Presentation of Petitioners Riot Games, Inc. & Valve Corp.

Case Nos. IPR2018-00129, -130, -131, -132

US Patent Nos. 5,822,523 & 6,226,686

Roadmap

Overview

- '523 and '686 Patents
- Combination of Aldred and RFC 1692

Independent Claim Disputes

- Patent Owner's Motivation to Combine Arguments
- Claim Construction of "Aggregated Payload" and "Aggregated Message"

Dependent Claim Disputes

- Group Messaging Claims
- Ulrich Combination Message Groups and Echo Suppression

Roadmap

Overview

- '523 and '686 Patents
- Combination of Aldred and RFC 1692

Independent Claim Disputes

- Patent Owner's Motivation to Combine Arguments
- Claim Construction of "Aggregated Payload" and "Aggregated Message"

Dependent Claim Disputes

- Group Messaging Claims
- Ulrich Combination Message Groups and Echo Suppression

Overview of the '523 and '686 Patents

United States Patent [19] Rothschild et al.			[11] [45]	Patent Number: Date of Patent:	5,822,523 Oct. 13, 1998	
[54]	[54] SERVER-GROUP MESSAGING SYSTEM FOR INTERACTIVE APPLICATIONS			Primary Examiner—William M. Treat Assistant Examiner—Zarni Maung		
[75]	Inventors:	Jeffrey J. Rothschild; Marc P. Kwiatkowski, both of Los Gatos; Daniel J. Samuel, Sunnyvale, all of Calif.	Attorney,	Agent, or Firm—II. C. a & Rosati ABSTRAC	Chan; Wison Sonsini	
[73]	Assignee:		network servers is	od for deploying interacts containing host computer s disclosed. The etwork architecture		
[21] [22]	Appl. No. Filed:	: 595,323 Feb. 1, 1996	work link messages		United Stat Rothschild et al.	

..... H04H 1/02

395/200.09

395/200 1 200 01

message groups maintai

For each message grou

maintain a list of all of

particular group. In its

Ex. 1001, Face

[58] Field of Search

tes Patent Rothschild et al.

(10) Patent No.: US 6,226,686 B1

(45) **Date of Patent:**

May 1, 2001

SERVER-GROUP MESSAGING SYSTEM FOR INTERACTIVE APPLICATIONS

(75) Inventors: Jeffrey Jackiel Rothschild, Los Gatos; Daniel Joseph Samuel, Sunnyvale;

Marc Peter Kwiatkowski, Los Gatos, all of CA (US)

(73) Assignee: HearMe, Mountain View, CA (US)

(*) Notice: Subject to any disclaimer, the term of this patent is extended or adjusted under 35

U.S.C. 154(b) by 0 days.

(21) Appl. No.: 09/407,371

Sep. 28, 1999

Related U.S. Application Data

Continuation of application No. 08/896,797, filed on Jul. 18, 1997, now Pat. No. 6,018,766, which is a continuation of application No. 08/595,323, filed on Feb. 1, 1996, now Pat. No. 5,822,523.

OTHER PUBLICATIONS

Ahuja, S.R., et al., "The Rapport Multimedia Conferencing System," Conference on Office Information Systems 1988, pp. 1-7.

(List continued on next page.)

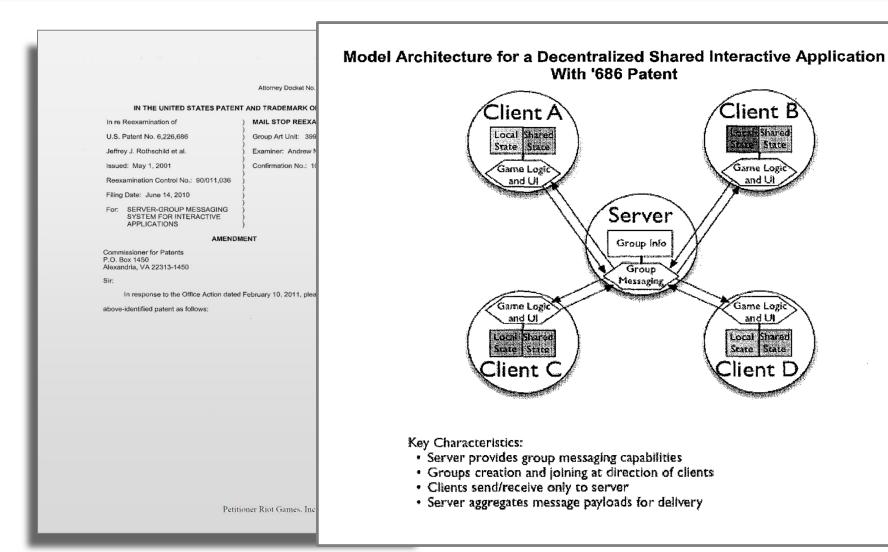
Primary Examiner—Zarni Maung (74) Attorney, Agent, or Firm-Sterne, Kessler, Goldstein, & Fox PLLC

ABSTRACT

A method for deploying interactive applications over a network containing host computers and group messaging servers is disclosed. The method operates in a conventional unicast network architecture comprised of conventional network links and unicast gateways and routers. The hosts send messages containing destination group addresses by unicast to the group messaging servers. The group addresses select message groups maintained by the group messaging servers. For each message group, the group messaging servers also maintain a list of all of the hosts that are members of the

Ex. 1002, Face

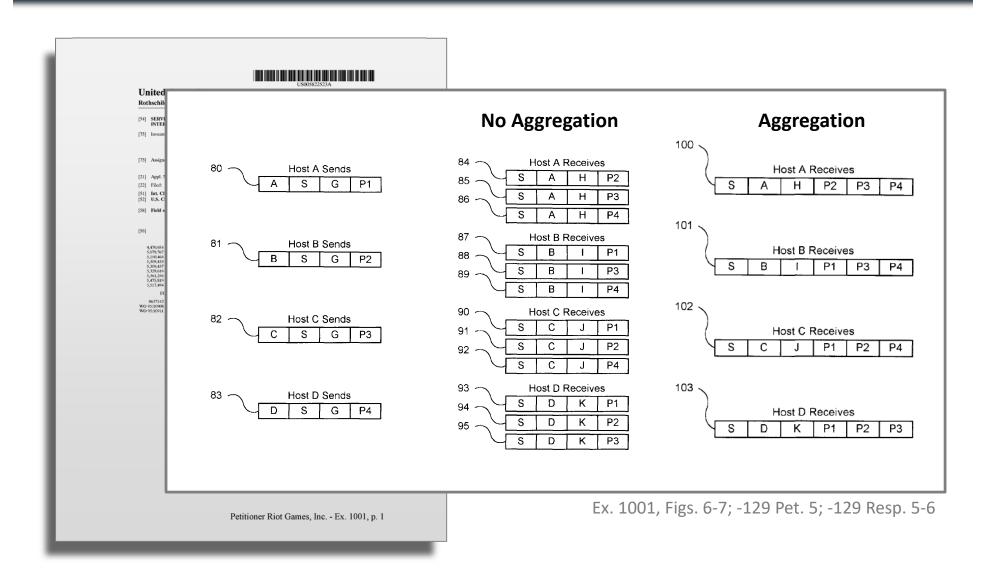
Overview of the '523 and '686 Patents



Ex. 1006, 223; -129 Pet. 8

State

Overview of the '523 and '686 Patents



'523 and '686 Patents - Ind. Claim 1

- 1. A method for providing group messages to a plurality of host computers connected over a unicast wide area communication network, comprising the steps of:
 - providing a group messaging server coupled to said network, said server communicating with said plurality of host computers using said unicast network and maintaining a list of message groups, each message group containing at least one host computer;
 - sending, by a plurality of host computers belonging to a first message group, messages to said server via said unicast network, said messages containing a payload portion and a portion for identifying said first message group;
 - aggregating, by said server in a time interval determined in accordance with a predefined criterion, said payload portions of said messages to create an aggregated payload;
 - forming an aggregated message using said aggregated payload; and
 - transmitting, by said server via said unicast network, said aggregated message to a recipient host computer belonging to said first message group.

- 1. A method for facilitating communications among a plurality of host computers over a network to implement a shared, interactive application, comprising the steps of:
 - (1) receiving a create message from one of the plurality of host computers, wherein said create message specifies a message group to be created;
 - (2) receiving join messages from a first subset of the plurality of host computers, wherein each of said join messages specifies said message group;
 - (3) receiving host messages from a second subset of said first subset of the plurality of host computers belonging to said message group, wherein each of said messages contains a payload portion and a portion that is used to identify said message group;
 - (4) aggregating said payload portions of said host messages received from said second subset of the plurality of host computers to create an aggregated payload;
 - (5) forming an aggregated message using said aggregated payload; and
 - (6) transmitting said aggregated message to said first subset of the plurality of host computers belonging to said message group;
 - wherein said aggregated message keeps the shared, interactive application operating consistently on each of said first subset of the plurality of host computers.

Independent Claim Disputes

- 1. A method for providing group messages to a plurality of host computers connected over a unicast wide area communication network, comprising the steps of:
 - providing a group messaging server coupled to said network, said server communicating with said plurality of host computers using said unicast network and maintaining a list of message groups, each message group containing at least one host computer;
 - sending, by a plurality of host computers belonging to a first message group, messages to said server via said unicast network, said messages containing a payload portion and a portion for identifying said first message group;
 - aggregating, by said server in a time interval determined in accordance with a predefined criterion, said payload portions of said messages to create an aggregated payload;
 - forming an aggregated message using said aggregated payload; and
 - transmitting, by said server via said unicast network, said aggregated message to a recipient host computer belonging to said first message group.

- 1. A method for facilitating communications among a plurality of host computers over a network to implement a shared, interactive application, comprising the steps of:
 - (1) receiving a create message from one of the plurality of host computers, wherein said create message specifies a message group to be created;
 - (2) receiving join messages from a first subset of the plurality of host computers, wherein each of said join messages specifies said message group;
 - (3) receiving host messages from a second subset of said first subset of the plurality of host computers belonging to said message group, wherein each of said messages contains a payload portion and a portion that is used to identify said message group;
 - (4) aggregating said payload portions of said host messages received from said second subset of the plurality of host computers to create an aggregated payload;
 - (5) forming an <u>aggregated message</u> using said aggregated payload; and
 - (6) transmitting said aggregated message to said first subset of the plurality of host computers belonging to said message group;
 - wherein said aggregated message keeps the shared, interactive application operating consistently on each of said first subset of the plurality of host computers.

Roadmap

Overview

- '523 and '686 Patents
- Combination of Aldred and RFC 1692

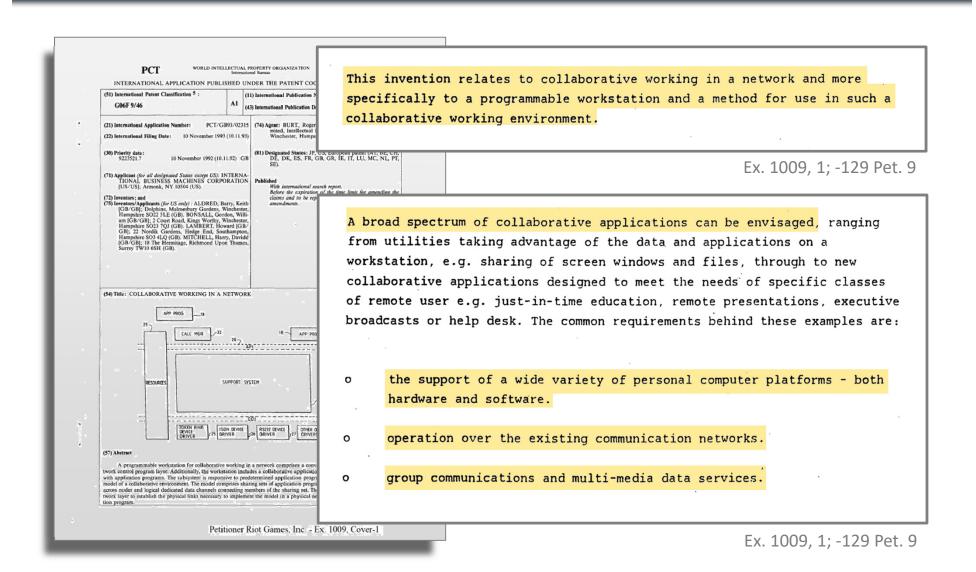
Independent Claim Disputes

- Patent Owner's Motivation to Combine Arguments
- Claim Construction of "Aggregated Payload" and "Aggregated Message"

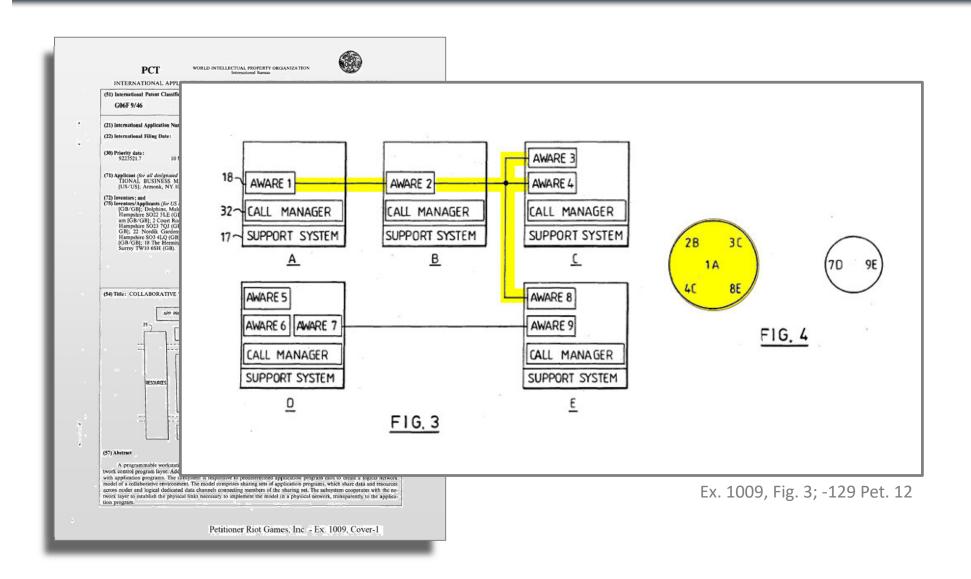
Dependent Claim Disputes

- Group Messaging Claims
- Ulrich Combination Message Groups and Echo Suppression

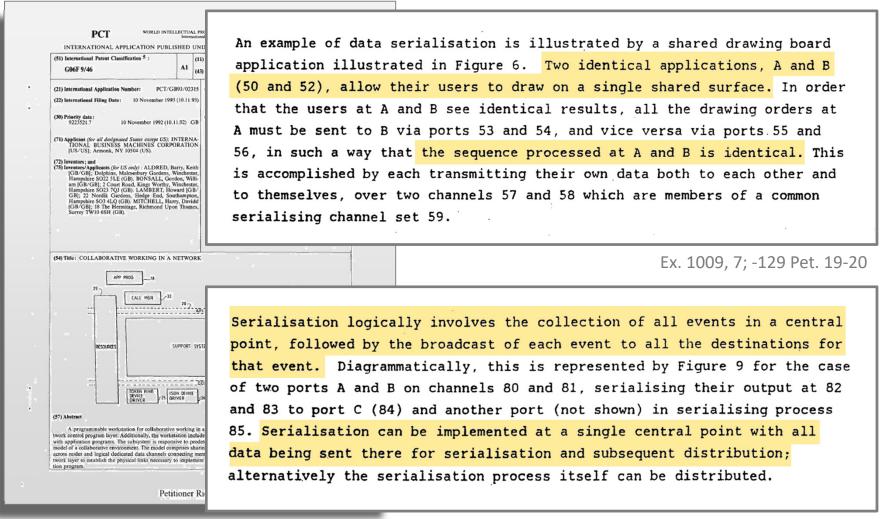
Aldred's Collaborative Working Environment



Aldred's Sharing Sets

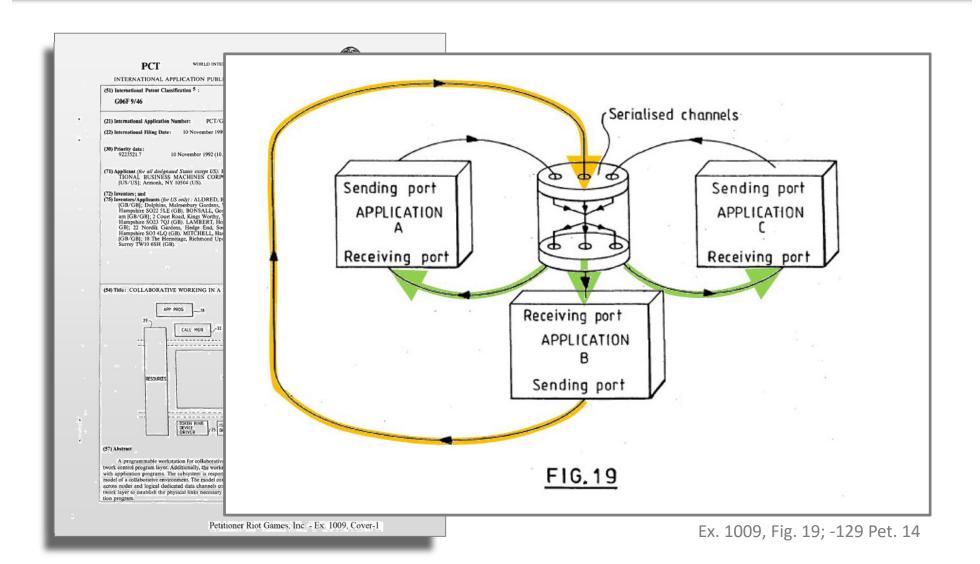


Aldred's Central Serialization Point

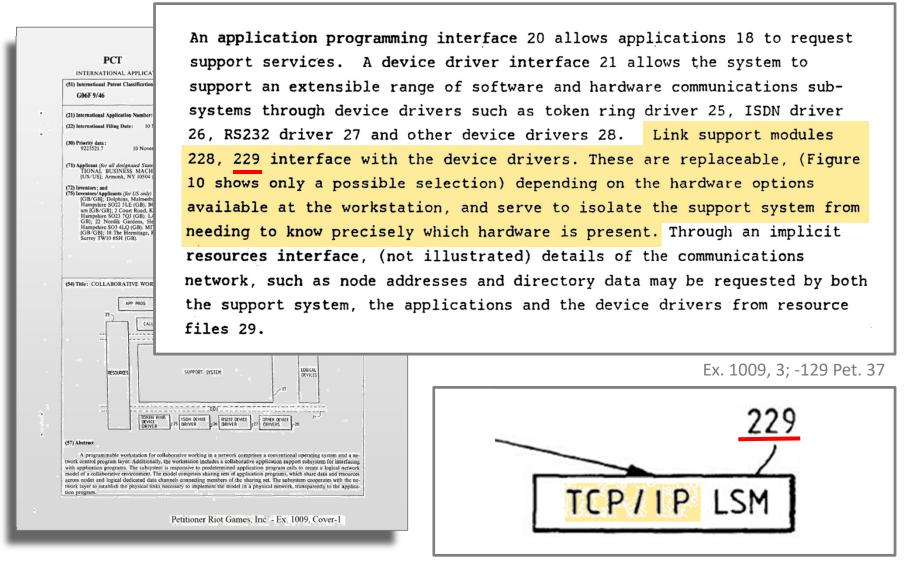


Ex. 1009, 9; -129 Pet. 20

Aldred's Central Serialization Point



Aldred's TCP/IP Networking Module



Ex. 1009, Fig. 10; -129 Pet. 37

RFC 1692 - Transport Multiplexing Protocol

TMux is designed to improve network utilization and reduce the interrupt load on hosts which conduct multiple sessions involving many short packets. It does this by multiplexing transport traffic Network Working Group onto a single IP datagram [2], thereby resulting in fewer, larger Request for Comments: 1692 Xvlogics. Category: Standards Track Sili packets. TMux is highly constrained in its method of accomplishing this task, seeking simplicity rather than sophistication. Transport Multiplexing Protocol (TMux) Ex. 1010, 2; -129 Pet. 36 Status of this Memo This document specifies an Internet standards tracks Internet community, and requests discussion and sug improvements. Please refer to the current edition Official Protocol Standards" (STD 1) for the standa and status of this protocol. Distribution of this Hence, a TMux message appears as: One of the problems with the use of terminal server | IP hdr | TM hdr | Tport segment | TM hdr | Tport segment | ... | number of small packets they can generate. Frequent packets are destined for only one or two hosts. which allows multiple short transport segments, ind application type, to be combined between a server a Where: Acknowledgments This specification is the result of the merger of t original TMux proposal which was the result of seve and related initiatives through IETF working groups originally proposed by Danny Cohen and Jon Postel i is a TMux mini-header and specifies the following TM hdr Applicability Statement Tport segment. The TMux protocol is intended to optimize the trans numbers of small data packets that are generated in many interactive Telnet and Rlogin sessions are co hosts on the network. In these situations, TMux ca network and host performance. TMux is not intended long streams composed of large blocks of data that Tport segment refers to the entire transport segment, including transmitted by such applications as FTP. The TMux protocol may be applicable to other situat packets are generated, but this was not considered transport headers. Ex. 1010, 3; -129 Pet. 36 Cameron, Crocker, Cohen & Postel TMux operates as an extension to the IP datagram protocol. Hence, it has no impact on most flow control mechanisms, since they operate at the transport layer -- above TMux. Petitioner Riot Gan Ex. 1010, 6; -129 Pet. 38

RFC 1692 - Message Construction

Network Working Group Request for Comments: 1692 Category: Standards Track

Xylo

Transport Multiplexing Protocol

Status of this Memo

This document specifies an Internet standards Internet community, and requests discussion a improvements. Please refer to the current ed Official Protocol Standards" (STD 1) for the and status of this protocol. Distribution of

Abstract

One of the problems with the use of terminal number of small packets they can generate. Pr packets are destined for only one or two host which allows multiple short transport segment application type, to be combined between a se

Acknowledgments

This specification is the result of the merge original TMux proposal which was the result o and related initiatives through IETF working originally proposed by Danny Cohen and Jon Po

Applicability Statement

The TMux protocol is intended to optimize the numbers of small data packets that are genera many interactive Telnet and Rlogin sessions a hoste on the network. In these situations, T network and host performance. TMux is not in long streams composed of large blocks of data transmitted by such applications as FTP.

The TMux protocol may be applicable to other packets are generated, but this was not const

Cameron, Crocker, Cohen & Postel

When a transport provider (e.g., TCP or UDP) sends a segment, TMux first removes the IP header (if present) and adds a TMux mini-header and the segment to the Multiplexed Message under construction for the host specified by the destination address of the segment.

When the first message to be transmitted is placed into the Multiplexed Message under construction, a timer is started. When the timer expires, the Multiplexed Message under construction is transmitted. This ensures that all segments available for sending before the timer expires are sent in a single Multiplexed Message. If, during construction of the Multiplexed Message, the buffer holding the message fills, the Multiplexed Message is transmitted immediately.

The delay time should be user configurable; a reasonable time is 20 to 30 milliseconds. The time period should be large enough to give a reasonable probability of sending multiple segments but not so large that the echo response time becomes a problem. This suggests that the upper limit for the timer is probably 1/10th second. As the cost of using timeouts on many systems is quite large, it is recommended that a single timer be used and that all TMux messages under construction are sent when the timer expires.

Petitioner Riot Games, Inc. - Ex. 1010, p. 1

Ex. 1010, 6; -129 Pet. 40

Obvious to use RFC 1692 in Aldred

UNITED STATES PATENT

BEFORE THE PATENT TRI

RIOT GAM Petitio

PALTALK HOI Patent (

Case No. IPR U.S. Patent No Issued: Octob Filed: Februa

Inventors: Jeffrey J. Rothschild, Marc

Title: SERVER-GROUP MESSAGIN APPLICA

PETITION FOR INTE

It would have been obvious to an Ordinary Artisan in 1995 to modify

Aldred's CSP to communicate with other nodes via RFC 1692's TMux protocol so as to "aggregat[e] ... said payload portions" as claimed. Ex. 1007, ¶145. Aldred and RFC 1692 are readily combinable; Aldred explains that "[t]he nature of the transport networks involved are totally hidden below the API." Ex. 1009, 28-29. "This means that an application is completely unaware of the network routing (eg direct or indirect), and the network types involved (eg single or multiple links, switched or fixed connections)." Id. Aldred also supports "TCP/IP," Id., Fig. 10 ("TCP/IP LSM 229"); see id., 3, and its exemplary networking software, IBM NetBIOS, could support TCP/IP. Id., 2-3; Ex. 1017.

-129 Pet. 36-37

Motivation to Combine Aldred with RFC 1692

UNITED STATES PATEN

BEFORE THE PATENT

RIOT C

PALTALK

Case No. U.S. Pater Issued: O Filed: Fe

Inventors: Jeffrey J. Rothschild,

Title: SERVER-GROUP MESSA APPI

PETITION FOR IN

RFC 1692 provides abundant motivation to incorporate TMux into Aldred.

small data packets," and "is designed to improve network utilization and reduce the

"The TMux protocol is intended to optimize the transmission of large numbers of

interrupt load on hosts which conduct multiple sessions involving many short

packets." Ex. 1010, 2; Ex. 1007, ¶147. RFC 1692 explains that "network and host

load could be greatly reduced if traffic from multiple users, destined for the same

host, could be sent in the same packet." Ex. 1010, 2 An Ordinary Artisan would

have therefore been motivated to use TMux in Aldred's scheme for these benefits.

Ex. 1007, ¶147.

-129 Pet. 38

Roadmap

Overview

- '523 and '686 Patents
- Combination of Aldred and RFC 1692

Independent Claim Disputes

- Patent Owner's Motivation to Combine Arguments
- Claim Construction of "Aggregated Payload" and "Aggregated Message"

Dependent Claim Disputes

- Group Messaging Claims
- Ulrich Combination Message Groups and Echo Suppression

Patent Owner's "Order" Argument

IPR2018-00129
Patent 5,822,523

UNITED STATES PATENT AND TRADEMARK

———
BEFORE THE PATENT TRIAL AND APPEAL B

———
RIOT GAMES, INC.,
Petitioner,

V.

PALTALK HOLDINGS, INC., Patent Owner.

> Case IPR2018-00129 Patent 5,822,523

PATENT OWNER'S RESPONSE TO PETITIO INTER PARTES REVIEW OF U.S. PATENT NO.

A POSITA would thus recognize that while the TMux timer has not expired, any packets removed from the queue at the CSP that should <u>not</u> be TMuxed, such as large packets, would be transmitted <u>immediately in a separate IP datagram</u> before the TMuxed message under construction that includes the small packets is transmitted. *Id.*, ¶ 75. This would cause the large packets to be transmitted <u>out of order</u> with respect to the small packets, because the small packets would be held back from transmission by the delay timer while the small packets reside in the TMux message under construction. *Id.* This would thus disrupt the required packet transmission order of the CSP as described in Aldred. *Id.*

-129 Resp. 25-26

Multiplexing "Large" Segments, e.g., FTP

Network Working Group Request for Comments: 1692 Category: Standards Track

network and host performance. TMux is not intended for multiplexing long streams composed of large blocks of data that are typically transmitted by such applications as FTP.

Silicon Graphics. Inc. Myricom J. Postel

Ex. 1010, 1; -129 Resp. 23

Transpor

Status of this Memo

This document specifies Internet community, and improvements. Please : Official Protocol Stan and status of this pro-

Abstract

One of the problems wit number of small packets packets are destined for which allows multiple s application type, to be

Acknowledgments

This specification is the result of the merger of two documents: the original TMux proposal which was the result of several discussions and related initiatives through IETF working groups; and IEN 90 [1]

originally proposed by

Applicability Statement

The TMux protocol is in numbers of small data p many interactive Telnet hosts on the network. network and host perfor long streams composed transmitted by such app

The TMux protocol may b packets are generated,

Cameron, Crocker, Cohen &

it is transmitted. It is also suggested that larger segments (e.g., those over 700 octets) should be sent as standard IP datagrams, and not multiplexed. This is to ensure that the delay caused by the TMux timer does not put a delay on those segments for which it is inadvisable. The size of the largest segments to be multiplexed should (if possible) be configurable.

Ex. 1010, 7; -129 Resp. 24

It is the responsibility of the sender to decide which segments should be TMux'd and which should not. For example, segments sent by FTP should not normally be multiplexed. In many situations, it may be sensible to restrict the sessions that can be multiplexed to just those involved in interactive traffic (Telnet and Rlogin) by examining the source and destination TCP port numbers. However, if a segment that would not normally be multiplexed is to be sent and a TMux message is already under construction, then the extra segment can be added to the TMux message under construction, and this complete message should be sent immediately, rather than waiting for the timer to expire.

Ex. 1010, 8-9; -129 Reply 3

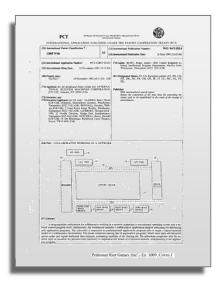
Dr. White's Reply Deposition

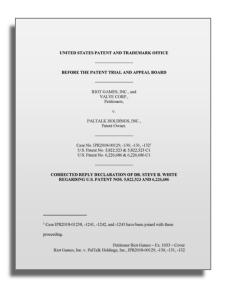
```
Page 1
        UNITED STATES PATENT AND TRADEMARK OFFICE
         BEFORE THE PATENT TRIAL AND APPEAL BOARD
4 RIOT GAMES, INC., and : Case IPR2018-00129
                              Case IPR2018-00130
    VALVE CORP.,
                        : 5,822,523 & 5,822,523 C1
              Petitioners, :
                            : Case IPR2018-00131
                           : 6,226,686 & 6,226,686 C1
    PALTALK HOLDINGS, INC.,: Case IPR2018-00132
             Patent Owner. : 6,226,686 & 6,226,686 C1
10
          ORAL DEPOSITION OF STEVE R. WHITE, PH.D.
                    New York, New York
                    December 19, 2018
13
14
15
18
19
    Reported by:
    Maureen Broderick, RPR
    JOB NO. 153003
                    TSG Reporting - Worldwide - 877-702-9580
                  PATENT OWNER PALTALK HOLDINGS - EX. 2005, PAGE 1
```

```
The premise of my question was, if the
packet received is equal to the largest message
size --
         And they're already messages in the
buffer.
         -- then wouldn't the larger packet --
wouldn't the combined packet exceed the maximum
transmittal unit?
          MR. BORDER: Objection. Form.
          THE WITNESS: If the message -- this is
     down in the weeds of a hypothetical that I
     haven't analyzed the detail, so I want to be
     very careful. So let me understand your
    premises.
          You have a message buffer -- a TMux
     message buffer. You have maximum transferable
     unit. The message buffer is less than or equal
     to, in size, the maximum transmissible unit.
     The buffer has messages in it. And the next
     message that you send is the size of the
     maximum transmissible unit.
          In that case, you couldn't multiplex the
    last message with the preceding messages,
     because it would be too large. And the, the
     ordering suggested here is that you would
     multiplex the smaller messages together, send
     those out, and then send your last message, the
     large message out after that.
```

Ex. 2005, 52:10-53:13; see -129 Sur-Reply 4

Aldred's Channels Maintain Packet Order





As illustrated in the schematic example of Figure 5, applications in a sharing set such as 40, 41 and 42 can establish data communication links with each other known as channels. Channels such as 43 and 44 are logically dedicated and uni-directional pipes, with application specified transmission characteristics. A channel is always defined by the sending application and it goes from a sending application to a receiving application. The ends of channels are known as ports; thus all channels have one sending port and one receiving port. A sending port such as 45 sends data packets down the channel; a receiving port such as 46 receives data packets from the channel in the order in which they were sent. There may be no direct mapping between the logical channel structure seen by the aware applications and the physical communication network in existence between the nodes.

Ex. 1009, 6; -129 Reply 5

channels would still provide their updates in order and maintain serialization. A person of ordinary skill in the art would read <u>Aldred</u>'s disclosure related to channels as maintaining the order of events at a logical level and regardless of the physical connection, which would resolve any out-of-order IP packet transmission.

Ex. 1053, ¶25; -129 Reply 5

RFC 793: TCP Reorders Out-of-Order Segments

RFC: 793

TRANSMISSION CONTROL PROTOCOL

DARPA INTERNET PROGRAM

September 1981

prepared for

Defense Advanced Research Projects Agency Information Processing Techniques Office 1400 Wilson Boulevard Arlington, Virginia 22209

by

Information Sciences Institute University of Southern California 4676 Admiralty Way Marina del Rey, California 90291

Reliability:

The TCP must recover from data that is damaged, lost, duplicated, or delivered out of order by the internet communication system. This is achieved by assigning a sequence number to each octet transmitted, and requiring a positive acknowledgment (ACK) from the receiving TCP. If the ACK is not received within a timeout interval, the data is retransmitted. At the receiver, the sequence numbers are used to correctly order segments that may be received out of order and to eliminate duplicates. Damage is handled by adding a checksum to each segment transmitted, checking it at the receiver, and discarding damaged segments.

As long as the TCPs continue to function properly and the internet system does not become completely partitioned, no transmission errors will affect the correct delivery of data. TCP recovers from internet communication system errors.

Ex. 1051, 4; -129 Reply 6

Petitioner Riot Games – Ex. 1051, Riot Games, Inc. v. PalTalk Holdings, Inc., IPR2018-00129, -130, -131, -132, Cover

Petition's Combination Uses TCP/IP

Paper No. 1

UNITED STATES PATENT AND TRADEMARK OFFICE

BEFORE THE PATENT TRIAL AND APPEAL BOARD

RIOT GAMES, INC., Petitioner,

v.

PALTALK HOLDINGS, INC., Patent Owner.

Case No. IPR2018-00129 U.S. Patent No. 5,822,523 Issued: October 13, 1998 Filed: February 1, 1996

Inventors: Jeffrey J. Rothschild, Marc P. Kwiatkowski, Daniel J. Samuel

Title: SERVER-GROUP MESSAGING SYSTEM FOR INTERACTIVE APPLICATIONS

PETITION FOR INTER PARTES REVIEW

It would have been obvious to an Ordinary Artisan in 1995 to modify

Aldred's CSP to communicate with other nodes via RFC 1692's TMux protocol so
as to "aggregat[e] ... said payload portions" as claimed. Ex. 1007, ¶145. Aldred
and RFC 1692 are readily combinable; Aldred explains that "[t]he nature of the
transport networks involved are totally hidden below the API." Ex. 1009, 28-29.
"This means that an application is completely unaware of the network routing (eg
direct or indirect), and the network types involved (eg single or multiple links,
switched or fixed connections)." Id. Aldred also supports "TCP/IP," Id., Fig. 10
("TCP/IP LSM 229"); see id., 3, and its exemplary networking software, IBM

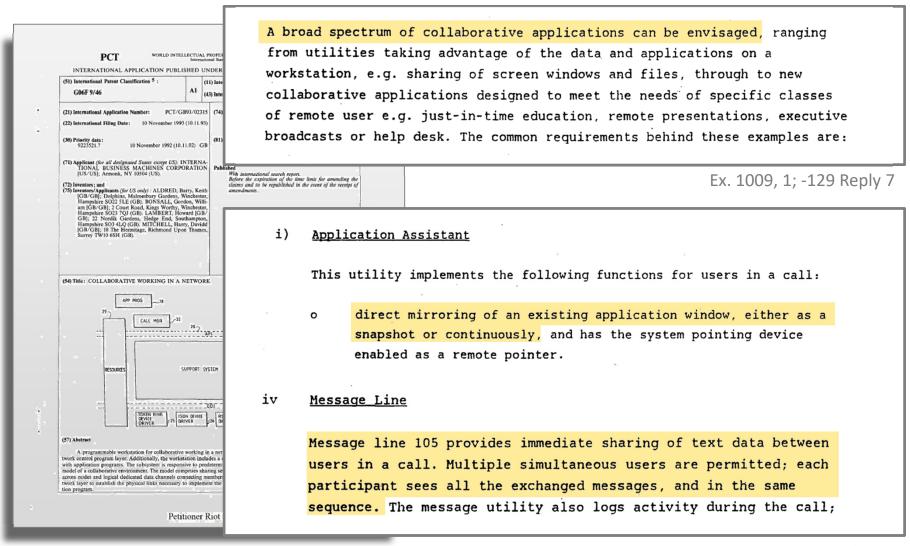
NetBIOS, could support TCP/IP. Id., 2-3; Ex. 1017.

-129 Pet. 36-37

As combined with RFC 1692, <u>Aldred</u>'s messages would be multiplexed via TMux ("aggregate[ed]") into a single packet ("aggregated payload") prior to transmission from the CSP to each member of the serialising channel set associated with the Sharing Set. Ex. 1007, ¶¶150-151. <u>RFC 1692</u> collects two or more TCP segments together into a multiplexed message. Ex. 1010, 2-3, 6. Those two segments are identified in the multiplexed message with TMux mini-headers. *Id.*, 3. When the multiplexed message is received by the receiver, the original two TCP segments and their respective headers are reconstructed from the information in the multiplexed message. *Id.*, 3. This would also meet Patent Owner's

-129 Pet. 39

Aldred Encompasses Small Packet Systems



Ex. 1009, 27-28; -129 Pet. 20; -129 Reply 7-8

Dr. White: RFC 1692 Reduces the Number of Packets

UNITED STATES PATENT AND TRADEMARK

BEFORE THE PATENT TRIAL AND APPEAL

RIOT GAMES, INC., Petitioner,

V.

PALTALK HOLDINGS, INC., Patent Owner.

Case Nos. IPR2018-00129, -00130 U.S. Patent No. 5,822,523

Case Nos. IPR2018-00131, -00132 U.S. Patent No. 6.226.686

DECLARATION OF DR. STEVE R. WHI' REGARDING U.S. PATENT NOS. 5,822,523 AND

Petitioner Riot Games, Inc.

that drawing orders and other events used to keep data consistent between applications, such as user input, would result in messages significantly smaller than the IP protocol supports, such that RFC 1692's methodology would improve Aldred's performance by reducing the number of packets. For example, a single character press in the conventional ASCII format of the time could be encoded in 8 bits (*i.e.*, one octet), meaning that a single aggregated message could contain numerous key presses or series of key presses from Aldred's nodes. As another example, a drawing order could be encoded in a pair of X and Y coordinates and some additional control information such as left-click or right-click, which one of ordinary skill would expect could be encoded in several octets.

Ex. 1007, ¶146; -129 Pet. 37-38

Combination = Known Elements + Known Functions

UNITED STATES PATENT AND TRADEMARK

BEFORE THE PATENT TRIAL AND APPEAL

RIOT GAMES, INC., Petitioner,

V.

PALTALK HOLDINGS, INC., Patent Owner.

Case Nos. IPR2018-00129, -00130 U.S. Patent No. 5,822,523

Case Nos. IPR2018-00131, -00132 U.S. Patent No. 6,226,686

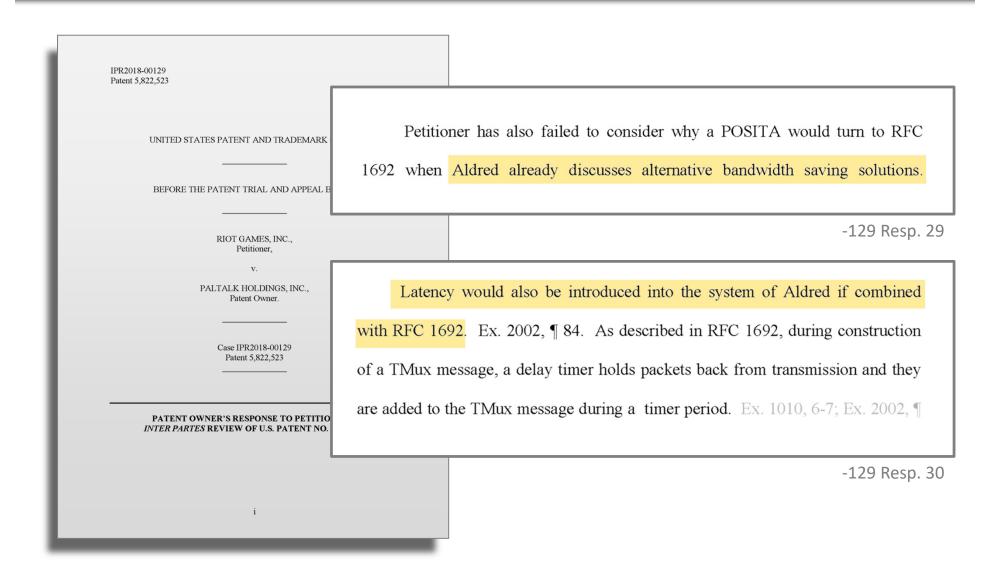
DECLARATION OF DR. STEVE R. WHI REGARDING U.S. PATENT NOS. 5,822,523 AND

Petitioner Riot Games, Inc

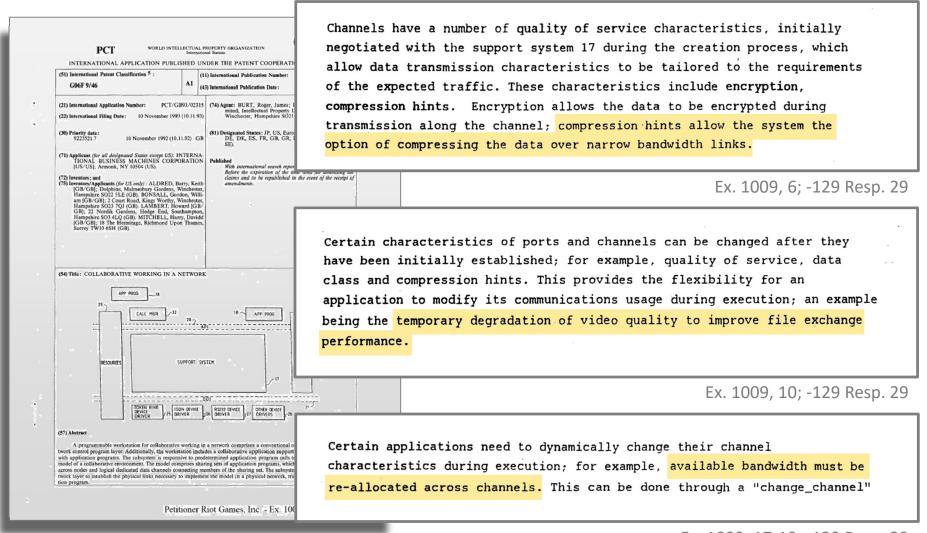
30. Incorporating the TMux protocol into <u>Aldred</u> would simply involve using the TMux-enhanced IP protocol for <u>Aldred</u>'s transport mechanism, and would have been well-within the abilities of an ordinary artisan. This arranges known elements, each performing the same function it had been known to perform (<u>Aldred</u>'s logical communication scheme using <u>RFC 1692</u>'s multiplexing IP protocol for underlying transport), to yield no more than one would expect from such an arrangement (RFC 1692's functionality multiplexing outgoing messages from Aldred's central serialisation point). Such a modification would also have been the product of ordinary skill and common sense, and would have been obvious to try because, as explained above, one of ordinary skill would have had good reason to pursue the known options within his or her technical grasp.

Ex. 1007, ¶149; -129 Pet. 38-39

Patent Owner's Non-Obviousness Arguments



Aldred's "Bandwidth Saving Techniques"



Ex. 1009, 17-18; -129 Resp. 29

Dr. White: Aldred's Quality of Service (QOS) Are Flexible

UNITED STATES PATENT AND TRADEMARK OFFICE

BEFORE THE PATENT TRIAL AND APPEAL BOARD

RIOT GAMES, INC., and VALVE CORP., Petitioners,

V

PALTALK HOLDINGS, INC., Patent Owner.

Case No. IPR2018-00129, -130, -131, -132¹
U.S. Patent No. 5,822,523 & 5,822,523 C1
U.S. Patent No. 6,226,686 & 6,226,686 C1

CORRECTED REPLY DECLARATION OF DR. STEVE R. WHITE REGARDING U.S. PATENT NOS. 5,822,523 AND 6,226,686

¹ Case IPR2018-01238, -1241, -1242, and -1243 have been joined with these proceeding.

Petitioner Riot Games – Ex. 1053 – Cover Riot Games, Inc. v. PalTalk Holdings, Inc., IPR2018-00129, -130, -131, -132

Dr. Almeroth suggests that "introducing the TMux protocol in the system of Aldred could cause quality of service characteristics for the logical channels to not be met." Ex.2002, ¶84. Initially, I note Aldred's quality of service characteristics are not fixed and can be "tailored" to each application's needs. See Ex. 1009, 6 (quality of service characteristics "negotiated" "which allow data transmission characteristics to be tailored to the requirements of the expected traffic"); id., 10 ("Connection is permitted between channels in different channel sets ... having different quality of service characteristics."); id., 16-18 (describing "flexible," "fixed," and "dynamic" quality of service needs). In fact, Aldred states that "an implicitly created channel will attempt to have characteristics such that it can transport satisfactorily any data packets expected to be sent down any one of the pre-existing channels from that port." Id., 19. Further, Aldred does not require specifying any particular QoS parameter. See, e.g., id., 18 ("An application can specify quality of service characteristics when creating a channel ... " (emphasis added)).

Ex. 1053, ¶35; -129 Reply 10

Dr. White: TMux Complements Other Technologies

```
Page 1
                             S. WHITE
          UNITED STATES PATENT AND TRADEMARK OFFICE
           BEFORE THE PATENT TRIAL AND APPEAL BOARD
                        RIOT GAMES, INC.
                          Petitioner
                     PALTALK HOLDINGS, INC.
                          Patent Owner
                       Case IPR2018-00129
                        Patent 5,822,523
                      Case IPR2018-00130
                       Patent 5,822,523
                      Case IPR2018-00131
                       Patent 6,226,686
                       Case IPR2018-00132
                        Patent 6,226,686
                DEPOSITION OF DR. STEVE WHITE
                        Washington, D.C.
23
                         July 24, 2018
          Reported by: Mary Ann Payonk; Job No. 144847
25 Job No. 144847
                     TSG Reporting - Worldwide 877-702-9580
                                     PATENT OWNER EXHIBIT 2004 - PAGE 1
```

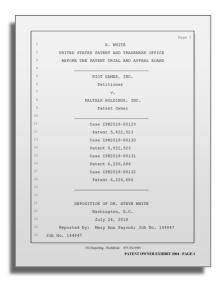
- Q. So would compression provide the same results as the TMux protocol?
 - A. No.
- Q. Why would one skilled in the art choose RFC 1692 over Denzer?

MR. DILLON: Objection to form.

A. It -- I don't think I would hold that one would choose one over the other. They're complementary technologies. If you have lots and lots of little packets that will all fit within a single IP packet, then multiplexing is an obvious thing to do. If you had a lot of redundancy in your data, then compression is an obvious thing to do. If you have both, then they're both obvious things to do.

Ex. 2004, 54:21-55:20; -129 Reply 9

TMux: an Engineering Tradeoff & Configurable



If you have a lot of -- if you have a whole lot of packets going on the network bandwidth, that can also adversely affect latency. So there's an engineering tradeoff between holding them back and sending them in larger packets, and sending them as soon as you possibly can.

Ex. 2004, 50:23-51:10; -129 Reply 10



The delay time should be user configurable; a reasonable time is 20 to 30 milliseconds. The time period should be large enough to give a reasonable probability of sending multiple segments but not so large that the echo response time becomes a problem. This suggests that the upper limit for the timer is probably 1/10th second. As the cost of using timeouts on many systems is quite large, it is recommended that a single timer be used and that all TMux messages under construction are sent when the timer expires.

Ex. 1010, 6; -129 Reply 9-10

Roadmap

Overview

- '523 and '686 Patents
- Combination of Aldred and RFC 1692

Independent Claim Disputes

- Patent Owner's Motivation to Combine Arguments
- Claim Construction of "Aggregated Payload" and "Aggregated Message"

Dependent Claim Disputes

- Group Messaging Claims
- Ulrich Combination Message Groups and Echo Suppression

"Aggregated Message" and "Aggregated Payload"

Term	Patent Owner's Construction	Petitioners' Construction
"aggregated message"	One or more messages containing a single transport layer message header, destination data, and data items from an aggregated payload	No "transport layer" header requirement, so no construction necessary.
"aggregated payload"	A collection of two or more data items that does not include transport layer headers	No "transport layer" header requirement, so no construction necessary.

-129 Resp. 4, 13; -129 Reply 11-12

Ex. 1010, 3; -129 Resp. 34

New "Transport Layer" Header Requirement

Term & Case	Microsoft Case (2007)	Sony Case & EPX (2010-11)	Riot D. Ct. Term Exch. (Jan. 2018)	Riot IPR Response (Feb. 2018)
"aggregated message"	No	No	No	Yes
"payload"	No	No	No	No
"aggregating"	No	No	No	No
"message"	No	No	No	No

-129 Reply 16

Context of '523 Patent, Claim 1

1. A method for providing group messages to a plurality of host computers connected over a unicast wide area communication network, comprising the steps of:

providing a group messaging server coupled to said network, said server communicating with said plurality of host computers using said unicast network and maintaining a list of message groups, each message group containing at least one host computer;

sending, by a plurality of host computers belonging to a first message group, messages to said server via said unicast network, said messages containing a payload portion and a portion for identifying said first message group;

aggregating, by said server in a time interval determined in accordance with a predefined criterion, said payload portions of said messages to create an aggregated payload;

forming an aggregated message using said aggregated payload; and

transmitting, by said server via said unicast network, said aggregated message to a recipient host computer belonging to said first message group.

'523 Patent – No "Layer" Requirement for Aggregating

1. A method for providing group messages to a plurality of host computers connected over a unicast wide area communication network, comprising the steps of:

. . .

aggregating, by said server in a time interval determined in accordance with a predefined criterion, said payload portions of said messages to create an aggregated payload;

forming an aggregated message using said aggregated payload; and

. . .

- 6. The method of claim 1 wherein said network is Internet and said server communicates with said plurality of host computers using a session layer protocol.
- 32. The method of claim 1, wherein said sending and said transmitting are performed by an upper-level protocol implemented above a transport layer protocol of said unicast network, wherein said transport layer protocol is TCP/IP.

"Payloads" Can Include Transport Headers





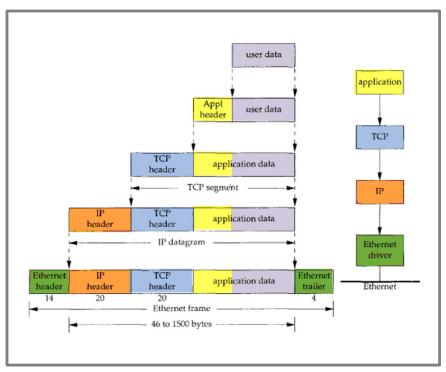
4 transport protocols TCP and UDP. UDP provides datagram delivery services to applications that does not guarantee reliable or in-order delivery of the datagrams. TCP is a connection oriented service to applications that does provide reliable delivery of a data stream. It handles division of the stream into packets and ensures reliable, in-order delivery. See the Internet Society RFCs: RFC-791 "Internet Protocol", RFC-793 "Transmission Control Protocol" and RFC-1180 "A TCP/IP Tutorial". IP, TCP and UDP are unicast protocols: packets, streams or datagrams are transmitted from a source to a single destination.

Ex. 1001, 3:24-50; -129 Reply 13

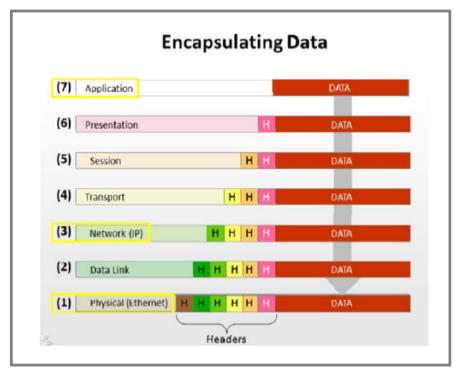
For example, a TCP module would call on the internet module to take a TCP segment (including the TCP header and user data) as the data portion of an internet datagram. The TCP module would provide the addresses and other parameters in the internet header to the internet module as arguments of the call. The internet module would then create an internet datagram and call on the local network interface to transmit the internet datagram.

Ex. 1011, 1; -129 Reply 13-14

Dr. Almeroth: "Payloads" Can Include Transport Headers



Ex. 1056, ¶68; -129 Reply 14



Ex. 1058, 3; -129 Reply 14-15

Dr. Almeroth: "Payloads" Can Include Transport Header

UNITED STATES PATENT AND TRADEMARK OFFICE

BEFORE THE PATENT TRIAL AND APPEAL BOARD

INTEL CORP. and CAVIUM, INC.,

Petitioners,

V.

ALACRITECH INC.,

Patent Owner.

Case IPR2017-01391 U.S. Patent 7,237,036

CORRECTED PATENT OWNER'S EXHIBIT 2026 DECLARATION OF KEVIN ALMEROTH, Ph.D.

Alacritech Exhibit

Petitioner Riot Games – Ex. 1056. Riot Games Inc. v. PalTalk Holdings Inc. IPR 2018-00129 -130 -131 -132 Cover

68. This process of adding a layer header to the data from the preceding layer is sometimes referred to as "encapsulation" because the data and layer header is treated as the data for the immediately following layer, which, in turn, adds its own layer header to the data from the preceding layer. Each layer is generally not aware of which portion of the data from the preceding layer constitutes the layer header or the user data; as such, each layer treats the data it receives from the preceding layer as some generic payload.

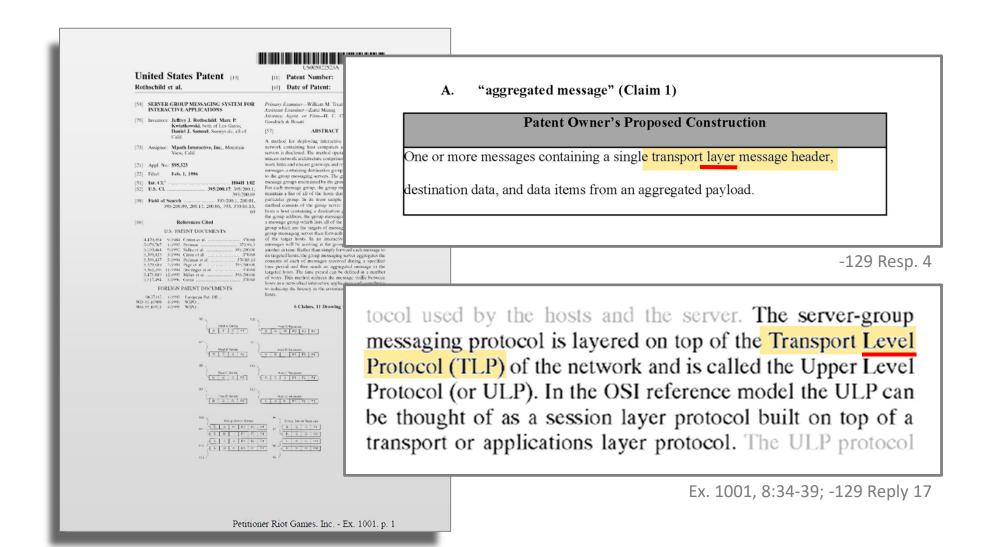
Ex. 1056, ¶68; Ex. 1053, ¶8; -129 Reply 15

receiving host). (Ex. 1001 at 4:34-39.) Each layer of the receiving host recognizes and manipulates only the headers associated with that layer, since to that layer the higher layer header data is included with and indistinguishable from the payload data. "Multiple interrupts, valuable central processing unit (CPU) processing time

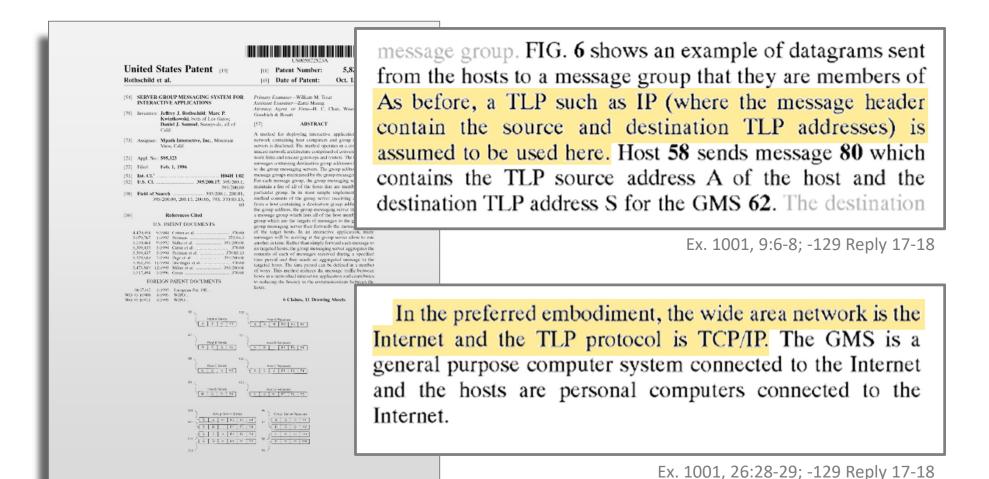
Ex. 1056, ¶69; Ex. 1053, ¶8; -129 Reply 15

Cavium, who filed a Petition in Case IPR2017-01718, has been joined as a petitioner in this proceeding.

"Transport Level Protocol": Coined Term in Patent



"Transport Level Protocol": "Such as IP" or "TCP/IP"



Petitioner Riot Games, Inc. - Ex. 1001, p. 1

"Transport <u>Level</u> Protocol": Dr. Almeroth Does Not Equate to Transport <u>Layer</u> Protocol

UNITED STATES PATENT AND TRADEMA Is the transport layer protocol described BEFORE THE PATENT TRIAL AND APP 0. RIOT GAMES, INC. here, is it the same as the transport level protocol Petitioner described in this patent? PALTALK HOLDINGS, INC Patent Owner I haven't tried to answer that question. Case IPR IPR 2018-00129 I'd have to look through the specification and see. Patent 5,822,523 Case IPR 2018-00130 Patent 5,822,523 Sitting here right now, I don't recall having an Case IPR 2018-00131 Patent 6,226,686 opinion one way or the other whether it's the same or Case IPR 2018-00132 Patent 6,226,686 not. Ex. 1052, 13:12-19; -129 Reply 19 DEPOSITION OF KEVIN C. ALMEROTH, Ph.D. Century City, California November 8, 2018 Reported by: Nikki Roy, CSR No. 3052 Job No. 150165 TSG Reporting - Worldwide 877-702-9580
Petitioner Riot Games - Ex. 1052. Riot Games, Inc. v. PalTalk Holdings, Inc., IPR2018-00129, -130, -131, -132, p. 1

"Transport <u>Level</u> Protocol": Dr. Almeroth Does Not Equate to Transport <u>Layer</u> Protocol

```
Page 1
            UNITED STATES PATENT AND TRADEMARK OFFICE
             BEFORE THE PATENT TRIAL AND APPEAU BOARD
                          RIOT GAMES, INC.
                             Petitioner
                       PALTALK HOLDINGS, INC.
                           Patent Owner
                      Case IPR IPR 2018-00129
                          Patent 5,822,523
                        Case IPR 2018-00130
                          Patent 5,822,523
14
                       Case IPR 2018-00131
15
                          Patent 6,226,686
                        Case IPR 2018-00132
                         Patent 6,226,686
19
              DEPOSITION OF KEVIN C. ALMEROTH, Ph.D.
                     Century City, California
                          November 8, 2018
    Reported by: Nikki Roy, CSR No. 3052
   Job No. 150165
                     TSG Reporting - Worldwide 877-702-9580
Petitioner Riot Games - Ex. 1052.
       Riot Games, Inc. v. PalTalk Holdings, Inc., IPR2018-00129, -130, -131, -132, p. 1
```

Q. All right. Well, go to column 9, line 5. It says (reading):

"As before, a transport layer protocol, such as IP, where the message header contains the source and destination TLP addresses is assumed to be used here."

So the patent certainly refers to IP as a

transport layer protocol, correct?

- A. It refers to it as a TLP and uses it as an example.
- Q. Okay. So you don't think the patent is incorrect here?
- A. I mean, the patent can characterize a TLP in any way it wants. IP is generally considered part of layer 3 in the network layer, and so is in the transport layer in the OSI stack. But you can define a transport layer as distinctive over an upper layer protocol, and you can do it the way that it's described here.

Ex. 1052, 79:21-80:15; -129 Reply 17-18

Patent Owner's "Disclaimer" Argument

IPR2018-00129 Patent 5,822,523

UNITED STATES PATENT AND TRADEMARK OFFICE

BEFORE THE PATENT TRIAL AND APPEAL BOARD

RIOT GAMES, INC., Petitioner,

V.

PALTALK HOLDINGS, INC., Patent Owner.

> Case IPR2018-00129 Patent 5,822,523

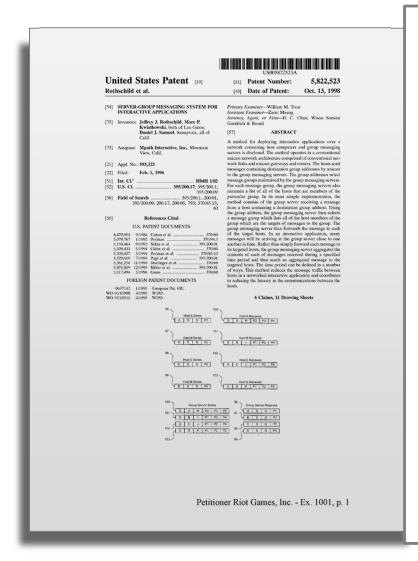
PATENT OWNER'S RESPONSE TO PETITION FOR INTER PARTES REVIEW OF U.S. PATENT NO. 5,822,523

The specification makes it clear that there is only one transport message header (including fields 123, 124, and 125) in an aggregated message, and clearly explains that data reduction is significant because there is only one transport message header. Ex. 2002, ¶¶ 53-54. "Where the general summary or description of the invention describes a feature of the invention . . . and criticizes other products . . . that lack that same feature, this operates as a clear disavowal . . ."

**Astrazeneca AB, Aktiebolaget Hassle, KBI-E, Inc. v. Mut. Pharm. Co., 384 F.3d 1333, 1340 (Fed. Cir. 2004). The '523 Patent both describes the advantages of transmitting a single transport layer message header, and criticizes other embodiments. Ex. 1001, 24:23-28, 10:40-44. To state that the specification does not limit aggregated payloads from including transport layer headers completely ignores the specific disclosure of the '523 Patent. Therefore, Patent Owner proposes adoption of this claim construction for "aggregated message."

-129 Resp. 12-13

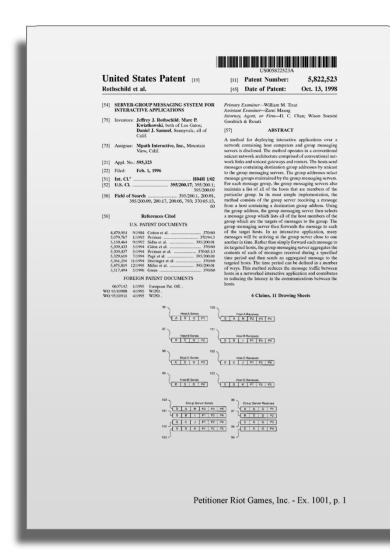
No "Disclaimer" - Three Possible Benefits of "Aggregation"



The effect of aggregation will be to greatly reduce the total message rate received by the hosts. A single message to a host will be able to carry multiple payload items received from the other hosts during the aggregation period. This fits very well the interactive applications of this invention where groups of hosts will be sending messages to all the other hosts in the group at a periodic rate. Aggregation will be very effective in collecting together all of the messages from all of the other hosts into a single message for each member of the group. The reduces processing at each receiving host since a single message will be received rather than many separate messages. Aggregation will also reduce the total data rate to the hosts since aggregation eliminates the need for separate message headers for each payload item. The savings will be significant for small payload items since there will be only one message header comprising fields 123, 124 and 125 for multiple payload items. In cases where a group of hosts are sending messages to the group at a periodic rate, it is often the case in many interactive applications that the data being sent by each host to the group is very similar to the messages sent by the other hosts. This affords the opportunity within an aggregated payload of multiple payload items to apply a data compression method across the multiple data elements of the payload elements. A

Ex. 1001, 24:12-35; -129 Reply 21

Any "Disclaimer" Not Specific to "Transport Layer"



```
1333, 1340 (Fed. Cir. 2004). The '523 Patent both describes the advantages of transmitting a single transport layer message header, and criticizes other embodiments. Ex. 1001, 24:23-28, 10:40-44. To state that the specification does
```

-129 Resp. 13

separate messages. Aggregation will also reduce the total data rate to the hosts since aggregation eliminates the need for separate message headers for each payload item. The savings will be significant for small payload items since there will be only one message header comprising fields 123, 124 and 125 for multiple payload items. In cases where a

Ex. 1001, 24:23-28; -129 Resp. 12; -129 Reply 22

sage. The received message is longer and contains multiple payloads, but this is a significant improvement over receiving multiple messages with the wasted overhead of multiple message headers and message processing time. Overall the invention has dramatically reduced the amount of data that must be sent and received by each host. Since the bit rate

Ex. 1001, 10:40-44; -129 Resp. 12; -129 Reply 22

Board Identified Headers in "Payload"

Trials@uspto.gov Tel: 571-272-7822 Paper 11 Entered: May 15, 2018

UNITED STATES PATENT AND TRADEN

BEFORE THE PATENT TRIAL AND APP

RIOT GAMES, INC., Petitioner.

V.

PALTALK HOLDINGS, INC Patent Owner.

Case IPR2018-00129 Patent 5,822,523 & 5,822,523 C

Before THU A. DANG and KARL D. EASTHOM, Administrative Patent Judges.

EASTHOM, Administrative Patent Judge.

DECISION
Institution of Inter Partes Review
35 U.S.C. 314(a)

As Patent Owner argues, Figure 9 of the specification does not include a TLP header in each payload packet of the aggregated payload. *See* Prelim. Resp. 7–8. Nonetheless, as noted above, headers, such as headers 117 and 118, or 120 and 122, appear in each payload. *See* Ex. 1001, 23:11–12 ("Each payload item in a message queue will contain a ULP source address, a data length and the data to be sent."). Even though an embodiment strips out a TLP header from a "message," it also looks up a TLP header of the

-129 Inst. Dec. 12

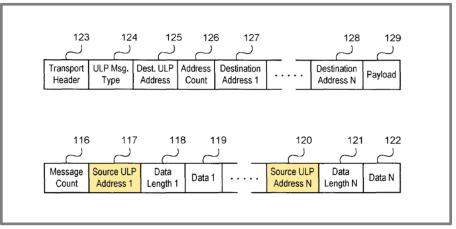
¹The Petition challenges original claims and claims issued pursuant to an *exparte* reexamination certificate. *See* Ex. 1001.

'523 Patent's Preferred Embodiment Includes Multiple "Transport Layer" Headers

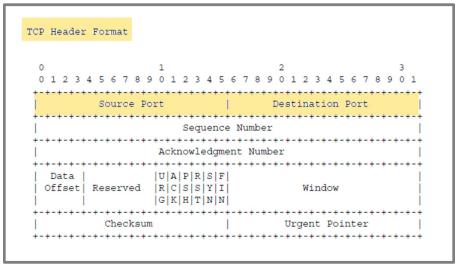


In the preferred embodiment, the wide area network is the Internet and the TLP protocol is TCP/IP. The GMS is a general purpose computer system connected to the Internet and the hosts are personal computers connected to the Internet.

TCP/IP provides an number of advantages that provide for a more efficient applications interface on the hosts 151. TCP/IP supports the concept of source and destination port numbers in its header. The ULP can make use of the port numbers to identify source and destination ULP connections. Most ULP send messages will be from hosts to a implicit ULP group addresses and most ULP receive messages will be from the implicit ULP addresses to the ULP host addresses. All of these and the ULP message type field can represented by source and destination port addresses within the TCP/IP header. This means that for most ULP messages, the ULP message encapsulated within the TCP/IP message need only contain the payload. There is the slight complication of the aggregated ULP receive messages sent from a ULP server process to a hosts. Here the destination port will be the host the source port will be for the implicit LJLP group address and the payload will still contain the source host ULP addresses in each the payload items.



Ex. 1001, Fig. 9; -129 Reply 19-20



Ex. 1051, 15; Ex. 1052, 43:7-44:13; -129 Reply 6, 19-20

Ex. 1001, 26:28-50; -129 Reply 19-20

Roadmap

Overview

- '523 and '686 Patents
- Combination of Aldred and RFC 1692

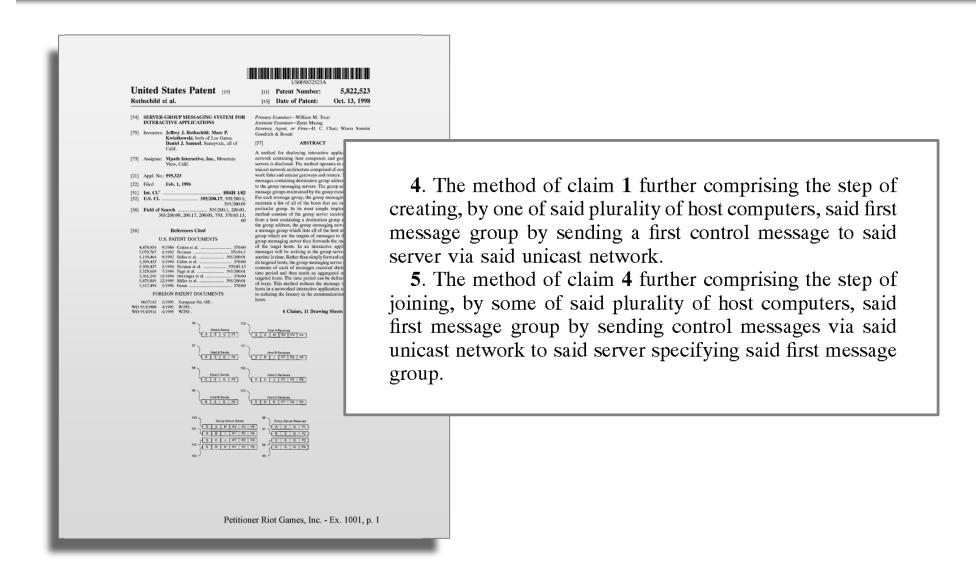
Independent Claim Disputes

- Patent Owner's Motivation to Combine Arguments
- Claim Construction of "Aggregated Payload" and "Aggregated Message"

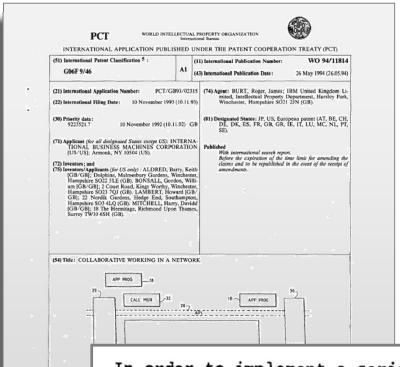
Dependent Claim Disputes

- Group Messaging Claims
- Ulrich Combination Message Groups and Echo Suppression

'523 Patent Claims 4-5



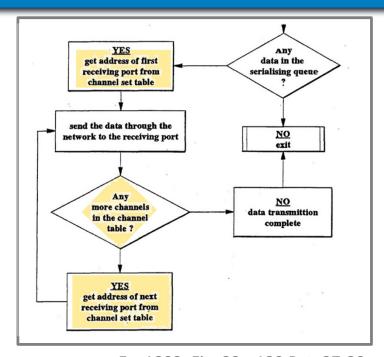
Aldred's CSP Maintains the Channel Set Table



(57) Abstract

A program twork control pr with application model of a colla

across nodes and twork layer to er



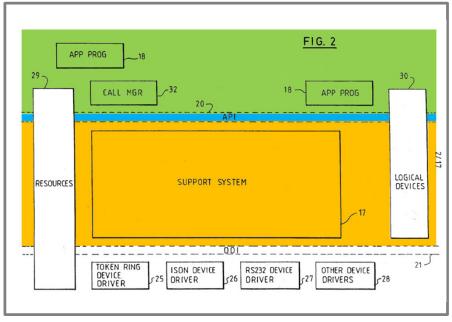
Ex. 1009, Fig. 22; -129 Pet. 27-29

In order to implement a serialised channel, as illustrated in the flow diagrams of Fig. 22, the subsystem uses a similar technique of establishing a channel set table with the addresses of all receiving ports. However, it also needs to maintain a serialising queue for the channel in which the data items to be serialised are loaded from the sending ports and held in the order in which it is desired to transmit them to all receiving ports.

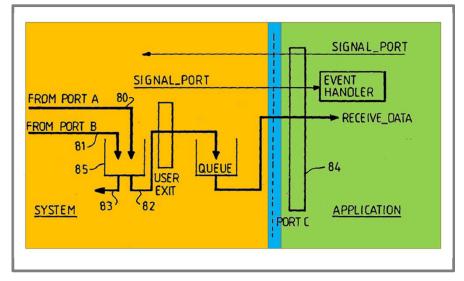
Ex. 1009, 51; -129 Pet. 27-29

Dr. White: Aldred's CSP is on a Sharing Set Node

"receive_data." One of ordinary skill in the art would understand this disclosure, in light of <u>Aldred</u>'s disclosure as a whole, as illustrating a node in <u>Aldred</u>'s scheme, with the serialising process 85 being located in the support "system" (orange) which interfaces with the local applications (green) on the node by way of an <u>API (blue)</u>:



Ex. 1007, ¶95; -129 Pet. 21-22



Ex. 1009, Fig. 9; -129 Pet. 21-22

Ex. 1009, Fig. 2; -129 Pet. 21-22

Petition: Aldred's CSP Manages Group Membership

Paper No. 1

UNITED STATES PATENT AND TRADEMARK OFFICE

BEFORE THE PATENT TRIAL AND APPEAL BOARD

RIOT GAMES, INC., Petitioner,

V.

PALTALK HOLDINGS, INC., Patent Owner.

Case No. IPR2018-00129 U.S. Patent No. 5,822,523 Issued: October 13, 1998 Filed: February 1, 1996

Inventors: Jeffrey J. Rothschild, Marc P. Kwiatkowski, Daniel J. Samuel

Title: SERVER-GROUP MESSAGING SYSTEM FOR INTERACTIVE APPLICATIONS

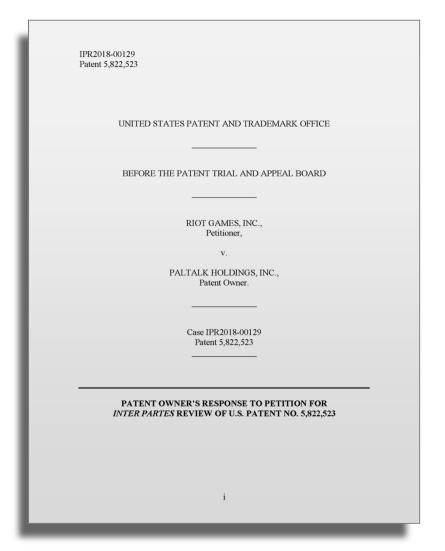
PETITION FOR INTER PARTES REVIEW

"Serialisation can be implemented at a **single central point** with all data being sent there for serialisation and subsequent distribution" Ex. 1009, 9 (emphasis added). In the context of <u>Aldred</u> (*see id.*, 2-6, Figs. 1, 3), this Central Serialisation Point ("CSP") is part of the support system of one of the participating nodes. *Id.*, Figs. 2, 9; Ex. 1007, ¶65-66. The CSP would "maintain a serialising queue for the channel in which the data items to be serialised are loaded from the sending ports and held in the order in which it is desired to transmit them to all receiving ports." Ex. 1009, 51. It is maintained and updated automatically: "New members may easily be added to the group with the necessary data channels being established and serialised automatically by the underlying system." *Id.*, 50, Fig.

22. Join requests and leave requests would modify this channel set in the same way as any other node in <u>Aldred</u>'s scheme. *Id.*, 5, 31-39; Ex. 1007, ¶67.

-129 Pet. 14-15

Patent Owner's Argument – "a CSP Is Not Involved"



Petitioner equates the CSP of Aldred with the group messaging server of the '523 Patent. Pet. at 25. However, Aldred does not disclose that "share app" or "unshare app" requests are sent to a CSP in order for a node/application to create, join, or leave a sharing set. Aldred also does not disclose that in response to a received "share_app" request, the CSP adds nodes to the sharing set, removes nodes from the sharing set, or stores or removes information regarding the nodes. While a CSP can operate on a node, "share app" and "unshare app" requests are not sent to the CSP to carry out any of the recited functions in the claims. Essentially, to create, join, leave, or remove a sharing set in Aldred, a host simply sends a request to one or more other hosts; a CSP is not involved in these processes of Aldred cited by Petitioner. Ex. 1009, 15; Ex. 2002, ¶ 137. Aldred further explains that, to add new members to a sharing set, the underlying system of the new node automatically establishes the channels to be used for the sharing set. Ex. 1009, 49. Aldred does not disclose that a server or CSP establishes the channels, or that the CSP adds or removes nodes in the sharing set, removes sharing sets, or stores or removes information regarding the nodes to create or edit a sharing set. Ex. 2002, ¶ 137.

-129 Resp. 53-54

Dr. Almeroth: Channel Set Table

```
Page 1
            UNITED STATES PATENT AND TRADEMARK OFFICE
             BEFORE THE PATENT TRIAL AND APPEAU BOARD
                          RIOT GAMES, INC.
                             Petitioner
                       PALTALK HOLDINGS, INC.
                           Patent Owner
                      Case IPR IPR 2018-00129
                          Patent 5,822,523
                        Case IPR 2018-00130
                          Patent 5,822,523
14
                        Case IPR 2018-00131
                          Patent 6,226,686
                        Case IPR 2018-00132
                         Patent 6,226,686
              DEPOSITION OF KEVIN C. ALMEROTH, Ph.D.
                     Century City, California
                          November 8, 2018
    Reported by: Nikki Roy, CSR No. 3052
25 Job No. 150165
                     TSG Reporting - Worldwide 877-702-9580 Petitioner Riot Games – Ex. 1052.
        Riot Games, Inc. v. PalTalk Holdings, Inc., IPR2018-00129, -130, -131, -132, p. 1
```

```
Q. What is the channel set table that you describe in paragraph 166?

A. So turn to page 51. It's talking about the serialized channel. It says (reading):

"As illustrated in the flow diagram of Figure 22, the subsystem uses a similar technique of establishing channel set table with the addresses of all the receiving ports."

So it appears to have all the addresses of the receiving ports I suspect it's described earlier as well.
```

Ex. 1052, 101:10-22; -129 Reply 24

```
Q. In the serialization process it sends to all members of that channel set?

A. That's what it describes. You send data in turn to each receiving port in the channel.

Q. And in order to send data to every member of the channel set table, you'd agree that the channel set table must exist, right?

A. If you're using the channel set table to identify the addresses of the receiving ports as it's described here, then you would need a channel set table.
```

Ex. 1052, 103:15-25; -129 Reply 24-25

Roadmap

Overview

- '523 and '686 Patents
- Combination of Aldred and RFC 1692

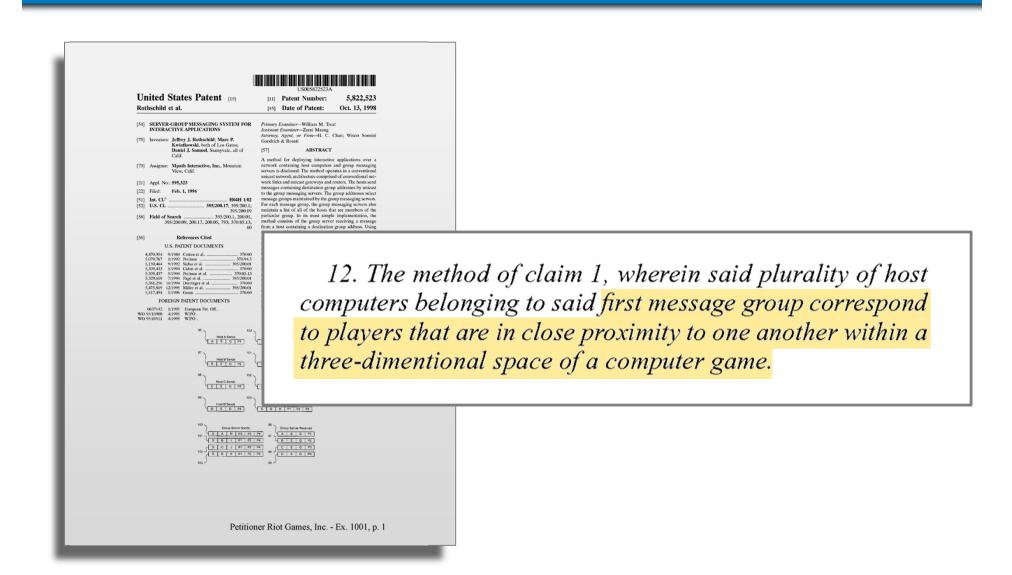
Independent Claim Disputes

- Patent Owner's Motivation to Combine Arguments
- Claim Construction of "Aggregated Payload" and "Aggregated Message"

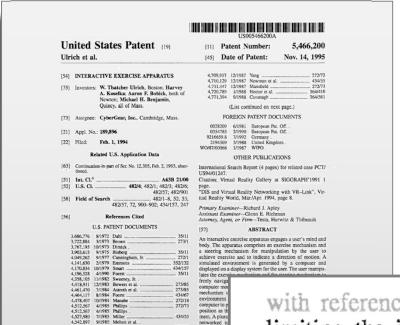
Dependent Claim Disputes

- Group Messaging Claims
- Ulrich Combination Message Groups and Echo Suppression

'523 Patent Claim 12



Overview of Ulrich



434/55

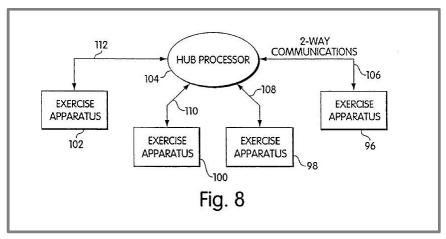
273/85

SWIJLATED 3-D ENVRONMENT

NETWORK

4,542,897 9/1985 Melton et al. .

4,572,509 2/1986 Sitrick 4,616,829 10/1986 Smack, Sr. et al. 4,630,817 12/1986 Buckley



Ex. 1012, Fig. 8; -130 Pet. 53

with reference to FIG. 7. The hub 104 is responsible for limiting the information directed to each apparatus in the large-scale direct network of FIG. 8. The hub 104 can ensure, for example, that each apparatus only gets (parameter) updates about other users in the same general area of the simulated environment.

Petitioner Riot Games, Inc. - Ex. 1012, p. 1

Ex. 1012, 8:64-9:10; -130 Pet. 53

Aldred and RFC 1692 in further view of Ulrich

UNITED STATES PATENT AND TRADEMARK OF

BEFORE THE PATENT TRIAL AND APPEAL BOA

RIOT GAMES, INC., Petitioner,

V.

PALTALK HOLDINGS, INC., Patent Owner.

Case No. IPR2018-00130 U.S. Patent No. 5,822,523 Issued: October 13, 1998 Filed: February 1, 1996

Inventors: Jeffrey J. Rothschild, Marc P. Kwiatkowski, Daniel J. Samuel

Title: SERVER-GROUP MESSAGING SYSTEM FOR INTEL APPLICATIONS

PETITION FOR INTER PARTES REVIEW

<u>Ulrich</u> renders this claim element obvious for two reasons. Ex. 1007, ¶246.

combined above) to execute <u>Ulrich</u>'s game simulation system as a shared

application on Aldred's system (i.e., "wherein said plurality of host computers

First, it would have been obvious to an Ordinary Artisan to modify Aldred (as

belonging to said first message group correspond to players ... within a three-

[dimensional] space of a computer game"). Ex. 1007, ¶247. Computer games

-130 Pet. 56

Second, it would also have been obvious to an Ordinary Artisan to modify

Aldred (as combined) so that the CSP maintains a database of locations and

selectively transmits game updates to "users in the same general area of the

simulated environment" ("close proximity to one another"), per the teachings of

<u>Ulrich</u>. Ex. 1012, 8:64-9:10; Ex. 1007, ¶252. <u>Ulrich</u> discloses a specific

-130 Pet. 58

Petition: Claim 12 Would Have Been Obvious

Paper No. 1

UNITED STATES PATENT AND TRA

BEFORE THE PATENT TRIAL ANI

RIOT GAMES, INC Petitioner,

V.

PALTALK HOLDINGS Patent Owner.

> Case No. IPR2018-00 U.S. Patent No. 5,822, Issued: October 13, 19 Filed: February 1, 19

Inventors: Jeffrey J. Rothschild, Marc P. Kwia

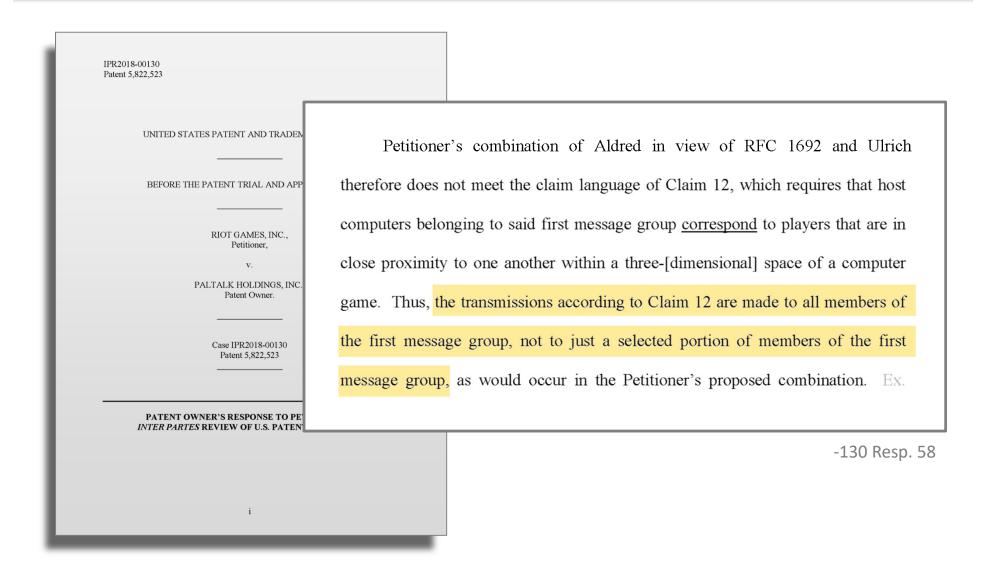
Title: SERVER-GROUP MESSAGING SYST APPLICATIONS

PETITION FOR INTER PART

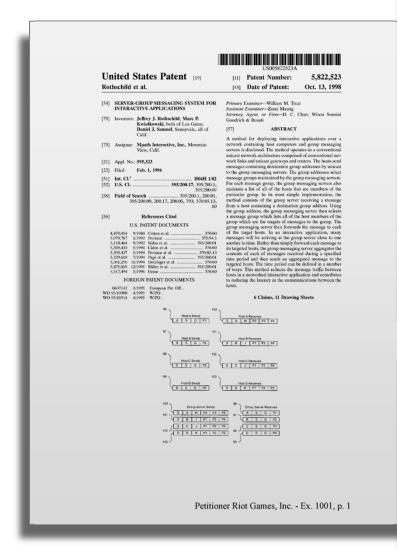
Id., 8:64-9:10. Instead of sending every location update to every host (as is done in Aldred and certain Ulrich embodiments), Ulrich's hub compares that host's location within the simulated environment to the pending updates and "limit[s] the information directed to each apparatus ... [such] that each apparatus only gets (parameter) updates about other users in the same general area of the simulated environment." Id., 8:64-9:10. Ulrich's updates thus implicitly direct the large-scale direct hub to re-transmit them to other machines in the same game whose players are "in the same general area of the simulated environment." Id., 9:5-9, 11:54-60; Ex. 1007, ¶252.

-130 Pet. 59

Patent Owner's Argument



'523 Patent's "Message Group"



implementation of networked interactive applications. Consider a computer game for multiple players that supports hundreds of players that are spread throughout a three dimensional space created by the game. At any time only a few players will be able to see and effect one another in the game since other players will be in other areas that are out of sight. Using conventional phone lines to carry the data from each players computer to the network, it will not be possible to send all actions of each player to all of the other players, but because only a few players will be in close proximity at any one time, it will not be necessary to do so. It is only necessary to send data between the players that are in close proximity to one another. These "groups" of players naturally map onto the message groups of the invention. As

Ex. 1001, 10:1-14; -130 Reply 25-26

could be created for each team within the game. To send a message to all the players within the area that were on one team, a ULP message would be sent to the ULP implicit message group for all the players in the area with an auxiliary address of the logical message group for all the players on the selected team. The GMS would perform the proper set intersection prior to sending the resulting messages to the targeted hosts. The result of this will be that the

Ex. 1001, 11:17-25; -130 Reply 25-26

Patent Owner's Construction of "Message Group"

IN THE UNITED STATES DISTRICT COURT FOR THE DISTRICT OF DELAWARE PALTALK HOLDINGS, INC., Plaintiff C.A. No. RIOT GAMES, INC... JURY TRIAL DEMANDED PLAINTIFF'S COMPLAINT FOR PATENT INFRINGEMENT Plaintiff Pall'alk Holdings, Inc. ("Pall'alk") brings this action against Riot Games, Inc. ("Riot Games"), and for its causes of action alleges as follows: THE PARTIES PalTalk is a Delaware corporation with its principal place of business at 500 North Broadway, Suite 259 Jericho, NY 11753. PalTalk was incorporated in 2001 and is the owner of various patents relating to methods and systems for communicating over networks. Upon information and belief, Riot Games, Inc. is a Delaware corporation having

place of business in Los Angeles. California and offering its products and services, incluaccused herein of infringement, to customers and or potential customers located in the

Delaware, Riot Games may be served with process through its registered agent: The C Trust Company (registered agent) Corporation Trust Center, 1209 Orange Street, Wilm A previously agreed construction for "message group" is:

A collection of one or more host computers that (1) have joined a particular group and (2) receive group messages addressed to that particular group.

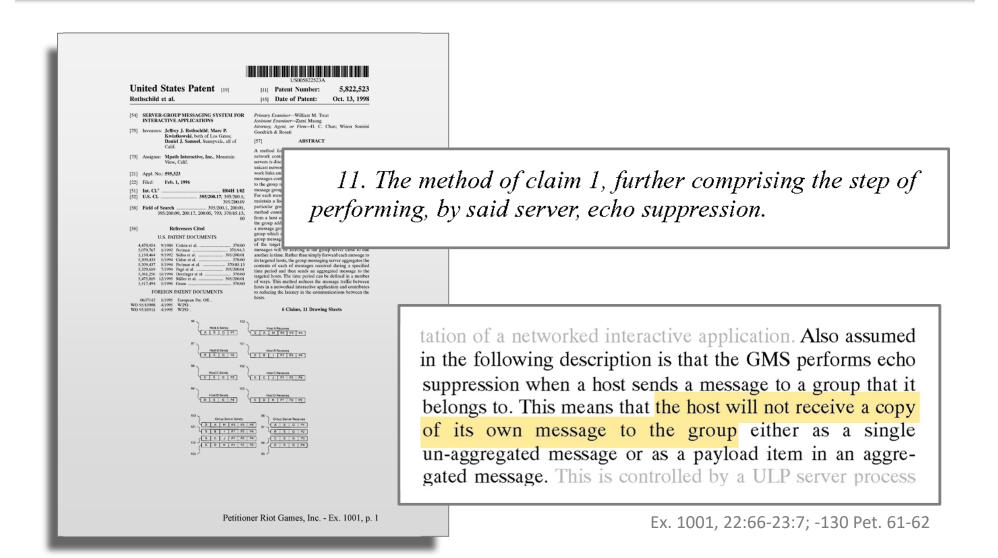
Ex. 1016, 90; -130 Reply 25-26

environment." *Id.*, 8:64-9:10. <u>Ulrich</u>'s updates thus implicitly direct the large-scale direct hub to re-transmit them to other machines in the same game whose players are "in the same general area of the simulated environment." *Id.*, 9:5-9,

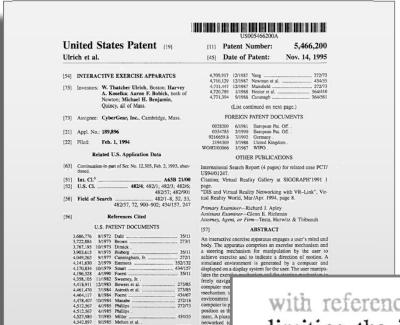
Petitioner Riot Games, Inc. - Ex. 1016, p. 1

-130 Pet. 59

'523 Patent Claim 11



Ulrich's "Echo Suppression"



434/55

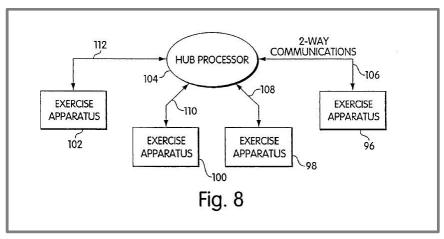
273/85

SWIJLATED 3-D ENVRONMENT

NETWORK

4,542,897 9/1985 Melton et al. .

4,572,509 2/1986 Sitrick 4,616,829 10/1986 Smack, Sr. et al. ... 4,630,817 12/1986 Buckley



Ex. 1012, Fig. 8; -130 Pet. 53

with reference to FIG. 7. The hub 104 is responsible for limiting the information directed to each apparatus in the large-scale direct network of FIG. 8. The hub 104 can ensure, for example, that each apparatus only gets (parameter) updates about other users in the same general area of the simulated environment.

Petitioner Riot Games, Inc. - Ex. 1012, p. 1

Ex. 1012, 9:5-10; -130 Pet. 62

Petition: Claim 12 Combination

Paper No

UNITED STATES PATENT AND TRADEMARK OFFICE

BEFORE THE PATENT TRIAL AND APPEAL BOARD

RIOT GAMES, INC., Petitioner,

V.

PALTALK HOLDINGS, INC., Patent Owner.

Case No. IPR2018-00130 U.S. Patent No. 5,822,523 Issued: October 13, 1998 Filed: February 1, 1996

Inventors: Jeffrey J. Rothschild, Marc P. Kwiatkowski, Daniel J. Samuel

Title: SERVER-GROUP MESSAGING SYSTEM FOR INTERACTIVE APPLICATIONS

PETITION FOR INTER PARTES REVIEW

Aldred discloses serialised channels having a CSP ("said server") that preserves the order of transmitted messages by collecting them centrally and transmitting the same stream of messages to each member. See, e.g., Ex. 1009, 9, 51, Fig. 22; Ex. 1007, ¶278. Aldred however, does not explicitly disclose that data is transmitted to a CSP but then "echo suppress[ed]." Aldred does disclose other channel arrangements where the transmitter does not automatically receive a copy of an update that it sends to a Sharing Set. See, e.g., Ex. 1009, Figs. 15, 16, 18. Ulrich also suggests performing "echo suppression" because it discloses directing

"updates about *other* users" to each apparatus. Ex. 1012, 9:5-10 (emphasis added).

-130 Pet. 62

It would have been additionally obvious to an Ordinary Artisan to modify

Aldred's CSP (as combined above) to perform "echo suppression." Ex. 1007,

¶280. For example, <u>Ulrich</u> suggests performing "echo suppression" because it discloses the direct hub directing "updates about other users" to each apparatus.

Ex. 1012, 9:5-10 (emphasis added). This would also reduce the amount of

-130 Pet. 62

Appendix

'686 Patent Claim 3

- 3. A method for facilitating communications among a plurality of host computers over a network to implement a shared, interactive application, comprising the steps of:
 - (1) receiving a create message from one of the plurality of host computers, wherein said create message specifies a message group to be created;
 - (2) receiving join messages from a first subset of the plurality of host computers, wherein each of said join messages specifies said message group;
 - (3) receiving host messages from a second subset of said first subset of the plurality of host computers belonging to said message group, wherein each of said messages contains a payload portion and a portion that is used to identify said message group;
 - (4) aggregating said payload portions of said host messages received from said second subset of the plurality of host computers to create an aggregated message;
 - (5) transmitting said aggregated message to said first subset of the plurality of host computers belonging to said message group;
 - wherein said aggregated message keeps the shared, interactive application operating consistently on each of said first subset of the plurality of host computers.

Ex. 1002, 28:11-34

'686 Patent Claim 7

- 7. A method for facilitating communications among a plurality of host computers over a network to implement a shared, interactive application, comprising the steps of:
 - (1) receiving messages from a subset of the plurality of host computers belonging to a message group, wherein each of said messages contains a payload portion and a portion that is used to identify said message group;
 - (2) aggregating said payload portions of said messages to create an aggregated payload; and
 - (3) transmitting said aggregated message to each of the plurality of host computers belonging to said message group;
 - wherein said aggregated message keeps the shared, interactive application operating consistently on each of the plurality of host computers belonging to said message group.

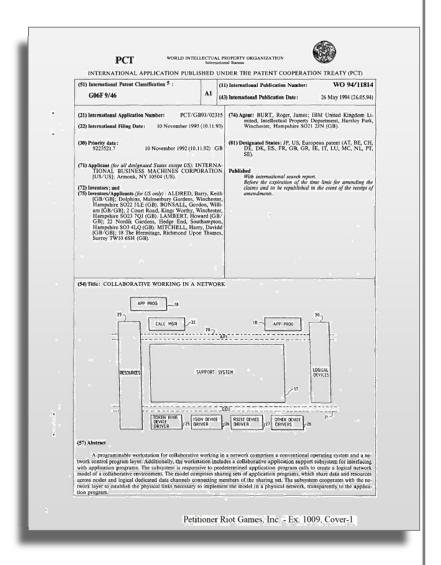
Ex. 1002, 28:58-29:7

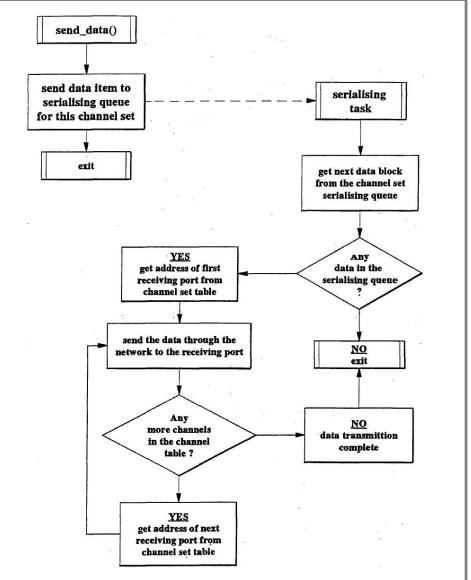
'686 Patent Claim 18

- 18. A method for facilitating communications among a plurality of host computers over a network to implement a shared, interactive application, comprising the steps of:
 - receiving a host message from one of the plurality of host computers belonging to a message group, wherein said host message contains a payload portion and a portion that is used to identify said message group;
 - (2) forming a server message by using said payload portion of said host message; and aggregating said payload portion with the payload portion of a second host message received from another of the plurality of host computers belonging to said message group
 - (3) transmitting said server message to each of the plurality of host computers belonging to said message group;
 - whereby said server message keeps the shared, interactive application operating consistently on each of the plurality of host computers belonging to said message group.

Ex. 1002, 30:18-39

Aldred's Figure 22





Ex. 1009, Fig. 22; -129 Pet. _

Claim 32 – "said transport layer protocol is TCP/IP"

Paper No. 1

UNITED STATES PATENT AND TRADEMARK OFFICE

BEFORE THE PATENT TRIAL AND APPEAL BOARD

RIOT GAMES, INC., Petitioner,

V.

PALTALK HOLDINGS, INC., Patent Owner.

Case No. IPR2018-00129 U.S. Patent No. 5,822,523 Issued: October 13, 1998 Filed: February 1, 1996

Inventors: Jeffrey J. Rothschild, Marc P. Kwiatkowski, Daniel J. Samuel

Title: SERVER-GROUP MESSAGING SYSTEM FOR INTERACTIVE APPLICATIONS

PETITION FOR INTER PARTES REVIEW

xii. <u>Claim 32</u>

Claim 32 recites "[t]he method of claim 1, wherein said sending and said transmitting are performed by an upper-level protocol implemented above a transport layer protocol of said unicast network, wherein said transport layer protocol is TCP/IP."

As explained above for claim 6 (§VI.A.vi), an Ordinary Artisan would understand Aldred's scheme is implemented on top of the underlying transport layer ("upper-level protocol implemented above a transport layer protocol"). Ex. 1017, 1338; Ex. 1007, ¶¶172-174. It would have been obvious to combine Aldred and RFC 1692 as expressed above. See §VI.A.i.e. The transport protocol used by the nodes in Aldred would therefore be RFC 1692's TMux-enhanced IP protocol, which uses TCP/IP ("wherein said transport layer protocol is TCP/IP"). See Ex. 1010, 6-7; Ex. 1007, ¶180.

Aldred as combined therefore renders this claim obvious.

-129 Pet. 52-53