OFF, and the corresponding inlet link IL2 and outlet link OL1 are not connected as denoted by the absence of thick circle at the intersection of inlet link IL2 and outlet link OL1. One-time programmable integrated circuit embodiments may be anti-fuse based programmable integrated circuit devices or mask programmable structured ASIC devices.

20 3) Integrated Circuit Placement and Route Embodiments:

All the embodiments disclosed in the current invention are useful in Integrated Circuit Placement and Route applications, for example in ASIC backend Placement and Route tools. FIG. 5A4 illustrates the detailed diagram 500A4 for the implementation of the diagram 500A1 in Integrated Circuit Placement and Route embodiments. In an

25 integrated circuit since the connections are known a-priori, the switch and crosspoints are actually virtual. However the concept of virtual switch and virtal crosspoint using the embodiments disclosed in the current invention reduces the number of required wires, wire length needed to connect the inputs and outputs of different netlists and the time required by the tool for placement and route of netlists in the integrated circuit.

Each virtual crosspoint is used to either to hardwire or provide no connectivity between the corresponding inlet link and outlet link. Specifically crosspoint CP(1,1) is implemented by direct connect point DCP(1,1) to hardwire (i.e., to permanently connect) inlet link IL1 and outlet link OL1 which is denoted by the thick circle at the intersection

- 5 of inlet link IL1 and outlet link OL1; crosspoint CP(2,2) is implemented by direct connect point DCP(2,2) to hardwire inlet link IL2 and outlet link OL2 which is denoted by the thick circle at the intersection of inlet link IL2 and outlet link OL2. The diagram 500A4 does not show direct connect point DCP(1,2) and direct connect point DCP(1,3) since they are not needed and in the hardware implementation they are eliminated.
- 10 Alternatively inlet link IL1 needs to be connected to outlet link OL1 and inlet link IL1 does not need to be connected to outlet link OL2. Also inlet link IL2 needs to be connected to outlet link OL2 and inlet link IL2 does not need to be connected to outlet link OL1. Furthermore in the example of the diagram 500A4, there is no need to drive the signal of inlet link IL1 horizontally beyond outlet link OL1 and hence the inlet link IL1 is
- 15 not even extended horizontally until the outlet link OL2. Also the absence of direct connect point DCP(2,1) illustrates there is no need to connect inlet link IL2 and outlet link OL1.

In summary in integrated circuit placement and route tools, the concept of virtual switches and virtual cross points is used during the implementation of the placement & 20 routing algorithmically in software, however during the hardware implementation cross points in the cross state are implemented as hardwired connections between the corresponding inlet link and outlet link, and in the bar state are implemented as no connection between inlet link and outlet link.

3) More Application Embodiments:

All the embodiments disclosed in the current invention are also useful in the design of SoC interconnects, Field programmable interconnect chips, parallel computer systems and in time-space-time switches.

Numerous modifications and adaptations of the embodiments, implementations, and examples described herein will be apparent to the skilled artisan in view of the disclosure.

<u>CLAIMS</u>

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What is claimed is:

1. An integrated circuit device comprising a plurality of sub-integrated circuit blocks and a routing network, and

Said each plurality of sub-integrated circuit blocks comprising a plurality of inlet links and a plurality of outlet links; and

Said routing network interconnects any one of said outlet link of one of said subintegrated circuit block to one or more said inlet links of one or more of said subintegrated circuit blocks; and

10 Said routing network comprising of a plurality of stages *y* , starting from the lowest stage to the highest stage; and

Said routing network comprising a plurality of switches of size $d \times d$, where $d \ge 2$, in each said stage and each said switch of size $d \times d$ having d inlet links and d outlet links; and

15 Said each sub-integrated circuit block comprising a plurality of said switches corresponding to each said stage; and

Said each sub-integrated circuit block comprising a plurality of forward connecting links connecting from switches in lower stage to switches in the immediate succeeding higher stage, and also comprising a plurality of backward connecting links connecting from switches in higher stage to switches in the immediate preceding lower stage; and

Said each sub-integrated circuit block comprising a plurality straight links in said forward connecting links from switches in lower stage to switches in the immediate succeeding higher stage and a plurality cross links in said forward connecting links from

25 switches in lower stage to switches in the immediate succeeding higher stage, and further comprising a plurality of straight links in said backward connecting links from switches

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in higher stage to switches in the immediate preceding lower stage and a plurality of cross links in said backward connecting links from switches in higher stage to switches in the immediate preceding lower stage.

The integrated circuit device of claim 1, wherein said all straight links are
 connecting from switches in each said sub-integrated circuit block are connecting to
 switches in the same said sub-integrated circuit block; and

said all cross links are connecting as either vertical or horizontal links between switches in two different said sub-integrated circuit blocks.

The integrated circuit device of claim 2, wherein said plurality of sub integrated circuit blocks arranged in a two-dimensional grid.

4. The integrated circuit device of claim 3, wherein said cross links in succeeding stages are connecting as alternative vertical and horizontal links between switches in said sub-integrated circuit blocks.

- 5. The integrated circuit device of claim 4, wherein said cross links from 15 switches in a stage in one of said sub-integrated circuit blocks are connecting to switches in the succeeding stage in another of said sub-integrated circuit blocks so that said cross links are either vertical links or horizontal and vice versa, and hereinafter such cross links are "shuffle exchange links").
- 6. The integrated circuit device of claim 5, wherein said all horizontal shuffle 20 exchange links between switches in any two corresponding said succeeding stages are substantially of equal length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are substantially of equal length in the entire said integrated circuit device.
- 7. The integrated circuit device of claim 6, wherein the shortest horizontal25 shuffle exchange links are connecting at the lowest stage and between switches in two

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nearest neighboring said sub-integrated circuit blocks, and length of the horizontal shuffle exchange links is doubled in each succeeding stage; and the shortest vertical shuffle exchange links are connecting at the lowest stage and between switches in two nearest neighboring said sub-integrated circuit blocks, and length of the vertical shuffle exchange links is doubled in each succeeding stage.

8. The integrated circuit device of claim 7, wherein $y \ge (\log_2 N)$ so that the length of the horizontal shuffle exchange links in the highest stage is equal to half the size of the horizontal size of said two dimensional grid of sub-integrated circuit blocks and the length of the vertical shuffle exchange links in the highest stage is equal to half the size of the vertical size of said two dimensional grid of sub-integrated circuit blocks.

9. The integrated circuit device of claim 8, wherein d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said subintegrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast Benes network with full bandwidth.

10. The integrated circuit device of claim 8, wherein d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said
20 routing network is strictly nonblocking for unicast Benes network and rearrangeably nonblocking for arbitrary fan-out multicast Benes network with full bandwidth.

The integrated circuit device of claim 8, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in
 each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast Benes network with

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full bandwidth.

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12. The integrated circuit device of claim 7, wherein $y \ge (\log_2 N)$ so that the length of the horizontal shuffle exchange links in the highest stage is equal to half the size of the horizontal size of said two dimensional grid of sub-integrated circuit blocks and the length of the vertical shuffle exchange links in the highest stage is equal to half the size of

said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks.

the vertical size of said two dimensional grid of sub-integrated circuit blocks, and

13. The integrated circuit device of claim 12, wherein d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said
10 forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast butterfly fat tree network with full bandwidth.

14. The integrated circuit device of claim 12, wherein d = 2 and there are at 15 least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast butterfly fat tree network and rearrangeably nonblocking for arbitrary fan-out multicast butterfly fat tree network with

20 full bandwidth.

15. The integrated circuit device of claim 12, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said

25 routing network is strictly nonblocking for arbitrary fan-out multicast butterfly fat tree network with full bandwidth.

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16. The integrated circuit device of claim 1, wherein said horizontal and vertical links are implemented on two or more metal layers.

17. The integrated circuit device of claim 1, wherein said switches comprising active and reprogrammable cross points and said each cross point is programmable by an
5 SRAM cell or a Flash Cell.

18. The integrated circuit device of claim 1, wherein said sub-integrated circuit blocks are of equal die size.

19. The integrated circuit device of claim 16, wherein said sub-integrated circuit blocks are Lookup Tables (hereinafter "LUTs") and said integrated circuit device
10 is a field programmable gate array (FPGA) device or field programmable gate array (FPGA) block embedded in another integrated circuit device.

20. The integrated circuit device of claim 16, wherein said sub-integrated circuit blocks are AND or OR gates and said integrated circuit device is a programmable logic device (PLD).

15 21. The integrated circuit device of claim 1, wherein said sub-integrated circuit blocks comprising any arbitrary hardware logic or memory circuits.

22. The integrated circuit device of claim 1, wherein said switches comprising active one-time programmable cross points and said integrated circuit device is a mask programmable gate array (MPGA) device or a structured ASIC device.

20 23. The integrated circuit device of claim 1, wherein said switches comprising passive cross points or just connection of two links or not and said integrated circuit device is a Application Specific Integrated Circuit (ASIC) device.

24. The integrated circuit device of claim 1, wherein said sub-integrated circuit blocks further recursively comprise one or more super-sub-integrated circuit
25 blocks and a sub-routing network.

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25. The integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and $y \ge (\log_2 N)$.

5 26. The integrated circuit device of claim 25, wherein *d* = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said subintegrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized multi-stage network with 10 full bandwidth.

27. The integrated circuit device of claim 25, wherein d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast generalized multi-stage network and rearrangeably nonblocking for arbitrary fan-out multicast generalized multi-stage network with full bandwidth.

28. The integrated circuit device of claim 25, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast generalized multi-stage network with full bandwidth.

29. The integrated circuit device of claim 5, wherein said all horizontal shuffle 25 exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and $y \ge (\log_2 N)$, and

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said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks.

30. The integrated circuit device of claim 29, wherein d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said
5 forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized butterfly fat tree network with full bandwidth.

31. The integrated circuit device of claim 29, wherein d = 2 and there are at
least two switches in each said stage in each said sub-integrated circuit block connecting
said forward connecting links and there are at least two switches in each said stage in
each said sub-integrated circuit block connecting said backward connecting links and said
routing network is strictly nonblocking for unicast generalized butterfly fat tree Network
and rearrangeably nonblocking for arbitrary fan-out multicast generalized butterfly fat

32. The integrated circuit device of claim 29, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said 20 routing network is strictly nonblocking for arbitrary fan-out multicast generalized butterfly fat tree network with full bandwidth.

33. The integrated circuit device of claim 1, wherein said straight links connecting from switches in each said sub-integrated circuit block are connecting to switches in the same said sub-integrated circuit block; and

25

said cross links are connecting as vertical or horizontal or diagonal links between two different said sub-integrated circuit blocks. 34. The integrated circuit device of claim 8, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said subintegrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast multi-link Benes network with full

5 network is rearrangeably nonblocking for unicast multi-link Benes network bandwidth.

35. The integrated circuit device of claim 8, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in
each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast multi-link Benes network and rearrangeably nonblocking for arbitrary fan-out multicast multi-link Benes network with full bandwidth.

36. The integrated circuit device of claim 8, wherein d = 4 and there are at 15 least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast multi-link Benes network with full bandwidth.

20 37. The integrated circuit device of claim 12, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said subintegrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast multi-link butterfly fat tree network

with full bandwidth.

38. The integrated circuit device of claim 12, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in

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each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast multi-link butterfly fat tree network and rearrangeably nonblocking for arbitrary fan-out multicast multi-link butterfly fat tree network with full bandwidth.

5 39. The integrated circuit device of claim 12, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast multi-link butterfly fat tree network with full bandwidth.

40. The integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and $y \ge (\log_2 N)$.

- 15 41. The integrated circuit device of claim 40, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said subintegrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized multi-link multi-stage
- 20 network with full bandwidth.

42. The integrated circuit device of claim 40, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said

25 routing network is strictly nonblocking for unicast generalized multi-link multi-stage network and rearrangeably nonblocking for arbitrary fan-out multicast generalized multilink multi-stage network with full bandwidth.

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43. The integrated circuit device of claim 40, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast generalized multi-link multi-stage network with full bandwidth.

44. The integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two
10 corresponding said succeeding stages are of different length and y ≥ (log₂ N), and said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks.

45. The integrated circuit device of claim 44, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said
15 forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized multi-link butterfly fat tree network with full bandwidth.

46. The integrated circuit device of claim 44, wherein d = 4 and there are at
20 least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast generalized multi-link butterfly fat tree Network and rearrangeably nonblocking for arbitrary fan-out multicast generalized multi-link butterfly fat tree network with full bandwidth.

47. The integrated circuit device of claim 44, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting

said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast generalized multilink butterfly fat tree network with full bandwidth.

- 5 48. The integrated circuit device of claim 1, wherein said plurality of forward connecting links use a plurality of buffers to amplify signals driven through them and said plurality of backward connecting links use a plurality of buffers to amplify signals driven through them; and said buffers can be inverting or non-inverting buffers.
- 49. The integrated circuit device of claim 1, wherein said wherein said all 10 switches of size $d \times d$ are either fully populated or partially populated.

VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS

Venkat Konda

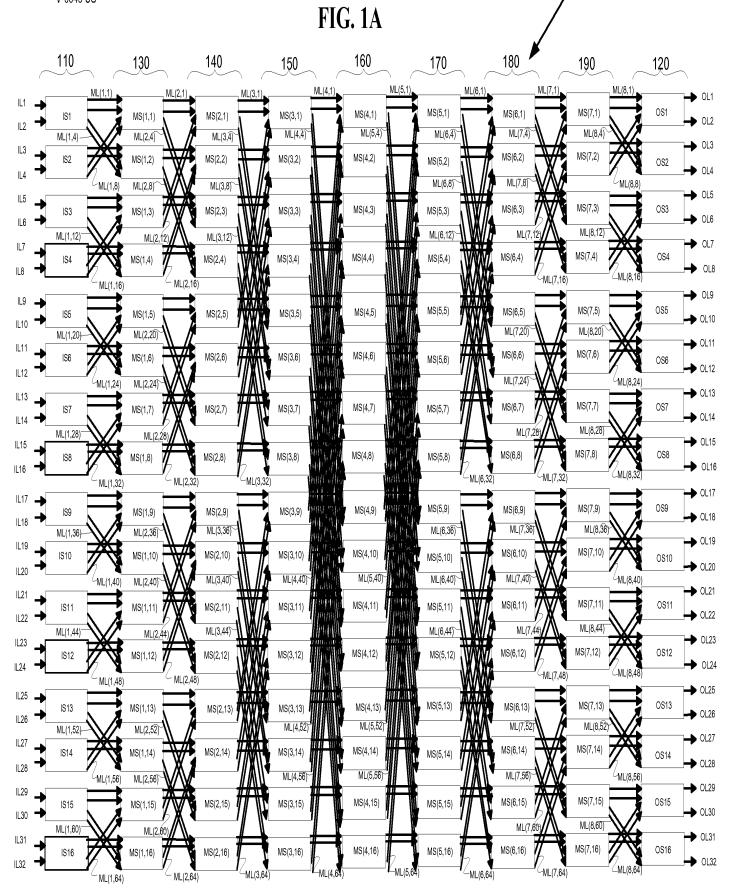
ABSTRACT OF DISCLOSURE

- In accordance with the invention, VLSI layouts of generalized multi-stage networks for broadcast, unicast and multicast connections are presented using only horizontal and vertical links. The VLSI layouts employ shuffle exchange links where outlet links of cross links from switches in a stage in one sub-integrated circuit block are connected to inlet links of switches in the succeeding stage in another sub-integrated circuit block so that said cross links are either vertical links or horizontal and vice versa.
- 10 In one embodiment the sub-integrated circuit blocks are arranged in a hypercube arrangement in a two-dimensional plane. The VLSI layouts exploit the benefits of significantly lower cross points, lower signal latency, lower power and full connectivity with significantly fast compilation.

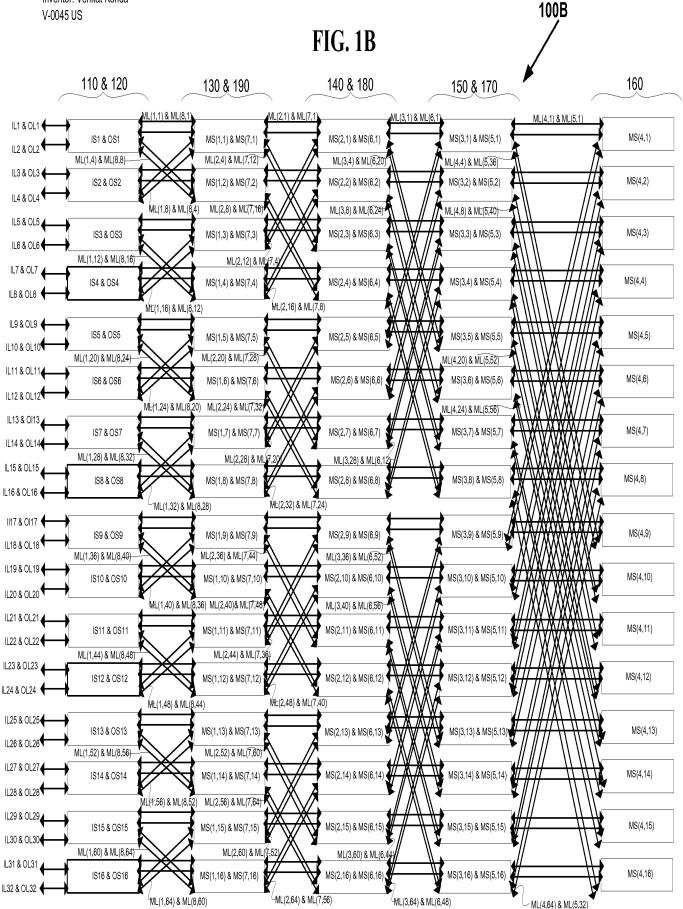
The VLSI layouts presented are applicable to generalized multi-stage networks 15 $V(N_1, N_2, d, s)$, generalized folded multi-stage networks $V_{fold}(N_1, N_2, d, s)$, generalized butterfly fat tree networks $V_{bft}(N_1, N_2, d, s)$, generalized multi-link multi-stage networks $V_{mlink}(N_1, N_2, d, s)$, generalized folded multi-link multi-stage networks $V_{fold-mlink}(N_1, N_2, d, s)$, generalized multi-link butterfly fat tree networks $V_{mlink-bft}(N_1, N_2, d, s)$, and generalized hypercube networks $V_{hcube}(N_1, N_2, d, s)$ for s =

20 1,2,3 or any number in general. The embodiments of VLSI layouts are useful in wide target applications such as FPGAs, CPLDs, pSoCs, ASIC placement and route tools, networking applications, parallel & distributed computing, and reconfigurable computing. Page 1 of 39

100A



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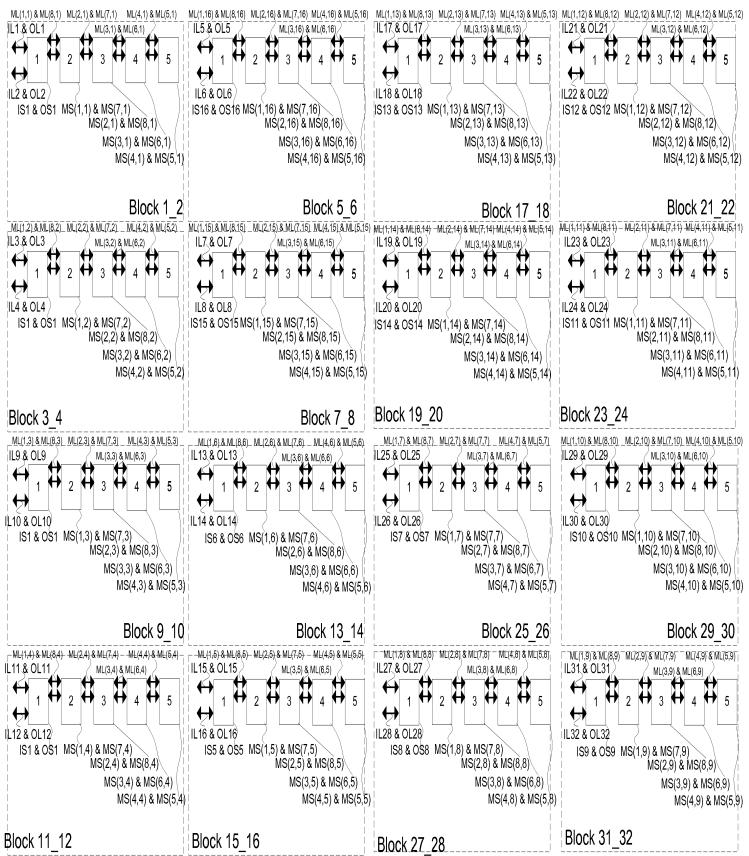


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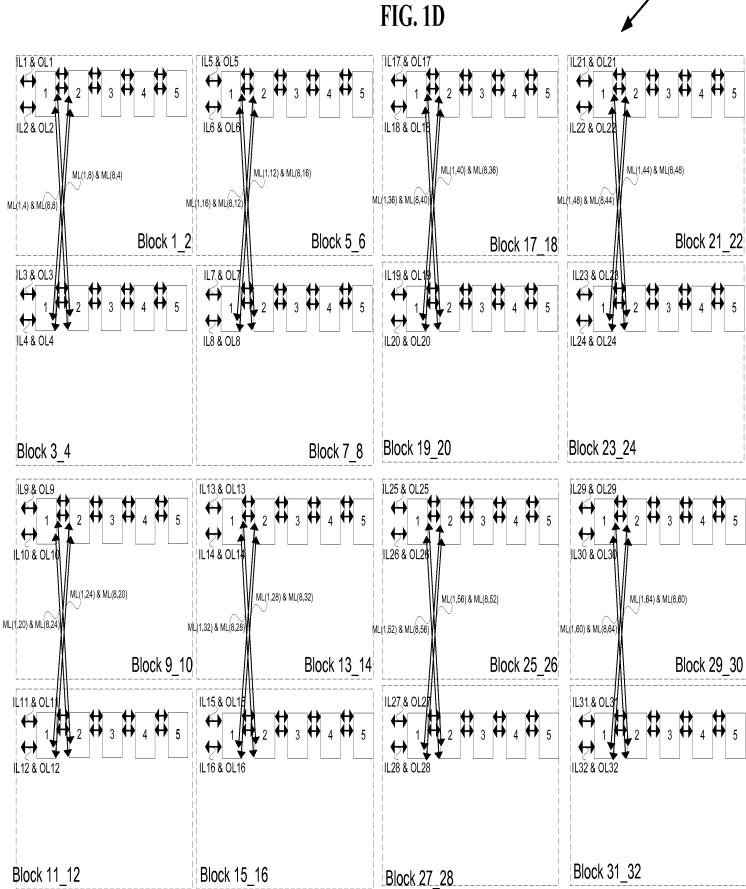
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FIG. 1C



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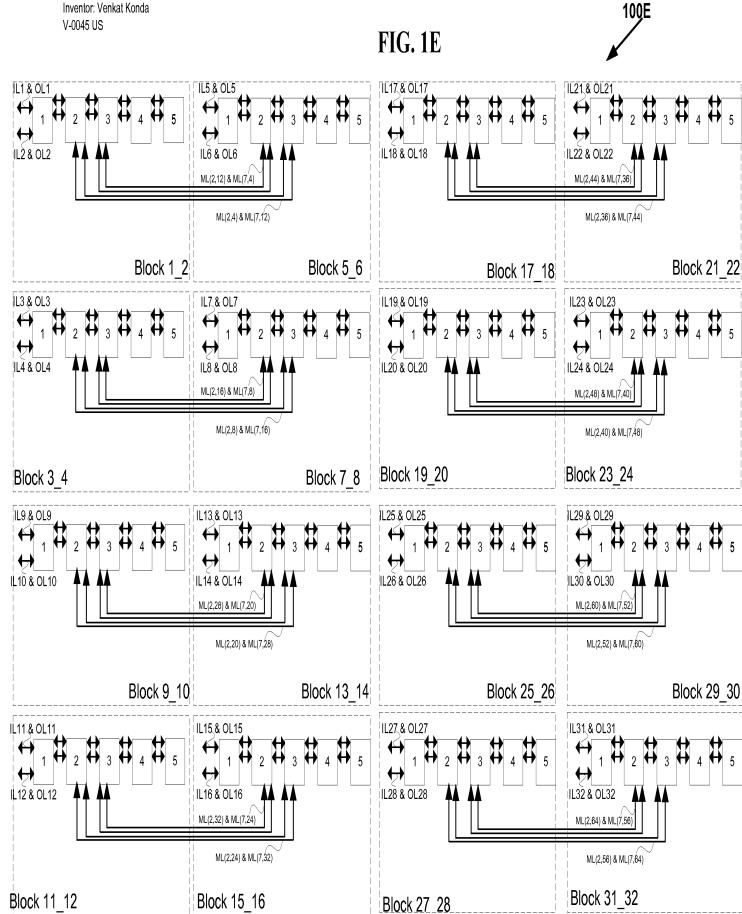




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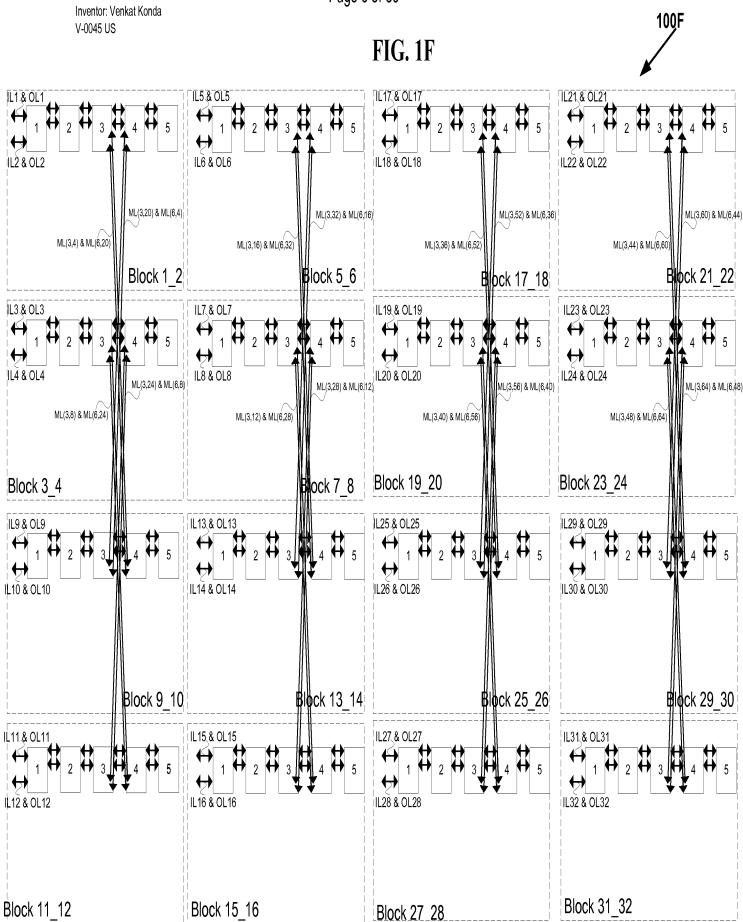
VLSI Layouts of Fully Connected Generalized Networks Inventor: Venkat Konda



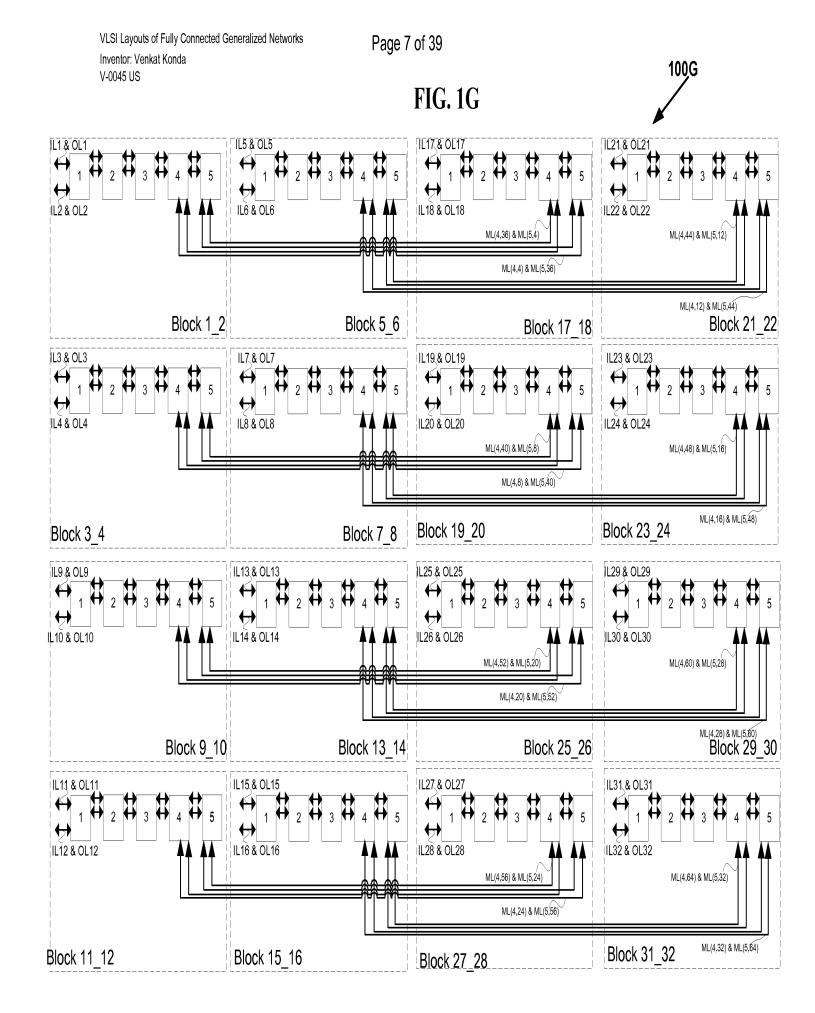
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VLSI Layouts of Fully Connected Generalized Networks

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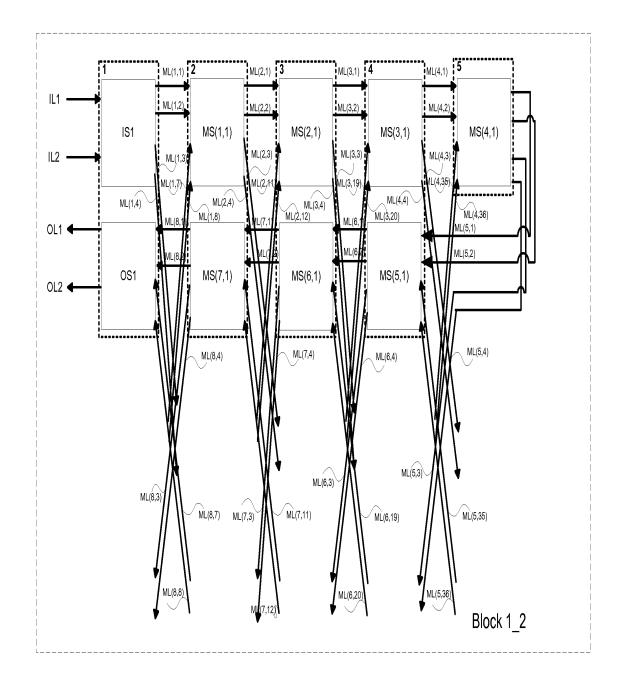
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VLSI Layouts of Fully Connected Generalized Networks Inventor: Venkat Konda V-0045 US		Page 8 of 39 FIG. 1H		100H		
1 2 3 4 5 6 7 1 1 2 3 4 5 6	129 & OL29 7 1 1 2 3 3 4 5 6 7 1130 & OL30	L27.&OL27 14 24 34 44 54 64 7 1L28 & OL28	11 21 31 41 51 61 7 12 8 0L2	H 7 & QL7 1 M 2 3 4 5 6 7 1L8 & OL3	129 & OL29 11 2 3 4 5 6 7 1230 & OL30	L27-& OL27 [11 2 3 4 5 6 7 [L28 & O'L28
Block 1 2. Block 5 6 13 & 013 14 24 34 45 66 7 1 14 24 34 45 66 14 & 014	Block 17 18 L3 & OL3 1 2 3 4 5 6 7 L32 & OL32	Block 21_22 125 & 0125 11 21 31 41 51 61 7 1126 & 0126		Block 69_70 15* 015 11 21 31 41 51 61 7 116 & 016	Block 81 82 137 & 0131 14 21 31 41 51 64 7 11 32 & 0132	Block 85_86 125 & 0125 1 14 24 34 44 54 64 7 126 & 0126
Block 3_4 15 & 019 14 24 34 45 66 7 1 14 24 34 45 66 11 10 & 01 10 11 10 & 01 10 11 10 & 01 10	Block 19 20 H-17 & OL 17 7 14 24 34 44 54 64 7 IL 18 & OL 18	Block 23 24 123 & 0123 11 24 31 41 51 66 7 1124 8 0124	Block 67_68	Block 71_72_ 115 & 0L15 11 21 31 41 51 61 7 1L16 & 0L16	Block 83_84	Block 87_88_ 123 & OL23
Block 9 10 Block 13 14 11 2 3 4 4 5 6 7 1 1 2 3 4 4 5 6 11 2 8 0L12	Block 25_26 119 & 019 11 21 31 41 51 61 7 120 & 0120	Block 29_30 121 8 01 21 11 21 31 41 51 61 7 122 8 01 22		Block 77_78_ 113&013 11 21 31 41 51 61 7 114 & 0114	Block 89 90 1519 80 19 14 21 31 41 51 61 7 1220 8 0120	Block 93 94
Block 11_12 Block 15_16	Block 27_28	Block 31_32	Block 75 76	Block 79 80	Block 91_92	Block 95_96
H1 & OL1 14 2 3 4 5 6 7 1 1 2 3 4 5 6 7 12 8 0L2 11 2 8 0L2	129 & OL29 14 24 34 44 54 64 7 1230 & OL30	127-8-0127 14 24 34 44 54 64 7 11.28 & 01.28	11 20 31 44 54 64 7 11 21 31 44 54 64 7	11 2 3 4 5 6 7 1L8 & OL3	123 & OL23 11 21 31 41 51 61 7 1L30 & OL30	122 3 3 4 5 6 7 1128 & O'28
Block 33_34 138 013 14 21 31 41 51 61 7 14 8 014 14 8 014 16 8 016		1 1 2 3 4 5 6 7		Block 101_102 H5 & CL5 11 21 31 41 51 61 7 H6 & OL6	Block 113 114 131 & 0131 14 21 31 41 51 61 7 11 32 & 0132	Block 117 118 125 & 0125 1 14 24 34 44 54 64 7 1126 & 0126
Block 35_36 15 & 0L9 14 21 31 41 51 61 7 11 21 31 41 51 61 7 11 10 & 0L10 12 16 & 0L16	Block 51_52_ 1.17 & 0.17 1.17 & 0.17 1.17 & 0.17 1.18 & 0.18	123 & OL23		Block 103_104 115&0L15 11 24 31 44 51 64 7 1116&0L16	Block 115_116 147.& 0L17 14.21.31.44.51.64.7 1L18 & 0L18	Block 119_120 123 & OL23 11 24 34 44 54 64 7 1L24 8 OL24
Block 41 42 H11&0L11 11 21 31 41 51 61 7 11 21 31 41 51 61 7 L12 & 0L12	Block. 57_58 1E19 & OE19 14 24 34 44 54 64 7 1L20 & OL20		111 20 30 40 50 60 7	Block 109 110 113 & 0L13 11 21 31 41 51 61 7 1L14 & 0L14	Block 121 122 198 0L19 11 21 31 41 51 61 7 120 8 0L20	Block 125_126 1218 0121 11 24 34 44 54 64 7 1122 8 0122
	Block 59_60	Block 63_64	Block 107_108	Block 111_112	Block 123_124	Block 127_128

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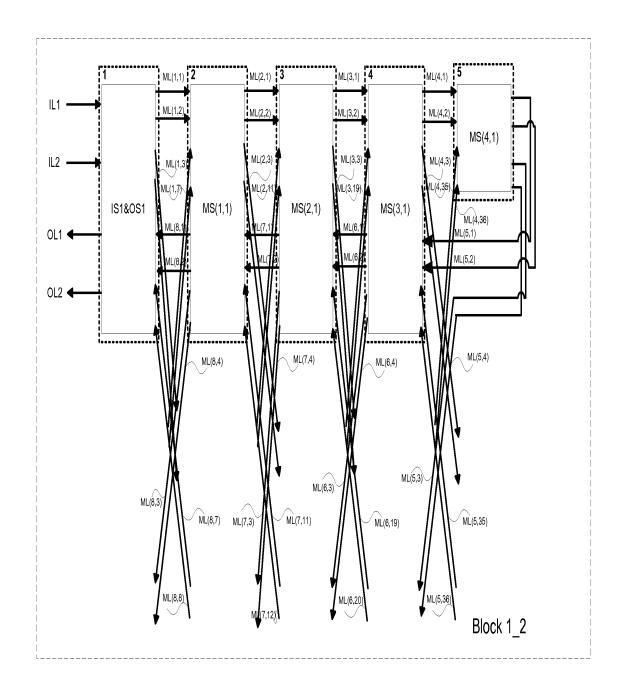






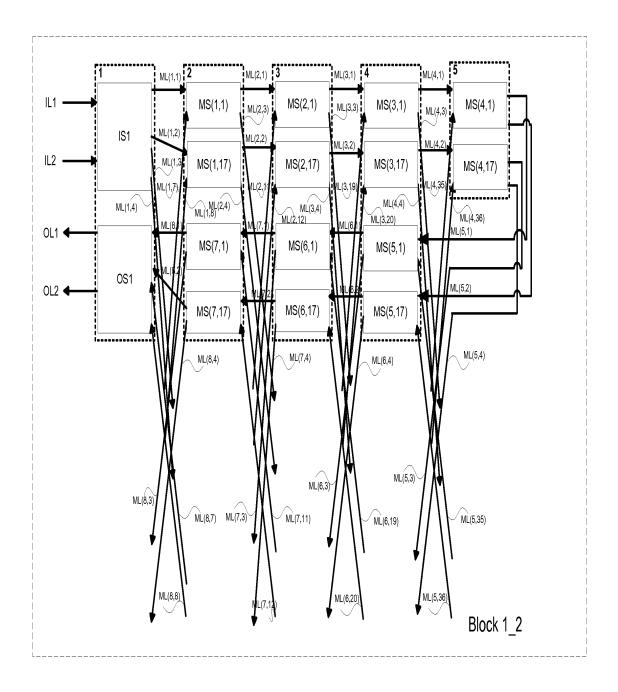






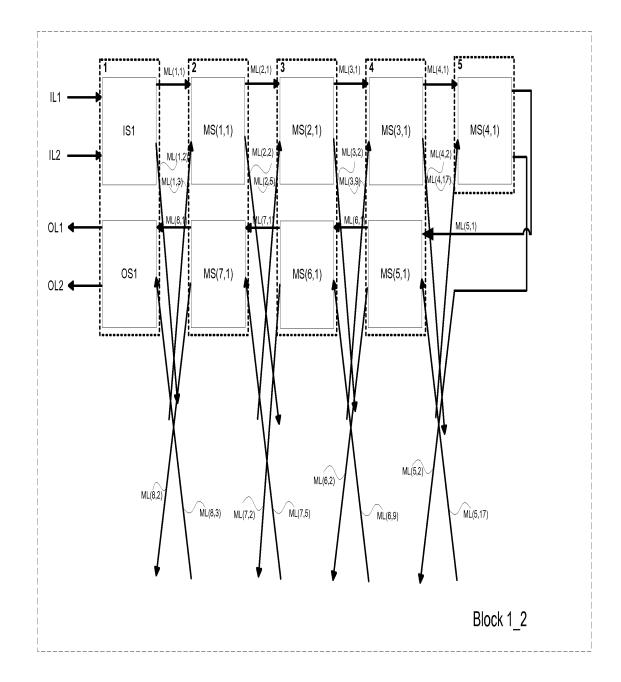






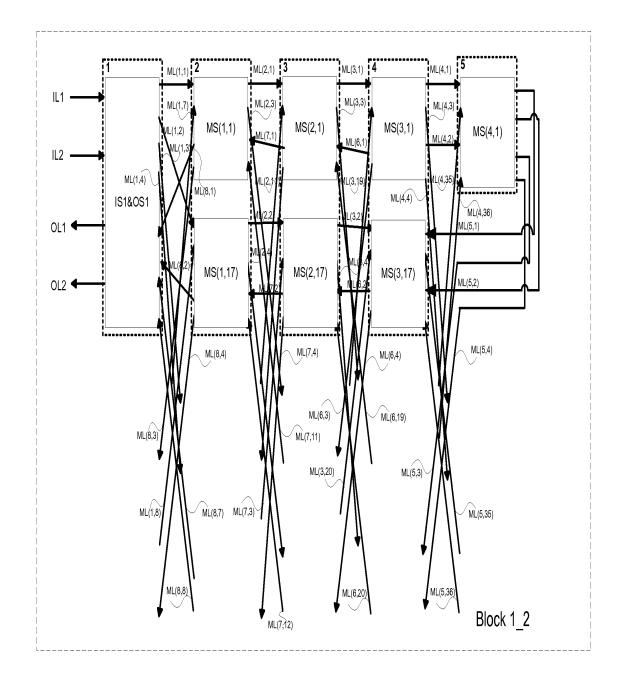
















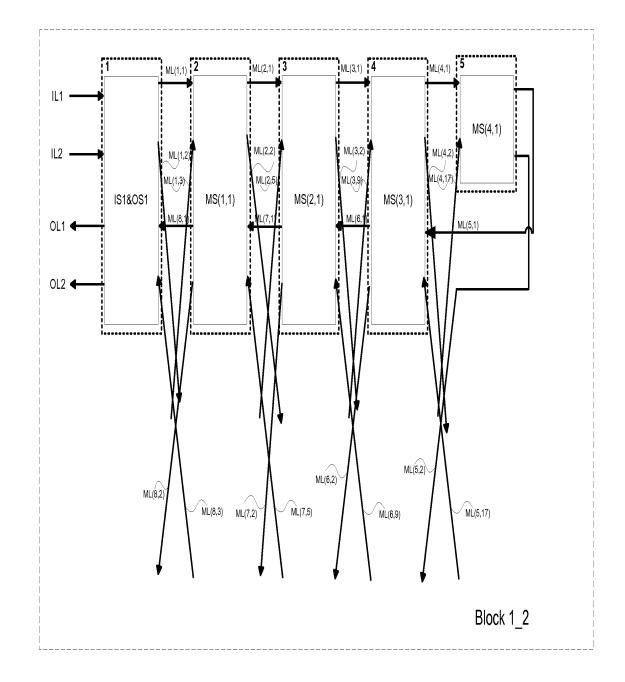
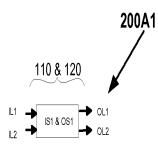
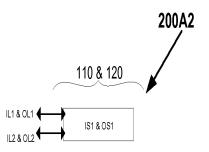
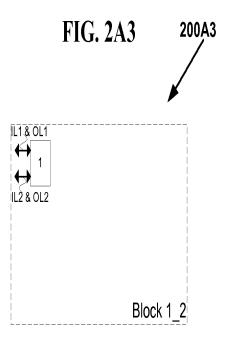


FIG. 2A1

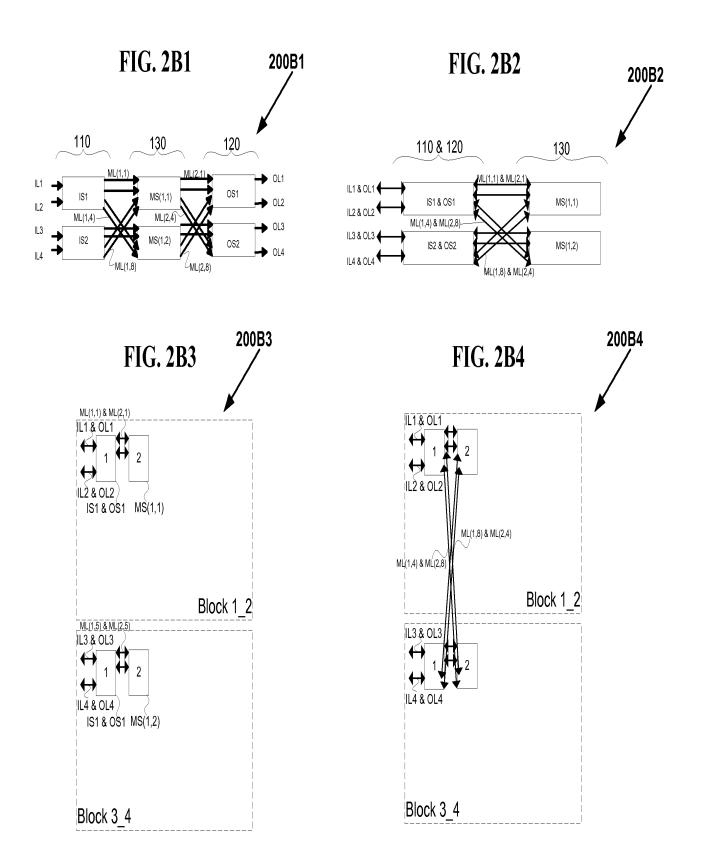
FIG. 2A2







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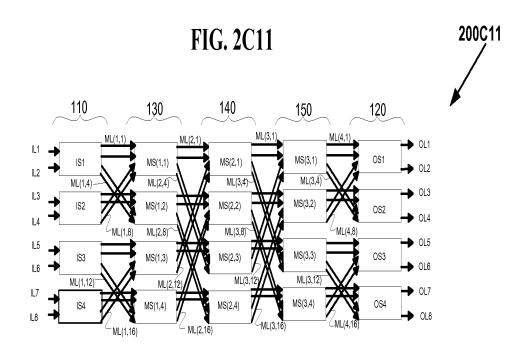
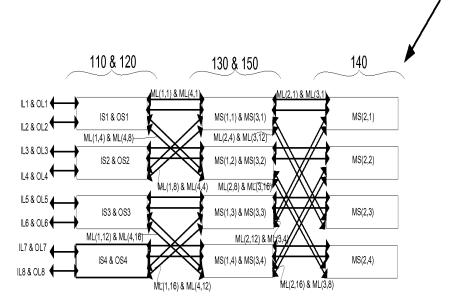
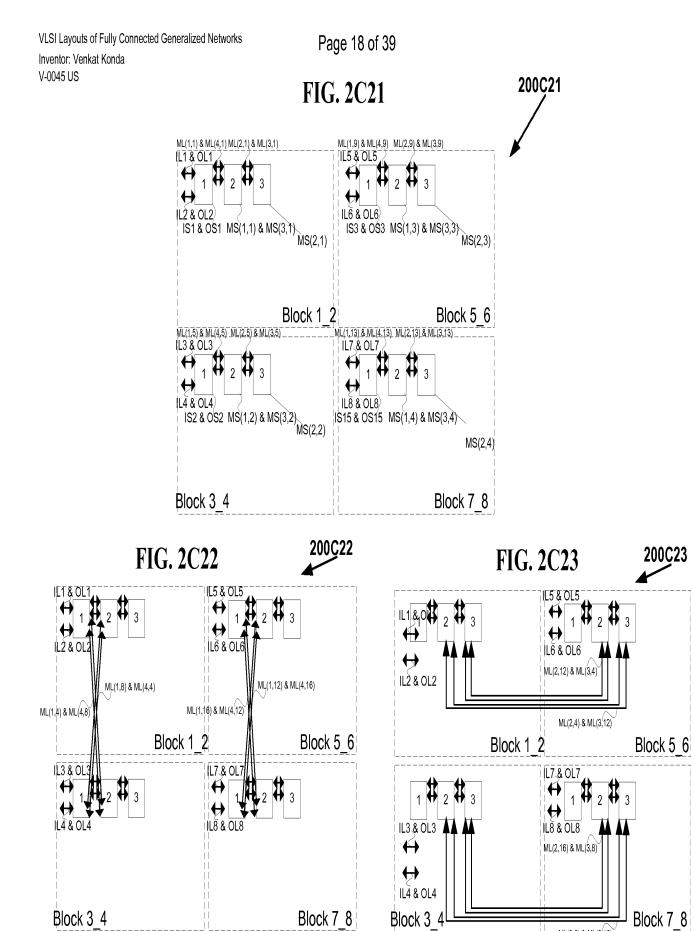


FIG. 2C12

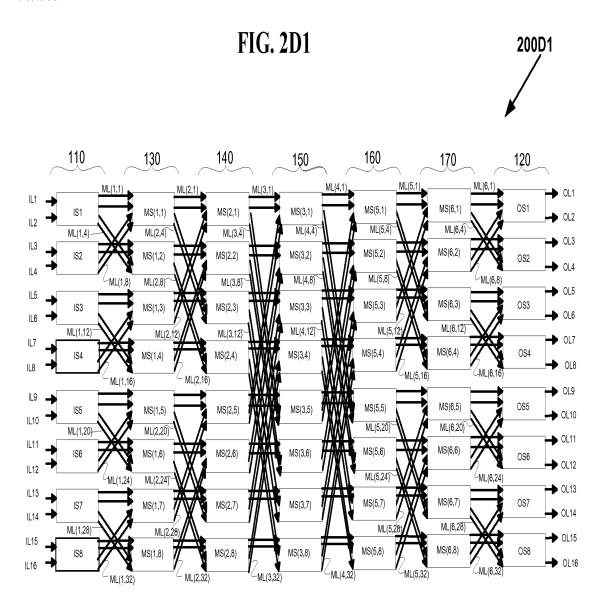
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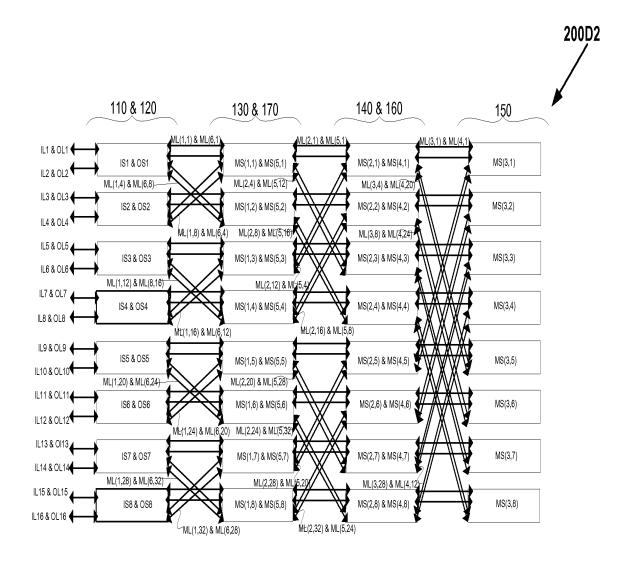
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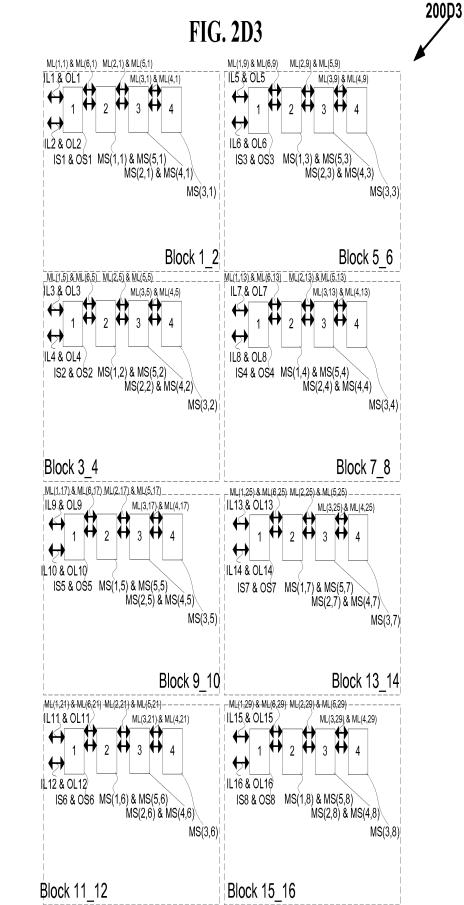


ML(2,8) & ML(3,16)

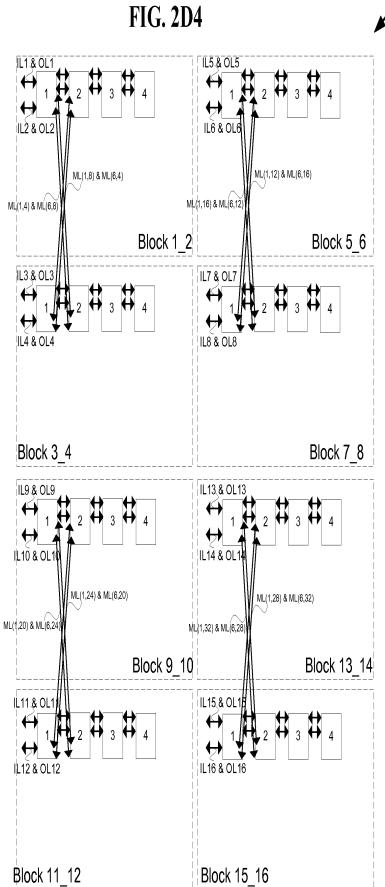


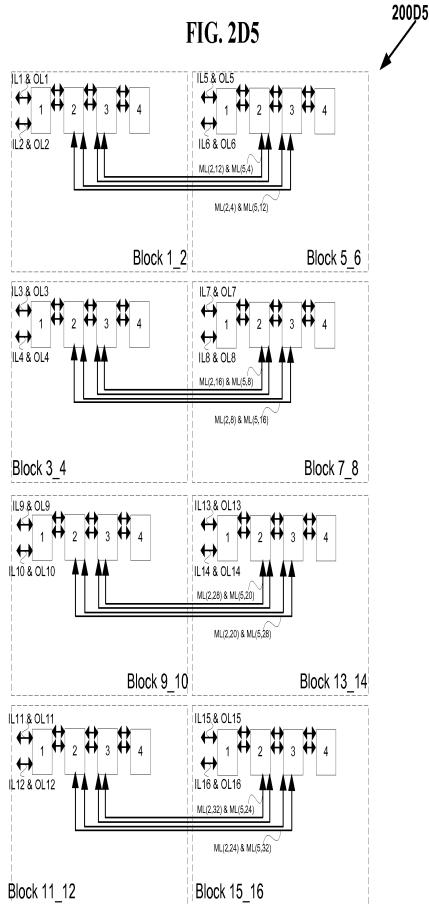




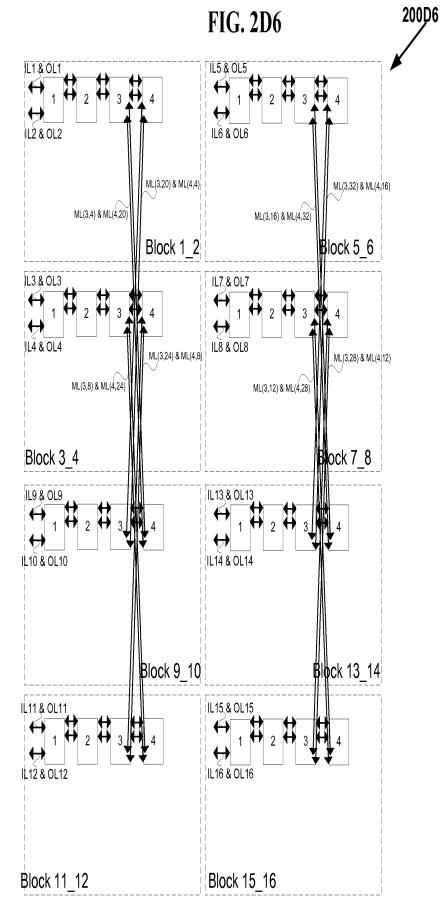


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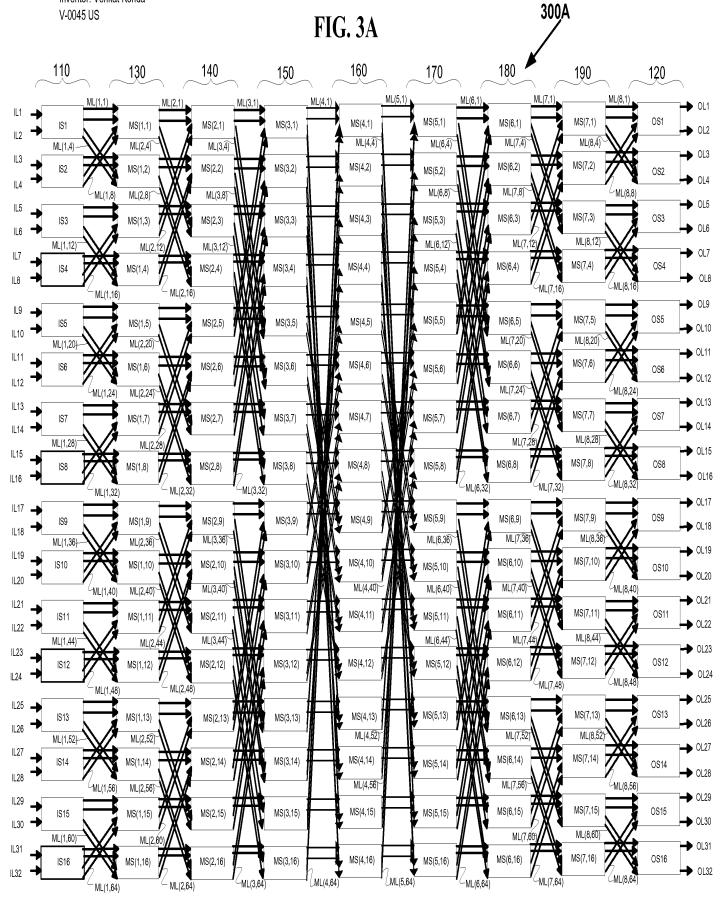
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110 & 120

IS1 & OS1

IS2 & OS2

IS3 & OS3

ML(1,4) & ML(8,8)-

IL1 & OL1

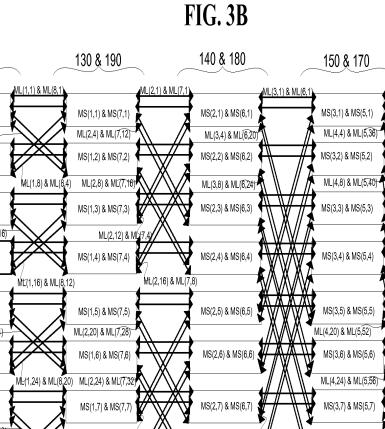
IL2 & OL2

IL3 & OL3

IL4 & OL4

IL5 & OL5

IL6 & OL6



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300B

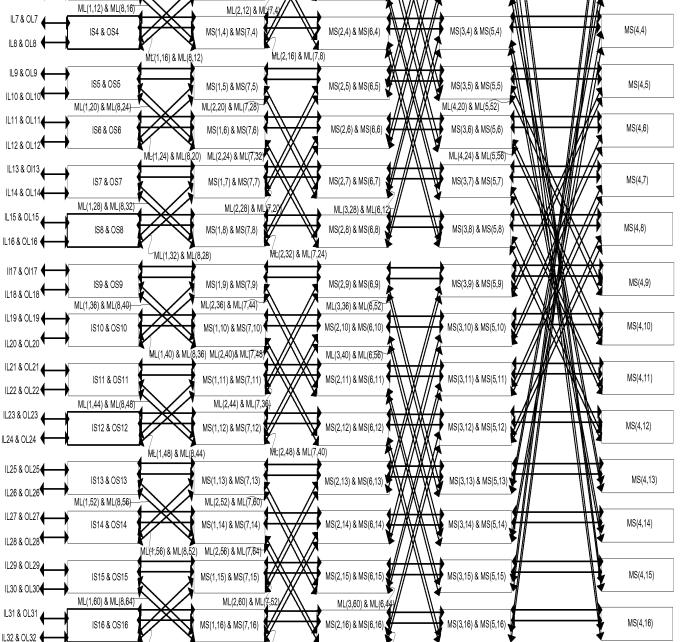
AL(4,1) & ML(5,1)

160

MS(4,1)

MS(4,2)

MS(4,3)



(ML(2,64) & ML(7,56)

(ML(3,64) & ML(6,48)

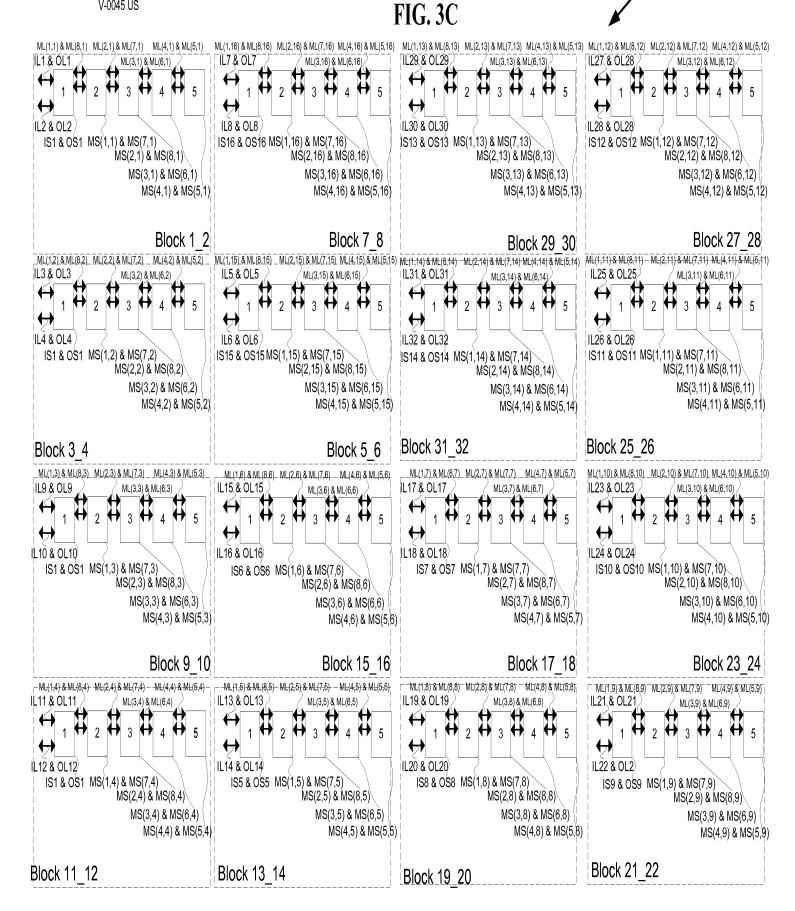
ML(1,64) & ML(8,60)

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300C

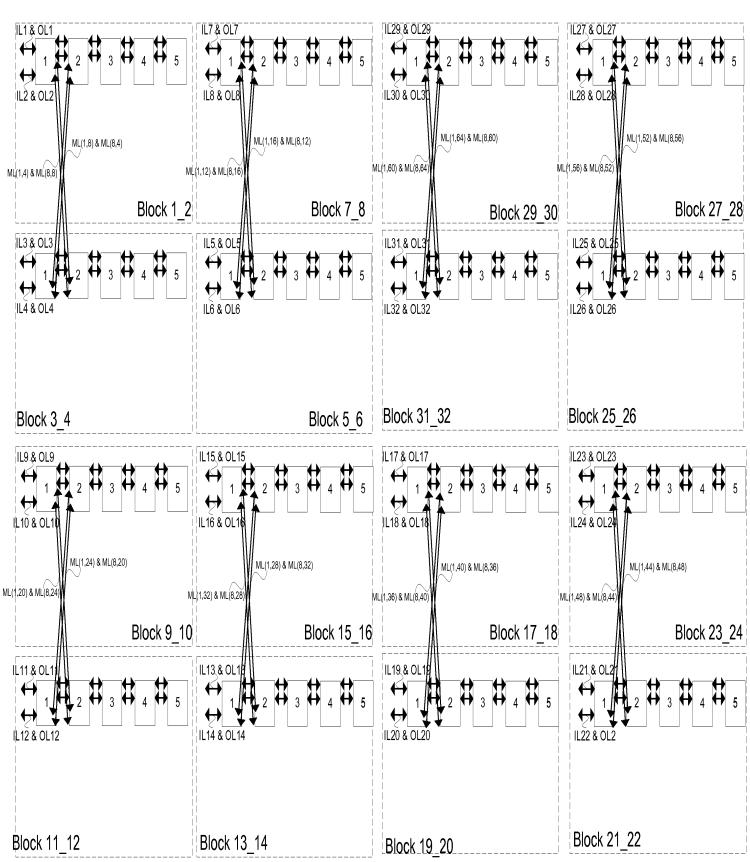
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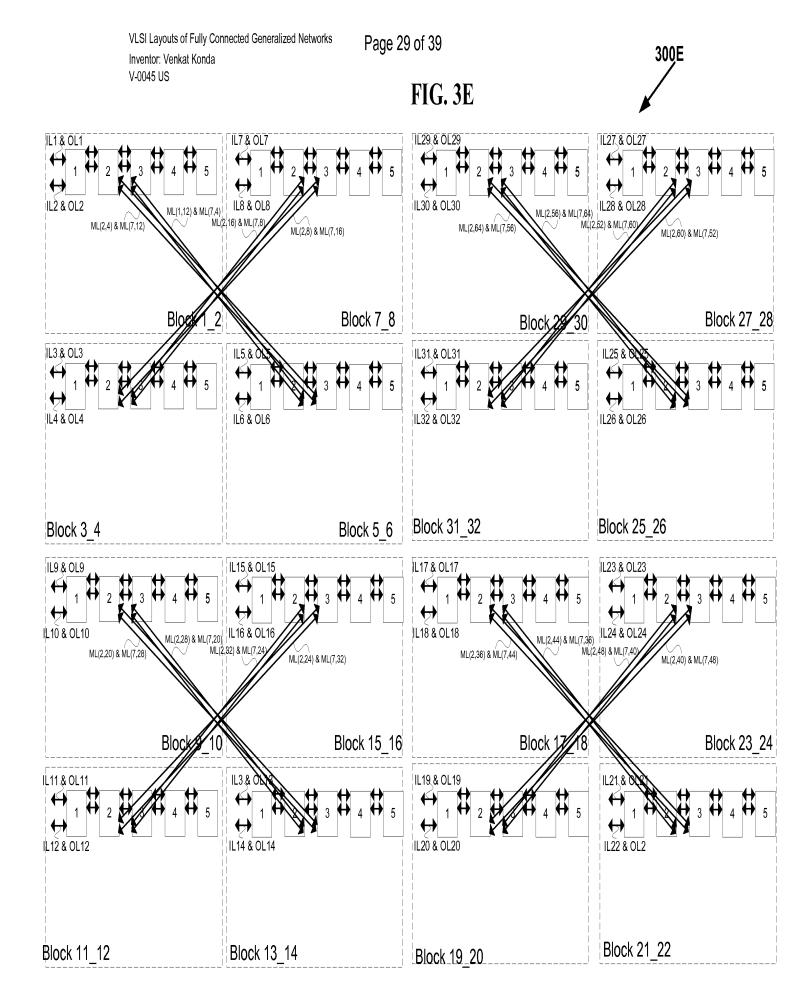
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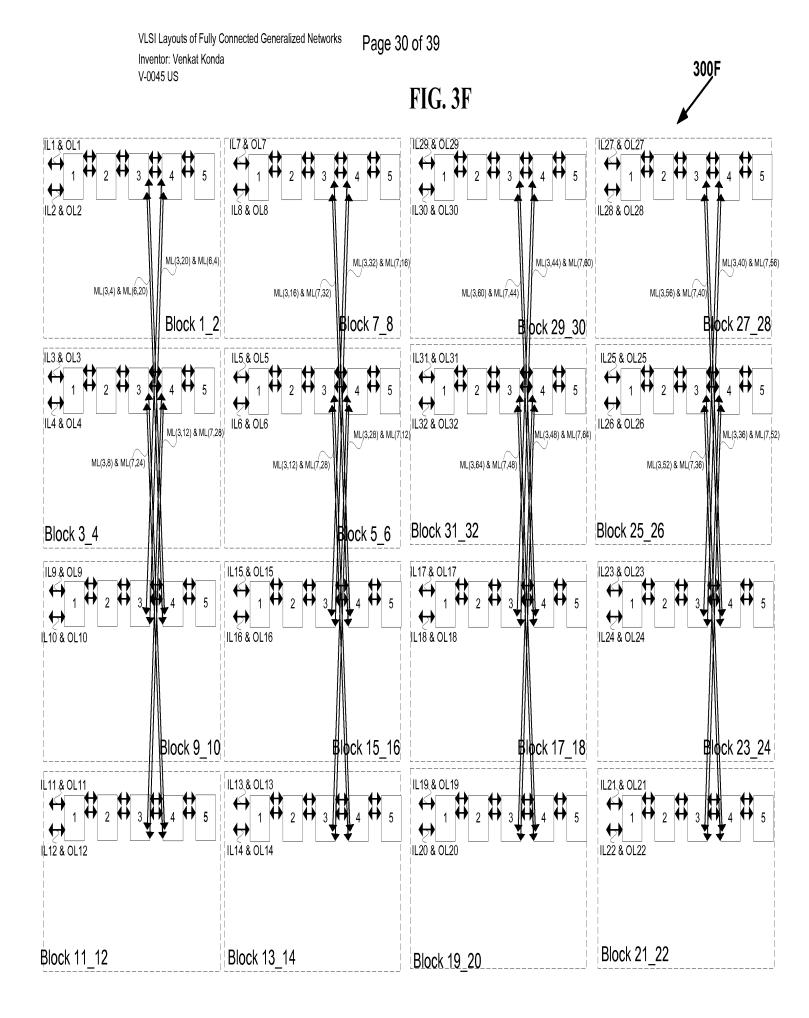
300D



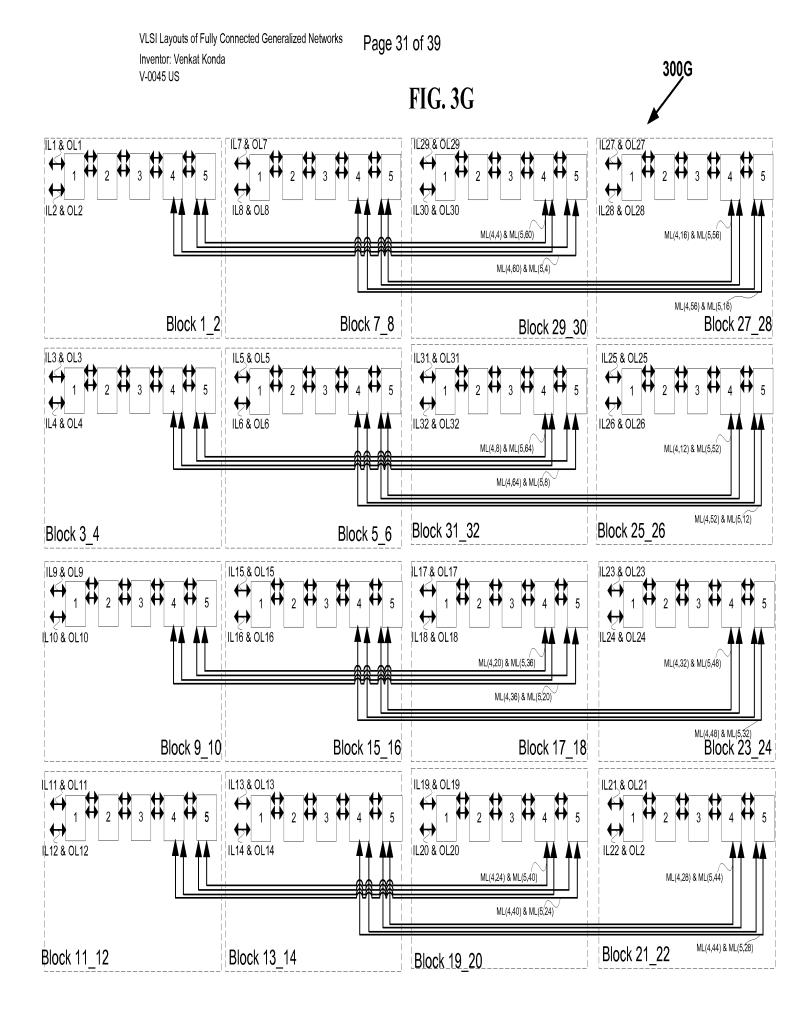
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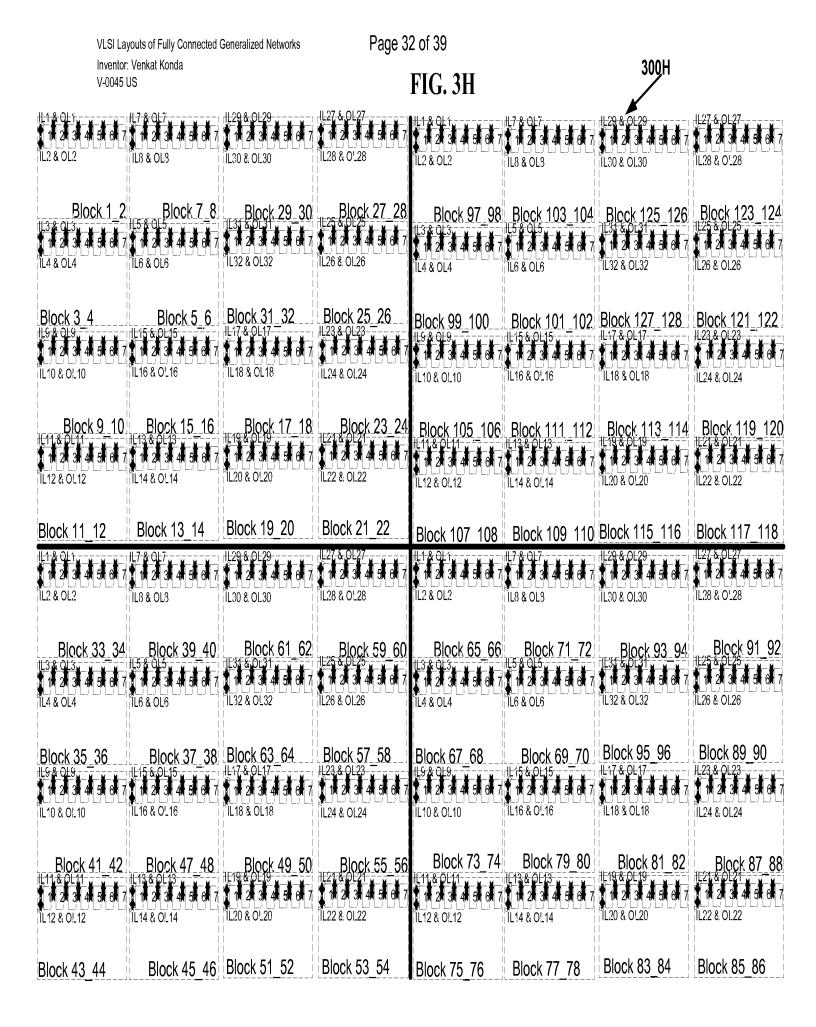


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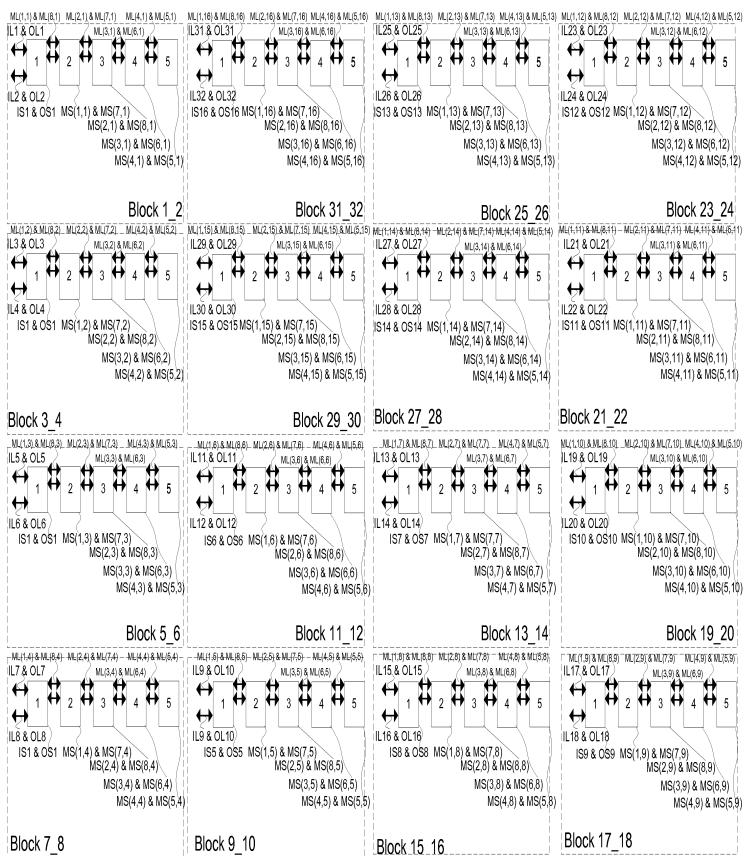
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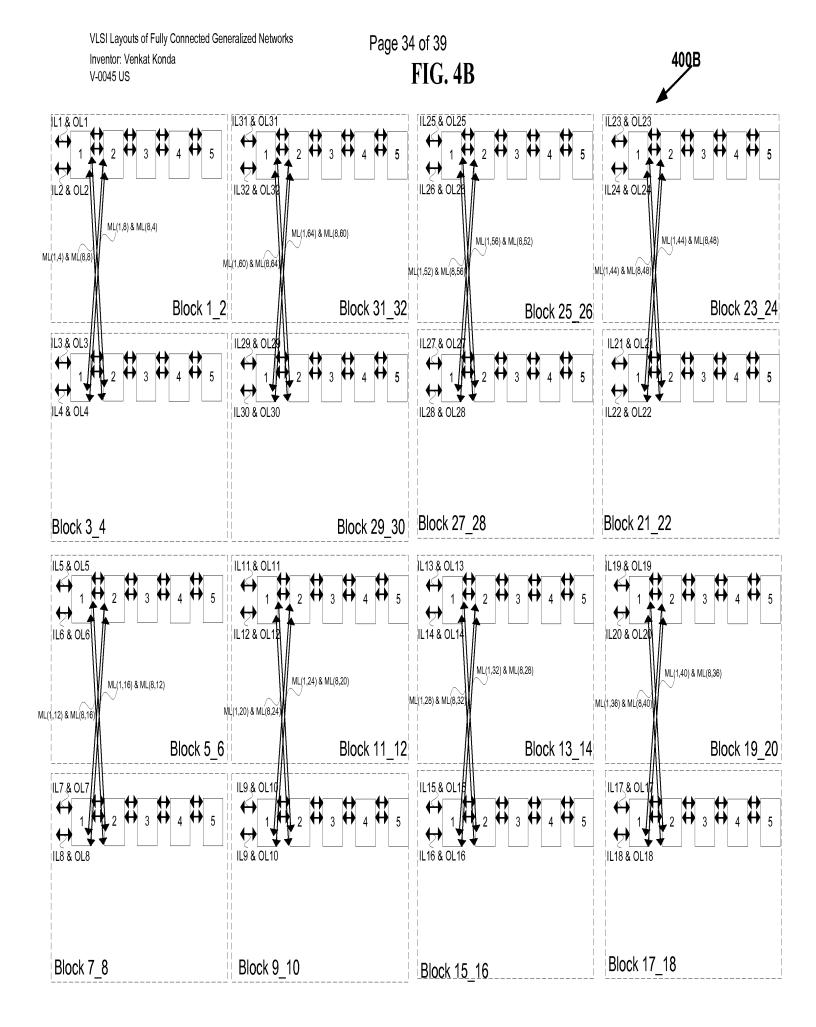




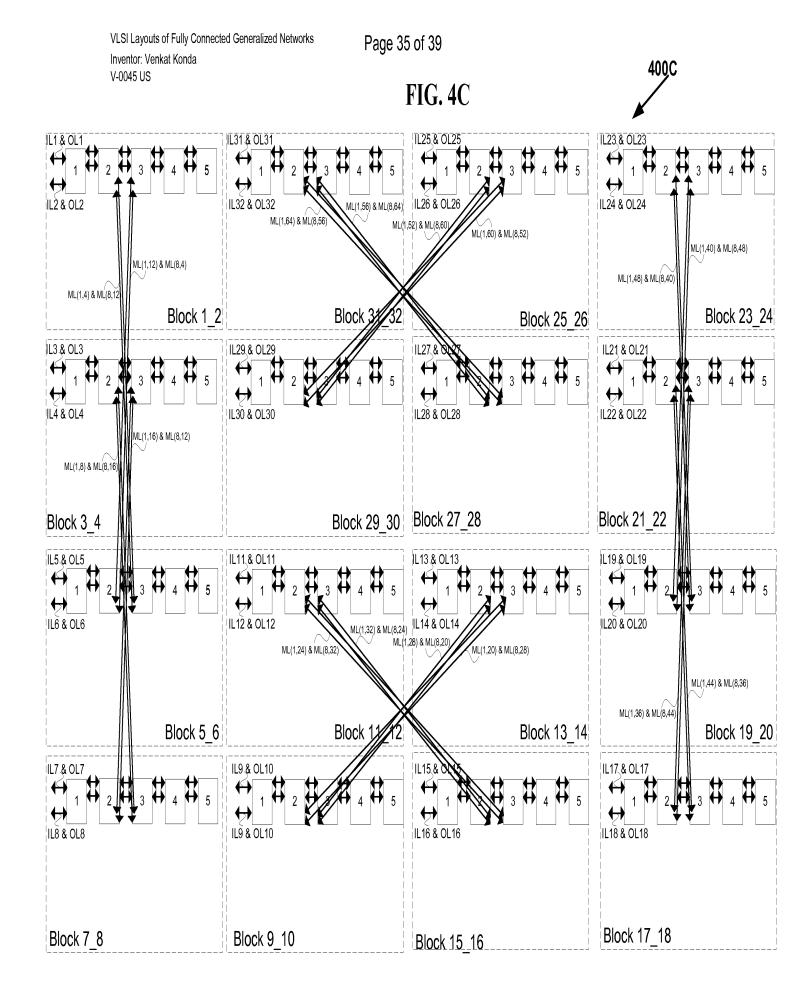


400A

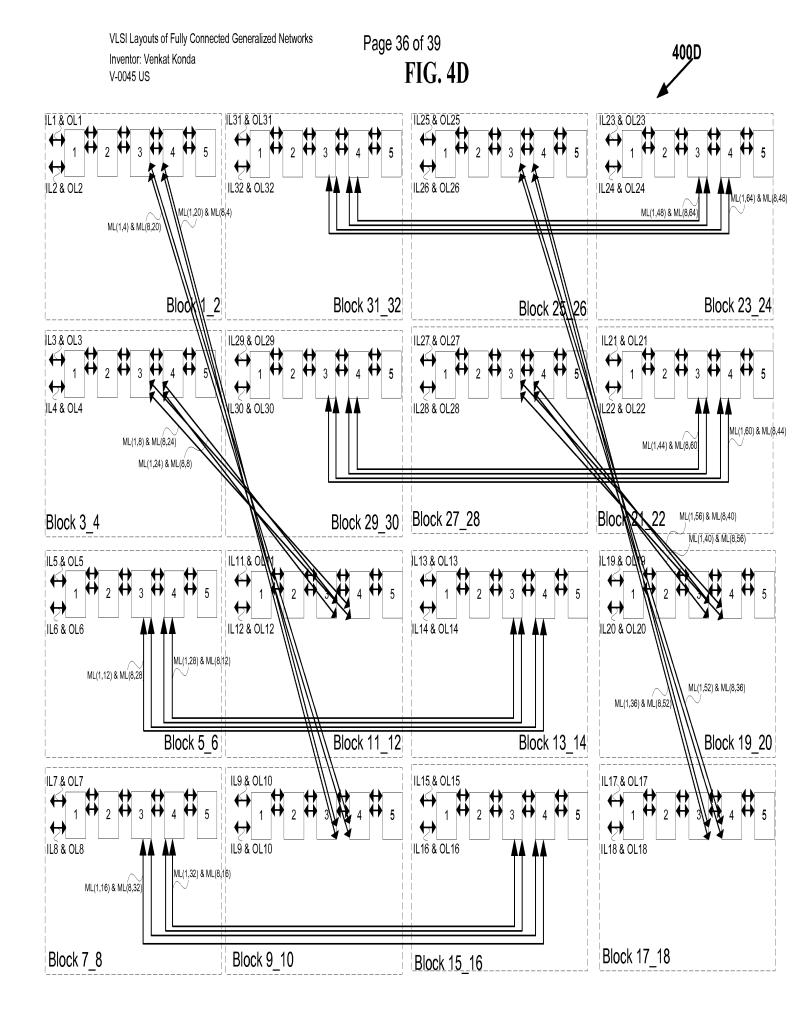


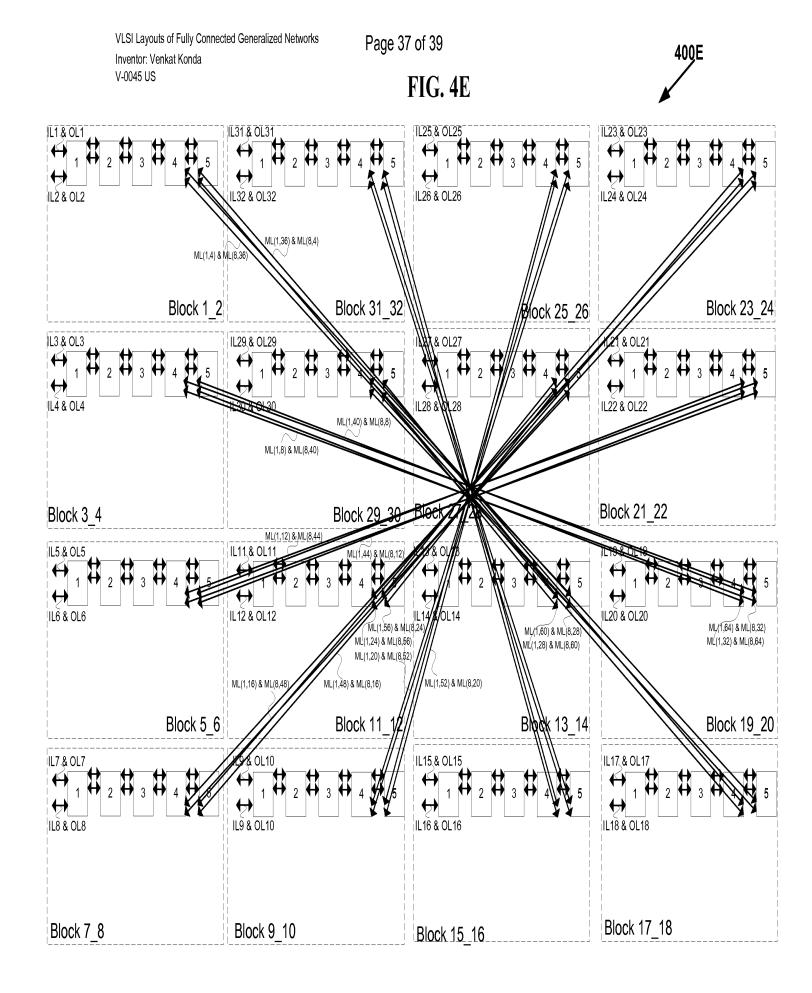


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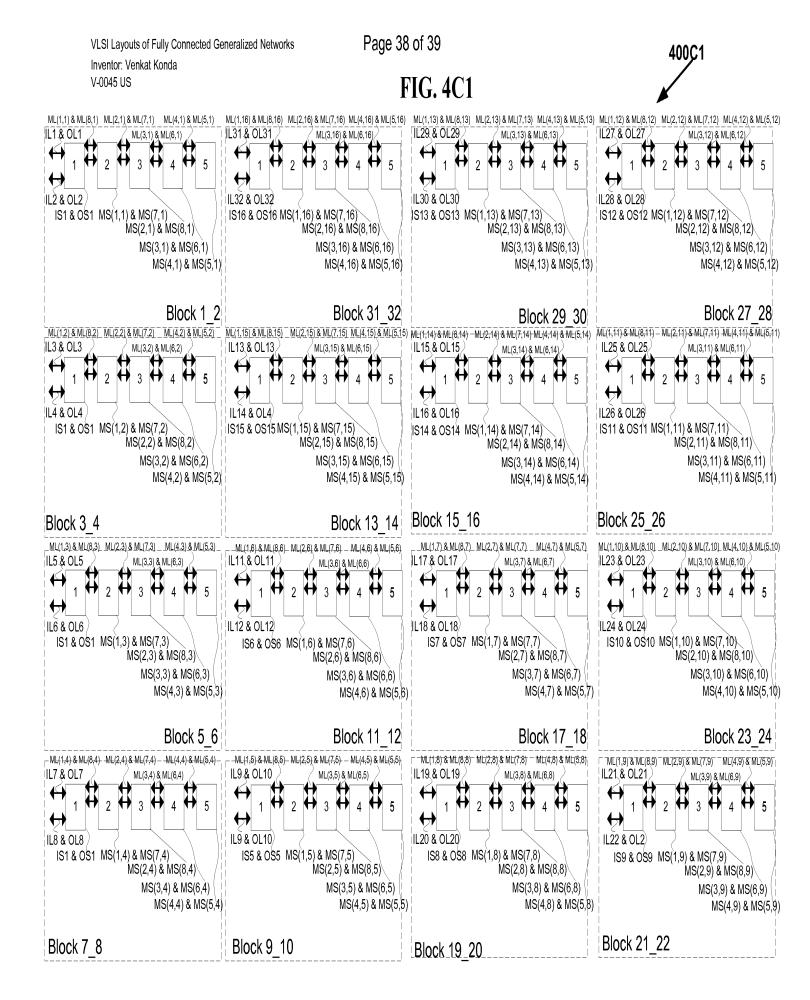


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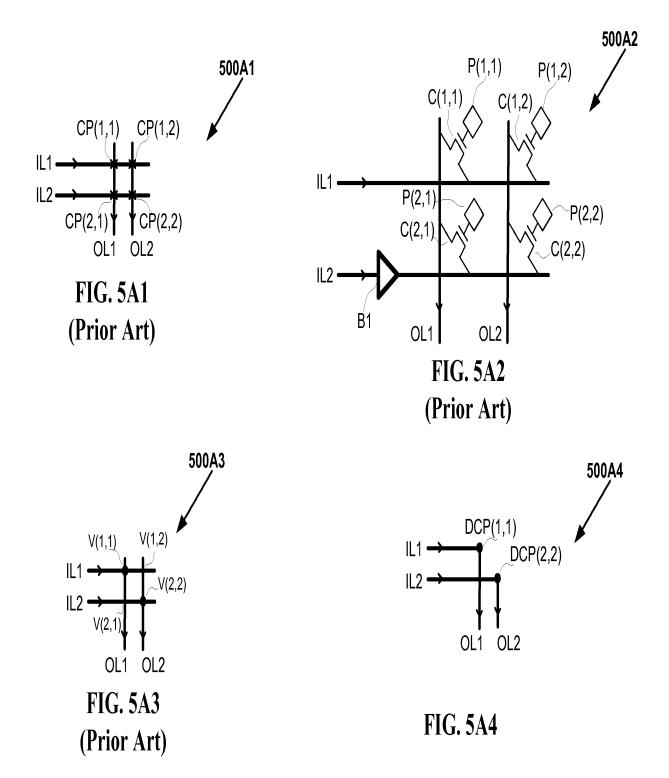


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VLSI Layouts of Fully Connected Generalized Networks Inventor: Venkat Konda V-0045 US

FIG. 5A



(19) World Intellectual Property Organization International Bureau

> (43) International Publication Date 4 December 2008 (04.12.2008)



PCT

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(71) Applicant and

- (72) Inventor: KONDA, Venkat [US/US]; 6278, Grand Oak Way, San Jose, CA 95135 (US).
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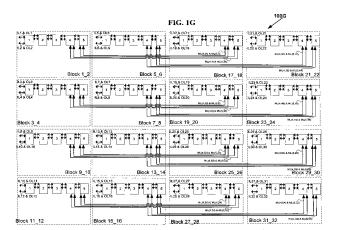
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(54) Title: VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS



(57) Abstract: In accordance with the invention, VLSI layouts of generalized multi-stage networks for broadcast, unicast and multicast connections are presented using only horizontal and vertical links. The VLSI layouts employ shuffle exchange links where outlet links of cross links from switches in a stage in one sub-integrated circuit block are connected to inlet links of switches in the succeeding stage in another sub- integrated circuit block so that said cross links are either vertical links or horizontal and vice versa. In one embodiment the sub- integrated circuit blocks are arranged in a hypercube arrangement in a two-dimensional plane. The VLSI layouts exploit the benefits of significantly lower cross points, lower signal latency, lower power and full connectivity with significantly fast compilation. 10

VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS Venkat Konda

5 CROSS REFERENCE TO RELATED APPLICATIONS

This application is Continuation In Part PCT Application to and incorporates by reference in its entirety the U.S. Provisional Patent Application Serial No. 60/940, 394 entitled "VLSI LAYOUTS OF FULLY CONNECTED GENERALIZED NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007.

This application is related to and incorporates by reference in its entirety the PCT Application Serial No. PCT/US08/56064 entitled "FULLY CONNECTED GENERALIZED MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed March 6, 2008, the U.S. Provisional Patent

- Application Serial No. 60/905,526 entitled "LARGE SCALE CROSSPOINT REDUCTION WITH NONBLOCKING UNICAST & MULTICAST IN ARBITRARILY LARGE MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed March 6, 2007, and the U.S. Provisional Patent Application Serial No. 60/940, 383 entitled "FULLY CONNECTED
- 20 GENERALIZED MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007.

This application is related to and incorporates by reference in its entirety the PCT Application Docket No. S-0038PCT entitled "FULLY CONNECTED GENERALIZED BUTTERFLY FAT TREE NETWORKS" by Venkat Konda assigned to the same

assignee as the current application, filed concurrently, the U.S. Provisional Patent
 Application Serial No. 60/940, 387 entitled "FULLY CONNECTED GENERALIZED
 BUTTERFLY FAT TREE NETWORKS" by Venkat Konda assigned to the same

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assignee as the current application, filed May 25, 2007, and the U.S. Provisional Patent Application Serial No. 60/940, 390 entitled "FULLY CONNECTED GENERALIZED MULTI-LINK BUTTERFLY FAT TREE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007

5 This application is related to and incorporates by reference in its entirety the PCT Application Docket No. S-0039PCT entitled "FULLY CONNECTED GENERALIZED MULTI-LINK MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed concurrently, the U.S. Provisional Patent Application Serial No. 60/940, 389 entitled "FULLY CONNECTED GENERALIZED

- 10 REARRANGEABLY NONBLOCKING MULTI-LINK MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007, the U.S. Provisional Patent Application Serial No. 60/940, 391 entitled "FULLY CONNECTED GENERALIZED FOLDED MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007 and
- 15 the U.S. Provisional Patent Application Serial No. 60/940, 392 entitled "FULLY CONNECTED GENERALIZED STRICTLY NONBLOCKING MULTI-LINK MULTI-STAGE NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed May 25, 2007.

This application is related to and incorporates by reference in its entirety the U.S. 20 Provisional Patent Application Serial No. 60/984, 724 entitled "VLSI LAYOUTS OF FULLY CONNECTED NETWORKS WITH LOCALITY EXPLOITATION" by Venkat Konda assigned to the same assignee as the current application, filed November 2, 2007.

This application is related to and incorporates by reference in its entirety the U.S. Provisional Patent Application Serial No. 61/018, 494 entitled "VLSI LAYOUTS OF

25 FULLY CONNECTED GENERALIZED AND PYRAMID NETWORKS" by Venkat Konda assigned to the same assignee as the current application, filed January 1, 2008. WO 2008/147928

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PCT/US2008/064605

BACKGROUND OF INVENTION

Multi-stage interconnection networks such as Benes networks and butterfly fat tree networks are widely useful in telecommunications, parallel and distributed computing. However VLSI layouts, known in the prior art, of these interconnection networks in an integrated circuit are inefficient and complicated.

Other multi-stage interconnection networks including butterfly fat tree networks, Banyan networks, Batcher-Banyan networks, Baseline networks, Delta networks, Omega networks and Flip networks have been widely studied particularly for self routing packet switching applications. Also Benes Networks with radix of two have been widely studied and it is known that Benes Networks of radix two are shown to be built with back to back

10 and it is known that Benes Networks of radix two are shown to be built with back to bac baseline networks which are rearrangeably nonblocking for unicast connections.

The most commonly used VLSI layout in an integrated circuit is based on a twodimensional grid model comprising only horizontal and vertical tracks. An intuitive interconnection network that utilizes two-dimensional grid model is 2D Mesh Network 15 and its variations such as segmented mesh networks. Hence routing networks used in VLSI layouts are typically 2D mesh networks and its variations. However Mesh Networks require large scale cross points typically with a growth rate of $O(N^2)$ where N is the number of computing elements, ports, or logic elements depending on the application.

- 20 Multi-stage interconnection with a growth rate of $O(N \times \log N)$ requires significantly small number of cross points. U.S. Patent 6,185,220 entitled "Grid Layouts of Switching and Sorting Networks" granted to Muthukrishnan et al. describes a VLSI layout using existing VLSI grid model for Benes and Butterfly networks. U.S. Patent 6,940,308 entitled "Interconnection Network for a Field Programmable Gate Array"
- 25 granted to Wong describes a VLSI layout where switches belonging to lower stage of Benes Network are layed out close to the logic cells and switches belonging to higher stages are layed out towards the center of the layout.

Due to the inefficient and in some cases impractical VLSI layout of Benes and butterfly fat tree networks on a semiconductor chip, today mesh networks and segmented mesh networks are widely used in the practical applications such as field programmable gate arrays (FPGAs), programmable logic devices (PLDs), and parallel computing

- 5 interconnects. The prior art VLSI layouts of Benes and butterfly fat tree networks and VLSI layouts of mesh networks and segmented mesh networks require large area to implement the switches on the chip, large number of wires, longer wires, with increased power consumption, increased latency of the signals which effect the maximum clock speed of operation. Some networks may not even be implemented practically on a chip
- 10 due to the lack of efficient layouts.

SUMMARY OF INVENTION

When large scale sub-integrated circuit blocks with inlet and outlet links are layed out in an integrated circuit device in a two-dimensional grid arrangement, (for example in
an FPGA where the sub-integrated circuit blocks are Lookup Tables) the most intuitive routing network is a network that uses horizontal and vertical links only (the most often used such a network is one of the variations of a 2D Mesh network). A direct embedding of a generalized multi-stage network on to a 2D Mesh network is neither simple nor efficient.

- In accordance with the invention, VLSI layouts of generalized multi-stage networks for broadcast, unicast and multicast connections are presented using only horizontal and vertical links. The VLSI layouts employ shuffle exchange links where outlet links of cross links from switches in a stage in one sub-integrated circuit block are connected to inlet links of switches in the succeeding stage in another sub-integrated circuit block so that said cross links are either vertical links or horizontal and vice versa.
- In one embodiment the sub-integrated circuit blocks are arranged in a hypercube arrangement in a two-dimensional plane. The VLSI layouts exploit the benefits of significantly lower cross points, lower signal latency, lower power and full connectivity with significantly fast compilation.

The VLSI layouts presented are applicable to generalized multi-stage networks $V(N_1, N_2, d, s)$, generalized folded multi-stage networks $V_{fold}(N_1, N_2, d, s)$, generalized butterfly fat tree networks $V_{bft}(N_1, N_2, d, s)$, generalized multi-link multi-stage networks $V_{mlink}(N_1, N_2, d, s)$, generalized folded multi-link multi-stage networks

5 $V_{fold-mlink}(N_1, N_2, d, s)$, generalized multi-link butterfly fat tree networks $V_{mlink-bft}(N_1, N_2, d, s)$, and generalized hypercube networks $V_{hcube}(N_1, N_2, d, s)$ for s = 1,2,3 or any number in general. The embodiments of VLSI layouts are useful in wide target applications such as FPGAs, CPLDs, pSoCs, ASIC placement and route tools, networking applications, parallel & distributed computing, and reconfigurable computing.

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BRIEF DESCRIPTION OF DRAWINGS

FIG. 1A is a diagram 100A of an exemplary symmetrical multi-link multi-stage network $V_{fold-mlink}(N,d,s)$ having inverse Benes connection topology of nine stages with N = 32, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention.

FIG. 1B is a diagram 100B of the equivalent symmetrical folded multi-link multi-stage network V_{fold-mlink}(N,d,s) of the network 100A shown in FIG. 1A, having inverse Benes connection topology of five stages with N = 32, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fanout multicast connections, in accordance with the invention.

FIG. 1C is a diagram 100C layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 1B, in one embodiment, illustrating the connection links belonging with in each block only.

FIG. 1D is a diagram 100D layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 1B, in one embodiment, illustrating the connection links ML(1,i) for i = [1, 64] and ML(8,i) for i = [1,64].

FIG. 1E is a diagram 100E layout of the network V_{fold-mlink} (N,d,s) shown in FIG.
5 1B, in one embodiment, illustrating the connection links ML(2,i) for i = [1, 64] and ML(7,i) for i = [1,64].

FIG. 1F is a diagram 100F layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 1B, in one embodiment, illustrating the connection links ML(3,i) for i = [1, 64] and ML(6,i) for i = [1,64].

10 FIG. 1G is a diagram 100G layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 1B, in one embodiment, illustrating the connection links ML(4,i) for i = [1, 64] and ML(5,i) for i = [1,64].

FIG. 1H is a diagram 100H layout of a network $V_{fold-mlink}(N,d,s)$ where N = 128, d = 2, and s = 2, in one embodiment, illustrating the connection links belonging with in 15 each block only.

FIG. 1I is a diagram 100I detailed connections of BLOCK 1_2 in the network layout 100C in one embodiment, illustrating the connection links going in and coming out when the layout 100C is implementing V(N,d,s) or $V_{fold}(N,d,s)$.

FIG. 1J is a diagram 100J detailed connections of BLOCK 1_2 in the network 20 layout 100C in one embodiment, illustrating the connection links going in and coming out when the layout 100C is implementing V(N,d,s) or $V_{fold}(N,d,s)$.

FIG. 1K is a diagram 100K detailed connections of BLOCK 1_2 in the network layout 100C in one embodiment, illustrating the connection links going in and coming out when the layout 100C is implementing V(N,d,s) or $V_{fold}(N,d,s)$.

FIG. 1K1 is a diagram 100M1 detailed connections of BLOCK 1_2 in the network layout 100C in one embodiment, illustrating the connection links going in and coming out when the layout 100C is implementing V(N, d, s) or $V_{fold}(N, d, s)$ for s = 1.

FIG. 1L is a diagram 100L detailed connections of BLOCK 1_2 in the network 5 layout 100C in one embodiment, illustrating the connection links going in and coming out when the layout 100C is implementing V(N,d,s) or $V_{fold}(N,d,s)$.

FIG. 1L1 is a diagram 100L1 detailed connections of BLOCK 1_2 in the network layout 100C in one embodiment, illustrating the connection links going in and coming out when the layout 100C is implementing V(N,d,s) or $V_{fold}(N,d,s)$ for s = 1.

FIG. 2A1 is a diagram 200A1 of an exemplary symmetrical multi-link multi-stage network V_{fold-mlink} (N, d, s) having inverse Benes connection topology of one stage with N = 2, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention. FIG. 2A2 is a diagram 200A2 of the equivalent symmetrical folded multi-link multi-stage network V_{fold-mlink} (N, d, s) of the network 200A1 shown in FIG. 2A1, having inverse Benes connection topology of one stage with N = 2, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention. FIG. 2A3 is a diagram 200A3 layout of the network V_{fold-mlink} (N, d, s) shown in FIG. 2A2, in one embodiment, illustrating all the connection

links.

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FIG. 2B1 is a diagram 200B1 of an exemplary symmetrical multi-link multi-stage network $V_{fold-mlink}(N,d,s)$ having inverse Benes connection topology of one stage with N = 4, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention. FIG. 2B2 is a diagram 200B2 of the equivalent symmetrical folded multi-link multi-stage network $V_{fold-mlink}(N,d,s)$ of the network 5

200B1 shown in FIG. 2B1, having inverse Benes connection topology of one stage with N = 4, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention. FIG. 2B3 is a diagram 200B3 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2B2, in one embodiment, illustrating the connection links belonging with in each block only. FIG. 2B4 is a diagram 200B4 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2B2, in one embodiment, illustrating the connection links belonging with in each block only. FIG. 2B4 is a diagram 200B4 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2B2, in one embodiment, illustrating the connection links belonging with in each block only. FIG. 2B4 is a diagram 200B4 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2B2, in one embodiment, illustrating the connection links belonging with in each block only. FIG. 2B4 is a diagram 200B4 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2B2, in one embodiment, illustrating the connection links belonging with in each block only. FIG. 2B4 is a diagram 200B4 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2B2, in one embodiment, illustrating the connection links belonging with in each block only.

FIG. 2C11 is a diagram 200C11 of an exemplary symmetrical multi-link multistage network V_{fold-mlink} (N, d, s) having inverse Benes connection topology of one stage with N = 8, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention. FIG. 2C12 is a diagram 200C12 of the equivalent symmetrical folded multi-link multi-stage network V_{fold-mlink} (N, d, s) of the network
200C11 shown in FIG. 2C11, having inverse Benes connection topology of one stage with N = 8, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention.

connection links ML(1,i) for i = [1, 8] and ML(2,i) for i = [1,8].

FIG. 2C21 is a diagram 200C21 layout of the network $V_{fold-mlink}(N,d,s)$ shown in 20 FIG. 2C12, in one embodiment, illustrating the connection links belonging with in each block only. FIG. 2C22 is a diagram 200C22 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2C12, in one embodiment, illustrating the connection links ML(1,i) for i = [1, 16] and ML(4,i) for i = [1,16]. FIG. 2C23 is a diagram 200C23 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2C12, in one embodiment, illustrating the connection 25 links ML(2,i) for i = [1, 16] and ML(3,i) for i = [1,16].

FIG. 2D1 is a diagram 200D1 of an exemplary symmetrical multi-link multi-stage network $V_{fold-mlink}(N,d,s)$ having inverse Benes connection topology of one stage with N = 16, d = 2 and s=2, strictly nonblocking network for unicast connections and 5

rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention.

FIG. 2D2 is a diagram 200D2 of the equivalent symmetrical folded multi-link multi-stage network $V_{fold-mlink}(N,d,s)$ of the network 200D1 shown in FIG. 2D1, having inverse Benes connection topology of one stage with N = 16, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention.

FIG. 2D3 is a diagram 200D3 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2D2, in one embodiment, illustrating the connection links belonging with in each block only.

FIG. 2D4 is a diagram 200D4 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2D2, in one embodiment, illustrating the connection links ML(1,i) for i = [1, 32] and ML(6,i) for i = [1,32].

FIG. 2D5 is a diagram 200D5 layout of the network V_{fold-mlink}(N,d,s) shown in
FIG. 2D2, in one embodiment, illustrating the connection links ML(2,i) for i = [1, 32] and ML(5,i) for i = [1,32].

FIG. 2D6 is a diagram 200D6 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 2D2, in one embodiment, illustrating the connection links ML(3,i) for i = [1, 32] and ML(4,i) for i = [1,32].

- FIG. 3A is a diagram 300A of an exemplary symmetrical multi-link multi-stage network $V_{hcube}(N,d,s)$ having inverse Benes connection topology of nine stages with N = 32, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fan-out multicast connections, in accordance with the invention.
- FIG. 3B is a diagram 300B of the equivalent symmetrical folded multi-link multistage network $V_{hcube}(N,d,s)$ of the network 300A shown in FIG. 3A, having inverse

Benes connection topology of five stages with N = 32, d = 2 and s=2, strictly nonblocking network for unicast connections and rearrangeably nonblocking network for arbitrary fanout multicast connections, in accordance with the invention.

FIG. 3C is a diagram 300C layout of the network V_{hcube}(N,d,s) shown in FIG.
3B, in one embodiment, illustrating the connection links belonging with in each block only.

FIG. 3D is a diagram 100D layout of the network $V_{hcube}(N,d,s)$ shown in FIG. 3B, in one embodiment, illustrating the connection links ML(1,i) for i = [1, 64] and ML(8,i) for i = [1,64].

FIG. 3E is a diagram 300E layout of the network V_{hcube} (N, d, s) shown in FIG.
3B, in one embodiment, illustrating the connection links ML(2,i) for i = [1, 64] and ML(7,i) for i = [1,64].

FIG. 3F is a diagram 300F layout of the network $V_{hcube}(N, d, s)$ shown in FIG. 3B, in one embodiment, illustrating the connection links ML(3,i) for i = [1, 64] and ML(6,i) 15 for i = [1,64].

FIG. 3G is a diagram 300G layout of the network $V_{hcube}(N,d,s)$ shown in FIG. 3B, in one embodiment, illustrating the connection links ML(4,i) for i = [1, 64] and ML(5,i) for i = [1,64].

FIG. 3H is a diagram 300H layout of a network V_{hcube}(N, d, s) where N = 128, d =
20 2, and s = 2, in one embodiment, illustrating the connection links belonging with in each block only.

FIG. 4A is a diagram 400A layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 1B, in one embodiment, illustrating the connection links belonging with in each block only.

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FIG. 4B is a diagram 400B layout of the network $V_{fold-mlink}(N, d, s)$ shown in FIG. 1B, in one embodiment, illustrating the connection links ML(1,i) for i = [1, 64] and ML(8,i) for i = [1,64].

FIG. 4C is a diagram 400C layout of the network V_{fold-mlink} (N, d, s) shown in FIG.
4C, in one embodiment, illustrating the connection links ML(2,i) for i = [1, 64] and ML(7,i) for i = [1,64].

FIG. 4D is a diagram 400D layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 4D, in one embodiment, illustrating the connection links ML(3,i) for i = [1, 64] and ML(6,i) for i = [1,64].

FIG. 4E is a diagram 400E layout of the network V_{fold-mlink} (N,d,s) shown in FIG.
4E, in one embodiment, illustrating the connection links ML(4,i) for i = [1, 64] and ML(5,i) for i = [1,64].

FIG. 4C1 is a diagram 400C1 layout of the network $V_{fold-mlink}(N,d,s)$ shown in FIG. 1B, in one embodiment, illustrating the connection links belonging with in each block only.

FIG. 5A1 is a diagram 500A1 of an exemplary prior art implementation of a two by two switch; FIG. 5A2 is a diagram 500A2 for programmable integrated circuit prior art implementation of the diagram 500A1 of FIG. 5A1; FIG. 5A3 is a diagram 500A3 for one-time programmable integrated circuit prior art implementation of the diagram 500A1 of FIG. 5A1; FIG. 5A4 is a diagram 500A4 for integrated circuit placement and route

20 of FIG. 5A1; FIG. 5A4 is a diagram 500A4 for integrated circuit placement and rout implementation of the diagram 500A1 of FIG. 5A1.

DETAILED DESCRIPTION OF THE INVENTION

The present invention is concerned with the VLSI layouts of arbitrarily large switching networks for broadcast, unicast and multicast connections. Particularly switching networks considered in the current invention include: generalized multi-stage

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networks $V(N_1, N_2, d, s)$, generalized folded multi-stage networks $V_{fold}(N_1, N_2, d, s)$, generalized butterfly fat tree networks $V_{bff}(N_1, N_2, d, s)$, generalized multi-link multistage networks $V_{mlink}(N_1, N_2, d, s)$, generalized folded multi-link multi-stage networks $V_{fold-mlink}(N_1, N_2, d, s)$, generalized multi-link butterfly fat tree networks

5 $V_{mlink-bft}(N_1, N_2, d, s)$, and generalized hypercube networks $V_{hcube}(N_1, N_2, d, s)$ for s = 1,2,3 or any number in general.

Efficient VLSI layout of networks on a semiconductor chip are very important and greatly influence many important design parameters such as the area taken up by the network on the chip, total number of wires, length of the wires, latency of the signals,

- 10 capacitance and hence the maximum clock speed of operation. Some networks may not even be implemented practically on a chip due to the lack of efficient layouts. The different varieties of multi-stage networks described above have not been implemented previously on the semiconductor chips efficiently. For example in Field Programmable Gate Array (FPGA) designs, multi-stage networks described in the current invention have
- 15 not been successfully implemented primarily due to the lack of efficient VLSI layouts. Current commercial FPGA products such as Xilinx Vertex, Altera's Stratix implement island-style architecture using mesh and segmented mesh routing interconnects using either full crossbars or sparse crossbars. These routing interconnects consume large silicon area for crosspoints, long wires, large signal propagation delay and hence
- 20 consume lot of power.

The current invention discloses the VLSI layouts of numerous types of multistage networks which are very efficient. Moreover they can be embedded on to mesh and segmented mesh routing interconnects of current commercial FPGA products. The VLSI layouts disclosed in the current invention are applicable to including the numerous generalized multi-stage networks disclosed in the following patent applications, filed

concurrently:

1) Strictly and rearrangeably nonblocking for arbitrary fan-out multicast and unicast for generalized multi-stage networks $V(N_1, N_2, d, s)$ with numerous connection

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topologies and the scheduling methods are described in detail in the PCT Application Serial No. PCT/US08/56064 that is incorporated by reference above.

2) Strictly and rearrangeably nonblocking for arbitrary fan-out multicast and unicast for generalized butterfly fat tree networks $V_{bft}(N_1, N_2, d, s)$ with numerous

5 connection topologies and the scheduling methods are described in detail in U.S. Provisional Patent Application Serial No. 60/940, 387 that is incorporated by reference above.

3) Rearrangeably nonblocking for arbitrary fan-out multicast and unicast, and strictly nonblocking for unicast for generalized multi-link multi-stage networks

10 $V_{mlink}(N_1, N_2, d, s)$ and generalized folded multi-link multi-stage networks $V_{fold-mlink}(N_1, N_2, d, s)$ with numerous connection topologies and the scheduling methods are described in detail in U.S. Provisional Patent Application Serial No. 60/940, 389 that is incorporated by reference above.

4) Strictly and rearrangeably nonblocking for arbitrary fan-out multicast and 15 unicast for generalized multi-link butterfly fat tree networks $V_{mlink-bft}(N_1, N_2, d, s)$ with numerous connection topologies and the scheduling methods are described in detail in U.S. Provisional Patent Application Serial No. 60/940, 390 that is incorporated by reference above.

5) Strictly and rearrangeably nonblocking for arbitrary fan-out multicast and 20 unicast for generalized folded multi-stage networks $V_{fold}(N_1, N_2, d, s)$ with numerous connection topologies and the scheduling methods are described in detail in U.S. Provisional Patent Application Serial No. 60/940, 391 that is incorporated by reference above.

6) Strictly nonblocking for arbitrary fan-out multicast for generalized multi-link
 multi-stage networks V_{mlink} (N₁, N₂, d, s) and generalized folded multi-link multi-stage
 networks V_{fold-mlink} (N₁, N₂, d, s) with numerous connection topologies and the scheduling

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methods are described in detail in U.S. Provisional Patent Application Serial No. 60/940, 392 that is incorporated by reference above.

7) VLSI layouts of numerous types of multi-stage networks with locality exploitation are described in U.S. Provisional Patent Application Serial No. 60/984, 724
that is incorporated by reference above.

8) VLSI layouts of numerous types of multistage pyramid networks are described in U.S. Provisional Patent Application Serial No. 61/018, 494 that is incorporated by reference above.

In addition the layouts of the current invention are also applicable to generalized 10 multi-stage pyramid networks $V_p(N_1, N_2, d, s)$, generalized folded multi-stage pyramid networks $V_{fold-p}(N_1, N_2, d, s)$, generalized butterfly fat pyramid networks $V_{bfp}(N_1, N_2, d, s)$, generalized multi-link multi-stage pyramid networks $V_{mlink-p}(N_1, N_2, d, s)$, generalized folded multi-link multi-stage pyramid networks $V_{fold-mlink-p}(N_1, N_2, d, s)$, generalized multi-link butterfly fat pyramid networks 15 March 20 March

15 $V_{mlink-bfp}(N_1, N_2, d, s)$, and generalized hypercube networks $V_{hcube}(N_1, N_2, d, s)$ for s = 1,2,3 or any number in general.

Symmetric RNB generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$:

Referring to diagram 100A in FIG. 1A, in one embodiment, an exemplary 20 generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with nine stages of one hundred and forty four switches for satisfying communication requests, such as setting up a telephone call or a data call, or a connection between configurable logic blocks, between an input stage 110 and output stage 120 via middle stages 130, 140, 150, 160, 170, 180 and 190 is shown where input stage 110

25 consists of sixteen, two by four switches IS1-IS16 and output stage 120 consists of sixteen, four by two switches OS1-OS16. And all the middle stages namely the middle

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stage 130 consists of sixteen, four by four switches MS(1,1) - MS(1,16), middle stage 140 consists of sixteen, four by four switches MS(2,1) - MS(2,16), middle stage 150 consists of sixteen, four by four switches MS(3,1) - MS(3,16), middle stage 160 consists of sixteen, four by four switches MS(4,1) - MS(4,16), middle stage 170 consists of

sixteen, four by four switches MS(5,1) - MS(5,16), middle stage 180 consists of sixteen, four by four switches MS(6,1) - MS(6,16), and middle stage 190 consists of sixteen, four by four switches MS(7,1) - MS(7,16).

As disclosed in U.S. Provisional Patent Application Serial No. 60/940,389 that is incorporated by reference above, such a network can be operated in rearrangeably non-

10 blocking manner for arbitrary fan-out multicast connections and also can be operated in strictly non-blocking manner for unicast connections.

In one embodiment of this network each of the input switches IS1-IS4 and output switches OS1-OS4 are crossbar switches. The number of switches of input stage 110 and of output stage 120 can be denoted in general with the variable $\frac{N}{d}$, where N is the total number of inlet links or outlet links. The number of middle switches in each middle stage 15 is denoted by $\frac{N}{d}$. The size of each input switch IS1-IS4 can be denoted in general with the notation d * 2d and each output switch OS1-OS4 can be denoted in general with the notation 2d * d. Likewise, the size of each switch in any of the middle stages can be denoted as 2d * 2d. A switch as used herein can be either a crossbar switch, or a 20 network of switches each of which in turn may be a crossbar switch or a network of switches. A symmetric multi-stage network can be represented with the notation $V_{{\it mlink}}(N,d,s)$, where N represents the total number of inlet links of all input switches (for example the links IL1-IL32), d represents the inlet links of each input switch or outlet links of each output switch, and s is the ratio of number of outgoing links from each input switch to the inlet links of each input switch. 25

Each of the $\frac{N}{d}$ input switches IS1 – IS16 are connected to exactly d switches in middle stage 130 through two links each for a total of $2 \times d$ links (for example input

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switch IS1 is connected to middle switch MS(1,1) through the links ML(1,1), ML(1,2), and also connected to middle switch MS(1,2) through the links ML(1,3) and ML(1,4)). The middle links which connect switches in the same row in two successive middle stages are called hereinafter straight middle links; and the middle links which connect

- 5 switches in different rows in two successive middle stages are called hereinafter cross middle links. For example, the middle links ML(1,1) and ML(1,2) connect input switch IS1 and middle switch MS(1,1), so middle links ML(1,1) and ML(1,2) are straight middle links; where as the middle links ML(1,3) and ML(1,4) connect input switch IS1 and middle switch MS(1,2), since input switch IS1 and middle switch MS(1,2) belong to two
- 10 different rows in diagram 100A of FIG. 1A, middle links ML(1,3) and ML(1,4) are cross middle links.

Each of the $\frac{N}{d}$ middle switches MS(1,1) – MS(1,16) in the middle stage 130 are

connected from exactly d input switches through two links each for a total of $2 \times d$ links (for example the links ML(1,1) and ML(1,2) are connected to the middle switch MS(1,1)

- from input switch IS1, and the links ML(1,7) and ML(1,8) are connected to the middle switch MS(1,1) from input switch IS2) and also are connected to exactly d switches in middle stage 140 through two links each for a total of 2×d links (for example the links ML(2,1) and ML(2,2) are connected from middle switch MS(1,1) to middle switch MS(2,1), and the links ML(2,3) and ML(2,4) are connected from middle switch MS(1,1)
- 20 to middle switch MS(2,3)).

Each of the
$$\frac{N}{d}$$
 middle switches MS(2,1) – MS(2,16) in the middle stage 140 are

connected from exactly d input switches through two links each for a total of $2 \times d$ links (for example the links ML(2,1) and ML(2,2) are connected to the middle switch MS(2,1) from input switch MS(1,1), and the links ML(1,11) and ML(1,12) are connected to the

25 middle switch MS(2,1) from input switch MS(1,3)) and also are connected to exactly d switches in middle stage 150 through two links each for a total of $2 \times d$ links (for example the links ML(3,1) and ML(3,2) are connected from middle switch MS(2,1) to middle switch MS(3,1), and the links ML(3,3) and ML(3,4) are connected from middle switch MS(2,1) to middle switch MS(2,1).

Each of the $\frac{N}{d}$ middle switches MS(3,1) – MS(3,16) in the middle stage 150 are connected from exactly d input switches through two links each for a total of $2 \times d$ links

(for example the links ML(3,1) and ML(3,2) are connected to the middle switch MS(3,1) from input switch MS(2,1), and the links ML(2,19) and ML(2,20) are connected to the

5 middle switch MS(3,1) from input switch MS(2,5)) and also are connected to exactly dswitches in middle stage 160 through two links each for a total of $2 \times d$ links (for example the links ML(4,1) and ML(4,2) are connected from middle switch MS(3,1) to middle switch MS(4,1), and the links ML(4,3) and ML(4,4) are connected from middle switch MS(3,1) to middle switch MS(4,9)).

Each of the N/d middle switches MS(4,1) – MS(4,16) in the middle stage 160 are connected from exactly d input switches through two links each for a total of 2×d links (for example the links ML(4,1) and ML(4,2) are connected to the middle switch MS(4,1) from input switch MS(3,1), and the links ML(4,35) and ML(4,36) are connected to the middle switch MS(4,1) from input switch MS(3,9)) and also are connected to exactly d switches in middle stage 170 through two links each for a total of 2×d links (for example the links ML(5,1) and ML(5,2) are connected from middle switch MS(4,1) to

middle switch MS(5,1), and the links ML(5,3) and ML(5,4) are connected from middle switch MS(4,1) to middle switch MS(5,9)).

Each of the
$$\frac{N}{d}$$
 middle switches MS(5,1) – MS(5,16) in the middle stage 170 are

- 20 connected from exactly d input switches through two links each for a total of $2 \times d$ links (for example the links ML(5,1) and ML(5,2) are connected to the middle switch MS(5,1) from input switch MS(4,1), and the links ML(5,35) and ML(5,36) are connected to the middle switch MS(5,1) from input switch MS(4,9)) and also are connected to exactly dswitches in middle stage 180 through two links each for a total of $2 \times d$ links (for
- 25 example the links ML(6,1) and ML(6,2) are connected from middle switch MS(5,1) to middle switch MS(6,1), and the links ML(6,3) and ML(6,4) are connected from middle switch MS(5,1) to middle switch MS(6,5)).

Each of the $\frac{N}{d}$ middle switches MS(6,1) – MS(6,16) in the middle stage 180 are connected from exactly *d* input switches through two links each for a total of $2 \times d$ links

(for example the links ML(6,1) and ML(6,2) are connected to the middle switch MS(6,1) from input switch MS(5,1), and the links ML(6,19) and ML(6,20) are connected to the

5 middle switch MS(6,1) from input switch MS(5,5)) and also are connected to exactly dswitches in middle stage 190 through two links each for a total of $2 \times d$ links (for example the links ML(7,1) and ML(7,2) are connected from middle switch MS(6,1) to middle switch MS(7,1), and the links ML(7,3) and ML(7,4) are connected from middle switch MS(6,1) to middle switch MS(7,3)).

Each of the N/d middle switches MS(7,1) – MS(7,16) in the middle stage 190 are connected from exactly d input switches through two links each for a total of 2×d links (for example the links ML(7,1) and ML(7,2) are connected to the middle switch MS(7,1) from input switch MS(6,1), and the links ML(7,11) and ML(7,12) are connected to the middle switch MS(7,1) from input switch MS(6,3)) and also are connected to exactly d switches in middle stage 120 through two links each for a total of 2×d links (for

example the links ML(8,1) and ML(8,2) are connected from middle switch MS(7,1) to middle switch MS(8,1), and the links ML(8,3) and ML(8,4) are connected from middle switch MS(7,1) to middle switch OS2).

Each of the
$$\frac{N}{d}$$
 middle switches OS1 – OS16 in the middle stage 120 are

20 connected from exactly *d* input switches through two links each for a total of $2 \times d$ links (for example the links ML(8,1) and ML(8,2) are connected to the output switch OS1 from input switch MS(7,1), and the links ML(8,7) and ML(7,8) are connected to the output switch OS1 from input switch MS(7,2)).

Finally the connection topology of the network 100A shown in FIG. 1A is known to be back to back inverse Benes connection topology.

Referring to diagram 100B in FIG. 1B, is a folded version of the multi-link multistage network 100A shown in FIG. 1A. The network 100B in FIG. 1B shows input stage

110 and output stage 120 are placed together. That is input switch IS1 and output switch OS1 are placed together, input switch IS2 and output switch OS2 are placed together, and similarly input switch IS16 and output switch OS16 are placed together. All the right going middle links (hereinafter "forward connecting links") {i.e., inlet links IL1 – IL32

- and middle links ML(1,1) ML(1,64)} correspond to input switches IS1 IS16, and all the left going middle links (hereinafter "backward connecting links") {i.e., middle links ML(8,1) ML(8,64) and outlet links OL1-OL32} correspond to output switches OS1 OS16.
- Middle stage 130 and middle stage 190 are placed together. That is middle
 switches MS(1,1) and MS(7,1) are placed together, middle switches MS(1,2) and
 MS(7,2) are placed together, and similarly middle switches MS(1,16) and MS(7,16) are
 placed together. All the right going middle links {i.e., middle links ML(1,1) ML(1,64)
 and middle links ML(2,1) ML(2,64)} correspond to middle switches MS(1,1) MS(1,16), and all the left going middle links {i.e., middle links ML(7,1) ML(7,64) and
- middle links ML(8,1) and ML(8,64)} correspond to middle switches MS(7,1) –MS(7,16).

Middle stage 140 and middle stage 180 are placed together. That is middle switches MS(2,1) and MS(6,1) are placed together, middle switches MS(2,2) and MS(6,2) are placed together, and similarly middle switches MS(2,16) and MS(6,16) are

- placed together. All the right going middle links {i.e., middle links ML(2,1) ML(2,64) and middle links ML(3,1) ML(3,64)} correspond to middle switches MS(2,1) MS(2,16), and all the left going middle links {i.e., middle links ML(6,1) ML(6,64) and middle links ML(7,1) and ML(7,64)} correspond to middle switches MS(6,1) MS(6,16).
- 25 Middle stage 150 and middle stage 170 are placed together. That is middle switches MS(3,1) and MS(5,1) are placed together, middle switches MS(3,2) and MS(5,2) are placed together, and similarly middle switches MS(3,16) and MS(5,16) are placed together. All the right going middle links {i.e., middle links ML(3,1) ML(3,64) and middle links ML(4,1) ML(4,64)} correspond to middle switches MS(3,1) -
- 30 MS(3,16), and all the left going middle links {i.e., middle links ML(5,1) ML(5,64) and

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middle links ML(6,1) and ML(6,64)} correspond to middle switches MS(5,1) - MS(5,16).

Middle stage 160 is placed alone. All the right going middle links are the middle links ML(4,1) - ML(4,64) and all the left going middle links are middle links ML(5,1) ML(5,64).

In one embodiment, in the network 100B of FIG. 1B, the switches that are placed together are implemented as separate switches then the network 100B is the generalized folded multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with nine stages as disclosed in U.S. Provisional Patent Application Serial No.

- 10 60/940,389 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a two by four switch and a four by two switch. For example the switch input switch IS1 and output switch OS1 are placed together; so input switch IS1 is implemented as two by four switch with the inlet links IL1 and IL2 being the inputs of the input switch IS1 and middle links ML(1,1)
- 15 ML(1,4) being the outputs of the input switch IS1; and output switch OS1 is implemented as four by two switch with the middle links ML(8,1), ML(8,2), ML(8,7) and ML(8,8) being the inputs of the output switch OS1 and outlet links OL1 OL2 being the outputs of the output switch OS1. Similarly in this embodiment of network 100B all the switches that are placed together in each middle stage are implemented as separate
 20 switches,

Hypercube Topology layout schemes:

Referring to layout 100C of FIG. 1C, in one embodiment, there are sixteen blocks namely Block 1_2, Block 3_4, Block 5_6, Block 7_8, Block 9_10, Block 11_12, Block

13_14, Block 15_16, Block 17_18, Block 19_20, Block 21_22, Block 23_24, Block
25_26, Block 27_28, Block 29_30, and Block 31_32. Each block implements all the switches in one row of the network 100B of FIG. 1B, one of the key aspects of the current invention. For example Block 1_2 implements the input switch IS1, output Switch OS1,

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middle switch MS(1,1), middle switch MS(7,1), middle switch MS(2,1), middle switch MS(6,1), middle switch MS(3,1), middle switch MS(5,1), and middle switch MS(4,1). For the simplification of illustration, Input switch IS1 and output switch OS1 together are denoted as switch 1; Middle switch MS(1,1) and middle switch MS(7,1) together are

5 denoted by switch 2; Middle switch MS(2,1) and middle switch MS(6,1) together are denoted by switch 3; Middle switch MS(3,1) and middle switch MS(5,1) together are denoted by switch 4; Middle switch MS(4,1) is denoted by switch 5.

All the straight middle links are illustrated in layout 100C of FIG. 1C. For example in Block 1_2, inlet links IL1 – IL2, outlet links OL1 - OL2, middle link

- 10 ML(1,1), middle link ML(1,2), middle link ML(8,1), middle link ML(8,2), middle link ML(2,1), middle link ML(2,2), middle link ML(7,1), middle link ML(7,2), middle link ML(3,1), middle link ML(3,2), middle link ML(6,1), middle link ML(6,2), middle link ML(4,1), middle link ML(4,2), middle link ML(5,1) and middle link ML(5,2) are illustrated in layout 100C of FIG. 1C.
- 15 Even though it is not illustrated in layout 100C of FIG. 1C, in each block, in addition to the switches there may be Configurable Logic Blocks (CLB) or any arbitrary digital circuit (hereinafter "sub-integrated circuit block") depending on the applications in different embodiments. There are four quadrants in the layout 100C of FIG. 1C namely top-left, bottom-left, top-right and bottom-right quadrants. Top-left quadrant implements
- Block 1_2, Block 3_4, Block 5_6, and Block 7_8. Bottom-left quadrant implements
 Block 9_10, Block 11_12, Block 13_14, and Block 15_16. Top-right quadrant
 implements Block 17_18, Block 19_20, Block 21_22, and Block 23_24. Bottom-right
 quadrant implements Block 25_26, Block 27_28, Block 29_30, and Block 31_32. There
 are two halves in layout 100C of FIG. 1C namely left-half and right-half. Left-half
 consists of top-left and bottom-left quadrants. Right-half consists of top-right and bottom-

Recursively in each quadrant there are four sub-quadrants. For example in top-left quadrant there are four sub-quadrants namely top-left sub-quadrant, bottom-left sub-quadrant, top-right sub-quadrant and bottom-right sub-quadrant. Top-left sub-quadrant of top-left quadrant implements Block 1_2. Bottom-left sub-quadrant of top-left quadrant

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implements Block 3_4. Top-right sub-quadrant of top-left quadrant implements Block 5_6. Finally bottom-right sub-quadrant of top-left quadrant implements Block 7_8. Similarly there are two sub-halves in each quadrant. For example in top-left quadrant there are two sub-halves namely left-sub-half and right-sub-half. Left-sub-half of top-left

- 5 quadrant implements Block 1_2 and Block 3_4. Right-sub-half of top-left quadrant implements Block 5_6 and Block 7_8. Finally applicant notes that in each quadrant or half the blocks are arranged as a general binary hypercube. Recursively in larger multistage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 > 32$, the layout in this embodiment in accordance with the current invention, will be such that the super-quadrants will also
- 10 be arranged in d-ary hypercube manner. (In the embodiment of the layout 100C of FIG. 1C, it is binary hypercube manner since d = 2, in the network $V_{fold-mlink}(N_1, N_2, d, s)$ 100B of FIG. 1B).

Layout 100D of FIG. 1D illustrates the inter-block links between switches 1 and 2 of each block. For example middle links ML(1,3), ML(1,4), ML(8,7), and ML(8,8) are

- 15 connected between switch 1 of Block 1_2 and switch 2 of Block 3_4. Similarly middle links ML(1,7), ML(1,8), ML(8,3), and ML(8,4) are connected between switch 2 of Block 1_2 and switch 1 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 100D of FIG. 1D can be implemented as vertical tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for example
- 20 middle links ML(1,4) and ML(8,8) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track).

Layout 100E of FIG. 1E illustrates the inter-block links between switches 2 and 3 of each block. For example middle links ML(2,3), ML(2,4), ML(7,11), and ML(7,12) are connected between switch 2 of Block 1_2 and switch 3 of Block 3_4. Similarly middle links ML(2,11), ML(2,12), ML(7,3), and ML(7,4) are connected between switch 3 of Block 1_2 and switch 2 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 100E of FIG. 1E can be implemented as horizontal tracks in one embodiment.

30 Also in one embodiment inter-block links are implemented as two different tracks (for

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example middle links ML(2,12) and ML(7,4) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(2,12) and ML(7,4) are implemented as a time division multiplexed single track).

5 Layout 100F of FIG. 1F illustrates the inter-block links between switches 3 and 4 of each block. For example middle links ML(3,3), ML(3,4), ML(6,19), and ML(6,20) are connected between switch 3 of Block 1_2 and switch 4 of Block 3_4. Similarly middle links ML(3,19), ML(3,20), ML(6,3), and ML(6,4) are connected between switch 4 of Block 1_2 and switch 3 of Block 3_4. Applicant notes that the inter-block links illustrated

- in layout 100F of FIG. 1F can be implemented as vertical tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(3,4) and ML(6,20) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track).
- 15 division multiplexed single track).

Layout 100G of FIG. 1G illustrates the inter-block links between switches 4 and 5 of each block. For example middle links ML(4,3), ML(4,4), ML(5,35), and ML(5,36) are connected between switch 4 of Block 1_2 and switch 5 of Block 3_4. Similarly middle links ML(4,35), ML(4,36), ML(5,3), and ML(5,4) are connected between switch 5 of

- Block 1_2 and switch 4 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 100G of FIG. 1G can be implemented as horizontal tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(4,4) and ML(5,36) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division
 multiplexed single track (for example middle links ML(4,4) and ML(5,36) are
 - implemented as a time division multiplexed single track).

The complete layout for the network 100B of FIG. 1B is given by combining the links in layout diagrams of 100C, 100D, 100E, 100F, and 100G. Applicant notes that in the layout 100C of FIG. 1C, the inter-block links between switch 1 and switch 2 of

30 corresponding blocks are vertical tracks as shown in layout 100D of FIG. 1D; the inter-

block links between switch 2 and switch 3 of corresponding blocks are horizontal tracks as shown in layout 100E of FIG. 1E; the inter-block links between switch 3 and switch 4 of corresponding blocks are vertical tracks as shown in layout 100F of FIG. 1F; and finally the inter-block links between switch 4 and switch 5 of corresponding blocks are

5 horizontal tracks as shown in layout 100G of FIG. 1G. The pattern is alternate vertical tracks and horizontal tracks. It continues recursively for larger networks of N > 32 as will be illustrated later.

Some of the key aspects of the current invention are discussed. 1) All the switches in one row of the multi-stage network 100B are implemented in a single block. 2) The blocks are placed in such a way that all the inter-block links are either horizontal tracks or vertical tracks; 3) Since all the inter-block links are either horizontal or vertical tracks, all the inter-block links can be mapped on to island-style architectures in current commercial FPGA's; 4) The length of the longest wire is about half of the width (or length) of the complete layout (For example middle link ML(4,4) is about half the width of the

15 complete layout).

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In accordance with the current invention, the layout 100C in FIG. 1C can be recursively extended for any arbitrarily large generalized folded multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ the sub-quadrants, quadrants, and super-quadrants are arranged in d-ary hypercube manner and also the inter-blocks are accordingly connected in d-ary hypercube topology. Even though all the embodiments in the current invention are illustrated for $N_1 = N_2$, the embodiments can be extended for $N_1 \neq N_2$.

Referring to layout 100H of FIG. 1H, illustrates the extension of layout 100C for the network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 128$; d = 2; and s = 2. There are four super-quadrants in layout 100H namely top-left super-quadrant, bottom-left super-

25 quadrant, top-right super-quadrant, bottom-right super-quadrant. Total number of blocks in the layout 100H is sixty four. Top-left super-quadrant implements the blocks from block 1_2 to block 31_32. Each block in all the super-quadrants has two more switches namely switch 6 and switch 7 in addition to the switches [1-5] illustrated in layout 100C of FIG. 1C. The inter-block link connection topology is the exactly the same between the

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switches 1 and 2; switches 2 and 3; switches 3 and 4; switches 4 and 5 as it is shown in the layouts of FIG. 1D, FIG. 1E, FIG. 1F, and FIG. 1G respectively.

Bottom-left super-quadrant implements the blocks from block 33_34 to block 63_64. Top-right super-quadrant implements the blocks from block 65_66 to block 95_96. And bottom-right super-quadrant implements the blocks from block 97_98 to block 127_128. In all these three super-quadrants also, the inter-block link connection topology is the exactly the same between the switches 1 and 2; switches 2 and 3; switches 3 and 4; switches 4 and 5 as that of the top-left super-quadrant.

Recursively in accordance with the current invention, the inter-block links connecting the switch 5 and switch 6 will be vertical tracks between the corresponding switches of top-left super-quadrant and bottom-left super-quadrant. And similarly the inter-block links connecting the switch 5 and switch 6 will be vertical tracks between the corresponding switches of top-right super-quadrant and bottom-right super-quadrant. The inter-block links connecting the switch 6 and switch 7 will be horizontal tracks between

- 15 the corresponding switches of top-left super-quadrant and top-right super-quadrant. And similarly the inter-block links connecting the switch 6 and switch 7 will be horizontal tracks between the corresponding switches of bottom-left super-quadrant and bottomright super-quadrant.
- Referring to diagram 100I of FIG. 1I illustrates a high-level implementation of
 Block 1_2 (Each of the other blocks have similar implementation) of the layout 100C of
 FIG. 1C which represents a generalized folded multi-link multi-stage network
 V_{fold-mlink} (N₁, N₂, d, s) where N₁ = N₂ = 32; d = 2; and s = 2. Block 1_2 in 100I illustrates
 both the intra-block and inter-block links connected to Block 1_2. The layout diagram
 100I corresponds to the embodiment where the switches that are placed together are
 implemented as separate switches in the network 100B of FIG. 1B. As noted before then
 the network 100B is the generalized folded multi-link multi-stage network
 V_{fold-mlink} (N₁, N₂, d, s) where N₁ = N₂ = 32; d = 2; and s = 2 with nine stages as disclosed
 in U.S. Provisional Patent Application Serial No. 60/940,389 that is incorporated by

reference above.

That is the switches that are placed together in Block 1_2 as shown in FIG. 1I are namely input switch IS1 and output switch OS1 belonging to switch 1, illustrated by dotted lines, (as noted before switch 1 is for illustration purposes only, in practice the switches implemented are input switch IS1 and output switch OS1); middle switch

5 MS(1,1) and middle switch MS(7,1) belonging to switch 2; middle switch MS(2,1) and middle switch MS(6,1) belonging to switch 3; middle switch MS(3,1) and middle switch MS(5,1) belonging to switch 4; And middle switch MS(4,1) belonging to switch 5.

Input switch IS1 is implemented as two by four switch with the inlet links IL1 and IL2 being the inputs of the input switch IS1 and middle links ML(1,1) – ML(1,4) being
the outputs of the input switch IS1; and output switch OS1 is implemented as four by two switch with the middle links ML(8,1) – ML(8,4) being the inputs of the output switch OS1 and outlet links OL1 – OL2 being the outputs of the output switch OS1.

Middle switch MS(1,1) is implemented as four by four switch with the middle links ML(1,1), ML(1,2), ML(1,7) and ML(1,8) being the inputs and middle links ML(2,1)
- ML(2,4) being the outputs; and middle switch MS(7,1) is implemented as four by four switch with the middle links ML(7,1), ML(7,2), ML(7,11) and ML(7,12) being the inputs and middle links ML(8,1) – ML(8,4) being the outputs. Similarly all the other middle switches are also implemented as four by four switches as illustrated in 100I of FIG. 1I.

- Now the VLSI layouts of generalized multi-link multi-stage network 20 $V_{mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 < 32$; d = 2; s = 2 and its corresponding version of folded generalized multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2$ < 32; d = 2; s = 2 are discussed. Referring to diagram 200A1 of FIG. 2A1 is generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 2$; d = 2. Diagram 200A2 of FIG. 2A2 illustrates the corresponding folded generalized multi-link multi-
- stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 2$; d = 2, version of the diagram 200A1 of FIG. 2A1. Layout 200A3 of FIG. 2A3 illustrates the VLSI layout of the network 200A2 of FIG. 2A2. There is only one block i.e., Block 1_2 comprising switch 1. Just like in the layout 100C of FIG. 1C, switch 1 consists of input switch IS1 and output switch OS1.

Referring to diagram 200B1 of FIG. 2B1 is generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 4$; d = 2; s = 2. Diagram 200B2 of FIG. 2B2 illustrates the corresponding folded generalized multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 4$; d = 2; s = 2, version of the diagram 200B1 of

- 5 FIG. 2B1. Layout 200B3 of FIG. 2B3 illustrates the VLSI layout of the network 200B2 of FIG. 2B2. There are two blocks i.e., Block 1_2 and Block 3_4 each comprising switch 1 and switch 2. Switch 1 in each block consists of the corresponding input switch and output switch. For example switch 1 in Block 1_2 consists of input switch IS1 and output switch OS1. Similarly switch 2 in Block 1_2 consists of middle switch (1,1). Layout
- 10 200B4 of FIG. 2B4 illustrates the inter-block links of the VLSI layout diagram 200B3 of FIG. 2B3. For example middle links ML(1,4) and ML(2,8). It must be noted that all the inter-block links are vertical tracks in this layout. (Alternatively all the inter-blocks can also be implemented as horizontal tracks).

Referring to diagram 200C11 of FIG. 2C11 is generalized multi-link multi-stage 15 network $V_{mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 8$; d = 2; s = 2. Diagram 200C12 of FIG. 2C12 illustrates the corresponding folded generalized multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 8$; d = 2; s = 2, version of the diagram 200C11 of FIG. 2C11. Layout 200C21 of FIG. 2C21 illustrates the VLSI layout of the network 200C12 of FIG. 2C12. There are four blocks i.e., Block 1_2, Block 3_4, Block 5_6, and

Block 7_8 each comprising switch 1, switch 2 and switch 3. For example switch 1 in
 Block 1_2 consists of input switch IS1 and output switch OS1; Switch 2 in Block 1_2
 consists of MS(1,1) and MS(3,1). Switch 3 in Block 1_2 consists of MS(2,1).

Layout 200C22 of FIG. 2C22 illustrates the inter-block links between the switch 1 and switch 2 of the VLSI layout diagram 200C21 of FIG. 2C21. For example middle

25 links ML(1,4) and ML(4,8) are connected between Block 1_2 and Block 3_4. It must be noted that all the inter-block links between switch 1 and switch 2 of all blocks are vertical tracks in this layout. Layout 200C23 of FIG. 2C23 illustrates the inter-block links between the switch 2 and switch 3 of the VLSI layout diagram 200C21 of FIG. 2C21. For example middle links ML(2,12) and ML(3,4) are connected between Block 1_2 and

Block 5_6. It must be noted that all the inter-block links between switch 2 and switch 3 of all blocks are horizontal tracks in this layout

Referring to diagram 200D1 of FIG. 2D1 is generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 16$; d = 2; s = 2. Diagram 200D2 of FIG.

- 2D2 illustrates the corresponding folded generalized multi-link multi-stage network
 V_{fold-mlink} (N₁, N₂, d, s) where N₁ = N₂ = 16; d = 2; s = 2, version of the diagram 200D1 of
 FIG. 2D1. Layout 200D3 of FIG. 2D3 illustrates the VLSI layout of the network 200D2 of FIG. 2D2. There are eight blocks i.e., Block 1_2, Block 3_4, Block 5_6, Block 7_8, Block 9 10, Block 11 12, Block 13 14 and Block 15 16 each comprising switch 1,
- switch 2, switch 3 and switch 4. For example switch 1 in Block 1_2 consists of input switch IS1 and output switch OS1; Switch 2 in Block 1_2 consists of MS(1,1) and MS(5,1). Switch 3 in Block 1_2 consists of MS(2,1) and MS(4,1), and switch 4 in Block 1_2 consists of MS(3,1).
- Layout 200D4 of FIG. 2D4 illustrates the inter-block links between the switch 1 and switch 2 of the VLSI layout diagram 200D3 of FIG. 2D3. For example middle links ML(1,4) and ML(6,8) are connected between Block 1_2 and Block 3_4. It must be noted that all the inter-block links between switch 1 and switch 2 of all blocks are vertical tracks in this layout. Layout 200D5 of FIG. 2D5 illustrates the inter-block links between the switch 2 and switch 3 of the VLSI layout diagram 200D3 of FIG. 2D3. For example
- 20 middle links ML(2,12) and ML(5,4) are connected between Block 1_2 and Block 5_6. It must be noted that all the inter-block links between switch 2 and switch 3 of all blocks are horizontal tracks in this layout. Layout 200D6 of FIG. 2D6 illustrates the inter-block links between the switch 3 and switch 4 of the VLSI layout diagram 200D3 of FIG. 2D3. For example middle links ML(3,4) and ML(4,20) are connected between Block 1_2 and
- 25 Block 9_10. It must be noted that all the inter-block links between switch 3 and switch 4 of all blocks are vertical tracks in this layout.

Generalized Multi-link Butterfly Fat Tree Network Embodiment:

In another embodiment in the network 100B of FIG. 1B, the switches that are placed together are implemented as combined switch then the network 100B is the

generalized multi-link butterfly fat tree network $V_{mlink-bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with five stages as disclosed in U.S. Provisional Patent Application Serial No. 60/940,390 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a six by

- 5 six switch. For example the input switch IS1 and output switch OS1 are placed together; so input switch IS1 and output OS1 are implemented as a six by six switch with the inlet links IL1, IL2, ML(8,1), ML(8,2), ML(8,7) and ML(8,8) being the inputs of the combined switch (denoted as IS1&OS1) and middle links ML(1,1), ML(1,2), ML(1,3), ML(1,4), OL1 and OL2 being the outputs of the combined switch IS1&OS1. Similarly in
- 10 this embodiment of network 100B all the switches that are placed together are implemented as a combined switch.

Layout diagrams 100C in FIG. 1C, 100D in FIG. 1D, 100E in FIG. 1E, 100F in FIG. 1G are also applicable to generalized multi-link butterfly fat tree network $V_{mlink-bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with five stages. The layout 100C in FIG. 1C can be recursively extended for any arbitrarily large generalized multilink butterfly fat tree network $V_{mlink-bft}(N_1, N_2, d, s)$. Accordingly layout 100H of FIG. 1H is also applicable to generalized multi-link butterfly fat tree network $V_{mlink-bft}(N_1, N_2, d, s)$.

Referring to diagram 100J of FIG. 1J illustrates a high-level implementation of 20 Block 1_2 (Each of the other blocks have similar implementation) of the layout 100C of FIG. 1C which represents a generalized multi-link butterfly fat tree network $V_{mlink-bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2. Block 1_2 in 100J illustrates both the intra-block and inter-block links. The layout diagram 100J corresponds to the embodiment where the switches that are placed together are implemented as combined 25 switch in the network 100B of FIG. 1B. As noted before then the network 100B is the generalized multi-link butterfly fat tree network $V_{mlink-bft}(N_1, N_2, d, s)$ where $N_1 = N_2 =$ 32; d = 2; and s = 2 with five stages as disclosed in U.S. Provisional Patent Application

Serial No. 60/940,390 that is incorporated by reference above.

That is the switches that are placed together in Block 1_2 as shown in FIG. 1J are namely the combined input and output switch IS1&OS1 belonging to switch 1, illustrated by dotted lines, (as noted before switch 1 is for illustration purposes only, in practice the switch implemented is combined input and output switch IS1&OS1); middle switch

5 MS(1,1) belonging to switch 2; middle switch MS(2,1) belonging to switch 3; middle switch MS(3,1) belonging to switch 4; And middle switch MS(4,1) belonging to switch 5.

Combined input and output switch IS1&OS1 is implemented as six by six switch with the inlet links IL1, IL2 and ML(8,1) - ML(8,4) being the inputs and middle links ML(1,1) - ML(1,4), and outlet links OL1 - OL2 being the outputs.

Middle switch MS(1,1) is implemented as eight by eight switch with the middle links ML(1,1), ML(1,2), ML(1,7), ML(1,8), ML(7,1), ML(7,2), ML(7,11) and ML(7,12) being the inputs and middle links ML(2,1) – ML(2,4) and middle links ML(8,1) – ML(8,4) being the outputs. Similarly all the other middle switches are also implemented as eight by eight switches as illustrated in 100J of FIG. 1J. Applicant observes that in
middle switch MS(1,1) any one of the right going middle links can be switched to any one of the left going middle links and hereinafter middle switch MS(1,1) provides U-turn links. In general, in the network V_{mlink-bft} (N₁, N₂, d, s) each input switch, each output switch and each middle switch provides U-turn links.

- In another embodiment, middle switch MS(1,1) (or the middle switches in any of
 the middle stage excepting the root middle stage) of Block 1_2 of
 V_{mlink-bft} (N₁, N₂, d, s) can be implemented as a four by eight switch and a four by four
 switch to save cross points. This is because the left going middle links of these middle
 switches are never setup to the right going middle links. For example, in middle switch
 MS(1,1) of Block 1_2 as shown FIG. 1J, the left going middle links namely ML(7,1),
 ML(7,2), ML(7,11), and ML(7,12) are never switched to the right going middle links
- ML(2,1), ML(2,2), ML(2,3), and ML(2,4). And hence to implement MS(1,1) two switches namely: 1) a four by eight switch with the middle links ML(1,1), ML(1,2), ML(1,7), and ML(1,8) as inputs and the middle links ML(2,1), ML(2,2), ML(2,3), ML(2,4), ML(8,1), ML(8,2), ML(8,3), and ML(8,4) as outputs and 2) a four by four

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switch with the middle links ML(7,1), ML(7,2), ML(7,11), and ML(7,12) as inputs and the middle links ML(8,1), ML(8,2), ML(8,3), and ML(8,4) as outputs are sufficient without loosing any connectivity of the embodiment of MS(1,1) being implemented as an eight by eight switch as described before.)

5 Generalized multi-stage network Embodiment:

In one embodiment, in the network 100B of FIG. 1B, the switches that are placed together are implemented as two separate switches in input stage 110 and output stage 120; and as four separate switches in all the middle stages, then the network 100B is the generalized folded multi-stage network $V_{fold}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2;

- 10 and s = 2 with nine stages as disclosed in U.S. Provisional Patent Application Serial No. 60/940,391 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a two by four switch and a four by two switch respectively. For example the switch input switch IS1 and output switch OS1 are placed together; so input switch IS1 is implemented as two by four
- 15 switch with the inlet links IL1 and IL2 being the inputs and middle links ML(1,1) ML(1,4) being the outputs; and output switch OS1 is implemented as four by two switch with the middle links ML(8,1), ML(8,4), ML(8,7) and ML(8,8) being the inputs and outlet links OL1 – OL2 being the outputs.
- The switches, corresponding to the middle stages that are placed together are 20 implemented as four two by two switches. For example middle switches MS(1,1), MS(1,17), MS(7,1), and MS(7,17) are placed together; so middle switch MS(1,1) is implemented as two by two switch with middle links ML(1,1) and ML(1,7) being the inputs and middle links ML(2,1) and ML(2,3) being the outputs; middle switch MS(1,17) is implemented as two by two switch with the middle links ML(1,2) and ML(1,8) being
- 25 the inputs and middle links ML(2,2) and ML(2,4) being the outputs; middle switch MS(7,1) is implemented as two by two switch with middle links ML(7,1) and ML(7,11) being the inputs and middle links ML(8,1) and ML(8,3) being the outputs; And middle switch MS(7,17) is implemented as two by two switch with the middle links ML(7,2) and ML(7,12) being the inputs and middle links ML(8,2) and ML(8,4) being the outputs;

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Similarly in this embodiment of network 100B all the switches that are placed together are implemented as separate switches.

Layout diagrams 100C in FIG. 1C, 100D in FIG. 1D, 100E in FIG. 1E, 100F in FIG. 1G are also applicable to generalized folded multi-stage network

- 5 $V_{fold}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with nine stages. The layout 100C in FIG. 1C can be recursively extended for any arbitrarily large generalized folded multistage network $V_{fold}(N_1, N_2, d, s)$. Accordingly layout 100H of FIG. 1H is also applicable to generalized folded multi-stage network $V_{fold}(N_1, N_2, d, s)$.
- Referring to diagram 100K of FIG. 1K illustrates a high-level implementation of 10 Block 1_2 (Each of the other blocks have similar implementation) of the layout 100C of FIG. 1C which represents a generalized folded multi-stage network $V_{fold} (N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2. Block 1_2 in 100K illustrates both the intra-block and inter-block links. The layout diagram 100K corresponds to the embodiment where the switches that are placed together are implemented as separate 15 switches in the network 100B of FIG. 1B. As noted before then the network 100B is the generalized folded multi-stage network $V_{fold} (N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2;

and s = 2 with nine stages as disclosed in U.S. Provisional Patent Application Serial No. 60/940,391 that is incorporated by reference above.

- That is the switches that are placed together in Block 1_2 as shown in FIG. 1K are
 namely the input switch IS1 and output switch OS1 belonging to switch 1, illustrated by
 dotted lines, (as noted before switch 1 is for illustration purposes only, in practice the
 switches implemented are input switch IS1 and output switch OS1); middle switches
 MS(1,1), MS(1,17), MS(7,1) and MS(7,17) belonging to switch 2; middle switches
 MS(2,1), MS(2,17), MS(6,1) and MS(6,17) belonging to switch 3; middle switches
 MS(3,1), MS(3,17), MS(5,1) and MS(5,17) belonging to switch 4; And middle switches
 - MS(4,1), and MS(4,17) belonging to switch 5.

Input switch IS1 and output switch OS1 are placed together; so input switch IS1 is implemented as two by four switch with the inlet links IL1 and IL2 being the inputs and

middle links ML(1,1) - ML(1,4) being the outputs; and output switch OS1 is implemented as four by two switch with the middle links ML(8,1), ML(8,4), ML(8,7) and ML(8,8) being the inputs and outlet links OL1 - OL2 being the outputs.

Middle switches MS(1,1), MS(1,17), MS(7,1), and MS(7,17) are placed together; 5 so middle switch MS(1,1) is implemented as two by two switch with middle links ML(1,1) and ML(1,7) being the inputs and middle links ML(2,1) and ML(2,3) being the outputs; middle switch MS(1,17) is implemented as two by two switch with the middle links ML(1,2) and ML(1,8) being the inputs and middle links ML(2,2) and ML(2,4) being the outputs; middle switch MS(7,1) is implemented as two by two switch with middle

10 links ML(7,1) and ML(7,11) being the inputs and middle links ML(8,1) and ML(8,3)being the outputs; And middle switch MS(7,17) is implemented as two by two switch with the middle links ML(7,2) and ML(7,12) being the inputs and middle links ML(8,2)and ML(8,4) being the outputs. Similarly all the other middle switches are also implemented as two by two switches as illustrated in 100K of FIG. 1K.

15 Generalized multi-stage network Embodiment with S = 1:

In one embodiment, in the network 100B of FIG. 1B (where it is implemented with s = 1), the switches that are placed together are implemented as two separate switches in input stage 110 and output stage 120; and as two separate switches in all the middle stages, then the network 100B is the generalized folded multi-stage network

- $V_{fold}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 1 with nine stages as disclosed in U.S. Provisional Patent Application Serial No. 60/940,391 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a two by two switch and a two by two switch. For example the switch input switch IS1 and output switch OS1 are placed together; so input
- 25 switch IS1 is implemented as two by two switch with the inlet links IL1 and IL2 being the inputs and middle links ML(1,1) - ML(1,2) being the outputs; and output switch OS1 is implemented as two by two switch with the middle links ML(8,1) and ML(8,3) being the inputs and outlet links OL1 – OL2 being the outputs.

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The switches, corresponding to the middle stages that are placed together are implemented as two, two by two switches. For example middle switches MS(1,1) and MS(7,1) are placed together; so middle switch MS(1,1) is implemented as two by two switch with middle links ML(1,1) and ML(1,3) being the inputs and middle links

- 5 ML(2,1) and ML(2,2) being the outputs; middle switch MS(7,1) is implemented as two by two switch with middle links ML(7,1) and ML(7,5) being the inputs and middle links ML(8,1) and ML(8,2) being the outputs; Similarly in this embodiment of network 100B all the switches that are placed together are implemented as two separate switches.
- Layout diagrams 100C in FIG. 1C, 100D in FIG. 1D, 100E in FIG. 1E, 100F in FIG. 1G are also applicable to generalized folded multi-stage network $V_{fold} (N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 1 with nine stages. The layout 100C in FIG. 1C can be recursively extended for any arbitrarily large generalized folded multistage network $V_{fold} (N_1, N_2, d, s)$. Accordingly layout 100H of FIG. 1H is also applicable to generalized folded multi-stage network $V_{fold} (N_1, N_2, d, s)$.
- Referring to diagram 100K1 of FIG. 1K1 illustrates a high-level implementation of Block 1_2 (Each of the other blocks have similar implementation) for the layout 100C of FIG. 1C when s = 1 which represents a generalized folded multi-stage network V_{fold} (N₁, N₂, d, s) where N₁ = N₂ = 32; d = 2; and s = 1 (All the double links are replaced by single links when s = 1). Block 1_2 in 100K1 illustrates both the intra-block and interblock links. The layout diagram 100K1 corresponds to the embodiment where the switches that are placed together are implemented as separate switches in the network 100B of FIG. 1B when s = 1. As noted before then the network 100B is the generalized folded multi-stage network V_{fold} (N₁, N₂, d, s) where N₁ = N₂ = 32; d = 2; and s = 1 with nine stages as disclosed in U.S. Provisional Patent Application Serial No. 60/940,391 that is incorporated by reference above.

That is the switches that are placed together in Block 1_2 as shown in FIG. 1K1 are namely the input switch IS1 and output switch OS1 belonging to switch 1, illustrated by dotted lines, (as noted before switch 1 is for illustration purposes only, in practice the switches implemented are input switch IS1 and output switch OS1); middle switches

MS(1,1) and MS(7,1) belonging to switch 2; middle switches MS(2,1) and MS(6,1) belonging to switch 3; middle switches MS(3,1) and MS(5,1) belonging to switch 4; And middle switch MS(4,1) belonging to switch 5.

Input switch IS1 and output switch OS1 are placed together; so input switch IS1 is
implemented as two by two switch with the inlet links IL1 and IL2 being the inputs and middle links ML(1,1) – ML(1,2) being the outputs; and output switch OS1 is implemented as two by two switch with the middle links ML(8,1) and ML(8,3) being the inputs and outlet links OL1 – OL2 being the outputs.

Middle switches MS(1,1) and MS(7,1) are placed together; so middle switch

- 10 MS(1,1) is implemented as two by two switch with middle links ML(1,1) and ML(1,3) being the inputs and middle links ML(2,1) and ML(2,2) being the outputs; And middle switch MS(7,1) is implemented as two by two switch with middle links ML(7,1) and ML(7,5) being the inputs and middle links ML(8,1) and ML(8,2) being the outputs. Similarly all the other middle switches are also implemented as two by two switches as
- 15 illustrated in 100K1 of FIG. 1K1.

Generalized Butterfly Fat Tree Network Embodiment:

In another embodiment in the network 100B of FIG. 1B, the switches that are placed together are implemented as two combined switches then the network 100B is the generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s

- 20 = 2 with five stages as disclosed in U.S. Provisional Patent Application Serial No.
 60/940,387 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a six by six switch. For example the input switch IS1 and output switch OS1 are placed together; so input output switch IS1&OS1 are implemented as a six by six switch with the inlet links IL1,
- IL2, ML(8,1), ML(8,2), ML(8,7) and ML(8,8) being the inputs of the combined switch (denoted as IS1&OS1) and middle links ML(1,1), ML(1,2), ML(1,3), ML(1,4), OL1 and OL2 being the outputs of the combined switch IS1&OS1.

The switches, corresponding to the middle stages that are placed together are implemented as two four by four switches. For example middle switches MS(1,1) and MS(1,17) are placed together; so middle switch MS(1,1) is implemented as four by four switch with middle links ML(1,1), ML(1,7), ML(7,1) and ML(7,11) being the inputs and

- 5 middle links ML(2,1), ML(2,3), ML(8,1) and ML(8,3) being the outputs; middle switch MS(1,17) is implemented as four by four switch with the middle links ML(1,2), ML(1,8), ML(7,2) and ML(7,12) being the inputs and middle links ML(2,2), ML(2,4), ML(8,2) and ML(8,4) being the outputs. Similarly in this embodiment of network 100B all the switches that are placed together are implemented as a two combined switches.
 - Layout diagrams 100C in FIG. 1C, 100D in FIG. 1D, 100E in FIG. 1E, 100F in FIG. 1G are also applicable to generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with five stages. The layout 100C in FIG. 1C can be recursively extended for any arbitrarily large generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$. Accordingly layout 100H of FIG. 1H is also applicable to
- 15 generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$.

Referring to diagram 100L of FIG. 1L illustrates a high-level implementation of Block 1_2 (Each of the other blocks have similar implementation) of the layout 100C of FIG. 1C which represents a generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$ where N₁ = N₂ = 32; d = 2; and s = 2. Block 1_2 in 100L illustrates both the intra-block and

20 inter-block links. The layout diagram 100L corresponds to the embodiment where the switches that are placed together are implemented as two combined switches in the network 100B of FIG. 1B. As noted before then the network 100B is the generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with five stages as disclosed in U.S. Provisional Patent Application Serial No. 60/940,387 that 25 is incorporated by reference above.

That is the switches that are placed together in Block 1_2 as shown in FIG. 1L are namely the combined input and output switch IS1&OS1 belonging to switch 1, illustrated by dotted lines, (as noted before switch 1 is for illustration purposes only, in practice the switch implemented is combined input and output switch IS1&OS1); middle switch

MS(1,1) and MS(1,17) belonging to switch 2; middle switch MS(2,1) and MS(2,17) belonging to switch 3; middle switch MS(3,1) and MS(3,17) belonging to switch 4; And middle switch MS(4,1) belonging to switch 5.

Combined input and output switch IS1&OS1 is implemented as six by six switch
with the inlet links IL1, IL2, ML(8,1), ML(8,2), ML(8,7) and ML(8,8) being the inputs and middle links ML(1,1) – ML(1,4) and outlet links OL1 – OL2 being the outputs.

Middle switch MS(1,1) is implemented as four by four switch with middle links ML(1,1), ML(1,7), ML(7,1) and ML(7,11) being the inputs and middle links ML(2,1), ML(2,3), ML(8,1) and ML(8,3) being the outputs; And middle switch MS(1,17) is

- 10 implemented as four by four switch with the middle links ML(1,2), ML(1,8), ML(7,2) and ML(7,12) being the inputs and middle links ML(2,2), ML(2,4), ML(8,2) and ML(8,4) being the outputs. Similarly all the other middle switches are also implemented as two four by four switches as illustrated in 100L of FIG. 1L. Applicant observes that in middle switch MS(1,1) any one of the right going middle links can be switched to any one of the
- 15 left going middle links and hereinafter middle switch MS(1,1) provides U-turn links. In general, in the network $V_{bft}(N_1, N_2, d, s)$ each input switch, each output switch and each middle switch provides U-turn links.

In another embodiment, middle switch MS(1,1) (or the middle switches in any of the middle stage excepting the root middle stage) of Block 1_2 of $V_{bit}(N_1, N_2, d, s)$ can

- be implemented as a two by four switch and a two by two switch to save cross points.
 This is because the left going middle links of these middle switches are never setup to the right going middle links. For example, in middle switch MS(1,1) of Block 1_2 as shown FIG. 1L, the left going middle links namely ML(7,1) and ML(7,11) are never switched to the right going middle links ML(2,1) and ML(2,3). And hence to implement MS(1,1)
- two switches namely: 1) a two by four switch with the middle links ML(1,1) and ML(1,7) as inputs and the middle links ML(2,1), ML(2,3), ML(8,1), and ML(8,3) as outputs and 2) a two by two switch with the middle links ML(7,1) and ML(7,11) as inputs and the middle links ML(8,1) and ML(8,3) as outputs are sufficient without loosing any

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connectivity of the embodiment of MS(1,1) being implemented as an eight by eight switch as described before.)

Generalized Butterfly Fat Tree Network Embodiment with S = 1:

- 5 In one embodiment, in the network 100B of FIG. 1B (where it is implemented with s = 1), the switches that are placed together are implemented as a combined switch in input stage 110 and output stage 120; and as a combined switch in all the middle stages, then the network 100B is the generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 1 with five stages as disclosed in
- 10 U.S. Provisional Patent Application Serial No. 60/940,387 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a four by four switch. For example the switch input switch IS1 and output switch OS1 are placed together; so input and output switch IS1&OS1 is implemented as four by four switch with the inlet links IL1, IL2, ML(8,1)
- and ML(8,3) being the inputs and middle links ML(1,1) ML(1,2) and outlet links OL1
 OL2 being the outputs

The switches, corresponding to the middle stages that are placed together are implemented as a four by four switch. For example middle switches MS(1,1) is implemented as four by four switch with middle links ML(1,1), ML(1,3), ML(7,1) and

20 ML(7,5) being the inputs and middle links ML(2,1), ML(2,2), ML(8,1) and ML(8,2) being the outputs..

Layout diagrams 100C in FIG. 1C, 100D in FIG. 1D, 100E in FIG. 1E, 100F in FIG. 1G are also applicable to generalized butterfly fat tree network

 $V_{bft}(N_1, N_2, d, s)$ where N₁ = N₂ = 32; d = 2; and s = 1 with five stages. The layout 100C in FIG. 1C can be recursively extended for any arbitrarily large generalized butterfly fat

tree network $V_{bfi}(N_1, N_2, d, s)$. Accordingly layout 100H of FIG. 1H is also applicable to generalized butterfly fat tree network $V_{bfi}(N_1, N_2, d, s)$.

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Referring to diagram 100L1 of FIG. 1L1 illustrates a high-level implementation of Block 1_2 (Each of the other blocks have similar implementation) for the layout 100C of FIG. 1C when s = 1 which represents a generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$ where N₁ = N₂ = 32; d = 2; and s = 1 (All the double links are replaced

- 5 by single links when s = 1). Block 1_2 in 100K1 illustrates both the intra-block and interblock links. The layout diagram 100L1 corresponds to the embodiment where the switches that are placed together are implemented as a combined switch in the network 100B of FIG. 1B when s = 1. As noted before then the network 100B is the generalized butterfly fat tree network $V_{bft}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 1 with
- 10 nine stages as disclosed in U.S. Provisional Patent Application Serial No. 60/940,387 that is incorporated by reference above.

That is the switches that are placed together in Block 1_2 as shown in FIG. 1L1 are namely the input and output switch IS1&OS1 belonging to switch 1, illustrated by dotted lines, (as noted before switch 1 is for illustration purposes only, in practice the switches implemented are input switch IS1 and output switch OS1); middle switch MS(1,1) belonging to switch 2; middle switch MS(2,1) belonging to switch 3; middle switch MS(3,1) belonging to switch 4; And middle switch MS(4,1) belonging to switch 5.

Input and output switch IS1&OS1 are placed together; so input and output switch IS1&OS1 is implemented as four by four switch with the inlet links IL1, IL2, ML(8,1)
and ML(8,3) being the inputs and middle links ML(1,1) – ML(1,2) and outlet links OL1 – OL2 being the outputs.

Middle switch MS(1,1) is implemented as four by four switch with middle links ML(1,1), ML(1,3), ML(7,1) and ML(7,5) being the inputs and middle links ML(2,1), ML(2,2), ML(8,1) and ML(8,2) being the outputs. Similarly all the other middle switches are also implemented as four by four switches as illustrated in 100L1 of FIG. 1L1.

In another embodiment, middle switch MS(1,1) (or the middle switches in any of the middle stage excepting the root middle stage) of Block 1_2 of $V_{mlink-bft}(N_1, N_2, d, s)$ can be implemented as a two by four switch and a two by two

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switch to save cross points. This is because the left going middle links of these middle switches are never setup to the right going middle links. For example, in middle switch MS(1,1) of Block 1_2 as shown FIG. 1L1, the left going middle links namely ML(7,1) and ML(7,5) are never switched to the right going middle links ML(2,1) and ML(2,2).

- And hence to implement MS(1,1) two switches namely: 1) a two by four switch with the middle links ML(1,1) and ML(1,3) as inputs and the middle links ML(2,1), ML(2,2), ML(8,1), and ML(8,2) as outputs and 2) a two by two switch with the middle links ML(7,1) and ML(7,5) as inputs and the middle links ML(8,1) and ML(8,2) as outputs are sufficient without loosing any connectivity of the embodiment of MS(1,1) being
- 10 implemented as an eight by eight switch as described before.)

Hypercube-like Topology layout schemes:

Referring to diagram 300A in FIG. 3A, in one embodiment, an exemplary generalized multi-link multi-stage network $V_{mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d =

2; and s = 2 with nine stages of one hundred and forty four switches for satisfying communication requests, such as setting up a telephone call or a data call, or a connection between configurable logic blocks, between an input stage 110 and output stage 120 via middle stages 130, 140, 150, 170, 170, 180 and 190 is shown where input stage 110 consists of sixteen, two by four switches IS1-IS16 and output stage 120 consists of sixteen, four by two switches OS1-OS16.

As disclosed in U.S. Provisional Patent Application Serial No. 60/940,389 that is incorporated by reference above, such a network can be operated in rearrangeably nonblocking manner for arbitrary fan-out multicast connections and also can be operated in strictly non-blocking manner for unicast connections.

25 The diagram 300A in FIG. 3A is exactly the same as the diagram 100A in FIG. 1A excepting the connection links between middle stage 150 and middle stage 160 as well as between middle stage 160 and middle stage 170.

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Each of the $\frac{N}{d}$ middle switches are connected to exactly d switches in middle stage 160 through two links each for a total of $2 \times d$ links (for example the links ML(4,1) and ML(4,2) are connected from middle switch MS(3,1) to middle switch MS(4,1), and the links ML(4,3) and ML(4,4) are connected from middle switch MS(3,1) to middle switch MS(4,15)).

Each of the
$$\frac{N}{d}$$
 middle switches MS(4,1) – MS(4,16) in the middle stage 160 are

connected from exactly d input switches through two links each for a total of $2 \times d$ links (for example the links ML(4,1) and ML(4,2) are connected to the middle switch MS(4,1) from input switch MS(3,1), and the links ML(4,59) and ML(4,60) are connected to the

10 middle switch MS(4,1) from input switch MS(3,15)) and also are connected to exactly d switches in middle stage 170 through two links each for a total of $2 \times d$ links (for example the links ML(5,1) and ML(5,2) are connected from middle switch MS(4,1) to middle switch MS(5,1), and the links ML(5,3) and ML(5,4) are connected from middle switch MS(4,1) to middle switch MS(5,15)).

Each of the $\frac{N}{d}$ middle switches MS(5,1) – MS(5,16) in the middle stage 170 are connected from exactly *d* input switches through two links each for a total of 2×*d* links (for example the links ML(5,1) and ML(5,2) are connected to the middle switch MS(5,1) from input switch MS(4,1), and the links ML(5,59) and ML(5,60) are connected to the middle switch MS(5,1) from input switch MS(4,15)).

- 20 Finally the connection topology of the network 100A shown in FIG. 1A is also basically back to back inverse Benes connection topology but with a slight variation. All the cross middle links from middle switches MS(3,1) MS(3,8) connect to middle switches MS(4,9) MS(4,16) and all the cross middle links from middle switches MS(3,9) MS(3,16) connect to middle switches MS(4,1) MS(4,8). Applicant makes a
- 25 key observation that there are many combinations of connections possible using this property. The difference in the connection topology between diagram 100A of FIG. 1A and diagram 300A of FIG. 3A is that the connections formed by cross middle links

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between middle stage 150 and middle stage 160 are made of two different combinations otherwise both the diagrams 100A and 300A implement back to back inverse Benes connection topology. Since these networks implement back to back inverse Benes topologies since there is difference in the connections of cross middle links between

5 middle stage 150 and middle stage 160, the same difference in the connections of cross middle links between 160 and middle stage 170 occurs.

Referring to diagram 300B in FIG. 3B, is a folded version of the multi-link multistage network 300A shown in FIG. 3A. The network 300B in FIG. 3B shows input stage 110 and output stage 120 are placed together. That is input switch IS1 and output switch OS1 are placed together, input switch IS2 and output switch OS2 are placed together, and similarly input switch IS16 and output switch OS16 are placed together. All the right going middle links {i.e., inlet links IL1 – IL32 and middle links ML(1,1) - ML(1,64)} correspond to input switches IS1 - IS16, and all the left going middle links {i.e., middle links ML(7,1) - ML(7,64) and outlet links OL1-OL32} correspond to output switches OS1 - OS16.

Just the same way there is difference in the connection topology between diagram 100A of FIG. 1A and diagram 300A of FIG. 3A in the way the connections are formed by cross middle links between middle stage 150 and middle stage 160 and also between middle stage 160 and middle stage 170, the exact similar difference is there between the

20 diagram 100B of FIG. 1B and the diagram 300B of FIG. 3B, i.e., in the way the connections are formed by cross middle links between middle stage 150 and middle stage 160 and also between middle stage 160 and middle stage 170.

In one embodiment, in the network 300B of FIG. 3B, the switches that are placed together are implemented as separate switches then the network 300B is the generalized 25 folded multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 32$; d = 2; and s = 2 with nine stages as disclosed in U.S. Provisional Patent Application Serial No. 60/940,389 that is incorporated by reference above. That is the switches that are placed together in input stage 110 and output stage 120 are implemented as a two by four switch and a four by two switch. For example the switch input switch IS1 and output switch OS1

30 are placed together; so input switch IS1 is implemented as two by four switch with the

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inlet links IL1 and IL2 being the inputs of the input switch IS1 and middle links ML(1,1) – ML(1,4) being the outputs of the input switch IS1; and output switch OS1 is implemented as four by two switch with the middle links ML(8,1), ML(8,2), ML(8,7) and ML(8,8) being the inputs of the output switch OS1 and outlet links OL1 – OL2 being the outputs of the output switch OS1. Similarly in this embodiment of network 300B all the

switches that are placed together are implemented as separate switches.

Referring to layout 300C of FIG. 3C, in one embodiment, there are sixteen blocks namely Block 1_2, Block 3_4, Block 5_6, Block 7_8, Block 9_10, Block 11_12, Block 13_14, Block 15_16, Block 17_18, Block 19_20, Block 21_22, Block 23_24, Block

10 25_26, Block 27_28, Block 29_30, and Block 31_32. Each block implements all the switches in one row of the network 300B of FIG. 3B, one of the key aspects of the current invention. For example Block 1_2 implements the input switch IS1, output Switch OS1, middle switch MS(1,1), middle switch MS(7,1), middle switch MS(2,1), middle switch MS(6,1), middle switch MS(3,1), middle switch MS(5,1), and middle switch MS(4,1).

15 For the simplification of illustration, Input switch IS1 and output switch OS1 together are denoted as switch 1; Middle switch MS(1,1) and middle switch MS(7,1) together are denoted by switch 2; Middle switch MS(2,1) and middle switch MS(6,1) together are denoted by switch 3; Middle switch MS(3,1) and middle switch MS(5,1) together are denoted by switch 4; And middle switch MS(4,1) is denoted by switch 5.

All the straight middle links are illustrated in layout 300C of FIG. 3C. For example in Block 1_2, inlet links IL1 – IL2, outlet links OL1 – OL2, middle link ML(1,1), middle link ML(1,2), middle link ML(8,1), middle link ML(8,2), middle link ML(2,1), middle link ML(2,2), middle link ML(7,1), middle link ML(7,2), middle link ML(3,1), middle link ML(3,2), middle link ML(6,1), middle link ML(6,2), middle link
ML(4,1), middle link ML(4,2), middle link ML(5,1) and middle link ML(5,2) are illustrated in layout 300C of FIG. 3C.

Even though it is not illustrated in layout 300C of FIG. 3C, in each block, in addition to the switches there may be Configurable Logic Blocks (CLB) or any arbitrary digital circuit or sub-integrated circuit block depending on the applications in different embodiments. There are four quadrants in the layout 300C of FIG. 3C namely top-left,

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bottom-left, top-right and bottom-right quadrants. Top-left quadrant implements Block 1_2, Block 3_4, Block 5_6, and Block 7_8. Bottom-left quadrant implements Block 9_10, Block 11_12, Block 13_14, and Block 15_16. Top-right quadrant implements Block 25_26, Block 27_28, Block 29_30, and Block 31_32. Bottom-right quadrant

- 5 implements Block 17_18, Block 19_20, Block 21_22, and Block 23_24. There are two halves in layout 300C of FIG. 3C namely left-half and right-half. Left-half consists of top-left and bottom-left quadrants. Right-half consists of top-right and bottom-right quadrants.
- Recursively in each quadrant there are four sub-quadrants. For example in top-left quadrant there are four sub-quadrants namely top-left sub-quadrant, bottom-left subquadrant, top-right sub-quadrant and bottom-right sub-quadrant. Top-left sub-quadrant of top-left quadrant implements Block 1_2. Bottom-left sub-quadrant of top-left quadrant implements Block 3_4. Top-right sub-quadrant of top-left quadrant implements Block 7 8. Finally bottom-right sub-quadrant of top-left quadrant implements Block 5_6.
- 15 Similarly there are two sub-halves in each quadrant. For example in top-left quadrant there are two sub-halves namely left-sub-half and right-sub-half. Left-sub-half of top-left quadrant implements Block 1_2 and Block 3_4. Right-sub-half of top-left quadrant implements Block 7_8 and Block 5_6. Recursively in larger multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where N₁ = N₂ > 32, the layout in this embodiment in accordance
- 20 with the current invention, will be such that the super-quadrants will also be arranged in a similar manner.

Layout 300D of FIG. 3D illustrates the inter-block links (in the layout 300C of FIG. 3C all the cross middle links are inter-block links) between switches 1 and 2 of each block. For example middle links ML(1,3), ML(1,4), ML(8,7), and ML(8,8) are connected
between switch 1 of Block 1_2 and switch 2 of Block 3_4. Similarly middle links ML(1,7), ML(1,8), ML(8,3), and ML(8,4) are connected between switch 2 of Block 1_2 and switch 1 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 100D of FIG. 1D can be implemented as vertical tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for example

30 middle links ML(1,4) and ML(8,8) are implemented as two different tracks); or in an

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alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track).

Layout 300E of FIG. 3E illustrates the inter-block links between switches 2 and 3
of each block. For example middle links ML(2,3), ML(2,4), ML(7,11), and ML(7,12) are connected between switch 2 of Block 1_2 and switch 3 of Block 3_4. Similarly middle links ML(2,11), ML(2,12), ML(7,3), and ML(7,4) are connected between switch 3 of Block 1_2 and switch 2 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 300E of FIG. 3E can be implemented as diagonal tracks in one embodiment.

10 Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(2,12) and ML(7,4) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(2,12) and ML(7,4) are implemented as a time division multiplexed single track).

15 Layout 300F of FIG. 3F illustrates the inter-block links between switches 3 and 4 of each block. For example middle links ML(3,3), ML(3,4), ML(6,19), and ML(6,20) are connected between switch 3 of Block 1_2 and switch 4 of Block 3_4. Similarly middle links ML(3,19), ML(3,20), ML(6,3), and ML(6,4) are connected between switch 4 of Block 1_2 and switch 3 of Block 3_4. Applicant notes that the inter-block links illustrated

20 in layout 300F of FIG. 3F can be implemented as vertical tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(3,4) and ML(6,20) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track.

Layout 300G of FIG. 3G illustrates the inter-block links between switches 4 and 5 of each block. For example middle links ML(4,3), ML(4,4), ML(5,35), and ML(5,36) are connected between switch 4 of Block 1_2 and switch 5 of Block 3_4. Similarly middle links ML(4,35), ML(4,36), ML(5,3), and ML(5,4) are connected between switch 5 of

30 Block 1_2 and switch 4 of Block 3_4. Applicant notes that the inter-block links illustrated

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in layout 300G of FIG. 3G can be implemented as horizontal tracks in one embodiment. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(4,4) and ML(5,36) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(4,4) and ML(5,36) are

implemented as a time division multiplexed single track).

The complete layout for the network 300B of FIG. 3B is given by combining the links in layout diagrams of 300C, 300D, 300E, 300F, and 300G. Applicant notes that in the layout 300C of FIG. 3C, the inter-block links between switch 1 and switch 2 are vertical tracks as shown in layout 300D of FIG. 3D; the inter-block links between switch 2 and switch 3 are horizontal tracks as shown in layout 300E of FIG. 3E; the inter-block links between switch 3 and switch 4 are vertical tracks as shown in layout 300F of FIG. 3F; and finally the inter-block links between switch 4 and switch 5 are horizontal tracks

15 tracks or diagonal tracks. It continues recursively for larger networks of N > 32 as will be illustrated later.

as shown in layout 300G of FIG. 3G. The pattern is either vertical tracks, horizontal

Some of the key aspects of the current invention related to layout diagram 300C of IFG. 3C are noted. 1) All the switches in one row of the multi-stage network 300B are implemented in a single block. 2) The blocks are placed in such a way that all the inter-

20 block links are either horizontal tracks, vertical tracks or diagonal tracks; 3) The length of the longest wire is about half of the width (or length) of the complete layout (For example middle link ML(4,4) is about half the width of the complete layout.);

The layout 300C in FIG. 3C can be recursively extended for any arbitrarily large generalized folded multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$. Referring to 25 layout 300H of FIG. 3H, illustrates the extension of layout 300C for the network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 = 128$; d = 2; and s = 2. There are four superquadrants in layout 300H namely top-left super-quadrant, bottom-left super-quadrant, top-right super-quadrant, bottom-right super-quadrant. Total number of blocks in the layout 300H is sixty four. Top-left super-quadrant implements the blocks from block 1_2

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to block 31_32. Each block in all the super-quadrants has two more switches namely switch 6 and switch 7 in addition to the switches [1-5] illustrated in layout 300C of FIG. 3C. The inter-block link connection topology is the exactly the same between the switches 1 and 2; switches 2 and 3; switches 3 and 4; switches 4 and 5 as it is shown in the layouts of FIG. 3D, FIG. 3E, FIG. 3F, and FIG. 3G respectively.

Bottom-left super-quadrant implements the blocks from block 33_34 to block 63_64. Top-right super-quadrant implements the blocks from block 65_66 to block 95_96. And bottom-right super-quadrant implements the blocks from block 97_98 to block 127_128. In all these three super-quadrants also, the inter-block link connection topology is the exactly the same between the switches 1 and 2; switches 2 and 3; switches

topology is the exactly the same between the switches 1 and 2; switches 2 and 3; switches3 and 4; switches 4 and 5 as that of the top-left super-quadrant.

Recursively in accordance with the current invention, the inter-block links connecting the switch 5 and switch 6 will be vertical tracks between the corresponding switches of top-left super-quadrant and bottom-left super-quadrant. And similarly the

- 15 inter-block links connecting the switch 5 and switch 6 will be vertical tracks between the corresponding switches of top-right super-quadrant and bottom-right super-quadrant. The inter-block links connecting the switch 6 and switch 7 will be horizontal tracks between the corresponding switches of top-left super-quadrant and top-right super-quadrant. And similarly the inter-block links connecting the switch 6 and switch 7 will be horizontal
- 20 tracks between the corresponding switches of bottom-left super-quadrant and bottomright super-quadrant.

Ring Topology layout schemes:

Layout diagram 400C of FIG. 4C is another embodiment for the generalized
folded multi-link multi-stage network V_{fold-mlink} (N₁, N₂, d, s) diagram 100B in FIG. 1B.
Referring to layout 400C of FIG. 4C, there are sixteen blocks namely Block 1_2,
Block 3_4, Block 5_6, Block 7_8, Block 9_10, Block 11_12, Block 13_14, Block 15_16,
Block 17_18, Block 19_20, Block 21_22, Block 23_24, Block 25_26, Block 27_28,

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Block 29_30, and Block 31_32. Each block implements all the switches in one row of the network 100B of FIG. 1B, one of the key aspects of the current invention. For example Block 1_2 implements the input switch IS1, output Switch OS1, middle switch MS(1,1), middle switch MS(7,1), middle switch MS(2,1), middle switch MS(6,1), middle switch

- MS(3,1), middle switch MS(5,1), and middle switch MS(4,1). For the simplification of illustration, Input switch IS1 and output switch OS1 together are denoted as switch 1; Middle switch MS(1,1) and middle switch MS(7,1) together are denoted by switch 2; Middle switch MS(2,1) and middle switch MS(6,1) together are denoted by switch 3; Middle switch MS(3,1) and middle switch MS(5,1) together are denoted by switch 4;
- 10 And middle switch MS(4,1) is denoted by switch 5.

All the straight middle links are illustrated in layout 400C of FIG. 4C. For example in Block 1_2, inlet links IL1 – IL2, outlet links OL1 – OL2, middle link ML(1,1), middle link ML(1,2), middle link ML(8,1), middle link ML(8,2), middle link ML(2,1), middle link ML(2,2), middle link ML(7,1), middle link ML(7,2), middle link

- ML(3,1), middle link ML(3,2), middle link ML(6,1), middle link ML(6,2), middle link ML(4,1), middle link ML(4,2), middle link ML(5,1) and middle link ML(5,2) are illustrated in layout 400C of FIG. 4C.
- Even though it is not illustrated in layout 400C of FIG. 4C, in each block, in addition to the switches there may be Configurable Logic Blocks (CLB) or any arbitrary digital circuit or sub-integrated circuit block depending on the applications in different embodiments. The topology of the layout 400C in FIG. 4C is a ring. For each of the neighboring rows in diagram 100B of FIG. 1B the corresponding blocks are also physically neighbors in layout diagram 400C of FIG. 4C. In addition the topmost row is also logically considered as neighbor to the bottommost row. For example Block 1_2
- (implementing the switches belonging to a row in diagram 100B of FIG. 1B) has Block
 3_4 as neighbor since Block 3_4 implements the switches in its neighboring row.
 Similarly Block 1_2 also has Block 31_32 as neighbor since Block 1_2 implements
 topmost row of switches and Block 31_32 implements bottommost row of switches in
 diagram 100B of FIG. 1B. The ring layout scheme illustrated in 400C of FIG. 4C can be

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generalized for a large multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 > 32$, in accordance with the current invention.

Layout 400B of FIG. 4B illustrates the inter-block links (in the layout 400A of FIG. 4A all the cross middle links are inter-block links) between switches 1 and 2 of each block. For example middle links ML(1,3), ML(1,4), ML(8,7), and ML(8,8) are connected between switch 1 of Block 1_2 and switch 2 of Block 3_4. Similarly middle links ML(1,7), ML(1,8), ML(8,3), and ML(8,4) are connected between switch 2 of Block 1_2 and switch 1 of Block 3_4. Applicant notes that the inter-block links illustrated in layout

10 tracks. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(1,4) and ML(8,8) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(1,4) and ML(8,8) are implemented as a time division multiplexed single track).

400B of FIG. 4B are implemented as vertical tracks or horizontal tracks or diagonal

- 15 Layout 400C of FIG. 4C illustrates the inter-block links between switches 2 and 3 of each block. For example middle links ML(2,3), ML(2,4), ML(7,11), and ML(7,12) are connected between switch 2 of Block 1_2 and switch 3 of Block 3_4. Similarly middle links ML(2,11), ML(2,12), ML(7,3), and ML(7,4) are connected between switch 3 of Block 1_2 and switch 2 of Block 3_4. Applicant notes that the inter-block links illustrated
- in layout 400C of FIG. 4C are implemented as vertical tracks or horizontal tracks or diagonal tracks. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(2,12) and ML(7,4) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(2,12) and ML(2,12) and ML(7,4)
 are implemented as a time division multiplexed single track).

Layout 400D of FIG. 4D illustrates the inter-block links between switches 3 and 4 of each block. For example middle links ML(3,3), ML(3,4), ML(6,19), and ML(6,20) are connected between switch 3 of Block 1_2 and switch 4 of Block 3_4. Similarly middle links ML(3,19), ML(3,20), ML(6,3), and ML(6,4) are connected between switch 4 of

30 Block 1_2 and switch 3 of Block 3_4. Applicant notes that the inter-block links illustrated

in layout 400D of FIG. 4D are implemented as vertical tracks or horizontal tracks or diagonal tracks. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(3,4) and ML(6,20) are implemented as two different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(3,4) and ML(6,20)

5 time division multiplexed single track (for example middle links ML(3,4) and ML(6,20) are implemented as a time division multiplexed single track).

Layout 400E of FIG. 4E illustrates the inter-block links between switches 4 and 5 of each block. For example middle links ML(4,3), ML(4,4), ML(5,35), and ML(5,36) are connected between switch 4 of Block 1_2 and switch 5 of Block 3_4. Similarly middle

10 links ML(4,35), ML(4,36), ML(5,3), and ML(5,4) are connected between switch 5 of Block 1_2 and switch 4 of Block 3_4. Applicant notes that the inter-block links illustrated in layout 400E of FIG. 4E are implemented as vertical tracks or horizontal tracks or diagonal tracks. Also in one embodiment inter-block links are implemented as two different tracks (for example middle links ML(4,4) and ML(5,36) are implemented as two

15 different tracks); or in an alternative embodiment inter-block links are implemented as a time division multiplexed single track (for example middle links ML(4,4) and ML(5,36) are implemented as a time division multiplexed single track).

The complete layout for the network 100B of FIG. 1B is given by combining the links in layout diagrams of 400A, 400B, 400C, 400D, and 400E.

20 Some of the key aspects of the current invention related to layout diagram 400A of FIG. 4A are noted. 1) All the switches in one row of the multi-stage network 100B are implemented in a single block. 2) The blocks are placed in such a way that all the interblock links are either horizontal tracks, vertical tracks or diagonal tracks; 3) Length of the different wires between the same two middle stages is not the same. However it gives an opportunity to implement the most connected circuits to place and route through the blocks which have shorter wires.

Layout diagram 400C1 of FIG. 4C1 is another embodiment for the generalized folded multi-link multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ diagram 100B in FIG. 1B. Referring to layout 400C1 of FIG. 4C1, there are sixteen blocks namely Block 1_2, Block

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3_4, Block 5_6, Block 7_8, Block 9_10, Block 11_12, Block 13_14, Block 15_16, Block 17_18, Block 19_20, Block 21_22, Block 23_24, Block 25_26, Block 27_28, Block 29_30, and Block 31_32. Each block implements all the switches in one row of the network 100B of FIG. 1B, one of the key aspects of the current invention. For example

- Block 1_2 implements the input switch IS1, output Switch OS1, middle switch MS(1,1), middle switch MS(7,1), middle switch MS(2,1), middle switch MS(6,1), middle switch MS(3,1), middle switch MS(5,1), and middle switch MS(4,1). For the simplification of illustration, Input switch IS1 and output switch OS1 together are denoted as switch 1; Middle switch MS(1,1) and middle switch MS(7,1) together are denoted by switch 2;
- Middle switch MS(2,1) and middle switch MS(6,1) together are denoted by switch 3;
 Middle switch MS(3,1) and middle switch MS(5,1) together are denoted by switch 4;
 And middle switch MS(4,1) is denoted by switch 5.

All the straight middle links are illustrated in layout 400C1 of FIG. 4C1. For example in Block 1 2, inlet links IL1 – IL2, outlet links OL1 – OL2, middle link

- ML(1,1), middle link ML(1,2), middle link ML(8,1), middle link ML(8,2), middle link ML(2,1), middle link ML(2,2), middle link ML(7,1), middle link ML(7,2), middle link ML(3,1), middle link ML(3,2), middle link ML(6,1), middle link ML(6,2), middle link ML(4,1), middle link ML(4,2), middle link ML(5,1) and middle link ML(5,2) are illustrated in layout 400C1 of FIG. 4C1.
- Even though it is not illustrated in layout 400C1 of FIG. 4C1, in each block, in addition to the switches there may be Configurable Logic Blocks (CLB) or any arbitrary digital circuit or sub-integrated circuit block depending on the applications in different embodiments. The topology of the layout 400C1 in FIG. 4C1 is another embodiment of ring layout topology. For each of the neighboring rows in diagram 100B of FIG. 1B the
 corresponding blocks are also physically neighbors in layout diagram 400C of FIG. 4C. In addition the topmost row is also logically considered as neighbor to the bottommost row. For example Block 1_2 (implementing the switches belonging to a row in diagram 100B of FIG. 1B) has Block 3_4 as neighbor since Block 3_4 implements the switches in its neighboring row. Similarly Block 1_2 also has Block 31_32 as neighbor since Block
- 30 1_2 implements topmost row of switches and Block 31_32 implements bottommost row

of switches in diagram 100B of FIG. 1B. The ring layout scheme illustrated in 400C of FIG. 4C can be generalized for a large multi-stage network $V_{fold-mlink}(N_1, N_2, d, s)$ where $N_1 = N_2 > 32$, in accordance with the current invention.

- All the layout embodiments disclosed in the current invention are applicable to 5 generalized multi-stage networks $V(N_1, N_2, d, s)$, generalized folded multi-stage networks $V_{fold}(N_1, N_2, d, s)$, generalized butterfly fat tree networks $V_{bft}(N_1, N_2, d, s)$, generalized multi-link multi-stage networks $V_{mlink}(N_1, N_2, d, s)$, generalized folded multilink multi-stage networks $V_{fold-mlink}(N_1, N_2, d, s)$, generalized multi-link butterfly fat tree networks $V_{mlink-bft}(N_1, N_2, d, s)$, and generalized hypercube networks $V_{hcube}(N_1, N_2, d, s)$
- 10 for s = 1,2,3 or any number in general, and for both $N_1 = N_2 = N$ and $N_1 \neq N_2$, and d is any integer.

Conversely applicant makes another important observation that generalized hypercube networks $V_{hcube}(N_1, N_2, d, s)$ are implemented with the layout topology being the hypercube topology shown in layout 100C of FIG. 1C with large scale cross point reduction as any one of the networks described in the current invention namely: generalized multi-stage networks $V(N_1, N_2, d, s)$, generalized folded multi-stage networks $V_{fold}(N_1, N_2, d, s)$, generalized butterfly fat tree networks $V_{bft}(N_1, N_2, d, s)$, generalized multi-link multi-stage networks $V_{mlink}(N_1, N_2, d, s)$, generalized folded multilink multi-stage networks $V_{fold-mlink}(N_1, N_2, d, s)$, generalized multi-link butterfly fat tree networks $V_{mlink-bft}(N_1, N_2, d, s)$ for s = 1,2,3 or any number in general, and for both $N_1 = N_2 = N$ and $N_1 \neq N_2$, and d is any integer.

Applications Embodiments:

All the embodiments disclosed in the current invention are useful in many varieties of applications. FIG. 5A1 illustrates the diagram of 500A1 which is a typical two by two switch with two inlet links namely IL1 and IL2, and two outlet links namely

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OL1 and OL2. The two by two switch also implements four crosspoints namely CP(1,1), CP(1,2), CP(2,1) and CP(2,2) as illustrated in FIG. 5A1. For example the diagram of 500A1 may the implementation of middle switch MS(1,1) of the diagram 100K of FIG. 1K where inlet link IL1 of diagram 500A1 corresponds to middle link ML(1,1) of

5 diagram 100K, inlet link IL2 of diagram 500A1 corresponds to middle link ML(1,7) of diagram 100K, outlet link OL1 of diagram 500A1 corresponds to middle link ML(2,1) of diagram 100K, outlet link OL2 of diagram 500A1 corresponds to middle link ML(2,3) of diagram 100K.

10 1) **Programmable Integrated Circuit Embodiments:**

All the embodiments disclosed in the current invention are useful in programmable integrated circuit applications. FIG. 5A2 illustrates the detailed diagram 500A2 for the implementation of the diagram 500A1 in programmable integrated circuit embodiments. Each crosspoint is implemented by a transistor coupled between the corresponding inlet link and outlet link, and a programmable cell in programmable

- 15 corresponding inlet link and outlet link, and a programmable cell in programmable integrated circuit embodiments. Specifically crosspoint CP(1,1) is implemented by transistor C(1,1) coupled between inlet link IL1 and outlet link OL1, and programmable cell P(1,1); crosspoint CP(1,2) is implemented by transistor C(1,2) coupled between inlet link IL1 and outlet link OL2, and programmable cell P(1,2); crosspoint CP(2,1) is
- 20 implemented by transistor C(2,1) coupled between inlet link IL2 and outlet link OL1, and programmable cell P(2,1); and crosspoint CP(2,2) is implemented by transistor C(2,2) coupled between inlet link IL2 and outlet link OL2, and programmable cell P(2,2).

If the programmable cell is programmed ON, the corresponding transistor couples the corresponding inlet link and outlet link. If the programmable cell is programmed

25 OFF, the corresponding inlet link and outlet link are not connected. For example if the programmable cell P(1,1) is programmed ON, the corresponding transistor C(1,1) couples the corresponding inlet link IL1 and outlet link OL1. If the programmable cell P(1,1) is programmed OFF, the corresponding inlet link IL1 and outlet link OL1 are not connected. In volatile programmable integrated circuit embodiments the programmable

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cell may be an SRAM (Static Random Address Memory) cell. In non-volatile programmable integrated circuit embodiments the programmable cell may be a Flash memory cell. Also the programmable integrated circuit embodiments may implement field programmable logic arrays (FPGA) devices, or programmable Logic devices (PLD),

5 or Application Specific Integrated Circuits (ASIC) embedded with programmable logic circuits or 3D-FPGAs.

FIG. 5A2 also illustrates a buffer B1 on inlet link IL2. The signals driven along inlet link IL2 are amplified by buffer B1. Buffer B1 can be inverting or non-inverting buffer. Buffers such as B1 are used to amplify the signal in links which are usually long.

10 2) One-time Programmable Integrated Circuit Embodiments:

All the embodiments disclosed in the current invention are useful in one-time programmable integrated circuit applications. FIG. 5A3 illustrates the detailed diagram 500A3 for the implementation of the diagram 500A1 in one-time programmable integrated circuit embodiments. Each crosspoint is implemented by a via coupled

- 15 between the corresponding inlet link and outlet link in one-time programmable integrated circuit embodiments. Specifically crosspoint CP(1,1) is implemented by via V(1,1) coupled between inlet link IL1 and outlet link OL1; crosspoint CP(1,2) is implemented by via V(1,2) coupled between inlet link IL1 and outlet link OL2; crosspoint CP(2,1) is implemented by via V(2,1) coupled between inlet link IL2 and outlet link OL1; and
- 20 crosspoint CP(2,2) is implemented by via V(2,2) coupled between inlet link IL2 and outlet link OL2.

If the via is programmed ON, the corresponding inlet link and outlet link are permanently connected which is denoted by thick circle at the intersection of inlet link and outlet link. If the via is programmed OFF, the corresponding inlet link and outlet link 25 are not connected which is denoted by the absence of thick circle at the intersection of inlet link and outlet link. For example in the diagram 500A3 the via V(1,1) is programmed ON, and the corresponding inlet link IL1 and outlet link OL1 are connected as denoted by thick circle at the intersection of inlet link IL1 and outlet link OL1; the via V(2,2) is programmed ON, and the corresponding inlet link IL2 and outlet link OL2 are

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connected as denoted by thick circle at the intersection of inlet link IL2 and outlet link OL2; the via V(1,2) is programmed OFF, and the corresponding inlet link IL1 and outlet link OL2 are not connected as denoted by the absence of thick circle at the intersection of inlet link IL1 and outlet link OL2; the via V(2,1) is programmed OFF, and the

5 corresponding inlet link IL2 and outlet link OL1 are not connected as denoted by the absence of thick circle at the intersection of inlet link IL2 and outlet link OL1. One-time programmable integrated circuit embodiments may be anti-fuse based programmable integrated circuit devices or mask programmable structured ASIC devices.

3) Integrated Circuit Placement and Route Embodiments:

- 10 All the embodiments disclosed in the current invention are useful in Integrated Circuit Placement and Route applications, for example in ASIC backend Placement and Route tools. FIG. 5A4 illustrates the detailed diagram 500A4 for the implementation of the diagram 500A1 in Integrated Circuit Placement and Route embodiments. In an integrated circuit since the connections are known a-priori, the switch and crosspoints are
- 15 actually virtual. However the concept of virtual switch and virtal crosspoint using the embodiments disclosed in the current invention reduces the number of required wires, wire length needed to connect the inputs and outputs of different netlists and the time required by the tool for placement and route of netlists in the integrated circuit.
- Each virtual crosspoint is used to either to hardwire or provide no connectivity 20 between the corresponding inlet link and outlet link. Specifically crosspoint CP(1,1) is implemented by direct connect point DCP(1,1) to hardwire (i.e., to permanently connect) inlet link IL1 and outlet link OL1 which is denoted by the thick circle at the intersection of inlet link IL1 and outlet link OL1; crosspoint CP(2,2) is implemented by direct connect point DCP(2,2) to hardwire inlet link IL2 and outlet link OL2 which is denoted
- 25 by the thick circle at the intersection of inlet link IL2 and outlet link OL2. The diagram 500A4 does not show direct connect point DCP(1,2) and direct connect point DCP(1,3) since they are not needed and in the hardware implementation they are eliminated. Alternatively inlet link IL1 needs to be connected to outlet link OL1 and inlet link IL1 does not need to be connected to outlet link OL2. Also inlet link IL2 needs to be
- 30 connected to outlet link OL2 and inlet link IL2 does not need to be connected to outlet

link OL1. Furthermore in the example of the diagram 500A4, there is no need to drive the signal of inlet link IL1 horizontally beyond outlet link OL1 and hence the inlet link IL1 is not even extended horizontally until the outlet link OL2. Also the absence of direct connect point DCP(2,1) illustrates there is no need to connect inlet link IL2 and outlet link OL1.

5 link OL1.

In summary in integrated circuit placement and route tools, the concept of virtual switches and virtual cross points is used during the implementation of the placement & routing algorithmically in software, however during the hardware implementation cross points in the cross state are implemented as hardwired connections between the

10 corresponding inlet link and outlet link, and in the bar state are implemented as no connection between inlet link and outlet link.

3) More Application Embodiments:

All the embodiments disclosed in the current invention are also useful in the design of SoC interconnects, Field programmable interconnect chips, parallel computer systems and in time-space-time switches.

Numerous modifications and adaptations of the embodiments, implementations, and examples described herein will be apparent to the skilled artisan in view of the disclosure.

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<u>CLAIMS</u>

What is claimed is:

1. An integrated circuit device comprising a plurality of sub-integrated circuit blocks and a routing network, and

Said each plurality of sub-integrated circuit blocks comprising a plurality of inlet links and a plurality of outlet links; and

Said routing network interconnects any one of said outlet link of one of said subintegrated circuit block to one or more said inlet links of one or more of said subintegrated circuit blocks; and

10 Said routing network comprising of a plurality of stages *y*, starting from the lowest stage to the highest stage; and

Said routing network comprising a plurality of switches of size $d \times d$, where $d \ge 2$, in each said stage and each said switch of size $d \times d$ having d inlet links and d outlet links; and

15 Said each sub-integrated circuit block comprising a plurality of said switches corresponding to each said stage; and

Said each sub-integrated circuit block comprising a plurality of forward connecting links connecting from switches in lower stage to switches in the immediate succeeding higher stage, and also comprising a plurality of backward connecting links

20 connecting from switches in higher stage to switches in the immediate preceding lower stage; and

Said each sub-integrated circuit block comprising a plurality straight links in said forward connecting links from switches in lower stage to switches in the immediate succeeding higher stage and a plurality cross links in said forward connecting links from

25 switches in lower stage to switches in the immediate succeeding higher stage, and further comprising a plurality of straight links in said backward connecting links from switches

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in higher stage to switches in the immediate preceding lower stage and a plurality of cross links in said backward connecting links from switches in higher stage to switches in the immediate preceding lower stage.

The integrated circuit device of claim 1, wherein said all straight links are
 connecting from switches in each said sub-integrated circuit block are connecting to
 switches in the same said sub-integrated circuit block; and

said all cross links are connecting as either vertical or horizontal links between switches in two different said sub-integrated circuit blocks.

3. The integrated circuit device of claim 2, wherein said plurality of sub-10 integrated circuit blocks arranged in a two-dimensional grid.

4. The integrated circuit device of claim 3, wherein said cross links in succeeding stages are connecting as alternative vertical and horizontal links between switches in said sub-integrated circuit blocks.

- 5. The integrated circuit device of claim 4, wherein said cross links from 15 switches in a stage in one of said sub-integrated circuit blocks are connecting to switches in the succeeding stage in another of said sub-integrated circuit blocks so that said cross links are either vertical links or horizontal and vice versa, and hereinafter such cross links are "shuffle exchange links").
- 6. The integrated circuit device of claim 5, wherein said all horizontal shuffle 20 exchange links between switches in any two corresponding said succeeding stages are substantially of equal length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are substantially of equal length in the entire said integrated circuit device.
- 7. The integrated circuit device of claim 6, wherein the shortest horizontal 25 shuffle exchange links are connecting at the lowest stage and between switches in two

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nearest neighboring said sub-integrated circuit blocks, and length of the horizontal shuffle exchange links is doubled in each succeeding stage; and the shortest vertical shuffle exchange links are connecting at the lowest stage and between switches in two nearest neighboring said sub-integrated circuit blocks, and length of the vertical shuffle exchange links is doubled in each succeeding stage.

8. The integrated circuit device of claim 7, wherein $y \ge (\log_2 N)$ so that the length of the horizontal shuffle exchange links in the highest stage is equal to half the size of the horizontal size of said two dimensional grid of sub-integrated circuit blocks and the length of the vertical shuffle exchange links in the highest stage is equal to half the size of the vertical size of said two dimensional grid of sub-integrated circuit blocks.

9. The integrated circuit device of claim 8, wherein d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said subintegrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast Benes network with full bandwidth.

10. The integrated circuit device of claim 8, wherein d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast Benes network and rearrangeably

nonblocking for arbitrary fan-out multicast Benes network with full bandwidth.

11. The integrated circuit device of claim 8, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in

25 each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast Benes network with full bandwidth.

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12. The integrated circuit device of claim 7, wherein $y \ge (\log_2 N)$ so that the length of the horizontal shuffle exchange links in the highest stage is equal to half the size of the horizontal size of said two dimensional grid of sub-integrated circuit blocks and the length of the vertical shuffle exchange links in the highest stage is equal to half the size of

5 the vertical size of said two dimensional grid of sub-integrated circuit blocks, and said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks.

13. The integrated circuit device of claim 12, wherein d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said
10 forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast butterfly fat tree network with full bandwidth.

14. The integrated circuit device of claim 12, wherein d = 2 and there are at
15 least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast butterfly fat tree network and rearrangeably nonblocking for arbitrary fan-out multicast butterfly fat tree network with
20 full bandwidth.

15. The integrated circuit device of claim 12, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said

25 routing network is strictly nonblocking for arbitrary fan-out multicast butterfly fat tree network with full bandwidth.

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16. The integrated circuit device of claim 1, wherein said horizontal and vertical links are implemented on two or more metal layers.

17. The integrated circuit device of claim 1, wherein said switches comprising active and reprogrammable cross points and said each cross point is programmable by an
5 SRAM cell or a Flash Cell.

18. The integrated circuit device of claim 1, wherein said sub-integrated circuit blocks are of equal die size.

19. The integrated circuit device of claim 16, wherein said sub-integrated circuit blocks are Lookup Tables (hereinafter "LUTs") and said integrated circuit device
10 is a field programmable gate array (FPGA) device or field programmable gate array (FPGA) block embedded in another integrated circuit device.

20. The integrated circuit device of claim 16, wherein said sub-integrated circuit blocks are AND or OR gates and said integrated circuit device is a programmable logic device (PLD).

15 21. The integrated circuit device of claim 1, wherein said sub-integrated circuit blocks comprising any arbitrary hardware logic or memory circuits.

22. The integrated circuit device of claim 1, wherein said switches comprising active one-time programmable cross points and said integrated circuit device is a mask programmable gate array (MPGA) device or a structured ASIC device.

20 23. The integrated circuit device of claim 1, wherein said switches comprising passive cross points or just connection of two links or not and said integrated circuit device is a Application Specific Integrated Circuit (ASIC) device.

24. The integrated circuit device of claim 1, wherein said sub-integrated circuit blocks further recursively comprise one or more super-sub-integrated circuit
25 blocks and a sub-routing network.

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25. The integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and $y \ge (\log_2 N)$.

5 26. The integrated circuit device of claim 25, wherein *d* = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said subintegrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized multi-stage network with 10 full bandwidth.

27. The integrated circuit device of claim 25, wherein d = 2 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast generalized multi-stage network and rearrangeably nonblocking for arbitrary fan-out multicast generalized multi-stage network with full bandwidth.

28. The integrated circuit device of claim 25, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting
20 said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast generalized multistage network with full bandwidth.

29. The integrated circuit device of claim 5, wherein said all horizontal shuffle 25 exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and $y \ge (\log_2 N)$, and

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said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks.

- 30. The integrated circuit device of claim 29, wherein d = 2 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said
 5 forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized butterfly fat tree network with full bandwidth.
- 31. The integrated circuit device of claim 29, wherein d = 2 and there are at
 least two switches in each said stage in each said sub-integrated circuit block connecting
 said forward connecting links and there are at least two switches in each said stage in
 each said sub-integrated circuit block connecting said backward connecting links and said
 routing network is strictly nonblocking for unicast generalized butterfly fat tree Network
 and rearrangeably nonblocking for arbitrary fan-out multicast generalized butterfly fat
 tree network with full bandwidth.
- 32. The integrated circuit device of claim 29, wherein d = 2 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said
 routing network is strictly nonblocking for arbitrary fan-out multicast generalized butterfly fat tree network with full bandwidth.

33. The integrated circuit device of claim 1, wherein said straight links connecting from switches in each said sub-integrated circuit block are connecting to switches in the same said sub-integrated circuit block; and

25 said cross links are connecting as vertical or horizontal or diagonal links between two different said sub-integrated circuit blocks.

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34. The integrated circuit device of claim 8, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said subintegrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast multi-link Benes network with full bandwidth.

35. The integrated circuit device of claim 8, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in
10 each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast multi-link Benes network and rearrangeably nonblocking for arbitrary fan-out multicast multi-link Benes network with full bandwidth.

36. The integrated circuit device of claim 8, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast multi-link Benes network with full bandwidth.

20 37. The integrated circuit device of claim 12, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast multi-link butterfly fat tree network with full bandwidth.

38. The integrated circuit device of claim 12, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in

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each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast multi-link butterfly fat tree network and rearrangeably nonblocking for arbitrary fan-out multicast multi-link butterfly fat tree network with full bandwidth.

5 39. The integrated circuit device of claim 12, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast multi-link butterfly fat tree network with full bandwidth.

40. The integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and $y \ge (\log_2 N)$.

15 41. The integrated circuit device of claim 40, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized multi-link multi-stage
20 network with full bandwidth.

42. The integrated circuit device of claim 40, wherein d = 4 and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said

25 routing network is strictly nonblocking for unicast generalized multi-link multi-stage network and rearrangeably nonblocking for arbitrary fan-out multicast generalized multilink multi-stage network with full bandwidth.

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43. The integrated circuit device of claim 40, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast generalized multilink multi-stage network with full bandwidth.

44. The integrated circuit device of claim 5, wherein said all horizontal shuffle exchange links between switches in any two corresponding said succeeding stages are of different length and said vertical shuffle exchange links between switches in any two

10 corresponding said succeeding stages are of different length and $y \ge (\log_2 N)$, and said each sub-integrated circuit block further comprising a plurality of U-turn links within switches in each of said stages in each of said sub-integrated circuit blocks.

45. The integrated circuit device of claim 44, wherein d = 4 and there is only one switch in each said stage in each said sub-integrated circuit block connecting said
15 forward connecting links and there is only one switch in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is rearrangeably nonblocking for unicast generalized multi-link butterfly fat tree network with full bandwidth.

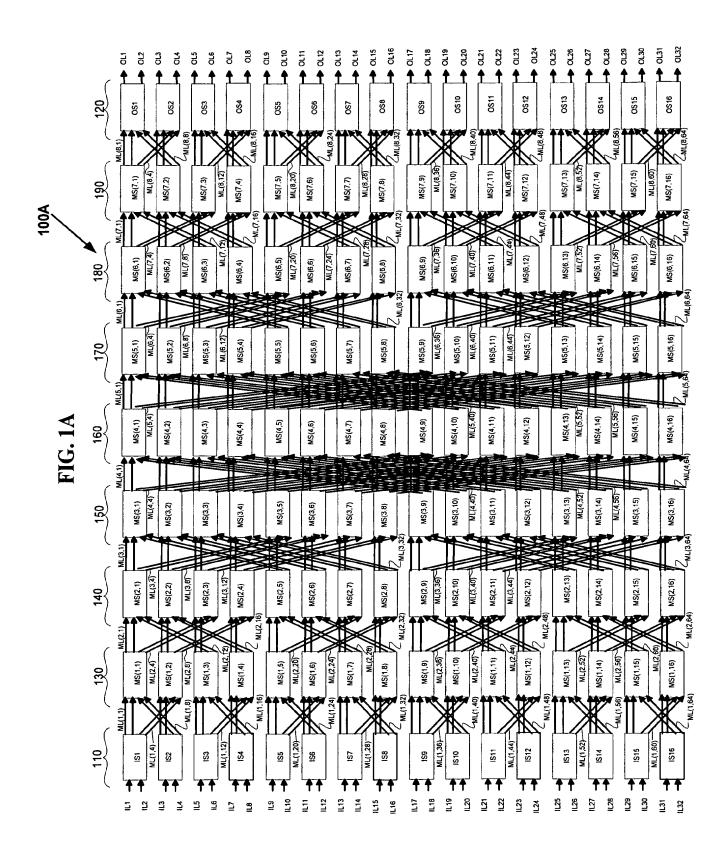
46. The integrated circuit device of claim 44, wherein d = 4 and there are at
20 least two switches in each said stage in each said sub-integrated circuit block connecting said forward connecting links and there are at least two switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for unicast generalized multi-link butterfly fat tree Network and rearrangeably nonblocking for arbitrary fan-out multicast generalized multi25 link butterfly fat tree network with full bandwidth.

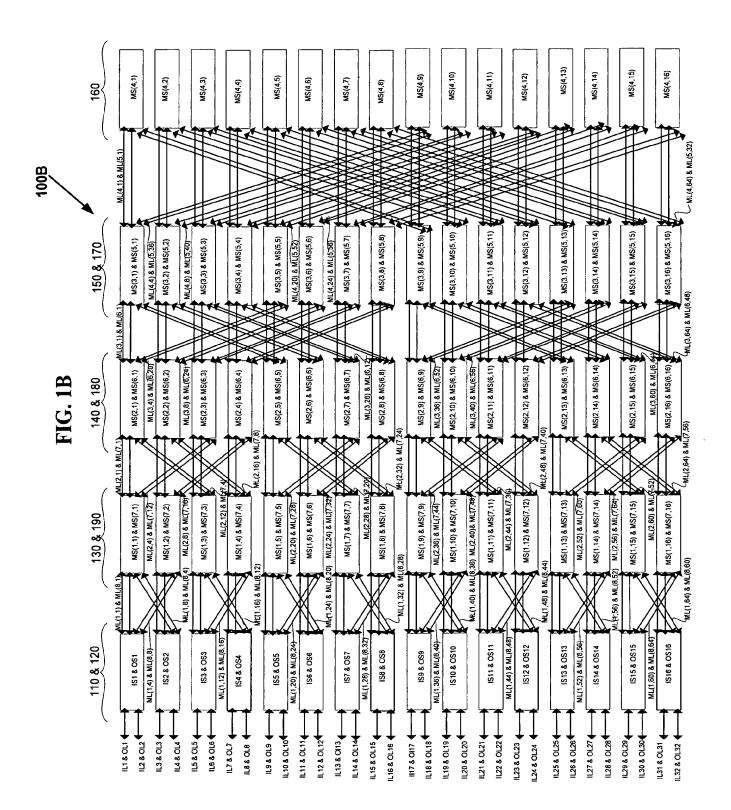
47. The integrated circuit device of claim 44, wherein d = 4 and there are at least three switches in each said stage in each said sub-integrated circuit block connecting

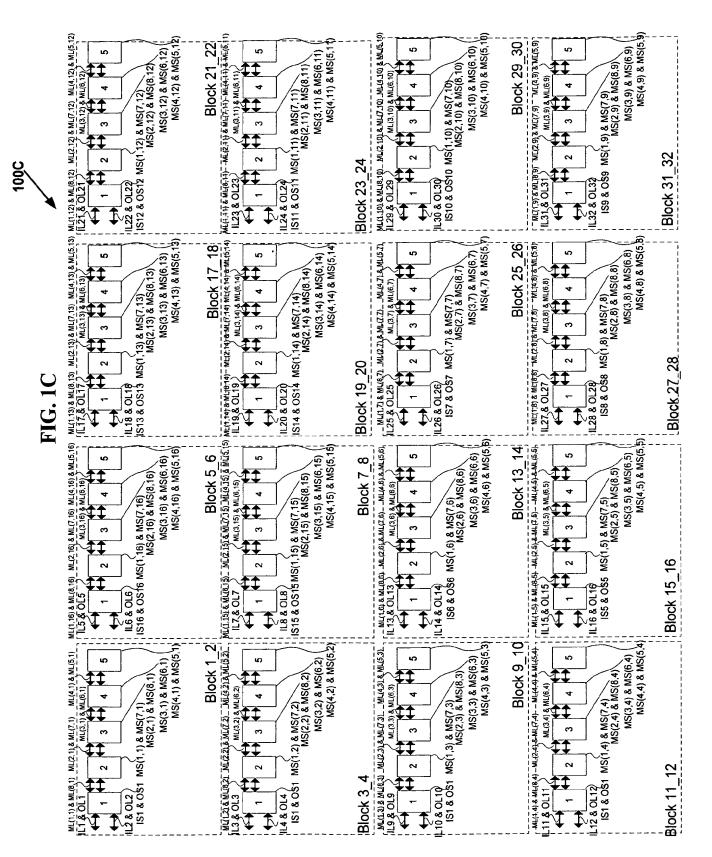
-66-

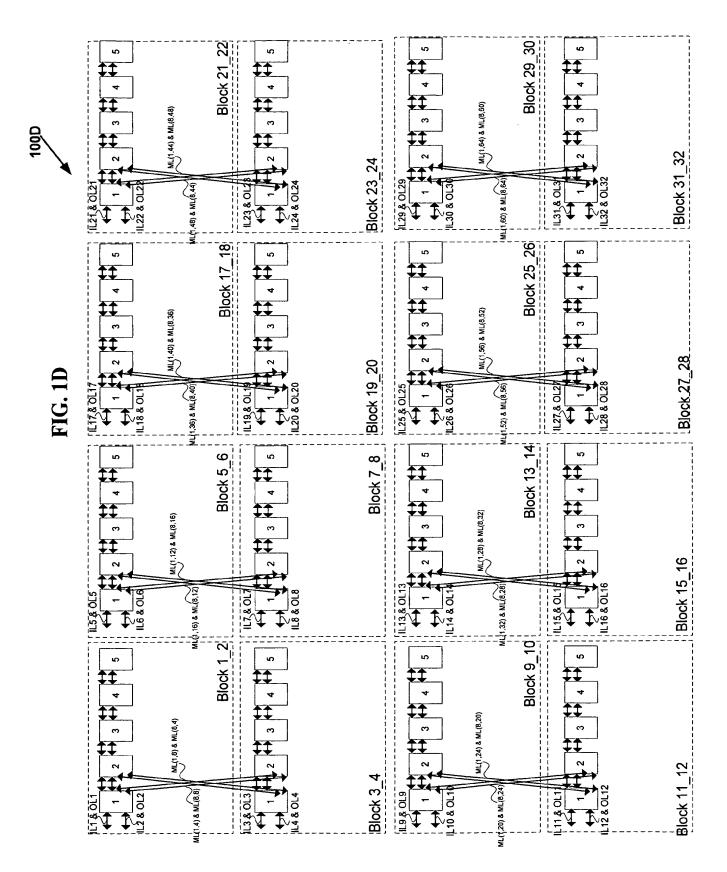
said forward connecting links and there are at least three switches in each said stage in each said sub-integrated circuit block connecting said backward connecting links and said routing network is strictly nonblocking for arbitrary fan-out multicast generalized multilink butterfly fat tree network with full bandwidth.

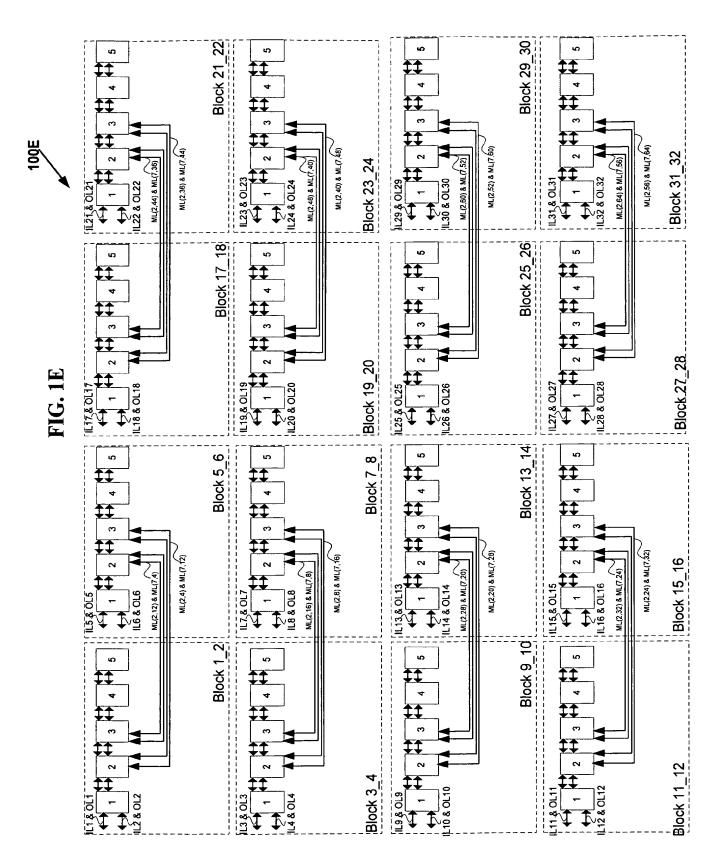
- 5 48. The integrated circuit device of claim 1, wherein said plurality of forward connecting links use a plurality of buffers to amplify signals driven through them and said plurality of backward connecting links use a plurality of buffers to amplify signals driven through them; and said buffers can be inverting or non-inverting buffers.
- 49. The integrated circuit device of claim 1, wherein said wherein said all 10 switches of size $d \times d$ are either fully populated or partially populated.





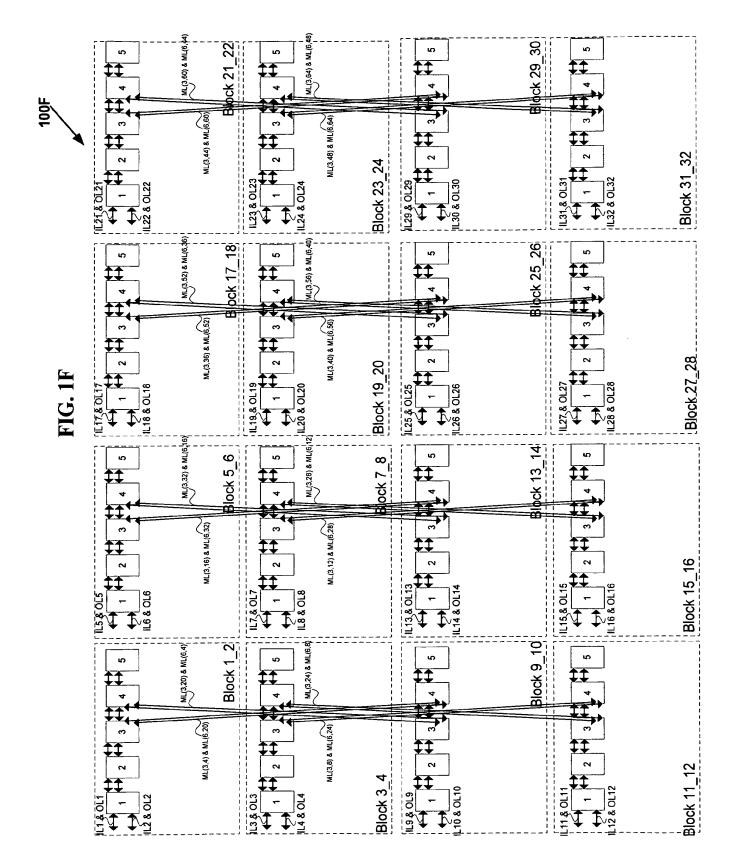


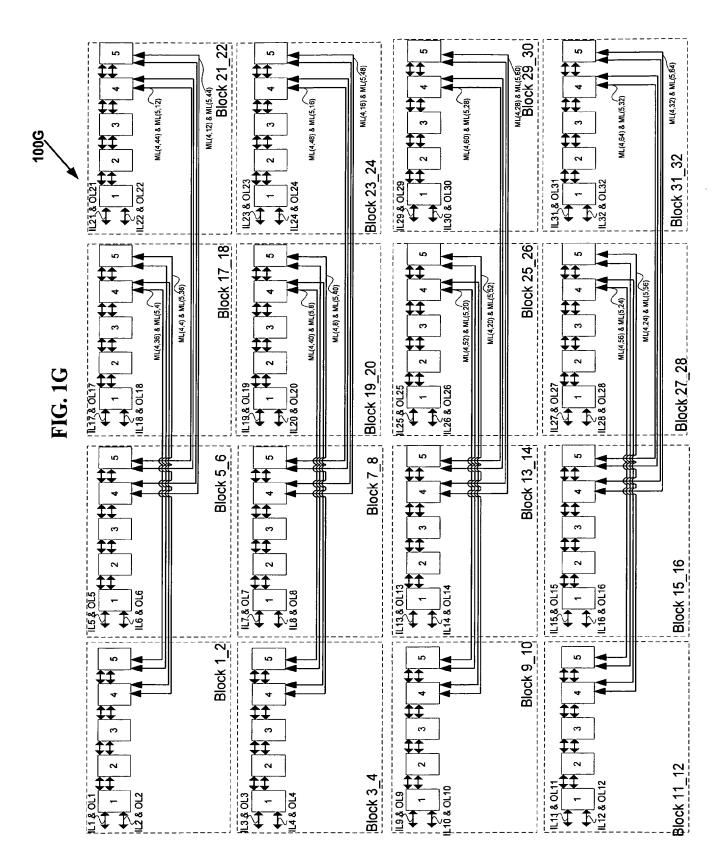




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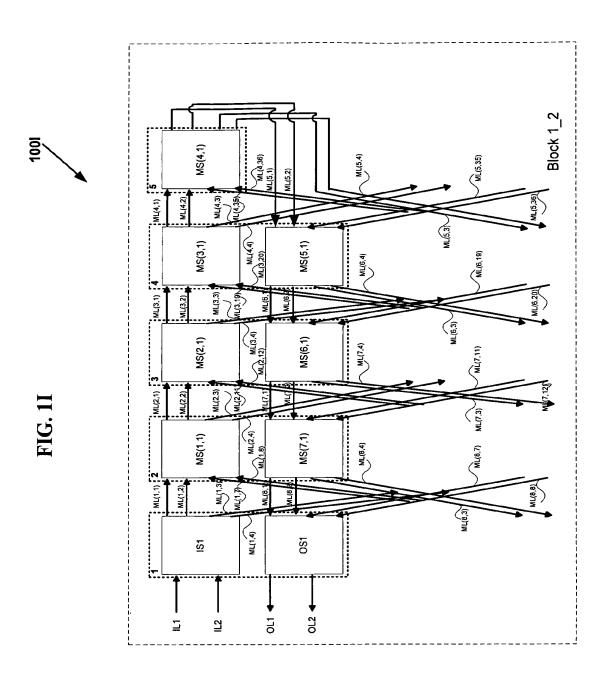






	27 4.0127	Block 35 86	Block 87_88 23 & 0.23 14 20 31 44 34 34 34 34 34 34 34 34 34 34 34 34	Block 93 94	Block 95 96	Block 117 118 (1) 21 3 4 3 8 7 26 8 0126	Block 119 120	Block 125 126	Block 127 128
100H		Block 81 82 8005 81 82 1 22 3 4 5 6 7 2	Block 83 84	Block 89 90	Block 91 92	Block 113 114	Block 115 116	Block 121 122	
			Block 71_72 115 60015 116 20 31 41 31 31 41 31 31 41 31 31 31 31 31 31 31 31 31 31 31 31 31	Block 77 _ 78	Block 79 80	Block 101 102	Block 103 104	Block 109 110 112 8 013 4 0 0 110	Block 107_108 Block 111_112 Block 123_124
FIG. 1H		Block 65, 66, B	Block 67_68	Block 73_74	Block 75 76	Block 97_98	Block 99, 100 10, 20, 39, 100 10, 20, 31, 41, 61, 7 1, 10, 8, 01, 10	Block 105_106	Block 107_108
	1128 & 0127	Block 21_22 1 1 2 3 3 4 5 6 6 7 1 26 8 0126	Block 23_24 123.8.023 11 21 31 31 41 61 7 11 24 31 01 24	Block 29_30	Block 31 32 1423 - 0127 1412 - 0127 1128 & 0128	Block 53_54	Block 55 56	Block 61_62 1 1 2 3 4 3 5 6 7	Block 63 64
		17_18	20	25_26					<u>30</u>
		BOCK 1423 BOCK 1423 34 4	Block 19 H147 8-0147- 1122 31 4 IL18 & 0L18	H-14 8005	Block 27 2	Block 49 1428 314 49 1428 31 413	Block 51_5 H47 & OL47_5 1 1 2 2 0 4 4 7 L18 8 0L18	BIOCK 5	Block 59
		Block 1_2 Block 1_2 Block 5_6 Block 5_6	Block 3 4. Block 7 8. Block 19 15 2 19 3 4 5 6 7 1 1 2 3 3 4 5 6 7 7 1 1 2 3 3 4 5 6 7 1 1 1 2 3 3 4 5 6 6 7 1 1 1 2 3 3 4 5 6 6 7 1 1 1 2 3 0 1 1 1 1 1 2 0 1 1 1 1 1 1 1 1 1 1 1 1	Block 9 10 Block 13 14 Block 2 al	Block 11 12 Block 15 16 Block 27 28	Block 33_34 Block 37_38 Block 49 13 0 3 14 15 0 15 0 15 1 38 15 50 14 14 24 34 54 54 54 54 54 54 54 54 54 54 54 54 54	Block 35 36 Block 39 40 Block 51 5 168 89 36 Block 39 40 Block 51 5 188 89 36 4 5 6 1 5 6 1 5 6 1 5 6 1 5 6 1 5 6 1 5 6 1 5 6 1 5 6 1 5 6 1 5 6 1 5 6 6 6 1 5 6 6 1 5 6 6 1 5 6 6 6 1 5 6 6 1 5 6 6 1 5 6 6 1 5 6 6 1 5 6 6 1 5 6 6 1 5 6 6 1 5 6 6 6 1 5 6 6 6 1 5 6 6 6 1 5 6 6 6 1 5 6 6 6 6	Riger 41 42 Block 45 46 Res Block 5 1 4 2 3 4 5 4 6 8 7 1 4 2 3 4 5 6 7 1 1 2 3 4 5	Block 47 48 Block 59 60

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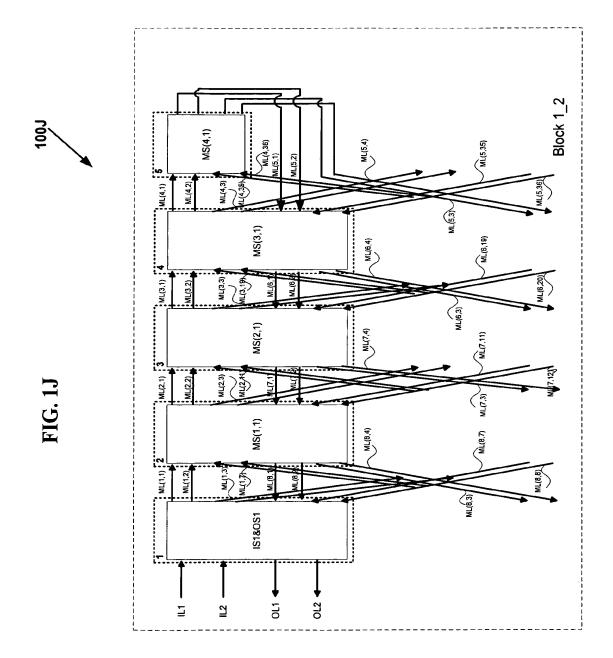
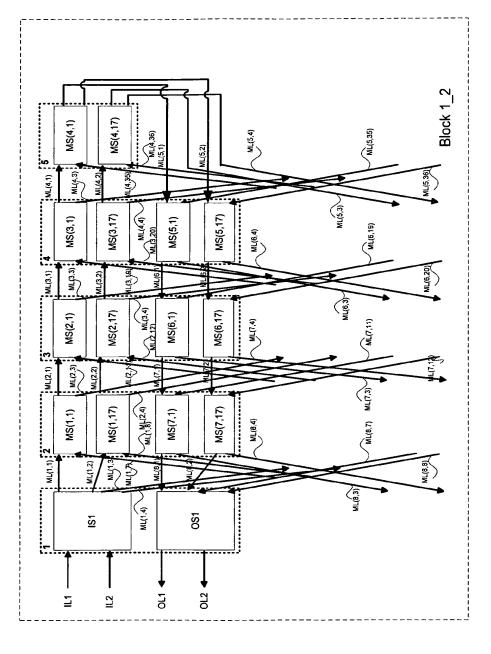
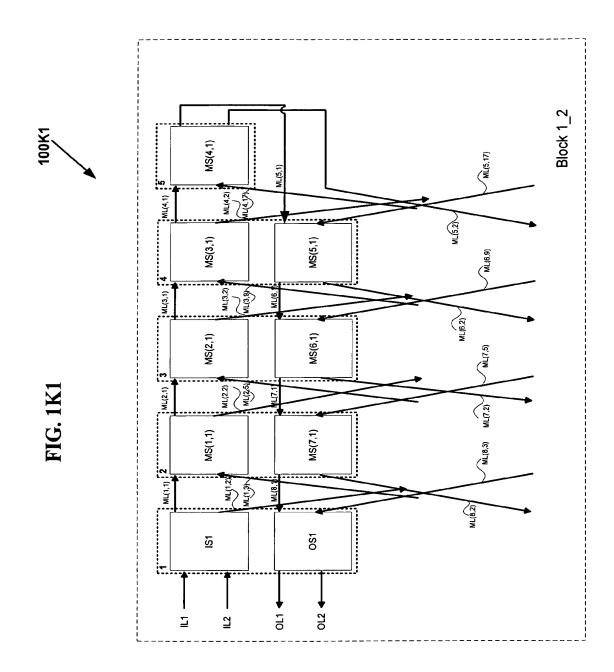


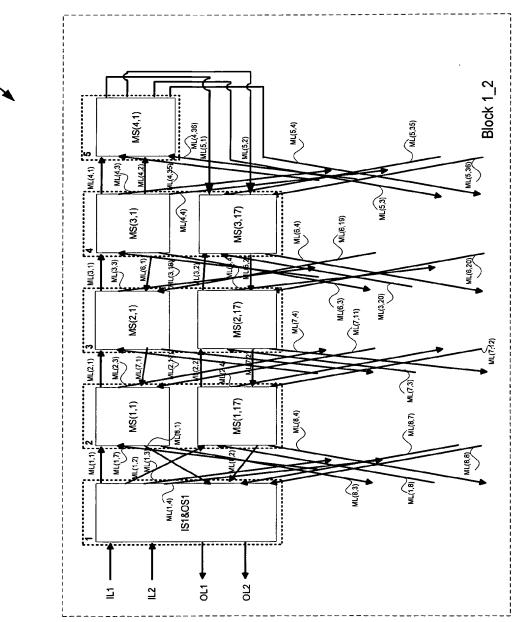




FIG. 1K



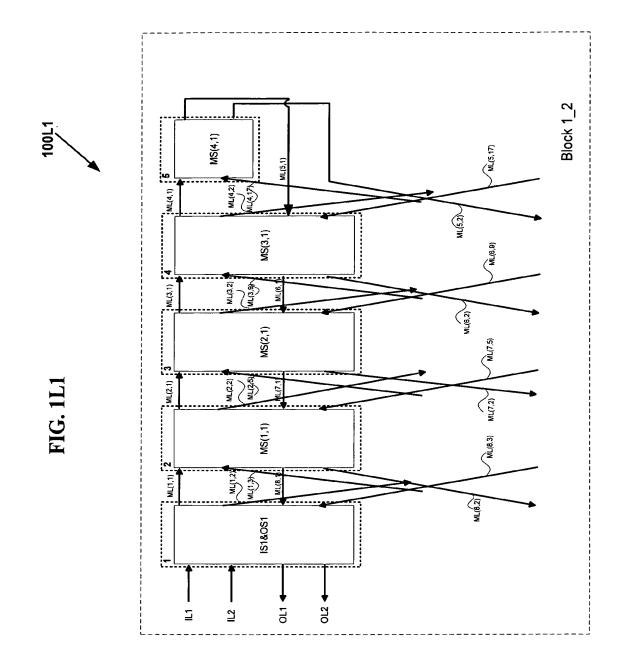


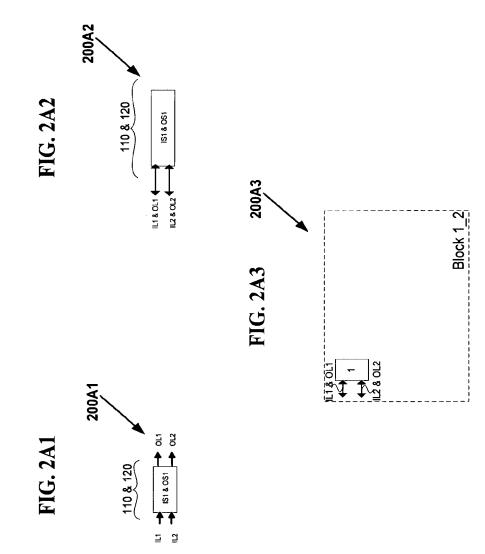


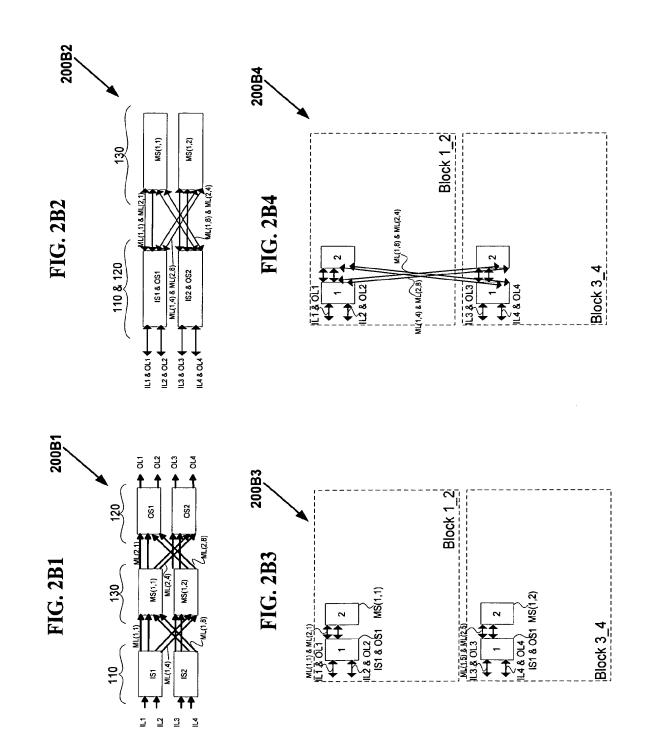
.1L

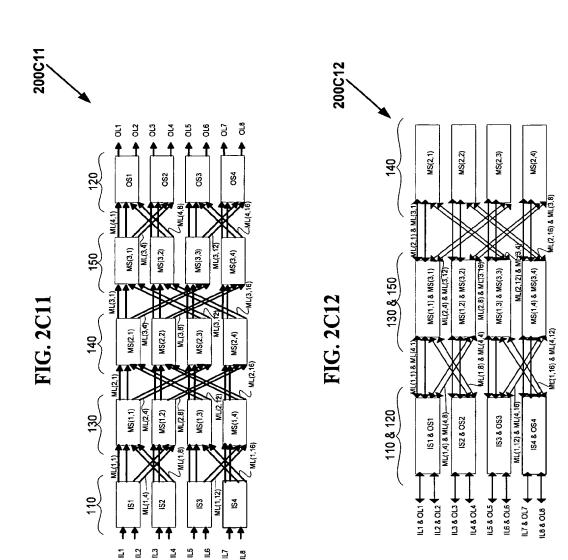
100L

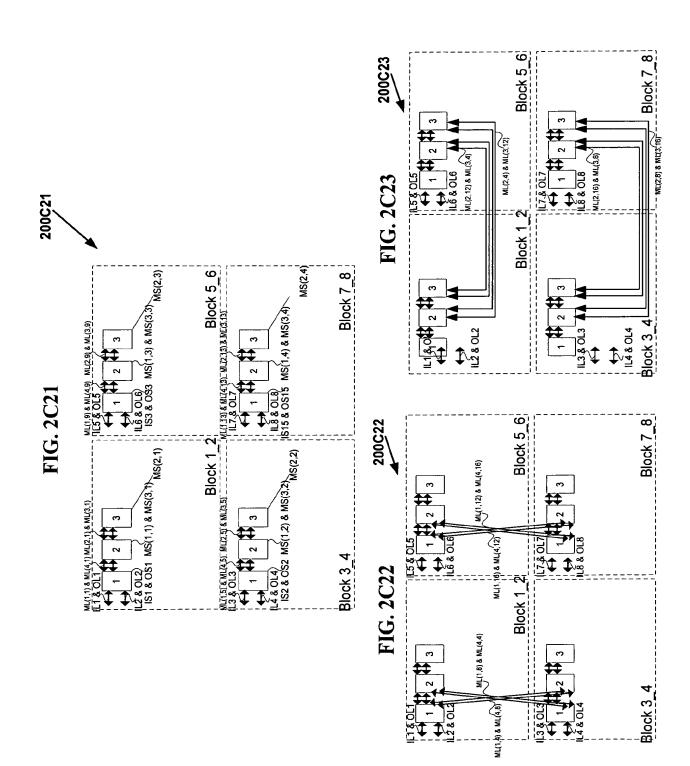
FIG. 1L

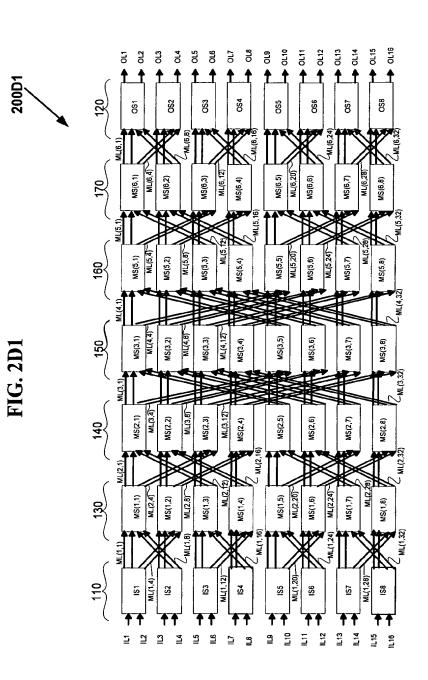












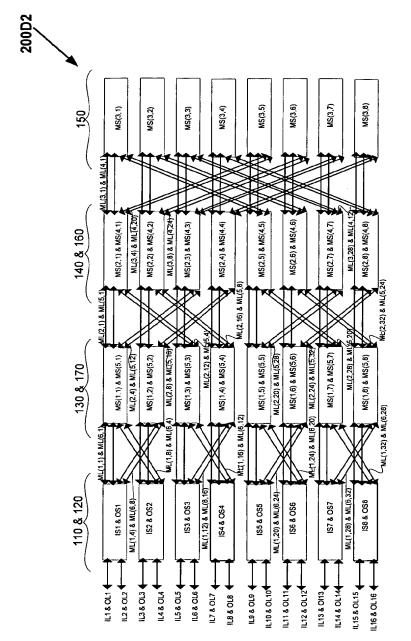
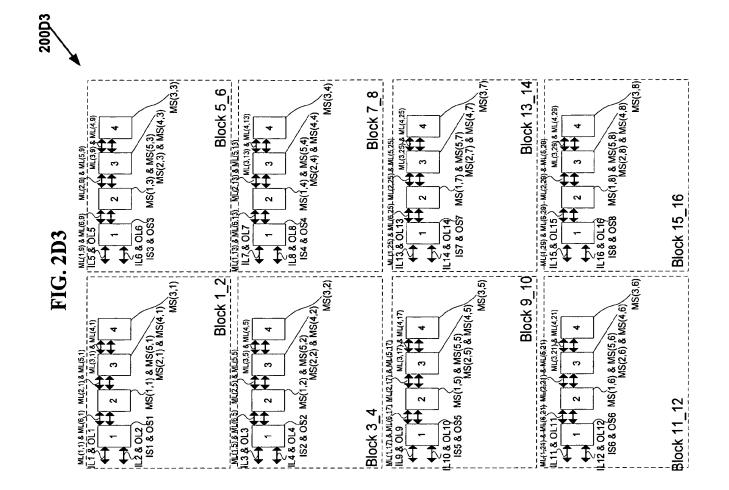
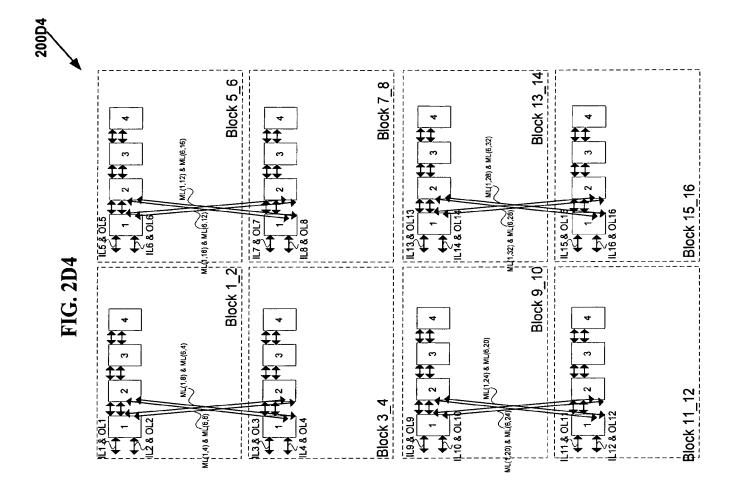
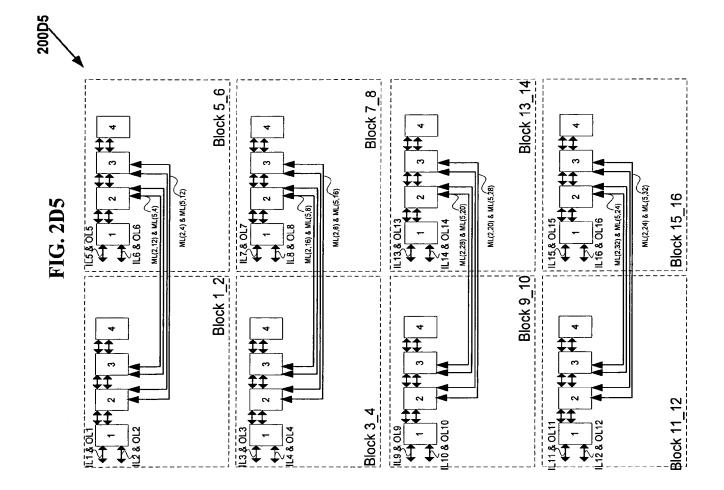


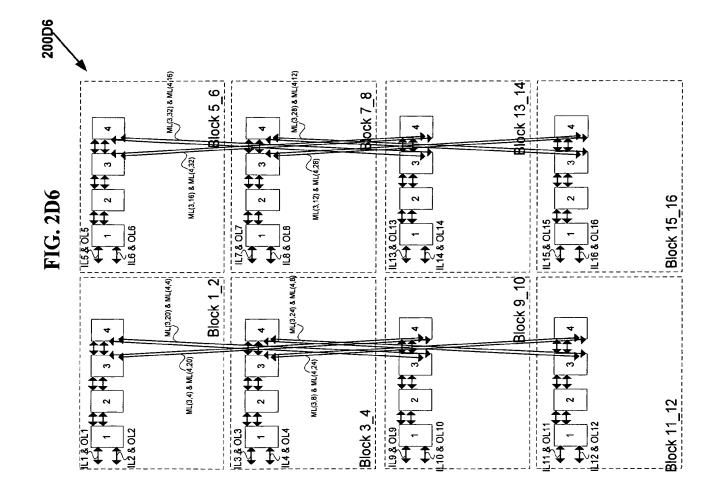
FIG. 2D2

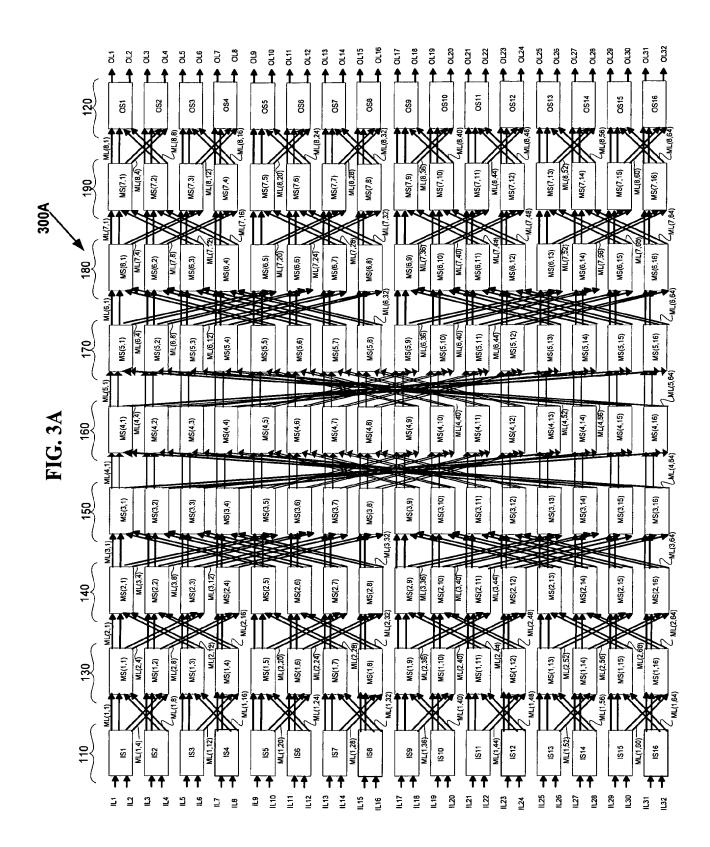


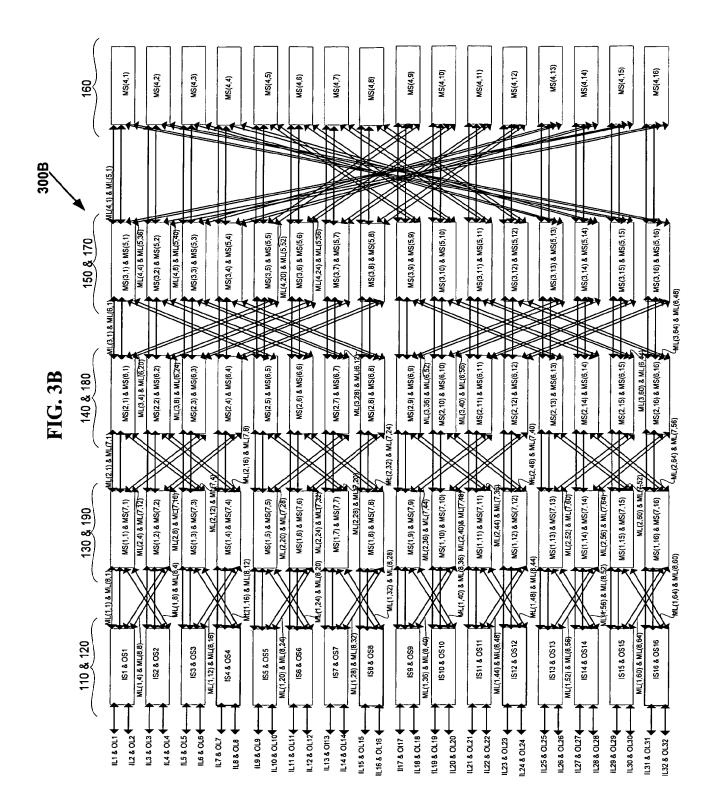




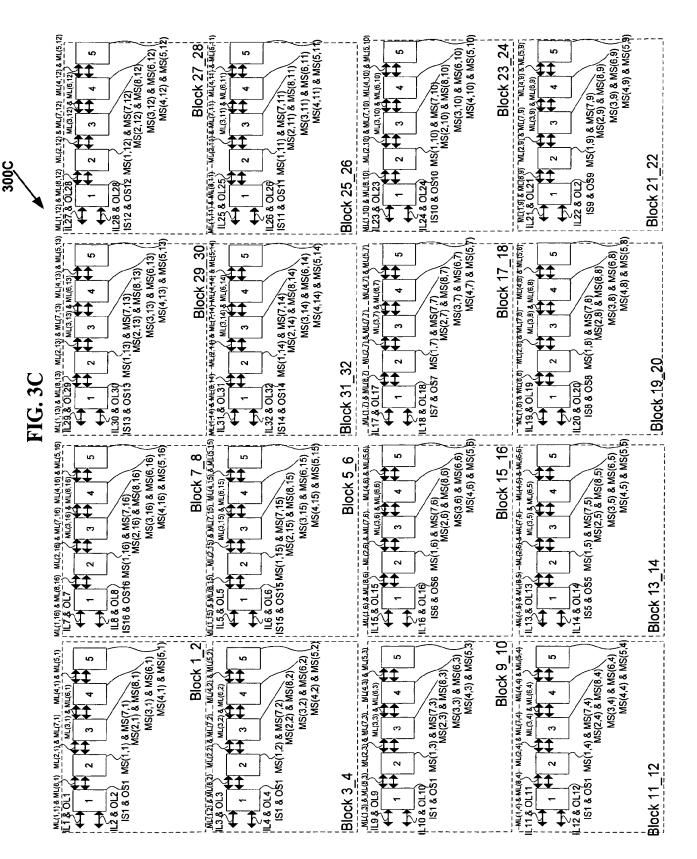
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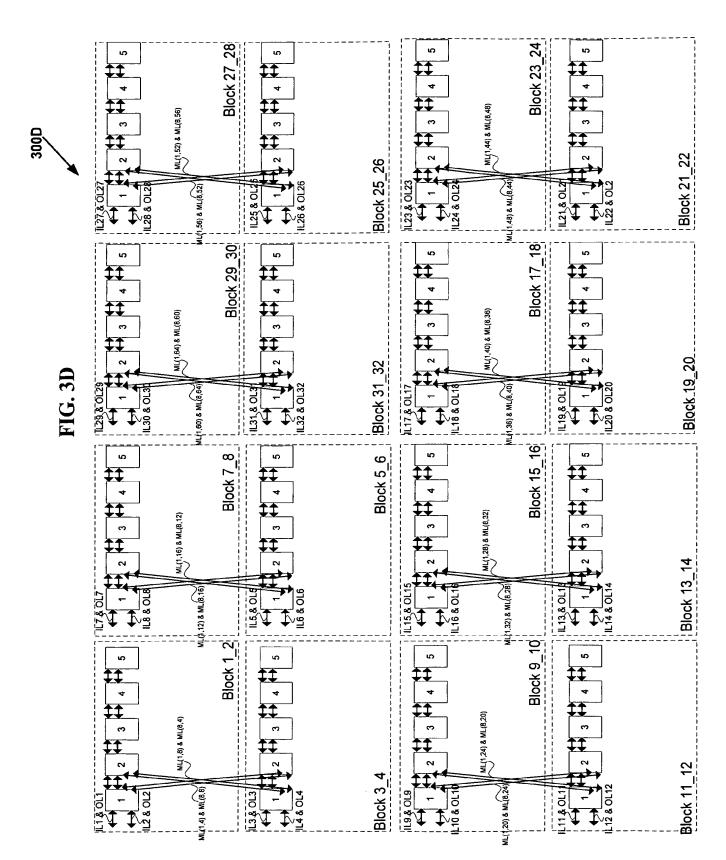


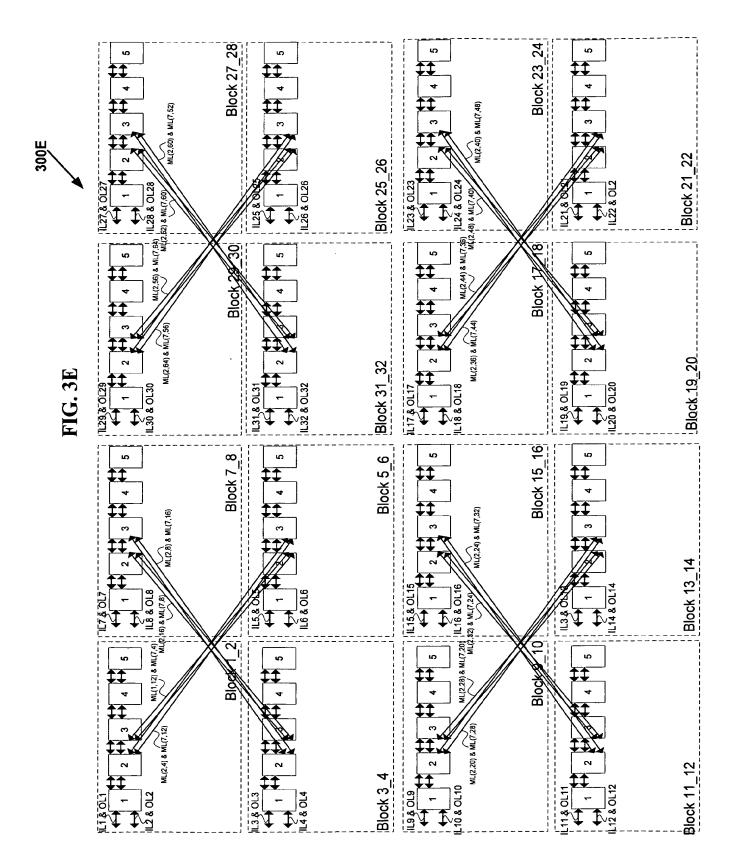




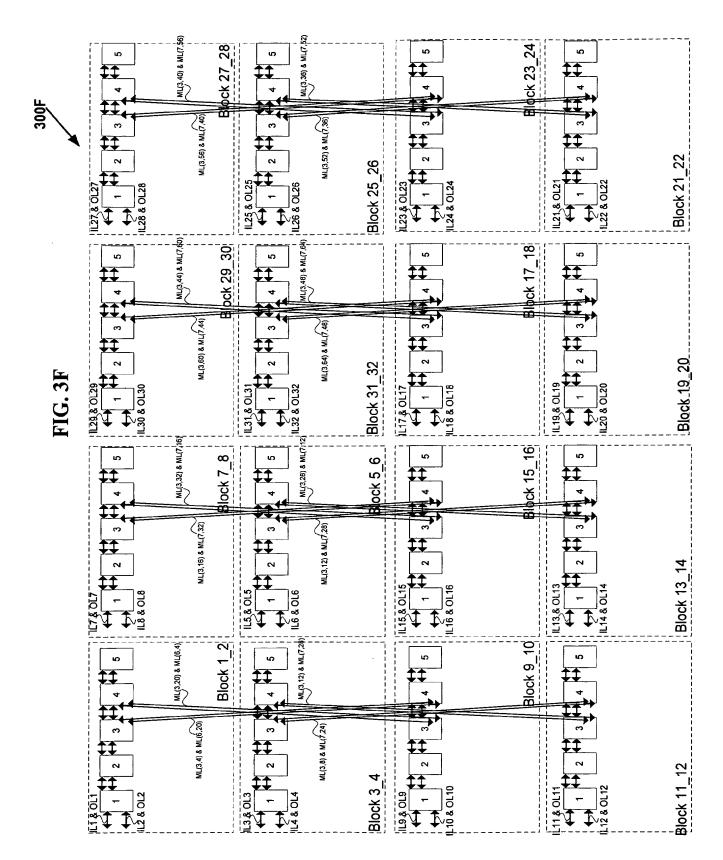


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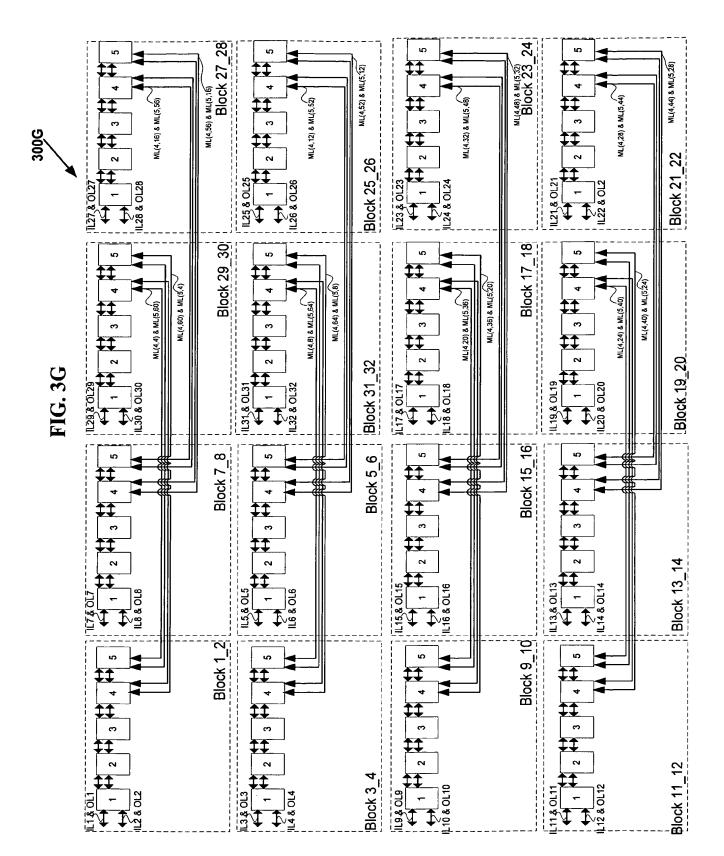




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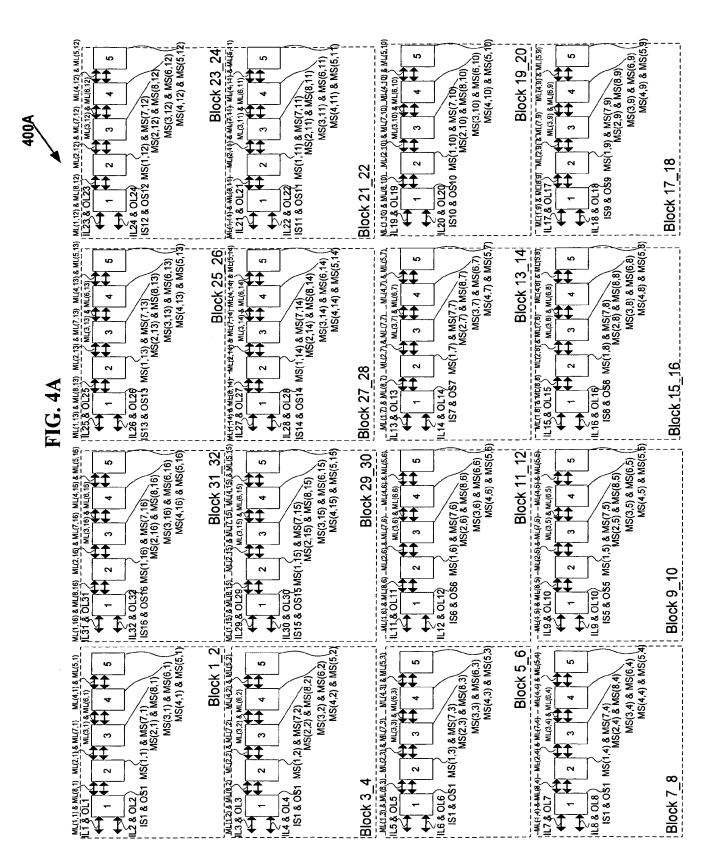


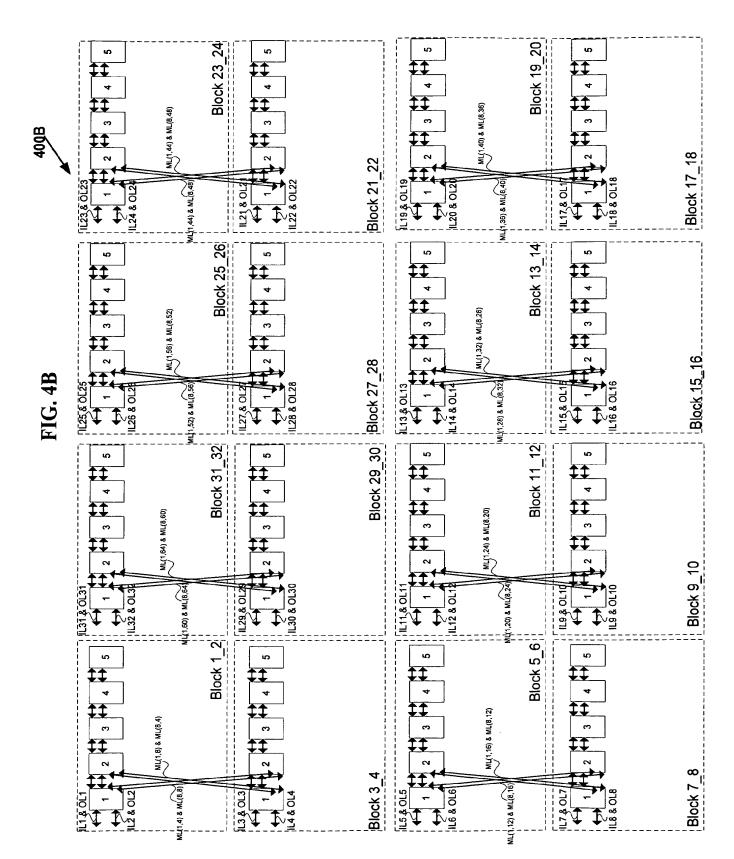


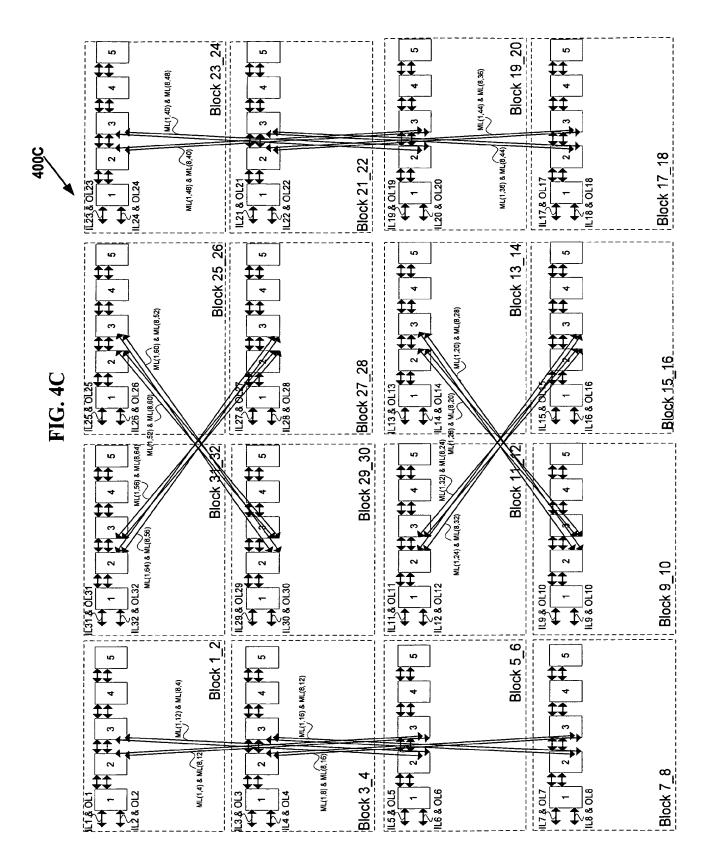
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	11-27 & 91-27	Block 123 124	Block 121_122 1.23 & 0.23 1.42 & 0.23 1.42 & 0.24	Block 119 120	Block 117 118		Block 91 92	Block 89 90 H23 & 023 1 1 2 3 4 3 4 3 4 3 4 3 4 3 4 3 4 3 4 3 4 3	Block 87 88	Block 85_86
300H	11-20-20-29-2	Block 125 126	Block 127_128	Block 113 114	Block 107 108 Block 109 110 Block 115 116 Block 117 118		Block 93_94	Block 95 96	Block 81 82	Block 83 84
		Block 103_104	Block 101_102	Block 111_112	Block 109 110		Block 71_72	Block 69_70	Block 79 80	Block 77 78
FIG. 3H		Block 97_98	Block 99_100	Block 105 106 F	Block 107 108		Block 65_66 Block 71	Block 67_68_ 1156099 11108.0110	Block 73 74	<u>Block 75_76</u>
	122 & 0127	Block 27_28	Block 25, 26	24 Block 23 24	22		Block 59_60	Block 57_58 N23 & 0123 1 1 2 1 3 1 4 3 1 6 7 1 2 1 3 1 4 5 1 6 7	Block 55 56	Block 53 54
		135 Block 29_30	Block.31.32 H478-047-32 1112-30-47 1118-30-118	1120 & 0120	Block 19 20		Block 61, 62	Block 63_64	120 & 0120	Block 51 52
		Block 1 2 Block 7 8	Block 5_6	HOCK 15 16	Block 13_14		Block 39_40	Block 37_38	Block 47 48	Block 45_46 Blo
	11-1-2-1	Block 1 2 Block 1 2	Block 3.4 Block 5.6 Block 1.6 Block 5.6 Block 1.47	Block 9 10 Block 15 16 Hock 15 16 Hock 15 16 Hock 11 16 Hock 11 16 Hock 11 16 Hock 15 16 Hock 11 16 Hock 15 16 Hock 11 16	Block 11_12	H1 8 0 2 1 3 4 4 1 1 1 1 1 1 2 1 3 4 0 2 1 2 1 3 4 4 5 1 2 1 3 4 4 5 1 2 1 3 4 4 5 1 2 1 3 4 4 5 1 2 1 3 4 4 5 1 3 1 3 1 3 1 3 1 3 1 3 1 3 1 3 1 3 1	Block 33_34 Block 39_40	Block 35 36 Block 37 38 Block 18 Block 37 38 Block 18 Block 19 Blo	No. Rigek 41 42 Block 47 48	Block 43_44

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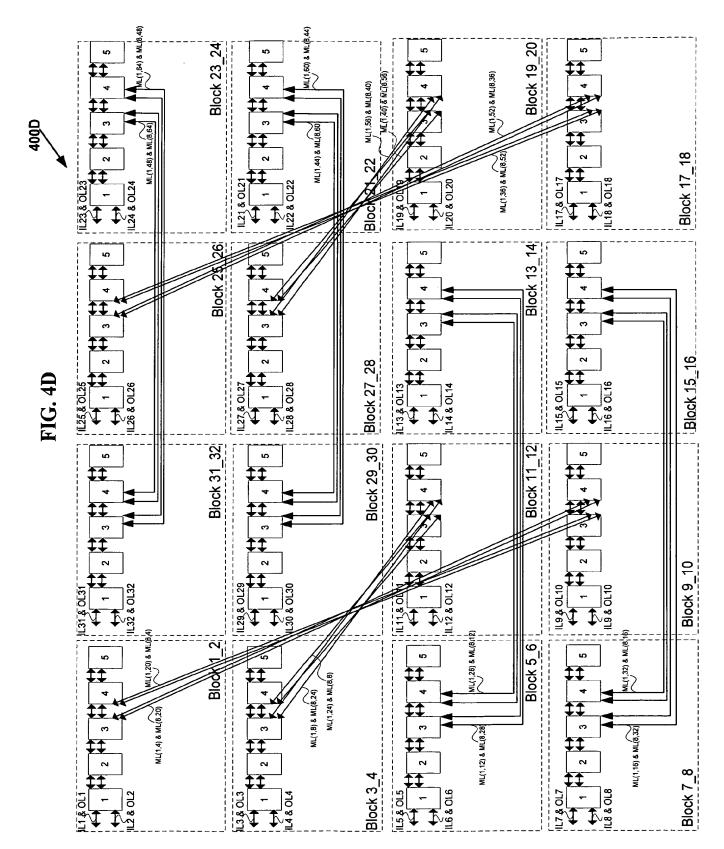


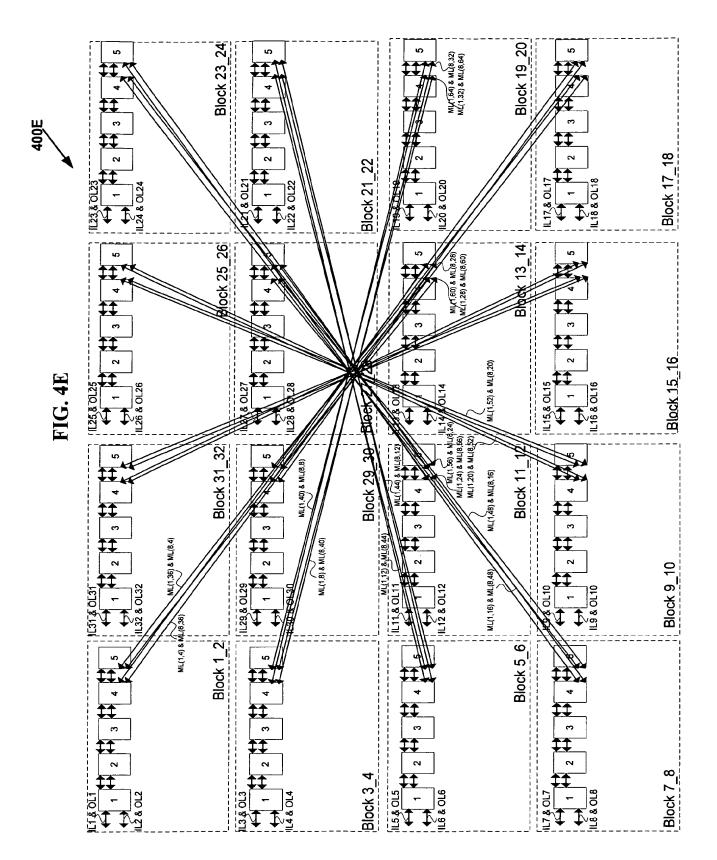




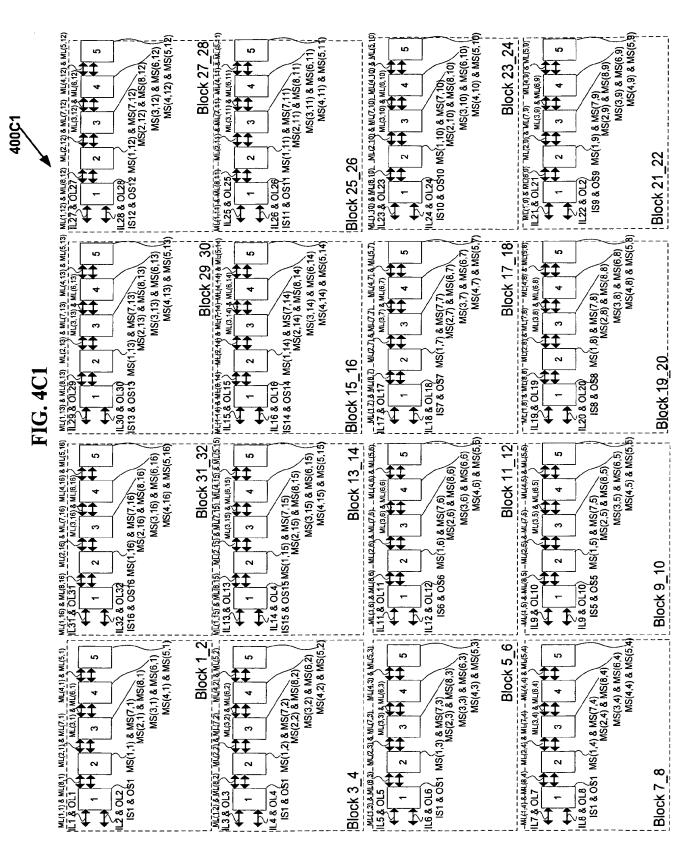
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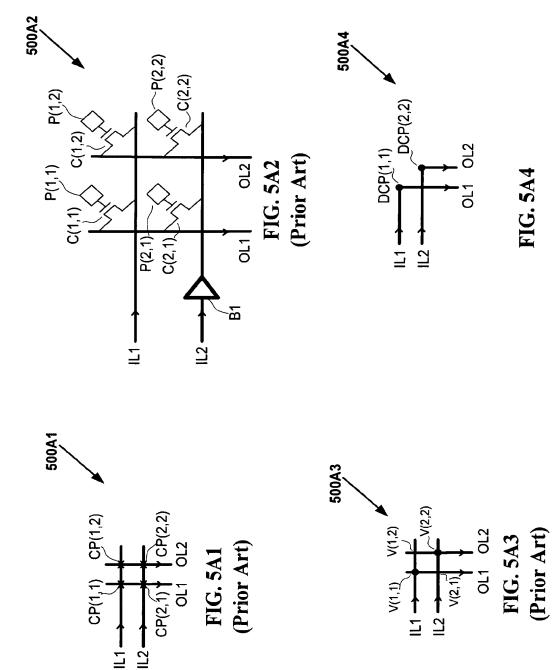


FIG. 5A