

## LSEARCH(3C)

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### NAME

`lsearch` – linear search and update

### SYNOPSIS

```
#include <stdio.h>
#include <search.h>

char *lsearch ((char *)key, (char *)base,
nelp, sizeof(*key), compar)
unsigned *nelp;
int (*compar)( );

char *lfind ((char *)key, (char *)base,
nelp, sizeof(*key), compar)
unsigned *nelp;
int (*compar)( );
```

### DESCRIPTION

*Lsearch* is a linear search routine generalized from Knuth (6.1) Algorithm S. It returns a pointer into a table indicating where a datum may be found. If the datum does not occur, it is added at the end of the table. *Key* points to the datum to be sought in the table. *Base* points to the first element in the table. *Nelp* points to an integer containing the current number of elements in the table. The integer is incremented if the datum is added to the table. *Compar* is the name of the comparison function which the user must supply (*strcmp*, for example). It is called with two arguments that point to the elements being compared. The function must return zero if the elements are equal and non-zero otherwise.

*Lfind* is the same as *Lsearch* except that if the datum is not found, it is not added to the table. Instead, a NULL pointer is returned.

### NOTES

The pointers to the key and the element at the base of the table should be of type pointer-to-element, and cast to type pointer-to-character.

The comparison function need not compare every byte, so arbitrary data may be contained in the elements in addition to the values being compared.

Although declared as type pointer-to-character, the value returned should be cast into type pointer-to-element.

### EXAMPLE

This fragment will read in less than TABSIZE strings of length less than ELSIZE and store them in a table, eliminating duplicates.

```
#include <stdio.h>
#include <search.h>

#define TABSIZE 30
#define ELSIZE 120

char line[ELSIZE], tab[TABSIZE][ELSIZE],
*lsearch( )
```

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```
unsigned nel = 0;
int strcmp( );
...
while (fgets(line, ELSIZE, stdin) != NULL &&
      nel < TABSIZE)
    (void) lsearch(line, (char *)tab, &nel,
                  ELSIZE, strcmp);
```

### SEE ALSO

bsearch(3C), hsearch(3C), string(3C), tsearch(3C).

### DIAGNOSTICS

If the searched for datum is found, both *lsearch* and *lfind* return a pointer to it. Otherwise, *lfind* returns NULL and *lsearch* returns a pointer to the newly added element.

### BUGS

Undefined results can occur if there is not enough room in the table to add a new item.

## MALLOC(3C)

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### NAME

malloc, free, realloc, calloc – main memory allocator

### SYNOPSIS

```
char *malloc (size)
unsigned size;

void free (ptr)
char *ptr;

char *realloc (ptr, size)
char *ptr;
unsigned size;

char *calloc (nelem, elsize)
unsigned nelem, elsize;

int mallopt (cmd, value)
int cmd, value;
```

### DESCRIPTION

*Malloc* and *free* provide a simple general-purpose memory allocation package. *Malloc* returns a pointer to a block of at least *size* bytes suitably aligned for any use.

The argument to *free* is a pointer to a block previously allocated by *malloc*; after *free* is performed this space is made available for further allocation, but its contents are left undisturbed.

Undefined results will occur if the space assigned by *malloc* is overrun or if some random number is handed to *free*.

*Malloc* allocates the first big enough contiguous reach of free space found in a circular search from the last block allocated or freed, coalescing adjacent free blocks as it searches. It calls *sbrk* (see *brk(2)*) to get more memory from the system when there is no suitable space already free.

*Realloc* changes the size of the block pointed to by *ptr* to *size* bytes and returns a pointer to the (possibly moved) block. The contents will be unchanged up to the lesser of the new and old sizes. If no free block of *size* bytes is available in the storage arena, then *realloc* will ask *malloc* to enlarge the arena by *size* bytes and will then move the data to the new space.

*Realloc* also works if *ptr* points to a block freed since the last call of *malloc*, *realloc*, or *calloc*; thus sequences of *free*, *malloc* and *realloc* can exploit the search strategy of *malloc* to do storage compaction.

*Calloc* allocates space for an array of *nelem* elements of size *elsize*. The space is initialized to zeros.

*Mallopt* provides for control over the allocation algorithm. The available values for *cmd* are:

**M\_MXFAST** Set *maxfast* to *value*. The algorithm allocates all blocks below the size of *maxfast* in large groups and then does them out

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- M\_NLBLKS** Set *numblks* to *value*. The above mentioned "large groups" each contain *numblks* blocks. *Numblks* must be greater than 0. The default value for *numblks* is 100.
- M\_GRAIN** Set *grain* to *value*. The sizes of all blocks smaller than *maxfast* are considered to be rounded up to the nearest multiple of *grain*. *Grain* must be greater than 0. The default value of *grain* is the smallest number of bytes which will allow alignment of any data type. Value will be rounded up to a multiple of the default when *grain* is set.
- M\_KEEP** Preserve data in a freed block until the next *malloc*, *realloc*, or *calloc*. This option is provided only for compatibility with the old version of *malloc* and is not recommended.

These values are defined in the `<malloc.h>` header file.

*Mallopt* may be called repeatedly, but may not be called after the first small block is allocated.

Each of the allocation routines returns a pointer to space suitably aligned (after possible pointer coercion) for storage of any type of object.

### SEE ALSO

`brk(2)`.

### DIAGNOSTICS

*Malloc*, *realloc* and *calloc* return a NULL pointer if there is not enough available memory. When *realloc* returns NULL, the block pointed to by *ptr* is left intact. If *mallopt* is called after any allocation, or if *cmd* or *value* are invalid, non-zero is returned. Otherwise, it returns zero.

### NOTE

Search time increases when many objects have been allocated; that is, if a program allocates but never frees, then each successive allocation takes longer.

## MATHERR (3M)

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### NAME

`matherr` – error-handling function

### SYNOPSIS

```
#include <math.h>
int matherr (x)
struct exception *x;
```

### DESCRIPTION

*Matherr* is invoked by functions in the Math Library when errors are detected. Users may define their own procedures for handling errors by including a function named *matherr* in their programs. *Matherr* must be of the form described above. A pointer to the exception structure *x* will be passed to the user-supplied *matherr* function when an error occurs. This structure, which is defined in the `<math.h>` header file, is as follows:

```
struct exception {
    int type;
    char *name;
    double arg1, arg2, retval;
};
```

The element *type* is an integer describing the type of error that has occurred, from the following list of constants (defined in the header file):

DOMAIN	domain error
SING	singularity
OVERFLOW	overflow
UNDERFLOW	underflow
TLOSS	total loss of significance
PLOSS	partial loss of significance

The element *name* points to a string containing the name of the function that had the error. The variables *arg1* and *arg2* are the arguments to the function that had the error. *Retval* is a double that is returned by the function having the error. If it supplies a return value, the user's *matherr* must return non-zero. If the default error value is to be returned, the user's *matherr* must return 0.

If *matherr* is not supplied by the user, the default error-handling procedures, described with the math functions involved, will be invoked upon error. These procedures are also summarized in the table below. In every case, *errno* is set to non-zero and the program continues.

### EXAMPLE

```
matherr(x)
register struct exception *x;
{
    switch (x->type) {
    case DOMAIN:
    case SING: /* print message and abort */
        fprintf(stderr, "domain error in %s\n", x->name);
```

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