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SEARCH-ON-THE-FLY WITH MERGE FUNCTION					
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Respectfully submitted,

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## SEARCH-ON-THE-FLY WITH MERGE FUNCTION

### Related Applications

This application is a continuation-in-part of Application Serial Number 09/513,340, filed February 25, 2000, entitled Search-On-The-Fly/Sort-On-The-Fly Search Engine, which is hereby incorporated by reference.

### Technical Field

The technical field is information management systems, interfaces, and mechanisms, and methods for searching one or more databases.

### Background

In the most general sense, a database is a collection of data. Various architectures have been devised to organize data in a computerized database. Typically, a computerized database includes data stored in mass storage devices, such as tape drives, magnetic hard disk drives and optical drives. Three main database architectures are termed hierarchical, network and relational. A hierarchical database assigns different data types to different levels of the hierarchy. Links between data items on one level and data items on a different level are simple and direct. However, a single data item can appear multiple times in a hierarchical database and this creates data redundancy. To eliminate data redundancy, a network database stores data in nodes having direct access to any other node in the database. There is no need to duplicate data since all nodes are universally accessible. In a relational database, the basic unit of data is a relation. A relation corresponds to a table having rows, with each row called a tuple, and columns, with each column called an attribute. From a practical standpoint, rows represent records of related data and columns identify individual data elements. The order in which the rows and columns appear in a table has no significance. In a relational database, one can add a new column to a table without having to modify older applications that access other columns in the table. Relational databases thus provide flexibility to accommodate changing needs.

All databases require a consistent structure, termed a schema, to organize and manage the information. In a relational database, the schema is a collection of tables. Similarly, for each

1 table, there is generally one schema to which it belongs. Once the schema is designed, a tool,  
2 known as a database management system (DBMS), is used to build the database and to  
3 operate on data within the database. The DBMS stores, retrieves and modifies data associated  
4 with the database. Lastly, to the extent possible, the DBMS protects data from corruption and  
5 unauthorized access.

6 A human user controls the DBMS by providing a sequence of commands selected from  
7 a data sublanguage. The syntax of data sublanguages varies widely. The American National  
8 Standards Institute (ANSI) and the International Organization for Standardization (ISO) have  
9 adopted Structured English Query Language (SQL) as a standard data sublanguage for  
10 relational databases. SQL comprises a data definition language (DDL), a data manipulation  
11 language (DML), and a data control language (DCL). The DDL allows users to define a  
12 database, to modify its structure and to destroy it. The DML provides the tools to enter,  
13 modify and extract data from the database. The DCL provides tools to protect data from  
14 corruption and unauthorized access. Although SQL is standardized, most implementations of  
15 the ANSI standard have subtle differences. Nonetheless, the standardization of SQL has  
16 greatly increased the utility of relational databases for many applications.

17 Although access to relational databases is facilitated by standard data sublanguages,  
18 users still must have detailed knowledge of the schema to obtain needed information from a  
19 database since one can design many different schemas to represent the storage of a given  
20 collection of information. For example, in an electronic commerce system, product information,  
21 such as product SKU, product name, product description, price, and tax code, may be stored  
22 in a single table within a relational database. In another electronic commerce system, product  
23 SKU, product name, description, and tax code may be stored in one table while product SKU  
24 and product price are stored in a separate table. In this situation, a SQL query designed to  
25 retrieve a product price from a database of the first electronic commerce system is not useful  
26 for retrieving the price for the same product in the other electronic system's database because  
27 the differences in schemas require the use of different SQL queries to retrieve product price.

1 As a consequence, developers of retail applications accessing product information from  
2 relational databases may have to adapt their SQL queries to each individual schema. This, in  
3 turn, prevents their applications from being used in environments where there are a wide variety  
4 of databases having different schemas, such as the World Wide Web.

5 A further problem with conventional search engines is a tendency to return very large  
6 amounts of data, or to require the search parameters to be narrowed. When large amounts of  
7 data are presented, the display may take many “pages” before all data is seen by the user. The  
8 time and expense involved in such a data review may be significant.

9 **Summary**

10 A Sort-on-the-Fly/Search-on-the-Fly search engine (hereafter, search-on-the-fly  
11 search engine) provides an intuitive means for searching databases, allowing a user to access  
12 data in the database without having to know anything about the database structure. A user  
13 selects a desired search term, and a database manager searches the database for all instances  
14 of the desired term, even if a specific file or table does not contain the instance. For example,  
15 if a user wants to search the database using the name of a specific individual as a database entry  
16 point, the database manager will search the database using the desired name, and will organize  
17 the search results so that all entries associated with that name are displayed. The database  
18 need not have a specific file (in a flat database) or a table (in a relational database) of names.  
19 The user may perform further on-the-fly searches to narrow or focus the search results, or for  
20 other reasons. For example, given search results for all names that include the name “Smith,”  
21 the user may then decide to search for all “Smiths” that include an association to an address in  
22 New Jersey. The search-on-the-fly search engine then conducts a further search using this  
23 criteria and produces a second search result. Further narrowing or broadening of the search  
24 are permitted, with the search-on-the-fly search engine returning results based on any new  
25 criteria.

26 In an embodiment, the search-on-the-fly search engine uses graphical user interfaces  
27 (GUIs) and one or more icons to make the search process as efficient as possible. The GUIs



1 may incorporate one or more pull down menus of available search terms. As a user selects an  
2 item from a first pulldown menu, a subsequent pulldown menu displays choices that are  
3 available for searching. The process continues until the search engine has displayed a discrete  
4 data entry from the database. The pulldown menus are not pre-formatted. Instead, the  
5 pulldown menus are created "on-the-fly" as the user steps through the search process. Thus,  
6 the search-on-the-fly search engine is inherently intuitive, and allows a user with little or no  
7 knowledge of the database contents, its organization, or a search engine search routine to  
8 execute comprehensive searches that return generally accurate results.

9 The search-on-the-fly search engine also searches on key words specified by the user.  
10 The search-on-the-fly search engine can be used to exclude certain items. The search-on-the-  
11 fly search engine incorporates other advanced features such as saving search results by  
12 attaching a cookie to a user's computer, and associating icons with the search results.

13 The search-on-the-fly search engine may be used with both internal and external  
14 databases. For example, the search-on-the-fly search engine may be used with a company  
15 internal database and one or more databases accessible through the Internet.

16 The search-on-the-fly search engine is user-friendly. With one interface, many different  
17 types of databases or database schemas may be searched or sorted.

18 Finally, the search-on-the-fly technique, and other techniques discussed above may be  
19 used in conjunction with a method of doing business, particularly a business method that uses  
20 the Internet as a communications backbone.

### 21 **Description of the Drawings**

22 The detailed description will refer to the following figures, in which like numerals refer  
23 to like objects, and in which:

24 Figure 1 is a block diagram of a system that uses a search-on-the-fly/sort-on-the-fly  
25 search engine;

26 Figure 2 is another overall block diagram of the system of Figure 1;

1           Figure 3 is a detailed block diagram of the search engine used with the system of  
2 Figure 2;

3           Figure 4 is an example of a search-on-the-fly using the search engine of Figure 3;

4           Figures 5 - 9 are detailed block diagrams of components of the search engine of  
5 Figure 3;

6           Figure 10 is another example of a search-on-the-fly using the search engine of Figure  
7 3;

8           Figures 11 - 15b are additional examples of a search-on-the-fly using the search  
9 engine of Figure 3;

10          Figures 16 - 20 are flow charts illustrating operations of the search engine of Figure  
11 3;

12          Figure 21 illustrates a further function of the search engine of Figure 3 in which  
13 results of more than one search are combined;

14          Figures 22 - 26 illustrate graphical user interfaces that may be displayed in  
15 conjunction with operation of the system of Figure 1;

16          Figure 27 is a flowchart illustrating an alternate operation of the query generator;

17          Figure 28 is a flowchart illustrating an alternate operation of the truncator;

18          Figures 29 - 36 illustrate user interfaces with search results from a search on the fly and  
19 a merge function;

20          Figures 37 - 39 illustrate a keyword search result form a search on the fly with the  
21 merge function; and

22          Figures 40-49 illustrate additional search results.

23           **Detailed Description**

24           Ordinary search engines place constraints on any search. In particular, a partial  
25 ordering of available search criteria limits application of the search engine only to certain search  
26 sequences. The user is given a choice of search sequences, and the order in which individual  
27 search steps in the search sequence become available limits the direction of the search. A user

1 who desires to take a vacation cruise may use an Internet search engine to find a desired  
2 vacation package. The search begins with presentation of a list of general categories, and the  
3 user clicks on "travel," which produces a list of subcategories. The user then clicks on  
4 "cruises" from the resulting list of subcategories, and so on in a cumulative narrowing of  
5 possibilities until the user finds the desired destination, date, cruise line, and price. The order  
6 in which choices become available amounts to a predefined "search tree," and the unspoken  
7 assumption of the search engine designer is that the needs and thought processes of any user  
8 will naturally conform to this predefined search tree.

9 To an extent, predefined constraints are helpful in that predefined constraints allow a  
10 search engine to logically and impersonally order the user's thoughts in such a way that if the  
11 user has a clear idea of what object the user wants, and if the object is there to be found, then  
12 the user is assured of finding the object. Indeed, the user may want to know that choosing any  
13 available category in a search sequence will produce an exhaustive and disjunctive list of  
14 subcategories from which another choice can be made. Unfortunately, an unnecessarily high  
15 cost is too often paid for this knowledge: The user is unnecessarily locked into a limited set of  
16 choice sequences, and without sufficient prior knowledge of the object being sought, this  
17 limitation can become a hindrance. Specifically, where prescribed search constraints are  
18 incompatible with the associative relationships in the user's mind, a conflict can arise between  
19 the thought processes of the user and the function of the search engine.

20 At one time, such conflicts were written off to the unavoidable differences between  
21 computers and the human mind. However, some "differences" are neither unavoidable nor  
22 problematic. In the case of search engine design, the solution is simple: upon selecting a  
23 category or entering a keyword, the user can be given not only a list of subcategories, but the  
24 option to apply previously available categories as well. In slightly more technical terms, the  
25 open topology of the search tree can be arbitrarily closed by permitting search sequences to  
26 loop and converge. Previous lists can be accessed and used as points of divergence from

1 which new sub-sequences branch off, and the attributes corresponding to distinct sub-  
2 sequences can later be merged.

3 A sort-on-the-fly/search-on-the-fly search engine (hereafter, search-on-the-fly search  
4 engine) provides an intuitive means for searching various types of databases, allowing a user  
5 to access data in the database without having to know anything about the database structure.  
6 A user selects a desired search term, and a database manager searches the database for all  
7 instances of the desired term, even if a specific file or table does not contain the instance. For  
8 example, if a user wants to search the database using the name of a specific individual as a  
9 database entry point, the database manager will search the database using the desired name,  
10 and will organize the search results so that all entries associated with that name are displayed.  
11 The database need not have a specific file (in a flat database) or a table (in a relational  
12 database) of names. The user may perform further on-the-fly searches to narrow the search  
13 results, or for other reasons. The search engine then conducts a further search using this criteria  
14 and produces a second search result. Further narrowing or broadening of the search are  
15 permitted, with the search engine returning results based on any new criteria.

16 Figure 1 is a block diagram of a system 10 that uses the search-on-the-fly search  
17 engine. In Figure 1, a database 12 is accessed using a hardware/software interface device 100  
18 to provide data to a user terminal 14. Additional databases 13 and 15 may also be accessed  
19 by the terminal 14 using the device 100. The databases 12, 13 and 15 may use different  
20 schemas, or may use a same schema. As will be described later, the device 100 may include  
21 the search-on-the-fly search engine. In an alternative embodiment, the search-on-the-fly search  
22 engine may be co-located with the terminal 14. In yet another embodiment, the search-on-the-  
23 fly search engine may be incorporated into the structure of one or more of the databases 12,  
24 13 and 15. The device 100 may interface with any one or more of the databases 12, 13 and  
25 15 using a network connection such as through the Internet, for example. Other  
26 communications mediums may also be used between the terminal 14, the device 100 and any  
27 one or more of the databases 12, 13 and 15. These mediums may include the public switched

1 telephone network (PSTN), cable television delivery networks, Integrated Services Digital  
2 Networks (ISDN), digital subscriber lines (DSL), wireless means, including microwave and  
3 radio communications networks, satellite distribution networks, and any other medium capable  
4 of carrying digital data.

5 The system shown in Figure 1 is but one of many possible variations. The search-on-  
6 the-fly search engine could also be incorporated within a single computer, such as a personal  
7 computer, a computer network with a host server and one or more user stations, an intranet,  
8 and an Internet-based system, as shown in Figure 2. Referring again to Figure 2, the terminal  
9 14 may be any device capable of displaying digital data including handheld devices, cellular  
10 phones, geosynchronous positioning satellite (GPS) devices, wrist-worn devices, interactive  
11 phone devices, household appliances, televisions, television set top boxes, handheld computers,  
12 and other computers.

13 Figure 3 is a detailed block diagram of an exemplary search-on-the-fly search engine  
14 125. The search engine 125 includes a request analyzer 130 that receives search requests 114  
15 from the terminal 14 (not shown in Figure 3) and sends out updated requests 115 to a query  
16 generator 150. A status control 140 receives a status update signal 116 and a request status  
17 control signal 118 and sends out a request status response 119 to the request analyzer 130.  
18 The status control 140 also keeps track of search cycles, that is, the number of search iterations  
19 performed. The query generator 150 receives the updated requests 115 from the request  
20 analyzer 130 and sends a database access signal 151 to a database driver 170. The query  
21 generator 150 receives results 153 of a search of the database 12 (not shown in Figure 3) from  
22 the database driver 170. The query generator 150 provides a display signal 175 to the terminal  
23 14. The database driver 170 sends a database access signal 171 to the database 12. Finally,  
24 a database qualifier 160 receives information 161 from the database driver 170 and provides  
25 a list 163 of available data fields from the database 12. As will be described later, the list of  
26 available data fields 163 may be displayed to a user at the terminal 14, and may be sorted and  
27 processed using the request analyzer 130 in conjunction with the database qualifier 160. The

1 database qualifier 160 also receives search information and other commands 131 from the  
2 request analyzer 130.

3 The search engine 125 may identify a database schema by simply using a trial and error  
4 process. Alternatively, the search engine 125 may use other techniques know in the art. Such  
5 techniques are described, for example, in U.S. Patent 5,522,066, "Interface for Accessing  
6 Multiple Records Stored in Different File System Formats," and U.S. Patent 5,974,407,  
7 "Method and Apparatus for Implementing a Hierarchical Database Management System  
8 (HDBMS) Using a Relational Database Management System (RDBMS) ad the Implementing  
9 Apparatus," the disclosures of which is hereby incorporated by reference.

10 The search engine 125 provides search-on-the-fly search capabilities and more  
11 conventional search capabilities. In either case, the search engine 125 may perform a  
12 preliminary database access function to determine if the user has access to the database 12.  
13 The search engine 125 also determines the database schema to decide if the schema is  
14 compatible with the user's data processing system. If the database schema is not compatible  
15 with the user's processing system, the search engine 125 may attempt to perform necessary  
16 translations so that the user at the terminal 14 may access and view data in the database 12.  
17 Alternatively, the search engine 125 may provide a prompt for the user indicating  
18 incompatibility between the terminal 14 and a selected database.

19 The search engine 125 may conduct a search using one or more search cycles. A  
20 search cycle includes receipt of a request 114, any necessary formatting of the request 114,  
21 and any necessary truncation steps. The search cycle ends when a result list 175 is provided  
22 to the terminal 14. The search engine 125 may retain a status of each past and current search  
23 cycle so that the user can modify the search at a later time. The user may also use this feature  
24 of retaining a status of past and current search cycles to combine results of multiple searches,  
25 using, for example, a Boolean AND function, a Boolean OR function, or other logic function.  
26 The above listed functions will be described in more detail later.

1           The search-on-the-fly function of the search engine 125 begins by determining available  
2 data fields of the database 12. The database 12 may have its data organized in one or more  
3 data fields, tables, or other structures, and each such data field may be identified by a data field  
4 descriptor. In many cases, the data field descriptor includes enough text for the user at the  
5 terminal 14 to determine the general contents of the data field. The list of data fields may then  
6 be presented at the terminal 14, for example, in a pull down list. An example of such a data  
7 field result list is shown in Figure 4, which is from a federal database showing data related to  
8 managed health care organizations. This database is available at  
9 <http://tobaccopapers.org/dnld.htm>. In Figure 4, the first data field listed is "PlanType," which  
10 is shown in result list 156. Other data field descriptors show the general categories of data in  
11 the database.

12           Using the terminal 14, the user may select one of the data field descriptors to be  
13 searched. For example, the user could select "city." If a number of entries, or records, in the  
14 city data field is short, a further result list of complete city names may be displayed. If the  
15 entries are too numerous to be displayed within a standard screen size, for example, the search  
16 engine 125 may, in an iterative fashion, attempt to reduce, or truncate, the result list until the  
17 result list may be displayed. In the example shown in Figure 4, entries in the city data field are  
18 so numerous (the database includes all U.S. cities that have a managed health care organization)  
19 that the search engine 125 has produced a result list 157 that shows only a first letter of the city.  
20 Based on the available database data fields, the user may then perform a further search-on-the-  
21 fly. In this case, the user may choose cities whose first initial is "N." The search engine 125  
22 then returns a result list 158 of cities whose names start with the letter "N." Because in this  
23 instance the result list 158 is short, no further truncation is necessary to produce a manageable  
24 list.

25           Figure 5 is a more detailed block diagram of the request analyzer 130. A protocol  
26 analyzer 133 receives the request 114 and provides an output 135 to a constraint collator 136.  
27 The protocol analyzer 133 examines the received request 114, determines a format of the

1 request 114, and performs any necessary translations to make the request format compatible  
2 with the database to be accessed. If the database to be accessed by the terminal 14 is part of  
3 a same computer system as the terminal 14, then the protocol analyzer 133 may not be  
4 required to perform any translations or to reformat the request 114. If the database to be  
5 accessed is not part of the same computer system as the terminal 14, then the protocol analyzer  
6 133 may be required to reformat the request 114. The reformatting may be needed, for  
7 example, when a request 114 is transmitted over a network, such as the Internet, to a database  
8 coupled to the network.

9 The constraint collator 136 provides the updated request 115 (which may be an initial  
10 request, or a subsequent request) to the query generator 150. The constraint collator 136 is  
11 responsible for interpreting the request 114. The constraint collator 136 performs this function  
12 by comparing the request 114 against information stored in the status control 140. In  
13 particular, the constraint collator 136 sends the request status control signal 118 to the status  
14 control 140 and receives the request status response 119. The constraint collator 136 then  
15 compares the request status response 119 to constraint information provided with the request  
16 114 to determine if the constraint status should be updated (e.g., because the request 114  
17 includes a new constraint). In an embodiment, the constraint collator 136 compares constraint  
18 information in a current request 114 to constraint information residing in the status control 140,  
19 and if the current request 114 includes a new constraint, such as a new narrowing request (for  
20 example, when the user clicks, touches or points over a field shown in a last search cycle), then  
21 the constraint collator 136 adds the updated information and sends the updated request 115  
22 to the query generator 150. If the constraint status should be updated, the constraint collator  
23 136 sends the status update 118 to the status control 140. If the request 114 is a refresh  
24 request, the constraint collator 136 sends a reset command 131 to the database qualifier 160.  
25 The updated request 115 (possibly with a new constraint) is then sent to the query analyzer 150  
26 for further processing.



1           Figure 6 is a block diagram of the query generator 150. The overall functions of the  
2 query generator 150 are to scan a database, such as the database 12, using the database driver  
3 170, and to collect search results based on constraints supplied by the request analyzer 130.  
4 The query generator 150 then returns the search results 175 to the terminal 14.

5           The query generator 150 includes a truncator 152 and a dispatcher 154. The truncator  
6 152 receives the updated request 115, including a new constraint, if applicable. The truncator  
7 152 creates new queries, based on new constraints, and applies the new requests 151 to the  
8 database 12 using the database driver 170. The truncator 152 may include a variable limit 155  
9 that is set, for example, according to a capacity of the terminal 14 to display the search results  
10 175. If data retrieved from the database 12 exceed the limit value, the truncator 152 adjusts  
11 a size (e.g., a number of entries or records) of the data until a displayable result list is achieved.  
12 One method of adjusting the size is by cycling (looping). Other methods may also be used to  
13 adjust the size of the result list. For example, the terminal 14 may be limited to displaying 20  
14 lines of data (entries, records) from the database 12. The truncator 152 will cycle until the  
15 displayed result list is at most 20 lines. In an embodiment, the truncation process used by the  
16 truncator 152 assumes that if the user requests all values in a particular data field from the  
17 database 12, and there are no other constraints provided with the request 114, and if the size  
18 of the resulting result list is larger than some numeric parameter related to a display size of the  
19 terminal 14, then the constraints may be modified by the truncator 152 so that the result list can  
20 accommodated (e.g., displayed on one page) by the terminal 14. For example, instead of a  
21 full name of a city, some part of the name - the first n letters - is checked against the database  
22 12 again, and n is reduced until the result list is small enough for the capacity of the terminal 14.  
23 If the maximum number of displayable results is three (3), and the database 12 contains the  
24 names of six cities "Armandia, Armonk, New Orleans, New York, Riverhead, Riverdale," then  
25 the first attempt to "resolve" the result list will stop after a result list display is created with the  
26 full name of the cities:  
27 Armandia, Armonk, New Orleans... (the limit was reached)

1 Try again with 7 characters:

2 Armandia, Armonk, New Orl, New Yor, (limit reached again)

3 Again with 5 characters:

4 Armandia, Armonk, New O, New Y, (limit reached again)

5 Again with 3 characters:

6 Arm (...), New (...), Riv (...). These results may now be displayed on the terminal 14. The  
7 display of Arm, New, Riv can then be used to conduct a further search-on-the-fly. For  
8 example, a user could then select Riv for a further search-on-the-fly. The result list returned  
9 would then list two cities, namely Riverhead and Riverdale.

10 In another embodiment, a fixed format is imposed such that all queries generated  
11 against a database will have preset limits corresponding to the capacity of the terminal 14.

12 In yet another embodiment, the truncator 152 may adjust the field size by division or  
13 other means. For example, if the display limit has been reached, the truncator 125 may reduce  
14 the field size, X by a specified amount. In an embodiment, X may be divided by two.  
15 Alternatively, X may be multiplied by a number less than 1, such as 3/4, for example. Adjusting  
16 the field size allows the search engine 125 to perform more focused searches and provides  
17 more accurate search results.

18 In still another embodiment, the user may select a limit that will cause the truncator 152  
19 to adjust the field size. For example, the user could specify that a maximum of ten entries  
20 should be displayed.

21 For certain data fields, a terminal 14, such as a hand-held device for example, may  
22 have a very limited display capacity. Alternatively a user may specify a limit on the number of  
23 entries for display. In these two illustrated cases, the search engine 125 may return a result list  
24 175 of the request 114 on multiple display pages, and the user may toggle between these  
25 multiple display pages. As an example, if the terminal 14 is limited to displaying a maximum of  
26 ten entries, and if the request 114 results in a return of a data field comprising the 400 largest  
27 cities in the United States, the truncator 152 will produce a list of 23 entries comprising 23

1 alphabetical characters (no cities that begin with Q, Y or Z - see Figure 4). The search engine  
2 125 may then display the results on three pages. Alternatively, the truncator 152 could  
3 produce a list of letter groups into which the cities would fall, such as A-D, E-G, H-M, N-R,  
4 and R-X, for example. In another alternative, the search engine 125 may send a notice to the  
5 terminal that the request 114 cannot be accommodated on the terminal 14 and may prompt the  
6 user to add an additional constraint to the request 114, so that a search result may be displayed  
7 at the terminal 14.

8           Adjusting the data field size also provides more convenient search results for the user.  
9 For example, if a user were to access an Internet-based database for books for sale, and were  
10 to request a list of all book titles beginning with the letter “F,” a common search engine might  
11 return several hundred titles or more, displaying perhaps twenty titles (entries) at a time. The  
12 user would then have to look through each of many pages to find a desired title. This process  
13 could be very time-consuming and expensive. Furthermore, if the search results were too large,  
14 the common search engine might return a notice saying the results were too large for display  
15 and might prompt the user to select an alternative search request. However, performing the  
16 same search using the search engine 125 allows the truncator 152 to reduce the size of the  
17 information displayed to a manageable level. In this example, if the request 114 includes the  
18 constraint “F,” the truncator 152 will loop through the data in a data field that includes book  
19 titles starting with the letter “F” until a list is available that can fit within the display limits of the  
20 terminal 14, or that fits within a limit set by the user, for example. The first list returned to the  
21 terminal 14 as a result of this request 114 may be a two letter combination with “F” as the first  
22 letter and a second letter of a book title as the second letter. For example, the first list may  
23 include the entries “Fa,” “Fe,” “Fi,” “Fo,” and “Fu,” all of which represent titles of books. The  
24 user could then select one of the entries “Fa,” “Fe,” “Fi,” “Fo,” and “Fu” to perform a further  
25 search, continuing the process until one or more desired titles are displayed. An example of  
26 a similar truncation result is shown in Figure 14.

1           When a parameter related to the search results is adequately truncated, the parameter  
2 is directed to the dispatcher 154, which retrieves the data from database 12 using the database  
3 driver 170. The dispatcher 154 then directs the final, truncated search results 175 back to the  
4 terminal 14 as a response to the request 114.

5           Figure 7 is a block diagram showing the status control 140, which is responsible for  
6 monitoring the status of a current search. Due to the nature of the search engine 125, the user  
7 can choose any combination of constraints, fields or keywords, including those from past and  
8 current search cycles. The status control 140 may keep track of all past cycles of the search,  
9 as well as all information necessary to return to any of those past search cycles. The status  
10 control 140 includes a status data module 142, and an index module 144. The status data  
11 module 142 contains data related to each such search cycle, including the constraint(s) entered  
12 during the search cycle, any truncation steps taken, and the results of such truncation, for  
13 example. The index module 144 provides access to these data. When the request 114 is being  
14 analyzed by the request analyzer 130, the constraint collator 136 sends a request status query  
15 116 to the index module 144. The status data module 142 contains information related to all  
16 past and current search cycles, which are referenced by the index module 144, and delivers  
17 a status response 119 for the most recent search cycle to the constraint collator 136. When a  
18 new constraint is sent to the query generator 150, the status data module 142 is updated 118  
19 by the constraint collator 136. Specific structures of the request 114, the request status query  
20 116, the status response 119 and the request status control 118 will be provided later.

21           The status data module 142 may be reset by the database qualifier 160 with all  
22 available fields when a refresh function is used. In an embodiment, the refresh function may be  
23 used to clear all past search cycles and the current search cycle from the status control 140.  
24 In such an event, the search results, such as the search results shown in Figure 4, will no longer  
25 be displayed at the terminal 14, and data related to the past and the current search cycles may  
26 not be used for future search cycles. In effect, the refresh function may cause the entire search  
27 to be discarded. The refresh function may be activated when a user selects a refresh button

1 (see Figure 4) on a displayed result list, or on another portion of a GUI. Alternatively, the  
2 refresh function may discard selected search cycles. In this alternative embodiment, the user  
3 may, for example, move a cursor to a desired result list from a past search cycle and activate  
4 a refresh, reset, back, or drop button. All data associated with search cycles subsequent to  
5 the selected search cycle, including all displayed result lists may then be discarded.

6 Figure 8 is a block diagram showing the database qualifier 160. The database qualifier  
7 160 provides data field information at the start of a search or when the search engine 125 is  
8 refreshed. A field assessor 162 access the database 12 using the database driver 170, and  
9 identifies and accesses discrete data fields and other information in the database 12. A field  
10 converter 164 structures the data field information into a usable (searchable/sortable) structure  
11 and sends 163 the formatted data field information to the status control 140. Techniques for  
12 identifying and accessing the data fields, and for formatting the data field information are well  
13 known in the art. Such techniques are described, for example, in U.S. Patent 5,222,066,  
14 Interface for Accessing Multiple Records Stored in Different File System Formats, the  
15 disclosure of which is hereby incorporated by reference.

16 Figure 9 is a block diagram of the database driver 170. The database driver 170 is  
17 the universal interface with the database 12, which can be a local or a remote database.

18 Figure 10 is an example of a search-on-the-fly using the search engine 125. In Figure  
19 10, a database 200 includes information related to a number of individuals. The information  
20 in the database 200 may be presented at the terminal 14 using a series of screens or menus 201  
21 - 230. The user first accesses the database 200 and is presented with a list 201 of the  
22 information or data fields contained in the database 200. The result list 201 is generated by the  
23 field assessor 162, and is provided for display at the terminal 14 by the query generator 150.  
24 As shown in Figure 10, a user has selected the data field "City" for display of information.  
25 However, the number of "cities" listed in the database 200 is too large to conveniently display  
26 at one time (i.e., on one page) at the terminal 14. Accordingly, the truncator 152 will loop a

1 required number of times until an adequate display is available. In Figure 10, the menu 203  
2 shows the results of the truncation with only the first letter of a city name displayed.

3 Using the menu 203, the user has selected cities beginning with the letter "A." The  
4 results are shown in menu 205. Now, the user elects to conduct another search-on-the-fly, by  
5 selecting the "sort-on-the-fly" option 206. The query generator 150 displays all the information  
6 fields available from the database 200, except for the information field already displayed,  
7 namely "City." The results are displayed in menu 207. The user then elects to further search  
8 on the data field "State." The query generator 150 returns the requested information as  
9 displayed in menu 209, listing five states by their common two-letter abbreviation. The user  
10 then chooses New York from the menu 209, and the query generator 150 returns a list of cities  
11 in New York, menu 211.

12 Next, the user elects to conduct another search-on-the-fly, option 212, and the query  
13 generator 150 returns only the remaining data fields for display in menu 215. From the menu  
14 215, the user selects "Address" for the next data field to search, and the query generator 150  
15 returns an menu 217 showing only first letters of the address. This signifies that the data field  
16 "Address" was too large to be easily displayed on the terminal 14. The user then elects to  
17 search on all addresses that begin with "C." The query generator 150 returns a list of  
18 addresses by displaying only street names, menu 219.

19 The user then elects to conduct a further search-on-the-fly, option 220, and the  
20 remaining two data fields, "Name" and "Phone" are displayed as options in menu 221. The  
21 user selects name, and the query generator returns a further breakdown of the data by last  
22 name and by first name, menu 223. This process continues, with further menus being used to  
23 select a last name and a first name from the database 200. When the final selection is made,  
24 information from the database 200 related to the individual is displayed in window 230.

25 In the example shown in Figure 10, the user could have refreshed the search engine  
26 125 at any time, and the search would have recommenced at the beginning. Alternatively, the  
27 user could, by simply selecting a prior menu, such as the menu 215, have changed the course

1 of the search. In this alternative, if the user had gone back to the menu 215 and instead of  
2 selecting "Address" selected "Phone," then the menus 217 - 229 would be removed from  
3 display at the terminal 14, and the search would begin over from the point of the menu 215.

4 Figures 11 - 15 illustrate exemplary searches of a remote database, such as the  
5 database 13 shown in Figure 1. The database in the illustrated example is for an Internet  
6 website 232 that sell books. The examples illustrated are based on a Barnes & Noble website.  
7 In Figure 11, the user has applied the search engine 125 to the website 232 database, and the  
8 query generator 150 has returned a list 233 of data fields from which the user may select to  
9 access data from the website 232 database. The list 233, and other lists described below, may  
10 be displayed as overlays on the website 232. In the example illustrated, the user selects "Title"  
11 for the first search cycle. Because the list of titles is too large to easily display at the terminal  
12 14, the truncator 152 loops until an alphanumeric list 234 is created. The list 234 is then  
13 returned to the terminal 14. For the next search cycle, the user selects titles that begin with the  
14 letter "C." Again, the data field contains too many entries to conveniently display at the terminal  
15 14, and the truncator 152 loops as appropriate until list 235 is created. The process continues  
16 with subsequent lists 236 and 237 being returned to the terminal 14.

17 Figures 12 - 15b illustrate alternate searches that may be completed using the website  
18 232 database.

19 For the search results shown in Figures 11 - 15a, the status control 140 may iterate as  
20 follows:

21 Status Control Started...

22 Key: Title1 Option: Title Level: 1 Filter: Field: Title

23 Key: A2 Option: A Level: 2 Filter: SUBSTRING([Title],1,1)='A' Field: Title

24 Key: AA3 Option: AA Level: 3 Filter: SUBSTRING([Title],1,2)='AA' AND  
25 SUBSTRING([Title],1,1) = 'A' Field: Title

26 Key: F4 Option: F Level: 4 Filter: SUBSTRING([Title],1,1)='F' Field: Title

1 Key: Fa5 Option: Fa Level: 5 Filter: SUBSTRING([Title],1,2)='Fa' AND  
 2 SUBSTRING([Title],1,1)='F' Field: Title  
 3 Key: Favo6 Option: Favo Level: 6 Filter: SUBSTRING([Title],1,4)='Favo'  
 4 AND SUBSTRING([Title],1,2)='Fa' AND SUBSTRING([Title],1,1)='F' Field: Title  
 5 Key: C7 Option: C Level: 7 Filter: SUBSTRING([Title],1,1)='C' Field: Title  
 6 Key: Ce8 Option: Ce Level: 8 Filter: SUBSTRING([Title],1,2)='Ce' AND  
 7 SUBSTRING([Title],1,1)='C' Field: Title  
 8 Key: Cells9 Option: Cells Level: 9 Filter: SUBSTRING([Title],1,5)='Cells'  
 9 AND SUBSTRING([Title],1,2)='Ce' AND SUBSTRING([Title],1,1)='C' Field: Title  
 10 Key: Cellula10 Option: Cellula Level: 10 Filter: SUBSTRING([Title],1,7)=  
 11 'Cellula' AND SUBSTRING([Title],1,2)='Ce' AND SUBSTRING([Title],1,1)=  
 12 'C' Field: Title  
 13 Key: CC11 Option: CC Level: 11 Filter: SUBSTRING([Title],1,2)='CC'  
 14 AND SUBSTRING([Title],1,1)='C' Field: Title  
 15 Status Control Terminated.

16 Figure 15b shows the results for a search for a low-fat cookbook using the search  
 17 engine 125 as applied to a remote database. In this example, the remote database is coupled  
 18 to a Barnes & Noble web page. The first query, and resulting message strings, are illustrated  
 19 by the following:

20 Query Analyzer  
 21 Message Received: ACK  
 22 Status Control: Refresh  
 23 Dispatcher  
 24 Message Sent: Categories~~~Title~~~Author~~~ISBN~SubTitle~Format~Date  
 25 P u b l i s h e d ~ S t o c k S t a t u s ~ R e c o m m e n d e d  
 26 A g e ~ P a g e s ~ R a t i n g s ~ P r i c e ~ R e t a i l ~ S a v i n g s ~ ~ ~ P u b l i s h e r  
 27 Query Analyzer



1 Message Received: CLK#0#1#Categories  
2 Status Control received an update:  
3 Key: Categories1 Option: Categories Level: 1 Filter: Field: Categories  
4 Query Generator  
5 Request is not cached, processing  
6 Generated Query: SELECT DISTINCT [Categories] FROM Books ORDER BY  
7 [Categories]  
8 Number of Matching Records: 2032  
9 Generated Query: SELECT DISTINCT SUBSTRING([Categories],1,82) FROM Books  
10 ORDER BY SUBSTRING([Categories],1,82)  
11 Number of Matching Records: 2022  
12 Generated Query: SELECT DISTINCT SUBSTRING([Categories],1,61) FROM Books  
13 ORDER BY SUBSTRING([Categories],1,61)  
14 Number of Matching Records: 1995  
15 Generated Query: SELECT DISTINCT SUBSTRING([Categories],1,45) FROM Books  
16 ORDER BY SUBSTRING([Categories],1,45)  
17 Number of Matching Records: 1751  
18 Generated Query: SELECT DISTINCT SUBSTRING([Categories],1,33) FROM Books  
19 ORDER BY SUBSTRING([Categories],1,33)  
20 Number of Matching Records: 1251  
21 Generated Query: SELECT DISTINCT SUBSTRING([Categories],1,24) FROM Books  
22 ORDER BY SUBSTRING([Categories],1,24)  
23 Number of Matching Records: 799  
24 Generated Query: SELECT DISTINCT SUBSTRING([Categories],1,18) FROM Books  
25 ORDER BY SUBSTRING([Categories],1,18)  
26 Number of Matching Records: 425

1 Generated Query: SELECT DISTINCT SUBSTRING([Categories],1,13) FROM Books  
2 ORDER BY SUBSTRING([Categories],1,13)  
3 Number of Matching Records: 319  
4 Generated Query: SELECT DISTINCT SUBSTRING([Categories],1,9) FROM Books  
5 ORDER BY SUBSTRING([Categories],1,9)  
6 Number of Matching Records: 147  
7 Generated Query: SELECT DISTINCT SUBSTRING([Categories],1,8) FROM Books  
8 ORDER BY SUBSTRING([Categories],1,8)  
9 Number of Matching Records: 111  
10 Generated Query: SELECT DISTINCT SUBSTRING([Categories],1,7) FROM Books  
11 ORDER BY SUBSTRING([Categories],1,7)  
12 Number of Matching Records: 78  
13 Generated Query: SELECT DISTINCT SUBSTRING([Categories],1,6) FROM Books  
14 ORDER BY SUBSTRING([Categories],1,6)  
15 Number of Matching Records: 44  
16 Generated Query: SELECT DISTINCT SUBSTRING([Categories],1,5) FROM Books  
17 ORDER BY SUBSTRING([Categories],1,5)  
18 Number of Matching Records: 26  
19 Truncator finished, took 15 seconds to make 13 iterations  
20 Caching this request...  
21 Dispatcher  
22 Message Sent: Afric~Art, ~Biogr~Busin~Compu~Cooki~Engin~Enter~Ficti~Histo~Home  
23 ~Horro~Kids!~Law:  
24 ~Medic~Mind,~Nonfi~Paren~Poetr~Refer~Relig~Scien~Small~Sport~Trave~Write~  
25 Query Analyzer  
26 Message Received: CLKCategories

1           In the example illustrated by Figure 15b and the above-listed message strings, an initial  
2 request would have returned 2032 book titles for cook books. This number of entries may be  
3 too large. Accordingly, the truncator 152, through 13 iterations, reduces the entries in a result  
4 list to 26. The entries in the truncated result list can then be easily reviewed by the user, and  
5 further searches may be performed to identify a desired book. As can be seen above, the user  
6 has selected "Categories" as a data field to search. As is also shown in Figure 15b, the search  
7 engine 125 may display other information windows, such as book availability, ordering and  
8 shipping information windows. With a simple drag-and-drop cursor operation, for example,  
9 the user may then order and pay for the desired book.

10           Figure 16 - 20 are flow charts illustrating operations of the search engine 125. Figure  
11 16 is a flowchart of an overall search routine 250. The process starts in block 251. The  
12 request analyzer 130 receives the request 114, block 252. The request 114 may be made  
13 using a hierarchical menu-based display or a graphical user interface, with one or more layers.  
14 Using either the menu or the GUI, the user may enter specific details by typing, selection of  
15 iconic symbols or pre-formatted text, and by using well-known data entry techniques, for  
16 example. The request 114 may also comprise a simple text or voice query. Use of voice  
17 recognition may be particularly useful in mobile environments, and to speed access to the  
18 database 12. Use of voice recognition may include simple commands, such as UP, DOWN,  
19 and SELECT, to select search terms from a pre-formatted list that is presented to the user at  
20 the terminal 14. More sophisticated use of voice recognition may include actually speaking  
21 letters or numbers, or full search terms, such as speaking a key word for a key word search,  
22 for example.

23           The protocol analyzer 133 provides an output 135 to the constraint collator 136, and  
24 the constraint collator 136 determines the nature of the request, block 254. If the request 114  
25 is a refresh request (i.e., a command to initiate the refresh function), the constraint collator 136  
26 sends a reset command 131 to the database qualifier 160. The updated request 115 (possibly  
27 with a new constraint) is then sent to the query analyzer 150 for further processing, including

1 analyzing the database 12, retrieving field descriptors, and formatting, block 256. The result  
2 of the data field descriptor retrieval and formatting are shown as an available data fields result  
3 list, block 258, and is returned to the terminal 14, block 260.

4 In block 254, if the request 114 is not a refresh request, the constraint collator 136  
5 provides the updated request 115 (which may be an initial request, or a subsequent request)  
6 to the query generator 150, block 264. The constraint collator 136 compares the request 114  
7 against information stored in the status control 140. In particular, the constraint collator 136  
8 sends the request status control signal 118 to the status control 140 and receives the request  
9 status response 119. The constraint collator 136 then compares the request status response  
10 119 to constraint information provided with the request 114 to determine if the constraint status  
11 should be updated (e.g., because the request 114 includes a new constraint). If the constraint  
12 status should be updated, the constraint collator 136 calls create new constraint subroutine  
13 270, and creates new constraints.

14 The create new constraints subroutine 270 is shown as a flowchart in Figure 17. The  
15 subroutine starts at 272. In block 274, the constraint collator 136 determines if the request is  
16 for a sort-on-the-fly operation. If sort-on-the-fly has been selected, field assessor 162  
17 prepares a new set of data fields, block 280. The new set of data fields are then sent to the  
18 query generator 150, block 284, and the subroutine 270 ends, block 286.

19 In block 274, if sort-on-the-fly was not selected, the request analyzer 130 may receive  
20 a key word constraint, block 276. The query generator 150 will then generate an input  
21 window in which the user may enter a desired key word, block 282. Alternatively, the query  
22 generator 150 may prompt the user to enter a key word using voice recognition techniques, or  
23 any other way of entering data. The process then moves to block 284. In block 276, if a key  
24 word search option was not selected, the constraint collator 136 enters the new constraint to  
25 the existing list of constraints, block 278. The process then moves to block 284.

26 Returning to Figure 16, the constraint collator 136 next updates the status control 140,  
27 block 290. In block 292, using the updated constraints, the query generator 150 generates a

1 next query of the database 12, block 292. The database driver 170 then extracts the result list  
2 from the database 12, according to the latest query, block 294. In block 296, the truncator  
3 152 determines if the result list may be displayed at the terminal 14. If the result list cannot be  
4 displayed, the process moves to block 298, and a truncation routine is executed. The process  
5 then returns to block 294. If the result list in block 296 is small enough, the result list is  
6 provided by the dispatcher 154 to the terminal 14, block 258.

7 As noted above, the request analyzer 130 determines the nature of the request,  
8 including any special commands. A special command may include a command to conduct a  
9 search-on-the-fly. Alternatively, the search engine 125 may adopt a search-on-the-fly  
10 mechanism as a default value. The search engine 125 also may incorporate other special  
11 search commands, such as a Boolean search, for example.

12 Figures 18 - 20 are flowcharts illustrating alternate truncation subroutines 298. In  
13 Figure 18, the subroutine 298 adjusts a size of a data field by decrementing a parameter TP  
14 related to entries in a selected data field. For example, if the data field comprises a list of U.S.  
15 cities by name, the parameter TP may be the number of alphabetical characters in a name. The  
16 results of such a truncation is shown in the example of Figure 4. The subroutine 298 starts at  
17 block 301. In block 303, the parameter TP is set to equal a size of the data field being  
18 searched. The truncator 152 then determines the list of records sized by the parameter TP,  
19 block 305. In block 307, the truncator 152 determines if the result list can be displayed at the  
20 terminal 14. If the result list cannot be displayed at the terminal 14, the truncator 152  
21 decrements the parameter TP, block 309. Processing then returns to block 305, and the  
22 truncator 152 gets a reduced result list using the truncated parameter TP. If the result list can  
23 be displayed at the terminal 14, the process moves to block 311 and the subroutine 298 ends.

24 Figure 19 is a flowchart illustrating an alternate truncation routine 298. The process  
25 starts in block 313. In block 315, the truncator 152 sets the parameter TP to a size of the data  
26 field being searched. In block 317, the truncator 152 determines the list of records sized by  
27 the parameter TP. In block 319, the truncator 152 determines if the result list can be displayed

1 at the terminal 14. If the result list cannot be displayed, the truncator 152 adjusts the size of  
2 the data field by dividing the parameter TP by a set amount, for example, by dividing the  
3 parameter TP by two, block 321. Processing then returns to block 317, and repeats. If the  
4 result list can be displayed at the terminal 14, the process moves to block 323 and the  
5 subroutine ends.

6 Figure 20 shows yet another alternative truncation subroutine 298. The process starts  
7 in block 325. In block 327, the truncator 152 sets the parameter TP to equal the size of the  
8 data field being searched. In block 329, the truncator 152 determines the list of records sized  
9 by the parameter TP. The truncator 152 then determines if the result list can be displayed at  
10 the terminal 14, block 331. If the result list cannot be displayed at the terminal 14, the  
11 truncator 152 determines if the parameter TP is less than ten, block 333. If the parameter TP  
12 is not less than ten, the truncator 152 adjusts the parameter TP by multiplying the parameter  
13 TP by a number less than one, block 337. In an embodiment, the number may be 3/4. The  
14 process then returns to block 329 and repeats. In block 333, if the value of the parameter TP  
15 is less than ten, the truncator 152 decrements the parameter TP by one, block 335. Processing  
16 then returns to block 329 and repeats. In block 331, if the list can be displayed at the terminal  
17 14, the process moves to block 339 and the subroutine 298 ends.

18 The examples illustrated in Figures 18 - 20 are but a few examples of the truncations  
19 subroutine. One of ordinary skill in the art could conceive of other methods to adjust the field  
20 size. In addition to using a truncation subroutine, the user may specify a limit for the field size.

21 As noted above, the search engine 125 may be used for multiple searches and may be  
22 used to search multiple databases, including databases with different schemas. The results of  
23 individual searches, including the control data provided in the status control 140, are saved.  
24 The search engine 125 may then be used to further sort (search), or otherwise operate on, the  
25 results of these multiple searches. In an embodiment, the search engine 125 may perform a  
26 Boolean AND operation on two search results. The result of the Boolean AND operation

1 would be a list of records, or entries, that are common to the two search results. Figure 21  
2 illustrates such a Boolean AND operation.

3 In Figure 21, a GUI 400 displays local database selections 410, including a database  
4 of recordings (compact discs - CDs) 412 and a database of contacts 414. The databases 412  
5 and 414 may be shown by text descriptions and an appropriate icon, for example. The  
6 database selections in this example are resident on a user's terminal, such as the terminal 14  
7 shown in Figure 1. Also displayed on the GUI 400 is a remote database selection 420 that  
8 represents databases, such as the databases 13 and 15 shown in Figure 1, that are located  
9 remotely from the terminal 14. In the example shown in Figure 21, the remote database  
10 selection 420 includes a database 422 for online record sales, which is represented by an icon  
11 (a CD) and a text title of the online retailer. The remote databases shown in the remote  
12 database selection 420 may include those databases for which the user has already established  
13 a link. In the example shown, the user may already have entered an Internet address for the  
14 online retailer. In addition to any returned web pages from the online retailer, the terminal 14  
15 may then display a representation of the database 422.

16 Continuing with the example, the user may use the search engine 125 to conduct a  
17 search-on-the-fly of the recordings database 412 and the virgin records database 422. The  
18 user may search both databases 412 and 422 for titles of recordings that are classified as  
19 "blues." The search engine 125 may return search results 416 and 424 for searches of both  
20 databases 412 and 422, respectively. The search results 416 and 424 may be displayed in a  
21 window section 430 of the GUI 400. The results 416 and 424 may also be represented by CD  
22 icons, such as the icons 432 and 434. The search results 416 and 424 may be stored as lists  
23 in one or more temporary databases, as represented by the windows 417 and 427. The search  
24 results 416 and 424 may also be stored in a scratch pad database 418. At this point, the user  
25 may wish to determine which recordings from the list 424 are contained in the list 416. The  
26 search engine may support this function by performing a Boolean AND operation of the lists  
27 416 and 424. The results of the Boolean AND operation are represented by the icon 436

1 displayed in the window 430. To execute the Boolean AND operation, the user may simply  
2 drag the icon 432 over the icon 434, and then select AND from a pop-up menu 438 that  
3 appears when the icons 432 and 434 intersect. Other techniques to execute the Boolean AND  
4 (or another Boolean function) may include typing in a command in a window, using voice  
5 recognition techniques, and other methods. In addition, other Boolean functions may be used.

6 The result represented by the icon 436 of the Boolean AND operation may then be  
7 stored in a database at the terminal 14, such as in the scratch pad database 418 or may be  
8 stored at another location. The result may then be subjected to further search-on-the-fly  
9 operations.

10 Also shown in Figure 21 is an online-purchase module 435 that may be used to  
11 consummate purchase of a product referenced in an online database such as the database 422.  
12 To initiate such a purchase, the user may drag an iconic or text representation of a desired  
13 product listed in the search result 424 over an icon 436 in the online-purchase module 435.  
14 This drag-and-drop overlaying these icon may initiate and complete the online purchase for the  
15 desired product.

16 Use of the search engine 125 may be facilitated by one or more GUIs that are  
17 displayed on the terminal 14. Figures 22 - 26 are examples of such GUIs. In Figure 22, a  
18 GUI 450 includes a display section 452 and one or more database sections such as local  
19 database section 470 and remote database section 460. The local database section 470  
20 includes databases local to the terminal 14. In the example shown, the local databases include  
21 a patients database 472, a general contacts database 474, a pharmacy database 476, a  
22 medicines database 478 and a scratch pad database 480. The remote databases include an  
23 Amazon.com database 462, an online record retailer database 464, a Physician's Desk  
24 Reference database 466 and an American Medical Association (AMA) online database 468.  
25 The remote and local databases may be represented by a text title and an icon, both contained  
26 in a small window as shown. A user may access one of the remote or local databases by  
27 moving a cursor over the desired window and then selecting the database. In the example



1 shown, the local medicines database 478 has been selected, and a list 490 of data fields in the  
2 medicines database 478 is displayed in the display section 452. Also included on the display  
3 section 452 is a keyword button 492 that may be used to initiate a key word search of the  
4 medicines database 478.

5 Figure 23 shows the GUI 450 with a user selection of a category data field from the  
6 list 490. The category data field is indicated as selected by an arrow adjacent to the data field  
7 name. When the category data field is selected, a category list 494 is displayed on display  
8 section 452. The category list 494 includes four entries, as shown.

9 The user may continue to search the medicines database 478 using key word  
10 techniques and search-on-the-fly techniques. Figure 24 shows the GUI 450 with results of  
11 several search cycles displayed.

12 Figure 25 illustrates a search of the PDR database 466. Such a search may be initiated  
13 by dragging a cursor to the window having the PDR 466 symbol (text or icon), and then  
14 operating a "select" button. Figure 26 shows a search of the Amazon database 462. This  
15 search may also be initiated by a "drag-and-drop" operation.

16 The SOTF search engine 125 may accommodate merging of one or more sets of search  
17 results. The multiple search results may be derived from a common database, or from more  
18 than one database. A search using the search engine 125 may be controlled through a user  
19 interface by one or more icons that can represent (1) filters or (2) the images of filters. Thus,  
20 the icon may represent spatial or temporal attributes, or sets of objects or procedures.  
21 Merging the icons thus has two interpretations corresponding to (1) and (2): either filters are  
22 added ("apply every filter in every icon to every image to which it can be applied"), or image  
23 sets are added. In an alternative embodiment, the addition (union or join) operator may be any  
24 other relational operator, e.g. divide, difference.

25 Use of the merge function may be explained by reference to Boolean lattices. A  
26 collection of entities can have attributes A or B or both. If {A} is the set of all A entities and  
27 {B} is the set of all B entities; the set whose elements all possess both attributes A and B may

1 now be written  $\{A \text{ and } B\}$ , and the set whose elements all possess either attribute A or  
2 attribute B or both may be written  $\{A \text{ or } B\}$ . The elements of  $\{A \text{ and } B\}$  can be considered  
3 to possess a new, less inclusive or specific attribute C, and the elements of  $\{A \text{ or } B\}$  to possess  
4 a new, more inclusive or general attribute D. In a lattice, the nodes are attributes; the most  
5 inclusive attribute (in this case D) is always at the top and is called the join of those attributes  
6 (nodes) immediately below it, and the most exclusive attribute (in this case C) is always at the  
7 bottom and is called the meet of those attributes (nodes) immediately above it. In other words,  
8 the OR operation yields the join of two attributes, while the AND operation yields their meet.  
9 Thus, the OR operator is upward or inductive (yielding the more inclusive join of the operands),  
10 while the AND operation is downward or deductive (yielding the more exclusive meet of the  
11 operands). The nodal attributes of such a lattice are analogous to filters; but since a principle  
12 called CF duality states that attributes and sets are to some extent interchangeable because  
13 every attribute characterizes a set and every set is characterized by an attribute, these attributes  
14 are logically equivalent to the sets they characterize.

15 In an example optical context, the downward AND operator corresponds to stacking  
16 colored filters, while the upward OR operator corresponds to mixing colored paints or filters.  
17 In color optics, stacking and unstacking colored lenses is called a subtractive process, while  
18 mixing or unmixing paints is called an additive process. Unfortunately, while combining or  
19 "adding" filters is subtractive with respect to the sets they characterize, it is additive with respect  
20 to the filters themselves, and adding sets is subtractive with respect to the filters. So it is better  
21 to refer to operations among attributes (filters, lenses, etc.) as "filtrative" or "infonegative", and  
22 to those among sets (paints, lights, etc.) as "constructive" or "infopositive". CF duality can now  
23 be rephrased as follows: every infonegative entity (attribute) descriptively characterizes an  
24 associated infopositive entity (set/object), and every infopositive entity instantiates or is  
25 descriptively characterized by an associated infonegative entity.

26 The search engine 125 includes iconization (iconic representation) of an algebra or  
27 calculus of relations defined on Boolean lattices. This representation begins with a set of

1 primitive icons extracted from base tables and defines new icons (derived tables, virtual  
2 databases) by means of simple user-executed operations. The icons can be effortlessly  
3 translated into lists of data corresponding to the icons, and it is these lists that comprise the real  
4 substance of any search procedure.

5 When search chains are branched into to chains A and B, the filters subsequently  
6 applied to each chain can be the same or different, and merging can signify any of two or more  
7 Boolean relationships (relational operations) defined on a relational database. Specifically,  
8 when chains merge, sets of filters can be added or intersected. Since filters are constraints,  
9 adding them amounts to intersecting their images, while adding their images amounts to  
10 intersecting the filters (infopositive-infonegative distinction). Equivalently, one may consider  
11 positive and negative filters effecting deduction and induction respectively; the filters are  
12 descriptive, while the images are substantive. The extent to which the images of filters can  
13 intersect depends on the commonality (predicative non-exclusivity) of domains. Icon algebras  
14 (of iconic operators) are "object-oriented" on the GUI level; they are UI extensions of the innate  
15 object-orientation of relational databases themselves, wherein the objects are records,  
16 attributes, tables, virtual databases and so on, and the operations are those of any relational  
17 algebra.

18 The looping and merging of search chains is to some extent algebraic. First, since  
19 actual topology is being changed, such transformations do not directly form a topological  
20 homeomorphisin group; the algebra remains Boolean, and the "homeomorphism" is defined on  
21 the operator graph of the Boolean algebra (of which the initial search tree is generally only a  
22 subspace). Icons representing sets of nested predicates are "Boolean objects"; when decision  
23 chains converge or diverge, objects merge or split, and these objects represent  
24 (combinatorially) unique search paths. Thus, operations among paths can be reduced to  
25 operations among objects; e.g., regress-diverge is just an object-splitting operation.  
26 Continuous looping applies "inverse deductive filters" to achieve induction by descriptive  
27 intersection of filter constraints, permitting the retrograde convergence of paths to identical

1 ancestral objects (inductive merging of objects), while inductive looping is just direct regression  
2 to an ancestral object preparatory to splitting it and thus effecting divergence of paths  
3 (deductive splitting of objects). Deductive convergence of paths is "natural" if iconic image sets  
4 intersect and "forced" if not; if natural, then there has been non-exclusivity of subobjects, and  
5 paths are not unique (even though identical filters can apply to divergent paths without  
6 impairing uniqueness). So all deductive merging is forced, and this entails a decision regarding  
7 which filters are to be conserved and which discarded. Any such operation will effectively  
8 "rewrite the paths", and doing this optimally is NP-complete.

9 More specifically, icons are subject to CF duality. The merge control thus has a  
10 "switch" toggling between "Qualities / Objects". When the switch is in the "qualities" position,  
11 merging icons performs a qualified deductive conjunction of filters and yields a set intersect;  
12 when it is in the "objects" position, merging the icons performs a disjunction of filters and an  
13 inductive union of sets, yielding a more general attribute (the general qualities created by the  
14 object-merge operation will be produced by sets of filters applied disjunctively). The search  
15 engine 125 is therefore capable of inductive and deductive information processing. A quality-  
16 merge in which filters do not cross the line between composite icons equates to an object  
17 merge; the set thus selected is characterized by a more general quality which amounts to the  
18 descriptive (filtrative) union. There is also a modified quality-merge in which filters in either icon  
19 applicable to both iconized sets are applied to both, thus crossing the line between icons. In  
20 this case, a true merging of paths occurs, as opposed to path icons. The search engine 125  
21 allows users to choose which filters are to cross the inter-icon line and which are not, resulting  
22 in complex Boolean expressions and the sets they characterize (determining consistency of  
23 complex expressions can amount to LSAT; sets of inconsistent expressions will simply yield a  
24 null return.

25 Icons may reside in the first menu box to appear, being transferred from menu to menu  
26 as the path is generated and filters are accumulated. When a direct regress occurs, the path  
27 is regarded as "complete" and is stored in a holding module. Prior to the merging operation,

1 the quality/object switch is set; and icon subfilters or subsets individually displayed. A "lattice  
2 navigator" will keep track of position and equivalence, folding the search graph in case a node  
3 of the original tree is inductively encountered in the course of an object-merge; otherwise, the  
4 icon remains in "internodal space" (which is to be regarded as a virtual space realized only in  
5 the event that the search tree is nondisjunctive in its nodes and therefore incomplete with  
6 respect to the semantic net generated by the tree).

7 Figure 27 is a flow chart illustrating an alternative operation 600 of the query generator  
8 150 of Figure 6. In the illustrated operation, the query generator 150 is adapted to receive  
9 multiple selections of items within a same menu function and within a same merge function. To  
10 provide this functionality of the query generator 150, the request analyzer 130 (see Figure 5)  
11 may be adapted to receive a collection of user choices.

12 The operation 600 begins in block 601. In block 603, the request analyzer 130  
13 receives constraints collected from the constraint collator 136, and the updated request 115,  
14 which may be an initial request or a subsequent request, is provided to the query generator  
15 150. In block 605, the query generator 150 determines if the constraints (the request 115) are  
16 in the same merge group. If the query generator 150 determines that the request 115 is in the  
17 same merge group, the process moves to block 607 and the query generator 150 generates  
18 the query with a Boolean AND. If the request is not in the same merge group, the query  
19 generator 150 generates the query with a Boolean OR, block 609.

20 In block 611, the items selected within the same unit are Or'ed and the default  
21 truncator may be used depending on the size of the returned items. In block 613, the generated  
22 query is executed. In block 615, the number of records to be displayed is checked. If the  
23 number is within a specified limit, the process moves to block 617 and the search results are  
24 returned for display. The operation 600 then ends, block 625. In block 6125, if the number  
25 of records to be displayed is too large, the process moves to block 621, and a truncation  
26 routine is executed.

1           The truncation routine may be any of the previously-described truncation routines.  
2           Figure 28 illustrates an alternate truncation routine 630. The routine 630 begins in block 631  
3           with the truncator 152 receiving the request 115. In block 633, the truncation is set to the size  
4           of the field being viewed on the GUI, and sets the False Flag. The query is then run against the  
5           database using the selected truncator, block 635. In block 635, the truncator 152 determines  
6           if the number of records that would be retrieved from the database can be displayed on the  
7           existing GUI. If the records can be displayed, the process moves to block 639, and the  
8           truncator 152 determines if the Flag is set False. If the Flag is set False, the process moves to  
9           block 653 and the records are returned (displayed on the GUI). The process then ends, block  
10          655. In block 637, if the number of records exceeds the display size of the GUI, the status of  
11          the Flag is checked as False. If false, the truncator is set to 1, and the flag is set to true, block  
12          647, and the process returns to block 635. If in block 637. If the flag is not set false, the  
13          process moves to block 651, and saved records are retrieved. The retrieved records are then  
14          displayed, block 653.

15           In block 639, if the Flag is not set to false, the retrieved records are saved, and the  
16          truncator 152 is incremented. The process then returns to block 635.

17           Figures 29 - 38 illustrate graphical user interfaces and search on the fly results using the  
18          search engine 125 with a merge function. In Figure 29, a search of a patent database has been  
19          executed to search for patents by primary examiner. The Primary Examiner results table lists  
20          the arabic numerals 0 - 7 and the letters A-Z, indicating that the database contains names of  
21          primary examiners beginning with these numerals/letters. To quickly narrow the search, the  
22          user selects the letter O, and results are returned listing last and first names all primary  
23          examiners whose last name begins with O. As can be seen by the returned results, the  
24          database lists several primary examiner instances of O'Dea. This could indicate an error in the  
25          database. The search engine 125 allows these errors to be detected and corrected. The  
26          correction may be made by selecting the incorrect instances, right-clicking the correct instance,  
27          and then choosing a 'correct all other's based on this instance" function.

1           Figure 30 shows how multiple-select capabilities of the search engine 125 may be used  
2           to enhance a search. In the illustrated example, the user searches for 3 M Company. Different  
3           versions of the company name are then displayed with the returned results. In this way, the  
4           user may select the different versions of the company that the user wants to use for the search.  
5           The pop-up pane shows a current status control for the GUI.

6           Figure 31 shows the results of subsequent menus showing the aggregation, or merge,  
7           of two previous constraints, “3m” and 3-M.” Figure 32 shows a merge execution. The user  
8           first selects the “3-M” and the “3M” company names using the check boxes in the previous  
9           menu. The user then selects the merge option, placing the menu on hold, and going to the “M”,  
10          “MP”, “MIN” and “MINNESOTA M” menus. The merge option is then selected on the menu  
11          and the merged menu is displayed showing the merge of searches between “3M” and  
12          “Minnesota Mining and Manufacturing Co.” Figures 32 - 36 show other search engine 125  
13          features including data mining and database correction.

14          Figures 37 - 39 show the results of a full text search of a patent database using the  
15          keyword “encryption” and searching on all fields. The initial search results are truncated to  
16          display by first letter/numeral of the patent title. From this intermediate search result menu, the  
17          user selects all patents whose title begins with the letter “E”, and a subsequent search result  
18          menu is displayed listing partial titles of all such patents. From the next intermediate list, the  
19          user selects the patent whose title begins “Electronic copy protection mechanis.” The search  
20          engine 125 then returns this specific patent, the first page of which is shown in Figure 39. The  
21          displayed patent includes the keyword “encryption” highlighted wherever it occurs. The display  
22          also indicates the number of instances of the keyword in the patent.

23          Figures 40-49 illustrates additional search results.

24          In specific embodiments, the search engine 125 is implemented as a program executed  
25          on a general purpose computer, such as a personal computer. The search engine may also be  
26          implemented as a routine attached to a database structure. In addition, the search engine may  
27          be implemented on any processor capable of executing the routines of the program. In

1 alternative embodiments, the search engine 125 may be implemented as a single special  
2 purpose integrated circuit (e.g., ASIC) having a main or central processor section for overall,  
3 system level control, and separate circuits dedicated to performing various different specific  
4 functions, computations and other processes under control of the central processor section.  
5 Those of ordinary skill in the art will appreciate that the search engine 125 may also be  
6 implemented using a plurality of separated dedicated or programmable integrated circuits, or  
7 other electronic circuits or devices (e.g., hardwired electronic or logic circuits such as discrete  
8 elements circuits, or programmable logic devices, such as PLDs, PLAs, or PALs). In general,  
9 any device or assembly of devices on which a finite state machine capable of implementing  
10 flowcharts similar to the flowcharts of Figures 16 - 20 and 27 and 28 can be used to implement  
11 the search engine 125.

12 The terms and descriptions used herein are set forth by way of illustration only and are  
13 not meant as limitations. Those skilled in the art will recognize that many variations are possible  
14 within the spirit and scope of the invention as defined in the following claims, and there  
15 equivalents, in which all terms are to be understood in their broadest possible sense unless  
16 otherwise indicated.

17



- 1 In the claims:
- 2 1. A method for searching databases, comprising:
- 3 determining a database schema for a database;
- 4 providing a list of database fields, wherein the list includes a descriptor indicating a data
- 5 category;
- 6 receiving a search selection for a database field on the provided list of database fields;
- 7 determining a quantity of entries in the selected database field;
- 8 if the quantity exceed a specified amount;
- 9 truncating data, and
- 10 displaying the truncated data; and
- 11 if the quantity does not exceed the specified amount, displaying contents of the
- 12 database field.
- 13 2. The method of claim 1, further comprising providing a key word search.
- 14 3. A method for searching a database, comprising:
- 15 generating a list of data fields;
- 16 receiving a first data field selection from the list of data fields;
- 17 determining a first quantity indicative of a number of entries of the selected data field;
- 18 if the first quantity exceeds a specified limit, reducing a size of data to be displayed from
- 19 the selected data field; and
- 20 displaying data from the selected data field.
- 21 4. The method of claim 3, wherein the specified limit is fixed.
- 22 5. The method of claim 3, wherein the specified limit is variable.
- 23 6. The method of claim 3, wherein the data are displayed on a terminal, and wherein the
- 24 specified limit is determined dynamically, based on a characteristic of the terminal.
- 25 7. The method of claim 3, wherein the specified limit is a user-determined limit.
- 26 8. The method of claim 3, wherein the method for reducing the size of the data to be
- 27 displayed from the selected data field comprises:

1 performing a truncation that reduces the size of the data to be displayed from the  
2 selected data field;  
3 comparing the reduced size to the specified limit; and  
4 if the reduced size exceeds the specified limit, repeating the truncation and comparing  
5 steps until the size of the data to be displayed from the selected data field is less than or equal  
6 to the specified limit.

7 9. The method of claim 8, wherein a parameter is related to the size of the data to be  
8 displayed from the selected data field, and wherein the truncation comprises decrementing the  
9 parameter.

10 10. The method of claim 9, wherein the parameter is decremented by a value of one.

11 11. The method of claim 8, wherein a parameter is related to the size of the data to be  
12 displayed from the selected data field, and wherein the truncation comprises dividing the  
13 parameter by a value.

14 12. The method of claim 11, wherein the value is two.

15 13. The method of claim 8, wherein a parameter is related to the size of the data to be  
16 displayed from the selected data field, and wherein the truncation comprises multiplying the  
17 parameter by a value.

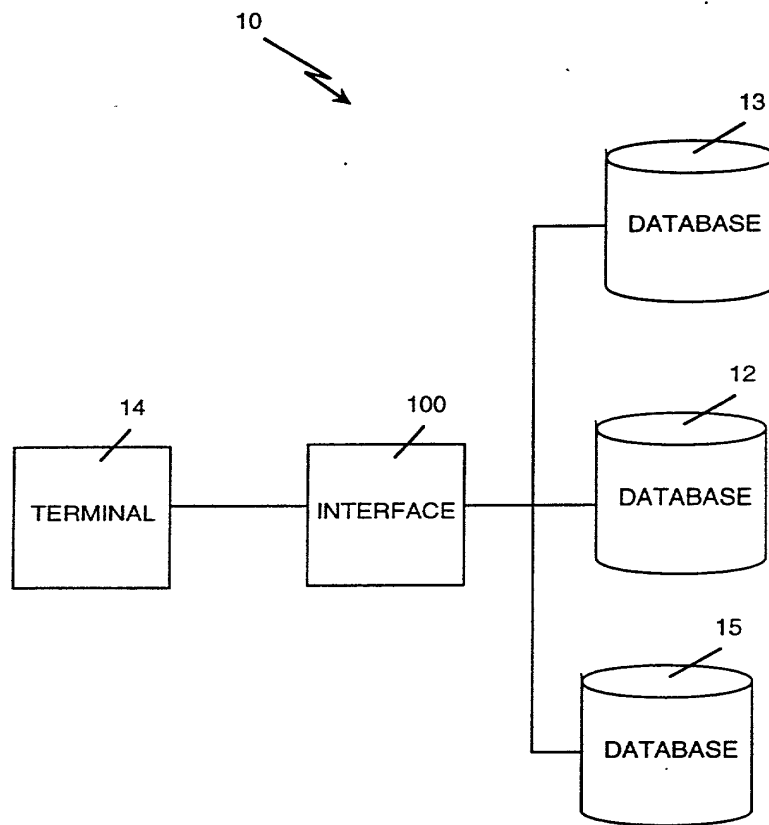
18 14. The method of claim 3, further comprising:

19 receiving a first constraint, wherein the first constraint is related to a data element in a  
20 data field; and

21 receiving one or more subsequent constraints, wherein search results are generated  
22 based on a combination of the first and the one or more subsequent constraints.

**ABSTRACT**

1  
2 A Sort-on-the-Fly/Search-on-the-Fly search engine provides an intuitive means for  
3 searching databases, allowing a user to access data in the database without having to know  
4 anything about the database structure. A user selects a desired search term, and the search  
5 engine searches the database for all instances of the desired term, even if a specific file or table  
6 does not contain the instance. The database need not have a specific file (in a flat database)  
7 or a table (in a relational database) of names. The user may specify other criteria, or  
8 constraints to narrow the search results, or for other reasons. The search engine then conducts  
9 a further search using this criteria and produces a second search result. Further narrowing or  
10 broadening of the search are permitted, with the search-on-the-fly search engine returning  
11 results based on any new constraints. If the returned data would be too large to be  
12 conveniently displayed at a terminal, the search engine executes a truncation routine so that the  
13 returned data is easily displayed.



*Fig. 1*

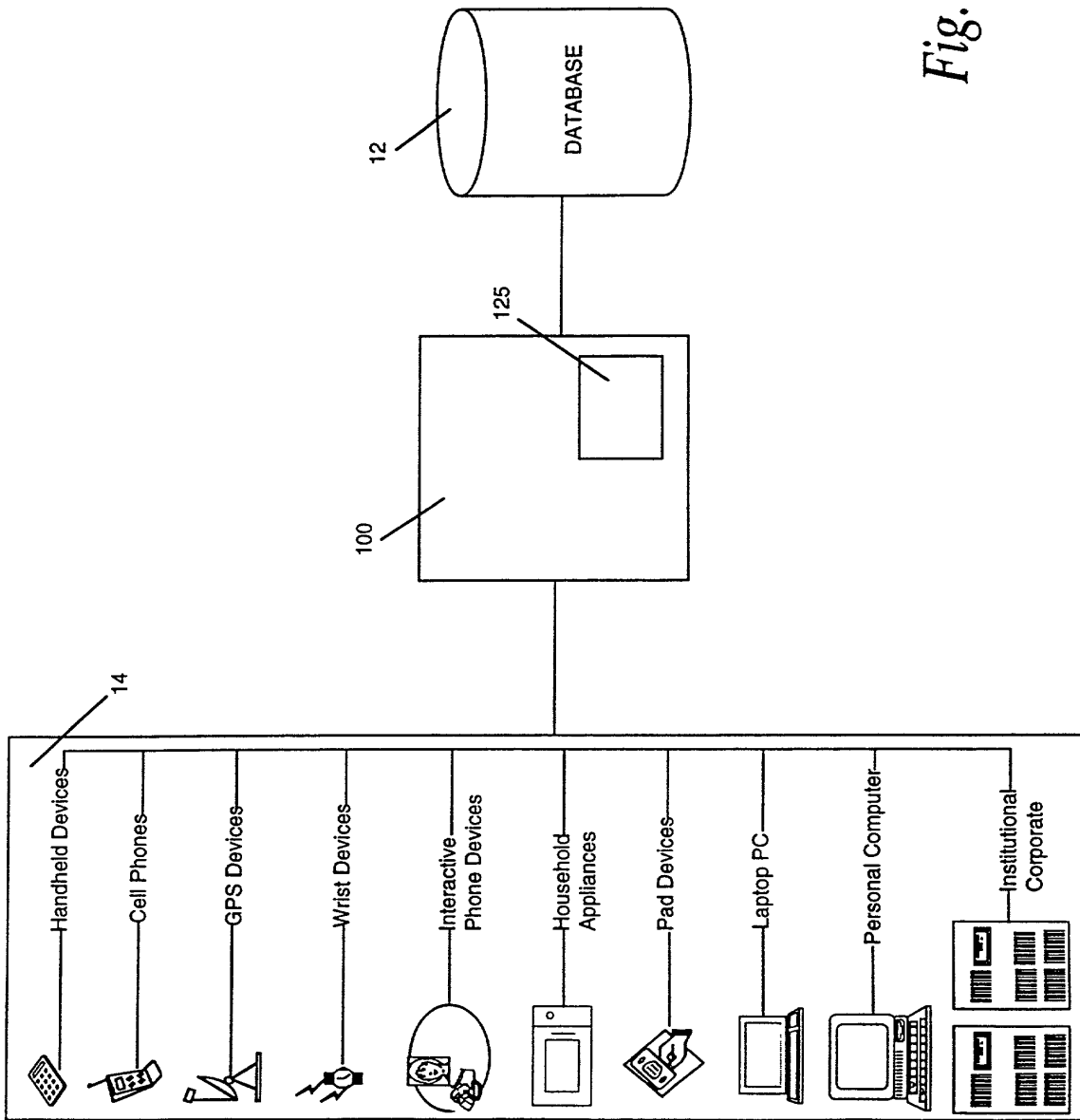


Fig. 2

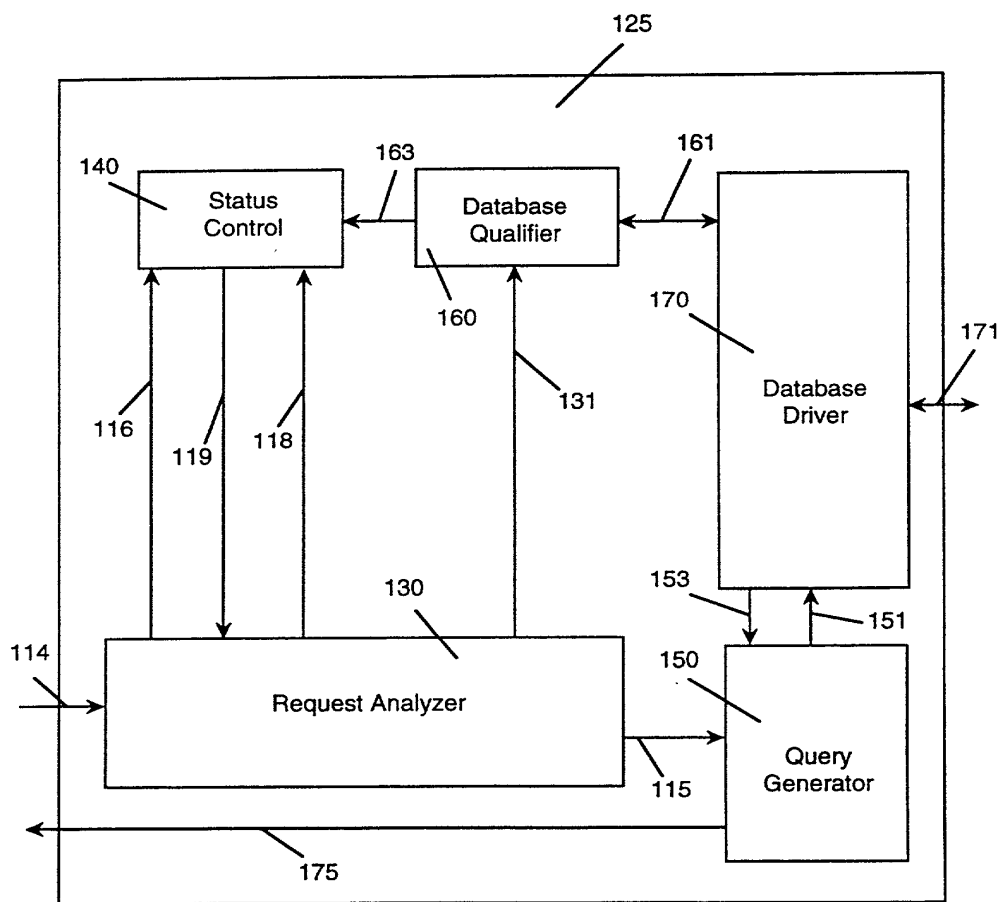
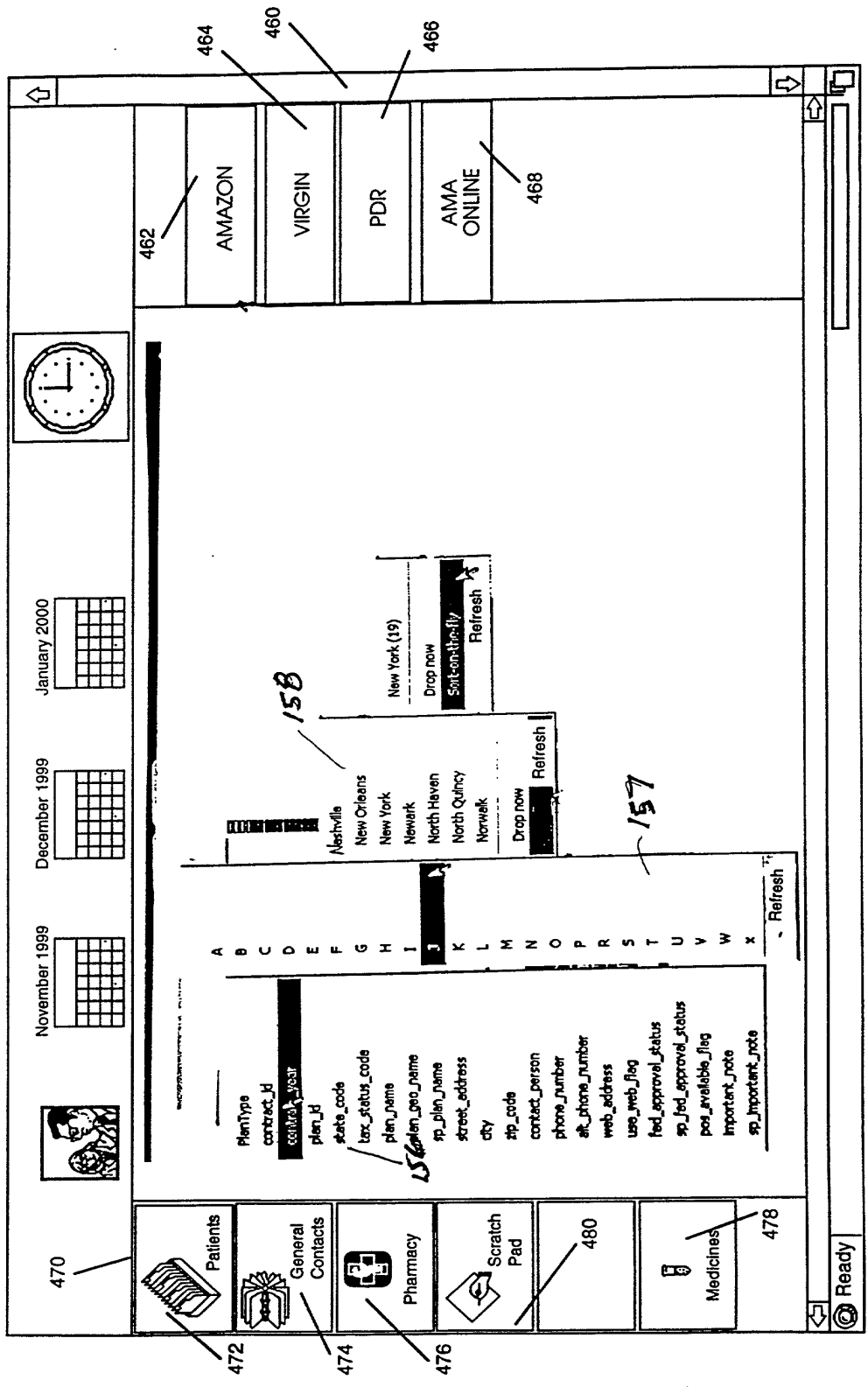
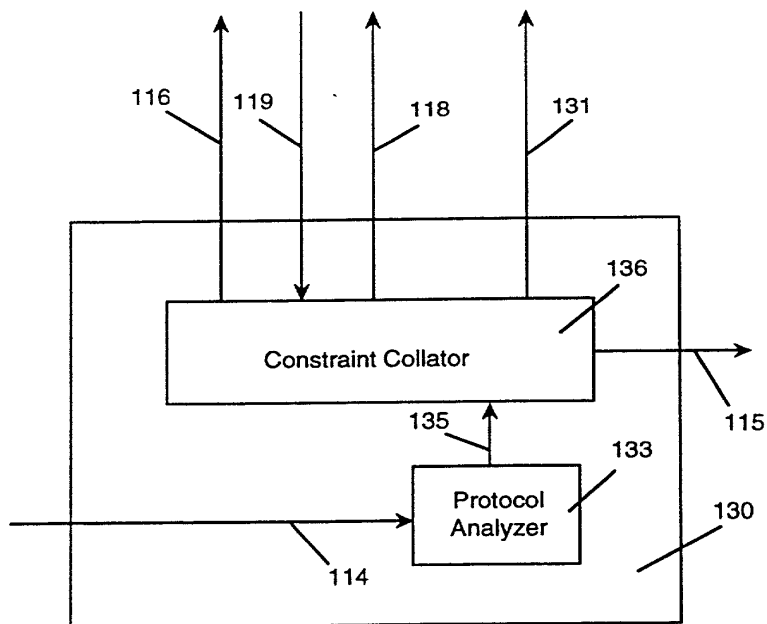


Fig. 3

450

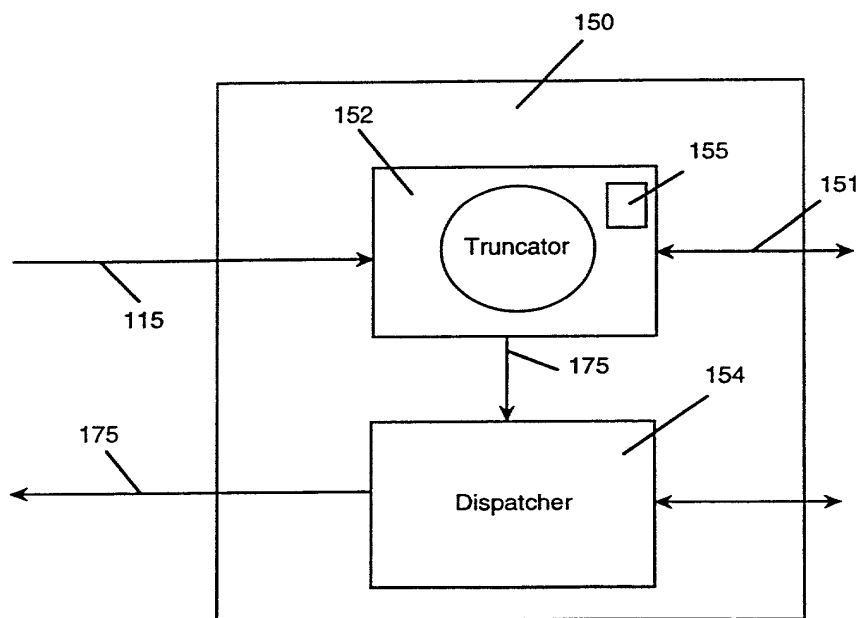
Fig. 4



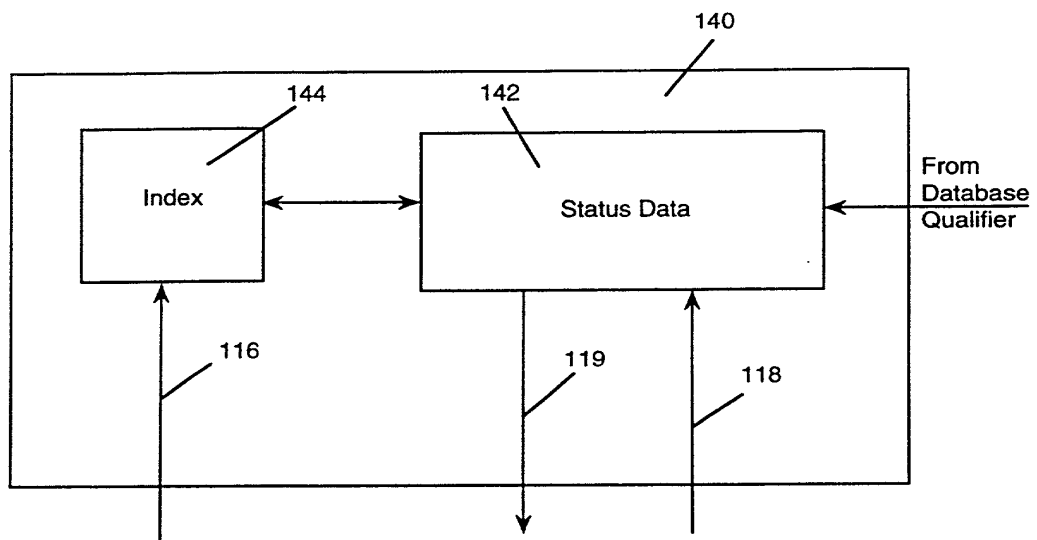


*Fig. 5*

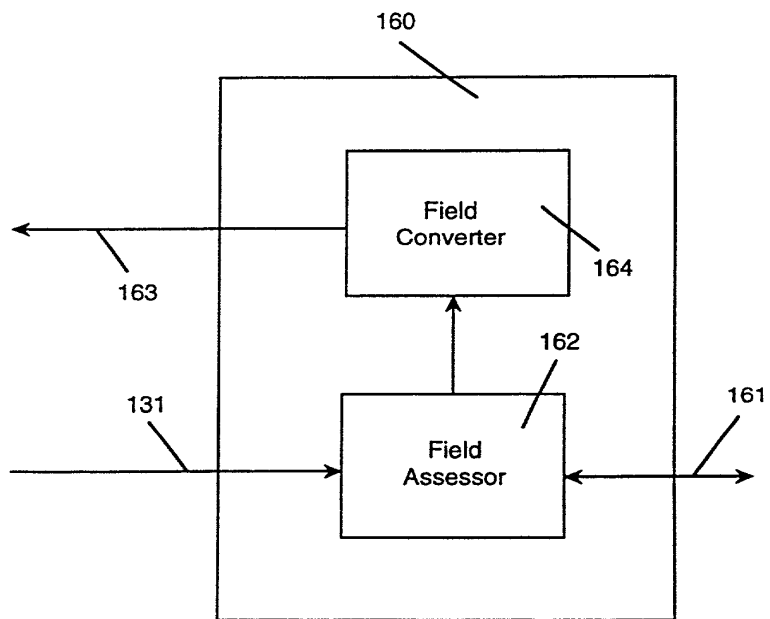




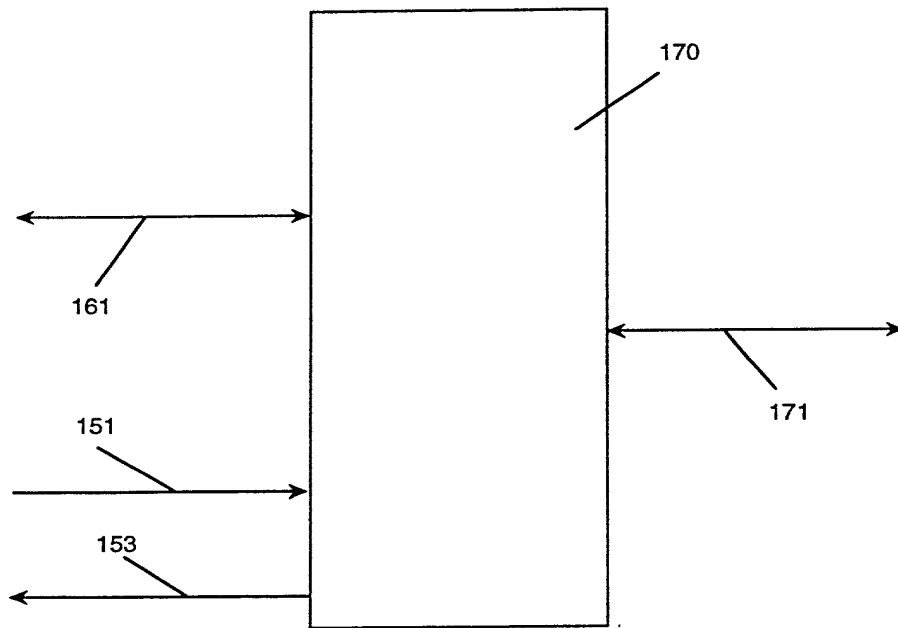
*Fig. 6*



*Fig. 7*



*Fig. 8*



*Fig. 9*

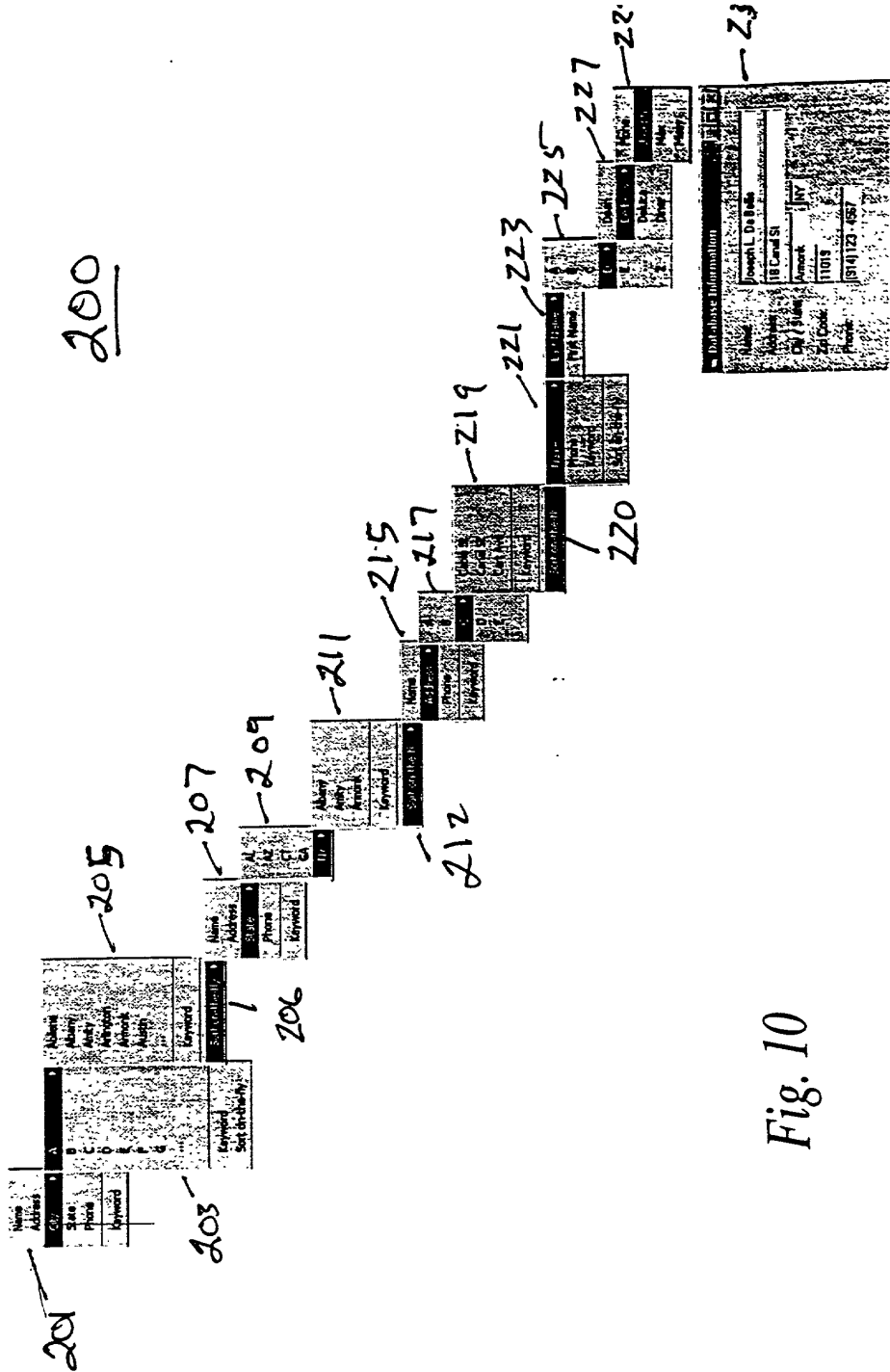


Fig. 10

Fig. 11

The screenshot shows the Barnesandnoble.com website interface. At the top, there is a navigation bar with 'bn.com BARNES & NOBLE' and a search bar. Below the navigation bar, there are several promotional banners and a list of search results. The search results are organized into a grid with columns for 'Title', 'Author', 'ISBN', 'Subtitle', 'Format', 'Date Published', 'Stock Status', 'Recommended Age', 'Pages', 'Ratings', 'Price', 'Retail', 'Savings', 'Publisher', 'Keyword...', and 'Print'. The search results are filtered by 'Cells' and include titles like 'Cells: A History of the Cell', 'Cells: Embryos and Evolution', and 'Cells: The Cell'. There are also links to 'Advanced Search' and 'Order Status'. The page number '237' is visible in the bottom right corner.

23

237

234 235 236



Fig. 13

233

www.barnesandnoble.com (www.bn.com) - Microsoft Internet Explorer

File Edit View Favorites Tools Help

Search History Favorites

bn.com BARNES & NOBLE

February 11, 2000  
 Welcome to the new sound of shopping!  
**bn.com**

Special Features  
 Oprah's Pick  
 bn.com Live  
 Award Winners  
 Bestsellers  
 Recommended  
 New Releases  
 bn.com Editors  
 bn.com Interests

Browse Subjects & Art, Architecture & Photography  
 Biography  
 Business  
 Computers  
 Cooking  
 Fiction & Literature  
 History  
 Home & Garden  
 Kids

Books  
 Browse Subjects  
 Categories

OTFT

1 2 3 4 5 6 7 8 9 A B C C+ C- C. Ca Cb Cc Cd Ce Cf Cg Ch Ci Cl Cm Cn Co Cp Cq Cr Cs Ct Cu Cz

Screen  
 y, pouncy" Tigger is starring in his own hit feature film  
 and the rest of the gang from the Hundred-Acre Wood.  
 of Tigger movie tie-in books.

Safe Shopping Guarantee  
 Order Status  
 Coming Soon  
 Place your orders now.  
 Deepak Chopra, How to Know God  
 Nora Roberts, Carolina Moon  
 Andrew Weil, Eating Well for Optimum Health  
 E. L. Doctorow, City of God  
 Helen Fielding, Bridget Jones: The Diary of a Young Woman

Barabara  
 Cite  
 Search

CC  
 CCD Arrays, Cameras and Displays, Vol. 57  
 CCD Astronomy  
 CCH Business Owner's Toolkit: Tax Guide 2000  
 CCIE Professional Development  
 CCIE Resource Kit  
 CCNA  
 CCNA Exam Notes  
 CCNA Routing and Switching Exam Crm  
 CCNP  
 CCNP Advanced Cisco Router Configuration Exam Crm  
 CCNP Cisco Intermediate Troubleshooting Exam Crm  
 CCNP Cisco LAN Switch Configuration Exam Crm  
 CCNP Exam Notes Advanced Cisco Ro  
 CCRN Certification Examination Review Course (Book+Set of 4 Audio Tapes)  
 CCTV  
 Cou-Card

233  
 IN OUT OF PRINT  
 Rare  
 Out of Print  
 Reconstruct  
 Black Histor  
 of Reconstri  
 original docu  
 look at  
 The Harlem Renaissance

234 239 240

My Computer



Fig. 14

The screenshot shows a search results page on Barnesandnoble.com. The search term is "Screenplay". The results are filtered to show 241 items. Handwritten annotations include "233" pointing to the search bar, "234" pointing to the "241" result count, and "242" pointing to the "Sort on-the-fly" dropdown menu. The page features a navigation bar with "bn.com" and "BARNES & NOBLE" logos, a search bar, and various filters like "Books", "Browse Subjects", and "Categories". The search results list includes titles such as "Screenplay: A Step-by-Step Guide to Writing, Producing, and Selling Your Screenplay" and "Screenplay: The Art of Writing for the Screen".





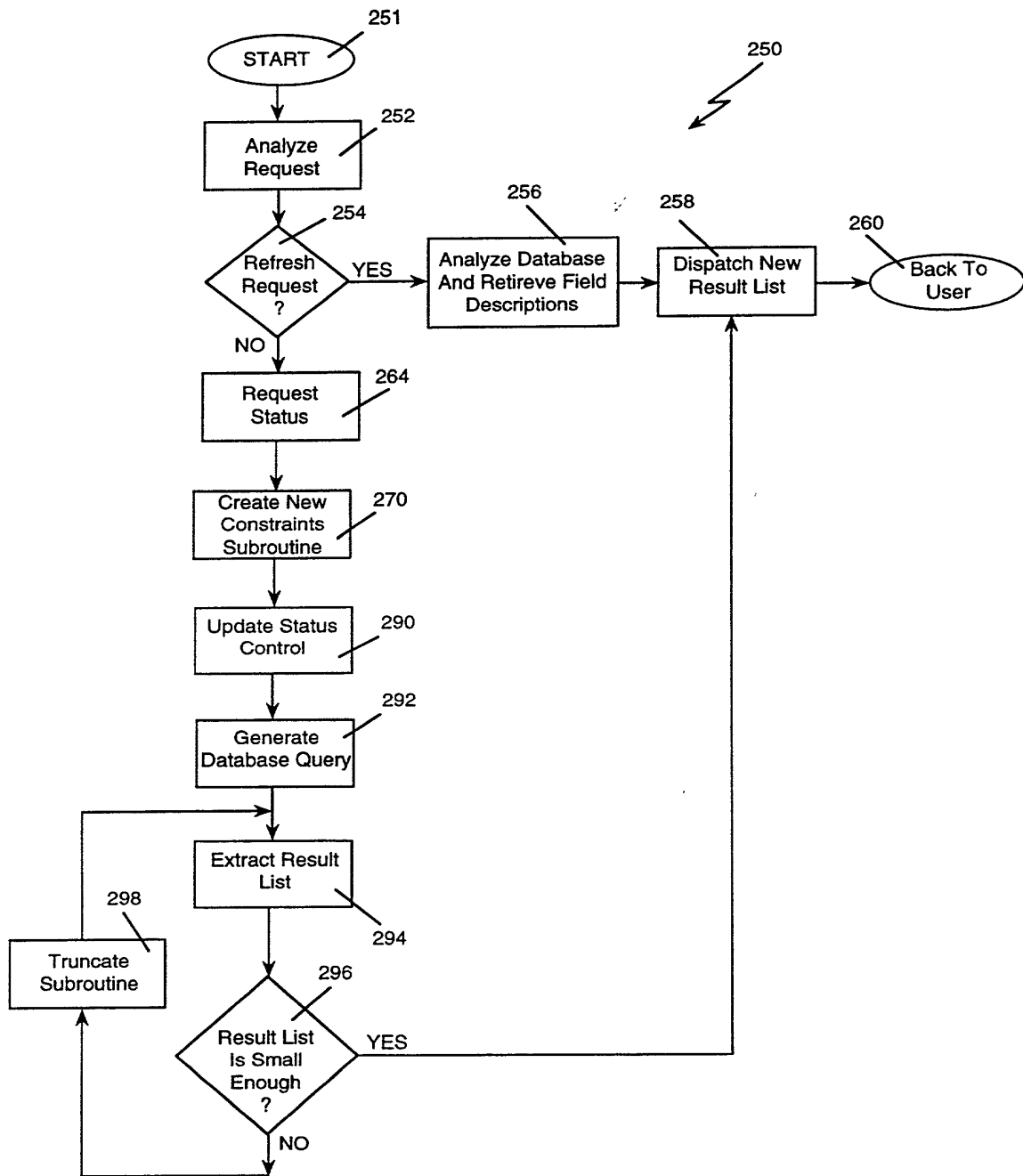


Fig. 16



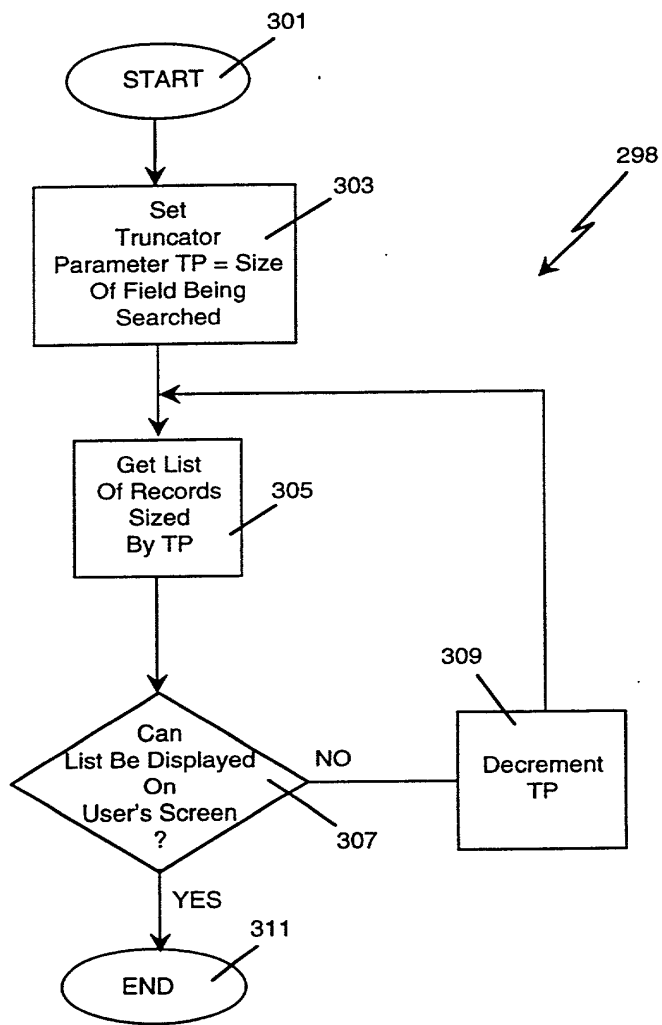


Fig. 18



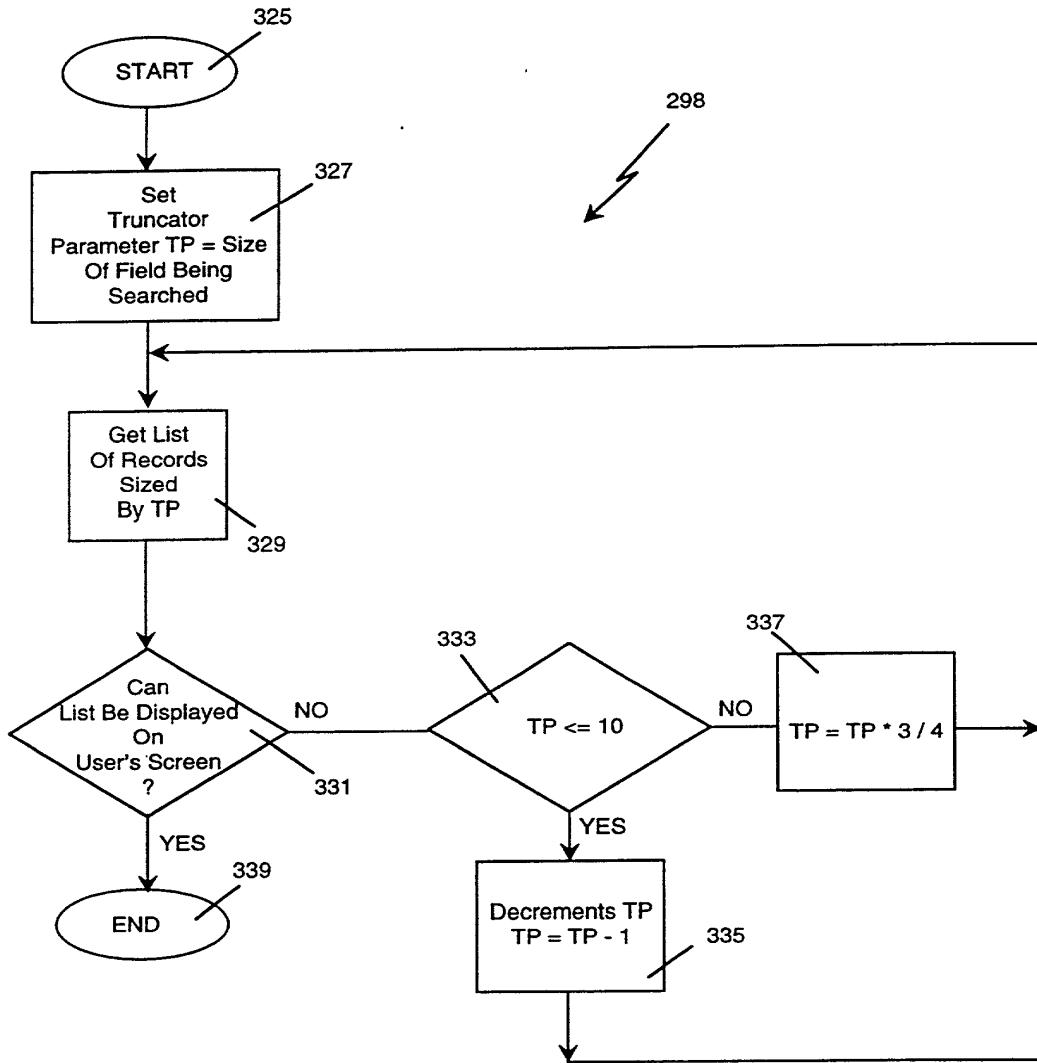


Fig. 20

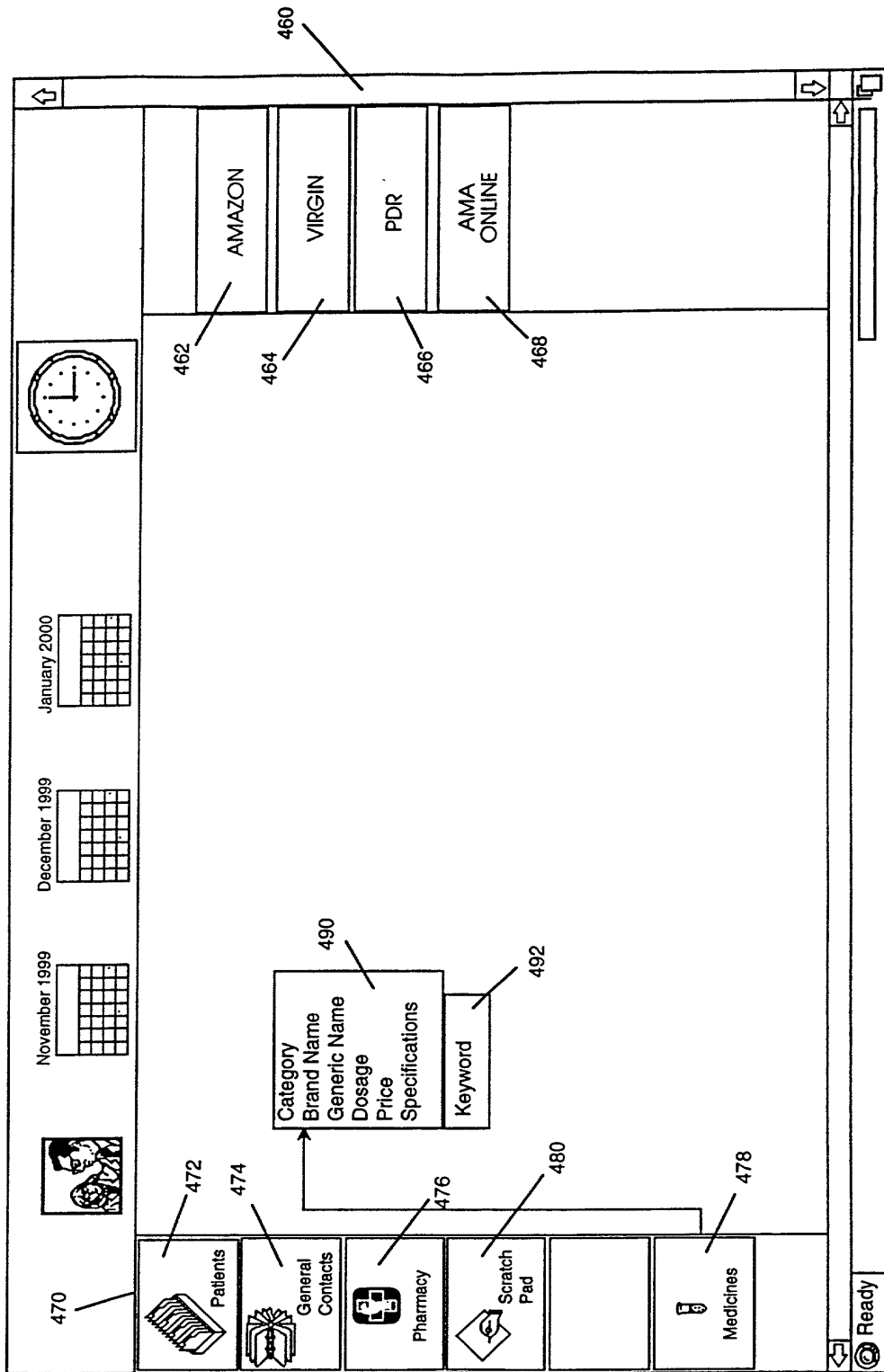




FIG. 22

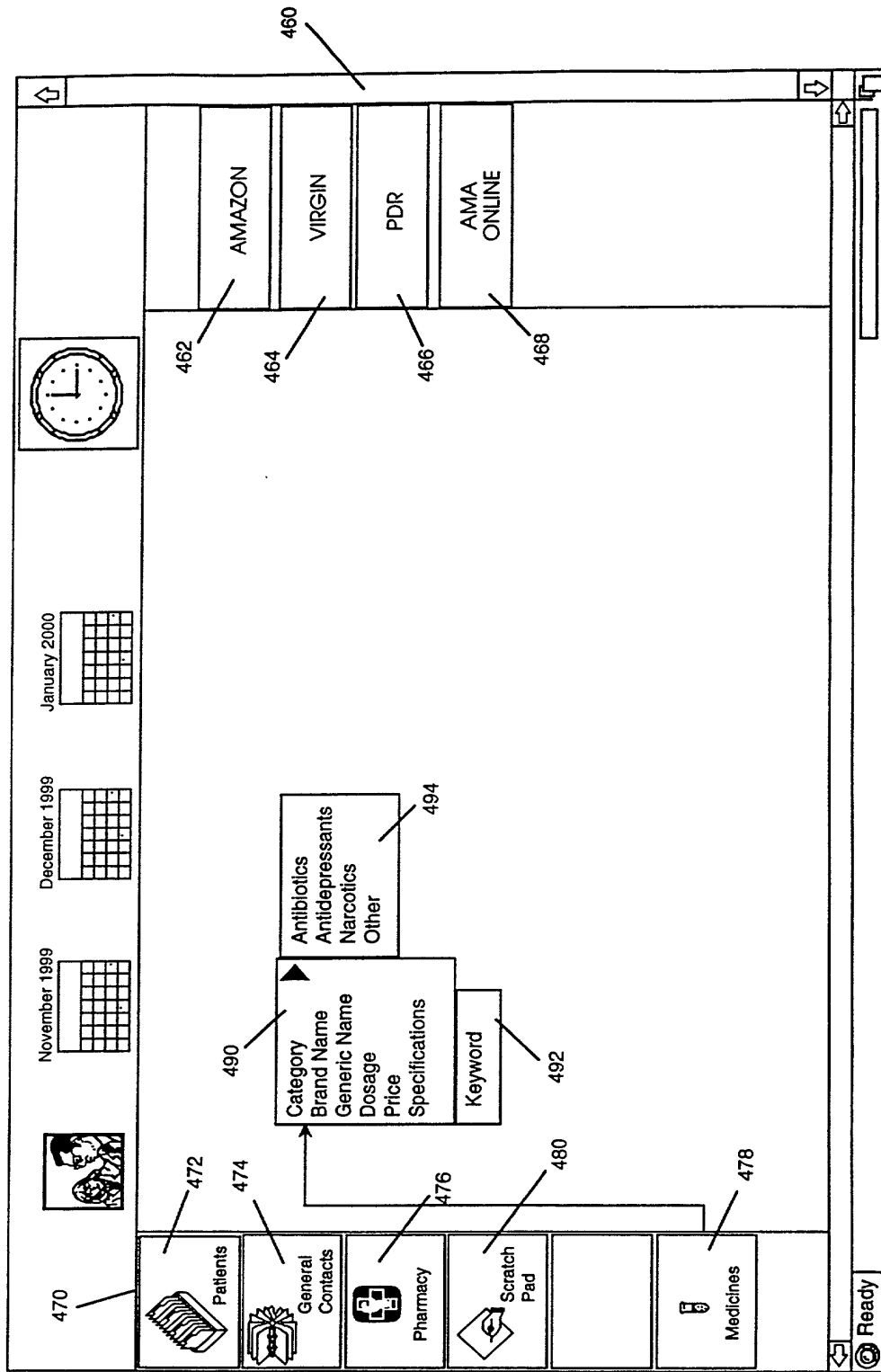
450

Fig. 22



450

Fig. 23

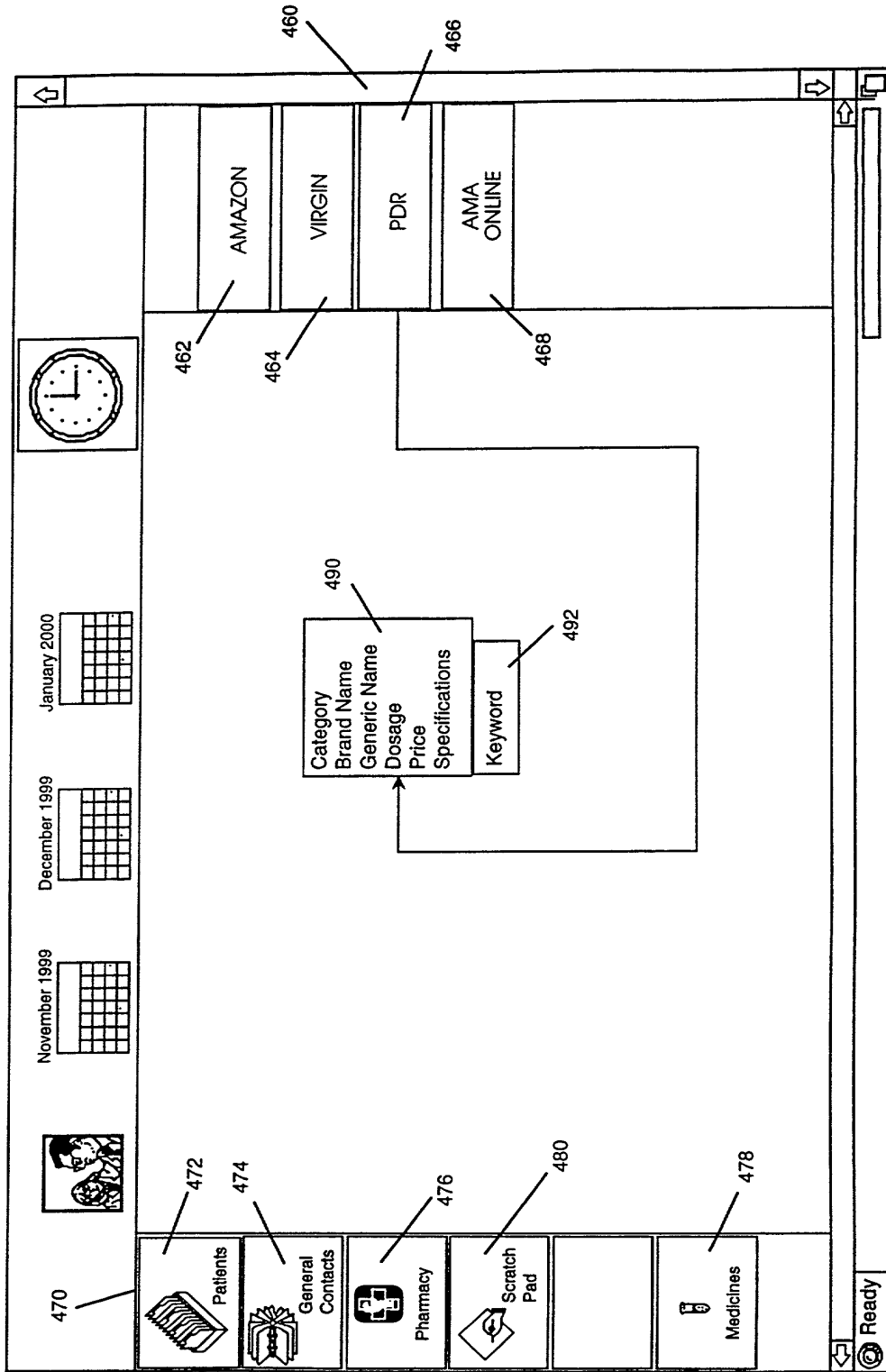




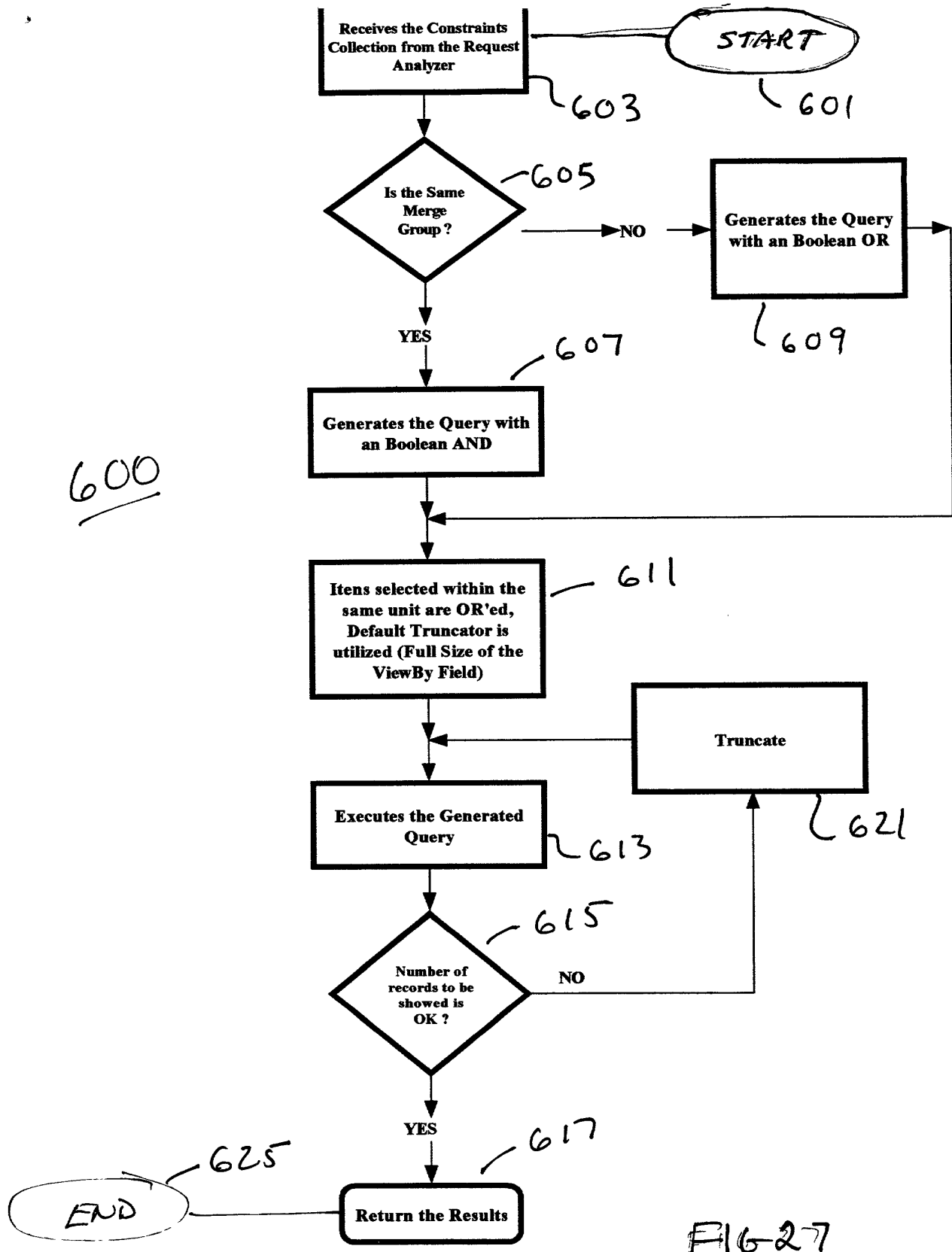
U.S. PATENT AND TRADEMARK OFFICE

450

Fig. 25







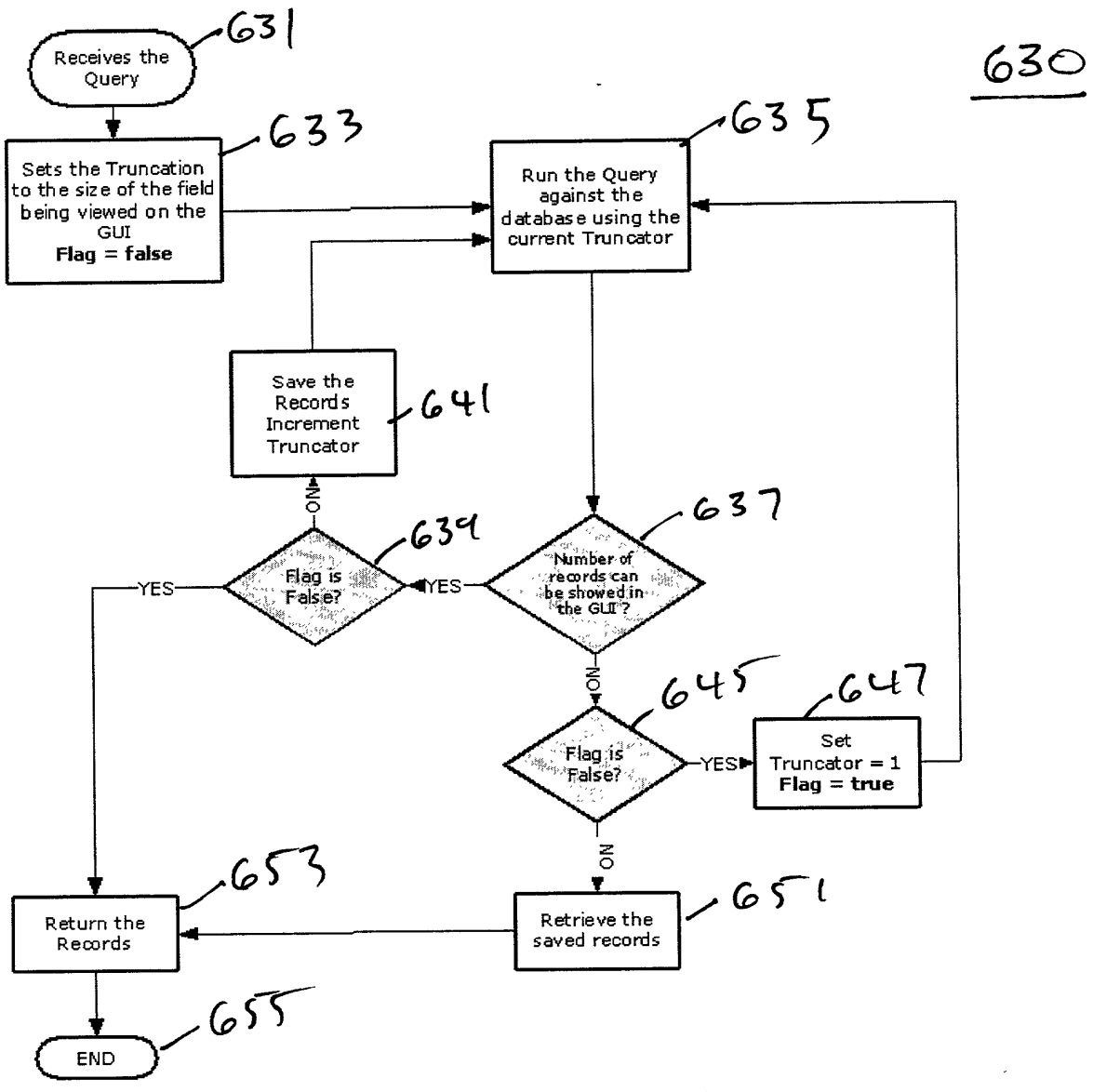


FIG. 28







MDI Form 1

S.O.F. DB Oper. 50F Setup

Assignee

- 3 COM Corporation (Santa Clara, CA)
- 3 I Research Exploitation Ltd. (GB)
- 3 S Becton B.V. (Greve, NL)
- 3 Sigma, Inc. (Covington, OH)
- 3 U Partners (Casper, WY)
- 3 W Industry Inc. (TVA)
- 3-D Diet, Inc. (Walnut Creek, CA)
- 3-Dimensional Pharmaceuticals, Inc. (Exton, PA)
- 3-M Company (St. Paul, MN)
- 314613 B.C. Ltd. (Vancouver, CA)
- 3244 Corporation (Oak Lawn, IL)
- 373249 B.C. Ltd. (Vancouver, CA)
- 378134 Alberta Ltd. (Calgary, CA)
- 379235 Ontario Ltd. (Punam, CA)
- 382534 Alberta Ltd. (Edmonton, CA)
- 3Com Corp. (Rolling Meadows, IL)
- 3Com Corp. (Santa Clara, CA)
- 3Com Corporation (Mountain View, CA)
- 3COM Corporation (Rolling Meadows, IL)
- 3Com Corporation (Santa Clara, CA)
- 3Com Technologies (Cayman Islands, VG)
- 3D International AIS (Oslo, NO)
- 3D Systems, Inc. (Valencia, CA)
- 3DPX Interactive, Incorporated (San Jose, CA)
- 3DGeo Development, Inc. (Mountain View, CA)
- 3Dlabs Ltd. (Eichen, GB)
- 3Dlabs, Ltd (Hamilton, MN, RW)
- 3DY Systems Ltd. (Noblesham, IL)
- 3E Corporation (West Lafayette, IN)
- 3E, Inc. (Mountain View, CA)
- 3M Company (St. Paul, MN)
- 3M Innovative Company (St. Paul, MN)
- 3M Innovative Properties (St. Paul, MN)
- 3M Innovative Properties Co (St. Paul, MN)
- 3M Innovative Properties Co. (St. Paul, MN)
- 3M Innovative Properties Company (Saint Paul, MN)
- 3M Innovative Properties Company (St. Paul, MN)
- 3M Innovative Properties Company (St. Paul, MN)
- 3M Innovative Properties Company (St. Paul, MN)
- 3M Management AB (Vellingby, SE)
- 3M S.p.A. (Turin, IT)
- 3U Partners (Casper, WY)
- 3V Enterprises Inc. (CA)

View By...

- Assignee
- 3-M Company (St. Paul, MN)
- 3M Company (St. Paul, MN)
- View By...
- Merge

View By...

- View By...
- Merge

LS References

Foreign References

Start

SQL Server Enterp... plot - Microsoft Vi... SQL Server Query ... Document1 - Micro... MDIForm1

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FIG. 31









Patent Title	Patents - Assignee
Anisotropic retardation layers for display devices	3M Innovative Properties Company (St. Paul, MN)
Dispenser roll cover and methods of preparation and use thereof	3M Company (St. Paul, MN)
Dumping unit for globular storage tank	3M Innovative Properties Co. (St. Paul, MN)
Dark acrylic pressure-sensitive adhesive	3M Innovative Properties Co. (St. Paul, MN)
Data accumulation system	3M Innovative Properties Company (St. Paul, MN)
Data cartridge with secondary tape guides	3M Innovative Properties Company (St. Paul, MN)
Data processing form	3M Innovative Properties Company (St. Paul, MN)
Data storage structure to enable subsequent computerized preiteration	3M Innovative Properties Company (St. Paul, MN)
DC Power supply for high power discharge devices	3M Innovative Properties Company (St. Paul, MN)
Decolorable imaging system	3M Innovative Properties Company (St. Paul, MN)
Decorative ribbon or sheet material	3M Innovative Properties Company (St. Paul, MN)
Demand and issued renewing imaging media	3M Innovative Properties Company (St. Paul, MN)
Dental filling composition utilizing zinc-containing inorganic filler	3M Innovative Properties Company (St. Paul, MN)
Dentin and enamel adhesive	3M Innovative Properties Company (St. Paul, MN)
Densitizer for ferromagnetic markers used with electromagnetic article surveillance systems	3M Innovative Properties Company (St. Paul, MN)
Densitizing dyes for photographic emulsions	3M Innovative Properties Company (St. Paul, MN)
Detachable adhesive disk	3M Innovative Properties Company (St. Paul, MN)
Detecting system	3M Innovative Properties Company (St. Paul, MN)
Detection of articles	3M Innovative Properties Company (St. Paul, MN)
Developer compositions for silver halide photographic materials comprising cyclic amino methylene diphosphonates	3M Innovative Properties Company (St. Paul, MN)
Developer compositions having layer of a pigment on the surface thereof	3M Innovative Properties Company (St. Paul, MN)
Developer material level sensor	3M Innovative Properties Company (St. Paul, MN)
Developer powder supply cartridge	3M Innovative Properties Company (St. Paul, MN)
Developing powder composition containing a fluorine-modified alkyl siloxane	3M Innovative Properties Company (St. Paul, MN)
Developing powder composition containing fluorosiloxane sulfonamide surface active agent	3M Innovative Properties Company (St. Paul, MN)
Device and method for applying flexible back to containers	3M Innovative Properties Company (St. Paul, MN)
Device for blocking but-weak between tubes	3M Innovative Properties Company (St. Paul, MN)
Device for cutting a support web for a radially expanded resilient sleeve	3M Innovative Properties Company (St. Paul, MN)
Device for exposing colorant to be transferred	3M Innovative Properties Company (St. Paul, MN)
Device for forming graphics	3M Innovative Properties Company (St. Paul, MN)
Device for fusing lengths of film over the open ends of cups	3M Innovative Properties Company (St. Paul, MN)
Device for restraining an object or objects therein	3M Innovative Properties Company (St. Paul, MN)
Device to slow solenoid actuation motion	3M Innovative Properties Company (St. Paul, MN)
Diagnostic radio-labeled polysaccharide derivatives	3M Innovative Properties Company (St. Paul, MN)
Diaper closure utilizing pressure-sensitive adhesive tape having textured foil backing	3M Innovative Properties Company (St. Paul, MN)
Dioxonium imaging system	3M Innovative Properties Company (St. Paul, MN)
Dielectric stress relief at a high voltage cable termination	3M Innovative Properties Company (St. Paul, MN)
Diffraction lens	3M Innovative Properties Company (St. Paul, MN)
Digital communications system with automatic frame synchronization and detector circuitry	3M Innovative Properties Company (St. Paul, MN)
Digital frame synchronizing circuit	3M Innovative Properties Company (St. Paul, MN)
Digital motor control system	3M Innovative Properties Company (St. Paul, MN)
Dimensionally-controlled cobalt-containing precision molded metal article	3M Innovative Properties Company (St. Paul, MN)
Direct positive silver halide emulsions containing quaternated merocyanine dyes	3M Innovative Properties Company (St. Paul, MN)
Directional reflection detector	3M Innovative Properties Company (St. Paul, MN)
Disk cartridge	3M Innovative Properties Company (St. Paul, MN)
Disk dispenser	3M Innovative Properties Company (St. Paul, MN)
Disinfectable dental sealant	3M Innovative Properties Company (St. Paul, MN)
Disinfecting method and compositions	3M Innovative Properties Company (St. Paul, MN)
Disk cartridge	3M Innovative Properties Company (St. Paul, MN)
Disk locking mechanism for disk cartridge	3M Innovative Properties Company (St. Paul, MN)
Dispersed imaging systems with tetra (hydrocarbyl) borate salts	3M Innovative Properties Company (St. Paul, MN)
Disk restraint	3M Innovative Properties Company (St. Paul, MN)
Disinfectable polypropylene adhesive-coated tape	3M Innovative Properties Company (St. Paul, MN)
Dispenser for a stack of note paper	3M Innovative Properties Company (St. Paul, MN)
Dispenser for adhesive coated sheet material	3M Innovative Properties Company (St. Paul, MN)
Dispenser for protected write-on labels	3M Innovative Properties Company (St. Paul, MN)
Dispenser package	3M Innovative Properties Company (St. Paul, MN)
Dispersed imaging systems with tetra (hydrocarbyl) borate salts	3M Innovative Properties Company (St. Paul, MN)
Drop wire connector	3M Innovative Properties Company (St. Paul, MN)
Dry magnetic pressure-fusible developing powder	3M Innovative Properties Company (St. Paul, MN)
Dry strip anthralin layer for photochromographic film	3M Innovative Properties Company (St. Paul, MN)
Dry transfer article	3M Innovative Properties Company (St. Paul, MN)
Dry transfer graphics article and methods of preparation and use thereof	3M Innovative Properties Company (St. Paul, MN)
Dry transfer graphics article method of preparation	3M Innovative Properties Company (St. Paul, MN)
Dual groove Fresnel lens for overhead projection	3M Innovative Properties Company (St. Paul, MN)
Dual particle population magnetic recording medium	3M Innovative Properties Company (St. Paul, MN)
Dual status magnetic marker having magnetically bleasable flux collectors for use	3M Innovative Properties Company (St. Paul, MN)
Durable glass elements	3M Innovative Properties Company (St. Paul, MN)
Durable melt-blown particle-loaded sheet material	3M Innovative Properties Company (St. Paul, MN)
Durable, polishable direct filling material	3M Innovative Properties Company (St. Paul, MN)
Durably stain-resistant and self-resistant ple fabric and process	3M Innovative Properties Company (St. Paul, MN)
Dust mop	3M Innovative Properties Company (St. Paul, MN)
Dust mop frame	3M Innovative Properties Company (St. Paul, MN)
Dyed aqueous air foams	3M Innovative Properties Company (St. Paul, MN)
Dyes suitable for sensitization of photoconductive systems	3M Innovative Properties Company (St. Paul, MN)
Electrical connector tape	3M Innovative Properties Company (St. Paul, MN)
Electrically conductive metal oxide coatings	3M Innovative Properties Company (St. Paul, MN)
Method for writing arbitrary index perturbations in a wave-guiding structure	3M Innovative Properties Company (St. Paul, MN)
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FIG-36



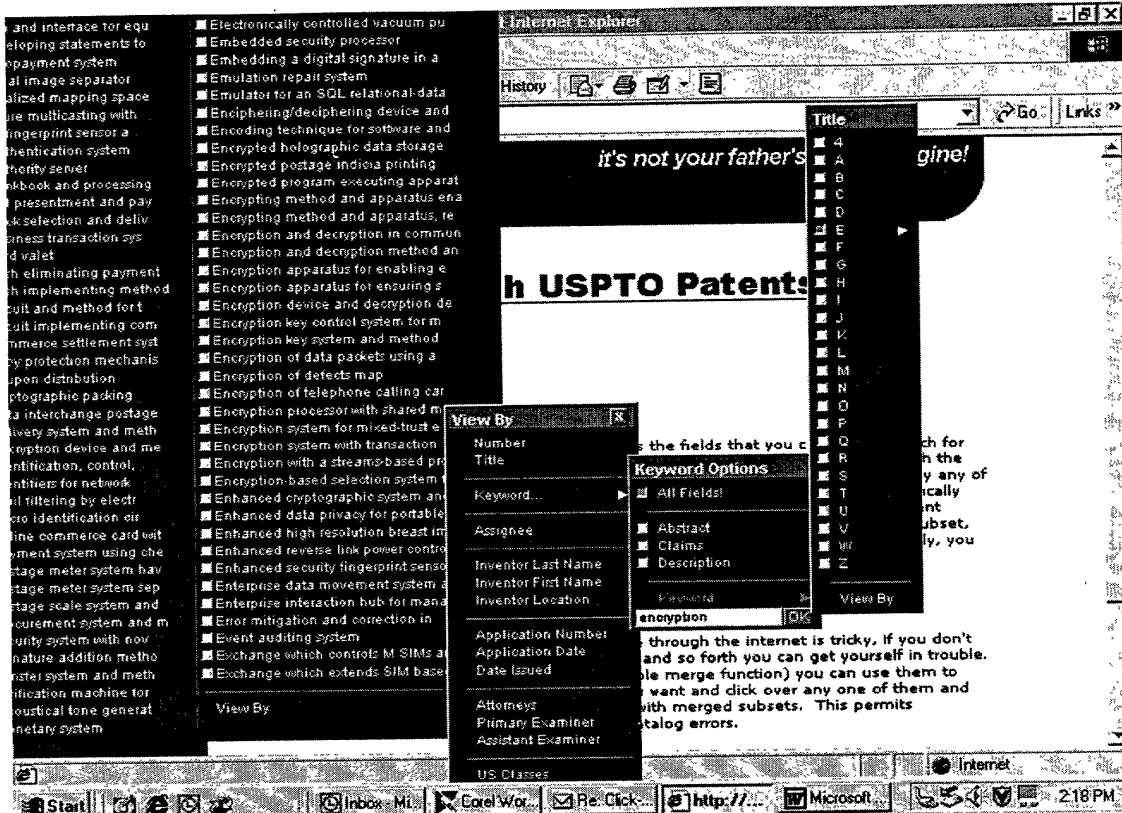


FIG 37

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Member Access

Home

Philosophy

Features

Technology Demos

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FAQ's

**Virtual Logistix Inc.**

- Efficient micropayment system
- Efficient spatial image separator
- Efficient virtualized mapping space
- Efficient, secure multicasting with
- Electric field fingerprint sensor a
- Electronic authentication system
- Electronic authority server
- Electronic bankbook and processing
- Electronic bill presentment and pay
- Electronic book selection and deliv
- Electronic business transaction sys
- Electronic card valet
- Electronic cash eliminating payment
- Electronic cash implementing method
- Electronic circuit implementing com
- Electronic commerce settlement syst
- Electronic copy protection mechanis
- Electronic coupon distribution
- Electronic cryptographic padding
- Electronic data interchange postage
- Electronic delivery system and meth
- Electronic encryption device and me
- Electronic identification, control,
- Electronic identifiers for network
- Electronic mail filtering by electr
- Electronic micro identification cir
- Electronic online commerce card out
- Electronic payment system using che
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- Electronic postage meter system sep
- Electronic postage scale system and
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- Electronic security system with nov
- Electronic signature addition metho
- Electronic transfer system and meth
- Electronic verification machine for
- Electronic, acoustical tone generat
- Electronic monetary system
- Embedding a digital signature in a
- Emulation repair system
- Emulator for an SQL relational data
- Enciphering/deciphering device and
- Encoding technique for software and
- Encrypted holographic data storage
- Encrypted postage indicia printing
- Encrypted program executing apparat
- Encrypting method and apparatus ena
- Encrypting method and apparatus, re
- Encryption and decryption in commun
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- Encryption apparatus for ensuring s
- Encryption device and decryption de
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- Encryption key system and method
- Encryption of data packets using a
- Encryption of defects map
- Encryption of telephone calling car
- Encryption processor with shared me
- Encryption system for mixed-trust e
- Encryption system with transaction
- Encryption with a stream-based pro
- Encryption-based selection system f
- Enhanced cryptographic system and m
- Enhanced data privacy for portable
- Enhanced high resolution breast ima
- Enhanced reverse link power control
- Enhanced security fingerprint senso
- Enterprise data movement system and
- Enterprise interaction hub for mana
- Error mitigation and correction in
- Event auditing system
- Exchange which controls M SIMs and
- Exchange which extends SIM based au

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Attorneys  
Primary Examiner  
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US Classes

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FIG. 38

# Patent: 5935246

## Electronic copy protection mechanism using challenge and response to prevent unauthorized execution of software

**Date Filed:**

4/11/1997

**Date Issued:**

8/10/1999

**Application Number:**

838620

**Go to USPTO.GOV**

[US PTO](#)

**Abstract:**

A copy protection mechanism for protecting software against copying, consists of a challenge mechanism embedded in each protected item of software. The challenge mechanism has no access to the customer's private keying material. In operation, the challenge mechanism sends a random challenge to the customer's signature server. The signature server signs the challenge, using the customer's private keying material and then returns the signed challenge to the challenge mechanism. The challenge mechanism then verifies the signed challenge, using the customer's public keying material, and prohibits the customer from using some or all of the protected item of software unless the verification is successful. The mechanism permits every customer to receive an identical copy of the copy protected program with the embedded challenge mechanism.

**Inventors:**

Benson, Glenn Stuart

**Inventor Location:**

Munich, DE

**Assignee:**

International Computers Limited (London, GB)

**US Classes:**

713/200  
713/201

**International Classes:**

**US References:**

4558176  
4926480  
4947430  
5109413  
5146575  
5224163  
5315657  
5371794  
5436972  
5568552  
5724425

**Foreign References:**

**Primary Examiner:**

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**Assistant Examiner:**

Mai, Rijue

**Attorney:**

Lee, Mann, Smith, McWilliams, Sweeney & Ohlson

**Claims:**

FIG. 39



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			Dority, John P., Cleaver, William	1622 Lake Johanna Blvd., Arden Hills, MN 55112
			Dorman, Ira S.	1925 Noble Dr., Minneapolis, MN 55422
			Dorman, William S.	2000 Argonne Dr., Minneapolis, MN 55421
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			Dorr, Carson, Sloan and Peters	3109 Clinton Ave. South, Minneapolis, MN 55408
			Dorr, Carson, Sloan, & Peterso	48 Woodland Gardens, London, N10 3UA, GB
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				Burnsberg, AU
				Burnsville, MN
				Chicago, IL
				Circle Pines, MN
				Darwin, MN
				Eagan, MN
				Edina, MN
				Embleton, AU
				Excelsior, MN
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				Kardinya, AU
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				Madison, WI
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				St. Paul, MN
				Stillwater, MN
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				Waseca, MN
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				Zimmerman, MN

FIG 41











S.O.F.

View By	Primary Examiner
	Clin
	Coan
	Cobi
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	Cohn
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	Comb
	Cong
	Cook
	Coop
	Corb
	Corc
	Core
	Cors
	Cosh
	Coug
	Cous
	Cove
	Cox,
	Crai
	Cram
	Cran
	Cras
	Cros
	Crow
	Croy
	Cuch
	Cuci
	Cuda
	Culb
	Culy
	Quom
	Curt
	Cusb
	Cusi
	Cust
	Czal
	Czar
	Czas
	Claw
	Clay

Primary Examiner  
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 Casaregola, Louis L.  
 Casaregola, Louis J.  
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FIG. 45

Primary Examiner

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- Dorner, Kenn
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- Doutreau, Le
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- Douglas, Win
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- Downey, Kenn
- Downey, Mary
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- Draper, Gern
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- Drezin, Norri
- Drodge, Jose
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- Drummond, Do
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- Dugan, James
- Duggs, Donoy
- Duggan, Dono
- Duggan, Dono
- Duzan, James
- Dwyer, James
- Dzilehzynski,
- Dzierzynski,

View By

- Title
- Patent Number
- Inventor Name
- Inventor's Location
- US References
- Foreign References
- US Classes
- International Classes
- Application Number
- Application Date
- Issue Date
- Primary Examiner
- Assistant Examiner
- Attorney
- Assignee

Primary Examiner

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- Dees, Josef
- Dees, Josee G.

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FIG. 46

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Ultricks; Charles D. (Newbury Park, CA)
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Ulinic France (Paris, FR)
Ulrich Baensch (Melle, DT)
Ulrich Luboschik (DE)
Ullstein Propeller AS (NO)
Ultimate Inc. (Quebec, CA)
Ultimate Window Coverings, Inc. (Lakewood, CO)
Ultra Light Arms, Inc. (Granville, WV)
Ultra Mortar, Inc. (Ogden, UT)
Ultra Plating Corporation (Green Bay, WI)
Ultra-Centrifuge Nederland N.V. (Almelo, NL)
Ultra-Mold Corporation (Willow Grove, PA)
Ultra-Precision Manufacturing, Ltd. (Birmingham, MI)
Ultra-Violet Products, Inc. (San Gabriel, CA)
Ultracentrifuge Nederland N.V. (The Hague, NL)
Ultracent Products, Inc. (Salt Lake City, UT)
Ultracent Products, Inc. (South Jordan, UT)
Ultrafibre, Inc. (Granville, OH)
Ultrakust Electronic GmbH (Ruhmannsfeiden, DE)
Ultralife Batteries, Inc. (Newark, NY)
Ultrasonic Embroidery Machine Company (North Haven, CT)
Ultrasonic Equipment Company (Addison, IL)
Ultramatrix, Inc. (Los Angeles, CA)
Ultramed Corporation ()
ULTRAMET (Pasadena, CA)
Ultrasonic Arrays, Inc. (Woodinville, WA)
Ultrasonic Energy Corporation (Riverdale, NY)
Ultrasonic Systems, Inc. (Farmingdale, NY)
Ultratec, Inc. (Madison, WI)
Ultratek International, Inc. (Concord, CA)
Ultrathermics (Redmond, WA)
Ultraviolet Purification Systems (Bedford Hills, NY)
Ultraviolet Purification Systems, Inc. (Bedford Hills, NY)
Ultraystems Defense and Space, Inc. (Irvine, CA)
Ultronic Systems Corporation (Mooresstown, NJ)
Ultrrox International (Santa Ana, CA)
Uly-Pak, Inc. (Carbondale, IL)
Uly-Pak, Inc. (Ulysses, KS)
Ulysses Corporation (St. George, UT)

Assignee	View By...
U	View By...
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Ua	View By...
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Uc	View By...
Ud	View By...
Ue	View By...
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Ug	View By...
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Ui	View By...
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Title	Title
Patent Number	Patent Number
Inventor Name	Inventor Name
Inventor's Location	Inventor's Location
US Referer	US Referer
Foreign Re	Foreign Re
US Classe	US Classe
International	International
Application	Application
Applicator	Applicator
Issue Date	Issue Date
Primary Exar	Primary Exar
Assistant Ex	Assistant Ex
Attorney	Attorney
Assignee	Assignee

Assignee	View By...
Ultrasonic Systems, Inc. (Farmingdale, NY)	

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FIG. 47



View By	Title	US Classes	View By...
Patent N	Artificial logs and log-making method and apparatus	440	View By...
	Fire log process and apparatus	441	View By...
	Fuel compacting apparatus	442	View By...
Inventor	Method and apparatus for recovering by-product still fines from a slurry thereof	443	View By...
Inventor	Method of charging solids into coal gasification reactor	444	View By...
	Process for making low-sulfur and low-ash fuels	445	View By...
US Refs		446	View By...
Foreign		447	View By...
US Classes		448	View By...
International Classes		449	View By...
Application Number		450	View By...
Application Date		451	View By...
Issue Date		452	View By...
Primary Examiner		453	View By...
Assistant Examiner		454	View By...
Attorney		455	View By...
Assignee		456	View By...

US Classes	View By...
441DA	View By...
441DB	View By...
441DC	View By...
441DD	View By...
441DE	View By...
441DF	View By...
441DH	View By...
441DJ	View By...
441DK	View By...
441DR	View By...
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4415R	View By...
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44160	View By...
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449/	View By...

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G	View By...

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FIG. 48

(4008054) Process for making low-sulfur and low-ash fuels

Title: Process for making low-sulfur and low-ash fuels  
 Assignee: Consolidation Coal Company (Pittsburgh, PA)  
 Attorney: Mikessell, Jr., William A., Fowler, Jr., D. Leigh, Price, Jr., Stanle  
 Examiner: Dees, Carl F.  
 Assistant Examiner:

Patent Number: 4008054  
 Application Number: 540310  
 Application Date: 1/10/1975  
 Issue Date: 2/15/1977

1. Field of the invention: This invention relates to a process for converting coal to low-sulfur and low-ash gaseous, liquid and solid fuels, and more particularly, to a process for supplying the energy requirements of a steel plant from an ash- and sulfur-containing coal. Description of the Prior Art: The primary source of energy for the steel industry continues to be coke for the blast furnace. The conventional method for coke manufacture, that is, by slot ovens, requires a blend of high and low volatile coals of proper swelling properties to produce a strong coke without damaging the ovens. Beyond these physical properties, there is a need for desirable chemical properties (i.e., low ash and sulfur content) to permit low-cost production of high quality hot metal. With the continued expansion of the world's productive capacity for steel, a growing shortage of good metallurgical coals is developing, particularly those having the essential low volatile coal ingredients. Low-sulfur coals also are

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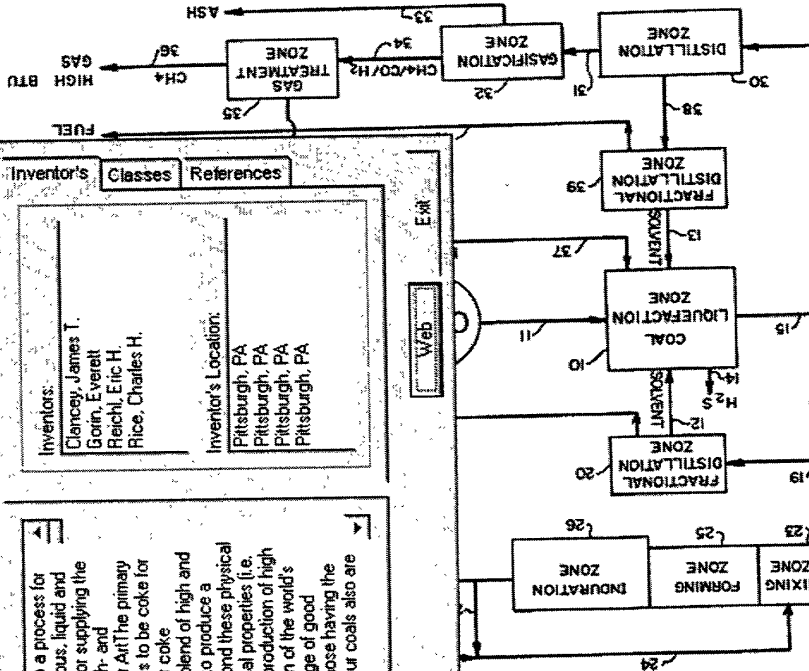


FIG. 49