

Patent Number:

Date of Patent:

[11]

[45]

US006151708A

6,151,708

Nov. 21, 2000

United States Patent [19]

Pedrizetti et al.

[54] **DETERMINING PROGRAM UPDATE** AVAILABILITY VIA SET INTERSECTION **OVER A SUB-OPTICAL PATHWAY**

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[21] Appl. No.: 08/994,594

Dec. 19, 1997 [22] Filed:

[51]

[52] [58]

395/707, 712, 709; 711/184; 382/166; 707/8; 709/221, 217; 717/11

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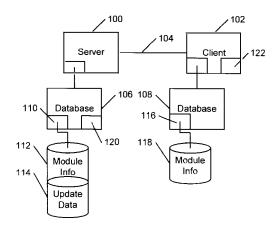
(List continued on next page.)

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ABSTRACT [57]

A set of software programs on a client computer is compared against a set of updates on a server computer to determine which updates are applicable and should be transferred from the server to the client. If the link between the client and server is slow, the listing of available updates must be represented in compact form. A many-to-one mapping function (e.g. a hash function) is applied to update identifiers to generate a table of single bit entries indicating the presence of particular updates on the server. This table is transferred to the client over the slow link. At the client, the same mapping function is applied to program identifiers, and corresponding entries of the transferred table are checked to determine whether the server has a potential update. If such a potential update is noted, a second transmission is requested by the client from the server—this one conveying additional data by which hash collisions can be identified by the client and disregarded. If availability of an actual update (versus a hash collision) is thereby confirmed, the client requests a third transmission from the server-this one conveying the actual update data. By this arrangement, optimized use is made of the low bandwidth link, with successively more information transferred as the likelihood of an applicable update is successively increased. (The same arrangement can be employed in reverse, with the bit table generated at the client and identifying program files available for possible updating, transferred to the server, etc.).

22 Claims, 13 Drawing Sheets



MICROSOFT CORP. EXHIBIT 1007



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FIG. 1

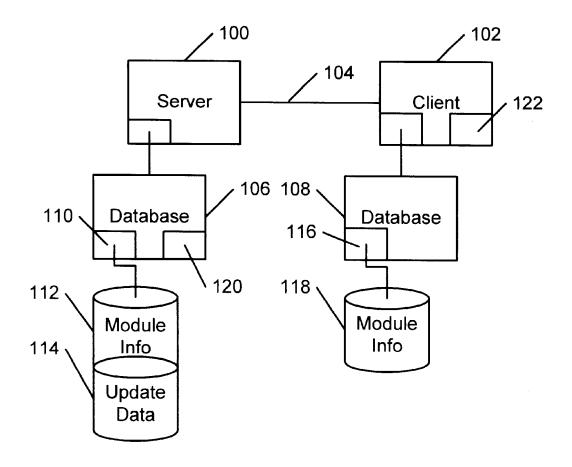




FIG. 2

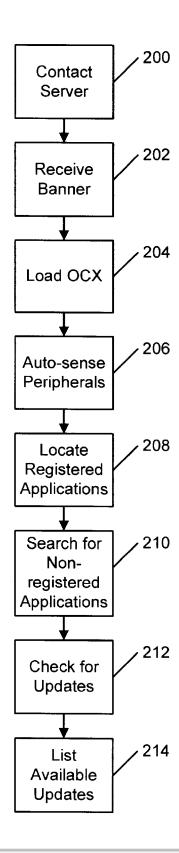
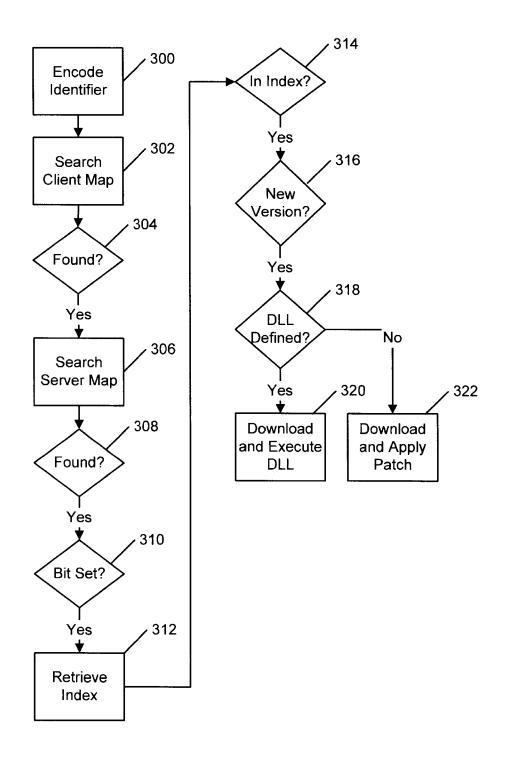




FIG. 3 — Step 212 of FIG. 2



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