# INTERNATIONAL STANDARD 

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## Information technology - Coding of moving pictures and associated audio for digital storage media at up to about 1,5 Mbit/s - <br> Part 3: <br> Audio

Technologies de l'information - Codage de l'image animée et du son associé pour les supports de stockage numérique jusqu'à environ 1,5 Mbit/s -
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## Foreword

ISO (the International Organization for Standardization) and IEC (the International Electrotechnical Commission) form the specialized system for worldwide standardization. National bodies that are members of ISO or IEC participate in the development of International Standards through technical committees established by the respective organization to deal with particular fields of technical activity. ISO and IEC technical committees collaborate in fields of mutual interest. Other international organizations, governmental and non-governmental, in liaison with ISO and IEC, also take part in the work.

In the field of information technology, ISO and IEC have established a joint technical committee, ISO/IEC JTC 1. Draft International Standards adopted by the joint technical committee are circulated to national bodies for voting. Publication as an International Standard requires approval by at least $75 \%$ of the national bodies casting a vote.

International Standard ISO/IEC 11172-3 was prepared by Joint Technical Committee ISO/IEC JTC 1, Information technology, Sub-Committee SC 29 , Coded representation of audio, picture, multimedia and hypermedia information.

ISO/IEC 11172 consists of the following parts, under the general title information technology - Coding of moving pictures and associated audio for digital storage media at up to about $1,5 \mathrm{Mbit/s}$ :

- Part 1: Systems
- Part 2: Video
- Part 3: Audio
- Part 4: Compliance testing

Annexes $A$ and $B$ form an integral part of this part of ISO/IEC 11172. Annexes C, D, E, F, G and H are for information only.

## Introduction

Note: Readers interested in an overview of MPEG Audio should read this Introduction and then proceed to annex A (Diagrams) and annex C (The encoding process) before reading the normative clauses 1 and 2 .

To aid in the understanding of the specification of the stored compressed bitstream and its decoding, a sequence of encoding, storage and decoding is described.

### 0.1 Encoding

The encoder processes the digital audio signal and produces the compressed bitstream for storage. The encoder algorithm is not standardized, and may use various means for encoding such as estimation of the auditory masking threshold, quantization, and scaling. However, the encoder output must be such that a decoder conforming to the specifications of clause 2.4 will produce audio suitable for the intended application.


Figure 1 .- Sketch of the basic structure of an encoder
Figure 1 illustrates the basic structure of a audio encoder. Input audio samples are fed into the encoder. The mapping creates a filtered and subsampled representation of the input audio stream. The mapped samples may be called either subband samples (as in Layer I or II, see below) or transformed subband samples (as in Layer III). A psychoacoustic model creates a set of data to control the quantizer and coding. These data are different depending on the actual coder implementation. One possibility is to use an estimation of the masking threshold to do this quantizer control. The quantizer and coding block creates a set of coding symbols from the mapped input samples. Again, this block can depend on the encoding system. The block 'frame packing' assembles the actual bitstream from the output data of the other blocks, and adds other information (e.g. error correction) if necessary.

There are four different modes possible, single channel, dual channel (two independent audio signals coded within one bitstream), stereo (left and right signals of a stereo pair coded within one bitstream), and Joint Stereo (left and right signals of a stereo pair coded within one bitstream with the stereo irrelevancy and redundancy exploited).

### 0.2 Layers

Depending on the application, different layers of the coding system with increasing encoder complexity and performance can be used. An ISO/IEC 11172-3 Audio Layer N decoder is able to decode bitstream data which has been encoded in Layer N and all layers below N .

## Layer I

This layer contains the basic mapping of the digital audio input into 32 subbands, fixed segmentation to format the data into blocks, a psychoacoustic model to determine the adaptive bit allocation, and quantization using block companding and formatting. The theoretical minimum encoding/decoding delay for Layer I is about 19 ms .

## Layer II

This layer provides additional coding of bit allocation, scalefactors and samples. Different framing is used. The theoretical minimum encoding/decoding delay for Layer II is about 35 ms .

## Layer III

This layer introduces increased frequency resolution based on a hybrid filterbank. It adds a different (nonuniform) quantizer, adaptive segmentation and entropy coding of the quantized values. The theoretical minimum encoding/decoding delay for Layer III is about 59 ms .

Joint Stereo coding can be added as an additional feature to any of the layers.

### 0.3 Storage

Various streams of encoded video, encoded audio, synchronization data, systems data and auxiliary data may be stored together on a storage medium. Editing of the audio will be easier if the edit point is constrained to coincide with an addressable point.

Access to storage may involve remote access over a communication system. Access is assumed to be controlled by a functional unit other than the audio decoder itself. This control unit accepts user commands, reads and interprets data base structure information, reads the stored information from the media, demultiplexes non-audio information and passes the stored audio bitstream to the audio decoder at the required rate.

### 0.4 Decoding

The decoder accepts the compressed audio bitstream in the syntax defined in 2.4.1, decodes the data elements according to 2.4.2, and uses the information to produce digital audio output according to 2.4.3.


Figure 2 -- Sketch of the basic structure of a decoder
Figure 2 illustrates the basic structure of a audio decoder. Bitstream data is fed into the decoder. The bitstream unpacking and decoding block does error detection if error-check is applied in the encoder (see 2.4.2.4). The bitstream data are unpacked to recover the various pieces of information. The reconstruction block reconstructs the quantized version of the set of mapped samples. The inverse mapping transforms these mapped samples back into uniform PCM.

# Information technology - Coding of moving pictures and associated audio for digital storage media at up to about $1,5 \mathrm{Mbit} / \mathrm{s}$ - 

## Part 3: <br> Audio

## Section 1: General

### 1.1 Scope

This part of ISO/IEC 11172 specifies the coded representation of high quality audio for storage media and the method for decoding of high quality audio signals. The input of the encoder and the output of the decoder are compatible with existing PCM standards such as standard Compact Disc and Digital Audio Tape.

This part of the ISO/IEC 11172 is intended for application to digital storage media providing a total continuous transfer rate of about $1,5 \mathrm{Mbits} / \mathrm{sec}$ for both audio and video bitstreams, such as CD, DAT and magnetic hard disc. The storage media may either be connected directly to the decoder, or via other means such as communication lines and the ISO/IEC 11172 multiplexed stream defined in ISO/IEC 11172-1. This part of ISO/IEC 11172 is intended for sampling rates of $32 \mathrm{kHz}, 44,1 \mathrm{kHz}$, and 48 kHz .

### 1.2 Normative references

The following International Standards contain provisions which, through reference in this text, constitute provisions of this part of ISO/IEC 11172. At the time of publication, the editions indicated were valid. All standards are subject to revision, and parties to agreements based on this part of ISO/IEC 11172 are encouraged to investigate the possibility of applying the most recent editions of the standards indicated below. Members of IEC and ISO maintain registers of currently valid International Standards.

ISO/IEC 11172-1:1993 Information technology - Coding of moving pictures and associated audio for digital storage media at up to about 1,5 Mbit/s - Part 1 Systems.

ISO/IEC 11172-2:1993 Information technology - Coding of moving pictures and associated audio for digital storage media at up to about $1,5 \mathrm{Mbit} / \mathrm{s}$ - Part 2: Video.

CCIR Recommendation 601-2 Encoding parameters of digital television for studios.
CCIR Report 624-4 Characteristics of systems for monochrome and colour television.
CCIR Recommendation 648 Recording of audio signals.
CCIR Report 955-2 Sound broadcasting by satellite for portable and mobile receivers, including Annex IV Summary description of Advanced Digital System II.

CCITT Recommendation J. 17 Pre-emphasis used on Sound-Programme Circuits.
IEEE Draft Standard P1180/D2 1990 Specification for the implementation of $8 \times 8$ inverse discrete cosine transform".

IEC publication 908:1987 CD Digital Audio System.

## Section 2: Technical elements

### 2.1 Definitions

For the purposes of ISO/IEC 11172, the following definitions apply. If specific to a part, this is noted in square brackets.
2.1.1 ac coefficient [video]: Any DCT coefficient for which the frequency in one or both dimensions is non-zero.
2.1.2 access unit [system]: In the case of compressed audio an access unit is an audio access unit. In the case of compressed video an access unit is the coded representation of a picture.
2.1.3 adaptive segmentation [audio]: A subdivision of the digital representation of an audio signal in variable segments of time.
2.1.4 adaptive bit allocation [audio]: The assignment of bits to subbands in a time and frequency varying fashion according to a psychoacoustic model.
2.1.5 adaptive noise allocation [audio]: The assignment of coding noise to frequency bands in a time and frequency varying fashion according to a psychoacoustic model.
2.1.6 alias [audio]: Mirrored signal component resulting from sub-Nyquist sampling.
2.1.7 analysis filterbank [audio]: Filterbank in the encoder that transforms a broadband PCM audio signal into a set of subsampled subband samples.
2.1.8 audio access unit [audio]: For Layers I and II an audio access unit is defined as the smallest part of the encoded bitstream which can be decoded by itself, where decoded means "fully reconstructed sound". For Layer III an audio access unit is part of the bitstream that is decodable with the use of previously acquired main information.
2.1.9 audio buffer [audio]: A buffer in the system target decoder for storage of compressed audio data.
2.1.10 audio sequence [audio]: A non-interrupted series of audio frames in which the following parameters are not changed:

> - ID
> - Layer
> - Sampling Frequency
> - For Layer I and II: Bitrate index
2.1.11 backward motion vector [video]: A motion vector that is used for motion compensation from a reference picture at a later time in display order.
2.1.12 Bark [audio]: Unit of critical band rate. The Bark scale is a non-linear mapping of the frequency scale over the audio range closely corresponding with the frequency selectivity of the human ear across the band.
2.1.13 bidirectionally predictive-coded picture; B-picture [video]: A picture that is coded using motion compensated prediction from a past and/or future reference picture.
2.1.14 bitrate: The rate at which the compressed bitstream is delivered from the storage medium to the input of a decoder.
2.1.15 block companding [audio]: Normalizing of the digital representation of an audio signal within a certain time period.
2.1.16 block [video]: An 8-row by 8 -column orthogonal block of pels.
2.1.17 bound [audio]: The lowest subband in which intensity stereo coding is used.
2.1.18 byte aligned: A bit in a coded bitstream is byte-aligned if its position is a multiple of 8-bits from the first bit in the stream.
2.1.19 byte: Sequence of 8 -bits.
2.1 .20 channel: A digital medium that stores or transports an ISO/IEC 11172 stream.
2.1.21 channel [audio]: The left and right channels of a stereo signal
2.1.22 chrominance (component) [video]: A matrix, block or single pel representing one of the two colour difference signals related to the primary colours in the manner defined in CCIR Rec 601. The symbols used for the colour difference signals are Cr and Cb .
2.1.23 coded audio bitstream [audio]: A coded representation of an audio signal as specified in this part of ISO/IEC 11172.
2.1.24 coded video bitstream [video]: A coded representation of a series of one more pictures as specified in ISO/IEC 11172-2.
2.1.25 coded order [video]: The order in which the pictures are stored and decoded. This order is not necessarily the same as the display order.
2.1.26 coded representation: A data element as represented in its encoded form.
2.1.27 coding parameters [video]: The set of user-definable parameters that characterize a coded video bitstream. Bitstreams are characterised by coding parameters. Decoders are characterised by the bitstreams that they are capable of decoding.
2.1.28 component [video]: A matrix, block or single pel from one of the three matrices (Iuminance and two chrominance) that make up a picture.
2.1.29 compression: Reduction in the number of bits used to represent an item of data.
2.1.30 constant bitrate coded video [video]: A compressed video bitstream with a constant average bitrate.
2.1.31 constant bitrate: Operation where the bitrate is constant from start to finish of the compressed bitstream.
2.1.32 constrained parameters [video]: The values of the set of coding parameters defined in 2.4.3.2 of ISO/IEC 11172-2.
2.1.33 constrained system parameter stream (CSPS) [system]: An ISO/IEC 11172 multiplexed stream for which the constraints defined in 2.4.6 of ISO/IEC 11172-1 apply.
2.1.34 CRC: Cyclic redundancy code.
2.1.35 critical band rate [audio]: Psychoacoustic function of frequency. At a given audible frequency it is proportional to the number of critical bands below that frequency. The units of the critical band rate scale are Barks.
2.1.36 critical band [audio]: Psychoacoustic measure in the spectral domain which corresponds to the frequency selectivity of the human ear. This selectivity is expressed in Bark.
2.1.37 data element: An item of data as represented before encoding and after decoding.
2.1.38 dc-coefficient [video]: The DCT coefficient for which the frequency is zero in both dimensions.
2.1.39 de-coded picture; D-picture [video]: A picture that is coded using only information from itself. Of the DCT coefficients in the coded representation, only the dc-coefficients are present.
2.1.40 DCT coefficient: The amplitude of a specific cosine basis function.
2.1.41 decoded stream: The decoded reconstruction of a compressed bitstream.
2.1.42 decoder input buffer [video]: The first-in first-out (FIFO) buffer specified in the video buffering verifier.
2.1.43 decoder input rate [video]: The data rate specified in the video buffering verifier and encoded in the coded video bitstream.
2.1.44 decoder: An embodiment of a decoding process.
2.1.45 decoding (process): The process defined in ISO/IEC 11172 that reads an input coded bitstream and produces decoded pictures or audio samples.
2.1.46 decoding time-stamp; DTS [system]: A field that may be present in a packet header that indicates the time that an access unit is decoded in the system target decoder.
2.1.47 de-emphasis [audio]: Filtering applied to an audio signal after storage or transmission to undo a linear distortion due to emphasis.
2.1.48 dequantization [video]: The process of rescaling the quantized DCT coefficients after their representation in the bitstream has been decoded and before they are presented to the inverse DCT.
2.1.49 digital storage media; DSM: A digital storage or transmission device or system.
2.1.50 discrete cosine transform; DCT [video]: Either the forward discrete cosine transform or the inverse discrete cosine transform. The DCT is an invertible, discrete orthogonal transformation. The inverse DCT is defined in annex A of ISO/IEC 11172-2.
2.1.51 display order [video]: The order in which the decoded pictures should be displayed. Normally this is the same order in which they were presented at the input of the encoder.
2.1.52 dual channel mode [audio]: A mode, where two audio channels with independent programme contents (e.g. bilingual) are encoded within one bitstream. The coding process is the same as for the stereo mode.
2.1.53 editing: The process by which one or more compressed bitstreams are manipulated to produce a new compressed bitstream. Conforming edited bitstreams must meet the requirements defined in this ISO/IEC 11172.
2.1.54 elementary stream [system]: A generic term for one of the coded video, coded audio or other coded bitstreams.
2.1.55 emphasis [audio]: Filtering applied to an audio signal before storage or transmission to improve the signal-to-noise ratio at high frequencies.
2.1.56 encoder: An embodiment of an encoding process.
2.1.57 encoding (process): A process, not specified in ISO/IEC 11172, that reads a stream of input pictures or audio samples and produces a valid coded bitstream as defined in ISO/IEC 11172.
2.1.58 entropy coding: Variable length lossless coding of the digital representation of a signal to reduce redundancy.
2.1.59 fast forward playback [video]: The process of displaying a sequence, or parts of a sequence, of pictures in display-order faster than real-time.

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2.1.60 FFT: Fast Fourier Transformation. A fast algorithm for performing a discrete Fourier transform (an orthogonal transform).
2.1.61 filterbank [audio]: A set of band-pass filters covering the entire audio frequency range.
2.1.62 fixed segmentation [audio]: A subdivision of the digital representation of an audio signal into fixed segments of time.
2.1.63 forbidden: The term "forbidden" when used in the clauses defining the coded bitstream indicates that the value shall never be used. This is usually to avoid emulation of start codes.
2.1.64 forced updating [video]: The process by which macroblocks are intra-coded from time-to-time to ensure that mismatch errors between the inverse DCT processes in encoders and decoders cannot build up excessively.
2.1.65 forward motion vector [video]: A motion vector that is used for motion compensation from a reference picture at an earlier time in display order.
2.1.66 frame [audio]: A part of the audio signal that corresponds to audio PCM samples from an Audio Access Unit.
2.1.67 free format [audio]: Any bitrate other than the defined bitrates that is less than the maximum valid bitrate for each layer.
2.1.68 future reference picture [video]: The future reference picture is the reference picture that occurs at a later time than the current picture in display order.
2.1.69 granules [Layer II] [audio]: The set of 3 consecutive subband samples from all 32 subbands that are considered together before quantization. They correspond to 96 PCM samples.
2.1.70 granules [Layer III] \{audio]: 576 frequency lines that carry their own side information.
2.1.71 group of pictures [video]: A series of one or more coded pictures intended to assist random access. The group of pictures is one of the layers in the coding syntax defined in ISO/IEC 11172-2.
2.1.72 Hann window [audio]: A time function applied sample-by-sample to a block of audio samples before Fourier transformation.
2.1.73 Huffman coding: A specific method for entropy coding.
2.1.74 hybrid filterbank [audio]: A serial combination of subband filterbank and MDCT.
2.1.75 IMDCT [audio]: Inverse Modified Discrete Cosine Transform.
2.1.76 intensity stereo [audio]: A method of exploiting stereo irrelevance or redundancy in stereophonic audio programmes based on retaining at high frequencies only the energy envelope of the right and left channels.
2.1.77 interlace [video]: The property of conventional television pictures where alternating lines of the picture represent different instances in time.
2.1.78 intra coding [video]: Coding of a macroblock or picture that uses information only from that macroblock or picture.
2.1.79 intra-coded picture; I-picture [video]: A picture coded using information only from itself.
2.1.80 ISO/IEC 11172 (multiplexed) stream [system]: A bitstream composed of zero or more elementary streams combined in the manner defined in ISO/IEC 11172-1.
2.1.81 joint stereo coding [audio]: Any method that exploits stereophonic irrelevance or stereophonic redundancy.
2.1.82 joint stereo mode [audio]: A mode of the audio coding algorithm using joint stereo coding.
2.1.83 layer [audio]: One of the levels in the coding hierarchy of the audio system defined in this part of ISO/IEC 11172.
2.1.84 layer [video and systems]: One of the levels in the data hierarchy of the video and system specifications defined in ISO/IEC 11172-1 and ISO/IEC 11172-2.
2.1.85 luminance (component) [video]: A matrix, block or single pel representing a monochrome representation of the signal and related to the primary colours in the manner defined in CCIR Rec 601. The symbol used for luminance is Y .
2.1.86 macroblock [video]: The four 8 by 8 blocks of luminance data and the two corresponding 8 by 8 blocks of chrominance data coming from a 16 by 16 section of the luminance component of the picture. Macroblock is sometimes used to refer to the pel data and sometimes to the coded representation of the pel values and other data elements defined in the macroblock layer of the syntax defined in ISO/IEC 11172-2. The usage is clear from the context.
2.1.87 mapping [audio]: Conversion of an audio signal from time to frequency domain by subband filtering and/or by MDCT.
2.1.88 masking [audio]: A property of the human auditory system by which an audio signal cannot be perceived in the presence of another audio signal .
2.1.89 masking threshold [audio]: A function in frequency and time below which an audio signal cannot be perceived by the human auditory system.
2.1.90 MDCT [audio]: Modified Discrete Cosine Transform.
2.1.91 motion compensation [video]: The use of motion vectors to improve the efficiency of the prediction of pel values. The prediction uses motion vectors to provide offsets into the past and/or future reference pictures containing previously decoded pel values that are used to form the prediction error signal.
2.1.92 motion estimation [video]: The process of estimating motion vectors during the encoding process.
2.1.93 motion vector [video]: A two-dimensional vector used for motion compensation that provides an offset from the coordinate position in the current picture to the coordinates in a reference picture.
2.1.94 MS stereo [audio]: A method of exploiting stereo irrelevance or redundancy in stereophonic audio programmes based on coding the sum and difference signal instead of the left and right channels.
2.1.95 non-intra coding [video]: Coding of a macroblock or picture that uses information both from itself and from macroblocks and pictures occurring at other times.
2.1.96 non-tonal component [audio]: A noise-like component of an audio signal.
2.1.97 Nyquist sampling: Sampling at or above twice the maximum bandwidth of a signal.
2.1.98 pack [system]: A pack consists of a pack header followed by one or more packets. It is a layer in the system coding syntax described in ISO/IEC 11172-1.
2.1.99 packet data [system]: Contiguous bytes of data from an elementary stream present in a packet.
2.1.100 packet header [system]: The data structure used to convey information about the elementary stream data contained in the packet data.
2.1.101 packet [system]: A packet consists of a header followed by a number of contiguous bytes from an elementary data stream. It is a layer in the system coding syntax described in ISO/IEC 11172-1.
2.1.102 padding [audio]: A method to adjust the average length in time of an audio frame to the duration of the corresponding PCM samples, by conditionally adding a slot to the audio frame.
2.1.103 past reference picture [video]: The past reference picture is the reference picture that occurs at an earlier time than the current picture in display order.
2.1.104 pel aspect ratio [video]: The ratio of the nominal vertical height of pel on the display to its nominal horizontal width.
2.1.105 pel [video]: Picture element.
2.1.106 picture period [video]: The reciprocal of the picture rate.
2.1.107 picture rate [video]: The nominal rate at which pictures should be output from the decoding process.
2.1.108 picture [video]: Source, coded or reconstructed image data. A source or reconstructed picture consists of three rectangular matrices of 8 -bit numbers representing the luminance and two chrominance signals. The Picture layer is one of the layers in the coding syntax defined in ISO/IEC 11172-2. Note that the term "picture" is always used in ISO/IEC 11172 in preference to the terms field or frame.
2.1.109 polyphase filterbank [audio]: A set of equal bandwidth filters with special phase interrelationships, allowing for an efficient implementation of the filterbank.
2.1.110 prediction [video]: The use of a predictor to provide an estimate of the pel value or data element currently being decoded.
2.1.111 predictive-coded picture; P-picture [video]: A picture that is coded using motion compensated prediction from the past reference picture.
2.1.112 prediction error [video]: The difference between the actual value of a pel or data element and its predictor.
2.1.113 predictor [video]: A linear combination of previously decoded pel values or data elements.
2.1.114 presentation time-stamp; PTS [system]: A field that may be present in a packet header that indicates the time that a presentation unit is presented in the system target decoder.
2.1.115 presentation unit; PU [system]: A decoded audio access unit or a decoded picture.
2.1.116 psychoacoustic model [audio]: A mathematical model of the masking behaviour of the human auditory system.
2.1.117 quantization matrix [video]: A set of sixty-four 8 -bit values used by the dequantizer.
2.1.118 quantized DCT coefficients [video]: DCT coefficients before dequantization. A variable length coded representation of quantized DCT coefficients is stored as part of the compressed video bitstream.
2.1.119 quantizer scalefactor [video]: A data element represented in the bitstream and used by the decoding process to scale the dequantization.
2.1.120 random access: The process of beginning to read and decode the coded bitstream at an arbitrary point.
2.1.121 reference picture [video]: Reference pictures are the nearest adjacent I- or P-pictures to the current picture in display order.
2.1.122 reorder buffer [video]: A buffer in the system target decoder for storage of a reconstructed Ipicture or a reconstructed P -picture.
2.1.123 requantization [audio]: Decoding of coded subband samples in order to recover the original quantized values.
2.1.124 reserved: The term "reserved" when used in the clauses defining the coded bitstream indicates that the value may be used in the future for ISO/IEC defined extensions.
2.1.125 reverse playback [video]: The process of displaying the picture sequence in the reverse of display order.
2.1.126 scalefactor band [audio]: A set of frequency lines in Layer III which are scaled by one scalefactor.
2.1.127 scalefactor index [audio]: A numerical code for a scalefactor
2.1.128 scalefactor [audio]: Factor by which a set of values is scaled before quantization.
2.1.129 sequence header [video]: A block of data in the coded bitstream containing the coded representation of a number of data elements.
2.1.130 side information: Information in the bitstream necessary for controlling the decoder.
2.1.131 skipped macroblock [video]: A macroblock for which no data are stored.
2.1.132 slice [video]: A series of macroblocks. It is one of the layers of the coding syntax defined in ISO/IEC 11172-2.
2.1.133 slot [audio]: A slot is an elementary part in the bitstream. In Layer I a slot equals four bytes, in Layers II and III one byte.
2.1.134 source stream: A single non-multiplexed stream of samples before compression coding.
2.1.135 spreading function [audio]: A function that describes the frequency spread of masking.
2.1.136 start codes [system and video]: 32 -bit codes embedded in that coded bitstream that are unique. They are used for several purposes including identifying some of the layers in the coding syntax.
2.1.137 STD input buffer [system]: A first-in first-out buffer at the input of the system target decoder for storage of compressed data from elementary streams before decoding.
2.1.138 stereo mode [audio]: Mode, where two audio channels which form a stereo pair (left and right) are encoded within one bitstream. The coding process is the same as for the dual channel mode.
2.1.139 stuffing (bits); stuffing (bytes) : Code-words that may be inserted into the compressed bitstream that are discarded in the decoding process. Their purpose is to increase the bitrate of the stream.
2.1.140 subband [audio]: Subdivision of the audio frequency band.
2.1.141 subband filterbank [audio]: A set of band filters covering the entire audio frequency range. In this part of ISO/IEC 11172 the subband filterbank is a polyphase filterbank.
2.1.142 subband samples [audio]: The subband filterbank within the audio encoder creates a filtered and subsampled representation of the input audio stream. The filtered samples are called subband samples.

From 384 time-consecutive input audio samples, 12 time-consecutive subband samples are generated within each of the 32 subbands.
2.1.143 syncword [audio]: A 12-bit code embedded in the audio bitstream that identifies the start of a frame.
2.1.144 synthesis filterbank [audio]: Filterbank in the decoder that reconstructs a PCM audio signal from subband samples.
2.1.145 system header [system]: The system header is a data structure defined in ISO/IEC 11172-1 that carries information summarising the system characteristics of the ISO/IEC 11172 multiplexed stream.
2.1.146 system target decoder; STD [system]: A hypothetical reference model of a decoding process used to describe the semantics of an ISO/IEC 11172 multiplexed bitstream.
2.1.147 time-stamp [system]: A term that indicates the time of an event.
2.1.148 triplet [audio]: A set of 3 consecutive subband samples from one subband. A triplet from each of the 32 subbands forms a granule.
2.1.149 tonal component [audio]: A sinusoid-like component of an audio signal.
2.1.150 variable bitrate: Operation where the bitrate varies with time during the decoding of a compressed bitstream.
2.1.151 variable length coding; VLC: A reversible procedure for coding that assigns shorter codewords to frequent events and longer code-words to less frequent events.
2.1.152 video buffering verifier; VBV [video]: A hypothetical decoder that is conceptually connected to the output of the encoder. Its purpose is to provide a constraint on the variability of the data rate that an encoder or editing process may produce.
2.1.153 video sequence [video]: A series of one or more groups of pictures. It is one of the layers of the coding syntax defined in ISO/IEC 11172-2.
2.1.154 zig-zag scanning order [video]: A specific sequential ordering of the DCT coefficients from (approximately) the lowest spatial frequency to the highest.

### 2.2 Symbols and abbreviations

The mathematical operators used to describe this International Standard are similar to those used in the C programming language. However, integer division with truncation and rounding are specifically defined. The bitwise operators are defined assuming twos-complement representation of integers. Numbering and counting loops generally begin from zero.

### 2.2.1 Arithmetic operators

| $+$ | Addition. |
| :---: | :---: |
| - | Subtraction (as a binary operator) or negation (as a unary operator). |
| ++ | Increment. |
| -- | Decrement. |
| * | Multiplication. |
| $\wedge$ | Power. |
| / | Integer division with truncation of the result toward zero. For example, 7/4 and -7/-4 are truncated to 1 and -7/4 and 7/-4 are truncated to -1 . |
| // | Integer division with rounding to the nearest integer. Half-integer values are rounded away from zero unless otherwise specified. For example $3 / / 2$ is rounded to 2 , and $-3 / / 2$ is rounded to $\mathbf{- 2}$. |
| DIV | Integer division with truncation of the result towards $-\infty$. |
| 11 | $\begin{array}{ll} \text { Absolute value. } \quad & \|x\|=x \text { when } x>0 \\ & \|x\|=0 \text { when } x==0 \\ & \|x\|=-x \text { when } x<0 \end{array}$ |
| \% | Modulus operator. Defined only for positive numbers. |
| Sign( ) | $\begin{array}{rll} \operatorname{Sign}(x) & =1 & x>0 \\ 0 & x==0 \\ -1 & x<0 \end{array}$ |
| NINT ( ) | Nearest integer operator. Returns the nearest integer value to the real-valued argument. Halfinteger values are rounded away from zero. |
| $\sin$ | Sine. |
| cos | Cosine. |
| exp | Exponential. |
| $\checkmark$ | Square root. |
| $\log _{10}$ | Logarithm to base ten. |
| $\log _{\mathrm{e}}$ | Logarithm to base e. |
| $\log _{2}$ | Logarithm to base 2. |
| 2.2.2 | gical operators |


| ॥ | Logical OR. |
| :--- | :--- |
| \&\& | Logical AND. |


| $l$ | Logical NOT |
| :--- | :--- |
| $\mathbf{2 . 2 . 3}$ | Relational operators |
| $>$ | Greater than. |
| $<$ | Greater than or equal to. |
| $=$ | Less than. |
| $=$ | Less than or equal to. |
| $=$ | Equal to. |
| $\max [\ldots .]$. | the maximum value in the argument list. |
| $\min [\ldots . .]$. | the minimum value in the argument list. |
| 2.2 .4 | Bitwise operators |

A twos complement number representation is assumed where the bitwise operators are used.

| $\&$ | AND |
| :--- | :--- |
| 1 | OR |
| $\gg$ | Shift right with sign extension. |
| $\ll$ | Shift left with zero fill. |
| $\mathbf{2 . 2 . 5}$ | Assignment |
| $=$ | Assignment operator. |
| 2.2.6 | Mnemonics |

The following mnemonics are defined to describe the different data types used in the coded bit-stream.

| bslbf | Bit string, left bit first, where "left" is the order in which bit strings are written in ISO/IEC 11172. Bit strings are written as a string of 1 s and 0 s within single quote marks, e.g. ' 10000001 '. Blanks within a bit string are for ease of reading and have no significance. |
| :---: | :---: |
| ch | Channel. If ch has the value 0 , the left channel of a stereo signal or the first of two independent signals is indicated. (Audio) |
| nct | Number of channels; equal to 1 for single_cbannel mode, 2 in other modes. (Audio) |
| gr | Granule of $3 * 32$ subband samples in audio Layer II, $18 * 32$ sub-band samples in audio Layer III. (Audio) |
| main_data | The main_data portion of the bitstream contains the scalefactors, Huffman encoded data, and ancillary information. (Audio) |
| main_data_beg | The location in the bitstream of the beginning of the main_data for the frame. The location is equal to the ending location of the previous frame's main_data plus one bit. It is calculated from the main_data_end value of the previous frame. (Audio) |
| part2_length | The number of main_data bits used for scalefactors. (Audio) |


| rpchof | Remainder polynomial coefficients, highest order first. (Audio) |
| :--- | :--- |
| sb | Subband. (Audio) |
| sblimit | The number of the lowest sub-band for which no bits are allocated. (Audio) |
| scfsi | Scalefactor selection information. (Audio) |
| switch_point_l | Number of scalefactor band (long block scalefactor band) from which point on window <br> switching is used. (Audio) |
| switch_point_s | Number of scalefactor band (short block scalefactor band) from which point on window <br> switching is used. (Audio) |
| uimsbf | Unsigned integer, most significant bit first. |
| viclbf | Variable length code, left bit first, where "left" refers to the order in which the VLC <br> codes are writen. |
| window | Number of the actual time slot in case of block_type $=2,0 \leq w i n d o w ~$$\quad$ (Audio) |

The byte order of multi-byte words is most significant byte first

### 2.2.7 Constants

```
\pi 3,14159265358...
```

```
e 2,71828182845\ldots
```


### 2.3 Method of describing bitstream syntax

The bitstream retrieved by the decoder is described in 2.4.1. Each data item in the bitstream is in bold type. It is described by its name, its length in bits, and a mnemonic for its type and order of transmission.

The action caused by a decoded data element in a bitstream depends on the value of that data element and on data elements previously decoded. The decoding of the data elements and definition of the state variables used in their decoding are described in 2.4.2. The following constructs are used to express the conditions when data elements are present, and are in normal type:

Note this syntax uses the 'C"-code convention that a variable or expression evaluating to a non-zero value is equivalent to a condition that is true.

```
while ( condition ) { If the condition is true, then the group of data elements occurs next
    data_element in the data stream. This repeats until the condition is not true.
)
do {
    data_element
        . . .
} while (condition)
if (condition) {
        data_element
            . . .
}
else { If the condition is not true, then the second group of data elements
    data_element
} in the data stream. This repeats until the condition is not true.
    The data element always occurs at least once.
    The data element is repeated until the condition is not true.
    If the condition is true, then the first group of data elements occurs
    next in the data stream.
    occurs next in the data stream.
```

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$$
\begin{aligned}
& \text { fo: (expr1; expr2; expr3) \{ expr1 is an expression specifying the initialization of the loop. Normally it } \\
& \text { data_element } \\
& \text { specifies the initial state of the counter. expr2 is a condition specifying a test } \\
& \text { made before each iteration of the loop. The loop terminates when the condition } \\
& \text { is not true. expr3 is an expression that is performed at the end of each iteration } \\
& \text { of the loop, normally it increments a counter. }
\end{aligned}
$$

Note that the most common usage of this construct is as follows:

$$
\begin{array}{ll}
\text { for }(i=0 ; i<n ; i++) & \begin{array}{l}
\text { The group of data elements occurs } n \text { times. Conditional constructs } \\
\text { within the group of data elements may depend on the value of the }
\end{array} \\
\text { data_element } \\
\text { loop control variable } i \text {, which is set to zero for the first occurrence, } \\
\text { incremented to one for the second occurrence, and so forth. }
\end{array}
$$

As noted, the group of data elements may contain nested conditional constructs. For compactness, the \{\} may be omitted when only one data element follows.

data_element [] \begin{tabular}{l}
data_element $\square$ is an array of data. The number of data elements is indicated by <br>
the context.

$\quad$

data_element $[n]$ is the $n+1$ th element of an array of data.
\end{tabular}

While the syntax is expressed in procedural terms, it should not be assumed that 2.4.3 implements a satisfactory decoding procedure. In particular, it defines a correct and error-free input bitstream. Actual decoders must include a means to look for start codes in order to begin decoding correctly.

## Definition of bytealigned function

The function bytealigned $O$ returns 1 if the current position is on a byte boundary, that is the next bit in the bitstream is the first bit in a byte. Otherwise it returns 0 .

## Definition of nextbits function

The function nextbits $O$ permits comparison of a bit string with the next bits to be decoded in the bitstream.

## Definition of next_start_code function

The next_start_code function removes any zero bit and zero byte stuffing and locates the next start code.

| S yntax | No. of bits | Mnemonic |
| :---: | :---: | :---: |
| ```next_start_code0 { while ( IbytealignedO) zero_bit while ( nextbits0 != '0000 00000000 00000000 0001') zero_byte``` |  |  |
|  | 1 | '0' |
|  | 8 | '00000000' |

This function checks whether the current position is bytealigned. If it is not, zero stuffing bits are present. After that any number of zero bytes may be present before the start-code. Therefore start-codes are always bytzaligned and may be preceded by any number of zero stuffing bits.

### 2.4 Requirements

### 2.4.1 Specification of the coded audio bitstream syntax

2.4.1.1 Audio sequence

| Syntax | No. of bits Mnemonic |
| :---: | :---: |
| ```audio sequence0 [ while (nextbits()==syncword) { frame() }``` |  |

2.4.1.2 Audio frame

| Syntax | No. of bits Mnemonic |
| :---: | :---: |
| frame( |  |
| \{ <br> header() <br> error_check() <br> audio_data0 <br> ancillary_data0 |  |

2.4.1.3 Header

| Syntax | No. of bits | Mnemonic |
| :---: | :---: | :---: |
| header() |  |  |
| syncword | 12 | bslbf |
| ID |  | bslbf |
| layer | 2 | bslbf |
| protection_bit | 1 | bslbf |
| bitrate_index | 4 | bsibf |
| sampling_frequency | 2 | bslbf |
| padding_bit | 1 | bsibf |
| private_bit | 1 | bslbf |
| mode | 2 | bslbf |
| mode_extension | 2 | bslbf |
| copyright | 1 | bslbf |
| original/copy | 1 | bslbf |
| emphasis | 2 | bslbf |

2.4.1.4 Error check

2.4.1.5 Audio data, Layer I

| Symtax | No. of bits | Mnemonic |
| :---: | :---: | :---: |
| ```audio_data0 ( for ( \(\mathrm{sb}=0\); \(\mathrm{sb}<\mathrm{bound} ; \mathrm{sb}++\) )``` |  |  |
|  |  |  |
| for (ch=0; ch<nch; ch + + ) |  | uimsbf |
| for (sb=bound; sb<32; sb++) | 4 | uimsbf |
| allocation[0][sb] | 4 | uimsbf |
| allocation $[1][\mathrm{sb}]=$ allocation $[0][\mathrm{sb}]$ |  |  |
| for ( $\mathrm{sb}=0 \cdot \mathrm{sb}<32 \cdot \mathrm{sb}++$ ) |  |  |
| for (ch $=0 ; \mathrm{ch}<\mathrm{nch}$; ch++) if (allocation[ch][sb]! $=0$ ) scalefactor[ch][sb] | 6 | uimsbf |
| for ( $s=0 ; s<12 ; s++$ ) <br> for ( $\mathrm{sb}=0$; $\mathrm{sb}<$ bound; $\mathrm{sb}++$ ) <br> for (ch=0; ch<nch; ch++) |  |  |
| sample[ch][sb][s] | $2 . .15$ | uimsbe |
| for ( $\mathrm{sb}=$ bound; $\mathrm{sb}<32 ; \mathrm{sb}++$ ) |  |  |
| if (allocation[0][sb]!=0) sample[0][sb][s] | $2 . .15$ | uimsbf |
| ) |  |  |

2.4.1.6 Audio data, Layer II


### 2.4.1.7 Audio data, Layer III

| Syntax | No. of bits | Mnemonic |
| :---: | :---: | :---: |
| audio_data0 f |  |  |
| main_data_begin <br> if (mode=_single_channel) | 9 | uimsbf |
| private_bits | 5 | bslbf |
| private_bits <br> for (ch=0; ch<nch; ch++) | 3 | bslbf |
| for (scfsi_band=0; scfsi_band<4; scfsi_band++) scfsi[ch][scfsi_band] | 1 | bsibf |
| for (gr=0; $\mathrm{gr}<2 ; \mathrm{gr}++$ ) <br> for (ch=0; ch<nch; ch++) \{ |  |  |
| part2_3_length [gr][ch] | 12 | uimsbf |
| big_values[gr][ch] | 9 | uimsbf |
| global_gain[gr][ch] | 8 | uimsbf |
| scalefac_compress[gr][ch] window_switching flag gr$][\mathrm{ch}]$ | 4 | bslbf |
| window_switching_flag[gr][ch] if (window_switching_flag[gr][ch]) \{ | 1 | bslbf |
| block_type[gr][ch] | 2 | bsibf |
| mixed_block_flag[gr][ch] for (region=0; region<2; region++) | 1 | uimsbe |
| table_select[gr][ch][region] | 5 | bslbf |
| for ( window $=0$; window $<3$; window++) subblock_gain[gr][ch][window] | 3 | uimsbf |
| else ( |  |  |
| for (region $=0$; region<3; region++) |  |  |
| table_select[gr][ch][region] | 5 | bslbf |
| region0_count[gr][ch] | 4 | bslbf |
| region 1_count[gr][ch] | 3 | bslbf |
| ) | 1 | bslbf |
| scalefac_scale[gr][ch] | 1 | bsibf |
| count1table_select[gr][ch] | 1 | bslbf |
| $\begin{gathered} \text { \} } \\ \text { main_data } 0 \end{gathered}$ |  |  |

The main data sitstream is defined below. The main_data field in the audio_data() syntax contains bytes from the main data bistream. However, because of the variable nature of Huffman coding used in Layer III, the main data for a frame does not generally follow the header and side information for that frame. The main_data for a frame starts at a location in the bitstream preceeding the header of the frame at a negative offset given by the value of main_data_beg in . (See definition of main_data_begin and figure A.7.a).

```
Syntax No. of bits Mnemonic
main_dataO
{
    for (gr=(1; gr<2; gr++) (
            for (ch=0; ch<nch; ch++) (
                if ((window_switching_flag[gr][ch]==1)
                                && (block_type[gr][ch]==2)) {
                if (mixed_block_flag[gr][ch]) {
                for (sfb=0; sfb<8; sfb++)
                    scalefac_l[gr][ch][sfb] 0..4 uimsbf
                        for (sfb=3; sfb<12; sfb++)
                                for (window=0; window<3; window++)
                                scalefac_s[gr][ch][sfb][window]
                                0..4 uimsbf
                    }
                    else {
                        for (sfb=0; sfb<12; sfb++)
                            for (window=0; window<3; window++)
                                    scalefac_s[gr][ch][sfb][window]
                                    0..4 uimsbf
                                    }
            l
            else {
                if ((scfsi[ch][0]==0) | (gr=0))
                        for(sfb=0;sfb<6;sfb++)
                            scalefac_l[gr][ch][sfb] 0..4 uimsbf
                if ((scfsi[ch][1]==0)|(gr== ))
                for(sfb=6;sfb<11;sfb++)
                        scalefac_l[gr][ch][sfb] 0..4 uimsbf
                if ((scfsi[ch][2]==0)|(gr== ))
                for(sfb=11;sfb<16;sfb++)
                        scalefac_l[gr][ch][sfb] 0..3 uimsbf
                if ((scfsi[ch][3]==0) il (gr == 0))
                for(sfb=16;sfb<21;sfb}++
                                    scalefac_l[gr][ch][sfb] 0..3 uimsbf
            }
            HuffmancodebitsO
        )
    }
    for (b=0; b<no_of_ancillary_bits; b++)
        ancillary_bit 1 bslbf
```

| Syntax | No. of bits | Mnemonic |
| :---: | :---: | :---: |
| Huffmancodebits0 \{ for ( $1=0 ; 1<$ big_values $* 2 ; 1+=2$ ) ( |  |  |
|  |  |  |
| hcod [\|x|][|y|] <br> if ( $\|x\|==15$ \& \& linbits>0) | $0 . .19$ | bslbf |
| $\begin{aligned} & \text { linbits } x \\ & \text { if }(x!=0) \end{aligned}$ | $1 . .13$ | uimsbf |
| $\operatorname{signx}$ | 1 | bslbf |
| linbitsy | $1 . .13$ | uimsbf |
| if ( $\mathrm{y}!=0$ ) |  |  |
| is signy | 1 | bsibf |
| is [1] $=x$ <br> is $[1+1]=y$ |  |  |
| $\}$ <br> for ( $1<$ big values* $2+$ count $1 * 4 ; 1+=4$ ) \{ |  |  |
| if ( $\mathrm{v}!=0$ ) |  |  |
| $\begin{aligned} & \quad \text { signv } \\ & \text { if }(w!=0) \end{aligned}$ | 1 | bslbf |
| $\operatorname{sign} w$ <br> if ( $\mathrm{x}:=0$ ) | 1 | bslbf |
| signx | 1 | bslbf |
| if ( y ! $=0$ ) |  |  |
| signy | 1 | bslbf |
| $\begin{aligned} & \text { is }[1]=\mathrm{v} \\ & \text { is }[1+1]=\mathrm{w} \end{aligned}$ |  |  |
| is $[1+2]=x$ |  |  |
| is $[1+3]=y$ |  |  |
| $\begin{aligned} & \text { for }(: 1<576 ; 1++) \\ & \quad i s[1]=0 \end{aligned}$ |  |  |
|  |  |  |
| $\square$ |  |  |
| 2.4.1.8 Ancillary data |  |  |
| Syntax | No. of bits | Mnemonic |
| ```ancillary_data0 ( if ((layer == 1)\| (layer == 2)) for (b=0;b<no_of_ancillary_bits; b++) ancillary_bit``` |  |  |
|  |  |  |
|  | 1 | bslbf |

### 2.4.2 Semantics for the audio bitstream syntax

### 2.4.2.1 Audio sequence general

frame - Layer I and Layer II: Part of the bitstream that is decodable by itself. In Layer I it contains information for 384 samples and in Layer II for 1152 samples. It starts with a syncword, and ends just before the next syncword. It consists of an integer number of slots (four bytes in Layer I, one byte in Layer II).
-- Layer III: $\quad$ Part of the bitstream that is decodable with the use of previously acquired main information. In Layer III it contains information for 1152 samples. Although the distance between the start of consecutive syncwords is an integer number of slots (one byte in Layer III), the audio information belonging to one frame is generally not contained between two successive syncwords.

### 2.4.2.2 Audio frame

header -- Part of the bitstream containing synchronization and state information.
error_check -- Part of the bitstream containing information for error detection.
audio_data -- Fart of the bitstream containing information on the audio samples.
ancillary_data -- Part of the bitstream that may be used for ancillary data.

### 2.4.2.3 Header

The first 32 bits (four bytes) are header information which is common to all layers.

## syncword -- The bit string '1111 1111 1111'.

ID -- One bit to indicate the ID of the algorithm. Equals ' 1 ' for ISO/IEC 11172-3 audio, ' 0 ' is reserved.
Layer -- 2 bits to indicate which layer is used, according to the following.

| Layer |  |
| :---: | :---: |
| '11' | Layer I |
| '10' | Layer II |
| '01' | Layer III |
| '00' | reserved |

To change the layer, a reset of the audio decoder may be required.
protection_bit .- One bit to indicate whether redundancy has been added in the audio bitstream to facilitate error detection and concealment. Equals ' 1 ' if no redundancy has been added, ' 0 ' if redundancy has been added.
bitrate_index -- Indicates the bitrate. The all zero value indicates the 'free format' condition, in which a fixed bitrate which does not need to be in the list can be used. Fixed means that a frame contains either N or $\mathrm{N}+1$ slots, depending on the value of the padding bit. The bitrate_index is an index to a table, which is different for the different layers.

The bitrate_index indicates the total bitrate irrespective of the mode (stereo, joint_stereo, dual_channel, single_channel).

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| bitrate_index | bitrate specified (kbits/s) |  |  |
| :---: | :---: | :---: | :---: |
|  | Layer I | Layer II | Layer III |
| '0000' | free | free | free |
| '0001' | 32 | 32 | 32 |
| '0010' | 64 | 48 | 40 |
| '0011' | 96 | 56 | 48 |
| '0100' | 128 | 64 | 56 |
| '0101' | 160 | 80 | 64 |
| '0110' | 192 | 96 | 80 |
| '0111' | 224 | 112 | 96 |
| '1000' | 256 | 128 | 112 |
| '1001' | 288 | 160 | 128 |
| '1010' | 320 | 192 | 160 |
| '1011' | 352 | 224 | 192 |
| '1100' | 384 | 256 | 224 |
| '1101' | 416 | 320 | 256 |
| '1110' | 448 | 384 | 320 |
| '1111' | forbidden | forbidden | forbidden |

In order to provide the smallest possible delay and complexity, the decoder is not required to support a continuously variable bitrate when in Layer I or II. Layer III supports variable bitrate by switching the bitrate_index. The switching of the bitrate_index can be used either to optimize storage requirements on DSM or to interpolate any mean data rate by switching between nearby values in the bitrate table. However, in free format, fixed bitrate is required. The decoder is also not required to support bitrates higher than $448 \mathrm{kbits} / \mathrm{s}, 384 \mathrm{kbits} / \mathrm{s}, 320 \mathrm{kbits} / \mathrm{s}$ in respect to Layer I, II and III when in free format mode.

For Layer II, not all combinations of total bitrate and mode are allowed. See the following table.

| bit rate (kbits/s) | Allowed modes |
| :---: | :---: |
| free format | all modes |
| 32 | single_channel |
| 48 | single_channel |
| 56 | single_channel. |
| 64 | all modes |
| 80 | single_channel |
| 96 | all modes |
| 112 | all modes |
| 128 | all modes |
| 160 | all modes |
| 192 | all modes |
| 224 | stereo, intensity stereo, dual channel |
| 256 | stereo, intensity stereo, dual channel |
| 320 | stereo, intensity stereo, dual channel |
| 384 | stereo, intensity stereo, dual channel |

sampling_frequency -- Indicates the sampling frequency, according to the following table.

| sampling frequency | frequency specified ( kHz ) |
| :---: | :---: |
| '00' | 44,1 |
| '01' | 48 |
| '10' | 32 |
| '11' | reserved |

A reset of the audio decoder may be required to change the sampling rate.
padding_bit -- If this bit equals ' 1 ', the frame contains an additional slot to adjust the mean bitrate to the sampling frequency, otherwise this bit will be ' 0 '. Padding is necessary with a sampling frequency of $44,1 \mathrm{kHz}$. Padding may also be required in free format.

$$
\begin{aligned}
& \text { Licensed to Venable LLP / Barbara Folensbee-Moore (lib01@venable,com) } \\
& \text { ISO Store Order: OP-149809 / Downloaded: } 2016-06-20 \\
& \text { Single user licence only, copying and networking prohibited. }
\end{aligned}
$$

The padding shisuld be applied to the bitstream such that the accumulated length of the coded frames, after a certain number of audio frames does not deviate more than ( $+0,-1$ slot) from the following computed value:

```
    current frame
accumilated frame length = (frame_size * bitrate/sampling frequency)
                        first frame
where rame_size = 384 for Layer I,
    1 152 for Layer II or III.
```

The following method can be used to determine whether or not to use padding:

```
for 1st audio frame:
    rest = 0;
    padding = no;
for eacis subsequent audio frame:
```

```
    if (Layer == 1) dif = (12 * bitrate) % sampling_frequency;
```

    if (Layer == 1) dif = (12 * bitrate) % sampling_frequency;
    else dif =(144 * bitrate) % sampling_frequency;
    else dif =(144 * bitrate) % sampling_frequency;
    rest = rest - dif;
    rest = rest - dif;
    if (rest < 0) {
    if (rest < 0) {
            padding = yes;
            padding = yes;
            rest = rest + sampling_frequency;
            rest = rest + sampling_frequency;
    }
    }
    else padding = no;
    ```
    else padding = no;
```

private_bit -- 13it for private use. This bit will not be used in the future by ISO/IEC.
mode -- Indicales the mode according to the following table. In Layer I and II the joint_stereo mode is intensity_stereo in Layer III it is intensity_stereo and/or ms_stereo.

| mode | mode specified |
| :---: | :---: |
| ${ }^{\prime} 00^{\prime}$ | stereo |
| '01' | joint_stereo (intensity_stereo and/or ms_stereo) |
| '10' | dual_channel |
| '11' | single_channel |

In Layer I, in all modes except joint stereo, the value of bound equals 32. In layer II, in all modes except joint_stereo, the value of bound equals sblimit. In joint_stereo mode the bound is determined by the mode_extension.
mode_extension -- These bits are used in joint_stereo mode. In Layer I and II they indicate which subbands are in intensity_stereo. All other subbands are coded in stereo.

| node_extension |  |
| :---: | :--- |
| '00' | subbands $4-31$ in intensity_stereo, bound $=4$ |
| '01' | subbands $8-31$ in intensity_stereo, bound $=8$ |
| ' 10 | subbands $12-31$ in intensity_stereo, bound $=12$ |
| '11' | subbands $16-31$ in intensity_stereo, bound $=16$ |

In Layer III they indicate which type of joint stereo coding method is applied. The frequency ranges over which the intensity_stereo and ms_stereo modes are applied are implicit in the algorithm. For more information see 2.4.3.4.

| mode_extension | intensity_stereo | ms_stereo |
| :---: | :---: | :---: |
| $00^{\prime}$ | off | off |
| ${ }^{\prime} 1^{\prime}$ | on | off |
| $' 10^{\prime}$ | off | on |
| $' 11^{\prime}$ | on | on |

Note that the mode "stereo" is used if the mode bits specify stereo or equivalently if the mode bits specify joint stereo and the mode_extension specifies intensity_stereo "off" and ms_stereo "off".
copyright -- If this bit equals ' 0 ', there is no copyright on the ISO/IEC $11172-3$ bitstream, ' 1 ' means copyright protected.
original/copy -- This bit equals ' 0 ' if the bitstream is a copy, ' 1 ' if it is an original.
emphasis -- Indicates the type of de-emphasis that shall be used.

| emphasis | emphasis specified |
| :---: | :---: |
| '00' | none |
| '01' | $50 / 15$ microseconds |
| '10' | reserved |
| '11' | CCITT J.17 |

### 2.4.2.4 Error check

cre_check -- A 16 bit parity-check word is used for optional error detection within the encoded bitstream.

### 2.4.2.5 Audio data, Layer I

allocation[ch][sb] -- Indicates the number of bits used to code the samples in subband sb of channel ch. For subbands in intensity_stereo mode the bitstream contains only one allocation data element per subband.

| allocation $[\mathbf{c h} \\| \mathbf{s b}]$ | bits per sample |
| :---: | :---: |
| 0 | 0 |
| 1 | 2 |
| 2 | 3 |
| 3 | 4 |
| 4 | 5 |
| 5 | 6 |
| 6 | 7 |
| 7 | 8 |
| 8 | 9 |
| 9 | 10 |
| 10 | 11 |
| 11 | 12 |
| 12 | 13 |
| 13 | 14 |
| 14 | 15 |
| 15 | forbidden |

Note: For code ' 0000 ' no samples are transferred.
scalefactor[ch][sb] -- Indicates the factor of subband sb of channel ch by which the requantized samples of subband sb in channel ch shall be multiplied. The six bits constitute an unsigned integer, index to table B. 1 "Layer I, II scalefactors".
sample[ch][sb][s] -- Coded representation of the s-th sample in subband sb of channel ch. For subbands in intensity_stereo mode the coded representation of the sample is valid for both channels.

### 2.4.2.6 Audio data, Layer 11

allocation[ch][sb] -- Contains information concerning the quantizers used for the samples in subband sb in channel ch, whether the information on three consecutive samples has been grouped to one code, and on the number of bits used to code the samples. The meaning and length of this field depends on the number of the subband, the bitrate, and the sampling frequency. The bits in this field form an unsigned integer used as an index to the relevant table in table B. 2 "Layer II bit allocation tables", which gives the number of levels used for quantization. For subbands in intensity_stereo mode the bitstream contains only one allocation data element per subband.
scfsi[ch][sb] -- Scalefactor selection information. This gives information on the number of scalefactors transferred for subband sb in channel ch and for which parts of the signal in this frame they are valid. The frame is divided into three equal parts of 12 subband samples each per subband.

| scfsi[sb] |  |
| :---: | :---: |
| three scalefactors transmitted, for parts $0,1,2$ respectively. |  |
| ' 01 ' | two scalefactors transmitted, first one valid for parts 0 and 1 , second one for part 2. |
| one scalefactor transmitted, valid for all three parts. |  |
| '11' | two scalefactors transmitted, first one valid for part 0 , the second one for parts 1 and 2. |

scalefactor[ch;[sb][p] -- Indicates the factor by which the requantized samples of subband sb in channel ch and of part $p$ of the frame should be multiplied. The six bits constitute an unsigned integer, index to table B. 1 "Layer I, II scalefactors'.
grouping[ch][sb] -- Is a function that determines, whether grouping is in effect for coding of samples in subband sb of ciannel ch . Grouping means, that three consecutive samples (a triplet) of the current subband sb in channel ch in the current granule gr are coded and transmitted using one common codeword and not using three separate codewords. Grouping[ch][sb] is true, if in the Bit Allocation table currently in use (see B.2) the value found under the sb (row) and the allocation[sb] (column) is either 3,5, or 9. Otherwise it is false. For subbands in intensity_stereo mode the grouping is valid for both channels.
samplecode[ch][sb][gr] -- Coded representation of the three consecutive samples in the granule gr in subband sb of channel ch. For subbands in intensity_stereo mode the coded representation of the samplecode is valid for both channels.
sample[ch][sb][s] -- Coded representation of the s-th sample in subband sb of channel ch. For subbands in intensity_stereo mode the coded representation of the sample is valid for both channels.

### 2.4.2.7 Audio data, Layer III

main_data_begin -- The value of main_data_begin is used to determine the location of the first bit of main data of a frame. The main_data_begin value specifies the location as a negative offset in bytes from the first byte of the audio sync word. The number of bytes belonging to the header and side information is not taken into account. For example, if main_data_begin $==0$, then main data starts after the side information. Examples are given in figure A. $\overline{7}$. a and figure A.7.b.
private_bits -- Bits for private use. These bits will not be used in the future by ISO/IEC. The number of private_bits depends on the number of channels. The number of bits allocated for private_bits is determined to equalize the total number of bits used for side-information.
scfsi[ch][scfsi_band] -- In Layer III, the scalefactor selection information works similarly to audio Layer II. The $\pi$ ain difference is the use of the variable scfsi_band to apply scfsi to groups of scalefactors instead of single scalefactors. The application of scalefactors to granules is controlled by scfsi.

| scfsilscfsi band |  |
| :---: | :--- |
| '0' | scalefactors are transmitted for each granule |
| ' 1 ' | scalefactors transmitted for granule 0 are also valid for granule 1 |

If short windows are switched on, i.e. block_type $=2$ for one of the granules, then scfsi is always 0 for this frame.
scfsi_band controls the use of the scalefactor selection information for groups of scalefactors (scfsi_bands).

| scfsi_band | scalefactor bands (see table B.8) |
| :---: | :---: |
| 0 | $0,1,2,3,4,5$ |
| 1 | $6,7,8,9,10$ |
| 2 | $11 \ldots 15$ |
| 3 | $16 \ldots 20$ |

part2_3_length[gr][ch] -- This value contains the number of main_data bits used for scalefactors and Huffman code data. Because the length of the side information is always the same, this value can be used to calculate the beginning of the main information for the next granule or the position of the ancillary information (if used). Note that single channel audio frames contain 17 bytes of side information and dual channel audio frames contain 32 bytes of side information (see 2.4.1.7 Audio Data, Layer III - syntax for audio_data().
big_values[gr][ch] -- The spectral values of each granule are coded with different Huffman code tables. The full frequency range from zero to the Nyquist frequency is divided into several regions, which then are coded using different tables. Partitioning is done according to the maximum quantized values. This is done with the assumption that values at higher frequencies are expected to have lower amplitudes or do not need to be coded at all. Starting at high frequencies, the pairs of quantized values equal to zero are counted. This number is named "rzero". Then, quadruples of quantized values with absolute value not exceeding 1 (i.e. only 3 possible quantization levels) are counted. This number is named "count1". Again an even number of values remain. Finally, the number of pairs of values in the region of the spectrum which extends down to zero is named "big_values". The maximum absolute value in this range is constrained to 8191. The following figure shows the partitioning:

```
xxxxxxxxxxxxx------------------000000000000000000000000000000
1 1 1 |
1 bigvalues*2 bigvalues*2+count1*4 iblen
```

The values 000 are all zero. Their number is a multiple of 2 .
The values -- are $-1,0$ or +1 . Their number is a multiple of 4 .
The values xxx are not bound. Their number is a multiple of 2 .
Iblen is 576 .
global gain[gr][ch] -- The quantizer step size information is transmitted in the side information variable global_gain. It is logarithmically quantized. For the application of global_gain, refer to the formula in 2.4.3.4, "Formula for requantization and all scaling".
scalefac_compress[gr][ch] .- Selects the number of bits used for the transmission of the scalefactors according to the following table:
if block_type is 0,1 , or 3 :
slen 1: length of scalefactors for the scalefactor bands 0 to 10
slen2: length of scalefactors for the scalefactor bands 11 to 20
if block_type is 2 and mixed_block_flag is 0:
slen 1: length of scalefactors for the scalefactor bands 0 to 5
slen2: length of scalefactors for the scalefactor bands 6 to 11
if block_type is 2 and mixed_block_flag is 1 :
slen 1: length of scalefactors for the scalefactor bands 0 to 7 (long window scalefactor band) and 3 to 5 (short window scalefactor band) Note: Scalefactor bands 0-7 are from the "long window
scalefactor band" table, and scalefactor bands 3 to 11 from the "short window scalefactor band" table. This combination of partitions is contiguous and spans the entire frequency spectrum. slen2: length of scalefactors for the scalefactor bands 6 to 11

| scalefac_compress[gr] | slen 1 | slen2 |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 1 | 0 | 1 |
| 2 | 0 | 2 |
| 3 | 0 | 3 |
| 4 | 3 | 0 |
| 5 | 1 | 1 |
| 6 | 1 | 2 |
| 7 | 2 | 3 |
| 8 | 2 | 1 |
| 9 | 2 | 2 |
| 10 | 3 | 3 |
| 11 | 3 | 1 |
| 12 | 4 | 2 |
| 13 | 4 | 3 |
| 14 | 2 | 3 |

window_switching_flag[gr][ch] -- Signals that the block uses an other than normal (type 0) window.

If window_switshing_flag is set, several other variables are set by default:

```
region0_count = 7 (in case of block_type==1 or block_type==3
    or block_type==2 and mixed_block_flag)
region0_count = 8 (in case of block_type==2 and not mixed_block_flag)
region1_count = 36 Thus all remaining values in the big_value region are contained in
                    region 1.
```

If window_switching_flag is not set, then the value of block_type is zero.
block_type[gr][ch] -- Indicates the window type for the granule (see description of the filterbank, Layer III).

| block_type[gr] |  |
| :---: | :---: |
| 0 | reserved |
| 1 | start block |
| 2 | 3 short windows |
| 3 | end block |

Block_type and mixed_block_flag give the information about assembling of values in the block and about length and count of the transforms (see figure A. 4 for a schematic, annex C for an analytic description). If vindow_switching_flag==1, then the mixed_block_flag indicates whether lower frequency polyphase filter subbands are coded using normal window type. The polyphase filterbank is described in 2.4.3.

In the case of Icng blocks (block_type not equal to 2 or in the lower subbands of block_type 2 if the mixed_block_flag is set) the IMDCT generates an output of 36 values every 18 input values. The output is windowed depeading on the block_type and the first half is overlapped with the second half of the block before. The resulting vector is the input of the synthesis part of the polyphase filterbank of one band.

In the case of short blocks (in the upper subbands of a type 2 block if the mixed_block_flag is set, or in all subbands of a type 2 block if mixed_block_flag is not set), three transforns are performed producing 12 output values etch. The three vectors are windowed and overlapped each. Concatenating 6 zeros on both ends of the resulting vector gives a vector of length 36 , which is processed like the output of a long transform.
mixed_block_flag[gr]ich] -- Indicates that lower frequencies are transformed with a window type that is different than that which is used at higher frequencies. If mixed_block_flag is zero, then all blocks are transformed as indicated by block_type[gr][ch]. If mixed_block_flag is one, then the frequency lines corresponding to the two lowest frequency polyphase subbands are transformed with normal window (block_type=0), while the remaining 30 subbands are transformed as block_type[gr][ch].
table_select[gr][ch][region] -- Different Huffman code tables are used depending on the maximum quantized value and the local statistics of the signal. There are a total of 32 possible tables given in table B. 7 .
subblock_gain[gr][ch][window] -- Indicates the gain offset (quantization: factor 4) from the global gain for one subblock. Used only with block type 2 (short windows). The values of the subblock have to be divided by $4^{\text {subblock_gain[window]) }}$ in the decoder. See 2.4.3.4 - Formula for requantization and all scaling.
region0_count[gr][ch] -- A further partitioning of the spectrum is used to enhance the performance of the Huffman coder. It is a subdivision of the region which is described by big_values. The purpose of this subdivision is to get better error robustness and better coding efficiency. Three regions are used, they are named: region 0,1 and 2 . Each region is coded using a different Huffman code table depending on the maximum quantized value and the local signal statistics.

The values region0_count and region1_count are used to indicate the boundaries of the regions. The region boundaries are aligned with the partitioning of the spectrum into scale factor bands.

The field region0_count contains one less than the number of scalefactor bands in region 0 . In the case of short blocks, each scale factor band is counted three times, once for each short window, so that a region0_count value of 8 indicates that region 1 begins at scalefactor band number 3 .

If block_type $=2$ and mixed_block_flag $==0$, the total amount of scalefactor bands for the granule in this case is $12 * 3=36$. If block_type $=2$ and mixed_block_flag $==1$, the amount of scalefactor bands is $8+9 * 3=35$. If block_type! $=2$, the amount of scalefactor bands is 21 .
region1_count[gr][ch] -- region1_count counts one less than the number of scalefactor bands in region 1. Again, if block_type $==2$ the scalefactor bands representing different time slots are counted separately.
preflag[gr][ch] -- This is a shortcut for additional high frequency amplification of the quantized values. If preflag is set, the values of a table are added to the scalefactors (see table B.6). This is equivalent to multiplication of the requantized scalefactors with table values. If block_type=m 2 (short blocks) preflag is never used.
scalefac_scale[gr][ch] -- The scalefactors are logarithmically quantized with a step size of 2 or ( $\sqrt{2}$ ) depending on scalefac_scale. The following table indicates the scale factor multiplier used in the requantization equation for each stepsize.

| scalefac_scale[gr] | scalefac_multiplier |
| :---: | :---: |
| 0 | 0,5 |
| 1 | 1 |

count1table_select[gr][ch] .- This flag selects one of two possible Huffman code tables for the region of quadruples of quantized values with magnitude not exceeding 1 .

| count1table_select[gr] |  |
| :---: | :--- |
| 0 | Table B.7-A |
| 1 | Table B.7-B |

scalefac_l[gr][ch][sfb], scalefac_s[gr][ch][sfb][window], is_pos[sfb] -- The scalefactors are used to colour the quantization noise. If the quantization noise is coloured with the right shape, it is masked completely. Unlike Layers I and II, the Layer III scalefactors say nothing about the local maximum of the quantized signal. In Layer III, scalefactors are used in the decoder to get division factors for groups of values. In the case of Layer III, the groups stretch over several frequency lines. These groups are called scalefactor bands and are selected to resemble critical bands as closely as possible.

The scalefac_ciompress table shows that the scalefactors $0 \ldots 10$ have a range of 0 to 15 (maximum length 4 bits) and the scalefactors $11 \ldots 21$ have a range of 0 to 7 (maximum length 3 bits).

If intensity_stereo is enabled (modebit_extension) the scalefactors of the "zero_part" of the difference (right) channel are used as intensity_stereo positions, is_pos[sfb] (see 2.4.3.4, MS_stereo mode). is_pos[sfb] is the intensity stereo position for scalefactor band sfb .

The subdivision of the spectrum into scalefactor bands is fixed for every block length and sampling frequency and stored in tables in the coder and decoder (see table B.8). The scale factor for frequency lines above the highest line in the tables is zero, which means that the actual multiplication factor is 1,0 .

The scalefactor: are logarithmically quantized. The quantization step is set with scalefac_scale.
huffmancodebits0 -- Huffman encoded data.
The syntax for huffmancodebits 0 shows how quantized values are encoded. Within the big_values partition, pairs of quantized values with an absolute value less than 15 are directly coded using a Huffman code. The codes are selected from huffman tables 0 through 31 in table B.7. Always pairs of values ( $x, y$ ) are coded. If quantized value; of magnitude greater than or equal to 15 are coded, the values are coded with a separate field following the huffman code. If one or both values of a pair is not zero, one or two sign bits are appended to the code word.

The Huffman tables for the big_values partition are comprised of three parameters:

| hcod $[\|x\|][\mid y]]$ | is the Huffman code table entry for values $x, y$. |
| :--- | :--- |
| hlen $[\|x\|][\|y\|]$ | is the Huffman length table entry for values $x, y$. |
| linbits | is the length of linbits or linbitsy when they are coded. |

The syntax for tiuffmancodebits contains the following fields and parameters:

| signv | is the sign of $v$ ( $(0$ if positive, 1 if negative). |
| :--- | :--- |
| $\operatorname{signv}$ | is the sign of $w$ (0 if positive, 1 if negative). |
| $\operatorname{signx}$ | is the sign of $x$ (0 if positive, 1 if negative). |
| $\operatorname{signy}$ | is the sign of $y$ ( 0 if positive, 1 if negative). |
| linbilsx | is used to encode the value of $x$ if the magnitude of $x$ is greater or |
|  | equal to 15 . This field is coded only if $\|x\|$ in hcod is equal to 15. If |
|  | linbits is zero, so that no bits are actually coded when $\|x\|==15$, then |
| the value linbitsx is defined to be zero. |  |
| linbilsy | is the same as linbitsx but for $y$. |
| is $[1]$ | Is the quantized value for frequency line number 1. |

The linbitsx or linbitsy fields are only used if a value greater or equal to 15 needs to be encoded. These fields are interpreted as unsigned integers and added to 15 to obtain the encoded value. The linbitsx and linbitsy fields are never used if the selected table is one for blocks with a maximum quantized value less than 15 . Note taat a value of 15 can still be encoded with a huffman table for which linbits is zero. In this case, the linbits \& or linbitsy fields are not actually coded, since linbits is zero.

Within the count1 partition, quadruples of values with magnatude less than or equal to one are coded. Again magnitude values are coded using a Huffman code from tables A or B in table B.7. Again, for each non-zero value, a sign bit is appended after the huffman code symbol.

The Huffman tables for the count1 partition are comprised of the following parameters:

$$
\begin{array}{ll}
\text { hcod }[|v|][|w|][|x|][|y|] & \text { is the Huffman code table entry for values } v, w, x, y . \\
\text { hlen }[|v|]||w|][|x|][|y|] & \text { is the Huffman length table entry for values } v, w, x, y .
\end{array}
$$

Huffman code table B is not really a 4-dimensional code because it is constructed from the trivial code: 0 is coded with a 1 , and 1 is coded with a 0 .

Quantized values above the count 1 partition are zero, so they are not encoded.
For clarity, the parameter "count1" is used in this document to indicate the number of Huffman codes in the countl region. However, unlike the bigvalues partition, the number of values in the countl partition is not
explicitly coded by a field in the syntax. The end of the count 1 partition is known only when all bits for the granule (as specified by part2_3_length), have been exhausted, and the value of count 1 is known implicitly after decoding the count I region.

The order of the Huffman data depends on the block_type of the granule. If block_type is 0,1 or 3 the Huffman encoded data is ordered in terms of increasing frequency.

If block_type $=2$ (short blocks) the Huffman encoded data is ordered in the same order as the scalefactor values for that granule. The Huffman encoded data is given for successive scalefactor bands, beginning with scalefactor band 0 . Within each scalefactor band, the data is given for successive time windows, beginning with window 0 and ending with window 2 . Within each window, the quantized values are then arranged in order of increasing frequency.

### 2.4.2.8 Ancillary data

Ancillary_bit -- User definable.
The number of ancillary bits (no_of_ancillary_bits) equals the available number of bits in an audio frame minus the number of bits actually used for header, error check and audio data. In Layer I and II the no_of_ancillary_bits corresponds to the distance between the end of the audio data and the beginning of the next header. In Layer III the no_of_ancillary_bits corresponds to the distance between the end of the Huffman_code_bits and the location in the bitstream where the next frame's main_data_begin pointer points to.

### 2.4.3 The audio decoding process

### 2.4.3.1 General

The first action is synchronization of the decoder to the incoming bitstream. Just after startup this may be done by searching in the bitstream for the 12 bit syncword. In some applications the ID , layer, and protection status are already known to the decoder, and thus the first 16 bits of the header should be regarded as a 16 bit syncword, thereby allowing a more reliable synchronization. The position of consecutive syncwords can 'ee calculated from the information provided by the seven bits after the protection_bit: the bitstrearn is subdivided in slots. The distance between the start of two consecutive syncwords is equal to " N " or " $\mathrm{N}+1$ " slots. The value of " N " depends on the layer.

For Layer I the following equation is valid:

$$
\mathrm{N}=12 * \frac{\text { bitrate }}{\text { sampling_frequency }}
$$

For Layers II and III the equation becomes:

$$
\mathrm{N}=1<4 * \frac{\text { bitrate }}{\text { sampling_frequency }}
$$

If this calculation does not give an integer number the result is truncated and 'padding' is required. In this case the number: of slots in a frame will vary between N and $\mathrm{N}+1$. The padding bit is set to ' 0 ' if the number of slots equals N , and to ' 1 ' otherwise. This knowledge of the position of consecutive syncwords greatly facilitates symchronization.

If the bitrate index equals ' 0000 ', the exact bitrate is not indicated. N can be determined from the distance between consecutive syncwords and the value of the padding bit.

The mode bits in the bitstream shall be read and if their value is ' 01 ' the mode_extension bits shall also be read. The mode_extension bits set the 'bound' as shown in 2.4.2.3 and thus indicate which subbands are coded in joint_stereo mode.

If the protection bit in the header equals ' 0 ', a CRC-check word has been inserted in the bitstream just after the header. The error detection method used is 'CRC-16' whose generator polynomial is:

$$
\mathrm{G}(\mathrm{X})=\mathrm{X}^{16}+\mathrm{X}^{15}+\mathrm{X}^{2}+1
$$

The bits included into the CRC-check are given by table B.5.
The method is cepicted in figure A. 9 "CRC-check diagram". The initial state of the shift register is '1111 $111111111111^{\prime}$. Then all the bits included into the CRC-check are input to the circuit shown in figure A. 9 "CRC-check diagram". After each bit is input the shift register is shifted by one bit. After the last shift operation, the outputs $\mathrm{b}_{15} \ldots \mathrm{~b}_{0}$ constitute a word to be compared with the CRC-check word in the bitstream. If the words are not identical, a transmission error has occurred in the protected field of the bitstream. To avoid annoying distortions, application of a concealment technique, such as muting of the acual frame or repetition of the previous frame, is recommended.

### 2.4.3.2 Layer I

After the part of the decoding which is common to all layers (see 2.4.3.1) the bit allocation information has to be read for all subbands, and the scalefactors read for all subbands with a nonzero bit allocation. The decoder flowchart is given in figure A. 1 "Layer I and II decoder flow chart".

### 2.4.3.2.1 Requantization of subband samples

From the bit allocation the number of bits nb that has to be read for the samples in each subband is known. The order of the samples is given in 2.4.1.5 for each mode. After the bits for one sample have been gathered from the bitstream, the first bit has to be inverted. The resulting number can be considered as a two's
complement fractional number, where the MSB represents the value -1 . The requantized value can be obtained by applying a linear formula :

$$
s^{\prime \prime}=\frac{2^{n b}}{2^{n b}-1} *\left(s^{\prime \prime \prime}+2^{-n b+1}\right)
$$

where

| $s^{\prime \prime \prime}$ | is the fractional number; |
| :--- | :--- |
| $\mathrm{s}^{\prime \prime}$ | is the requantized value; |
| nb | is the number of bits allocated to samples in the subband. |

Samples in subbands which are in intensity_stereo mode must be copied to both channels. The requantized value has to be rescaled. The multiplication factor can be found in the table B. 1 "Layer I, II scalefactors". The rescaled value s' is calculated as :

$$
s^{\prime}=\text { factor }{ }^{*} s^{\prime \prime}
$$

### 2.4.3.2.2 Synthesis subband filter

If a subband has no bits allocated to it, the samples in that subband are set to zero. Each time the subband samples for all 32 subbands of one channel have been calculated, they can be applied to the synthesis subband filter and 32 consecutive audio samples can be calculated. The actions in flow diagram figure A. 2 "Synthesis subband filter flow chart" show the reconstruction operation. The coefficients $\mathrm{N}_{\mathrm{ik}}$ for the neatrixing operation are given by

$$
N_{i k}=\cos \left[(16+i)(2 k+1) \frac{\pi}{64}\right] \quad 0 \leq i \leq 63,0 \leq k \leq 31
$$

The coefficients $D_{i}$ for the windowing operation can be found in table B. 3 "Coefficients $D_{i}$ of the synthesis v/indow". The coefficients have been derived by numerical optimization. One frame contains $12 * 32=384$ subband samples, which result, after filtering, in 384 audio samples.

### 2.4.3.3 Layer II

Layer II is a more efficient but more complex coding scheme than Layer I. The flowchart in figure A. 1 "Layer I and II decoder flow chart" applies to both Layers I and II. The first step is to perform the decoding which is common to all three layers (see 2.4.3.1).

### 2.4.3.3.1 Bit allocation decoding

For different combinations of bitrate and sampling frequency different bit allocation tables exist (table B. 2 "Layer II bit allocation tables"). Note that the bitrates given in the table headers are per channel. If the mode is not single_channel, the bitrate should be divided by two to obtain the bitrate per channel. The decoding of the bit allocation table is done in a three-step approach. The first step consists of reading 'nbal' ( 2,3, or 4 ) bits of information for one subband from the bitstream. The value of 'nbal' is given in the second column of the relevant table B. 2 "Layer II bit allocation tables". These bits shall be interpreted as an unsigned integer number. The second step uses this number and the number of the subband as indices to point to a value in the table. This value represents the number of levels 'nlevels' used to quantize the samples in the subband. As a third step, using table B. 4 "Layer II classes of quantization", the number of bits used to code the quantized samples, the requantization coefficients, and whether the codes for three consecutive subband samples have been grouped to one code can be determined. It can be seen from the bit allocation tables that some of the highest subbands will never have bits allocated. The number of the lowest subband that wil not have bits allocated to it is assigned to the identifier 'sblimit'.

### 2.4.3.3.2 Scalefactor selection information decoding

The 36 samples in one subband within a frame are divided in three equal parts of 12 subband samples. Each part can have its own scalefactor. The number of scalefactors that has to be read from the bitstream depends on scfsi[sb]. The scalefactor selection information scfsi[sb] is read from the bitstream for the subbands that have a nonzero bit allocation. If scfsi[sb] equals ' 00 ' three scalefactors are transmitted, for parts $0,1,2$ respectively. If scfsi[sb] equals '01' two scalefactors are transmitted, the first one valid for parts 0 and 1 , the
second one for part 2. If scfsi[sb] equals ' 10 ' one scalefactor is transmitted, valid for all three parts. If scfsi[sb] equals '11' two scalefactors are transmitted, the first one valid for part 0 , the second one for parts 1 and 2.

### 2.4.3.3.3 Scalefactor decoding

For every subbend with a nonzero bit allocation the coded scalefactors for that subband are read from the bitstream. The number of coded scalefactors and the part of the subband samples they refer to is defined by $\mathrm{scfsi}[\mathrm{sb}]$. The 6 bits of a coded scalefactor should be interpreted as an unsigned integer index to table B. 1 "Layer I, II scalefactors". This table gives the scalefactor by which the relevant subband samples should be multiplied after requantization.

### 2.4.3.3.4 Requantization of subband samples

Next the coded samples are read. As can be seen from 2.4.1.6, the coded samples appear as triplets, the code contains three consecutive samples at a time. From table B. 4 "Layer II classes of quantization" it is known how many bits save to be read for one triplet from the bitstream for each subband. Also from table B. 4 "Layer II classes of quantization", it is known whether this code consists of three consecutive separable codes for each sample or of one combined code for the three samples (grouping). In the last case degrouping must be performed. The combined code has to be regarded as an unsigned integer, called ' $c$ '. The following algorithm will supply the three separate codes $s[0], s[1], s[2]$.

```
for (i=3; i<3;i++) {
    s[i]= c % nlevels
    c = c DIV nlevels
}
```

where nlevels is: the number of steps as shown in table B. 2 "Layer II bit allocation table".
The first bit of $\epsilon$ ach of the three codes has to be inverted, and the resulting numbers should be regarded as two's complement fractional numbers, where the MSB represents the value -1 . The requantized values can be obtained by applying a linear formula:

$$
s^{\prime \prime}=C *\left(s^{\prime \prime \prime}+D\right)
$$

where

| $s^{\prime \prime \prime}$ | is the fractional number; |
| :--- | :--- |
| $s^{\prime \prime}$ | is the requantized value. |

The values of the constants C and D are given in table B. 4 "Layer II classes of quantization". The requantized values have to be rescaled. The multiplication factors can be found in the table B. 1 "Layer I, II scalefactors". as described above. The rescaled value s' is calculated as :

$$
s^{\prime}=\text { factor } * s^{\prime \prime}
$$

### 2.4.3.3.5 Synthesis subband filter

If a subband ha; no bits allocated to it, the samples in that subband are set to zero. Each time the subband samples for all 32 subbands of one channel have been calculated, they can be applied to the synthesis subband filter and 32 consecutive audio samples can be calculated. For that purpose, the actions in the flow diagram of figure A. 2 "Synthesis subband filter flow chart" have to be carried out. The coefficients $\mathrm{N}_{\mathrm{ik}}$ for the matrixing operation are given by

$$
N_{i k}=\cos \left[(16+i)(2 k+1) \frac{\pi}{64}\right] \quad 0 \leq i \leq 63,0 \leq k \leq 31
$$

The coefficient: $D_{i}$ for the windowing operation can be found in table $B .3$ "Coefficients $D_{i}$ of the synthesis window". One frame contains $36 * 32=1152$ subband samples, which result after filtering in 1152 audio samples.

### 2.4.3.4 Layer III

Additional frequency resolution is provided by the use of an hybrid filterbank. Every band is split into 18 frequency lines by use of an MDCT. The window length of the MDCT is 36. Adaptive window switching is used to control time artifacts (pre-echoes), see the description in annex C. The frequency above whicia shorter blocks (better time resolution) are used can be selected. Parts of the signal below a frequency depending on "mixed_block_flag" are coded with better frequency resolution, parts of the signal above are coded with better time resolution.

The frequency components are quantized using a nonuniform quantizer and coded using a Huffman encoder. The Huffman coder uses one of 18 different tables (see table B.7). A buffer is used to help enhance the coding efficiency of the Huffman coder and to help in the case of pre-echo conditions (see the description in annex C). The size of the input buffer is the size of one frame at the bitrate of $160 \mathrm{kbits} / \mathrm{s}$ per channel for Layer III. The short term buffer technique used is called "bit reservoir" because it uses short-term variable bitrate with a maximum integral offset from the mean bitrate.
liach frame holds the data from 2 granules. The audio data in a frame is allocated in the following way:

- main_data_begin pointer
- side info for both granules (scfsi)
- side info granule 1
- side info granule 2

The header and this part of the audio data constitute the side information stream.

- scalefactors and Huffiman code data granule 1
- scalefactors and Huffman code data granule 2
- ancillary data

These data constitute the main data stream. The main_data_begin pointer specifies a negative offset from the position of the first byte of the header.

### 2.4.3.4.1 Decoding

The first action is the synchronization of the decoder to the incoming bitstream. This is done as in the other layers. The header information (first 32 bits including syncword) is read in just as in the other layers. The information about sampling frequency is used to select the scalefactor_band table (see annex B.8).

### 2.4.3.4.2 Side information

The side information must be extracted from the bitstream and stored for use during the decoding of the associated frame. The table select information is used to select the Huffman decoder table and the number of ESC-bits (linbits), according to table B.7.

### 2.4.3.4.3 Start of main_data

The main_data (scalefactors, Huffman coded data and ancillary information) are not necessarily located adjacent to the side information. This is described in figure A.7.a and figure A.7.b. The beginning of the main data part is located by using the main_data_begin pointer of the current frame. The allocation of the main data is done in a way that all main data are resident in the input buffer when the header of the next flame is arriving in the input buffer. The decoder has to skip Header and side information when decoding he main data. It knows their positions from the bitrate_index and padding_bit. The length of the Header is always 4 bytes, the length of the side information is 17 bytes in mode single_channel and 32 bytes in the other modes. Main data can span more than one block of header and side information (see figure A.7.b).

### 2.4.3.4.4 Buffer considerations

The following rule can be used to calculate the maximum number of bits used for one granule:
The buffer length is 7680 bits. This value is used as the maximum buffer at every bitrate. At the highest possible bitrate of Layer III ( $320 \mathrm{kbits} / \mathrm{s}$ per stereo signal) and sampling frequency 48 kHz the mean frame length is $(320000 / 48000)^{*} 1152=7680$ bits. Therefore the frames must be of constant length at this bitrate and sampling frequency. At $64 \mathrm{kbits} / \mathrm{s}(128 \mathrm{kbits} / \mathrm{s}$ stereo) the mean granule length is
( $64000 / 48000 ; 576=768$ bit at 48 kHz sampling frequency. This means that there is a maximum deviation (short time buffer) of $7680-4 * 768=4608$ bits is allowed at $64 \mathrm{kbits} / \mathrm{s}$. The actual deviation is equal to the number of bytes denoted by the main_data_begin offset pointer. The actual maximum deviation is $2 * * 9 * 8$ bit $=4096$ bits. For intermediate bitrates the delay and buffer length can be calculated accordingly. The exchange of buffer between the left and right channel in a stereo bitstream is allowed without restrictions. Because of the constraint on the buffer size main_data_begin is always set to 0 in the case of bitrate_i 1 de $x==14$, i.e. data rate $320 \mathrm{kbits} / \mathrm{s}$ per stereo signal. In this case all data are allocated between adjacen: header words.

At sampling frequencies lower than 48 kHz the buffer should be constrained such that the same physical buffer size is sulficient as the one calculated for the 48 kHz case above.

### 2.4.3.4.5 Scalefactors

The scalefactors are decoded according to the slen 1 and slen 2 which themselves are determined from the values of scalefac_compress. The decoded values can be used as entries into a table or used to calculate the factors for each :icalefactor band directly. When decoding the second granule, the scfsi has to be considered. For the bands in which the corresponding scfsi is set to 1 , the scalefactors of the first granule are also used for the second granule, therefore they are not transmitted for the second granule.

The number of bits used to encode scalefactors is called part2_length, and is calculated as follows.

$$
\begin{aligned}
& \text { For blozk_type }=x=0,1 \text {, or } 3 \text { (long blocks): } \\
& \text { part } 2 \text { _length }=11 * \text { slen } 1+10 * \text { slen } 2 . \\
& \text { For block_type }==2(\text { short blocks) and mixed_block_flag }=0 \text { : } \\
& \text { part2_length }=18 * \text { slen } 1+18 * \text { slen } 2 \text {. }
\end{aligned}
$$

$$
\text { For block_type }==2 \text { (short blocks) and mixed_block_flag }=1:
$$

$$
\text { part } 2 \text { _ length }=17 * \text { slen } 1+18 * \text { slen } 2 \text {. }
$$

These formulas arre valid if $\mathrm{gr}==0$ or if $\mathrm{gr}==1$ and scfsi[ch][scfsi_band] $==0$ for all scfsi_bands, i.e. scalefactor selection information is not used.

### 2.4.3.4.6 Huffman decoding

All necessary information including the table which realizes the Huffman code tree can be generated from the tables in table B.7. First the big_values data are decoded, using the tables with the number table_select[gr][ich][region]. The frequency lines in region 0, region 1 and region 2 are Huffman decoded in pairs until big_values number of line-pairs have been decoded. The remaining Huffmancodebits are decoded using the table according to count1table_select[gr][ch]. Decoding is done until all Huffman code bits have been decoded or until quantized values representing 576 frequency lines have been decoded, whichever comes first. If there are more Huffman code bits than necessary to decode 576 values they are regarded as stuffing bits and discarded. The variable countl is implicitely derived as the number of quadruples of decoded values using count1table_select.

### 2.4.3.4.7 Requantizer

The nonuniform quantizer uses a power law. For each output value, "is", from the Huffman decoder, "lis $\left.\right|^{4 / 3 "}$ is calculated. This can be done either by table lookup or by explicit calculation.

### 2.4.3.4.7.1 Formula for requantization and all scaling

One complete formula describes all the processing from the Huffman decoded values to the input of the synthesis filterbank. All necessary scaling factors are contained within this formula. The output data are reconstructed frcm requantized samples. Global gain and subblock gain values affect all values within one time window (in the case of block_type=2). Scalefactors and preflag further adjust the gain within each scalefactor band. An illustration can be found in figure A.8.

The following is the requantization equation for short windows. The Huffman decoded value at buffer index $i$ is called $i_{i}$, the input to the synthesis filterbank at index $i$ is called $x r_{\mathrm{j}}$ :

$$
\begin{aligned}
\mathrm{xr}_{\mathrm{i}} & =\operatorname{sign}\left(\mathrm{is}_{\mathrm{i}}\right) *\left|i s_{\mathrm{i}}\right|^{\frac{4}{3}} * 2^{\frac{1}{4}(\text { global_gain[gr]-210-8* subblock_gain[window][gr]) }} \\
& * 2^{-(\mathrm{scalefac}} \text { _multiplier*scalefac_s[gr][ch][sfb][window])}
\end{aligned}
$$

For long blocks, the formula is:

$$
\begin{aligned}
\mathrm{xr}_{\mathrm{i}} & =\operatorname{sign}\left(\mathrm{is}_{\mathrm{i}}\right) *\left|\mathrm{is}_{\mathrm{i}}\right|^{\frac{4}{3}} * 2^{\frac{1}{4}(\text { global_gain }[\mathrm{gr}]-210)} \\
& \left.* 2^{-(\mathrm{scalefac}} \mathrm{g}_{-} \text {multiplier*}(\text { scalefac_ } 1[\mathrm{sfb}][\mathrm{ch}][g r]+\text { preflag }[g r] * \text { pretab }[\mathrm{sfb}])\right)
\end{aligned}
$$

Pretab[sfb] is a value given in the preemphasis table B.6. The constant 210 in this formula is needed tc scale the output appropriately. It is a system constant. The synthesis filterbank is assumed to be implemented according to the formulas below. The range of the output values of the decoder (PCM samples) is between $-1,0$ and $+1,0$.

### 2.4.3.4.8 Reordering

If short blocks are used (block_type=-2), the rescaled data $\operatorname{xr[scf}$ _band][window][freq_line] (as described in huffmancodebits( in 2.4.1.7) shall be reordered in subband order, xr[subband[window][freq_line], prior to, the IMDCT operation.

### 2.4.3.4.9 Stereo Processing

A.fter requantization, the reconstructed values are processed for MS or intensity stereo modes or both, before going to the synthesis filterbank. In MS_stereo mode, both channels of a granule must have the same block_type.

### 2.4.3.4.9.1 MS_stereo mode

This mode switch (found in the header: mode_extension) allows switching from "independent stereo" to MS_stereo. If MS_stereo is enabled but intensity stereo is not enabled the entire spectrum is decoded in MS_stereo. If both MS_Stereo and intensity stereo are enabled, the upper bound of the scalefactor bands decoded in MS_stereo is derived from the "zero_part" of the difference (right) channel. In this case the sicalefactor band in which the last non-zero (right channel) frequency line occurs is the last scalefactor band to which the MS_stereo equations apply. Above this bound intensity stereo may be applied if enabled in the header. The "zero_part" of the difference channel is the part of the spectrum from
"'jigvalues ${ }^{*} 2+$ count1 * 4 " (see 2.4 .2 .7 ) to the Nyquist rate.

### 2.4.3.4.9.2 MS matrix

In MS stereo mode the values of the normalized middle/side channels $\mathrm{M}_{\mathrm{i}} / \mathrm{S}_{\mathrm{i}}$ are transmitted instead of the left/right channel values $\mathrm{L}_{\mathrm{i}} / \mathrm{R}_{\mathbf{i}}$. Thus $\mathrm{L}_{\mathbf{i}} / \mathrm{R}_{\mathrm{i}}$ are reconstructed using

$$
L_{i}=\frac{M_{i}+S_{i}}{\sqrt{2}} \text { and } R_{i}=\frac{M_{i}-S_{i}}{\sqrt{2}} .
$$

The values $M_{i}$ are transmitted in the left, values $S_{i}$ are transmitted in the right channel.
If window switching occurs, then the M and S channels must switch synchronously.

### 2.4.3.4.9.3 Intensity stereo mode

This mode switch (found in the header: mode_extension) allows switching from 'normal stereo' to intensity stereo. In Layer III, intenisty stereo is not done using a pair of scalefactor s as in Layers I and II, but by specifying the magnitude (via the scalefactors of the right channel as normal) and a stereo position is_pos $s_{s b}[s f b]$. is_pos ${ }_{s b}[\mathrm{sfb}]$ is transmitted instead of scalefactors for the right channels. The stereo position is used to derive the left and right channel signals according to the formulas below. The lower bound of the scalefactor bands decoded in intensity stereo is derived from the "zero_part" of the right

[^0]channel. Above this bound decoding of intensity stereo is applied using the scalefactors of the right channel as intensity stereo positions. An intensity stereo position of 7 in one scalefactor band indicates that this scalefactor band is not decoded as intensity stereo.


For each scalefactor band (sb) coded in intensity stereo, the following steps are executed:

1) the intensity stereo position is_possb is read from the scalefactor of the right channel.
2) if (iis_pos ${ }_{s b}==7$ ) do not perform the following steps (illegal is_pos).
3) is_ catio $=\tan \left(\right.$ is_ $\left._{-} \operatorname{pos}_{s b} * \frac{\pi}{12}\right)$.
4) $L_{i}:=L_{i} * \frac{i s_{1} \text { ratio }}{1+i s_{-} \text {ratio }}$ for all indices $i$ within the actual scalefactor band $s b$.
5) $R_{i}:=L_{i} * \frac{1}{1+i s_{-} \text {ratio }}$ for all indices $i$ within the actual scalefactor band $s b$.

### 2.4.3.4.10 Synthesis filterbank

Figure A.4. shows a block diagram including the synthesis filterbank. The frequency lines are preprocessed by the "alias recluction" scheme (see the block diagrams in in figure A.5 and in table B.9, for the coefficients) and fed into the IMDCT matrix, each 18 into one transform block. The first half of the output values are addel to the stored overlap values from the last block. These values are new output values and are input values for the polyphase filterbank. The second half of the output values is stored for overlap with the next data granule. For every second subband of the polyphase filterbank every second input value is multiplied by -1 to correct for the frequency inversion of the polyphase filterbank.

### 2.4.3.4.10.1 Alias reduction

For long block_type granules (block_type $!=2$ ) the input to the synthesis filterbank is processed for alias reduction before processing by the IMDCT. The following pseudo code describes the alias reduction computation:

```
for (sb=1; sb<32; sb++)
    for' (i=0; i<8; i++) {
        xar[18*sb-1-i] = xr[18*sb-1-i]Cs[i] - xr[18*sb+i]Ca[i]
        xar[18*sb+i] =xr[18*sb+i]Cs[i] +xr[18*sb-1-i]Ca[i]
    }
```

The indices of arrays $\times \operatorname{ar}[]$ and $\times r[]$ label the frequency lines in a granule, arranged in order of lowest frequency to highest frequency, with zero being the index of the lowest frequency line, and 575 being the index of the highest. The coefficients: $\mathrm{Cs}[\mathrm{i}]$ and $\mathrm{Ca}[\mathrm{i}]$ can be found in table B.9. Figures A. 5 and A. 6 illustrate the alias reduction computation.

Alias reduction is not applied for granules with block_type $==2$ (short block).

### 2.4.3.4.10.2 IMDCT

In the following, n is the number of windowed samples (for short blocks n is 12 , for long blocks n is 36 ). In the case of a block of type "short", each of the three short blocks is transformed separately. n/2 values $\mathrm{X}_{\mathrm{k}}$ are transformed to n values $\mathrm{x}_{\mathrm{i}}$. The analytical expression of the IMDCT is:

$$
x_{i}=\sum_{i=0}^{\frac{n}{i}-1} X_{k} \cos \left(\frac{\pi}{2 n}\left(2 i+1+\frac{\pi}{2}\right)(2 k+1)\right) \quad \text { for } \mathrm{i}=0 \text { to } \mathrm{n}-1
$$

### 2.4.3.4.10.3 Windowing

Depending on the block_type different shapes of windows are used.

$$
\begin{aligned}
& \text { a) block_type }=0 \text { (normal window) } \\
& z_{i}=x_{i} \sin \left(\frac{\pi}{36}\left(i+\frac{1}{2}\right)\right) \quad \text { for } \mathrm{i}=0 \text { to } 35 \\
& \text { b) } \quad \text { block_type }=I \text { (start block) } \\
& \mathrm{z}_{\mathrm{i}}= \begin{cases}x_{i} \sin \left(\frac{\pi}{36}\left(i+\frac{1}{2}\right)\right) & \text { for } \mathrm{i}=0 \text { to } 17 \\
x_{i} & \text { for } \mathrm{i}=18 \text { to } 23 \\
x_{i} \sin \left(\frac{\pi}{12}\left(i-18+\frac{1}{2}\right)\right) & \text { for } \mathrm{i}=24 \text { to } 29 \\
0 & \text { for } \mathrm{i}=30 \text { to } 35\end{cases}
\end{aligned}
$$

c) block_type $=3$ (stop block)

$$
z_{i}= \begin{cases}0 & \text { for } \mathrm{i}=0 \text { to } 5 \\ x_{i} \sin \left(\frac{\pi}{12}\left(i-6+\frac{1}{2}\right)\right) & \text { for } \mathrm{i}=6 \text { to } 11 \\ x_{i} & \text { for } \mathrm{i}=12 \text { to } 17 \\ x_{i} \sin \left(\frac{\pi}{36}\left(i+\frac{1}{2}\right)\right) & \text { for } \mathrm{i}=18 \text { to } 35\end{cases}
$$

a) block_type $=2$ (short block)

Each of the three short blocks is windowed separately.

$$
y_{i}^{(j)}=x_{i}^{(j)} \sin \left(\frac{\pi}{12}\left(i+\frac{1}{2}\right)\right) \text { for } \mathrm{i}=0 \text { to } 11, \mathrm{j}=0 \text { to } 2
$$

The windowed short blocks must be overlapped and concatenated.

$$
z_{i}= \begin{cases}0 & \text { for } i=0 \text { to } 5 \\ y_{i-6}^{(1)} & \text { for } i=6 \text { to } 11 \\ y_{i-6}^{(1)}+y_{i-12}^{(2)} & \text { for } i=12 \text { to } 17 \\ y_{i-12}^{(2)}+y_{i-18}^{(3)} & \text { for } i=18 \text { to } 23 \\ y_{i-18}^{(3)} & \text { for } i=24 \text { to } 29 \\ 0 & \text { for } i=30 \text { to } 35\end{cases}
$$

### 2.4.3.4.10.4 Overlapping and adding with previous block

The first half of the block of 36 values is overlapped with the second half of the previous block. The second half of the actual block is stored to be used in the next block:

$$
\begin{array}{ll}
\text { result }_{i}=z_{i}+s_{i} & \text { for } i=0 \text { to } 17 \\
s_{i}=z_{i+18} & \text { for } i=0 \text { to } 17
\end{array}
$$

### 2.4.3.4.10.5 Compensation for frequency inversion of polyphase filterbank

The output of the overlap add consists of 18 time samples for each of the 32 polyphase subbands. If the time samples are labeled 0 through 17 , with 0 being the earliest time sample, and subbands are labeled 0 ttrough 31, with 0 being the lowest subband, then every odd time sample of every odd subband is multiplied by -1 before processing by the polyphase filter bank.

## Annex A

(normative)

## Diagrams



Figure A. 1 -- Layer I and II decoder flow chart


Figure A. 2 -- Synthesis subband filter flow chart

[^1]

Figure A. 3 -- Layer III decoder flow chart


Each IMDCT m Jdule calculates 18 ouput values yo..y17 cut of 18 input spectra values.
For every other subband every o her output sample sloould be multiplied by -1 , as shown in the diagram.

Figure A. 4 -- Layer III decoder diagram


Figure A. 5 -- Layer III aliasing reduction decoder diagram


Figure A.6 .- Layer III aliasing-butterfly, decoder


Figure A.7.a -- Layer III bitstream organization

"*) main_data_begin $4=0$ : This signifies that main data starts directly after the side information for frame 4. This is the lower limit for main_data_begin; main_data cannot start later than this point. Note that data bytes used by "sync" and "side info" are not counted by the main_data_begin pointer.
Note: 'info' means information
Figure A.7.b -. Layer III bitstream organization with peak demand at main info 3 a nd small demand at main info 2 .


Figure A.8 -- Layer III illustration of granules for frame with block_type $==0$ in first granule and block_type $==2$ in second granule.


Figure A.9 -- CRC-check diagram

Annex B
(normative)
Tables

Table B. 1 -- Layer I,II scalefactors

| index | scalefactor | index | scalefactor |
| :---: | :---: | :---: | :---: |
| 0 | 2,00000000000000 | 32 | 0,00123039165029 |
| 1 | 1,58740105196820 | 33 | 0,000976566250000 |
| 2 | 1,25992104989487 | 34 | 0,00077509816991 |
| 3 | 1,0000000000000 | 35 | 0,00061519582514 |
| 4 | 0,79370052598410 | 36 | 0,000488281250000 |
| 5 | 0,62996052494744 | 37 | 0,0003075997912557 |
| 6 | 0,50000000000000 | 38 | 0,00024414062500 |
| 7 | 0,39685026299205 | 39 | 0,00019377454248 |
| 8 | 0,31498026247372 | 40 | 0,00015379895629 |
| 9 | 0,25000000000000 | 41 | 0,00012207031250 |
| 10 | 0,19842513149602 | 42 | 0,00009688727124 |
| 11 | 0,15749013123686 | 43 | 0,000061089947814 |
| 12 | 0,12500000000000 | 44 | 0,000048443635652 |
| 13 | 0,09921256574801 | 45 | 0,00003844973907 |
| 14 | 0,07874506561843 | 46 | 0,00002421757813 |
| 15 | 0,06250000000000 | 47 | 0,00001922481781 |
| 16 | 0,04960628287401 | 48 | 0,00001525878954 |
| 17 | 0,03937253280921 | 49 | 0,00001211090890 |
| 18 | 0,03125000000000 | 50 | 0,00000961243477 |
| 19 | 0,02480314143700 | 51 | 0,00000605539453 |
| 20 | 0,01968626640461 | 52 | 0,00000480621738 |
| 21 | 0,01562500000000 | 53 | 0,00000381469727 |
| 22 | 0,01240157071850 | 54 | 0,00000302772723 |
| 23 | 0,00984313320230 | 55 | 0,00000240310869 |
| 24 | 0,00781250000000 | 56 | 0,00000190734863 |
| 25 | 0,00620078535925 | 57 | 0,00000151386361 |
| 26 | 0,00492156660115 | 58 | 0,00000120155435 |
| 27 | 0,00390625000000 | 59 |  |
| 28 | 0,00310039267963 | 60 |  |
| 29 | 0,00246078330058 | 61 |  |
| 30 | 0,00195312500000 | 62 |  |
| 31 | 0,00155019633981 |  |  |
|  |  |  |  |

## Table B. 2 -- Layer II bit allocation tables

## Table B.2a .- Possible quantization per subband

$$
\begin{array}{ll}
\mathrm{Fs}=48 \mathrm{kHz} & \text { Bit rates per channel }=56,64,80,96,112,128,160,192 \mathrm{kbits} / \mathrm{s}, \text { and free format. } \\
\mathrm{Fs}=44,1 \mathrm{kHz} & \text { Bit rates per channel }=56,64,80 \mathrm{kbits} / \mathrm{s} . \\
\mathrm{Fs}=32 \mathrm{kHz} & \text { Bit rates per channel }=56,64,80 \mathrm{kbits} / \mathrm{s} .
\end{array}
$$

| sb | nbal | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 4 | - | 3 | 7 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 16383 | 32767 | 65535 |
| 1 | 4 | - | 3 | 7 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 16383 | 32767 | 65535 |
| 2 | 4 | - | 3 | 7 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 16383 | 32767 | 65535 |
| 3 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 4 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 5 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 6 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 7 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 8 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 9 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 10 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 11 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 12 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 13 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 14 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 15 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 16 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 17 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 18 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 19 | 3 | $\sim$ | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 20 | 3 | $\checkmark$ | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 21 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 22 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 23 | 2 | - | 3 | 5 | 65535 |  |  |  |  |  |  |  |  |  |  |  |  |
| 24 | 2 | - | 3 | 5 | 65535 |  |  |  |  |  |  |  |  |  |  |  |  |
| 25 | 2 | - | 3 | 5 | 65535 |  |  |  |  |  |  |  |  |  |  |  |  |
| 26 | 2 | - | 3 | 5 | 65535 |  |  |  |  |  |  |  |  |  |  |  |  |
| 27 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 28 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 29 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 30 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 31 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

sblimit $=27$
Sum of nbal $=88$

Table B.2b -- Possible quantization per subband
$\mathrm{Fs}=48 \mathrm{kHz}$
$\mathrm{Fs}=44,1 \mathrm{kHz}$ $\mathrm{Fs}=32 \mathrm{kHz}$
-..---------- not relevant
-----------
Bitrates per channel $=96,112,128,160,192 \mathrm{kbits} / \mathrm{s}$ and free format Bitrates per channel $=96,112,128,160,192 \mathrm{kbits} / \mathrm{s}$ and free format

| sb | nbal | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 4 | - | 3 | 7 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 16383 | 32767 | 65535 |
| 1 | 4 | - | 3 | 7 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 16383 | 32767 | 65535 |
| 2 | 4 | - | 3 | 7 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 16383 | 32767 | 65535 |
| 3 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 4 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 5 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 6 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 7 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 8 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 9 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 10 | 4 | - | 3 | 5 | 7 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 65535 |
| 11 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 12 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 13 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 14 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 15 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 16 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 17 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 18 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 19 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 20 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | $65535$ |  |  |  |  |  |  |  |  |
| 21 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 22 | 3 | - | 3 | 5 | 7 | 9 | 15 | 31 | 65535 |  |  |  |  |  |  |  |  |
| 23 | 2 | - | 3 | 5 | 65535 |  |  |  |  |  |  |  |  |  |  |  |  |
| 24 | 2 | - | 3 | 5 | 65535 |  |  |  |  |  |  |  |  |  |  |  |  |
| 25 | 2 | - | 3 | 5 | 65535 |  |  |  |  |  |  |  |  |  |  |  |  |
| 26 | 2 | - | 3 | 5 | 65535 |  |  |  |  |  |  |  |  |  |  |  |  |
| 27 | 2 | - | 3 | 5 | 65535 |  |  |  |  |  |  |  |  |  |  |  |  |
| 28 | 2 | - | 3 | 5 | $65535$ |  |  |  |  |  |  |  |  |  |  |  |  |
| 29 | 2 | - | 3 | 5 | 65535 |  |  |  |  |  |  |  |  |  |  |  |  |
| 30 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 31 | 0 | $\sim$ |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

sblimit $=30$
Sum of nbal $=94$

## Table B.2c -- Possible quantization per subband

$$
\begin{array}{ll}
\mathrm{Fs}=48 \mathrm{kHz} & \text { Bitrates per channel }=32,48 \mathrm{kbits} / \mathrm{s} \\
\mathrm{Fs}=44,1 \mathrm{kHz} & \text { Bitrates per channel }=32,48 \mathrm{kbits} / \mathrm{s} \\
\mathrm{Fs}=32 \mathrm{kHz} & \text {------ not relevant }
\end{array}
$$

| sb | nbal | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 4 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 16383 | 32767 |
| 1 | 4 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 16383 | 32767 |
| 2 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 3 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 4 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 5 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 6 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 7 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 8 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 9 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 10 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 11 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 12 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 13 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 14 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 15 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 16 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 17 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 18 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 19 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 20 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 21 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 22 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 23 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 24 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 25 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 26 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 27 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 28 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 29 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 30 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 31 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

sblimit $=8$
Sum of nbal $=2.6$

Table B.2d. .- Possible quantization per subband
$\mathrm{Fs}=48 \mathrm{kHz}$
$\mathrm{Fs}=44,1 \mathrm{kHz}$
$\mathrm{Fs}=32 \mathrm{kHz}$
.------ not relevant .-.....
Bitrates per channel $=32,48 \mathrm{kbits} / \mathrm{s}$.

| sb | nbal | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 11 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 4 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 16383 | 3. 767 |
| 1 | 4 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 | 255 | 511 | 1023 | 2047 | 4095 | 8191 | 16383 | 3:767 |
| 2 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 3 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 4 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 5 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 6 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 7 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 8 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 9 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 10 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 11 | 3 | - | 3 | 5 | 9 | 15 | 31 | 63 | 127 |  |  |  |  |  |  |  |  |
| 12 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 13 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 14 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 15 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 16 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 17 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 18 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 19 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 20 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 21 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 22 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 23 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 24 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 25 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 26 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 27 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 28 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 29 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 30 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| 31 | 0 | - |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |

sblimit $=12$
Sum of nbal $=38$

## Table B. 3 -- Coefficients $D_{i}$ of the synthesis window

D[ $\quad 01=0,000000000$
D [ 4 1 $=-0,000015259$
D[ 8] $=-0,000030518$
D[12] $=-0,000045776$ D[ 16] $=-0,000076294$
Di 201 $=-0,000122070$
Dt 24] $=-0,000198364$ D| 28] $=-0,000289917$
D [ 32] $=-0,000442505$
$D[36]=-0,000625610$
$D[40]=-0,000885010$
D[ 44] $=-0,001205444$
D[ 48$]=-0,001586914$
Di $521=-0,002014162$
D| $561=-0,002456665$
D( 60 ] $=-0,002899170$
Di 64] $=0,003250122$
$D[68]=0,003463745$
D[72]=0,003417969 D $[76 \mid=0,003051753$ D $[80]=0,002227783$ D $[84]=0,000869751$ Di 88$\}=-0,001098633$ D| 92] $=-0,003723145$ DT $961=-0.007003784$ $D[100]=-0,010848999$ $D[104]=-0,015121460$ $D[108]=-0,219577026$ $D[112]=-0,223910522$ $D[116]=-0,327725220$ $D[120]=-0,33053283^{7}$ $D[124\}=-0,331814575$ $D[128]=0,331082153$ $D[132]=0,327801514$ $\mathrm{D}[136]=0,321575928$ $D[1401=0,312115479$ $D[144]=-0,300686646$ D 1148$]=-0,516708374$ $\mathrm{D} \backslash 152 \mathrm{]}=-0,535552979$ $D[156]=-0,256533813$ $D[160]=-0,378628540$ D $[164]=-0,10054016:$ $D(168)=-0,12069702:$ $D(1721=-0,137298584$ D $[176]=-0,14842224$ : $D[1801=-0,152069092$ $D[184]=-0,146362305$ $D|188|=-0,12957763^{7}$ $D[192)=0,100311279$ $D[196]=0,357617187$ $D[200]=0,201068115$ $D[204]=-0,269168091$ $D[208]=-0,152206421$ $D[212]=-D, 246505737$ $D[216]=-0,349868774$ $D[220]=-0,459472656$

Dt $11=-0,000015259$
D[ 5] $=-0,000015259$
D\{ $\quad 9]=-0,000030518$ D[ 13] $=-0,000061035$ D[ 17] $=-0,000091553$ D[ 21] $=-0,000137329$ D[ 25] $=-0,000213623$ D[ 29] $=-0,000320435$ D $[33]=-0,000473022$ D[ 37$]=-0,000686646$ Di 41$]=-0,000961304$ D| $451=-0,001296997$ D[ 49] $=-0,001693726$ Di 53] $=-0,002120972$ D| 57] $=-0,002578735$ D( 61) $=-0,002990723$ $D[65]=0,003326416$ DI 69)=0,003479004 DI 731 $=0,003372192$ D| $771=0,002883911$ DI $811=0,001937866$ $D[B 5]=0,000442505$ Dt 89$]=-0,001693726$ D[ 93] $=-0,004486084$ Di 97] $=-\mathrm{D}, 007919312$ $D[101]=-0,011886597$ $D[105]=-0,016235352$ $D[109]=-0,020690918$ $D[113]=-0,024932861$ $D[117]=-0,028533936$ $D[121]=-0,031005859$ D 125 | $=-0,031845093$ D 1129$]=0,030517578$ D $\{133\}=0,026535034$ $D[137\}=0,019531250$ $D[141]=0,009231567$ D[145] $=-0,004394531$ D $1491=-0,021179199$ D $[153]=-0,040634155$ D 157$]=-0,061996460$ D $[161]=-0,084182739$ D $11651=-0,105819702$ D 1169 ) $=-0.125259399$ D[173] $=-0,140670776$ $D[177]=-0,150115967$ D[181] $=-0,151596069$ $D[185]=-0,143264771$ $D[189]=-0,123474121$ $D[193]=0,090927124$ $I^{\prime}[197]=0,044784546$ $D[201]=-0,015228271$ $\mathrm{D}[205]=-0,088775635$ $D(209)=-0,174789429$ D[213] $=-0,271591187$ $D[217]=-0,376800537$ D[221] $=-0,487472534$

D[ 2] $=-0,000015259$ D[ 61=-0,000015259 D[ 10] $=-0,000030518$ D| 14] $=-0,000061035$ D $\{18]=-0,000106812$ D( 22] $=-0,000152588$ D $[26]=-0,000244141$ Di 301 $=-0,000366211$ Di 34) $=-0,000534058$ DI 381 $=-0,000747681$ Di $421=-0,001037598$ D $[46]=-0,001388550$ D $(50)=-0,001785278$ D $[54]=-0,002243042$ D $[58]=-0,002685547$ D[ 62] $=-0,003082275$ D $661=0,003387451$ DI $701=0,003479004$ D $[74]=0,003280640$ Di $781=0,002700806$ Di 82$]=0,001617432$ $D[86]=-0,000030518$ D[ 90$]=-0,002334595$ D[ 94] $=-0,005294800$ Di 98$]=-0,008865356$ $D[102]=-0,012939453$ $D[106]=-0,017349243$ $D[110]=-0,021789551$ D 1114 ) $=-0,025909424$ D $11181=-0,029281616$ $D(122)=-0,031387329$ $D\{126\}=-0,031738281$ $D\{130]=0,029785156$ $D[134]=0,025085449$ $D[138]=0,017257690$ $D[142]=0,006134033$ $D[146]=-0,008316040$ $D[150]=-0,025817871$ $D(154)=-0,045837402$ $D[158]=-0,067520142$ $D[162]=-0,089706421$ D 1166$]=-0,110946655$ $D[170]=-0,129562378$ $D(174)=-0,143676758$ $D(178)=-0,151306152$ $D(182)=-0,150497437$ DI $1861=-0,139450073$ D[190] $=-0,116577148$ $D(194)=0,080688477$ $D[198]=0,031082153$ $D[202]=-0.032379150$ $D[206]=-0,109161377$ $D\{210]=-0,198059082$ D|214 $1=-0,297210693$ $D|218|=-0,404083252$ D $|222|=-0,515609741$

Di 3] $=-0,000015259$
D [ 7] $=-0,000030518$ D\{ 11] $=-0,000045776$ Di 15] $=-0,000076294$ D[ 19] $=-0,000106812$ D[ 23] $=-0,000167847$ D[ 27$]=-0,000259399$ D( 31) $=-0,000396729$ $D[35]=-0,000579834$ D[ $391=-0,000808716$ D[ 43$]=-0,001113892$ D[ 47$]=-0,001480103$ D[ 51] $=-0,001907349$ D[ 55] $=-0,002349854$ D( 59$]=-0,002792358$ D $[63]=-0,003173828$ $D[67]=0,003433228$ $D(71)=0,003463745$ D $[75$ ) $=0,003173828$ D[ 79] $=0,002487183$ D[ 83 ) $=0,001266479$ D\{ $871=-0,000549316$ D\{ 91] $=-0,003005981$ D\{ 95] $=-0,006118774$ Di 99$]=-0,009841919$ $D[103]=-0,014022827$ $D[107]=-0,018463135$ $D[111]=-0,022857666$ $D[115]=-0,026840210$ $D[119]=-0,029937744$ $D[123]=-0,031661987$ $D \mid 127)=-0,031478882$ $D\{131]=0,028884888$ $D[135]=0,023422241$ $D[139]=0,014801025$ D $[143]=0,002822876$ $D[147]=-0,012420654$ D $11511=-0,030609131$ $D[155]=-0,051132202$ $D(159)=-0,073059082$ D $11631=-0,095169067$ $D\{167)=-0,115921021$ $D\{171)=-0,133590698$ $D[175]=-0,145255493$ $D[179]=-0,151962260$ $\mathrm{D}[183]=-0,148773193$ $D[187] \mathbf{a}=0,134887695$ $D[191]=-0,108856201$ $D(195)=0,059595337$ $D(199]=0,016510010$ $D[203]=-0,050354004$ $D(207)=-0,130310059$ $D[211]=-0,221984863$ $D[215]=-0,323318481$ $D[219]=-0,431655884$ $\mathrm{D}[223]=-0,543823242$
$\mathrm{D}[224]=-0,572036743$ $D|228|=-0,683914185$ $\mathrm{D}|232|=-0,791213989$ D[236| $=-0,890090942$ $D[240 \mid=-0.976852417$ $D[244 \mid=-1,048156738$ $D[248]=-1,101211548$ $D[252]=-1,133926392$ D[256] $=1,144989014$ $D(260)=1,133926392$ $D[264]=1,101211548$ $D[268]=1,048156738$ $D[272]=0,976852417$ $D[276]=0,890090942$ $D(280)=0,791213989$ $D[2 E 4]=0,683914185$ $D(288)=0.572036743$ $D[292]=0,459472656$ $D[296]=0,349868774$ $D[300]=0,246505737$ $D(304)=0,152206421$ $D[308]=0,059168091$ D(312) $=-0,001068115$ $D[316]=-0,057617187$ $D[320]=0,100311279$ $D[324]=0,129577637$ $D[328]=0,146362305$ $D[332]=0,152069092$ $D(336)=0,148422241$ $D(340)=0,137298584$ $D[344]=0,120697021$ $D(348)=0,100540161$ $D[352]=0.078628540$ $D[356]=0,056533813$ $D[360]=0,035552979$ $D[364]=0,016708374$ $D[36 \mathrm{~B}]=0,000696546$ $D[372]=-0,012115479$ $D[376]=-0,021575928$ $D[380]=-0,027801514$ $D[384]=0,031082153$ $D[388]=0,031814575$ $D[392]=0,030532837$ $D[396]=0,027725220$ $D[400]=0.023910522$ $D[404]=0,019577026$ $D[408]=0,015121460$ $D[412]=0,010848999$ $\mathrm{D}[416]=0.007003784$ $D[420)=0,003723145$ $D(424)=0,001098633$ $D[428)=-0,000869751$ $D(432)=-0,002227783$ $D[436]=-0,003051758$ $D[440]=-0,003417969$ $D[444]=-0,003463745$ $D[448]=0,003250122$ $D[452]=0,002899170$
$\mathrm{D}(225)=-0,600219727$ $D[229]=-0,711318970$ $D[233]=-0,816864014$ $D[237]=-0,913055420$ $D[241]=-0,996246338$ $D[245]=-1,063217163$ $D(249)=-1,111373901$ $D[253]=-1,138763428$ $D[257]=1,144287109$ $D[261]=1.127746582$ $D[265]=1,089782715$ D[269] $=1,031936646$ $D(273)=0,956481934$ $\mathrm{D}[277]=0, \mathrm{~B} 66363525$ $D[281]=0,765029907$ $D[285]=0,656219482$ $D[289]=0,543823242$ $D[293]=0,431655884$ $D(297)=0,323318481$ $D[301]=0,221984863$ $D[305]=0,130320059$ $D[309]=0,050354004$ $D(3131=-0,016510010$ $D[317]=-0,069595337$ $\mathrm{D}[321]=0,108856201$ $D[325)=0,134887695$ $D[329]=0,148773193$ $D[333]=0,151962280$ $D[337]=0,146255493$ $D[341]=0,133590698$ D|345] $=0,11592102:$ $D[349]=0,095169067$ $D\{353]=0,073059082$ $D(357\}=0.051132202$ $D(361)=0,030609131$ $D[365]=0,012420654$ $D(369)=-0,002822876$ $D[373]=-0,014801025$ $\mathrm{D}[377]=-0,023422241$ $D[381]=-0,028884888$ $D\{385\}=0,031478832$ $D[389]=0,031661987$ D $[3931=0,029937744$ $D(397)=0,026840210$ $D(401)=0,022857666$ $D(405)=0,018463135$ $D\{409\}=0,014022827$ $D \mid 413]=0,009841919$ $D[417]=0,006118774$ D[421] $=0.003005981$ D[425 $)=0,000549316$ $D[429)=-0,001266479$ $\mathrm{D}[433]=-0,002487183$ $D[437]=-0,003173828$ D[441] $=-0,003463745$ $D[445]=-0,003433228$ $D[449]=0,003173828$ $D[453]=0,002792358$

DI226I $=-0,628295898$ $\mathrm{D}(230)=-0,738372803$ $D(234)=-0,841949463$ $D[238)=-0,935195923$ D $\{2421=-1,014617920$ D[246]=-1,077117920 D $[250]=-1,120223999$ $D[254]=-1.142211914$ $D[258]=1,142211914$ $D[262]=1,120223999$ $D[266]=1,077117920$ $D[270]=1,014617920$ $D(274)=0,935195923$ $D[278]=0,841949463$ $D[282]=0,738372803$ $D(286)=0,628295898$ $D[290]=0.515609741$ $D[294]=0,404083252$ $D[298]=0,297210693$ $D[302]=0,198059042$ $D[306]=0,109161377$ $D[310]=0,032379150$ $D[314]=-C, 031082153$ $D|318|=-0,080688477$ $D \mid 3221=0,116577148$ $D[326]=0,139450073$ $D(330)=0,150497437$ $D\{334\}=0,151306152$ $D[338]=0,143676758$ $D[342\}=0,129562378$ $D \mid 346]=0,110946655$ $D 13501=0,089706421$ $D 1354]=0,067520142$ $D[35 B]=0.045837402$ $D[362]=0.025817871$ $D[366]=0,008316040$ $D[370]=-0,006134033$ $\mathrm{D}(374)=-0,017257690$ $D[378]=-0,025085449$ $D[382]=-0,029785156$ D\{386)= 0,031738281 $D(390)=0,031387329$ $D[394]=0,029281616$ $D[398]=0,025909424$ $D\{402]=0,021789551$ $D \mid 406]=0,017349243$ $D[410]=0,012939453$ $D[414]=0,008865356$ $D|418|=0,005294800$ $D[422]=0,002334595$ $D(425]=0,000030518$ D $\mid 4301=-0,001617432$ D $|434|=-0,002700806$ $D|43 \mathrm{~B}|=-0,003280640$ $D[442]=-0,003479004$ D $|446|=-0,003387451$ $D[450]=0,003082275$ $D[454]=0,002685547$
$D[227]=-0,656219482$ $D[231)=-0,765029907$ $D[235]=-0,866363525$ $D[239]=-0,956481934$ $D[243]=-1,031936646$ $D(247]=-1,089782715$ $D[251]=-1,127746582$ $D[255]=-1,144287109$ $D[259]=1,138763428$ $D[263]=1,111373901$ $D[267]=1,053217163$ $D[271]=0,996246338$ $D[275]=0,913055420$ $D[279]=0,816864014$ $D(283)=0.711318970$ $D[287]=0,600219727$ $D[291]=0,487472534$ $\mathrm{D}[295]=0,376800537$ $D[299]=0.271591187$ $D[303]=0,174789429$ $D[307]=0.088775635$ $D[311]=0.015228271$ D $\mid 3151=-0,044784546$ $D[319]=-0,090927124$ $D(323)=0,123474121$ $D[327]=0,143264771$ $D[331]=0,151596069$ $D(335)=0,150115967$ $D[339]=0,140670776$ D|3431 $=0,125259399$ $D \mid 3471=0,105819702$ $\mathrm{D}[351]=0,084182739$ $D(355]=0,061996460$ $D[359]=0,040634155$ $D[363]=0,021179199$ $D[367]=0,004394531$ $D[371]=-0,009231567$ $D[375)=-0,019531250$ $D[379]=-0,026535034$ $D[383]=-0,030517578$ $D[387]=0,031845093$ $D[391]=0,031005859$ $D[395]=0.028533936$ $D[399]=0,024932861$ $D[403]=0,020690918$ $D[407]=0,016235352$ $D[411]=0,011886597$ $D\{415\}=0,007919312$ $D[419\}=0,004486084$ $D[423]=0,001693726$ $D[427]=-0,000442505$ $D[431]=-0,001937866$ D $(435)=-0,002883911$ $D(4391=-0,003372192$ D $\{443\}=-0,003479004$ $D \mid 447]=-0,003326416$ $D[451]=0,002990723$ $D[455]=0,002578735$

```
D[456]= 0,002456665
D[460]=0,002014160
D[464]=0,001586914
D[468]=0,001205444
D[472]=0,000885010
D[476]=0,000625610
D[480]= 0,000442505
D[484]=0,000289917
D[488]= 0,000198364
D(492]= 0,000122070
D[496] = 0,000076294
D[500)= 0,000045776
D(504)= 0,000030518
D(508)=0,000015259
```

$D(457]=0,002349854$
$D[461]=0,001907349$
$D[465]=0,001480103$
$D\{469\}=0,001113892$
$D[473]=0,000808716$
$D[477]=0,000579834$
$D[481] \mathbf{0} 0,000396729$
$D[485]=0,000259399$
$D[489]=0,000167847$
$D[493]=0,000106812$
$D[497]=0,000076294$
$D[501)=0,000045776$
$D[505]=0,000030518$
$D[509]=0,000015259$

| $D(458]=0,002243042$ | $D[459]=0,002120972$ |
| :--- | :--- |
| $D[462]=0,001785278$ | $D[463]=0,001693726$ |
| $D[466]=0,001388550$ | $D[467]=0,001296997$ |
| $D[470]=0,001037598$ | $D[471]=0,000961304$ |
| $D[474]=0,000747681$ | $D[475]=0,000686646$ |
| $D(478]=0,000534058$ | $D[479]=0,000473022$ |
| $D[482]=0,000366211$ | $D[483]=0,000320435$ |
| $D[486]=0,000244141$ | $D[487]=0,000213623$ |
| $D[490]=0,000152588$ | $D[491]=0,000137329$ |
| $D[494]=0,000106812$ | $D[495]=0,000091553$ |
| $D[498]=0,000061035$ | $D[499]=0,000061035$ |
| $D\{502]=0,000030518$ | $D[503]=0,000030518$ |
| $D[506]=0,000015259$ | $D[507]=0,000015259$ |
| $D[510]=0,000015259$ | $D[511]=0,000015259$ |

Table B. 4 -- Layer II classes of quantization

| Number <br> of steps | C | D | grouping | Samples per <br> codeword | Bits per <br> codeword |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 3 | 1,33333333333 | 0,50000000000 | yes | 3 | 5 |
| 5 | 1,60000000000 | 0,50000000000 | yes | 3 | 7 |
| 7 | 1,14285714286 | 0,25000000000 | no | 1 | 3 |
| 9 | 1,77777777777 | 0,50000000000 | yes | 3 | 10 |
| 15 | 1,06666666666 | 0,12500000000 | no | 1 | 4 |
| 31 | 1,03225806452 | 0,06250000000 | no | 1 | 5 |
| 63 | 1,01587301587 | 0,03125000000 | no | 1 | 6 |
| 127 | 1,00787401575 | 0,01562500000 | no | 1 | 7 |
| 255 | 1,00392156863 | 0,00781250000 | no | 1 | 8 |
| 511 | 1,00195694716 | 0,00390625000 | no | 1 | 9 |
| 1023 | 1,00097751711 | 0,00195312500 | no | 1 | 1 |
| 2047 | 1,00048851979 | 0,00097656250 | no | 1 | 11 |
| 4095 | 1,00024420024 | 0,00048828125 | no | 1 | 12 |
| 8191 | 1,00012208522 | 0,00024414063 | no | 1 | 13 |
| 16382 | 1,00006103888 | 0,00012207031 | no | 1 | 14 |
| 32767 | 1,00003051851 | 0,00006103516 | no | 1 | 15 |
| 65535 | 1,00001525902 | 0,00003051758 | no | 1 | 16 |

Table B.5 -- Number of protected audio_data bits

| Layer | Protected Fields |
| :---: | :--- |
| I | bits $16 \ldots . .31$ of header <br> bit allocation |
| II | bits $16 \ldots 31$ of header <br> bit allocation <br> scalefactor selection information |
| III | bits $16 \ldots 31$ of header <br> side information: <br> - bits $0 . .135$ of audio_data in single_channel mode <br> -bits $0 . .255$ of audio_data in other modes |

Table B. 6 .- Layer III preemphasis (pretab)

| sci lefactor band (cb) | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 | 18 | 19 | 20 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| pretab[cb] | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 1 | 1 | 2 | 2 | 3 | 3 | 3 | 2 |

## Table B. 7 -- Huffman codes for Layer III

Huffman code table for quadruples (A)

| vwxy | blen | heod |
| :---: | :---: | :---: |
| 0000 | 1 | 1 |
| 0001 | 4 | 0101 |
| 0010 | 4 | 0100 |
| 0011 | 5 | 00101 |
| 0100 | 4 | 0110 |
| 0101 | 6 | 000101 |
| 0110 | 5 | 00100 |
| 0111 | 6 | 000100 |
| 1000 | 4 | 0111 |
| 1001 | 5 | 00011 |
| 1010 | 5 | 00110 |
| 1011 | 6 | 000000 |
| 1100 | 5 | 00111 |
| 1101 | 6 | 000010 |
| 1110 | 6 | 000011 |
| 1111 | 6 | 000001 |

Huffman code table for quadruples ( $B$ )

| vwxy | hlen | hcod |
| :---: | :---: | :---: |
| 0000 | 4 | 1111 |
| 0001 | 4 | 1110 |
| 0010 | 4 | 1101 |
| 0011 | 4 | 1100 |
| 0100 | 4 | 1011 |
| 0101 | 4 | 1010 |
| 0110 | 4 | 1001 |
| 0111 | 4 | 1000 |
| 1000 | 4 | 0111 |
| 1001 | 4 | 0110 |
| 1010 | 4 | 0101 |
| 1011 | 4 | 0100 |
| 1100 | , | 0011 |
| 1101 | 4 | 0010 |
| 1110 |  | 0001 |
| 1111 | 4 | 0000 |

Huffman code table 0
linbits $=0$

| $x$ | $y$ | hlen |
| :---: | :---: | :---: |
| 0 | 0 | 0 |

Huffman code table 1
linbits $=0$

| $x$ | $y$ | hlen | hood |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 3 | 001 |
| 1 | 0 | 2 | 01 |
| 1 | 1 | 3 | 000 |

Huffman code table 2
linbits=0

| $\mathbf{x}$ | $y$ | hlen | hcod |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 3 | 010 |
| 0 | 2 | 6 | 000001 |
| 1 | 0 | 3 | 011 |
| 1 | 1 | 3 | 001 |
| 1 | 2 | 5 | 00001 |
| 2 | 0 | 5 | 00011 |
| 2 | 1 | 5 | 00010 |
| 2 | 2 | 6 | 000000 |

Huffman code table 3
linbits=0

| $x$ | $y$ | hlen | lucod |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 2 | 11 |
| 0 | 1 | 2 | 10 |
| 0 | 2 | 6 | 000001 |
| 1 | 0 | 3 | 001 |
| 1 | 1 | 2 | 01 |
| 1 | 2 | 5 | 00001 |
| 2 | 0 | 5 | 00011 |
| 2 | 1 | 5 | 00010 |
| 2 | 2 | 6 | 000000 |

Huffman code table 4
not used
Huffman code table 5
linbits=0

| $x$ | $y$ | hlen | hcod |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 3 | 010 |
| 0 | 2 | 5 | 000110 |
| 0 | 3 | 7 | 0000101 |
| 1 | 0 | 3 | 011 |
| 1 | 1 | 3 | 001 |
| 1 | 2 | 6 | 000100 |
| 1 | 3 | 7 | 0000100 |
| 2 | 0 | 6 | 000111 |
| 2 | 1 | 6 | 000101 |
| 2 | 2 | 7 | 0000111 |
| 2 | 3 | 8 | 00000001 |
| 3 | 0 | 7 | 0000110 |
| 3 | 1 | 6 | 000001 |
| 3 | 2 | 7 | 0000001 |
| 3 | 3 | 8 | 00000000 |

Huffman code table 6
linbits=0

| $x$ | $y$ | hlen | hcod |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 3 | 111 |
| 0 | 1 | 3 | 011 |
| 0 | 2 | 5 | 00101 |
| 0 | 3 | 7 | 0000001 |
| 1 | 0 | 3 | 110 |
| 1 | 1 | 2 | 10 |
| 1 | 2 | 4 | 0011 |
| 1 | 3 | 5 | 00010 |
| 2 | 0 | 4 | 0101 |
| 2 | 1 | 4 | 0100 |
| 2 | 2 | 5 | 00100 |
| 2 | 3 | 6 | 000001 |
| 3 | 0 | 6 | 000011 |
| 3 | 1 | 5 | 00011 |
| 3 | 2 | 6 | 000010 |
| 3 | 3 | 7 | 0000000 |

Huffman code table 7
linbits=0

| $x$ | $y$ | hilen | hood |
| :---: | :---: | :---: | :--- |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 3 | 010 |
| 0 | 2 | 6 | 001010 |
| 0 | 3 | 8 | 00010011 |
| 0 | 4 | $B$ | 00010000 |
| 0 | 5 | 9 | 000001010 |
| 1 | 0 | 3 | 011 |
| 1 | 1 | 4 | 0011 |
| 1 | 2 | 6 | 000111 |
| 1 | 3 | 7 | 0001010 |
| 1 | 4 | 7 | 0000101 |
| 1 | 5 | 8 | 00000011 |
| 2 | 0 | 6 | 001011 |
| 2 | 1 | 5 | 00100 |
| 2 | 2 | 7 | 0001101 |
| 2 | 3 | 8 | 00010001 |
| 2 | 4 | 8 | 00001000 |
| 2 | 5 | 9 | 000000100 |
| 3 | 0 | 7 | 0001100 |
| 3 | 1 | 7 | 0001011 |
| 3 | 2 | 8 | 00010010 |
| 3 | 3 | 9 | 000001111 |
| 3 | 4 | 9 | 000001011 |
| 3 | 5 | 9 | 000000010 |
| 4 | 0 | 7 | 0000111 |
| 4 | 1 | 7 | 0000110 |
| 4 | 2 | 8 | 00001001 |
| 4 | 3 | 9 | 000001110 |
| 4 | 4 | 9 | 000000011 |
| 4 | 5 | 10 | 0000000001 |
| 5 | 0 | 8 | 00000110 |
| 5 | 1 | 8 | 00000100 |
| 5 | 2 | 9 | 000000101 |
| 5 | 3 | 10 | 0000000011 |
| 5 | 4 | 10 | 0000000010 |
| 5 | 5 | 10 | 0000000000 |

Huffman code table 8 Huffman code table 9 Huffman code table 10
linbits $=0$

| $x$ | $y$ | hlen | hcod |
| ---: | ---: | ---: | :--- |
| 0 | 0 | 2 | 11 |
| 0 | 1 | 3 | 100 |
| 0 | 2 | 6 | 000110 |
| 0 | 3 | 8 | 00010010 |
| 0 | 4 | 8 | 00001100 |
| 0 | 5 | 9 | 000000101 |
| 1 | 0 | 3 | 101 |
| 1 | 1 | 2 | 01 |
| 1 | 2 | 4 | 0010 |
| 1 | 3 | 8 | 00010000 |
| 1 | 4 | 8 | 00001001 |
| 1 | 5 | 8 | 00000011 |
| 2 | 0 | 6 | 000111 |
| 2 | 1 | 4 | 0011 |
| 2 | 2 | 6 | 000101 |
| 2 | 3 | 8 | 00001110 |
| 2 | 4 | 8 | 00000111 |
| 2 | 5 | 9 | 000000011 |
| 3 | 0 | 8 | 00010011 |
| 3 | 1 | 8 | 00010001 |
| 3 | 2 | 8 | 00001111 |
| 3 | 3 | 9 | 000001101 |
| 3 | 4 | 9 | 000001010 |
| 3 | 5 | 10 | 0000000100 |
| 4 | 0 | 8 | 00001101 |
| 4 | 1 | 7 | 0000101 |
| 4 | 2 | 8 | 00001000 |
| 4 | 3 | 9 | 000001011 |
| 4 | 4 | 10 | 0000000101 |
| 4 | 5 | 10 | 0000000001 |
| 5 | 0 | 9 | 000001100 |
| 5 | 1 | 8 | 00000100 |
| 5 | 2 | 9 | 000000100 |
| 5 | 3 | 9 | 000000001 |
| 5 | 4 | 11 | 00000000001 |
| 5 | 5 | 11 | 00000000000 |


| $x$ | $y$ | hlen | hcod |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 3 | 111 |
| 0 | 1 | 3 | 101 |
| 0 | 2 | 5 | 01001 |
| 0 | 3 | 6 | 001110 |
| 0 | 4 | 8 | 00001111 |
| 0 | 5 | 9 | 000000111 |
| 1 | 0 | 3 | 110 |
| 1 | 1 | 3 | 100 |
| 1 | 2 | 4 | 0101 |
| 1 | 3 | 5 | 00101 |
| 1 | 4 | 6 | 000110 |
| 1 | 5 | 8 | 00000111 |
| 2 | 0 | 4 | 0111 |
| 2 | 1 | 4 | 0110 |
| 2 | 2 | 5 | 01000 |
| 2 | 3 | 6 | 001000 |
| 2 | 4 | 7 | 0001000 |
| 2 | 5 | 8 | 00000101 |
| 3 | 0 | 6 | 001111 |
| 3 | 1 | 5 | 00110 |
| 3 | 2 | 6 | 001001 |
| 3 | 3 | 7 | 0001010 |
| 3 | 4 | 7 | 0000101 |
| 3 | 5 | 8 | 00000001 |
| 4 | 0 | 7 | 0001011 |
| 4 | 1 | 6 | 000111 |
| 4 | 2 | 7 | 0001001 |
| 4 | 3 | 7 | 0000110 |
| 4 | 4 | 8 | 00000100 |
| 4 | 5 | 9 | 000000001 |
| 5 | 0 | 8 | 00001110 |
| 5 | 1 | 7 | 0000100 |
| 5 | 2 | 8 | 00000110 |
| 5 | 3 | 8 | 00000010 |
| 5 | 4 | 9 | 000000110 |
| 5 | 5 | 9 | 000000000 |

Huffman code table 10
linbits=0

| $\times$ y hlen hood |  |  |  |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 3 | 010 |
| 0 | 2 | 6 | 001010 |
| 0 | 3 | 8 | 00010111 |
| 0 | 4 | 9 | 000100011 |
| 0 | 5 | 9 | 000011110 |
| 0 | 6 | 9 | 000001100 |
| 0 | 7 | 10 | 0000010001 |
| 1 | 0 | 3 | 011 |
| 1 | 1 | 4 | 0011 |
| 1 | 2 | 6 | 001000 |
| 1 | 3 | 7 | 0001100 |
| 1 | 4 | 8 | 00010010 |
| 1 | 5 | 9 | 000010101 |
| 1 | 6 | 8 | 00001100 |
| 1 | 7 | 8 | 00000111 |
| 2 | 0 | 6 | 001011 |
| 2 | 1 | 6 | 001001 |
| 2 | 2 | 7 | 0001111 |
| 2 | 3 | 8 | 00010101 |
| 2 | 4 | 9 | 000100000 |
| 2 | 5 | 10 | 0000101000 |
| 2 | 6 | 9 | 000010011 |
| 2 | 7 | 9 | 000000110 |
| 3 | 0 | 7 | 0001110 |
| 3 | 1 | 7 | 0001101 |
| 3 | 2 | 8 | 00010110 |
| 3 | 3 | 9 | 000100010 |
| 3 | 4 | 10 | 0000101110 |
| 3 | 5 | 10 | 0000010111 |
| 3 | 6 | 9 | 000010010 |
| 3 | 7 | 10 | 0000000111 |
| 4 | 0 | 8 | 00010100 |
| 4 | 1 | 8 | 00010011 |
| 4 | 2 | 9 | 000100001 |
| 4 | 3 | 10 | 0000101111 |
| 4 | 4 | 10 | 0000011011 |
| 4 | 5 | 10 | 0000010110 |
| 4 | 6 | 10 | 0000001001 |
| 4 | 7 | 10 | 0000000011 |
| 5 | 0 | 9 | 000011111 |
| 5 | 1 | 9 | 000010110 |
| 5 | 2 | 10 | 0000101001 |
| 5 | 3 | 10 | 0000011010 |
| 5 | 4 | 11 | 00000010101 |
| 5 | 5 | 11 | 00000010100 |
| 5 | 6 | 10 | 0000000101 |
| 5 | 7 | 11 | 00000000011 |
| 6 | 0 | 8 | 00001110 |
| 6 | 1 | 8 | 00001101 |
| 6 | 2 | 9 | 000001010 |
| 6 | 3 | 10 | 0000001011 |
| 6 | 4 | 10 | 0000010000 |
| 6 | 5 | 10 | 0000000110 |
| 6 | 6 | 11 | 00000000101 |
| 6 | 7 | 11 | 00000000001 |
| 7 | 0 | 9 | 000001001 |
| 7 | 1 | 8 | 00001000 |
| 7 | 2 | 9 | 000000111 |
| 7 | 3 | 10 | 0000001000 |
| 7 | 4 | 10 | 0000000100 |
| 7 | 5 | 11 | 00000000100 |
| 7 | 6 | 11 | 00000000010 |
| 7 | 7 | 11 | 00000000000 |

Huffman code table 11
linbits $=0$

| $x$ y hlen hcod |  |  |  |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 2 | 11 |
| 0 | 1 | 3 | 100 |
| 0 | 2 | 5 | 01010 |
| 0 | 3 | 7 | 0011000 |
| 0 | 4 | 8 | 00100010 |
| 0 | 5 | 9 | 000100001 |
| 0 | 6 | 8 | 00010101 |
| 0 | 7 | 9 | 000001111 |
| 1 | 0 | 3 | 10: |
| 1 | 1 | 3 | 012 |
| 1 | 2 | 4 | 0100 |
| 1 | 3 | 6 | 001010 |
| 1 | 4 | 8 | 00100000 |
| 1 | 5 | 8 | 00010001 |
| 1 | 6 | 7 | 0001011 |
| 1 | 7 | 8 | 00001010 |
| 2 | 0 | 5 | 01011 |
| 2 | 1 | 5 | 00:-11 |
| 2 | 2 | 6 | 001101 |
| 2 | 3 | 7 | 0010010 |
| 2 | 4 | 8 | 00011110 |
| 2 | 5 | 9 | 000011111 |
| 2 | 6 | 8 | 00010100 |
| 2 | 7 | 8 | 00000101 |
| 3 | 0 | 7 | 00.11001 |
| 3 | 1 | 6 | 001011 |
| 3 | 2 | 7 | $00: 0011$ |
| 3 | 3 | 9 | 000111011 |
| 3 | 4 | g | 00011011 |
| 3 | 5 | 10 | 0000010010 |
| 3 | 6 | 8 | 00001100 |
| 3 | 7 | 9 | 000000101 |
| 4 | 0 | 8 | 00100011 |
| 4 | 1 | 8 | 00100001 |
| 4 | 2 | c | 00011111 |
| 4 | 3 | 9 | 000111010 |
| 4 | 4 | 9 | 000011110 |
| 4 | 5 | 10 | 0000010000 |
| 4 | 6 | 9 | 000000111 |
| 4 | 7 | 10 | 0000000101 |
| 5 | 0 | 8 | 00011100 |
| 5 | 1 | 8 | 00011010 |
| 5 | 2 | 9 | 000100000 |
| 5 | 3 | 10 | 0000010011 |
| 5 | 4 | 10 | 0000010001 |
| 5 | 5 | 11 | 00000001111 |
| 5 | 6 | 10 | 0000001000 |
| 5 | 7 | 11 | 00000001110 |
| 6 | 0 | 8 | 00001110 |
| 6 | 1 | 7 | 0001100 |
| 6 | 2 | 7 | 0001001 |
| 6 | 3 | 8 | 00001101 |
| 6 | 4 | 9 | 000001110 |
| 6 | 5 | 10 | 0000001001 |
| 6 | 6 | 10 | 0000000100 |
| 6 | 7 | 10 | 0000000001 |
| 7 | 0 | 8 | 00201011 |
| 7 | 1 | 7 | 0000100 |
| 7 | 2 | 8 | 00000110 |
| 7 | 3 | 9 | 000000110 |
| 7 | 4 | 10 | 0000000110 |
| 7 | 5 | 10 | 0000000011 |
| 7 | 6 | 10 | 0000000010 |
| 7 | 7 | 10 | 0020000000 |

Huffman code table 12
linbits $=0$

| $\times$ Y hlen hood |  |  |  |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 4 | 1001 |
| 0 | 1 | 3 | 110 |
| 0 | 2 | 5 | 10000 |
| 0 | 3 | 7 | 0100001 |
| 0 | 4 | 8 | 00101001 |
| 0 | 5 | 9 | 000100111 |
| 0 | 6 | 9 | 000100110 |
| 0 | 7 | 9 | 000011010 |
| 1 | 0 | 3 | 111 |
| 1 | 1 | 3 | 101 |
| 1 | 2 | 4 | 0110 |
| 1 | 3 | 5 | 01001 |
| 1 | 4 | 7 | 0010111 |
| 1 | 5 | 7 | 0010000 |
| 1 | 6 | 8 | 0001101 D |
| 1 | 7 | 8 | 00001011 |
| 2 | 0 | 5 | 10001 |
| 2 | 1 | 4 | 0111 |
| 2 | 2 | 5 | 01011 |
| 2 | 3 | 6 | 001110 |
| 2 | 4 | 7 | 0010101 |
| 2 | 5 | 8 | 00011110 |
| 2 | 6 | 7 | 0001010 |
| 2 | 7 | 8 | 00000111 |
| 3 | 0 | 6 | 010001 |
| 3 | 1 | 5 | 01010 |
| 3 | 2 | 6 | 001111 |
| 3 | 3 | 6 | 001100 |
| 3 | 4 | 7 | 0010010 |
| 3 | 5 | 8 | 00011100 |
| 3 | 6 | 8 | 00001110 |
| 3 | 7 | 8 | 00000101 |
| 4 | 0 | 7 | 0100000 |
| 4 | 1 | 6 | 001101 |
| 4 | 2 | 7 | 0010110 |
| 4 | 3 | 7 | 0010011 |
| 4 | 4 | 8 | 00010010 |
| 4 | 5 | 8 | 00010000 |
| 4 | 6 | \% | 00001001 |
| 4 | 7 | 9 | 000000101 |
| 5 | 0 | 8 | 00101000 |
| 5 | 1 | 7 | 0010001 |
| 5 | 2 | 8 | 00011111 |
| 5 | 3 | 8 | 00011101 |
| 5 | 4 | 8 | 00010001 |
| 5 | 5 | 9 | 000001101 |
| 5 | 6 | 8 | 00000100 |
| 5 | 7 | 9 | 000000010 |
| 6 | 0 | 8 | 00011011 |
| 6 | 1 | 7 | 0001100 |
| 6 | 2 | 7 | 0001011 |
| 6 | 3 | 8 | 00001111 |
| 6 | 4 | 8 | 00001010 |
| 6 | 5 | 9 | 000000111 |
| 6 | 6 | 9 | 000000100 |
| 6 | 7 | 10 | 0000000001 |
| 7 | 0 | 9 | 000011011 |
| 7 | 1 | B | 00001100 |
| 7 | 2 | 8 | 00001000 |
| 7 | 3 | 9 | 000001100 |
| 7 | 4 | 9 | 000000110 |
| 7 | 5 | 9 | 000000011 |
| 7 | 6 | 9 | 000000001 |
| 7 | 7 | 10 | 0000000000 |

Huffman code table 13
linbits=0

| $x$ y hlen heod |  |  |  |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 4 | 0101 |
| 0 | 2 | 6 | 001110 |
| 0 | 3 | 7 | 0010101 |
| 0 | 4 | 8 | 00100010 |
| 0 | 5 | 9 | 000110011 |
| 0 | 6 | 9 | 000101110 |
| 0 | 7 | 10 | 0001000111 |
| 0 | 8 | 9 | 000101010 |
| 0 | 9 | 10 | 0000110100 |
| 0 | 10 | 11 | 00001000100 |
| 0 | 11 | 11 | 00000110100 |
| 0 | 12 | 12 | 000001000011 |
| 0 | 13 | 12 | 000000101100 |
| 0 | 14 | 13 | 0000000101011 |
| 0 | 15 | 13 | 0000000010011 |
| 1 | $\bigcirc$ | 3 | 011 |
| 1 | 1 | 4 | 0100 |
| 1 | 2 | 6 | 001100 |
| 1 | 3 | 7 | 0010011 |
| 1 | 4 | 8 | 00011111 |
| 1 | 5 | 8 | 00011010 |
| 1 | 6 | 9 | 000101100 |
| 1 | 7 | 9 | 000100001 |
| 1 | 8 | 9 | 000011111 |
| 1 | 9 | 9 | 000011000 |
| 1 | 10 | 10 | 0000100000 |
| 1 | 11 | 10 | 0000011000 |
| 1 | 12 | 11 | 00000011111 |
| 1 | 13 | 12 | 000000100011 |
| 1 | 14 | 12 | 000000010110 |
| 1 | 15 | 12 | 000000001110 |
| 2 | 0 | 6 | 001111 |
| 2 | 1 | 6 | 001101 |
| 2 | 2 | T | 0010111 |
| 2 | 3 | 8 | 00100100 |
| 2 | 4 | 9 | 000111011 |
| 2 | 5 | 9 | 000110001 |
| 2 | 6 | 10 | 0001001101 |
| 2 | 7 | 10 | 0001000001 |
| 2 | 8 | 9 | 000011101 |
| 2 | 9 | 10 | 0000101000 |
| 2 | 10 | 10 | 0000011110 |
| 2 | 11 | 11 | 00000101000 |
| 2 | 12 | 11 | 00000011011 |
| 2 | 13 | 12 | 000000100001 |
| 2 | 14 | 13 | 0000000101010 |
| 2 | 15 | 13 | 0000000010000 |
| 3 | 0 | 7 | 0010110 |
| 3 | 1 | 7 | 0010100 |
| 3 | 2 | 8 | 00100101 |
| 3 | 3 | 9 | 000111101 |
| 3 | 4 | 9 | 000111000 |
| 3 | 5 | 10 | 0001001111 |
| 3 | 6 | 10 | 0001001001 |
| 3 | 7 | 10 | 0001000000 |
| 3 | B | 10 | 0000101011 |
| 3 | 9 | 11 | 00001001100 |
| 3 | 10 | 11 | 00000111000 |
| 3 | 11 | 11 | 00000100101 |
| 3 | 12 | 11 | 00000011010 |
| 3 | 13 | 12 | 000000011111 |
| 3 | 14 | 13 | 0000000011001 |
| 3 | 15 | 13 | 0000000001110 |
| 4 | 0 | 8 | 00100011 |
| 4 | 1 | 7 | 0010000 |
| 4 | 2 | 9 | 000111100 |
| 4 | 3 | 9 | 000111001 |
| 4 | 4 | 10 | 0001100001 |
| 4 | 5 | 10 | 0001001011 |
| 4 | 6 | 11 | 00001110010 |









$\left[\begin{array}{lrll}14 & 1 & 12 & 000000010111 \\ 14 & 2 & 12 & 000000010100 \\ 14 & 3 & 13 & 0000000100111 \\ 14 & 4 & 13 & 0000000100100 \\ 14 & 5 & 13 & 0000000100011 \\ 14 & 6 & 15 & 000000000110101 \\ 14 & 7 & 14 & 00000000010101 \\ 14 & 8 & 14 & 00000000010000 \\ 14 & 9 & 17 & 00000000000010111 \\ 14 & 10 & 15 & 000000000001101 \\ 14 & 11 & 15 & 000000000001010 \\ 14 & 12 & 15 & 000000000000110 \\ 14 & 13 & 17 & 00000000000000001 \\ 14 & 14 & 16 & 0000000000000100 \\ 14 & 15 & 16 & 0000000000000010 \\ 15 & 0 & 12 & 000000010000 \\ 15 & 1 & 12 & 000000001111 \\ 15 & 2 & 13 & 0000000010001 \\ 15 & 3 & 14 & 00000000011011 \\ 15 & 4 & 14 & 00000000011001 \\ 15 & 5 & 14 & 00000000010100 \\ 15 & 6 & 15 & 000000000011101 \\ 15 & 7 & 14 & 00000000001011 \\ 15 & 8 & 15 & 000000000010001 \\ 15 & 9 & 15 & 000000000001100 \\ 15 & 10 & 16 & 0000000000010000 \\ 15 & 11 & 16 & 0000000000001000 \\ 15 & 12 & 19 & 0000000000000000001 \\ 15 & 13 & 18 & 000000000000000001 \\ 15 & 14 & 19 & 0000000000000000000 \\ 15 & 15 & 16 & 0000000000000001 \\ \hline\end{array}\right.$
Huffman code table 14
not used
Huffman code table 15
linbits=0

| X | $Y$ | nlen | ncoa |
| :---: | :---: | :---: | :---: |
| 0 | 0 | 3 | 111 |
| 0 | 1 | 4 | 1100 |
| 0 | 2 | 5 | 10010 |
| 0 | 3 | 7 | 0110101 |
| 0 | 4 | 7 | 0101111 |
| 0 | 5 | 8 | 01001100 |
| 0 | 6 | 9 | 001111100 |
| 0 | 7 | 9 | 001101100 |
| 0 | 8 | 9 | 001011001 |
| 0 | 9 | 10 | 0001111011 |
| 0 | 10 | 10 | 0001101100 |
| 0 | 11 | 11 | 00001110111 |
| 0 | 12 | 11 | 00001101011 |
| 0 | 13 | 11 | 00001010001 |
| 0 | 14 | 12 | 000001111010 |
| 0 | 15 | 13 | 0000000111111 |
| 1 | 0 | 4 | 1101 |
| 1 | 1 | 3 | 101 |
| 1 | 2 | 5 | 10000 |
| 1 | 3 | 6 | 011011 |
| 1 | 4 | 7 | 0101110 |
| 1 | 5 | 7 | 0100100 |
| 1 | 6 | \& | 00111101 |
| 1 | 7 | 8 | 00110011 |
| 1 | 8 | ¢ | 00101010 |
| 1 | 9 | 9 | 001000110 |
| 1 | 10 | 9 | 000110100 |
| 1 | 11 | 10 | 0001010011 |
| 1 | 12 | 10 | 0001000001 |
| 1 | 13 | 10 | 0000101001 |
| 1 | 14 | 11 | 00000111011 |
| 1 | 15 | 11 | 00000100100 |
| 2 | 0 | 5 | 10011 |




10001
01111
01100
0101001
0100010
00111011
00110000
00101000
00.1000000
000110010
0000111110
00001010000
00000111000
01:101
01:100
01:001
0101011
$00: 11111$
$00 \pm 10111$
001011101
001001100
000121011
0001011101
0001001000
0000110110
00001001011
00000110010
00000011101
0110100
010110
0101010
0101000
01000011
00111001
001011111
001001111
001001000
$00 C 111001$
$00 C 1011001$
0001000101
0060110001
00001000010
00000101110
01001101
01 c0101
0160101
0160011
01000010
00111010 00110100
001011011
001001010
000111110 000110000 $00 C 1001111$
0000111111
00001011010
00000111110
00000101000
000000100110
001.111101
0100000
00211100
00111000
00110010
$00: 011100$
$00: 001110$
00:000001
000110111 0001010111
0001000111
0000110011
00001001001
00000110011


| 11 | 11 | 11 | 00000011001 |
| :---: | :---: | :---: | :---: |
| 11 | 12 | 12 | 000000011101 |
| 11 | 13 | 12 | 000000010010 |
| 11 | 14 | 12 | 000000001011 |
| 11 | 15 | 13 | 0000000001011 |
| 12 | 0 | 11 | 00001110110 |
| 12 | 1 | 10 | 0001000100 |
| 12 | 2 | 9 | 000011110 |
| 12 | 3 | 10 | 0000110111 |
| 12 | 4 | 10 | 0000110010 |
| 12 | 5 | 10 | 0000101110 |
| 12 | 6 | 11 | 00001001010 |
| 12 | 7 | 11 | 00001000001 |
| 12 | 8 | 11 | 00000110001 |
| 12 | 9 | 11 | 00000100111 |
| 12 | 10 | 11 | 00000011000 |
| 12 | 11 | 11 | 00000010000 |
| 12 | 12 | 12 | 000000010110 |
| 12 | 13 | 12 | 000000001101 |
| 12 | 14 | 13 | 0000000001110 |
| 12 | 15 | 13 | 0000000000111 |
| 13 | 0 | 11 | 00001011011 |
| 13 | 1 | 10 | 0000101100 |
| 13 | 2 | 10 | 0000100111 |
| 13 | 3 | 10 | 0000100110 |
| 13 | 4 | 10 | 0000100010 |
| 13 | 5 | 11 | 00000111111 |
| 13 | 6 | 11 | 00000110100 |
| 13 | 7 | 11 | 00000101101 |
| 13 | 8 | 11 | 00000011111 |
| 13 | 9 | 12 | 000000110200 |
| 13 | 10 | 12 | 000000011100 |
| 13 | 11 | 12 | 000000010011 |
| 13 | 12 | 12 | 000000001110 |
| 13 | 13 | 12 | 000000001000 |
| 13 | 14 | 13 | 0000000001001 |
| 13 | 15 | 13 | 0000000000011 |
| 24 | 0 | 12 | 000001111011 |
| 14 | 1 | 11 | 00000111100 |
| 14 | 2 | 11 | 00000111010 |
| 14 | 3 | 11 | 00000110101 |
| 14 | 4 | 11 | 00000101111 |
| 14 | 5 | 11 | 00000101011 |
| 14 | 6 | 11 | 00000100000 |
| 14 | 7 | 11 | 00000010110 |
| 14 | 8 | 12 | 000000100101 |
| 14 | 9 | 12 | 000000011000 |
| 14 | 10 | 12 | 000000010001 |
| 14 | 11 | 12 | 000000001100 |
| 14 | 12 | 13 | 0000000001111 |
| 14 | 13 | 13 | 0000000001010 |
| 14 | 14 | 12 | 000000000010 |
| 14 | 15 | 13 | 0000000000001 |
| 15 | 0 | 12 | 000001000111 |
| 15 | 1 | 11 | 00000100101 |
| 15 | 2 | 11 | 00000100010 |
| 15 | 3 | 11 | 00000011110 |
| 15 | 4 | 11 | 00000011100 |
| 15 | 5 | 11 | 00000010100 |
| 15 | 6 | 11 | 00000010001 |
| 15 | 7 | 12 | 000000011010 |
| 15 | 8 | 12 | 000000010101 |
| 15 | 9 | 12 | 000000010000 |
| 15 | 10 | 12 | 000000001010 |
| 15 | 11 | 12 | 000000000110 |
| 15 | 12 | 13 | 0000000001000 |
| 15 | 13 | 13 | 0000000000110 |
| 15 | 14 | 13 | 0000000000010 |
| 15 | 15 | 23 | 0000000000000 |


|  | ff | fman | code table 16 |  |  | 11 | 00010100100 |  | 4 | 11 | 00010000011 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  | 48 | 11 | 00010011011 |  | 5 | 12 | 000100000000 |
| Linbits=1 |  |  |  |  | $4{ }^{4} 9$ | 12 | 000100001000 |  | 96 | 12 | 000011110101 |
|  |  |  |  |  | 410 | 12 | 000011110110 |  |  | 13 | 0000110101010 |
| $x$ y hlen hood |  |  |  |  | 411 | 12 | 000011100010 |  | 98 | 13 | 0000110010110 |
|  |  |  |  |  | 412 | 13 | 0000110001011 |  |  | 13 | 0000110001010 |
|  | 0 | 1 | 1 |  | 413 | 13 | 0000101111110 |  |  | 13 | 0000110000000 |
|  | 1 | 14 | 0101 |  | 414 | 13 | 0000101101010 |  |  | 14 | 00001011011111 |
|  | 2 | 6 | 001110 |  | 415 | 9 | 000001001 |  | 912 | 13 | 0000101100111 |
|  | 3 | 8 | 00101100 |  | 5 | 9 | 001000010 |  | 913 | 14 | 00001011000110 |
|  | 4 | 9 | 001001010 |  | 5 | 8 | 00011110 |  | 914 | 13 | 0000101100000 |
| 3 | 5 | 9 | 000111111 |  | 52 | 9 | 000111011 |  | 915 | 11 | 00000001011 |
|  | 6 | 10 | 0001101110 |  | 53 | 9 | 000111000 | 10 | 0 | 11 | 00010001011 |
| 3 | 7 | 10 | 0001011101 |  | 54 | 10 | 0001100110 | 10 | 0 | 11 | 00010000001 |
| 3 | 8 | 11 | 00010101100 |  | 55 | 11 | 00010111001 | 10 | 2 | 10 | 0001000011 |
| , | 9 | 11 | 00010010101 |  | 56 | 11 | 00010101101 | 10 | - 3 | 11 | 00001111101 |
| 3 | 10 | 11 | 00010001010 |  | 57 | 12 | 000100001001 | 10 | 0 | 12 | 000011110111 |
| 13 | 11 | 12 | 000011110010 |  | 58 | 11 | 00010001110 | 10 | 5 | 12 | 000011101001 |
| ' 1 | 12 | 12 | 000011100001 |  | 59 | 12 | 000011111101 | 10 | 6 | 12 | 000011100101 |
| 12 | 13 | 12 | 000011000011 |  | 510 | 12 | 000011101000 | 10 | , | 12 | 000011011011 |
| $\begin{aligned} & 19 \\ & 1 \end{aligned}$ | 14 | 13 | 0000101111000 |  | 511 | 13 | 0000110010000 | 10 | , | 13 | 0000110001001 |
| , | 15 | 9 | 000010001 |  | 512 | 13 | 0000110000100 | 10 | 9 | 14 | 00001011100111 |
| 1) | 0 | 3 | 011 |  | 513 | 13 | 0000101111010 | 10 | 10 | 14 | 00001011100001 |
| $\therefore$ | 1 | 4 | 0100 |  | 514 | 14 | 00000110111101 | 10 | 11 | 14 | 00001011010000 |
| $\vdots$ | 2 | 6 | 001100 |  | 515 | 10 | 0000010000 | 10 | 12 | 15 | 000001101110101 |
| : | 3 | 7 | 0010100 |  | 60 | 10 | 0001101111 | 10 | 13 | 15 | 000001101110010 |
| 1 | 4 | 8 | 00100011 | 6 | 6 | 9 | 000110110 | 10 | 14 | 14 | 00000110110111 |
| 1 | 5 | 9 | 000111110 |  | 52 | 9 | 000110100 | 10 | 15 | 10 | 0000000100 |
| 1 | 6 | 9 | 000110101 |  | 63 | 10 | 0001100100 | 11 | 0 | 12 | 000011110011 |
| 1 1 | 7 | 9 | 000101111 |  | 64 | 11 | 00010111000 | 11 | 1 | 11 | 00001111000 |
| $\begin{aligned} & 1 \\ & 1 \end{aligned}$ | 8 | 10 | 0001010011 |  | 5 | 11 | 00010110010 | 11 | 2 | 11 | 00001110110 |
| 1 1 | 9 | 10 | 0001001011 | 6 | 6 | 11 | 00010100000 | 11 | 3 | 11 | 00001110011 |
| 1 3 | 10 | 10 | 0001000100 | 6 | 7 | 11 | 00010000101 | 11 | 4 | 12 | 000011100011 |
| 1 1 | 11 | 11 | 00001110111 | 6 | 8 | 12 | 000100000001 | 11 | 5 | 12 | 000011011111 |
| 1 | 12 | 12 | 000011001001 | 6 | 9 | 12 | 000011110100 | 11 | 5 | 13 | 0000110001100 |
|  | 13 | 11 | 00001101011 | 6 | 10 | 12 | 000011100100 | 11 | 7 | 14 | 00001011101010 |
| 1 | 14 | 12 | 000011001111 | 6 | 11 | 12 | 000011011001 | 11 | 8 | 14 | 00001011100110 |
| 1 1 | 15 | 8 | 00001001 | 6 | 12 | 13 | 0000110000001 | 11 | 9 | 14 | 00001011100000 |
| 1 | 0 | 6 | 001111 | 6 | 13 | 13 | 0000101101110 | 11 | 10 | 14 | 00001011010001 |
| 2 | 1 | 6 | 001101 | 6 | 14 | 14 | 00001011001011 | 11 | 11 | 14 | 00001011001000 |
| 2 | 2 | 7 | 0010111 | 6 | 15 | 10 | 0000001010 | 11 | 12 | 14 | 00001011000010 |
| 2 2 2 | 3 | 8 | 00100110 | 7 | 0 | 10 | 0001100010 | 11 | 13 | 13 | 0000011011111 |
| 2 | 4 | 9 | 001000011 | 7 | 1 | 9 | 000110000 | 11 | 14 | 14 | 00000110110100 |
| 2 | 5 | 9 | 000111010 | 7 | 2 | 10 | 0001011011 | 11 | 15 | 11 | 00000000110 |
| 2 2 | 6 | 10 | 0001100111 | 7 | 3 | 10 | 0001011000 | 12 | 0 | 12 | 000011001010 |
| 2 | 7 | 10 | 0001011010 | 7 | 4 | 11 | 00010100101 | 12 | 1 | 12 | 000011100000 |
| 2 | 8 | 11 | 00010100001 | 7 | 5 | 11 | 00010011101 | 12 | 2 | 12 | 000011011110 |
| 2 | 9 | 10 | 0001001000 | 7 | 6 | 11 | 00010010100 | 12 | 3 | 12 | 000011011010 |
| 2 | 10 | 11 | 00001111111 | 7 | 7 | 12 | 000100000101 | 12 | 4 | 12 | 000011011000 |
| 2 | 11 | 11 | 00001110101 | 7 | 8 | 12 | 000011111000 | 12 | 5 | 13 | 0000110000101 |
| 2 | 12 | 11 | 00002101110 | 7 | 9 | 13 | 0000110010111 | 12 | 6 | 13 | 0000110000010 |
| 2 | 13 | 12 | 000011010001 | 7 | 10 | 13 | 0000110001101 | 12 | 7 | 13 | 0000101111101 |
| 2 | 14 | 12 | 000011001110 | 7 | 11 | 13 | 0000101110100 | 12 | 8 | 13 | 0000101101100 |
| 2 | 15 | 9 | 000010000 | 7 | 12 | 13 | 0000101111100 | 12 | 9 | 15 | 000001101111000 |
| 3 | 0 | 8 | 00101101 | 7 | 13 | 15 | 000001101111001 | 12 | 10 | 14 | 00000110111011 |
| 3 | 1 | 7 | 0010101 | 7 | 14 | 15 | 000001101110100 | 12 | 11 | 14 | 00001011000011 |
|  | 2 | 8 | 00100111 | 7 | 15 | 10 | 0000001000 | 12 | 12 | 14 | 00000110111000 |
|  | 3 | 9 | 001000101 | 8 | 0 | 10 | 0001010101 | 12 | 13 | 14 | 00000110110101 |
| 3 | 4 | 9 | 001000000 | B | 1 | 10 | 0001010100 | 12 | 14 | 16 | 0000011011000000 |
| 3 | 5 | 10 | 0001110010 | 8 | 2 | 10 | 0001010001 | 12 | 15 | 11 | 00000000100 |
| 3 3 3 3 | 6 | 10 | 0001100011 | ${ }^{8}$ |  | 11 | 00010011111 | 13 | 0 | 14 | 00001011101011 |
| 3 3 3 3 | 7 | 10 | 0001010111 | B | 4 | 11 | 00010011100 | 13 | 1 | 12 | 000011010011 |
| 3333 | 8 | 11 | 00010011110 | 8 | 5 | 11 | 00010001111 | 13 | 2 | 12 | 000011010010 |
|  | 9 | 11 | 00010001100 | 8 | 6 | 12 | 000100000100 | 13 | 3 | 12 | 000011010000 |
| 3 3 | 10 | 12 | 000011111100 | 8 | 7 | 12 | 000011111001 | 13 | 4 | 13 | 0000101110010 |
| 3 | 11 | 12 | 000011010100 | 8 | 8 | 13 | 0000110101011 | 13 | 5 | 13 | 0000101111011 |
| 3 3 | 12 | 12 | 000011000111 | 8 | 9 | 13 | 0000110010001 | 13 | 6 | 14 | 00001011011110 |
| 3 | 13 | 13 | 0000110000011 | 8 | 10 | 13 | 0000110001000 | 13 | 7 | 14 | 00001011010011 |
| 3 | 14 | 13 | 0000101101101 | ${ }^{5}$ | 11 | 13 | 0000101111111 | 13 | 8 | 14 | 00001011001010 |
| 3 | 15 | 10 | 0000011010 | E |  | 14 | 00001011010111 | 13 | 9 | 16 | 0000011011000111 |
| 4 | - | 9 | 001001011 | 8 |  | 14 | 00001011001001 | 13 | 10 | 15 | 000001101110011 |
| 4 | 1 | 8 | 00100100 | 8 | 14 | 14 | 00001011000100 | 13 | 11 | 15 | 000001101101101 |
| 4 | 2 | 9 | 001000100 | 旦 | 15 | 10 | 0000000111 | 13 | 12 | 15 | 000001101101100 |
| 4 | 3 |  | 001000001 | 9 | 0 | 11 | 00010011010 | 13 | 13 | 17 | 00000110110000011 |
| 4 | 4 | 10 | 0001110011 | 9 | 1 | 10 | 0001001100 | 13 | 14 | 15 | 000001101100001 |
| 4 | 5 | 10 | 0001100101 | 9 | 2 | 10 | 0001001001 | 13 | 15 | 11 | 00000000010 |
| 4 | 6 | 11 | 00010110011 | 9 | 3 | 11 | 00010001101 | 14 | 0 | 13 | 0000101111001 |


| 14 | 1 | 13 | 0000101110001 |
| ---: | ---: | ---: | :--- |
| 14 | 2 | 11 | 00001100110 |
| 14 | 3 | 12 | 000010111011 |
| 14 | 4 | 14 | 00001011010110 |
| 14 | 5 | 14 | 00001011010010 |
| 14 | 6 | 13 | 0000101100110 |
| 14 | 7 | 14 | 00001011000111 |
| 14 | 8 | 14 | 00001011000101 |
| 14 | 9 | 15 | 000001101100010 |
| 14 | 10 | 16 | 0000011011000110 |
| 14 | 11 | 15 | 000001101100111 |
| 14 | 12 | 17 | 00000110110000010 |
| 14 | 13 | 15 | 000001101100110 |
| 14 | 14 | 14 | 00000110110010 |
| 14 | 15 | 11 | 00000000000 |
| 15 | 0 | 9 | 000001100 |
| 15 | 1 | 8 | 00001010 |
| 15 | 2 | 8 | 00000111 |
| 15 | 3 | 9 | 000001011 |
| 15 | 4 | 9 | 000001010 |
| 15 | 5 | 10 | 0000010001 |
| 15 | 6 | 10 | 0002001011 |
| 15 | 7 | 10 | 0000001001 |
| 15 | 6 | 11 | 00000001101 |
| 15 | 9 | 11 | 00000001100 |
| 15 | 10 | 11 | 00000001010 |
| 15 | 11 | 11 | 00000000111 |
| 15 | 12 | 11 | 00000000101 |
| 15 | 13 | 11 | 00000000011 |
| 15 | 14 | 11 | 00000000001 |
| 15 | 15 | 8 | 00000011 |

Huffman code table 17
same as table 16 , but linbits $=2$
Huffman code table 18
same as table 16, but linbits=3
Huffman code table 19
same as table 16, but linbits=4
Huffman code table 20
same as table 16 , but linbits=6
Huffman code table 21
same as table 16, but linbits $=8$
Huffman code table 22
same as table 16 , but linbits $=10$
Huffman code table 23
same as table 16 , but linbits=13


| $\varsigma$ | 4 | 9 | 010111101 | 141 | 10 | 0100011000 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 9 | 5 | 9 | 010110101 | 142 | 10 | 0100010110 |
| 9 | 6 | 9 | 010101110 | 143 | 10 | 0100010010 |
| 5 | 7 | 10 | 0101001101 | 144 | 10 | 0100001011 |
| 9 | 8 | 10 | 0101000001 | 145 | 10 | 0100001000 |
| 9 | 9 | 10 | 0100110001 | 146 | 10 | 0100000011 |
| 9 | 10 | 10 | 0100100001 | 147 | 11 | 00101111110 |
| 9 | 11 | 10 | 0100010011 | 14 8 | 11 | 00101111010 |
| 9 | 12 | 11 | 01000001001 | 149 | 11 | 00101110100 |
| 9 | 13 | 11 | 00101111011 | 1410 | 11 | 00101101111 |
| 9 | 14 | 11 | 00101110011 | 1411 | 11 | 00101101011 |
| 9 | 15 | 8 | 00001011 | $14 \quad 12$ | 11 | 00101101000 |
| 10 | 0 | 11 | 01010011100 | 1413 | 11 | 00101100110 |
| 10 | 1 | 9 | 010111000 | 1414 | 11 | 00101100100 |
| 10 | 2 | 9 | 010110111 | 1415 | 8 | 00000000 |
| 10 | 3 | 9 | 010110011 | 150 | 8 | 00101011 |
| 10 | 4 | 9 | 010101111 | 151 | 7 | 0010100 |
| 10 | 5 | 10 | 0101011000 | $15 \quad 2$ | 7 | 0010011 |
| 10 | 6 | 10 | 0101001011 | $15 \quad 3$ | 7 | 0010001 |
| 10 | 7 | 10 | 0100111010 | 154 | 7 | 0001111 |
| 10 | 8 | 10 | 0100110000 | 155 | 7 | 0001101 |
| 10 | 9 | 10 | 0100100010 | 156 | 7 | 0001011 |
| 10 | 10 | 10 | 0100010101 | 157 | 7 | 0001001 |
| 10 | 11 | 11 | 01000010010 | 158 | 7 | 0000111 |
| 10 | 12 | 11 | 00101111111 | 159 | 7 | 0000110 |
| 10 | 13 | 11 | 00101110101 | 1510 | 7 | 0000100 |
| 10 | 14 | 11 | 00101101110 | 1511 | 8 | 00000111 |
| 10 | 15 | 8 | 00001010 | 1512 | 8 | 00000101 |
| 11 | 0 | 11 | 01010001100 | 1513 | 8 | 00000011 |
| 11 | 1 | 10 | 0101011010 | 1514 | 8 | 00000001 |
| 11 | 2 | 9 | 010101011 | 1515 | 4 | 0011 |
| 11 | 3 | 9 | 010101000 |  |  |  |
| 11 | 4 | 9 | 010100100 |  |  |  |
| 11 | 5 | 10 | 0100111110 | Hufim | an | code table 25 |
| 11 | 6 | 10 | 0100110101 |  |  |  |
| 11 | 7 | 10 | 0100101011 | same as table 24 , but linbits $=5$ |  |  |
| 11 | 8 | 10 | 0100011111 |  |  |  |
| 11 | 9 | 10 | 0100010100 | Huffman code table 26 |  |  |
| 11 | 10 | 10 | 0100000111 |  |  |  |
| 11 | 11 | 11 | 01000000001 | same as table 24 , but linbits=6 |  |  |
| 11 | 12 | 11 | 00101110111 |  |  |  |
| 11 | 13 | 11 | 00101110000 |  |  |  |
| 11 | 14 | 11 | 00101101010 | Huffman code table 27 |  |  |
| 11 | 15 | 8 | 00000110 |  |  |  |
| 12 | 0 | 11 | 01010001000 | same as table 24 , but linbits $=7$ |  |  |
| 12 | 1 | 10 | 0101000010 |  |  |  |
| 12 | 2 | 10 | 0100111100 | Huffman code table 28 |  |  |
| 12 | 3 | 10 | 0100111000 |  |  |  |
| 12 | 4 | 10 | 0100110011 |  |  |  |
| 12 | 5 | 10 | 0100101110 | same as table 24 , but linbits $=8$ |  |  |
| 12 | 6 | 10 | 0100100100 |  |  |  |
| 12 | 7 | 10 | 0100011100 | Huffman code table 29 |  |  |
| 12 | 8 | 10 | 0100001101 |  |  |  |
| 12 | 9 | 10 | 0100000101 |  |  |  |
| 12 | 10 | 11 | 01000000000 | same as table 24 , but linbits $=9$ |  |  |
| 12 | 11 | 11 | 00101111000 |  |  |  |
| 12 | 12 | 11 | 00101110010 | Huffman code table 30 |  |  |
| 12 | 13 | 11 | 00101101100 |  |  |  |
| 12 | 14 | 11 | 00101100111 |  |  |  |
| 12 | 15 | 8 | 00000100 | same as table 24 , but linbits $=11$ |  |  |
| 13 | 0 | 11 | 01001101100 |  |  |  |
| 13 | 1 | 10 | 0100101100 | Huffman code table 31 |  |  |
| 13 | 2 | 10 | 0100101000 |  |  |  |
| 13 | 3 | 10 | 0100100110 |  |  |  |
| 13 | 4 | 10 | 0100100000 | same as table 24 , but linbits=13 |  |  |
| 13 | 5 | 10 | 0100011010 |  |  |  |
| 13 | 6 | 10 | 0100010001 |  |  |  |
| 13 | 7 | 10 | 0100001010 |  |  |  |
| 13 | 8 | 11 | 01000000011 |  |  |  |
| 13 | 9 | 11 | 00101111100 |  |  |  |
| 13 | 10 | 11 | 00101110110 |  |  |  |
| 13 | 11 | 11 | 00101110001 |  |  |  |
| 13 | 12 | 11 | 00101101101 |  |  |  |
| 13 | 13 | 11 | 00101101001 |  |  |  |
| 13 | 14 | 11 | 00101100101 |  |  |  |
| 13 | 15 | 8 | 00000010 |  |  |  |
| 14 | 0 | 12 | 010000001001 |  |  |  |

## Table B. 8 -- Layer III scalefactor bands

These tables list the width of each scalefactor band. There are 21 bands at each sampling frequency for long (type 0,1 or 3 ) windows and 12 bands each for short windows.

Table B.8a. .- 32 kHz sampling rate
long blocks:

| scale factor band | width of band | index of start | index of end |
| :---: | :---: | :---: | :---: |
| 0 | 4 | 0 | 3 |
| 1 | 4 | 4 | 7 |
| 2 | 4 | 8 | 11 |
| 3 | 4 | 12 | 15 |
| 4 | 4 | 16 | 19 |
| 5 | 4 | 20 | 23 |
| 6 | 6 | 24 | 29 |
| 7 | 6 | 36 | 35 |
| 8 | 8 | 44 | 43 |
| 9 | 10 | 54 | 53 |
| 10 | 12 | 66 | 65 |
| 11 | 16 | 82 | 81 |
| 12 | 20 | 102 | 101 |
| 13 | 24 | 126 | 125 |
| 14 | 30 | 156 | 155 |
| 15 | 38 | 194 | 193 |
| 16 | 46 | 240 | 239 |
| 17 | 56 | 296 | 295 |
| 18 | 68 | 364 | 363 |
| 19 | 84 | 448 | 447 |
| 20 | 102 |  | 549 |

short blocks:

| scale factor band | width of band | index of start | index of end |
| :---: | :---: | :---: | :---: |
| 0 | 4 | 0 | 3 |
| 1 | 4 | 4 | 7 |
| 2 | 4 | 8 | 11 |
| 3 | 4 | 12 | 15 |
| 4 | 6 | 16 | 21 |
| 5 | 8 | 22 | 29 |
| 6 | 12 | 30 | 41 |
| 7 | 16 | 42 | 57 |
| 8 | 20 | 58 | 77 |
| 9 | 26 | 78 | 103 |
| 10 | 34 | 104 | 137 |
| 11 | 42 | 138 | 179 |

Table B.8b. .- $44,1 \mathrm{kHz}$ sampling rate
long blocks:

| scale factor band | width of band | index of start | index of end |
| :---: | :---: | :---: | :---: |
| 0 | 4 | 0 | 3 |
| 1 | 4 | 4 | 7 |
| 2 | 4 | 8 | 11 |
| 3 | 4 | 12 | 15 |
| 4 | 4 | 16 | 19 |
| 5 | 4 | 20 | 23 |
| 6 | 6 | 24 | 29 |
| 7 | 6 | 36 | 35 |
| 8 | 8 | 44 | 43 |
| 9 | 8 | 52 | 51 |
| 10 | 10 | 62 | 61 |
| 11 | 12 | 74 | 73 |
| 12 | 16 | 90 | 89 |
| 14 | 20 | 110 | 109 |
| 15 | 24 | 134 | 133 |
| 16 | 28 | 196 | 161 |
| 17 | 34 | 238 | 195 |
| 18 | 42 | 288 | 237 |
| 19 | 50 | 342 | 287 |
| 20 | 54 |  | 341 |

short blocks:

| scale factor band | width of band | index of start | index of end |
| :---: | :---: | :---: | :---: |
| 0 | 4 | 0 | 3 |
| 1 | 4 | 4 | 7 |
| 2 | 4 | 8 | 11 |
| 3 | 4 | 12 | 15 |
| 4 | 6 | 16 | 21 |
| 5 | 8 | 22 | 29 |
| 6 | 10 | 30 | 39 |
| 7 | 12 | 40 | 51 |
| 8 | 14 | 62 | 65 |
| 9 | 18 | 84 | 83 |
| 10 | 22 | 106 | 105 |
| 11 | 30 | 135 |  |

Table B.8c. -- 48 kHz sampling rate
long blocks:

| scale factor band | width of band | index of start | index of end |
| :---: | :---: | :---: | :---: |
| 0 | 4 | 0 | 3 |
| 1 | 4 | 4 | 7 |
| 2 | 4 | 8 | 11 |
| 3 | 4 | 12 | 15 |
| 2 | 4 | 16 | 19 |
| 5 | 4 | 20 | 23 |
| 6 | 6 | 24 | 29 |
| 7 | 6 | 30 | 35 |
| 8 | 6 | 46 | 41 |
| 9 | 8 | 50 | 49 |
| 10 | 10 | 60 | 59 |
| 11 | 12 | 72 | 71 |
| 12 | 16 | 88 | 87 |
| 13 | 18 | 106 | 105 |
| 14 | 22 | 128 | 127 |
| 15 | 28 | 156 | 155 |
| 16 | 34 | 190 | 189 |
| 17 | 40 | 230 | 229 |
| 18 | 46 | 276 | 275 |
| 19 | 54 | 330 | 329 |
| 20 | 54 |  | 383 |

short blocks:

| scale fac or band | width of band | index of start | index of end |
| :---: | :---: | :---: | :---: |
| 0 | 4 | 0 | 3 |
| 1 | 4 | 4 | 7 |
| 2 | 4 | 8 | 11 |
| 3 | 4 | 12 | 15 |
| 4 | 6 | 16 | 21 |
| 5 | 6 | 22 | 27 |
| 6 | 10 | 28 | 37 |
| 7 | 12 | 38 | 49 |
| 8 | 14 | 60 | 63 |
| 9 | 16 | 80 | 79 |
| 10 | 20 | 100 | 99 |
| 11 | 26 |  | 125 |

Table B.9 -- Layer III coefficients for aliasing reduction:

| (i) | $\mathbf{c}_{\mathbf{i}}$ |
| :--- | :--- |
| 0 | $-0,6$ |
| $\mathbf{1}$ | $-0,535$ |
| $\mathbf{2}$ | $-0,33$ |
| 3 | $-0,185$ |
| 4 | $-0,095$ |
| 5 | $-0,041$ |
| 6 | $-0,0142$ |
| 7 | $-0,0037$ |

The butterfly coefficients $\mathrm{cs}_{\mathrm{i}}$ and $\mathrm{ca}_{\mathrm{i}}$ are calculated as follows:

$$
\mathrm{cs}_{i}=\frac{1}{\sqrt{1+c_{i}^{2}}}, \mathrm{ca}_{i}=\frac{\mathrm{c}_{\mathrm{i}}}{\sqrt{1+\mathrm{c}_{\mathrm{i}}^{2}}}
$$

# Annex C 

(informative)

## The encoding process

## C. 1 Encoider

## C.1.1 Overview

For each of the layers, an example of one suitable encoder with the corresponding flow-diagram is given in this annex. In subsequent clauses the analysis subband filter and the layer-specific encoding techniques are described. In annex D two examples of psychoacoustic models, which are common to all layers, are described. A short introduction describes the overall philosophy.

## C.1.1.1 Introduction

The ISO/IEC 11172-3 (MPEG-Audio) algorithm is a psychoacoustic algorithm. The figure C. 1 shows the primary parts of a psychoacoustic algorithm.


Figure C. 1 .- ISO/IEC 11172-3 (MPEG-audio) encoder block diagram
The four primary parts of the psychoacoustic encoder are:

## C.1.1.1.1 The filterbank

The filterbank cloes a time to frequency mapping. There are two filterbanks used in the ISO/IEC 11172-3 (MPEG-Audio) algorithm, a polyphase filterbank and a hybrid polyphase/MDCT filterbank. Each provides a specific mapping in time and frequency. These filterbanks are critically sampled (i.e. there are as many samples in the analyzed domain as there are in the time domain). These filterbanks provide the primary frequency separation for the encoder, and the reconstruction filters for the decoder. The output samples of the filterbank are quantized.

## C.1.1.1.2 The psychoacoustic model

The psychoacoustic model calculates a just noticable noise-level for each band in the filterbank. This noise level is used in the bit or noise allocation to determine the actual quantizers and quantizer levels. There are two psychoacoustic models presented in annex D. While they can both be applied to any layer of the ISO/IEC 11172-3 (MPEG-Audio) algorithm, in practice Model 1 has been used for Layers I and II, and Model 2 for Layer III. In both psychoacoustic models, the final output of the model is a signal-to-mask ratio (SMR) for each band (Layers I and II) or group of bands (Layer III).

## C.1.1.1.3 Bit or noise Allocation

The allocator looks at both the output samples from the filterbank and the SMR's from the psychoacousic model, and adjusts the bit allocation (Layers I and II) or noise allocation (Layer III) in order simultaneously to) meet both the bitrate requirements and the masking requirements. At low bitrates, these methods atternpt to spend bits in a fashion that is psychoacousticly inoffensive when they cannot meet the psychoacoustic demand at the required bitrate.

## C.1.1.1.4 The bitstream formatter

The bitstream formatter takes the quantized filterbank outputs, together with the bit allocation (Layers I and II) or noise allocation (Layer III) and other required side information, and encodes and formats that information in an efficient fashion. In the case of Layer III, the Huffman codes are also inserted at this point.

## C.1.1.2 The filterbank

In Layers I and II, a filterbank with 32 subbands is used. In each subband, 12 or 36 samples are grouped for processing. In Layer III, the filterbank has a signal-dependent resolution, where there are either $6 \times 32$ or $18 \times 32$ frequency bands. In the case where there are $6 \times 32$ frequency samples, the 3 sets of each frequency are quantized separately.

## C.1.1.3 Bit or noise allocation method

There are two different bitrate control methods explained in this annex. In Layers I and II this method is a bit allocation process, i.e. a number of bits is assigned to each sample (or group of samples) in each subband. The method for Layer III is a noise-allocation loop, where the quantizers are varied in an organized fashion, and the variable to be controlled is actually the injected noise. In either case, the result is a set ol quantization parameters and quantized output samples that are given to the bitstream formatter.

## C.1.1.4 Bitstream formatting

The bitstream formatter varies from layer to layer. In Layers I and II, a fixed PCM code is used for each subband sample, with the exception that in Layer II quantized samples may be grouped. In Layer III, Huffman codes are used to represent the quantized frequency samples. These Huffman codes are variablelength codes that allow for more efficient bitstream representation of the quantized samples at the cost of additional complexity.

## C.1.2 Input high-pass filter

The encoding algorithms provide a frequency response down to d.c. However, in applications where this is not a requirement, it is recommended that a high-pass filter be included at the input of the encoder. The cutoff frequency should be in the range of 2 to 10 Hz .

The application of such a high-pass filter avoids an unneccessarily high bitrate requirement for the lowest subband and increases the overall audio quality.

## C.1.3 Analysis subband filter

Ar analysis subband filterbank is used to split the broadband signal with sampling frequency $\mathrm{f}_{\mathrm{s}}$ into 32 equally spaced subbands with sampling frequencies $\mathrm{f}_{\mathrm{S}} / 32$. The flow chart of this process with the appropriate formulas is given in figure C. 4 "Analysis Subband Filter Flow Chart". The analysis subband filtering includes the following steps:

- Input 32 audio samples.
- Build an input sample vector X of 512 elements. The 32 audio samples are shifted in at positions 0 t) 31 , the most recent one at position 0 , and the 32 oldest elements are shifted out.
- Window vector X by vector C . The coefficients are to be found in table C.1.
- Calculate the 64 values $Y_{i}$ according to the formula given in the flow chart.
- Calculate the 32 subband samples $\mathrm{S}_{\mathrm{i}}$ by matrixing. The coefficients for the matrix can be calculated by the following formula:

$$
M_{i k}=\cos [(2 i+1)(k-16) \pi / 64], \quad \text { for } i=0 \text { to } 31 \text {, and } k=0 \text { to } 63 .
$$

[^2]Table C. 1 .- Coefficients $C_{i}$ of the Analysis Window

| C[ $\quad 0]=0,000000000$ | C[ 1] 1 - 0,000000477 | C[ 2$]=-0,000000477$ | C[ $31=-0,000000477$ |
| :---: | :---: | :---: | :---: |
| C[ 4] $\quad$ - 0, 000000477 | C[ 5] $=-0,000000477$ | ci 6] $=-0,000000477$ | c] $71=-0,000000954$ |
|  | Cl $\quad 91=-0,000000954$ | c $[10]=-0,000000954$ | c 111$]=-0,000001431$ |
| $C[12]=-0,000001431$ | C[ $131=-0,000001907$ | C[ 14] $=-0,000001907$ | c( 15] $=-0,000002384$ |
| C[ 16] $=-0,000002384$ | C( 17] $=-0,000002861$ | C $[18]=-0,000003338$ | C $[19]=-0,000003338$ |
| C[ 20] $=-0,000003815$ | C( 21] $=-0,000004292$ | c\| 221=-0,000004768 | c( 23) $=-0,000005245$ |
| c! 24] $=-0,000006199$ | c[ 25$]=-0,000006676$ | C[ 26$]=-0,000007629$ | C\| $271=-0,000008106$ |
| C $\ 281=-0,000009060$ | C[ 29] $=-0,000010014$ | C\| $301=-0,000011444$ | C [ 31] $=-0,000012398$ |
| C( 32) $=-0,0 \subset 0013828$ | $C[33]=-0,000014782$ | ci 34$]=-0,000016689$ | c] 35$]=-0,000018120$ |
| C[ 36] $=-0$, oc 0019550 | $C[37]=-0,000021458$ | c $(38)=-0,000023365$ | c] 39$]=-0,000025272$ |
| C[ 40$]=-0,0 C 0027657$ | C[ 41$]=-0,000030041$ | C $(42]=-0,000032425$ | C [ 43] $=-0,000034809$ |
| C[ 44] $=-0$, OC 0037670 | C $\mid 45]=-0,000040531$ | c $[46]=-0,000043392$ | c\| 47$]=-0,000046253$ |
| C[ 48$]=-0,0 C 0049591$ | C\| 49$]=-0,000052929$ | C [ 50$]=-0,000055790$ | c) 51$]=-0,000059605$ |
| C[ 52$]=-0,000062943$ | C\| $531=-0,000066280$ | C $1541=-0,000070095$ | ci 55$]=-0,000073433$ |
| C[ 56] $=-0,000076771$ | cl $571=-0,000080585$ | C $[58]=-0,000083923$ | c\| 59] $=-0,000087261$ |
| C[ 60] $=-0,000090599$ | c) $611=-0,000093460$ | C[ 62] $]=-0,000096321$ | C\| 63$]=-0,000099182$ |
| CI 64] $=0,000101566$ | c\| $651=0,000103951$ | c $[66]=0,000105358$ | c\| $671=0,000107288$ |
| C $\mid 681=0,000108242$ | ci $691=0,000108719$ | c $[70]=0,000108719$ | c $711=0,000108242$ |
| c $(72)=0,000105812$ | ci $731=0,000105381$ | c\| $741=0,000102520$ | $C(75)=0,000099182$ |
| c\| 76] $=0,000095367$ | $\mathrm{C}(77)=0,000090122$ | c\| $781=0,000084400$ | c $[79]=0,000077724$ |
| Cl 80) $=0,000069618$ | c $[81]=0,000060558$ | c\| 82$]=0,000050545$ | $c(83)=0,000039577$ |
| C\| 84 $1=0.000027180$ | $c(85)=0,000013828$ | cl 86$]=-0,000000954$ | c $\|87\|=-0,000017166$ |
| C! 88] $=-0,000034332$ | C[ 89$]=-0,000052929$ | C\| 90] $=-0,000072955$ | c [ 91$]=-0,000093937$ |
| C! 92] $=-0,000116348$ | C [ 93] $=-0,000140190$ | c\| 94 ] $=-0,000165462$ | c\| 951=-0,000191212 |
| C! 96] $=-0,000218868$ | C[ 97] $=-0,000247478$ | C 198$]=-0,000277042$ | ci $991=-0,000307560$ |
| $\mathrm{C}\{100]=-0,000339031$ | $C[101]=-0,000371456$ | $c \mid 102]=-0,000404353$ | $C[103]=-0,000438213$ |
| $C[104]=-0,000472546$ | $C[105]=-0,000507355$ | $C[106]=-0,000542164$ | $C\{107\}=-0,000576973$ |
| $\mathrm{C}[108]=-0,000611782$ | $C[109]=-0,000646591$ | $C(110)=-0,000680923$ | $c\{111]=-0,000714302$ |
| $C[112]=-0,000747204$ | $c \mid 113)=-0,000779152$ | $C[114]=-0,000809669$ | $c[115]=-0,000838757$ |
| $c[116]=-0,000866413$ | $C\|117\|=-0,000891685$ | $c\|118\|=-0,000915051$ | $C(119)=-0,000935555$ |
| $C[120]=-0,000954151$ | C $\|121\|=-0,000968933$ | $\mathrm{C}\|122\|=-0,000980854$ | $C(123)=-0,000989437$ |
| $C[124]=-0.000994205$ | C\| $1251=-0,000995159$ | $C[126)=-0,000991821$ | $c(127)=-0,000983715$ |
| $C[128]=0,000971317$ | $c \mid 129]=0,000953674$ | $C[130]=0,000930786$ | $c[131)=0,000902653$ |
| $C[132]=0,000868797$ | C[133] $=0,000829220$ | $C\{134)=0,000783920$ | $c(135)=0,000731945$ |
| $C[136]=0,000674248$ | $\mathrm{c}\|137\|=0,000610352$ | C[138\|-0,000539303 | $C[139)=0,000462532$ |
| $C[140)=0,000378609$ | $c\|141\|=0,000288436$ | $C[142 \mid=0,000191689$ | $C(143)=0,000088215$ |
| $C[144]=-0,000021458$ | $C[145)=-0,000137329$ | $C[146 \mid=-0,000259876$ | $C[147)=-0,000388145$ |
| $C[148]=-0,000522137$ | $c \mid 149)=-0,000661850$ | C $[1501=-0,000806808$ | $c(151)=-0,000956535$ |
| $C[152]=-0,001111031$ | C $\mid 1531=-0,001269817$ | C[154] $=-0,001432419$ | $C[155)=-0,001597881$ |
| $C[156]=-0,001766682$ | c\|157] $=-0,001937389$ | C11581=-0,002110004 | $C(159)=-0,002283096$ |
| $C[160]=-0,002457142$ | $\mathrm{c} \mid 1611=-0,002630711$ | $c \mid 162]=-0,002803326$ | $C[163)=-0,002974033$ |
| $c[164]=-0,003141880$ | c\| $1651=-0,003306865$ | $C[1661=-0,003467083$ | $C(167)=-0,003622532$ |
| $C[168]=-0,003771782$ | $C \mid 1691=-0.003914356$ | $C[170]=-0,004048824$ | $c(171)=-0.004174709$ |
| $C[172]=-0,004290581$ | c $11731=-0,004395962$ | $c[174]=-0,004489899$ | $C[175)=-0,004570484$ |
| $C[1761=-0,004638195$ | c\|177]=-0,004691124 | C $[1781=-0,004728317$ | c $\mid 179\}=-0,004748821$ |
| $C[180]=-0,004752159$ | $c \mid 181]=-0,004737377$ | $\mathrm{c}[182]=-0,004703045$ | $C[183)=-0,004649162$ |
| $C[184]=-0,004573822$ | $C \mid 185]=-0,004477024$ | $C[186]=-0,004357815$ | $C[187)=-0,004215240$ |
| $c[188]=-0,004049301$ | C\|189] $=-0,003858566$ | C $\{190 \mid=-0,003643036$ | $C[191)=-0,003401756$ |
| $C[192]=0,003134727$ | $\mathrm{C} \mid 193]=0,002841473$ | $c[194]=0,002521515$ | $c(195)=0,002174854$ |
| $C[196]=0,001800537$ | $c \mid 1971=0,001399517$ | $C[198]=0,000971317$ | $c(199)=0,000515938$ |
| $C[200]=0,000033379$ | C $\mid 2011=-0,000475883$ | $C[202]=-0,001011848$ | $c(203)=-0,001573563$ |
| $C[204]=-0,002162503$ | $C \mid 205)=-0,002774239$ | $C[206 \mid=-0,003411293$ | $C(207)=-0,004072189$ |
| $C[208]=-0,004756451$ | $c \mid 209)=-0,005462170$ | $C[210]=-0,006189346$ | $\mathrm{C}(211)=-0,006937027$ |
| $C[212]=-0,007703304$ | $c \mid 213]=-0,008487225$ | $C \mid 214]=-0,009287834$ | $C[215]=-0,010103703$ |
| $C[216]=-0,010933399$ | $c[217]=-0,011775017$ | $c(218)=-0,012627602$ | $c[219]=-0,013489246$ |
| $c[220]=-0,014358521$ | $c(221)=-0,015233517$ | $c(222)=-0,016112804$ | $c[223]=-0,016994476$ |
| $C[224]=-0,017876148$ | $C(225)=-0,018756866$ | $C[226]=-0,019634247$ | $c[227]=-0,020506859$ |
| $\mathrm{C} \mid 22 \mathrm{~B}]=-0.021372318$ | $C(229)=-0.022228718$ | $c \mid 230]=-0,023074150$ | $C[231]=-0,023907185$ |
| $C[232]=-0,024725437$ | $\mathrm{C}[233)=-0,025527000$ | $c[234]=-0,026310921$ | $c(235)=-0,027073860$ |
| $C[236]=-0,027815342$ | $C(237)=-0,028532982$ | $C[238]=-0,029224873$ | $c(239)=-0,029890060$ |
| $\mathrm{C}(240)=-0,030526638$ | $\mathrm{C} \mid 241]=-0,031132698$ | $C[242]=-0.031706810$ | $\mathrm{c}[243]=-0,032248020$ |
| $c\{244\}=-0,032754898$ | $C(245)=-0,033225536$ | $C(246) \mathbf{n - 0 , 0 3 3 6 5 9 9 3 5}$ | C[247] $=-0,034055710$ |
| $\mathrm{c}\|248\|=-0,034412861$ | $\mathrm{c} \mid 249]=-0,034730434$ | $c(250)=-0,035007000$ | c\|251] $=-0,035242081$ |
| $c\|252\|=-0,035435200$ | $c[253]=-0,035586357$ | cโ2541 $=-0,035894122$ | $c[255)=-0,035758972$ |
| $c\|256\|=0,035730907$ | $c[257]=0,035758972$ | $C 12581=0,035694122$ | $c(259)=0,035585357$ |
| $c \mid 2601=0,035435200$ | c[261] $=0,035242061$ | $\tau[262]=0,035007000$ | $\mathrm{C}(253)=0.034730434$ |
| $c[264]=0.034412861$ | $c[265]=0,034055710$ | $c[255]=0.033659935$ | $c(257)=0,033225536$ |
| $C[268]=0,032754898$ | $c[269]=0,032248020$ | $c[270)=0.031706810$ | $C(271)=0.031132698$ |
| $\mathrm{c}[272]=0,020526638$ | c[273] $=0,029890060$ | $c[274]=0.029224873$ | $c(275)=0,028532982$ |
| $c[276]=0,027815342$ | $\mathrm{c}\{277\}=0,027073860$ | $C[278]=0,026310921$ | $C[279]=0,025527000$ |
| $C[2 B D]=0,024725437$ | $c\{2 B 1\}=0,023907185$ | $c[282]=0,023074150$ | $C[283)=0,022228718$ |
| $\mathrm{c}[284]=0,021372318$ | c $28851=0.020506859$ | $C[286]=0.019634247$ | $C[287]=0,018756866$ |
| $C(286)=0,017876148$ | $c[289]=0.016994476$ | $C\{290]=0,016112804$ | $c[291]=0,015233517$ |


| $C[292]=0,014358521$ | $C[293]=0,013489246$ | $C(294)=0,012627602$ | $C[295]=0,011775017$ |
| :---: | :---: | :---: | :---: |
| $C[296]=0,010933399$ | $C[297]=0,010103703$ | $C(298)=0,009287834$ | $c[299]=0,008487225$ |
| $C[300]=0,007703304$ | $C[301]=0,005937027$ | $C(302)=0,006189346$ | $\mathrm{c}[303]=0,005462170$ |
| $C[304]=0,004756451$ | $C[305]=0,004072189$ | $C[306]=0,003411293$ | $\mathrm{c}[307]=0,002774239$ |
| $C[308]=0,002161503$ | $c(309)=0,001573563$ | $\mathrm{C}[310]=0,001011848$ | $\mathrm{c}[311]=0,000475883$ |
| $C[312]=-0,000033379$ | $C\{313\}=-0,000515938$ | $c[314)=-0,000971317$ | $\mathrm{c}[315]=-0,001399517$ |
| $C[316]=-0,001800537$ | $C(317)=-0,002174854$ | $C[318]=-0,002521515$ | $C[319)=-0,002841473$ |
| $C[320]=0,003134727$ | $C[321]=0,003401756$ | $C[322]=0,003643036$ | $C[323]=0,003858566$ |
| $\mathrm{C}(324)=0,004049301$ | $C[325]=0,004215240$ | $C[326]=0,004357815$ | $c[327]=0,004477024$ |
| $c[328]=0,004573822$ | $C[329]=0,004649162$ | $C[330]=0,004703045$ | $c[331)=0,004737377$ |
| $C[332]=0,004752159$ | $C[333]=0,004748821$ | $\mathrm{c}[334)=0,004728317$ | $\mathrm{c}[335]=0,004691124$ |
| $C[336]=0,004638195$ | $C[337]=0,004570484$ | $\mathrm{c}[33 \mathrm{~B}]=0,004489899$ | $c[339]=0,004395962$ |
| $C[340]=0,004290581$ | $C[341)=0,004174709$ | $C[342]=0,004048824$ | $\mathrm{C}[343]=0,003914356$ |
| $C(344)=0,003771782$ | $C[345]=0,003622532$ | $c(345)=0,003467083$ | $\mathrm{c}[3471=0,003306866$ |
| $C[348]=0,003141880$ | $C[349]=0,002974033$ | $c[350]=0,002803326$ | $c[351]=0,002630711$ |
| $C[352]=0,002457142$ | $C[353]=0,002283096$ | $\mathrm{c}(354)=0,002110004$ | $C[355]=0,001937389$ |
| $C[356]=0,001766682$ | $C[357)=0,001597881$ | $C(358)=0,001432419$ | $c[359]=0,001269817$ |
| $C[360]=0,001111031$ | $c[361]=0,000956535$ | $C(362)=0,000806808$ | $C[363]=0,000661850$ |
| $\mathrm{c}[364]=0,000522137$ | $c(365)=0,000388145$ | $c[366]=0,000259876$ | $c(367)=0,000137329$ |
| $c\{368\}=0,000021458$ | $C\{369\}=-0,000083215$ | $c[370]=-0,000191689$ | $c[371]=-0,000288486$ |
| $\mathrm{c} \mid 3721=-0,000378609$ | $\mathrm{C}(373)=-0,000462532$ | $C[374)=-0,000539303$ | $C(375)=-0,000610352$ |
| $C(376]=-0,000674248$ | $\mathrm{c}[377]=-0,000731945$ | $C[378]=-0,000783920$ | $C[379)=-0,000829220$ |
| $C[380]=-0,000868797$ | $C[381]=-0,000902653$ | $C[382]=-0.000930786$ | $C(383)=-0.000953674$ |
| $\mathrm{C} \mid 3841=0,000971317$ | $c[385]=0,000983715$ | $c[386)=0,000991821$ | $c(387)=0,000995159$ |
| $C(388)=0,000994205$ | $C(389)=0,000989437$ | $c[390]=0,000980854$ | $C(391)=0,000968933$ |
| $C[392]=0,000954151$ | $c[393]=0,000935555$ | $c\|394\|=0,000915051$ | $C[395]=0,000891685$ |
| $c[396]=0,000866413$ | $C[397]=0,000839757$ | $C[398)=0,000809669$ | $C[399]=0,000779152$ |
| $C(400)=0,000747204$ | $C[401]=0,000714302$ | $C[402]=0,000680923$ | $C[403]=0,000646591$ |
| $C(404)=0,000611782$ | $C[405]=0,000576973$ | $C \mid 406)=0,000542164$ | $C[407]=0,000507355$ |
| $C[408]=0,000472546$ | $C[409]=0,000438213$ | $C[410]=0,000404358$ | $c \mid 411]=0,000371456$ |
| $C(412)=0,000339031$ | $c[413]=0,000307560$ | $c \mid 414]=0,000277042$ | $c[415]=0,000247478$ |
| $C[416]=0,000218868$ | $\mathrm{c}\|417\|=0,000191212$ | $C \mid 418]=0,000165462$ | $c \mid 419]=0,000140190$ |
| C[420) 0,000116348 | $c\{421)=0,000093937$ | $C[422]=0,000072956$ | $C \mid 423]=0,000052929$ |
| $C[424]=0,000034332$ | $c\{425\}=0,000017166$ | $c[426]=0,000000954$ | $\mathrm{c} \mid 427]=-0,000013828$ |
| $C[428]=-0,000027180$ | $c \mid 429]=-0,000039577$ | $C[430]=-0,000050545$ | $C\|431\|=-0,000060558$ |
| $C[432]=-0,000069618$ | $c(433)=-0,000077724$ | $c[434]=-0,000084400$ | $c(435]=-0,000090122$ |
| $C[436]=-0,000095367$ | $C[437]=-0,000099182$ | $C[438]=-0,000102520$ | $C[439]=-0,000105381$ |
| $C[440]=-0,000106812$ | $c \mid 441)=-0,000108242$ | $C[442]=-0,000108719$ | $\mathrm{C}\{443]=-0,000108719$ |
| $c[444]=-0,000108242$ | $c[445]=-0,000107288$ | $C[446]=-0,000105858$ | $C(447)=-0,000103951$ |
| $c[448]=0,000101566$ | $c(449]=0,000099182$ | C[450]口 0,000096321 | $C(451)=0,000093460$ |
| $C[452]=0.000090599$ | $C\{453\}=0,000087261$ | $C[454]=0,000083923$ | $C[455]=0,000080585$ |
| $c[456]=0.000076771$ | $\mathrm{c} \mid 457\}=0,000073433$ | $C[458]=0,000070095$ | $c(459)=0,000066280$ |
| $C[460 \mid=0,000062943$ | c $\{461\}=0,000059605$ | $c(462)=0,000055790$ | $c(463)=0,000052929$ |
| $C[464]=0,000049591$ | $c(465)=0,000046253$ | $c[466]=0,000043392$ | $c(467)=0,000040531$ |
| $c[468]=0,000037670$ | $C \mid 469\}=0,000034809$ | $C[470]=0.000032425$ | $C(471)=0,000030041$ |
| $C[472]=0,000027657$ | $C[473)=0,000025272$ | $C[474]=0,000023365$ | $C \mid 475\}=0,000021453$ |
| $C[476]=0,000019550$ | $c(477)=0.000018120$ | $C[478)=0,000016689$ | $c(479)=0,000014782$ |
| $C[480]=0,000013828$ | $c \mid 481]=0,000012398$ | $C[482]=0,000011444$ | $C[483)=0,000010014$ |
| $C[484\}=0,000009060$ | $c \mid 485]=0,000008106$ | $C(486)=0,000007629$ | $C[487]=0,000006676$ |
| $C[488]=0,000006199$ | $c[489]=0,000005245$ | $C[490)=0,000004768$ | $C[491]=0,000004292$ |
| c[492] $=0,000003815$ | $c(493)=0,000003338$ | $C[494]=0,000003338$ | $C[495]=0,000002861$ |
| $C[496]=0,000002384$ | $\mathrm{c}\|497\|=0,000002384$ | $C[498]=0,000001907$ | $c(499)=0,000001907$ |
| $c[500]=0,000001431$ | $c\|501\|=0,000001431$ | $C[502]=0,000000954$ | $C[503]=0,000000954$ |
| $c[504]=0,000000954$ | $c \mid 505]=0,000000954$ | $C[506]=0,000000477$ | $C(507)=0.000000477$ |
| $C[508)=0,000000477$ | $C \mid 509]=0,000000477$ | $C[510]=0,000000477$ | $\mathrm{C}[511]=0,000000477$ |

## C.1.4 Psychoacoustic models

Two examples of psychoacoustic models are presented in annex D, "Psychoacoustic models".

## C.1.5 Encoding

## C.1.5.1 Layer I encoding

## C.1.5.1.1 Introduction

This clause describes a possible Layer I encoding method. The description is made with reference to figure C.5, "Layer I, II Encoder Flow Chart".

## C.1.5.1.2 Psychoacoustic model

The calculation of the psychoacoustic parameters can be done either with Psychoacoustic Model 1 described in clause D.1. or with Psychoacoustic Model 2 as described in D.2. The FFT shiftlength equals 384 samples. Either model provides the signal-to-mask ratio for every subband.

## C.1.5.1.3 Analysis subband filtering

The subband analysis is described in the clause C.1.3, "Analysis subband filter".

## C.1.5.1.4 Scalefactor calculation

The calculation of the scalefactor for each subband is performed every 12 subband samples. The maximum of the absolute value of these 12 samples is determined. The lowest value in table B.1, "Layer I, II Scalefactors", which is larger than this maximum is used as the scalefactor.

## C.1.5.1.5 Coding of scalefactors

The index in the table B.1, "Layer I, II Scalefactors" is represented by 6 bits, MSB first. The scalefactor is transmitted only if a non-zero number of bits has been allocated to the subband.

## C.1.5.1.6 Bit allocation

Before adjustment to a fixed bitrate, the number of bits that are available for coding the samples and the scalefactors must be determined. This number can be obtained by subtracting from the total number of bits available "cb", the number of bits needed for the header "bhdr" ( 32 bits), the CRC checkword "bcre" if used ( 16 bits), the bit allocation "bbal", and the number of bits required for ancillary data "banc":
$\mathrm{adb}=\mathrm{cb}-(\mathrm{bhdr}+\mathrm{bcrc}+\mathrm{bbal}+\mathrm{banc})$
The resulting number of bits can be used to code the subband samples and the scalefactors. The principle used in the allocation procedure is minimization of the total noise-to-mask ratio over the frame with the constraint that the number of bits used does not exceed the number of bits available for that frame. The possible number of bits allocated to one sample can be found in the table in 2.4.2.5 of the main part of the audio standard (Audio data, Layer I); the range is $0 . . .15$ bits, excluding an allocation of 1 bit.

The allocation procedure is an iterative procedure, where in each iteration step the number of levels of the subband samples of greatest benefit is increased.

First the mask-to-noise ratio "MNR" for each subband is calculated by subtracting from the signal-to-noiseratio "SNR" the signal-to-mask-ratio "SMR":
$\mathrm{MNR}=\mathrm{SNR}-\mathrm{SMR}$
The signal-to-noise-ratio can be found in the table C.2, "Layer I Signal-to-Noise Ratio". The signal-to-mask-ratio is the output of the psychoacoustic model.

Then zero bits are allocated to the samples and the scalefactors. The number of bits for the samples "bspl " arid the number of bits for the scalefactors "bscf" are set to zero. Next an iterative procedure is started. Eas:h iteration loop contains the following steps :

- Determination of the minimal MNR of all subbands.
- The accuracy of the quantization of the subband with the minimal MNR is increased by using the next higher number of bits.
- The new MNR of this subband is calculated.
- bspl is updated according to the additional number of bits required. If a non-zero number of bits is assigned to a subband for the first time, bscf has to be incremented by 6 bits. Then adb is calculated again using the formula:

$$
\mathrm{adb}=\mathrm{cb}-(\mathrm{bhdr}+\mathrm{bcrc}+\mathrm{bbal}+\mathrm{bscf}+\mathrm{bspl}+\mathrm{banc})
$$

The iterative procedure is repeated as long as adb is not less than any possible increase of bspl and bscf within one loop.

## C.1.5.1.7 Quantization and encoding of subband samples

A linear quantizer with a symmetric zero representation is used to quantize the subband samples. This representation prevents small value changes around zero from quantizing to different levels. Each of the suibband samples is normalized by dividing its value by the scalefactor to obtain $X$, and quantized using the following formula :

- Calculate $\mathrm{AX}+\mathrm{B}$
- Take the N most significant bits.
- Invert the MSB.

A and B can be found in table C.3, "Layer I Quantization Coefficients". N represents the necessary numbe of bits to encode the number of steps. The inversion of the most significant bit (MSB) is done in order to avoid the all ' 1 ' representation of the code, because the all ' 1 ' code is used for the synchronization word.

## C.1.5.1.8 Coding of bit allocation

The 4-bit code for the allocation is given in 2.4.2.5, "Audio data Layer I", of the main part of the audio standard.

## C.1.5.1.9 Ancillary data

The Audio standard provides a number of bits for the inclusion and transmission of variable length ancillay $y$ dala with the audio bitstream. The ancillary data will reduce the number of bits available for audio, which may result in a degradation of audio quality.

The presence of a bit pattern in the ancillary data matching the syncword may hamper synchronization. Tbis problem is more likely to occur when the free format is used.

## C.1.5.1.10 Formatting

The encoded subband information is transferred in frames (See also 2.4.1.2, 2.4.1.3, 2.4.1.5 and 2.4.1.8). The number of slots in a frame varies with the sample frequency (Fs) and bitrate. Each frame contains information on 384 samples of the original input signal, so the frame rate is Fs/384.

| Fs (kHz) | Frame size <br> $(\mathrm{ms})$ |
| :---: | :---: |
| 48 | 8 |
| 44,1 | $8,7074 \ldots$ |
| 32 | 12 |

A frame may carry audio information from one or two channels.

The length of a slot in Layer $I$ is 32 bits. The number of slots in a frame can be computed by this formula:
Number of slots/frame $(N)=\frac{\text { bitrate }}{\text { Fs }} * 12$
If this does not give an integer number the result is truncated and 'padding' is required. This means that the number of slots may vary between N and $\mathrm{N}+1$.

An overview of the Layer I format is given in figure C.2:


Figure C. 2 -- Layer I Format
Table C. 2 -. Layer I Signal-to-Noise Ratios

| No. of steps | SNR (dB) |
| :---: | :---: |
| 0 | 0,00 |
| 3 | 7,00 |
| 7 | 16,00 |
| 15 | 25,28 |
| 31 | 31,59 |
| 63 | 37,75 |
| 127 | 43,84 |
| 255 | 49,89 |
| 511 | 55,93 |
| 1023 | 61,96 |
| 2047 | 67,98 |
| 4095 | 74,01 |
| 8191 | 80,03 |
| 16383 | 96,05 |
| 32767 |  |

Table -- C. 3 Layer I Quantization Coefficients

| No. of steps | $\mathbf{A}$ | $\mathbf{B}$ |
| :---: | :---: | :---: |
| 3 | 0,750000000 | $-0,250000000$ |
| 7 | 0,875000000 | $-0,125000000$ |
| 15 | 0,937500000 | $-0,062500000$ |
| 31 | 0,968750000 | $-0,031250000$ |
| 63 | 0,984375000 | $-0,015625000$ |
| 127 | 0,992187500 | $-0,007812500$ |
| 255 | 0,996093750 | $-0,001906250$ |
| 511 | 0,998046875 | $-0,000976563$ |
| 1023 | 0,999023438 | $-0,000488281$ |
| 2047 | 0,999511719 | $-0,000244141$ |
| 4095 | 0,999755859 | $-0,000122070$ |
| 8191 | 0,999877930 | $-0,000061035$ |
| 16383 | 0,999938965 | $-0,000030518$ |
| 32767 |  |  |

## C.1.5.2 Layer II encoding

## C.1.5.2.1 Introduction

This clause describes a possible Layer II encoding method. The description is made according to figure C. 5 , "Layer I, II encoder flow chart".

## C.1.5.2.2 Psychoacoustic model

The calculation of the psychoacoustic parameters can be done either with Psychoacoustic Model 1 describixd in clause D.1. or with Psychoacoustic Model 2 described in clause D.2. If Psychoacoustic Model 1 is use d to calculate the psychoacoutic parameters, the FFT shifulength is 1152 samples. If Psychoacoustic Mod:l 2 is used, the calculation is performed twice with a shiftlength of 576 samples and the largest of each pain of signal to mask ratios is used. Either model provides the signal-to-mask ratio for every subband.

## C.1.5.2.3 Analysis subband filter

The analysis subband filter is described in clause C.1.3. "Analysis subband filter".

## C.1.5.2.4 Scalefactor calculation

The calculation of the scalefactor for each subband is performed every 12 subband samples. The maximurr of the absolute value of these 12 samples is determined. The lowest value in table B.1, "Layer I, II Scalefactors", which is larger than this maximum is used as the scalefactor.

## C.1.5.2.5 Coding of scalefactors

A frame corresponds to 36 subband samples and therefore contains three scalefactors per subband. Define 'scr ${ }^{\text {as }}$ the index in table B.1, "Layer I, II Scalefactors". First, the two differences dscf $f_{1}$ and dscf $f_{2}$ of the successive scalefactor indices $\mathrm{scf}_{1}, \mathrm{scf}_{2}$ and $\mathrm{scf}_{3}$ are calculated:

$$
\begin{aligned}
& \operatorname{dscf}_{1}=\operatorname{scf} f_{1}-\operatorname{scf}_{2} \\
& \operatorname{dscf}_{2}=\operatorname{scf}_{2}-\operatorname{scf}_{3}
\end{aligned}
$$

The class of each of the differences is determined as follows:

| class. | dscf |
| :---: | :---: |
| 1 | dscf $<=-3$ |
| 2 | $-3<$ dscf $<0$ |
| 3 | dscf $=0$ |
| 4 | $0<$ dscf $<3$ |
| 5 | dscf $>=3$ |

The pair of classes of differences indicate the entry point in table C.4, "Layer II Scalefactors Transmission Patterns". The column labelled "scalefactor used in encoder" gives the three scalefactors which are actually used. " 1 ", " 2 " and " 3 " mean respectively the first, second and third scalefactor within a frame, " 4 " means the maximum of the three scalefactors. If, after this adjusting of scalefactors two or three are the same, not all scalefactors need to be transmitted for a certain subband within one frame. Only the scalefactors indicated in the "transmission pattern" column are transmitted. The information describing the number and the position of the scalefactors in each subband is called "scalefactor selection information".

## C.1.5.2.6 Coding of scalefactor selection information

The "scalefactor selection information" (scfsi) is coded by a two bit word, which is also to be found in tab e C.4, "Layer II scalefactor transmission patterns". Only the scfsi for the subbands which will get a nonzero bit allocation are transmitted.

## C.1.5.2.7 Bit allocation

Before adjustment to a fixed bitrate, the number of bits ,"adb", that are available for coding the samples an 1 the scalefactors must be determined. This number can be obtained by subtracting from the total number of
available bits "cb", the number of bits needed for the header "bhdr" ( 32 bits), the CRC checkword "bcrc" if used ( 16 bits), the bit allocation "bbal", and the number of bits "banc" required for ancillary data:
$\mathrm{adb}=\mathrm{cb}-$ (bhdr $+\mathrm{bcrc}+$ bbal + banc)
The resulting number can be used to code the subband samples and the scalefactors. The principle used in the allocation procedure is minimization of the total noise-to-mask ratio over the frame with the constraint that the number of bits used does not exceed the number of bits available for that frame. Use is made of table B.2, "Layer II Possible Quantization per subband" that indicates for every subband the number of steps that may be used to quantize the samples. The number of bits required to represent these quantized samples can be derived from table B.4, "Layer II Classes of Quantization".

The allocation procedure is an iterative procedure where, in each iteration step the number of levels of the subband that has the greatest benefit is increased.

First the mask-to-noise ratio "MNR" for each subband is calculated by subtracting from the signal-to-noiseratio "SNR" the signal-to-mask-ratio "SMR":
$\mathrm{MNR}=\mathrm{SNR}-\mathrm{SMR}$
The signal-to-noise-ratio can be found in table C. 5 "Layer II Signal-to-Noise Ratios". The signal-to-maskratio is the output of the psychoacoustic model.

Then zero bits are allocated to the samples and the scalefactors. The number of bits for the samples "bspl" and the number of bits for the scalefactors "bscf" are set to zero. Next an iterative procedure is started. Each iteration loop contains the following steps :

- Determination of the minimal MNR of all subbands.
- The accuracy of the quantization of the subband with the minimal MNR is increased by using the next higher entry in the relevant table B.2, "Layer II Possible Quantization per Subband".
- The new MNR of this subband is calculated.
- bspl is updated according to the additional number of bits required. If a non-zero number of bits is assigned to a subband for the first time, bsel has to be updated, and bscf has to be updated according to the number of scalefactors required for this subband. Then adb is calculated again using the formula :
$\mathrm{adb}=\mathrm{c} \mathrm{b}-(\mathrm{bhdr}+\mathrm{bcrc}+\mathrm{bbal}+\mathrm{bsel}+\mathrm{bscf}+\mathrm{bspl}+\mathrm{banc})$
The iterative procedure is repeated as long as adb is not less than any possible increase of bspl, bsel and bscf within one loop.


## C.1.5.2.8 Quantization and encoding of subband samples

Each of the 12 subband samples is normalized by dividing its value by the scalefactor to obtain X and quantized using the following formula:

- Calculate $\mathrm{A}^{*} \mathrm{X}+\mathrm{B}$
- Take the N most significant bits.
- Inver the MSB

A and B can be found in the table C.6, "Layer II Quantization Coefficients". N represents the necessary number of bits to encode the number of steps. The inversion of the MSB is done in order to avoid the all ' 1 ' code that is used for the synchronization word.

Given the numter of steps that the samples will be quantized to, table B.4, "Layer II Classes of Quantization" shows whether grouping will be used. If grouping is not required, the three samples are coded with individual codewords.

If grouping is required, three consecutive samples are coded as one codeword. Only one value $v_{m}$, MSB first, is transmitted for this triplet. The relationships between the coded value $v_{m}(m=3,5,9)$ and the three consecutive subband samples $\mathrm{x}, \mathrm{y}, \mathrm{z}$ are:

$$
\begin{array}{rlr}
\mathrm{v}_{3}=9 \mathrm{z}+3 \mathrm{y}+\mathrm{x} & \left(\mathrm{v}_{3} \text { in } 0 \ldots 26\right) \\
\mathrm{v}_{5}=25 \mathrm{z}+5 \mathrm{y}+\mathrm{x} & \left(\mathrm{v}_{5} \text { in } 0 \ldots 124\right) \\
\mathrm{v}_{9}=81 \mathrm{z}+9 \mathrm{y}+\mathrm{x} & & \left(\mathrm{v}_{9} \text { in } 0 \ldots 728\right)
\end{array}
$$

## C.1.5.2.9 Coding of bit allocation

For the purpose of a more efficient coding, only a limited number of possible quantizations, which may re different for each subband, are allowed. Only the index with wordlength "nbal" in the relevant table B.2, "Layer II Possible Quantizations per Subband" is transmitted, MSB first.

## C.1.5.2.10 Ancillary data

The Audio standard provides a number of bits for the inclusion and transmission of variable length ancillary data with the audio bitstream. The ancillary data will reduce the number of bits available for audio, which may result in a degradation of audio quality.

The presence of a bit pattern in the ancillary data matching the syncword may hamper synchronization. T iis problem is more likely to occur when the free format is used.

## C.1.5.2.11 Formatting

An overview of the Layer II format can be seen in figure C.3.
Frame


Figure C. 3 -- Layer II Format
The differences compared to the Layer I format are:

- The length of a slot equals 8 bits.
- A new block scfsi containing the scalefactor selection information has been introduced.
- The bit allocation information, scalefactors and samples have been subject to further coding (see the relateds).

The details can be found in 2.4.1.

Table C. 4 -- Layer II scalefactor transmission patterns

| Class $_{1}$ | $\mathrm{Class}_{2}$ | Scalefactors used in encoder | Transmission pattern | Selection Information |
| :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 123 | 123 | 0 |
| 1 | 2 | 122 | 12 | 3 |
| 1 | 3 | 122 | 12 | 3 |
| 1 | 4 | 133 | 13 | 3 |
| 1 | 5 | 123 | 123 | 0 |
| 2 | 1 | 113 | 13 | 1 |
| 2 | 2 | $1 \begin{array}{lll}1 & 1 & 1\end{array}$ | 1 | 2 |
| 2 | 3 | $\begin{array}{lll}1 & 1 & 1\end{array}$ | 1 | 2 |
| 2 | 4 | 444 | 4 | 2 |
| 2 | 5 | 113 | 13 | 1 |
| 3 | 1 | $1 \begin{array}{lll}1 & 1 & 1\end{array}$ | 1 | 2 |
| 3 | 2 | $1 \begin{array}{lll}1 & 1 & 1\end{array}$ | 1 | 2 |
| 3 | 3 | $1 \begin{array}{lll}1 & 1 & 1\end{array}$ | 1 | 2 |
| 3 | 4 | $\begin{array}{llll}3 & 3 & 3\end{array}$ | 3 | 2 |
| 3 | 5 | $\begin{array}{llll}1 & 1 & 3\end{array}$ | 13 | 1 |
| 4 | 1 | 222 | 2 | 2 |
| 4 | 2 | $\begin{array}{lll}2 & 2 & 2\end{array}$ | 2 | 2 |
| 4 | 3 | 222 | 2 | 2 |
| 4 | 4 | 333 | 3 | 2 |
| 4 | 5 | 123 | 123 | 0 |
| 5 | 1 | 1223 | 123 | 0 |
| 5 | 2 | 1222 | 12 | 3 |
| 5 | 3 | 122 | 12 | 3 |
| 5 | 4 | $1 \begin{array}{lll}1 & 3 & 3\end{array}$ | 13 | 3 |
| 5 | 5 | 123 | 123 | 0 |

Table C. 5 -- Layer II Signal-to-Noise Ratios

| No. of steps | SNR (dB) |
| :---: | :---: |
| 0 | 0,00 |
| 3 | 7,00 |
| 5 | 11,00 |
| 7 | 16,00 |
| 9 | 20,84 |
| 15 | 25,28 |
| 31 | 31,59 |
| 63 | 37,75 |
| 127 | 43,84 |
| 255 | 49,89 |
| 511 | 55,93 |
| 1023 | 61,96 |
| 2047 | 67,98 |
| 4095 | 74,01 |
| 8191 | 80,03 |
| 16383 | 86,05 |
| 32767 | 92,01 |
| 65535 | 98,01 |

Table C. 6 -. Layer II quantization coefficients

| No. of steps | $\mathbf{A}$ | $\mathbf{B}$ |
| :---: | :---: | :---: |
| 3 | 0,750000000 | $-0,2500000000$ |
| 5 | 0,625000000 | $-0,375000000$ |
| 7 | 0,875000000 | $-0,125000000$ |
| 9 | 0,562500000 | $-0,437500000$ |
| 15 | 0,937500000 | $-0,062500000$ |
| 31 | 0,968750000 | $-0,031250000$ |
| 63 | 0,984375000 | $-0,015625000$ |
| 127 | 0,992187500 | $-0,003812500$ |
| 255 | 0,996093750 | $-0,00195350$ |
| 511 | 0,998046875 | $-0,000976563$ |
| 1023 | 0,999023438 | $-0,000488281$ |
| 2047 | 0,999511719 | $-0,000244141$ |
| 4095 | 0,999755859 | $-0,000122070$ |
| 8191 | 0,999877930 | $-0,000061035$ |
| 16383 | 0,999938965 | $-0,000030518$ |
| 32767 | 0,999969482 |  |
| 5535 | 0,999984741 |  |



Figure C. 4 -- Analysis subband filter flow chart


Figure C. 5 .- Layer I, II Encoder flow chart

## C.1.5.3 Layer III encoding

## C.1.5.3.1 Introduction

This clause describes a possible Layer III encoding method. The basic data flow is described by the general psychoacoustic coder block diagram. The basic blocks are described in more detail and below.

## C.1.5.3.2 F'sychoacoustic model

The calculation of the psychoacoustic parameters can be done either with Psychoacoustic Model I described in clause D.1. or with Psychoacoustic Model 2 described in clause D.2. A description of modifications to Psychoacoustic Model 2 for use with Layer III can be found below. The model is run twice per block, using a shift length of 576 samples. A signal-to-mask-ratio is provided for every scalefactor band.

## C.1.5.3.2.1 Adaptation of psychoacoustic model II for Layer III

Psychoacoustic Model 2 (clause D.2) is modified as described below for the use with Layer III encoding.

## General considerations:

The model is calculated twice in parallel. One computation is done with a shift length iblen of 192 samples (to be used with short blocks), the other is done with a shift length of 576 samples. For the shift length of 192 samples the block length of the FFT is changed to 256 , and the parameters changed accordingly.

Change to unpredictability calculation:
The calculation of the unpredictability metric in Psychoacoustic Model 2 is changed.

- Calculation of the unpredictability:

The unpredictability cw is calculated for the first 206 spectral lines. For the other spectral lines, the unpredictability is set to 0,4 .
The unpredictability for the first 6 lines is calculated from the long FFT (window length $=1024$, shiften $=576$ ). For the spectral lines 6 up to 205 , the unpredictability is calculated from the short FFT (window length 256 , shifulen $=192$ ):

$$
\mathrm{cw}(w)=\left\{\begin{array}{lll}
\mathrm{cw} \_\mathrm{l}(\mathrm{w}) & \text { for } & 0 \leq w<6 \\
\mathrm{cw} s((w+2) D I V 4) & \text { for } & 6 \leq w<206 \\
0,4 & \text { for } & w \geq 206
\end{array}\right.
$$

cw_l is the unpredictability calculated from the long FFT, $\mathrm{cw} \_$s is the unpredictability calculated from the second short block out of three short blocks within one granule.

- The spreading function has been replaced:

$$
\begin{array}{ll}
\text { If } \mathrm{j} \geq \mathrm{i} & \text { tmpy }=3,0(\mathrm{j}-\mathrm{i}) \\
\text { else } & \text { tmpy }=1,5(\mathrm{j}-\mathrm{i}) \text { is used. }
\end{array}
$$

Only values of the spreading function greater than $10^{-6}$ are used. All other values are set to zero.

- For converting the unpredictability the parameters

$$
\begin{aligned}
& \text { conv } 1=-0,299 \\
& \text { conv2 }=-0,43
\end{aligned}
$$

are used.

- The parameter NMT (noise masking tone) is set to $6,0 \mathrm{~dB}$ for all threshold calculation partions. The parameter 7 MN (tone masking noise) is set to $29,0 \mathrm{~dB}$ for all partitions.
For minval see table "threshold calculation partitions" (table C.7).
- The psychoacoustic entropy is estimated from the ratio thr/eb, where thr is the threshold and eb is he energy:
$\mathrm{pe}=-\sum\left(\right.$ cbwidth $\left._{\mathrm{k}} \cdot \log \left(\mathrm{th}_{\mathrm{k}} /\left(\mathrm{eb}_{\mathrm{k}}+1.\right)\right)\right)$
where $k$ indexes the threshold calculation partitions and cbwidth is the width of the threshold calculation partition (see tables).
- pre-echo control

The following constants are used for the control of pre-echo's (see block diagram):

$$
\text { rpelev }=2
$$

$$
\text { rpelev2 }=16
$$

- The threshold is not spread over the FFT lines. The threshold calculation partitions are converted directly to scalefactor bands. The first partition which is added to the scalefactor band is weighted wi $h$ w1, the last with w2 (see table C. 8 "Converting Threshold Calculation Partitions to Scalefactor Bands"). The table contains also the number of partitions (cbw) converted to one scalefactor band (excluding the first and the last partition).

The parameters bo and bu are shown in table C.8. They are used for converting threshold calculation partitions to scalefactor bands.

- For short blocks a simplified version of the threshold calculation (constant signal to noise ratio) is used. The constants can be found in the columns labelled "SNR (dB)" in table C.7(def) below.

Table C. 7 .- Threshold calculation partitions with following parameters width, minval, threshold in quiet, norm and bval:

Table C.7.a .- Sampling_frequency $=48 \mathbf{~ k H z}$ long blocks

| no. | FFT-Iines | minval | gthr | norm | bval |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 24,5 | 4,532 | 0.970 | 0,000 |
| 1 | 1 | 24,5 | 4,532 | 0.755 | 0,469 |
| 2 | 1 | 24,5 | 4,532 | 0,738 | 0.937 |
| 3 | 1 | 24.5 | 0,904 | 0,730 | 1,406 |
| 4 | 1 | 24,5 | 0,904 | 0,724 | 1,875 |
| 5 | 1 | 20 | 0,090 | 0,723 | 2,344 |
| 6 | 1 | 20 | 0,090 | 0,723 | 2,812 |
| 7 | 1 | 20 | 0,029 | 0,723 | 3,281 |
| 8 | 1 | 20 | 0.029 | 0,718 | 3.750 |
| 9 | 1 | 20 | 0.009 | 0,690 | 4,199 |
| 10 | 1 | 20 | 0,009 | 0,660 | 4,625 |
| 11 | 1 | 18 | 0.009 | 0,641 | 5,047 |
| 12 | 1 | 18 | 0,009 | 0,600 | 5,437 |
| 13 | 1 | 18 | 0,009 | 0,584 | 5,828 |
| 14 | 1 | 12 | 0,009 | 0,531 | 6,187 |
| 15 | 1 | 12 | 0,009 | 0,537 | 6.522 |
| 16 | 2 | 6 | 0,018 | 0,857 | 7,174 |
| 17 | 2 | 6 | 0,018 | 0.858 | 7.800 |
| 18 | 2 | 3 | 0,018 | 0.853 | 8,402 |
| 19 | 2 | 3 | 0,018 | 0.824 | 8.966 |
| 20 | 2 | 3 | 0.018 | 0,778 | 9,483 |
| 21 | 2 | 3 | 0.018 | 0,740 | 9.966 |
| 22 | 2 | 0 | 0,018 | 0,709 | 10,426 |
| 23 | 2 | 0 | 0.018 | 0,676 | 10,866 |
| 24 | 2 | 0 | 0,018 | 0,632 | 11,279 |
| 25 | 2 | 0 | 0,018 | 0,592 | 11,669 |
| 26 | 2 | 0 | 0.018 | 0,553 | 12,042 |
| 27 | 2 | 0 | 0.018 | 0,510 | 12,386 |
| 28 | 2 | 0 | 0,018 | 0,513 | 12,721 |
| 29 | 3 | 0 | 0.027 | 0,608 | 13,115 |
| 30 | 3 | 0 | 0.027 0.027 | 0,673 0,636 | 13,561 13,983 |
| 32 | 3 | 0 | 0,027 | 0.586 | 14,371 |
| 33 | 3 | 0 | 0,027 | 0,571 | 14,741 |
| 34 | 4 | 0 | 0,036 | 0,616 | 15,140 |
| 35 | 4 | 0 | 0.036 | 0,640 | 15,562 |
| 36 | 4 | 0 | 0,036 | 0,597 | 15,962 |
| 37 | 4 | 0 | 0.036 | 0,538 | 16,324 |
| 38 | 4 | 0 | 0,036 | 0,512 | 16,665 |
| 39 | 5 | 0 | 0.045 | 0,528 | 17,020 |
| 40 | 5 | 0 | 0.045 | 0.516 | 17,373 |
| 41 | 5 | 0 | 0,045 | 0,493 | 17,708 |
| 42 | 6 | 0 | 0,054 | 0,499 | 18,045 |
| 43 | 7 | 0 | 0,063 | 0,525 | 18,398 |
| 44 | 7 | 0 | 0,063 | 0,541 | 18,762 |
| 45 | 8 | 0 | 0,072 | 0,528 | 19,120 |
| 46 | 8 | 0 | 0,072 | 0,510 | 19,466 |
| 47 | 8 | 0 | 0.072 | 0,506 | 19.807 |
| 48 | 10 | 0 | 0,180 | 0.525 | 20.159 20.522 |
| 49 50 | 10 10 | 0 | 0.180 0.180 | 0,536 0.518 | 20,522 20.873 |
| 51 | 13 | 0 | 0,372 | 0,501 | 21,214 |
| 52 | 13 | 0 | 0,372 | 0,496 | 21,553 |
| 53 | 14 | 0 | 0,400 | 0,497 | 21,892 |
| 54 | 18 | 0 | 1,628 | 0,495 | 22,231 |
| 55 | 18 | 0 | 1,628 | 0,494 | 22,569 |
| 56 | 20 | 0 | 1.808 | 0.497 | 22,909 |
| 57 | 25 | 0 | 22,607 | 0,494 | 23.248 |
| 58 | 25 | 0 | 22,607 | 0.487 | 23,583 |
| 59 | 35 | 0 | 31,650 | 0.483 | 23,915 |
| 60 | 67 | 0 | 605,867 | 0,482 | 24,246 |
| 61 | 67 | 0 | 605,867 | 0,524 | 24,576 |

Table C.7.b Sampling_frequency $=44,1 \mathbf{k H z}$ long blocks

| no. | FFT-lines | minval | qthr | norm | bval |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 24.5 | 4,532 | 0.951 | 0,000 |
| 1 | 1 | 24.5 | 4,532 | 0,700 | 0,431 |
| 2 | 1 | 24.5 | 4,532 | 0,681 | 0,861 |
| 3 | 1 | 24.5 | 0,904 | 0,675 | 1,292 |
| 4 | , | 24,5 | 0,904 | 0,667 | 1,723 |
| 5 | 1 | 20 | 0,090 | 0,665 | 2,153 |
| 6 | 1 | 20 | 0,090 | 0,664 | 2,584 |
| 7 | 1 | 20 | 0,029 | 0,664 | 3.015 |
| 8 | 1 | 20 | 0,029 | 0,664 | 3,445 |
| 9 | 1 | 20 | 0,029 | 0,655 | 3.876 |
| 10 | 1 | 20 | 0,009 | 0,616 | 4,279 |
| 11 | 1 | 20 | 0,009 | 0,597 | 4,670 |
| 12 | 1 | 18 | 0,009 | 0,578 | 5,057 |
| 13 | t | 18 | 0,009 | 0.541 | 5,415 |
| 14 | 1 | 18 | 0,009 | 0,575 | 5,774 |
| 15 | 2 | 12 | 0,018 | 0,856 | 6,422 |
| 16 | 2 | 6 | 0,018 | 0,846 | 7,026 |
| 17 | 2 | 6 | 0,018 | 0,840 | 7.609 |
| 18 | 2 | 3 | 0,018 | 0,822 | 8,168 |
| 19 | 2 | 3 | 0,018 | 0,800 | 8.710 |
| 20 | 2 | 3 | 0,018 | 0,753 | 9,207 |
| 21 | 2 | 3 | 0,018 | 0,704 | 9.662 |
| 22 | 2 | 0 | 0,018 | 0.674 | 10,099 |
| 23 | 2 | 0 | 0,018 | 0,640 | 10,515 |
| 24 | 2 | 0 | 0.018 | 0,609 | 10.917 |
| 25 | 2 | 0 | 0.018 | 0.566 | 11.293 |
| 26 | 2 | 0 | 0.018 | 0.535 | 11.652 |
| 27 | 2 | 0 | 0.018 | 0.531 | 11.997 |
| 28 | 3 | 0 | 0,027 | 0.615 | 12,394 |
| 29 | 3 | 0 | 0,027 | 0,686 | 12,850 |
| 30 | 3 | 0 | 0,027 | 0,650 | 13,277 |
| 31 | 3 | 0 | 0,027 | 0.611 | 13.681 |
| 32 | 3 | 0 | 0,027 | 0,567 | 14,062 |
| 33 | 3 | 0 | 0,027 | 0,520 | 14,411 |
| 34 | 3 | 0 | 0,027 | 0,513 | 14,751 |
| 35 | 4 | 0 | 0.036 | 0,557 | 15,119 |
| 36 | 4 | 0 | 0,036 | 0.584 | 15.508 |
| 37 | 4 | 0 | 0.036 | 0,570 | 15.883 |
| 38 | 5 | 0 | 0,045 | 0,579 | 16,263 |
| 39 | 5 | 0 | 0.045 | 0.585 | 16,654 |
| 40 | 5 | 0 | 0,045 | 0,548 | 17,020 |
| 41 | 6 | 0 | 0,054 | 0,536 | 17.374 |
| 42 | 6 | 0 | 0,054 | 0,550 | 17,744 |
| 43 | 7 | 0 | 0,063 | 0,532 | 18,104 |
| 44 | 7 | 0 | 0,063 | 0,504 | 18,447 |
| 45 | 7 | 0 | 0,063 | 0,496 | 18,781 |
| 46 | 9 | 0 | 0.081 | 0.516 | 19.130 |
| 47 | 9 | 0 | 0,081 | 0,527 | 19,487 |
| 48 | 9 | 0 | 0.081 | 0.516 | 19.838 |
| 49 | 10 | 0 | 0,180 | 0,497 | 20.179 |
| 50 | 10 | 0 | 0,180 | 0,489 | 20.510 |
| 51 | 11 | 0 | 0,198 | 0,502 | 20.852 |
| 52 | 14 | 0 | 0,400 | 0,502 | 21.196 |
| 53 | 14 | 0 | 0,400 | 0,491 | 21.531 |
| 54 | 15 | 0 | 0,429 | 0.497 | 21.870 |
| 55 | 20 | 0 | 1,808 | 0.504 | 22.214 |
| 56 | 20 | 0 | 1.808 | 0,504 | 22.558 |
| 57 | 21 | 0 | 1,899 | 0.495 | 22.898 |
| 58 | 27 | 0 | 24,415 | 0,486 | 23.232 |
| 59 | 27 | 0 | 24,415 | 0,484 | 23.564 |
| 60 | 36 | 0 | 32,554 | 0,483 | 23.897 |
| 61 | 73 | 0 | 660,124 | 0,475 | 24.229 |
| 62 | 18 | 0 | 162,770 | 0,515 | 24,542 |

Table C.7.c .- Sampling_frequency $=32 \mathbf{k H z}$ long blocks

| no. | FFT-lines | minval | qthr | norm | bval |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 2 | 24.5 | 9,064 | 0,997 | 0,312 |
| 1 | 2 | 24,5 | 9,064 | 0,893 | 0,937 |
| 2 | 2 | 24.5 | 1.808 | 0.881 | 1,562 |
| 3 | 2 | 20 | 0,181 | 0,873 | 2,187 |
| 4 | 2 | 20 | 0,181 | 0,872 | 2,812 |
| 5 | 2 | 20 | 0,057 | 0,871 | 3,437 |
| 6 | 2 | 20 | 0,018 | 0,860 | 4,045 |
| 7 | 2 | 20 | 0,018 | 0,839 | 4,625 |
|  | 2 | 18 | 0,018 | 0.812 | 5.173 |
| 9 | 2 | 18 | 0,018 | 0,784 | 5.698 |
| 10 | 2 | 12 | 0.018 | 0,741 | 6,184 |
| 11 | 2 | 12 | 0,018 | 0,697 | 6,634 |
| 12 | 2 | 6 | 0,018 | 0,674 | 7,070 |
| 13 | 2 | 6 | 0,018 | 0,651 | 7.492 |
| 14 | 2 | 6 | 0.018 | 0,633 | 7.905 |
| 15 | 2 | 3 | 0,018 | 0,611 | 8,305 |
| 16 | 2 | 3 | 0.018 | 0.589 | 8.695 |
| 17 | 2 | 3 | 0.018 | 0.575 | 9.064 |
| 18 | 3 | 3 | 0.027 | 0.654 | 9.483 |
| 19 | 3 | 3 | 0.027 | 0,724 | 9.966 |
| 20 | 3 | ${ }^{\circ}$ | 0,027 | 0.701 | 10.425 |
| 21 | 3 | ${ }^{\circ}$ | 0,027 | 0.673 | 10.866 |
| 22 23 | 3 3 3 | ${ }_{0}^{0}$ | 0,027 0,027 | 0.631 0.592 | 11,279 11,669 |
| 24 | 3 | 0 | 0.027 | 0,553 | 12,042 |
| 25 | 3 | 0 | 0.027 | 0.510 | 12,386 |
| 26 | 3 | 0 | 0.027 | 0,505 | 12,721 |
| 27 | 4 | 0 | 0.036 | 0.562 | 13.091 |
| 28 | 4 | 0 | 0,036 | 0,598 | 13,488 |
| 29 | 4 | - | 0,036 | 0.589 | 13.873 |
| 30 | 5 | 0 | 0,045 | 0,607 | 14,268 |
| 31 32 | 5 | 0 | 0,045 | 0,620 | 14,679 |
| 32 | 5 | 0 | 0,045 | 0,580 | 15,067 |
| $\begin{array}{r}33 \\ 34 \\ \hline\end{array}$ | 5 | 0 | 0,045 | 0.532 | 15,424 |
| 34 35 | 5 6 | $\bigcirc$ | 0.045 0,054 | 0,517 0,517 | 15,771 16,120 |
| 36 | 6 | 0 | 0.054 | 0,509 | 16,466 |
| 37 | 6 | 0 | 0.054 | 0,506 | 16,807 |
| 38 | 8 | 0 | 0,072 | 0,522 | 17,158 |
| 39 | 8 | 0 | 0.072 | ${ }^{0.531}$ | 17.518 |
| 40 | 8 ${ }_{\text {10 }}$ | ${ }_{0}^{0}$ | 0,072 0,090 | 0,519 0,512 | 17,869 18,215 |
| 42 | 10 | 0 | 0,090 | 0,509 | 18,562 |
| 43 | 10 | 0 | 0,090 | 0,497 | 18,902 |
| 44 | 12 | 0 | 0,108 | 0.494 | 19,239 |
| 45 | 12 | 0 | 0.108 | 0,501 | 19,579 |
| 46 | 13 | $\bigcirc$ | 0.117 | -0.507 | 19.925 20.269 |
| 47 48 | 14 | 0 | 0.252 | 0,502 | 20,269 |
| 48 49 | 14 | 0 | 0.252 | 0.493 0.497 | 20,606 |
| 50 | 16 | ${ }_{0}$ | 0,572 | 0,506 | 21,288 |
| 51 | 20 | 0 | 0.572 | 0,510 | 21,635 |
| 52 | 27 | $\bigcirc$ | 0,658 | 0,504 | 21,979 |
| 53 | 27 | 0 | 2,441 | 0.496 | 22,319 |
| 54 55 | 27 | $\bigcirc$ | 2,441 2,894 | 0,493 0.490 | 22,656 22.993 |
| 55 56 | 32 37 | 0 | 2,894 33,458 | 0.490 0.483 | 22,993 23,326 |
| 57 | 37 | 0 | 33,458 | 0.458 | 23,656 |
| 58 | 12 | 0 | 10.851 | 0,500 | 23.937 |

Table C.7.d .- Sampling_frequency $=48 \mathrm{kHz}$ short blocks

| no. | FFT-lines | qthr | norm | SNR (db) | bval |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 4,532 | 0,970 | -8,240 | 0,000 |
| 1 | 1 | 0.904 | 0,755 | -8,240 | 1,875 |
| 2 | 1 | 0,029 | 0,738 | -8,240 | 3.750 |
| 3 | 1 | 0,009 | 0,730 | -8.240 | 5,437 |
| 4 | 1 | 0.009 | 0,724 | -8,240 | 6,857 |
| 5 | 1 | 0.009 | 0.723 | -8,240 | 8,109 |
| 6 | 1 | 0.009 | 0,723 | -8,240 | 9,237 |
| 7 | I | 0,009 | 0,723 | -8,240 | 10,202 |
| 8 | 1 | 0,009 | 0,718 | -8,240 | 11,083 |
| 9 | 1 | 0,009 | 0,690 | -8,240 | 11,864 |
| 10 | 1 | 0,009 | 0,660 | -7,447 | 12,553 |
| 11 | 1 | 0,009 | 0.641 | -7,447 | 13,195 |
| 12 | 1 | 0,009 | 0,600 | -7,447 | 13,781 |
| 13 | 1 | 0,009 | 0,584 | -7,447 | 14,309 |
| 14 | 1 | 0,009 | 0,532 | -7,447 | 14,803 |
| 15 | 1 | 0,009 | 0,537 | -7,447 | 15,250 |
| 16 | 1 | 0,009 | 0,857 | -7,447 | 15,667 |
| 17 | 1 | 0,009 | 0.858 | -7,447 | 16,068 |
| 18 | 1 | 0,009 | 0,853 | $-7,447$ | 16,409 |
| 19 | 2 | 0,018 | 0.824 | -7,447 | 17,044 |
| 20 | 2 | 0.018 | 0.778 | -6,990 | 17,607 |
| 21 | 2 | 0.018 | 0,740 | -6,990 | 18,097 |
| 22 | 2 | 0.018 | 0,709 | -6,990 | 18.528 |
| 23 | 2 | 0,018 | 0,676 | -6,990 | 18.930 |
| 24 | 2 | 0,018 | 0,632 | -6,990 | 19.295 |
| 25 | 2 | 0,018 | 0,592 | -6,990 | 19,636 |
| 26 | 3 | 0,054 | 0,553 | -6,990 | 20,038 |
| 27 | 3 | 0,054 | 0,510 | -6,990 | 20,486 |
| 28 | 3 | 0,054 | 0,513 | -6,990 | 20,900 |
| 29 | 4 | 0,114 | 0,608 | -6.990 | 21,305 |
| 30 | 4 | 0,114 | 0,673 | -6,020 | 21.722 |
| 31 | 5 | 0.452 | 0,637 | -6.020 | 22,128 |
| 32 | 5 | 0,452 | 0,586 | $-6,020$ $-6,020$ | 22,512 22,877 |
| 33 | 5 | 0,452 | 0,571 | $-6,020$ $-5,229$ | 22,877 23,241 |
| 34 35 | 7 | 6,330 6,330 | 0,616 | $-5,229$ $-5,229$ | 23,241 23,616 |
| 35 36 | 7 11 | 6,330 9,947 | 0,640 0,597 | $-5,229$ $-5,229$ | 23,616 23,974 |
| 37 | 17 | 153,727 | 0,538 | -5,229 | 24,312 |

Table C.7.e .- Sampling_frequency $=44,1 \mathbf{k H z}$ short blocks

| no. | FFT-lines | qthr | norm | SNR (db) | bval |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 4,532 | 0,952 | -8,240 | 0,000 |
| 1 | 1 | 0,904 | 0,700 | -8,240 | 1,723 |
| 2 | 1 | 0,029 | 0,681 | -8,240 | 3,445 |
| 3 | 1 | 0,009 | 0,675 | -8,240 | 5,057 |
| 4 | 1 | 0,009 | 0,667 | -8,240 | 6,422 |
| 5 | 1 | 0,009 | 0,665 | -8,240 | 7,609 |
| 6 | 1 | 0,009 | 0,664 | -8,240 | 8.710 |
| 7 | 1 | 0,009 | 0,664 | -8,240 | 9,662 |
| 8 | 1 | 0,009 | 0,664 | -8,240 | 10,515 |
| 9 | 1 | 0,009 | 0.655 | -8,240 | 11,293 |
| 10 | 1 | 0,009 | 0.616 | -7,447 | 12,009 |
| 11 | 1 | 0,009 | 0,597 | -7.447 | 12,625 |
| 12 | 1 | 0,009 | 0,578 | -7,447 | 13,210 |
| 13 | 1 | 0,009 | 0,541 | -7,447 | 13,748 |
| 14 | 1 | 0,009 | 0,575 | -7,447 | 14,241 |
| 15 | 1 | 0,009 | 0.856 | -7.447 | 14,695 |
| 16 | 1 | 0,009 | 0.846 | -7,447 | 15,125 |
| 17 | 1 | 0,009 | 0.840 | -7.447 | 15.508 |
| 18 | 1 | 0.009 | 0.822 | -7.447 | 15,891 |
| 19 | 2 | 0,018 | 0.800 | -7.447 | 16.537 |
| 20 | 2 | 0,018 | 0,753 | -6.990 | 17,112 |
| 21 | 2 | 0,018 | 0,704 | -6,990 | 17,620 |
| 22 | 2 | 0,018 | 0,674 | -6,990 | 18,073 |
| 23 | 2 | 0,018 | 0,640 | -6.990 | 18,470 |
| 24 | 2 | 0,018 | 0,609 | -6.990 | 18,849 |
| 25 | 3 | 0,027 | 0.566 | -6,990 | 19,271 |
| 26 | 3 3 | 0,027 | 0,535 | -6,990 | 19,741 |
| 27 | 3 | 0,054 | 0,531 | -6,990 | 20,177 |
| 28 29 | 3 | 0,054 | 0,615 | $-6,990$ $-6,990$ | 20,576 |
| 29 30 | 4 | 0,054 | 0,686 | $-6,990$ $-6,020$ | 20,950 |
| 30 31 | 4 4 | 0,114 0,114 | 0,650 0,612 | $-6,020$ $-6,020$ | 21,316 21,699 |
| 32 | 5 | 0.452 | 0,567 | -6,020 | 22,078 |
| 33 | 5 | 0.452 | 0,520 | -6,020 | 22.438 |
| 34 | 5 | 0,452 | 0.513 | -5,229 | 22,782 |
| 35 | 7 | 6.330 | 0.557 | -5,229 | 23.133 |
| 36 | 7 | 6,330 | 0,584 | -5,229 | 23,484 |
| 37 | 7 | 6,330 | 0,570 | -5,229 | 23,828 |
| 38 | 19 | 171.813 | 0,578 | -4,559 | 24,173 |

Table C.7.f -- Sampling_frequency $=32 \mathbf{k H z}$ short blocks

| no. | FFT-lines | gthr | norm | SNR (db) | bral |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 4,532 | 0,997 | -8,240 | 0,000 |
| 1 | 1 | 0,904 | 0,893 | -8,240 | 1,250 |
| 2 | 1 | 0,090 | 0,881 | -8,240 | 2,500 |
| 3 | 1 | 0.029 | 0,873 | -8,240 | 3.750 |
| 4 | 1 | 0.009 | 0.872 | -8,240 | 4,909 |
| 5 | 1 | 0,009 | 0.871 | -8,240 | 5.958 |
| 6 | $t$ | 0,009 | 0,860 | -8,240 | 6.857 |
| 7 | 1 | 0,009 | 0,839 | -8,240 | 7,700 |
| 8 | 1 | 0,009 | 0,812 | -8,240 | 8,500 |
| 9 | 1 | 0,009 | 0,784 | -8,240 | 9.237 |
| 10 | 1 | 0,009 | 0,741 | -7,447 | 9.895 |
| 11 | 1 | 0,009 | 0,697 | -7,447 | 10.500 |
| 12 | 1 | 0.009 | 0,674 | -7.447 | 11,083 |
| 13 | 1 | 0,009 | 0,651 | -7.447 | 11.604 |
| 14 | 1 | 0.009 | 0.633 | -7.447 | 12,107 |
| 15 | 1 | 0,009 | 0,611 | -7,447 | [2,554 |
| 16 | 1 | 0,009 | 0,589 | -7,447 | 13,000 |
| 17 | 1 | 0,009 | 0.575 | -7,447 | 13.391 |
| 18 | 1 | 0,009 | 0,654 | -7,447 | 13,781 |
| 19 | 2 | 0,018 | 0,724 | -7,447 | 14.474 |
| 20 | 2 | 0.018 | 0.701 | -6,990 | 15,096 |
| 21 | 2 | 0,018 | 0.673 | -6.990 | 15.667 |
| 22 | 2 | 0.018 | 0.631 | -6.990 | 16.177 |
| 23 | 2 | 0,018 | 0.592 | -6,990 | 16.636 |
| 24 | 2 | 0,018 | 0,553 | -6,990 | [7,057 |
| 25 | 2 | 0,018 | 0,510 | -6,990 | 17,429 |
| 26 | 2 | 0,018 | 0,506 | -6,990 | 17,786 |
| 27 78 | 3 | 0,027 | 0,562 | $-6,990$ $-6,990$ | 18,177 |
| 28 | 3 | 0,027 | 0,598 | $-6,990$ $-6,990$ | 18,597 18,994 |
| 29 30 | 3 | 0,027 0,027 | 0,589 0,607 | $-6,990$ $-6,020$ | 18,994 19,352 |
| 31 | 3 | 0,027 | 0,620 | -6,020 | 19,693 |
| 32 | 4 | 0,072 | 0.580 | -6,020 | 20,066 |
| 33 | 4 | 0,072 | 0.532 | -6,020 | 20,461 |
| 34 | 4 | 0,072 | 0.517 | -5,229 | 20,841 |
| 35 | 5 | 0, 143 | 0,517 | -5,229 | 21,201 |
| 36 | 5 | 0,143 | 0,509 | -5,229 | 21,549 |
| 37 | 6 | 0,172 | 0,506 | -5,229 | 21,911 |
| 38 | 7 | 0,633 | 0,522 | -4,559 | 22,275 |
| 39 | 7 | 0,633 | 0,531 | -4,559 | 22,625 |
| 40 | 8 | 0,723 9,043 | 0,519 | $-3,980$ $-3,980$ | $22,971$ |
| 41 | 10 | 9,043 | 0,512 | -3,980 | 23,321 |

Table C. 8 .- Tables for converting threshold calculation partitions to scalefactor bands

Table C.8.a -- Sampling_frequency $=48 \mathrm{kHz}$ long blocks

| no. sb | cbw | bu | bo | w 1 | w2 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 3 | 0 | 4 | 1,000 | 0,056 |
| 1 | 3 | 4 | 7 | 0,944 | 0,611 |
| 2 | 4 | 7 | 11 | 0,389 | 0,167 |
| 3 | 3 | 11 | 14 | 0,833 | 0,722 |
| 4 | 3 | 14 | 17 | 0,278 | 0,639 |
| 5 | 2 | 17 | 19 | 0,361 | 0,417 |
| 6 | 3 | 19 | 22 | 0,583 | 0,083 |
| 7 | 2 | 22 | 24 | 0,917 | 0,750 |
| 8 | 3 | 24 | 27 | 0,250 | 0,417 |
| 9 | 3 | 27 | 30 | 0,583 | 0,648 |
| 10 | 3 | 30 | 33 | 0,352 | 0,611 |
| 11 | 3 | 33 | 36 | 0,389 | 0,625 |
| 12 | 4 | 36 | 40 | 0.375 | 0,144 |
| 13 | 3 | 40 | 43 | 0,856 | 0,389 |
| 14 | 3 | 43 | 46 | 0,611 | 0,160 |
| 15 | 3 | 46 | 49 | 0,840 | 0,217 |
| 16 | 3 | 49 | 52 | 0,783 | 0,184 |
| 17 | 2 | 52 | 54 | 0,816 | 0,886 |
| 18 | 3 | 54 | 57 | 0,114 | 0,313 |
| 19 | 2 | 57 | 59 | 0,687 | 0,452 |
| 20 | 1 | 59 | 60 | 0,548 | 0.908 |

Table C.8.b -- Sampling_frequency $=44,1 \mathbf{k H z}$ long blocks

| no. sb | cbw | bu | bo | w 1 | w2 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 3 | 0 | 4 | 1,000 | 0,056 |
| 1 | 3 | 4 | 7 | 0,944 | 0,611 |
| 2 | 4 | 7 | 11 | 0,389 | 0,167 |
| 3 | 3 | 11 | 14 | 0,833 | 0,722 |
| 4 | 3 | 14 | 17 | 0,278 | 0,139 |
| 5 | 3 | 17 | 18 | 0,861 | 0,917 |
| 6 | 3 | 18 | 21 | 0,083 | 0,583 |
| 7 | 3 | 21 | 24 | 0,417 | 0,250 |
| 8 | 3 | 27 | 27 | 0,750 | 0,805 |
| 9 | 3 | 30 | 30 | 0,194 | 0,574 |
| 10 | 43 | 33 | 36 | 0,426 | 0,537 |
| 11 | 3 | 40 | 40 | 0,463 | 0,819 |
| 12 | 3 | 43 | 43 | 0,180 | 0,100 |
| 13 | 3 | 46 | 46 | 0,900 | 0,468 |
| 14 | 3 | 49 | 52 | 0,532 | 0,623 |
| 15 | 3 | 52 | 55 | 0,376 | 0,450 |
| 16 | 2 | 55 | 57 | 0,448 | 0,552 |
| 17 | 2 | 57 | 59 | 0,403 |  |
| 18 | 2 | 59 | 61 | 0,597 | 0,643 |
| 19 |  |  |  | 0,278 | 0,722 |
| 20 |  |  |  |  | 0,960 |

Table C.8.c .- Sampling_frequency $=32 \mathbf{~ k H z}$ long blocks

| no. sb | cbw | bu | bo | w 1 | w 2 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 1 | 0 | 2 | 1,000 | 0,528 |
| 1 | 2 | 2 | 4 | 0,472 | 0,305 |
| 2 | 2 | 4 | 6 | 0,694 | 0,083 |
| 3 | 1 | 6 | 7 | 0,917 | 0,861 |
| 4 | 2 | 7 | 9 | 0,139 | 0,639 |
| 5 | 2 | 9 | 11 | 0,361 | 0,417 |
| 6 | 3 | 11 | 14 | 0,583 | 0,083 |
| 7 | 2 | 14 | 16 | 0,917 | 0,750 |
| 8 | 3 | 16 | 19 | 0,250 | 0,870 |
| 9 | 3 | 19 | 22 | 0,130 | 0,833 |
| 10 | 4 | 22 | 26 | 0,167 | 0,389 |
| 11 | 4 | 26 | 30 | 0,611 | 0,478 |
| 12 | 4 | 30 | 34 | 0,522 | 0,033 |
| 13 | 3 | 34 | 37 | 0,967 | 0,917 |
| 14 | 4 | 37 | 41 | 0,083 | 0,617 |
| 15 | 3 | 41 | 44 | 0,383 | 0,995 |
| 16 | 4 | 44 | 48 | 0,005 | 0,274 |
| 17 | 3 | 48 | 51 | 0,726 | 0,480 |
| 18 | 3 | 51 | 54 | 0,519 | 0,261 |
| 19 | 2 | 54 | 56 | 0,739 | 0,884 |
| 20 | 2 | 56 | 58 | 0,116 | 1,000 |

Table C.8.d -- Sampling_frequency $=48 \mathrm{kHz}$ short blocks

| no. sb | cbw | bu | bo | w | 1 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 2 | 0 | 3 | 1,000 | 0,167 |
| 1 | 2 | 3 | 5 | 0,833 | 0,833 |
| 2 | 3 | 5 | 8 | 0,167 | 0,500 |
| 3 | 3 | 8 | 11 | 0,500 | 0,167 |
| 4 | 4 | 11 | 15 | 0,833 | 0,167 |
| 5 | 4 | 15 | 19 | 0,833 | 0,583 |
| 6 | 3 | 19 | 22 | 0,417 | 0,917 |
| 7 | 4 | 22 | 26 | 0,083 | 0,944 |
| 8 | 4 | 36 | 30 | 0,055 | 0,042 |
| 9 | 3 | 32 | 32 | 0,958 | 0,567 |
| 10 | 3 | 35 | 0,433 | 0,167 |  |
| 11 | 2 | 35 | 0,833 | 0,618 |  |

Table C.8.e -- Sampling_frequency $=44,1 \mathrm{kHz}$ short blocks

| no. sb | cbw | bu | bo | w 1 | w 2 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 2 | 0 | 3 | 1,000 | 0,167 |
| 1 | 2 | 3 | 5 | 0,833 | 0,833 |
| 2 | 3 | 5 | 8 | 0,167 | 0,500 |
| 3 | 3 | 8 | 11 | 0,500 | 0,167 |
| 4 | 4 | 11 | 15 | 0,833 | 0,167 |
| 5 | 5 | 15 | 20 | 0,833 | 0,250 |
| 6 | 3 | 20 | 23 | 0,750 | 0,583 |
| 7 | 4 | 23 | 27 | 0,417 | 0,055 |
| 8 | 3 | 27 | 30 | 0,944 | 0,375 |
| 9 | 3 | 30 | 33 | 0,625 | 0,300 |
| 10 | 3 | 33 | 36 | 0,700 | 0,167 |
| 11 | 2 | 36 | 38 | 0,833 | 1,000 |

Table C.8.f -- Sampling_frequency $=32 \mathrm{kHz}$ short blocks

| no. sb | cbw | bu | bo | w 1 | w 2 |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 2 | 0 | 3 | 1,000 | 0,167 |
| 1 | 2 | 3 | 5 | 0,833 | 0,833 |
| 2 | 3 | 5 | 8 | 0,167 | 0,500 |
| 3 | 4 | 8 | 11 | 0,500 | 0,167 |
| 4 | 5 | 11 | 15 | 0,833 | 0,167 |
| 5 | 4 | 20 | 20 | 0,833 | 0,250 |
| 6 | 5 | 24 | 24 | 0,750 | 0,250 |
| 7 | 4 | 29 | 33 | 0,750 | 0,055 |
| 8 | 43 | 37 | 0,944 | 0,375 |  |
| 9 | 3 | 37 | 40 | 0,625 | 0,472 |
| 10 | 1 | 40 | 0,528 | 0,937 |  |
| 11 |  |  |  |  |  |



Figure C.6.a -- Block diagram psychoacoustic model 2, Layer III: calculate thresholc|


Figure C.6.b -- Block diagram psychoacoustic model 2, Layer III: calculate threshold (part 1)


Figure C.6.c -- Block diagram psychoacoustic model 2, Layer III: calculate threshold (part 2)


Figure C.6.d .. Block diagram psychoacoustic model 2, Layer III: calculate threshold for short blocks

## Window switching decision:

The decision whether the filterbank should be switched to short windows is derived from the calculation (if the masking threshold by calculating the estimate of the psychoacoustic entropy (PE) and switching when the PE exceeds the value 1800 . If this condition is met, the sequence start (block_type=1), short (block_type=2), short, stop (block_type=3) is started. Figure C. 7 shows the possible state changes :or the window switching logic.


Figure C. 7 -- Window Switching State Diagram

## C.1.5.3.3 Analysis part of the hybrid filterbank

The subband analysis of the polyphase filterbank is described in clause C.1.3, "Subband analysis filter". The output of the polyphase filterbank is the input to the subdivision using the MDCT. According to the output of the psychoacoustic model (variables blocksplit_flag and block_type) the window and transform types normal, start, short or stop are used.

18 consecutive output values of one granule and 18 output values of the granule before are assembled to one blcick of 36 samples.

Block type "normal"

$$
z_{i}=x_{i}^{\prime} \sin \left(\frac{\pi}{36}\left(i+\frac{1}{2}\right)\right) \quad \text { for } i=0 \text { to } 35
$$

## Block type "start"

$$
z_{i}= \begin{cases}x_{i}^{\prime} \sin \left(\frac{\pi}{36}\left(i+\frac{1}{2}\right)\right) & \text { for } i=0 \text { to } 17 \\ x_{i}^{\prime} & \text { for } i=18 \text { to } 23 \\ x_{i}^{\prime} \sin \left(\frac{\pi}{12}\left(i-18+\frac{1}{2}\right)\right) & \text { for } i=24 \text { to } 29 \\ 0 & \text { for } i=30 \text { to } 35\end{cases}
$$

Block type "stop"

$$
z_{i}= \begin{cases}0 & \text { for } i=0 \text { to } 5 \\ x_{i}^{\prime} \sin \left(\frac{\pi}{12}\left(i-6+\frac{1}{2}\right)\right) & \text { for } i=6 \text { to } 11 \\ x_{i}^{\prime} & \text { for } i=12 \text { to } 17 \\ x_{i}^{\prime} \sin \left(\frac{\pi}{36}\left(i+\frac{1}{2}\right)\right) & \text { for } i=18 \text { to } 35\end{cases}
$$

## Block type "short"

The block of 36 samples is divided into three overlapping blocks:

$$
\begin{aligned}
& y_{i}^{(0)}=x_{i+6}^{\prime} \text { for } \mathrm{i}=0 \text { to } 11 \\
& y_{i}^{(1)}=x_{i+12}^{\prime} \text { for } \mathrm{i}=0 \text { to } 11 \\
& y_{i}^{(2)}=x_{i+18}^{\prime} \text { for } \mathrm{i}=0 \text { to } 11
\end{aligned}
$$

Each of the three small blocks is windowed separately:

$$
z_{i}^{(k)}=y_{i}^{(k)} \sin \left(\frac{\pi}{12}\left(i+\frac{1}{2}\right)\right) \text { for } i=0 \text { to } 11, \text { for } k=0 \text { to } 2
$$

## MDCT:

In the following n is the number of windowed samples. For short blocks n is 12 , for long blocks n is 36 . The analytical expression of the MDCT is:

$$
x_{i}=\sum_{k=0}^{n-1} z_{k} \cos \left(\frac{\pi}{2 n}\left(2 k+1+\frac{n}{2}\right)(2 i+1)\right) \quad \text { for } i=0 \text { to } \frac{n}{2}-1
$$

## Aliasing-Butterfly, Encoder:

The calculation of aliasing reduction in the encoder is performed as in the decoder. The general procedure is shown in figure A.5. The butterfly definition to be used in the encoder is shown in figure C.8. The coefficients $\mathrm{ca}_{\mathrm{i}}$ and $\mathrm{cs}_{\mathrm{i}}$ can be found in table B.9.


Figure C. 8 -- Encoder Butterfly Definition

## C.1.5.3.4 Calculation of average available bits

The average number of bits per granule is calculated from the frame size. The bitrate $64 \mathrm{kbits} / \mathrm{s}$ is used a: an example. At bitrate $64 \mathrm{kbjits} / \mathrm{s}$ at 48000 samples per second,
( $64000^{*}(1152 / 48000)$ bits per frame) / ( 2 granules per frame) $=768$ bits per granule.
As the header takes 32 bits and the side information takes 17 bytes ( 136 bits) in single_channel mode, the average amount of available bits for the main_data for a granule is given by
mean_bits $=768$ bits per granule $-(32+136$ bits per frame $) /(2$ granules per frame $)=684$ bits per granule.

## Bit reservoir:

The bit reservoir can provide additional bits which may be used for the granule. The number of bits whic. are provided is determined within the iteration loops.

## C.1.5.3.5 Quantization and encoding of frequency domain samples

The frequency domain data are quantized and coded within two nested iteration loops. Sublause C.1.5.4 contains a detailed description of these iteration loops.

## C.1.5.3.6 Ancillary data

The Audio Standard provides a number of bits for the inclusion and transmission of variable length ancillary dala with the audio bitstream. The ancillary data will reduce the number of bits available for audio, which may result in a degradation of audio quality.

The presence of a bit pattern in the ancillary data matching the syncword may hamper synchronization. This problem is more likely to occur when the free format is used.

## C.1.5.3.7 Formatting

The details about the Layer III bitstream format can be found in 2.4.4. The formatting of the Huffman coce words is described below:

The Huffiman code words are in sequence from low to high frequencies. In the iteration loops the following variables have been calculated and are used in encoding the Huffiman code words:

$$
\begin{array}{ll}
\begin{array}{l}
\text { is(i), i=0...575 } \\
\text { table_selectregion] } \\
\text { region_adress1 } \\
\text { region_adress2 } \\
\text { max_value[region] }
\end{array} & \begin{array}{l}
\text { quantized frequency domain values } \\
\text { Huffman code table used for regions (region }=0,1,2) \\
\text { defines the border between region 0 and 1 } \\
\text { defines the border between region 1 and 2 }
\end{array} \\
\text { maximum absolute value of quantized data in regions (region }=0,1,2 \text { ) } \\
\text { The data are written to the bitstream according to the Huffman code syntax described in } 2.4 .2 .7
\end{array}
$$

```
for region number from 0 to 2
    if table_select for this region is 0
        nothing to do, all values in region are zero
    else
        if table_select for this region is > 15
            an ESC-table is used: look up linbits value connected to the table used
            for i = begin of region to end of region, count in pairs
                x= is(i), y= is(i+1)
            if }x>1
                linbitsx = x - 15,x=15
            end if
            signx = sign(x), x = abs(x)
            if y>14
                linbitsy = y - 15, y = 15
            end if
            signy = sign(y), y = abs(y)
            look for codeword = hcod([x][y]]) in table table_select
            write hcod([x][y]), beginning with the leftmost bit, number of bits is hlen([x][y])
            if }x>1
                    write linbitsx to the bitstream, number of bits is linbits
            end if
            if x != 0
                write signx to bitstream
            end if
            if y>14
                write linbitsy to the bitstream, number of bits is linbits
            end if
            if y !=0
                    write signy to bitstream
            end if
        end do
        else
            no ESC-words are used in this region:
            for i= begimning of region to end of region, count in pairs
            x= is(i), y= is(i+1)
            signx = sign(x), x = abs(x)
            signy = sign(y), y=abs(y)
            look for codeword =hcod([x][y]) in table table_seletct
            write hcod([x][y]), beginning with the leftmost bit, number of bits is hlen([x][y])
            if }x!=
                    write signx to bitstream
            end if
            if y !=0
                write signy to bitstream
            end if
        end do
        end if
    end if
end for
```

A possible application for the private_bits is to use them as frame counter.

## C.1.5.4 Layer III iteration loops

## C.1.5.4.1 Introduction

The description of the Layer III loop module is subdivided into three levels. The top level is called "loops frame program". The loops frame program calls a subroutine named "outer iteration loop" which calls the subroutine "inner iteration loop". For each level a corresponding flow diagram is shown.

The loops module quantizes an input vector of spectral data in an iterative process according to several demands. The inner loop quantizes the input vector and increases the quantizer step size until the output vector can be coded with the available number of bits. After completion of the inner loop an outer loop checks the distortion of each scalefactor band and, if the allowed distortion is exceeded, amplifies the $\mathrm{s} a \mathrm{alefactor}$ band and calls the inner loop again.

Layer III loops module input:
(1) vector of the magnitudes of the spectral values $\operatorname{xr}(0 . .575)$.
(2) $\quad \mathrm{xmin}(\mathrm{sb})$, the allowed distortion of the scalefactor bands. $\mathrm{xmin}=\mathrm{ratio}(\mathrm{sb})^{*} \mathrm{en}(\mathrm{sb}) / \mathrm{bw}(\mathrm{sb})$.
(3) window_switching_flag which, in conjunction with mixed_block_flag and block_type, determines the number of scalefactor bands.
(4) mean_bits (bit available for the Huffman coding and the coding of the scalefactors).
(5) more_bits, the number of bits in addition to the average number of bits, as demanded by the val ue of the psychoacoustic entropy for the granule:
more_bits $=3.1 *$ PE - (average number of bits)
Layer III loops module output:
(1) vector of quantized values ix(0..575).
(2) scalefac_1(sb) or scalefac_s(sb) depending on window_switching_flag, block_type and mixed_block_flag.
(3) global gain (quantizer step size information)
global_gain $=$ qquant + system_constant.
system_constant includes all the scaling operations of the encoder and an offset to achieve the correct output with the decoding process described in the main part.
(4) number of unused bits available for later use.
(5) preflag (loops preemphasis on/off).
(6) Huffman code related side information

- big_values (number of pairs of Huffman coded values, excluding "count1")
- count1 table_select (Huffman code table of absolut values $<=1$ at the upper end of the spectrum
- table_select[0..2](Huffman code table of regions)
- region0_count, region1_count (used to calculate boundaries between regions)
- part2_3_length


## C.1.5.4.2 Preparatory steps

## C. 1.5.4.2.1 Reset of all iteration variables

The scalefactors of the coder partitions, scalefac_1[sb] or scalefac_s[sb], are respectively set to zero.
The counter qquant for the quantizer step size is reset to zero.
Prifflag is reset to zero.
Scalefac_scale is reset to zero.
The inital value of quantanf is set as follows:
quantanf $=$ system_const $* \log ($ (sfm) ,
where sfm is the spectral flatness measure and quantanf depends on the computational implementation of the encoder.

The spectral flatness measure sfm is given by

$$
\operatorname{sfm}=\frac{e^{\frac{1}{n}\left(\sum_{i=0}^{0-1} \log x r(i)^{2}\right)}}{\frac{1}{n} \sum_{i=0}^{n-1} x r(i)^{2}}
$$

The value of system_const is chosen so that for all signals the first iteration of the inner loop for all signals comes out with a bit sum higher than the desired bitsum. By that it is ensured that the first call of the inner loop results in the solution which uses as many of the available bits as possible. In order to spare computing time it is desirable to minimize the number of iterations by adapting the value of quantanf to the bitrate and the signal statistics.

## C.1.5.4.2.2 Bit reservoir controf

Bits are saved to the reservoir when fewer than the mean_bits are used to code one granule. If bits are saved for a frame, the value of main_data_end is increased accordingly. See figure A.7.a.

The number of bits which are made available for the main_data (called "max_bits") is derived from the actual estimated threshold (the PE as calculated by the psychoacoustic model), the average number of bits (mean_bits) and the actual content of the bit reservoir. The number of bytes in the bit reservoir is given by main_data_end.

The actual rules for the control of the bit reservoir are given below:

- If a number of bytes available to the inner iteration loop is not used for the Huffman encoding or other main_data, the number is added to the bit reservoir.
- If the bit reservoir contains more than 0,8 times the maximum allowed content of the bit reservoir, all bytes exceeding this number are made available for main_data (in addition to mean_bits)
- If more_bits is greater than 100 bits, then $\max$ (more_bits/ $8,0,6 *$ main_data_end) bytes are taken from the bit reservoir and made available for main_data (in addition to mean_bits).
- After the actual loops computations have been completed, the number of bytes not used for main_data is added to the bit reservoir.
- If after the step above the number of bytes in the bit reservoir exceeds the maximum allowed content, stuffing bits are written to the bitstream and the content of the bit reservoir is adjusted accordingly.


## C.1.5.4.2.3 Calculation of the scalefactor selection information (scfsi)

The scfsi contains the information, which scalefactors (grouped in the scfsi_bands) of the first granule can also be used for the second granule. These scalefactors are therefore not transmitted, the bits gained can be used for the Hufíman coding.
To determine the usage of the scfsi, the following information of each granule must be stored:
a) The block type
b) The total energy of the granule:

$$
\text { en_tot }=\text { int }\left\{\log 2\left(\sum_{i=1}^{n}|x r(i)|^{2}\right)\right\}
$$

where n is the total number of spectral values.
c) The energy of each scalefactor band:

$$
\operatorname{en}(s b)=\operatorname{int}\left\{\log _{2}\left(\sum_{i=l \mathrm{bl}(\mathrm{sb})}^{\mathrm{lbl}(\mathrm{sb})+\mathrm{b} w(\mathrm{sb})-1}|\mathrm{xr}(\mathrm{i})|^{2}\right)\right\}
$$

where $\mathrm{lbl}(\mathrm{sb})$ is the number of the first coefficient belonging to scalefactor band sb and $\mathrm{bw}(\mathrm{sb})$ is the number of coefficients within scalefactor band sb.
d) The allowed distortion of each scalefactor band:

$$
\mathrm{xm}(\mathrm{sb})=\text { int }\left\{\log _{2}(\mathrm{xmin}(\mathrm{i}))\right\}
$$

$\mathrm{x} \min (\mathrm{sb})$ is calculated by the psychoacoustic model.

The scalefactors of the first granule are always transmitted. When coding the second granule, the information of the two granules is compared. There are four criteria to determine if the scfsi can be used in general. If one of the four is not fulfilled, the scfsi is disabled (that means it is set to 0 in all scfsi_bands). The criteria are (index 0 means first, index 1 second granule):
a) The spectral values are not all zero
b) None of the granules contains short blocks
c)

$$
\left|e n \_t o t_{0}-e^{2} \_t o t_{1}\right|<e^{2} \_ \text {tot }{ }_{\text {krit }}
$$

d)

$$
\sum_{\text {all scalefactor bands }}\left|\mathrm{en}(\mathrm{sb})_{0}-\mathrm{en}(\mathrm{sb})_{1}\right|<\text { en_dif }{ }_{\text {krit }}
$$

If the scfsi is not disabled after the tests above, there are two criteria for each scfsi_band, which have both to be fulfilled to enable scfsi (that means to set it to 1 in this scfsi_band):
a)

b)
$\sum_{\text {ala.boiostri_baxd }}\left|\mathrm{xm}(\mathrm{sb})_{0}-\mathrm{xm}(\mathrm{sb})_{1}\right|<\mathrm{xm}\left(\mathrm{scfsi} \mathrm{I}_{1} \text { band }\right)_{\text {krit }}$
The constants (with the index krit) have to be chosen so, that the scfsi is only enabled in case of similar energy/distortion.
Suggested values are:

| en_tot krit | $=$ | 10 |  |  |
| :--- | :--- | :--- | :--- | :--- |
| en_dif $_{\text {krit }}$ | $=$ | 100 |  |  |
| en(scfsi_band) | krit | $=$ | 10 | for each scfsi_band |
| xm(scfsi_band) | krit | $=$ | 10 | for each scfsi_band |

## C.1.5.4.3 Outer iteration loop (distortion control loop)

The outer iteration loop controls the quantization noise which is produced by the quantization of the frequency domain lines within the inner iteration loop. The colouring of the noise is done by multiplication of the lines within scalefactor bands with the actual scalefactors before doing the quantization. The following pseudo-code illustrates the multiplication.
do for each scalefactor band:
do from lower index to upper index of scale factor band $\operatorname{xr}(\mathrm{i})=\operatorname{xr}(\mathrm{i}) * \sqrt{ }(2)^{\wedge}((1+$ scalefac_scale $) *$ scalefac(sb) )
end do
end do
Where scalefac is either scalefac_1 or scalefac_s as appropriate.

In the actual system the multiplication is done incrementally with just the increase of the scalefactors applied in each distortion control loop. This is described in C.1.5.4.3.5 below.

The distortion loop is always started with scalefac_scale $=0$. If after some iterations the maximum length of the scalefactors would be exceeded (see scalefac_compress table in 2.4.2.7 and C.1.5.4.3.5 below), then scalefac_scale is increased to the value 1 thus increasing the possible dynamic range of the scalefactors. In this case the actual scalefactors and frequency lines have to be corrected accordingly.

## C.1.5.4.3.1 Saving of the scalefactors

The scalefactors of all scalefactor bands, scalefac_1(sb) or scalefac_s(sb), as well as the quantizer step size qquant are saved. If the computation of the outer loop is cancelled without having reached a proper result, this value together with the quantized spectrum give an approximation and can be transmitted.

## C.1.5.4.3.2 Call of inner iteration loop

For each outer iteration loop (distortion control loop) the inner iteration loop (rate control loop) is called. The parameters are the frequency domain values (hybrid filterbank output) with the scalefactors applied to the values within the scalefactor bands and the number of bits which are available to the rate control loop. The result is the number of bits actually used and the quantized frequency lines ix(i).

## C.1.5.4.3.3 Calculation of the distortion of the scalefactor bands

For each scalefactor band the actual distortion is calculated according to:

$$
\operatorname{xfsf}(s b)=\sum_{\mathrm{lbl}(\mathrm{sb})}^{\mathrm{i}=\mathrm{lbl}(\mathrm{sb})+\mathrm{bw}(\mathrm{sb})-1} \frac{\left(|\operatorname{xr}(\mathrm{i})|-\mathrm{ix}(\mathrm{i})^{4 / 3} * \sqrt[4]{2}^{\text {qquant+quantanf }}\right)^{2}}{\text { bandwidth }(\mathrm{sb})}
$$

where lbl(sb) is the number of the coefficient representing the lowest frequency in a scalefactor band and $\mathrm{bw}(\mathrm{sb})$ is the nu:nber of coefficients within this band.

## C. 1.5.4.3.4 Preemphasis

The preemphasi: option (switched on by setting preflag to a value of 1) provides the possibility to amplifiy the upper part of the spectrum according to the preemphasis tables, table B.6.

```
if (preflag=1) {
    ifqstep = 2^ (0,5*(1+scalefac_scale))
    xmin}(\textrm{j})=x\operatorname{min}(\textrm{j}) *ifqstep ^ (2* prefact(j))
    for (i=lower limit of scalefactor band j; i <=upper limit of scalefactor band j; i++) {
        xr(i) = xr(i) * ifqstepprefact(j)
    }
)
```

The condition to switch on the preemphasis is up to the implementation. For example preemphasis could be switched on $i f$ in all of the upper 4 scalefactor bands the actual distortion exceeds the threshold after the first call of the inner loop.

If the second grenule is being coded and scfsi is active in at least one scfsi_band, the preemphasis in the second granule is set equal to the setting in first granule.

## C.1.5.4.3.5 Amplification of scalefactor bands which violate the masking threshold

All spectral values of the scalefactor bands which have a distortion that exceeds the allowed distortion are amplified by a factor of ifqstep. The value of ifqstep is transmitted by scalefac_scale.

```
if ( \((x \min -x f s f)\) of scalefactor band \(\mathrm{j}<0)\) \{
    \(x \min (\mathrm{j})=\mathrm{xmin}(\mathrm{j}) *\) ifqstep \({ }^{\wedge} 2\)
    if \(q(j)=i f q(j)+1\)
    for ( \(\mathrm{i}=\) =lower limit of scalefactor band; \(\mathrm{i}<=\) upper limit of scalefactor band; \(\mathrm{i}++\) ) (
        \(\operatorname{xr}(\mathrm{i})=\mathrm{xr}(\mathrm{i}) *\) ifqstep
    \}
\}
```

If the second granule is being coded and scfsi is active in at least one scfsi_band, the following steps have to be done:
a) ifqstep has to be set similar to the first granule
b) If it is the first iteration, the scalefactors of scalefactor bands in which scfsi is enabled bave to be taken over from the first granule. The corresponding spectral values have to be amplified:

```
if ( scfsi according to scalefactor band j == 1) {
    ifq(j)=ifq(j) first granule
    for (i=lower limit of scalefactor band; i <=upper limit of scalefactor band; i++) {
                xr(i)= xr(i) * ifqstep ^ scalefac(j)
    }
}
```

Where scalefac is either scalefac_1 0 or scalefac_s 0 as appropriate.
c) If it is not the first iteration, the amplification must be prevented for scalefactor bands in which scfsi is enabled.

## C.1.5.4.3.6 Conditions for the termination of the loops processing

Normally the loops processing terminates if there is no scalefactor band with more than the allowed distortion. However this is not always possible to obtain. In this case there are other conditions to terminate the outer loop. If
a) All scalefactor beands are already amplified, or
b) The amplification of at least one band exceeds the upper limit which is determined by the transmission format of the scalefactors. The upper limit is a scalefactor of 15 for scalefactor bands 0 through 10 and 7 for scalefactors 11 through 20. In the case of block_type $==2$ and mixed_block_flag $==0$, the upper limit is 15 for scalefactors 0 through 18. In the case of block_type $==2$ and mixed_block_flag $==1$, the upper limit is 15 for scalefactors 0 through 17. The upper limit is 7 for other scalefactors.

The loop processing stops, and by restoring the saved sacalefac_1(sb) or scalefac_s(sb) a useful output is available. For realtime implementation, there might be a third condition added which terminates the loops in case of a lack of computing time.

## C.1.5.4.4 Inner iteration loop (rate control loop)

The inner iteration loop does the actual quantization of the frequency domain data and prepares the formatting. The table selection, subdivision of the big_values range into regions and the selection of the quantizer step size takes place here.

## C.1.5.4.4.1 Quantization

The quantization of the complete vector of spectral values is done according to

$$
\operatorname{ix}(\mathrm{i})=\operatorname{nint}\left(\left(\frac{|\operatorname{xr}(\mathrm{i})|}{\sqrt{2}^{\text {qquant }+ \text { quantanf }}}\right)^{0.75}-0,0946\right)
$$

## C.1.5.4.4.2 Test of the maximum of the quantized values

The maximum allowed quantized value is limited. This limit is set to constraint the table size if a tablelookup is used to requantize the quantized frequency lines. The limit is given by the possible values of the length identifier, "linbits", of values flagged with an ESC-code. Therefore before any bit counting is done the quantizer st:psize is increased by
qquant $=$ qquant +1
until the maximum of the quantized values is within the range of the largest Huffman code table.

## C.1.5.4.4.3 Calculation of the run length of zeros

The run length rzero of pairs of spectral coefficients quantized to zero on the upper end of the spectrum is counted and called "rzero".

## C.1.5.4.4.4 Calculation of the run length of values less or equal one

The run length of quadrupels of spectral coefficients quantized to one or zero, following the rzero pairs of zeros, is calculated and called "count1".

## C.1.5.4.4.5 Counting the bits necessary to code the values less or equal one

One Huffman code word is used to code one of the "count1" quadrupels. There are two different Huffman code books with corresponding code length tables (table A and table B in clause B.7). The number of bits to code all the countl quadrupels is given by:
bitsum_count $1=\min ($ bitsum_table 0, bitsum_table1 $)$
where count1table_0 is used to point to table A

and count1table_1 is used to point to table B
$\mathrm{k}=$ firstcount $1+$ count $1-1$
bitsum_table1 $==\quad \sum_{k=\text { firstcount1 }}$ count1table_1 $(\mathrm{ix}(4 \mathrm{k})+2 * \mathrm{ix}(4 \mathrm{k}+1)+4 * \mathrm{ix}(4 \mathrm{k}+2)+8 * \mathrm{ix}(4 \mathrm{k}+3))$
Count1table_0 as well as count1table_1 have to include the number of bits necessary to encode the sign bits.

The information which table is used is transmitted by countitable_select, which is " 0 " for table A or " 1 " for table $B$, respectively.

## C.1.5.4.4.6 Call of subroutine SUBDIVIDE

The number of pairs of quantized values not counted in "count1" or "rzero" is called bigvalues. SUBDIVIDE splits the scalefactor bands corresponding to this values into three groups. The last one, incomplete generally, counts as a complete one. The number of scalefactor bands in the first and second regions are contained in (region0_count+1) and (region 1_count+1) respectively. The number of scalefactor bands in the third region can be calculated using bigvalues. The split strategy is up to the implementation. A very simple one for instance is to assign $1 / 3$ of the scalefactor bands to the first and $1 / 4$ to the last region.

Subdivide in case of blocksplit is done analoguously but there are only two subregions. Region1_count is set to a default in this case. This default is 8 in the case of split_point $=0$ and 9 in the case of split_point $=1$. Both these values point to the same absolute frequency.

## C.1.5.4.4.7 Calculation of the code book for each subregion

There are 32 different Huffman code tables available for the coding of pairs of quantized values. They diff $: r$ from each other in the maximum value that can be coded and in the signal statistics for which they are optimized. Only codes for values $<16$ are in the table. For values $>=16$ there are two tables provided, wl ere the largest value 15 is an escape character. In this case the value 15 is coded in an additional word using i. linear PCM code with a word length called linbits.

A simple way to choose a table is to use the maximum of the quantized values in a subregion. tables which have the same size are optimized for different signal statistics. Therefore additional coding gain car be achieved for example by trying all of these tables.

## C.1.5.4.4.8 Counting of the bits necessary to code the values in the subregions

The number of bits necessary to code the quantized values of a subregion is given by:

$$
\begin{aligned}
\operatorname{bitsum}(j) & =\sum_{k=0}^{k=n n(j)-1} \operatorname{bitz}(\text { tableselect }(j), \min (15, \operatorname{ix}(2 k+f e(j))), \min (15, \operatorname{ix}(2 k+f e(j)+1))) \\
& +\sum_{k=0}^{n p(j)-1}(s(i x(2 k+f e(j))-15)+s(i x(2 k+f e(j)+1)-15)) * \operatorname{linbits}(j)
\end{aligned}
$$

$\mathrm{np}(\mathrm{j})$ : number of pairs in a sub region
fe(j): number of the first quantized value in a sub-region
bitz: table with Huffman code length
$s(\ldots) \quad$ step function: if $x>=0 \quad s(x)=1$
if $x<0 \quad s(x)=0$
Note that the Huffman code length tables have to include the number of bits necessary to encode the sign bits.


Figure C.9.a .- Layer III iteration loop


Figure C.9.b -. Layer III outer iteration loop


Figure C.9.c -- Layer III inner iteration loop

## Annex D

(informative)

## Psychoacoustic models

## D.1. Psychoacoustic model 1

The calculation of the psychoacoustic model has to be adapted to the corresponding layer. This example i: valid for Layers I and II. The model can be adapted to Layer III.

There is no principal difference in the application of psychoacoustic model 1 to Layer I or II.
Layer I: A new bit allocation is calculated for each block of 12 subband or 384 input PCM samples.

Layer I: A new bit allocation is calculated for three blocks totaling 36 subband samples corresponding to $3^{*} 384$ (1 152) input PCM samples.

The bit allocation of the 32 subbands is calculated on the basis of the signal-to-mask ratios of all the subbands. Therefore, it is necessary to determine for each subband, the maximum signal level and the minimum masking threshold. The minimum masking threshold is derived from an FFT of the input PCMI signal, followed by a psychoacoustic model calculation.

The FFT in parallel with the subband filter compensates for the lack of spectral selectivity obtained at lou frequencies by the subband filterbank. This technique provides both a sufficient time resolution for the codsd audio signal (Polyphase filter with optimized window for minimal pre-echoes) and a sufficient spectral resolution for the calculation of the masking thresholds. The frequencies and levels of aliasing distortions can be calculated. This is necessary for calculating a minimum bitrate for those subbands which need som: bits to cancel the aliasing components in the decoder. The additional complexity to calculate the better frequency resolution is necessary only in the encoder, and introduces no additional delay or complexity in the decoder.

The calculation of the signal-to-mask-ratio is based on the following steps:
Step 1

- Calculation of the FFT for time to frequency conversion.

Step 2

- Determination of the sound pressure level in each subband.

Step 3

- Determination of the threshold in quiet (absolute threshold).

Step 4

- Finding of the tonal (more sinusoid-like) and non-tonal (more noise-like) components of the audio signa .

Step 5

- Decimation of the maskers, to obtain only the relevant maskers.

Step 6

- Calculation of the individual masking thresholds.

Step 7

- Determination of the global masking threshold.

Step 8

- Determination of the minimum masking threshold in each subband.

Step 9

- Calculation of the signal-to-mask ratio in each subband.

These steps will be further discussed. A sampling frequency of 48 kHz is assumed. For the other two sampling frequeacies all frequencies mentioned should be scaled accordingly.

## Step 1: $\mathrm{FFT}^{-}$Analysis

The masking threshold is derived from an estimate of the power density spectrum that is calculated by a 512 -point FFT for Layer I, or by a 1024 -point FFT for Layer II. The FFT is calculated directly from the input PCM signal, windowed by a Hann window.

For a coincidence in time between the bit allocation and the corresponding subband samples, the PCMsamples entering: the FFT have to be delayed:
a) The delay of the analysis subband filter is 256 samples, corresponding to $5,3 \mathrm{~ms}$ at the 48 kHz sampling rate. A window shift of 256 samples is required to compensate for the delay in the analysis subband iliter.
b) The Hann window must coincide with the subband samples of the frame. For Layer I this amounts to an additional window shift of 64 samples. For Layer II an additional window shift of minus 64 samples is required.

Technical data of the FFT:

|  | Layer I | Layer II |
| :--- | :---: | :---: |
| - transform length | 512 samples | 1024 samples |
| Windo'v size if fs $=48 \mathrm{kHz}$ | $10,67 \mathrm{~ms}$ | $21,3 \mathrm{~ms}$ |
| Windo'v size if fs $=44,1 \mathrm{kHz}$ | $11,6 \mathrm{~ms}$ | $23,2 \mathrm{~ms}$ |
| Window size if fs $=32 \mathrm{kHz}$ | 16 ms | 32 ms |
| Frequency resolution | sampling_frequency / 512 | sampling_frequency / 1024 |

- Hann window, h(i):
$h(i)=\sqrt{8 / 3} * 0,5 *\{1-\cos [2 * \pi *(i) / N]\} \quad 0<=\mathbf{i}<=N-1$
- power density spectrum $X(\mathrm{k})$ :
$X(k)=10 * \log _{10}\left|\frac{1}{N} \sum_{l=0}^{N-1} h(1) * s(1) * e^{(-j * k * 1 * 2 * \pi / N)}\right|^{2} d B \quad k=0 \ldots N / 2$,
where :(1) is the input signal.
A normalization to the reference level of 96 dB SPL (Sound Pressure Level) has to be done in such a way that the maximum value corresponds to 96 dB .


## Step 2: Determination of the sound pressure level

The sound pressure level $L_{s b}$ in subband $n$ is computed by:
$\mathrm{L}_{\mathrm{sb}}(\mathrm{n})=\operatorname{MAX}\left[\underset{\mathrm{M}}{\mathrm{X}(\mathrm{k}), 20 * \log ) \text { in subband } \mathrm{n}}\left(\mathrm{scf}_{\max }(\mathrm{n}) * 32768\right)-10\right] \mathrm{dB}$
where $X(k)$ is the sound pressure level of the spectral line with index $k$ of the FFT with the maximum amplitude in the frequency range corresponding to subband $n$. The expression scfmax $(n)$ is in Layer I the scalefactor, and in Layer II the maximum of the three scalefactors of subband $n$ within a frame. The " -10 $d B "$ term correc s for the difference between peak and RMS level. The sound pressure level $L_{s b}(n)$ is computed for every subband n .

The following alternative method of calculating $\mathrm{L}_{\mathrm{sb}}(\mathrm{n})$ offers a potential for better encoder performance, but this technique has not been subjected to a formal audio quality test.

The alternative sound pressure level $L_{s b}$ in subband $n$ is computed by:

```
\(\mathrm{L}_{\mathrm{s} \cdot}(\mathrm{n})=\quad \operatorname{MAX}\left[\mathrm{X}_{\mathrm{spl}}(\mathrm{n}), 20 * \log \left(\mathrm{scf}_{\max }(\mathrm{n}) * 32768\right)-10\right] \mathrm{dB}\)
with
```


where $\mathbf{X}_{\text {spl }}(\mathrm{n})$ is the alternative sound pressure level corresponding to subband n .

## Siep 3: Considering the threshold in quiet

The threshold in quiet $\mathrm{LT}_{\mathrm{q}}(\mathbf{k})$, also called absolute threshold, is available in the tables "Frequencies, critic al band rates and absolute threshold" (tables D.1a, D.1b, D.1c for Layer I; tables D.1d, D.1e, D.1f for Layer II). These tables depend on the sampling rate of the input PCM signal. Values are available for eacs sample in the frequency domain where the masking threshold is calculated. An offset depending on the overall bit rate is used for the absolute threshold. This offset is -12 dB for bit rates $>=96 \mathrm{kbits} / \mathrm{s}$ and 0 dB for bit rates $<96 \mathrm{kbits} / \mathrm{s}$ per channel.

## Step 4: Finding of tonal and non-tonal components

The tonality of a masking component has an influence on the masking threshold. For this reason, it is wcrthwhile to discriminate between tonal and non-tonal components. For calculating the global masking threshold it is necessary to derive the tonal and the non-tonal components from the FFT spectrum.

This step starts with the determination of local maxima, then extracts tonal components (sinusoids) and calculates the intensity of the non-tonal components within a bandwidth of a critical band. The boundaries of the critical bands are given in the tables "Critical band boundaries" (tables D.2a, D.2b, D. 2 c for Layer ;; tat les D.2d, D.2e, D. 2 f for Layer II).

The bandwidth of the critical bands varies with the center frequency with a bandwidth of about only $0,1 \mathrm{kF}[\mathrm{z}$ at low frequencies and with a bandwidth of about 4 kHz at high frequencies. It is known from psychoacoustic experiments that the ear has a better frequency resolution in the lower than in the higher frequency region. To determine if a local maximum may be a tonal component, a frequency range df arouni the local maximum is examined. The frequency range df is given by:

Sanpling rate: 32 kHz

| Laver I: | df $=125 \mathrm{~Hz}$ <br> df $=187,5 \mathrm{~Hz}$ <br> df $=375 \mathrm{~Hz}$ |
| :--- | :--- |
| Layer II: | df $=62,5 \mathrm{~Hz}$ |
|  | df $=93,75 \mathrm{~Hz}$ <br>  <br> df $=187,5 \mathrm{~Hz}$ <br> df $=375 \mathrm{~Hz}$ |


| 0 kHz | $<\mathrm{f}<=$ | $4,0 \mathrm{kHz}$ |
| :---: | :---: | :---: |
| $4,0 \mathrm{kHz}$ | $<\mathrm{f}<$ | $8,0 \mathrm{kHz}$ |
| $8,0 \mathrm{kHz}$ | $<\mathrm{f}$ <= | $15,0 \mathrm{kHz}$ |
| 0 kHz | $<\mathrm{f}<$ | $3,0 \mathrm{kHz}$ |
| $3,0 \mathrm{kHz}$ | $<\mathrm{f}<$ | $6,0 \mathrm{kHz}$ |
| $6,0 \mathrm{kHz}$ | < f | $12,0 \mathrm{kHz}$ |
| $12,0 \mathrm{kHz}$ | $<\mathrm{f}<$ | $24,0 \mathrm{kHz}$ |

Sanpling rate: $44,1 \mathrm{kHz}$


Sampling rate: 48 kHz

| Layer I: | $\begin{aligned} & \text { df }=187,5 \mathrm{~Hz} \\ & \text { df }=281,25 \mathrm{~Hz} \\ & \text { df }=562,50 \mathrm{~Hz} \end{aligned}$ | $\begin{array}{r} 0 \mathrm{kHz} \\ 6,0 \mathrm{kHz} \\ 12,0 \mathrm{kHz} \end{array}$ | $\begin{aligned} & <\mathrm{f}<= \\ & <\mathrm{f}<= \\ & <\mathrm{f}<= \end{aligned}$ | $\begin{array}{r} 6,0 \mathrm{kHz} \\ 12,0 \mathrm{kHz} \\ 24,0 \mathrm{kHz} \end{array}$ |
| :---: | :---: | :---: | :---: | :---: |
| Layer II: | df $=93,750 \mathrm{~Hz}$ | 0 kHz | $<\mathrm{f}<=$ | $3,0 \mathrm{kHz}$ |
|  | df $=140,63 \mathrm{~Hz}$ | $3,0 \mathrm{kHz}$ | $<\mathrm{f}<=$ | 6,0 kHz |
|  | df $=281,25 \mathrm{~Hz}$ | $6,0 \mathrm{kHz}$ | $<\mathrm{f}<$ | $12,0 \mathrm{kHz}$ |
|  | df $=562,50 \mathrm{~Hz}$ | $12,0 \mathrm{kHz}$ | $<\mathrm{f}<=$ | $24,0 \mathrm{kHz}$ |

To make lists of the spectral lines $\mathbf{X}(\mathbf{k})$ that are tonal or non-tonal, the following three operations are performed:

## a) Labelling of local maxima

A spectral line $X(k)$ is labelled as a local maximum if

$$
X(k)>X(k-1) \text { and } X(k)>=X(k+1)
$$

b) Listing of tonal components and calculation of the sound pressure level

A local maximum is put in the list of tonal components if

$$
X(k)-X(k+j)>=7 d B,
$$

where j is chosen according to
Layer I:
$\mathrm{j}=-2,+2$
$\mathrm{j}=-3,-2,+2,+3$
$\mathrm{j}=-6, \ldots,-2,+2, \ldots,+6$
for $2<k<63$
$\mathrm{j}=-3,-2,+2,+3$
for $63<=k<127$
for $127<=k<=250$
Layer II:

| $\mathrm{j}=-2,+2$ | for $\quad 2<k<63$ |  |
| :--- | :--- | ---: |
| $\mathrm{j}=-3,-2,+2,+3$ | for $\quad 63<=k<127$ |  |
| $\mathrm{j}=-6, \ldots,-2,+2, \ldots,+6$ | for | $127<=k<255$ |
| $\mathrm{j}=-12, \ldots,-2,+2, \ldots,+12$ | for | $255<=k<=500$ |

If $\mathrm{X}(\mathrm{k})$ is found to be a tonal component, then the following parameters are listed:
Index number $k$ of the spectral line.

- Sound pressure level $\mathrm{X}_{\mathrm{tm}}(\mathrm{k})=10^{*} \log _{10}\left\{10^{\frac{X(\mathrm{k}-1)}{10}}+10^{\frac{X(\mathrm{k})}{10}}+10^{\frac{X(\mathrm{k}+1)}{10}}\right\}$, in dB
- Tonal flag.

Next, all spectral lines within the examined frequency range are set to $-\infty \mathrm{dB}$.

## c) Listing of non-tonal components and calculation of the power

The non-tonal (noise) components are calculated from the remaining spectral lines. To calculate the nontonal components from these spectral lines $\mathbf{X}(\mathrm{k})$, the critical bands $\mathbf{z}(\mathrm{k})$ are determined using the tables, "Critical band boundaries" (tables D.2a, D.2b, D.2c for Layer I; tables D.2d, D.2e, D. 2 f for Layer II). In Layer I, 23 critical bands are used for the sampling rate of $32 \mathrm{kHz}, 24$ critical bands for $44,1 \mathrm{kHz}$ and 25 critical bands are used for 48 kHz . In Layer II, 24 critical bands are used for 32 kHz sampling rate, and 26 critical bands are used for $44,1 \mathrm{kHz}$ and 48 kHz sampling rate. Within each critical band, the power of the spectral lines (remaining after the tonal components have been zeroed) are summed to form the sound pressure level of the new non-tonal component $X_{n m}(k)$ corresponding to that critical band.

The following parameters are listed:

- Index number $k$ of the spectral line nearest to the geometric mean of the critical band.
- Sound pressure level $X_{\mathrm{nm}}(\mathrm{k})$ in dB.
- Non-tonal flag.


## Step 5: Decimation of tonal and non-tonal masking components

Decimation is a procedure that is used to reduce the number of maskers which are considered for the calculation of the global masking threshold.
a) Tonal $\mathrm{X}_{\mathrm{tm}}(\mathrm{k})$ or non-tonal components $\mathrm{X}_{\mathrm{nm}}(\mathrm{k})$ are considered for the calculation of the masking threshold only if:

$$
\mathrm{X}_{\mathrm{tm}}(\mathrm{k})>=\mathrm{LT}_{\mathrm{q}}(\mathrm{k}) \quad \text { or } \quad \mathrm{X}_{\mathrm{nm}}(\mathrm{k})>=\mathrm{LT}_{\mathrm{q}}(\mathrm{k})
$$

In this expression, $\mathrm{LT}_{q}(\mathrm{k})$ is the absolute threshold (or threshold in quiet) at the frequency of index k . These values are given in tables D.1a, D.1b, D.1c for Layer I; tables D.1d, D.1e, D.1f for Layer II.
b) Decimation of two or more tonal components within a distance of less then 0,5 Bark: Keep the component with the highest power, and remove the smaller component(s) from the list of tonal components. For this operation, a sliding window in the critical band domain is used with a width of 0,5 Bark.

In the following, the index $j$ is used to indicate the relevant tonal or non-tonal masking components from the combined decimated list.

## Step 6: Calculation of individual masking thresholds

Of the original $N / 2$ frequency domain samples, indexed by $k$, only a subset of the samples, indexed by $i$, are considered for the global masking threshold calculation. The samples used are shown in tables D.1a, D.1b, D.1c for Layer I; tables D.1d, D.1e, D. 1 f for Layer II.

## Layer I:

For the frequency lines corresponding to the frequency region which is covered by the first six subbands no subsampling is used. For the frequency region corresponding to the next six subbands every second spectral line is considered. Finally, in the case of $44,1 \mathrm{kHz}$ and 48 kHz sampling rates, in the frequency region corresponding to the remaining subbands, every fourth spectral line is considered up to 20 kHz . In the case of 32 kHz sampling rate, in the frequency region corresponding to the remaining subbands, every fourth spectral line is considered up to 15 kHz (see also tables D.1a, D.1b, D.1c for Layer I).

## Layer II:

For the frequency lines corresponding to the frequency region which is covered by the first three subbands no subsampling is used. For the frequency region which is covered by next three subbands every second spectral line is considered. For the frequency region corresponding to the next six subbands every fourth spectral line is considered. Finally, in the case of $44,1 \mathrm{kHz}$ and 48 kHz sampling rates, in the remaining subbands every eighth spectral line is considered up to 20 kHz . In the case of 32 kHz sampling rate, in the frequency region corresponding to the remaining subbands, every eighth spectral line is considered up to 15 kHz:. (See also tables D.1d, D.1e, D.1f for Layer II).

The number of samples, n , in the subsampled frequency domain is different depending on the sampling rate; and layers.
32 kHz sampling rate:

$$
\begin{aligned}
& \mathrm{n}=108 \text { for Layer I } \\
& \mathrm{n}=106 \text { for Layer I } \\
& \mathrm{n}=102 \text { for Layer } \mathrm{I}
\end{aligned}
$$

and
$\mathrm{n}=132$ for Layer II
$44,1 \mathrm{kHz}$ sampling rate:
and
$\mathrm{n}=130$ for Layer II
48 kHz sampling rate:
and
$\mathrm{n}=126$ for Layer II

Every tonal and non-tonal component is assigned the value of the index i that most closely corresponds to the frequency of the original spectral line $X(k)$. This index $i$ is given in tables D.1a, D.1b, D.1c for Layer I; tables D.1d, D.1e, D.1f for Layer II.

The individual masking thresholds of both tonal and non-tonal components are given by the following expression:

$$
\begin{gathered}
\text { Licensed to Venable LLP / Barbara Folensbee-Moore (Iib01@venable.corm) } \\
\text { ISO Store Order; OP-149809 / Downloaded: } 2016-06-20 \\
\text { Single user licence only, copying and networking prohibited. }
\end{gathered}
$$

| $\mathrm{LT}_{t m}[\mathrm{z}(\mathrm{j}), \mathrm{z}(\mathrm{i})]$ |  |
| :--- | :--- |
| $\mathrm{LT}_{\mathrm{nm}}[\mathrm{z}(\mathrm{j}), \mathrm{z}(\mathrm{i})]$ | $=\mathrm{X}_{\mathrm{tm}}[\mathrm{z}(\mathrm{j})]+\mathrm{av}_{\mathrm{tm}}[\mathrm{z}(\mathrm{j})]+\mathrm{vf}[\mathrm{z}(\mathrm{j}), \mathrm{z}(\mathrm{i})] \mathrm{dB}$ |
| $=$ | $\mathrm{X}_{\mathrm{nm}}[\mathrm{z}(\mathrm{j})]+\mathrm{av}_{\mathrm{nm}}[\mathrm{z}(\mathrm{j})]+\mathrm{vf}[\mathrm{z}(\mathrm{j}), \mathrm{z}(\mathrm{i})] \mathrm{dB}$ |

In this formula, $\mathrm{LT}_{\mathrm{tm}}$ and $\mathrm{LT}_{\mathrm{nm}}$ are the individual masking thresholds at critical band rate z in Bark of the masking component at the critical band rate of the masker $\mathrm{z}_{\mathrm{m}}$ in Bark. The values in dB can be either positive or negative. The term $\mathrm{X}_{\mathrm{tm}}[\mathrm{z}(\mathrm{j})]$ is the sound pressure level of the masking component with the index number $j$ at the corresponding critical band rate $z(j)$. The term av is called the masking index and vf the masking function of the masking component $X_{t m}[z(j)]$. The masking index av is different for tonal and non-tonal maskers ( $a v_{\mathrm{tm}}$ and $a v_{\mathrm{nm}}$ ).

For tonal maskers it is given by
$a v_{t m}=-1,525-0,275 * z(j)-4,5 d B$,
and for non-tonal maskers
$a v_{n m}=-1,525-0,175 * z(j)-0.5 d B$.
The masking function vf of a masker is characterized by different lower and upper slopes, which depend on the distance in Bark $d z=z(i)-z(j)$ to the masker. In this expression $i$ is the index of the spectral line at which the masking function is calculated and $j$ that of the masker. The critical band rates $z(j)$ and $z(i)$ can be found in tables D.1a, D.1b, D.1c for Layer I; tables D.1d, D.1e, D.1f for Layer II. The masking function, which is the same for tonal and non-tonal maskers, is given by:
$\mathrm{vf}=17$ * $(\mathrm{dz}+1)-\left(0,4^{*} \mathrm{X}[\mathrm{z}(\mathrm{j})]+6\right) \mathrm{dB}$
for $-3<=d z<-1$ Bark
$\mathrm{vf}=(0,4 * X[\mathrm{z}(\mathrm{j})]+6) * \mathrm{dz} \mathrm{dB}$
for $-1<=\mathrm{dz}<0$ Bark
$\mathrm{vf}=-17^{*} \mathrm{dz} \quad \mathrm{dB}$
for $0<=\mathrm{dz}<1$ Bark
$\mathrm{vf}=-(\mathrm{dz}-1)^{*}(17-0,15 * X[\mathrm{z}(\mathrm{j})])-17 \mathrm{~dB}$
for $1<=\mathrm{dz}<8$ Bark
In these expressions $\mathrm{X}[\mathrm{z}(\mathrm{j})]$ is the sound pressure level of the $\mathrm{j}^{\prime}$ th masking component in dB . For reasons of implementation complexity, the masking is no longer considered ( $\mathrm{LT}_{\mathrm{tm}}$ and $\mathrm{LT}_{\mathrm{nm}}$ are set to $-\infty \mathrm{dB}$ outside this range) if $\mathrm{dz}<-3$ Bark, or $\mathrm{dz}>=8$ Bark.

## Step 7: Calculation of the global masking threshold LTg

The global masking threshold $\mathrm{LT}_{\mathrm{g}}(\mathrm{i})$ at the $\mathrm{i}^{\prime} t \mathrm{th}$ frequency sample is derived from the upper and lower slopes of the individual masking thresholds of each of the $j$ tonal and non-tonal maskers and from the threshold in quiet $\mathrm{LT}_{\mathrm{q}}(\mathrm{i})$. This is also given in tables D.1a, D.1b, D.1c for Layer I; tables D.1d, D.1e, D.1f for Layer II. The global masking threshold is found by summing the powers corresponding to the individual masking thresholds and the threshold in quiet.
$L T_{g}(\mathrm{i})=10 \log _{10}\left(10^{L T_{q}(i) / 10}+\sum_{j=1}^{m} 10^{L T_{t m}(z(j), z(i)) / 10}+\sum_{j=1}^{n} 10^{L T_{n m}(z(j), z(i)) / 10}\right)$
The total number of tonal maskers is given by m, and the total number of non-tonal maskers is given by $n$. For a given $i$, the range of $j$ can be reduced to just encompass those masking components that are within -8 to +3 Bark from i . Outside of this range $\mathrm{LT}_{\mathrm{tm}}$ and $\mathrm{LT}_{\mathrm{nm}}$ are $-\infty \mathrm{dB}$.

## Step 8: Determination of the minimum masking threshold

The minimum masking level $\operatorname{LTmin}(n)$ in subband $n$ is determined by the following expression:
$L T_{\text {min }}(\mathrm{n})=\underset{\mathrm{f}}{\mathrm{i})} \underset{\operatorname{Min}\left[L T_{\mathrm{g}}(\mathrm{i})\right] \mathrm{dB}}{ }$ $f(i)$ in subband $n$
where $f(i)$ is the frequency of the $i$ th frequency sample. The $f(i)$ are tabulated in tables D.1a, D.1b, D.1c for Layer I; tables D.1d, D.1e, D.1f for Layer II. A minimum masking level LT $_{\min }(\mathrm{n})$ is computed for every subband.

Step 9: Calculation of the signal-to-mask-ratio
The signal-to-mask ratio
$\operatorname{SMR}_{\mathrm{sb}}(\mathrm{n})=\mathrm{L}_{\mathrm{sb}}(\mathrm{n})-\mathrm{LT}_{\min }(\mathrm{n}) \mathrm{dB}$
is computed for every subband n .


| Index Number $i$ | Frequency $=[\mathrm{Hz}]$ | $\begin{gathered} \text { Crit. Band } \\ \text { Rate } \\ {[z]} \\ \hline \end{gathered}$ | $\begin{gathered} \text { Absolute } \\ \text { Thresh. } \\ \text { [dB] } \\ \hline \end{gathered}$ | 71 72 73 | $\begin{array}{ll}5 & 875,00 \\ 6 & 000,00 \\ 6 & 250,00\end{array}$ | $\begin{aligned} & 19,482 \\ & 19,606 \\ & 19,847 \end{aligned}$ | $\begin{aligned} & 1,95 \\ & 2,08 \\ & 2,33 \end{aligned}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 62,50 | 0,617 | 33.44 | 74 | 6500,00 | 20,079 | 2,59 |
| 2 | 125,00 | 1,232 | 19.20 | 75 | 6 750,00 | 20,300 | 2,86 |
| 3 | 187.50 | 1,842 | 13.87 | 76 | 7 7 000,00 | 20,513 | 3,17 |
| 4 | 250,00 | 2,445 | 11,01 | 77 | 7 250,00 | 20,717 | 3,51 |
| 5 | 312,50 | 3,037 | 9,20 | 78 | 7500,00 | 20,912 | 3,89 |
| 6 | 375,00 | 3,618 | 7,94 | 79 | 7750.00 | 21,098 | 4,31 |
| 7 | 437,50 | 4,185 | 7,00 | 80 | 8 000,00 | 21,275 | 4.79 |
| 8 | 500,00 | 4,736 | 6,28 | 81 | 8 \% 250,00 | 21,445 | 5,31 |
| 9 | 562,50 | 5.272 | 5,70 | 82 | 8500,00 | 21,606 | 5,88 |
| 10 | 625,00 | 5.789 | 5,21 | 83 84 | 8 8 750,00 | 21,760 | 6,50 |
| 11 | 687,50 | 6,289 | 4.80 | 84 | 9 l 900,00 | 21,906 | 7,19 |
| 12 | 750,00 | 6.770 | 4,45 | 85 86 | 9 9 9 500,00 | 22,046 | 7.93 8.75 |
| 13 | 812,50 | 7,233 | 4,14 | 86 87 | 9500,00 9750,00 | 22,178 22,304 | 8,75 |
| 14 | 875,00 937,50 | 7,677 8,103 | 3,86 | 88 | 100000 | 22,304 22,424 | 9,63 10,58 |
| 16 | 100000 | 8,511 | 3,61 3,37 | 89 | 10 250,00 | 22,538 | 11,60 |
| 17 | 1062,50 | 8.901 | 3.15 | 90 | 10 500,00 | 22,646 | 12,71 |
| 18 | 1 125,00 | 9.275 | 2,93 | 91 | 10 750,00 | 22,749 | 13,90 |
| 19 | 1 187.50 | 9,632 | 2,73 | 92 | 11 000,00 | 22,847 | 15,18 |
| 20 | 1250,00 | 9,974 | 2,53 | 93 | 11 250,00 | 22,941 | 16,54 |
| 21 | 1312,50 | 10,301 | 2,32 | 94 | $\begin{array}{ll}11 & 500,00 \\ \text { 11 } & 750,00\end{array}$ | 23,030 | 18,01 |
| 22 | 1375,00 | 10,614 | 2,12 | 95 | 11 <br> 12 <br> 1750,00 | 23,114 | 19,57 |
| 23 | 1437,50 | 10,913 | 1,92 | 96 97 | $\begin{aligned} & 12 \\ & 12 \\ & 12\end{aligned} 2500,000$ | 23.195 23.272 | 21.23 23.01 |
| 24 | 1500,00 | 11,199 | 1,71 | 98 | 12 12 500,00 | 23,272 23,345 | 23,01 24,90 |
| 25 | 1562,50 | 11.474 | 1.49 1.27 | 99 | 12 750,00 | 23,415 | 24,90 26,90 |
| 26 27 | 16625,00 1687,50 | 11.736 11.988 | 1.27 1.04 | 100 | 13000,00 | 23,482 | 26,90 29,03 |
| 28 | 1 1750,00 | 11,988 12,230 | 1.04 0.80 | 101 | 13 250,00 | 23.546 | 31,28 |
| 29 | 1812,50 | 12,461 | 0,55 | 102 | 13 500,00 | 23,607 | 33,67 |
| 30 | 1875,00 | 12,684 | 0,29 | 103 | 13 750,00 | 23,666 | 36,19 |
| 31 | 1937,50 | 12,898 | 0,02 | 104 | $14,000,00$ | 23,722 | 38,86 |
| 32 | 2000,00 | 13,104 | -0,25 | 105 | $\begin{array}{ll}14 & 250,00 \\ 14 & 500,00\end{array}$ | 23,775 | 41,67 |
| 33 | 2062,50 | 13,302 | -0,54 | 106 | 14500,00 <br> 14 <br> 150,00 | 23.827 | 44,63 |
| 34 | 2 125,00 | 13,493 | -0,83 | 107 | $14 \quad 750,00$ 15000,00 | 23,876 23.923 | 47,76 51,04 |
| 35 | $\begin{array}{ll}2 & 187,50\end{array}$ | 13,678 | -1,12 | 108 | 15000,00 | 23.923 | 51,04 |
| 36 | 2 250,00 | 13.855 | -1,43 |  |  |  |  |
| 37 | 2312,50 | 14.027 | -1.73 |  |  |  |  |
| 38 | 2375,00 | 14,193 | -2,04 |  |  |  |  |
| 39 | 2 437,50 | 14,354 | -2,34 |  |  |  |  |
| 40 | 2500,00 | 14,509 | -2,64 |  |  |  |  |
| 41 | 2562,50 | 14,660 | -2,93 |  |  |  |  |
| 42 | 2 625,00 | 14,807 | -3,22 |  |  |  |  |
| 43 | 2 687,50 | 14,949 | -3,49 |  |  |  |  |
| 44 | 2750.00 | 15.087 | -3,74 |  |  |  |  |
| 45 | 2812,50 | 15,221 | -3,98 |  |  |  |  |
| 47 | 2 2 937,50 | 15,478 | -4,20 |  |  |  |  |
| 48 | 3 000,00 | 15,602 | -4,57 |  |  |  |  |
| 49 | 3 125,00 | 15,841 | -4,82 |  |  |  |  |
| 50 | 3 250,00 | 16,069 | -4,96 |  |  |  |  |
| 51 | 3 375,00 | 16,287 | -4,97 |  |  |  |  |
| 52 | 3500,00 | 16,496 | -4,86 |  |  |  |  |
| 53 | 3 625,00 | 16,697 | -4,63 |  |  |  |  |
| 54 | 3 750,00 | 16.891 | -4,29 |  |  |  |  |
| 55 | 3 875,00 | 17,078 | -3,87 |  |  |  |  |
| 56 | 4 4000,00 | 17,259 | -3,39 |  |  |  |  |
| 57 | 4 125,00 | 17,434 | -2,86 |  |  |  |  |
| 58 | 4 250,00 | 17,605 | -2,31 |  |  |  |  |
| 59 | 4 375,00 | 17.770 | -1,77 |  |  |  |  |
| 60 | 4500,00 | 17.932 | -1.24 |  |  |  |  |
| 61 | 4 625,00 | 18.089 | -0,74 |  |  |  |  |
| 62 | $4.750,00$ | 18,242 | -0,29 |  |  |  |  |
| 63 | 4.875,00 | 18,392 | 0,12 |  |  |  |  |
| 64 65 | 5000,00 $¢ \quad 125,00$ | 18,539 18,682 | 0.48 0.79 |  |  |  |  |
| 66 | $\Sigma$ 250,00 | 18,823 | 1,06 |  |  |  |  |
| 67 | 5375,00 | 18,960 | 1,29 |  |  |  |  |
| 68 | 5 500,00 | 19.095 | 1,49 |  |  |  |  |
| 69 | 5625.00 | 19.226 | 1.66 |  |  |  |  |
| 70 | 5750.00 | 19.356 | 1.81 |  |  |  |  |

Table D.1b. -- Frequencies, critical band rates and absolute threshold Table is valid for Layer $I$ at a sampling rate of $44,1 \mathrm{kHz}$.

| Index Number i | Frequency $[\mathrm{Hz}$ | $\begin{gathered} \hline \text { Crit.Band } \\ \text { Rate } \\ {[z]} \\ \hline \hline \end{gathered}$ | Absolute Thresh. [dB] | $\begin{aligned} & 71 \\ & 72 \\ & 73 \end{aligned}$ | $\begin{array}{ll}8 & 096,48 \\ 8 & 268,75 \\ 8 & 613,28\end{array}$ | $\begin{aligned} & 21,342 \\ & 21,457 \\ & 21,677 \end{aligned}$ | $\begin{aligned} & 4,98 \\ & 5,35 \\ & 6,15 \end{aligned}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 86,13 | 0,850 | 25.87 | 74 | 8 957,81 | 21,882 | 7,07 |
| 2 | 172,27 | 1,694 | 14.85 | 75 | 9 302,34 | 22,074 | 8,10 |
| 3 | 258,40 | 2,525 | 10.72 | 76 | 9646,88 | 22,253 | 9,25 |
| 4 | 344,53 | 3,337 | 8,50 | 77 | 9 991,41 | 22,420 | 10,54 |
| 5 | 430,66 | 4,124 | 7,10 | 78 | $\begin{array}{ll}10 & 335,94\end{array}$ | 22,576 | 11.97 |
| 6 | 516,80 | 4,882 | 6,11 | 79 | 10680,47 | 22,721 | 13,56 |
| 7 | 602,93 | 5,608 | 5,37 | 80 | $\begin{array}{ll}11 & 025,00\end{array}$ | 22,857 | 15,31 |
| 8 | 689,06 | 6,301 | 4,79 | 81 | 11369,53 | 22,984 | 17,23 |
| 9 | 775,20 | 6,959 | 4,32 | 82 | 117814,06 | 23,102 | 19,34 |
| 10 | 861,33 | 7.581 | 3,92 | 83 | 12 058,59 | 23,213 | 21,64 |
| 21 | 947.46 | 8,169 | 3,57 | 84 | $12.403,13$ | 23,317 | 24,15 |
| 12 | 1033.59 | 8,723 | 3,25 | 85 | $\begin{array}{ll}12 & 747,66 \\ 13 & 092,19\end{array}$ | 23,415 | 26,88 |
| 13 | 1119,73 | 9.244 | 2,95 | 86 | $\begin{array}{ll}13 & 092,19 \\ 13 & 436,72\end{array}$ | 23,506 | 29.84 |
| 14 | 1205,86 | 9.734 | 2,67 | 87 88 | $\begin{array}{ll}13 & 436,72 \\ 13 & 781,25\end{array}$ | 23,592 | 33.05 |
| 15 | 1291,99 | 10,195 | 2,39 | 88 | $\begin{array}{ll}13 & 781,25 \\ 14 & 125,78\end{array}$ | 23,673 23,749 | 36.52 |
| 16 | $\begin{aligned} & 1 \\ & 1 \\ & 1\end{aligned} 378,13$ | 10,629 | 2,11 | 99 | $\begin{array}{ll}14 & 125,78 \\ 14 & 470,31\end{array}$ | 23,749 23,821 | 40,25 44,27 |
| 18 | 1550,39 | 11,421 | 1,83 1,53 | 91 | $14.814,84$ | 23,888 | 48,59 |
| 19 | 1636,52 | 11,783 | 1,23 | 92 | 15 159,38 | 23,952 | 53,22 |
| 20 | 1 722,66 | 12,125 | 0.90 | 93 | 15 503,91 | 24,013 | 58,18 |
| 21 | 1808.79 | 12,448 | 0,56 | 94 | 15848,44 | 24,070 | 63.49 |
| 22 | 1894.92 | 12,753 | 0,21 | 95 | $\begin{array}{ll}16 & 192,97 \\ 16 & 537\end{array}$ | 24,125 | 68,00 |
| 23 | 1981,05 | 13,042 | -0,17 | 96 | 16537,50 | 24,176 | 68,00 |
| 24 | 2067,19 | 13,317 | -0,56 | 97 98 | 16 1682,03 | 24,225 | 68,00 |
| 25 | 2153,32 | 13,578 | -0,96 | 98 99 | 17 17 17 571,09 | 24.271 24.316 | 58,00 |
| 26 | $\begin{aligned} & 2 \\ & 2\end{aligned} 239,45$ | 13,826 | -1,38 | 109 | 17 <br> 17 <br> 1815,63 | 24.316 24.358 | 58,00 58,00 |
| 27 | $\begin{array}{ll}2 & 325.59 \\ 2 & 411.72\end{array}$ | 14,062 | -1.79 | 101 | 17 18 18 260,16 | 24,358 24,398 | 58,00 58,00 |
| 28 29 | $\begin{array}{ll}2 & 411,72 \\ 2 & 497,85\end{array}$ | 14,288 14.504 | $-2,21$ $-2,63$ | 102 | 18604,69 | 24,398 24,436 | 58,00 58,00 |
| 30 | $\begin{aligned} & 2 \\ & 2\end{aligned} 583,98$ | 14,514 | $-2,03$ $-3,03$ | 103 | 18949,22 | 24,473 | 58,00 |
| 31 | 2 670,12 | 14,909 | $-3,41$ | 104 | 19293,75 | 24,508 | 58,00 |
| 32 | 2 756,25 | 15,100 | -3,77 | $\begin{aligned} & 105 \\ & 106 \end{aligned}$ | $\begin{aligned} & 19 \\ & 19\end{aligned} 638,28182.81$ | 24,542 24.574 | $58,00$ $58.00$ |
| 33 | 2842,38 | 15,284 | -4,09 | 106 | 19 982.81 | 24,574 | 38,00 |
| 34 | 2928,52 | 15,460 | -4,37 |  |  |  |  |
| 35 | 3 014,65 | 15,631 | -4,60 |  |  |  |  |
| 36 | 3 100,78 | 15,796 | -4,78 |  |  |  |  |
| 37 | 3186,91 | 15,955 | -4,91 |  |  |  |  |
| 38 | 3 273,05 | 16,110 | -4,97 |  |  |  |  |
| 39 | 3 359,18 | 16.260 | -4,98 |  |  |  |  |
| 40 | $\begin{array}{ll}3 & 445,31\end{array}$ | 16,406 | -4,92 |  |  |  |  |
| 41 | $\begin{array}{ll}3 & 531,45 \\ 3 & 617\end{array}$ | 16,547 16.685 | -4,81 |  |  |  |  |
| 43 | 3 703,71 | 16,820 | -4,43 |  |  |  |  |
| 44 | 3789.84 | 16,951 | -4,17 |  |  |  |  |
| 45 | 3875,98 | 17,079 | -3,87 |  |  |  |  |
| 45 | 3 962,11 | 17,205 | -3,54 |  |  |  |  |
| 47 | 4 048,24 | 17.327 | -3.19 |  |  |  |  |
| 48 | 4134,38 | 17.447 | -2,82 |  |  |  |  |
| 49 | 4 306,64 | 17.680 | -2,06 |  |  |  |  |
| 50 | 4 478,91 | 17,905 | -1,32 |  |  |  |  |
| 51 | 4 651,17 | 18.121 | -0,64 |  |  |  |  |
| 52 | 4823,44 | 18,331 | -0,04 |  |  |  |  |
| 53 | 4 4 995,70 | 18,534 | 0,47 |  |  |  |  |
| 54 | $\begin{array}{lll}5 & 167,97\end{array}$ | 18,731 | 0,89 |  |  |  |  |
| 55 | $\begin{array}{ll}5 & 340,23\end{array}$ | 18.922 | 1,23 |  |  |  |  |
| 56 | 5 512,50 | 19,108 | 1.51 |  |  |  |  |
| 57 | ${ }_{5}^{5} 684,77$ | 19,289 | 1.74 |  |  |  |  |
| 58 59 | $\begin{array}{ll}5 & 857,03 \\ 6 & 029,30\end{array}$ | 19,464 19,635 | 1.93 2,11 |  |  |  |  |
| 60 | 6 6 201,56 | 19,635 19,801 | 2,28 |  |  |  |  |
| 61 | 6 373,83 | 19,963 | 2,46 |  |  |  |  |
| 62 | 6 546,09 | 20,120 | 2,63 |  |  |  |  |
| 63 | 6718,36 | 20,273 | 2,82 |  |  |  |  |
| 64 | 6890,63 | 20,421 | 3,03 |  |  |  |  |
| 65 | 7062,89 | 20,565 | 3,25 |  |  |  |  |
| 66 | 7 235,16 | 20,705 | 3.49 |  |  |  |  |
| 67 | 7407.42 | 20,840 | 3,74 |  |  |  |  |
| 68 | 7579.69 | 20,972 | 4,02 |  |  |  |  |
| 69 | 7751.95 | 21,099 | 4,32 |  |  |  |  |
| 70 | 7924,22 | 21,222 | 4,64 |  |  |  |  |

Table D.1c. -- Frequencies, critical band rates and absolute threshold Table is valid for Layer $I$ at a sampling rate of 48 kHz .

| Index Number $i$ | Frequency $[\mathrm{Hz}]$ | $\begin{gathered} \hline \text { Crit.Band } \\ \text { Rate } \\ \|z\| \\ \hline \end{gathered}$ | Absolute Thresh. [dB] | 71 72 73 | $\begin{array}{ll} 8 & 812,50 \\ 9 & 000,00 \\ 9 & 375,00 \end{array}$ | $\begin{aligned} & 21,797 \\ & 21,906 \\ & 22,113 \end{aligned}$ | $\begin{aligned} & 6,67 \\ & 7,19 \\ & 8,33 \end{aligned}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 93,75 | 0,925 | 24,17 | 74 | 9 750,00 | 22,304 | 9,63 |
| 2 | 187,50 | 1.842 | 13,87 | 75 | $10 \quad 125,00$ | 22,482 | 11,08 |
| 3 | 281,25 | 2,742 | 10,01 | 76 | 10 500,00 | 22,646 | 12,71 |
| 4 | 375,00 | 3,618 | 7.94 | 77 | 10875,00 | 22,799 | 14,53 |
| 5 | 468.75 | 4,463 | 6,62 | 78 | $\begin{array}{ll}11 & 250,00 \\ 11 & 625,00\end{array}$ | 22,941 | 16,54 |
| 6 | 562,50 | 5,272 | 5,70 | 79 | $11.625,00$ | 23,072 | 18,77 |
| 7 | 656,25 | 6,041 | 5,00 | 80 | 12 000,00 | 23,195 | 21,23 |
| 8 | 750,00 | 6,770 | 4,45 | 81 | 12375,00 | 23,309 | 23,94 |
| 9 | 843,75 | 7,457 | 4,00 | 82 | 12 750,00 | 23,415 | 26,90 |
| 10 | 937,50 | 8,103 | 3,61 | 83 | 13 125,00 | 23,515 | 30,14 |
| 11 | 1 031,25 | 8,708 | 3,26 | 84 | 13 500,00 | 23,607 | 33,67 |
| 12 | 1 125,00 | 9.275 | 2,93 | 85 | 13 875,00 | 23,694 | 37,51 |
| 13 | 1218,75 | 9.805 | 2,63 | 86 | 14 250,00 | 23.775 | 41.67 |
| 14 | 1312,50 | 10,301 | 2,32 2,32 | 87 | 14 625,00 | 23,852 | 46,17 |
| 15 | 1 406,25 | 10,765 | 2.02 | 88 | 15000,00 | 23,923 | 51.04 |
| 16 | 1500,00 | 11.199 | 1.71 | 89 | $\begin{array}{ll}15 & 375,00 \\ 15 & 750,00\end{array}$ | 23,991 | 56,29 |
| 17 | 1593,75 | 11,606 | 1,38 | 90 | 15 750,00 | 24,054 | 61,94 |
| 18 | 1 687,50 | 11,988 | 1,04 | 91 | 16 125,00 | 24,114 | 68,00 |
| 19 | 1 781,25 | 12,347 | 0,67 | 92 | 16 500,00 | 24.171 | 68,00 |
| 20 | 1875,00 | 12,684 | 0,29 | 93 | 16 875,00 | 24,224 | 68,00 |
| 21 | 1968.75 | 13,002 | -0,11 | 94 | 17 250,00 | 24.275 | 68,00 |
| 22 | 2062,50 | 13,302 | -0.54 | 95 | 17 <br> 18 <br> 18 <br> 025,000 | 24,322 24,368 | 68,00 68,00 |
| 23 | 2156,25 | 13,586 | -0.97 | 96 | 18 18 18 375,00 | 24,368 24,411 | 68,00 68,00 |
| 24 | 2 250,00 | 13,855 | -1.43 | 98 | 18 18750,00 | 24,451 | 68,00 |
| 25 | $\begin{array}{ll}2 & 343.75 \\ 2 & 437.50\end{array}$ | 14,111 | $-1.88$ | 98 | 18 19 1950,00 125,00 | 24,491 | 68,00 |
| 26 27 | $\begin{array}{ll}2 & 437,50 \\ 2 & 531,25\end{array}$ | 14,354 14,585 | $-2,34$ $-2,79$ | 100 | 19 500,00 | 24,528 | 68,00 |
| 28 | 2 625,00 | 14,807 | $-2,79$ $-3,22$ | 101 | 19875,00 | 24,564 | 68,00 |
| 29 | 2718,75 | 15,018 | -3,62 | 102 | 20 250,00 | 24,597 | 68,00 |
| 30 | 2812,50 | 15,221 | -3,98 |  |  |  |  |
| 31 | 2 906,25 | 15,415 | -4,30 |  |  |  |  |
| 32 | 3000,00 | 15,602 | -4,57 |  |  |  |  |
| 33 | 3093,75 | 15,783 | -4,77 |  |  |  |  |
| 34 | 3187,50 | 15,956 | -4,91 |  |  |  |  |
| 35 | 3 281,25 | 16,124 | -4,98 |  |  |  |  |
| 36 | 3 375,00 | 16,287 | -4,97 |  |  |  |  |
| 37 | 3 468,75 | 16,445 | -4,90 |  |  |  |  |
| 38 | 3 562,50 | 16,598 | -4,76 |  |  |  |  |
| 39 | 3 656,25 | 16,746 | -4,55 |  |  |  |  |
| 40 | 3750,00 | 16,891 | -4,29 |  |  |  |  |
| 41 | 3843.75 | 17.032 | -3,99 |  |  |  |  |
| 42 | 3937.50 | 17,169 | -3,64 |  |  |  |  |
| 43 | 4031.25 | 17,303 | -3,26 |  |  |  |  |
| 44 | 4 125,00 | 17,434 | -2,86 |  |  |  |  |
| 45 | 4218,75 | 17,563 | -2,45 |  |  |  |  |
| 46 | 4 312,50 | 17,688 | -2,04 |  |  |  |  |
| 47 | 4 406,25 | 17,811 | -1,63 |  |  |  |  |
| 48 | 4500,00 | 17,932 | -1,24 |  |  |  |  |
| 49 | 4 687,50 | 18,166 | -0,51 |  |  |  |  |
| 50 | 4875,00 | 18,392 | 0,12 |  |  |  |  |
| 51 | 5062,50 | 18,611 | 0,64 |  |  |  |  |
| 52 | 5 250,00 | 18,823 | 1.06 |  |  |  |  |
| 53 | 5 437,50 | 19.028 | 1.39 |  |  |  |  |
| 54 | 5 625,00 | 19,226 | 1,66 |  |  |  |  |
| 55 | 5812,50 | 19,419 | 1.88 |  |  |  |  |
| 56 | 6000,00 | 19.606 | 2,08 |  |  |  |  |
| 57 | 6187,50 | 19,788 | 2,27 |  |  |  |  |
| 58 | 6375.00 | 19,964 | 2,46 |  |  |  |  |
| 59 | 6 562,50 | 20,135 | 2.65 |  |  |  |  |
| 60 | 6750,00 | 20,300 | 2,86 |  |  |  |  |
| 61 | 6 937,50 | 20,461 | 3.09 |  |  |  |  |
| 62 | 7125,00 | 20.616 | 3,33 |  |  |  |  |
| 63 | $\begin{array}{ll}7 & 312,50 \\ 7 & 500,00\end{array}$ | 20,766 20,912 | 3,60 3,89 |  |  |  |  |
| 64 | 7800,00 $7 \quad 687,50$ | 20,912 21,052 | 3,89 4,20 |  |  |  |  |
| 66 | 7875,00 | 21,188 | 4,54 |  |  |  |  |
| 67 | 8062.50 | 21,318 | 4,91 |  |  |  |  |
| 68 | 8250,00 | 21,445 | 5,31 |  |  |  |  |
| 69 | 8437,50 | 21,567 | 5.73 |  |  |  |  |
| 70 | 8625,00 | 21,684 | 6,18 |  |  |  |  |

Table D.Id. .- Frequencies, critical band rates and absolute threshold Table is valid for Layer II at a sampling rate of 32 kHz .


Table D.Ie. -. Frequencies, Critical Band Rates and Absolute Threshold Table is valid for Layer II at a sampling rate of $44,1 \mathrm{kHz}$.

| Index Number $i$ | Frequency $[\mathrm{Hz}]$ | $\begin{gathered} \hline \text { Crit.Band } \\ \text { Rate } \\ {[z]} \\ \hline \hline \end{gathered}$ | Absolute Thresh. [dB] | 71 72 73 | $\begin{array}{ll}4 & 048,24 \\ 4 & 134,38 \\ 4 & 306,64\end{array}$ | $\begin{aligned} & 17,327 \\ & 17,447 \\ & 17,680 \end{aligned}$ | $\begin{aligned} & -3,19 \\ & -2,82 \\ & -2,06 \end{aligned}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2 | 43.07 | 0,425 | 45,05 | 74 | 4 478,91 | 17.905 | -1,32 |
| 2 | 86,13 | 0,850 | 25,87 | 75 | 4651,17 | 18,121 | -0,64 |
| 3 | 129.20 | 1,273 | 18,70 | 76 | 4823,44 | 18,331 | -0,04 |
| 4 | 172,27 | 1,694 | 14,85 | 77 | 4 995,70 | 18,534 | 0,47 |
| 5 | 215,33 | 2,112 | 12,41 | 78 | $5 \begin{array}{lll}5 & 167.97\end{array}$ | 18,731 | 0,89 |
| 6 | 258.40 | 2,525 | 10,72 | 79 | 5340,23 | 18,922 | 1.23 |
| 7 | 301,46 | 2,934 | 9.47 | 80 | 5 512,50 | 19.108 | 1,51 |
| 8 | 344,53 | 3,337 | 8,50 | 81 | 5684,77 | 19,289 | 1,74 |
| 9 | 387.60 | 3,733 | 7.73 | 82 | 5857,03 | 19.464 | 1,93 |
| 10 | 430,66 | 4,124 | 7,10 | 83 | 6029.30 | 19.635 | 2,11 |
| 11 | 473,73 | 4,507 | 6,56 | 84 | 6 201,56 | 19.801 | 2,28 |
| 12 | 516,80 | 4,882 | 6,11 | 85 | 6 373,83 | 19,963 | 2,46 |
| 13 | 559,86 | 5.249 | 5,72 | 86 | 6 546,09 | 20,120 | 2,63 |
| 14 | 602,93 | 5,608 | 5,37 | 87 | 6 6 718,36 | 20,273 | 2,82 |
| 15 | 646,00 | 5,959 | 5.07 | 88 | 6 890,63 | 20.421 | 3,03 |
| 16 | 689,06 | 6.301 | 4,79 | 89 | 7 062,89 | 20,565 | 3,25 |
| 17 | 732,13 | 6,634 | 4,55 | 90 | 7 235,16 | 20,705 | 3,49 |
| 18 | 775,20 | 6,959 | 4,32 | 91 | 7407.42 | 20,840 | 3,74 |
| 19 | 818.26 | 7,274 | 4,11 | 92 | 7 579,69 | 20,972 | 4,02 |
| 20 | 861,33 | 7,581 | 3,92 | 93 | 7 751,95 | 21,099 | 4,32 |
| 21 | 904,39 | 7,879 | 3,74 | 94 | 7924.22 | 21,222 | 4,64 |
| 22 | 947,46 | 8,169 | 3,57 | 95 | 8 096,48 | 21,342 | 4,98 |
| 23 | 990,53 | 8,450 | 3,40 | 96 | 8 268,75 | 21,457 | 5.35 |
| 24 | 1033,59 | 8,723 | 3,25 | 97 | 8 613,28 | 21,677 | 6,15 |
| 25 | 1 076,66 | 8,987 | 3,10 | 98 | 8957.81 | 21,882 | 7,07 |
| 26 | I 119,73 | 9,244 | 2,95 | 99 100 | $\begin{array}{ll}9 & 302,34 \\ 9 & 646,88\end{array}$ | 22,074 | 8.10 |
| 27 | l 162,79 | 9,493 | 2,81 | 100 | ${ }^{9} 646,88$ | 22,253 | 9.25 |
| 28 | 1205.86 | 9.734 | 2,67 | 101 | $9 \mathrm{991.41}$ | 22,420 | 10,54 |
| 29 | 1248.93 | 9,968 | 2,53 | 102 | 10335,94 | 22,576 | 11.97 |
| 30 | 1 291,99 | 10.195 | 2,39 | 103 | 10 680,47 | 22,721 | 13,56 |
| 31 | I 335,06 | 10,416 | 2,25 | 104 | $11.025,00$ | 22,857 | 15,31 |
| 32 | l 378,13 | 10,629 | 2,11 | 105 | $\begin{array}{ll}11 & 369,53 \\ 11 & 714,06\end{array}$ | 22,984 | 17.23 |
| 33 | 1 <br> 1 | 10,836 | 1,97 | 106 | $\begin{array}{ll}11 & 714,06 \\ 12 & 058\end{array}$ | 23,102 | 19,34 |
| 34 | 1 464,26 | 11.037 | 1.83 | 107 | 12 058,59 | 23,213 | 21.64 |
| 35 | 1 507,32 | 11,232 | 1,68 | 108 | 12 403,13 | 23,317 | 24,15 |
| 36 | 1550,39 | 11,421 | 1,53 | 109 | 12747.66 | 23,415 | 26,88 |
| 37 | 1593.46 | 11,605 | 1,38 | 110 | 13092,19 | 23,506 | 29,84 |
| 38 | 1 1 636,52 | 11,783 | 1,23 | 111 | $\begin{array}{ll}13 & 436.72 \\ 13 & 781.25\end{array}$ | 23,592 | 33,05 |
| 39 | 1679,59 | 11,957 | 1,07 | 112 | $\begin{array}{ll}13 & 781,25\end{array}$ | 23.673 | 36,52 |
| 40 | - 722,66 | 12,125 | 0,90 | 113 114 | $\begin{array}{lll}14 & 125,78 \\ 14 & 470,31\end{array}$ | 23,749 | 40,25 |
| 41 | I 765,72 | 12,289 | 0,74 | 114 | 14 <br> 14 <br> 14 <br> 150,31 | 23,821 23,888 | 44,27 |
| 42 | - 808,79 | 12,448 | 0,56 | 115 | $\begin{array}{ll}14 & 814,84 \\ 15 & 159\end{array}$ | 23.888 | 48,59 |
| 43 | 1851.86 | 12,603 | 0,39 | 116 | 15 159,38 | 23.952 | 53,22 |
| 44 | - 894,92 | 12,753 | 0,21 | 117 | 15 503,91 | 24,013 | 58,18 |
| 45 | 1937,99 | 12,900 | 0,02 | 118 | $\begin{array}{ll}15 & 848,44 \\ 16 & 192,97\end{array}$ | 24,070 | 63,49 |
| 46 | 1 981,05 | 13,042 | -0,17 | 119 | $\begin{array}{lll}16 & 192,97 \\ 16 & 537\end{array}$ | 24,125 | 68,00 |
| 47 | 2024,12 | 13,181 | -0,36 | 120 | 16 <br> 16 <br> 1637,50 <br> 882 | 24,176 | 68.00 |
| 48 | 2067,19 | 13,317 | -0,56 | 121 | 16882,03 | 24,225 | 68,00 |
| 49 | 2153,32 | 13,578 | -0,96 | 122 | 17 226,56 | 24,271 | 68.00 |
| 50 | 2 2 239,45 | 13,826 | -1,38 |  | 17 <br> 17 <br> 17 <br> 181509 | 24,316 | 68,00 |
| 51 | 2325,59 | 14,062 | -1,79 | 124 | $\begin{array}{ll}17 & 915,63 \\ 18 & 260,16\end{array}$ | 24,358 24,398 | 68,00 |
| 52 | 2 411,72 | 14,288 | $-2,21$ | 125 | 18 <br> 18 <br> 18 <br> 604,69 | 24,398 24,436 | 68,00 |
| 53 54 | $\begin{array}{ll}2 & 497.85 \\ 2 & 583.88\end{array}$ | 14,504 | $-2,63$ $-3,03$ | 126 | 18 <br> 18 <br> 18 <br> 649,69 | 24,436 24,473 | 68,00 68,00 |
| 54 55 | $\begin{array}{ll}2 & 583,98 \\ 2 & 670,12\end{array}$ | 14,711 14,909 | $-3,03$ $-3,41$ | 127 128 | $\begin{array}{ll}18 & 949,22 \\ 19 & 293,75\end{array}$ | 24,473 24,508 | 68,00 68,00 |
| 56 | 2 756,25 | 14,909 15,100 | $-3,41$ $-3,77$ | 129 | 19 638,28 | 24,542 | 68,00 |
| 57 | 2842,38 | 15,284 | -4,09 | 130 | 19 982,81 | 24,574 | 68,00 |
| 58 | 2 928,52 | 15,460 | -4,37 |  |  |  |  |
| 59 | 3014,65 | 15,631 | -4,60 |  |  |  |  |
| 60 | 3100.78 | 15,796 | -4,78 |  |  |  |  |
| 61 | 3186,91 | 15,955 | -4,91 |  |  |  |  |
| 62 | ${ }^{3}$ 273,05 | 16,110 | -4.97 |  |  |  |  |
| 63 | 3 359.18 | 16,260 | -4,98 |  |  |  |  |
| 64 | 3 445,31 | 16,406 | -4,92 |  |  |  |  |
| 65 | 3531,45 | 16,547 | -4,81 |  |  |  |  |
| 66 | 3617,58 | 16,685 | -4,65 |  |  |  |  |
| 67 | 3 703,71 | 16,820 | -4,43 |  |  |  |  |
| 68 | 3789.84 | 16,951 | -4,17 |  |  |  |  |
| 69 | 3 875,98 | 17,079 | -3,87 |  |  |  |  |
| 70 | 3962,11 | 17.205 | -3,54 |  |  |  |  |

Table D.1f. -- Frequencies, critical band rates and absolute threshold Table is valid for Layer II at a sampling rate of $48 \mathbf{~ k H z}$

| $\begin{gathered} \text { Index } \\ \text { Nuinber } \\ \text { i } \end{gathered}$ | Frequency $[\mathrm{Hz}]$ | $\begin{gathered} \hline \text { Crit.Band } \\ \text { Rate } \\ {[z]} \\ \hline \end{gathered}$ | Absolute Thresh. <br> [dB] | 71 72 73 | $\begin{array}{ll}4 & 406,25 \\ 4 & 500,00 \\ 4 & 687,50\end{array}$ | $\begin{aligned} & 17,811 \\ & 17,932 \\ & 18,166 \end{aligned}$ | $\begin{aligned} & -1,63 \\ & -1,24 \\ & -0,51 \end{aligned}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 46,88 | 0,463 | 42,10 | 74 | 4875.00 | 18,392 | 0,12 |
| 2 | 93,75 | 0,925 | 24,17 | 75 | 5 062,50 | 18,611 | 0,64 |
| 3 | 140,63 | 1,385 | 17,47 | 76 | 5250,00 | 18,823 | 1,06 |
| 4 | 187,50 | 1,842 | 13.87 | 77 | $\begin{aligned} & 5 \\ & 5 \\ & 5\end{aligned} 627,50$ | 19,028 | 1,39 |
| 5 | 234,38 | 2,295 | 11.60 | 78 | $\begin{aligned} & 5 \\ & 5 \\ & 5\end{aligned} 825,00$ | 19.226 | 1,66 |
| 6 | 281,25 | 2,742 | 10,01 | 79 | 5812,50 | 19,419 | 1,88 |
| 7 | 328,13 | 3,184 | 8,84 | 80 | 6000,00 | 19,606 | 2,08 |
| 8 | 375,00 | 3,618 | 7,94 | 81 | 6187,50 | 19,788 | 2,27 |
| 9 | 421,88 | 4,045 | 7.22 | 82 | 6 375,00 | 19,964 | 2,46 |
| 10 | 468.75 | 4,463 | 6,62 | 83 | 6 562,50 | 20,135 | 2,65 |
| 11 | 515,63 | 4,872 | 6.12 | 84 | 6750,00 | 20,300 | 2,86 |
| 12 | 562,50 | 5,272 | 5.70 | 85 | 6 937,50 | 20,461 | 3,09 |
| 13 | 609,38 | 5,661 | 5,33 | 86 87 | $\begin{array}{ll}7 & 125,00 \\ 7 & 31250\end{array}$ | 20,616 | 3,33 |
| 14 | 656,25 | 6,0411 | 5,00 | 87 88 | 7 <br> 7 <br> 7 | 20.766 | 3.60 |
| 15 | 703,13 | 6,411 | 4,71 | 88 | 7500,00 7 68750 | 20,912 | 3.89 |
| 16 | 750,00 | 6,770 | 4,45 | 89 90 | 7687,50 78875,00 | 21,052 21.188 | 4,20 |
| 17 .8 | 796,88 | 7.119 | 4,21 | 90 | 7875,00 8062,50 | 21,188 | 4,54 |
| 88 | 843.75 | 7.457 | 4,00 | 92 | 88 | 21,318 | 4,91 5,31 |
| :0 | 890,63 937,50 | 7.785 8.103 | 3,79 3,61 | 93 | 8437,50 | 21,567 | 5,73 |
| $\therefore 1$ | 984,38 | 8,410 | 3,43 | 94 | 8625,00 | 21,684 | 6,18 |
| $: 2$ | 1 031,25 | 8,708 | 3,26 | 95 | 8812,50 | 21,797 | 6,67 |
| $: 13$ | 1078,13 | 8,996 | 3,09 | 96 | 9 9000,00 | 21,906 | 7.19 |
| ! $!4$ | 1 125,00 | 9,275 | 2,93 | 97 | 9375,00 | 22,113 | 8,33 |
| :5 | 1 171,88 | 9,544 | 2,78 | 8 | 9 125, | 22,304 | 9.63 |
| ! $!6$ | 1218.75 | 9.805 | 2,63 | 100 | 10 500,00 | 22,646 | 11,08 |
| 8 | 1 265,63 | 10,057 | 2,47 | 101 | 10875,00 | 22,799 | 14.53 |
| 18 -9 | 1312,50 | 10,301 | 2,32 | 102 | 11 250,00 | 22,941 | 14,53 |
| 19 80 | 1 1 359,38 | 10.537 | 2,17 | 103 | 11 625,00 | 23,072 | 18,77 |
| $\bigcirc 1$ | $\begin{aligned} & 1 \\ & 1 \\ & 1\end{aligned} 453,13$ | 10,765 | 2,02 1,86 | 104 | 12 000,00 | 23,195 | 21,23 |
| $\pm 2$ | 1500,00 | 11,199 | 1,71 | 105 | 12 375,00 | 23,309 | 23,94 |
| $\bigcirc 3$ | 1546,88 | 11,406 | 1,55 | 106 | 12 750,00 | 23,415 | 26,90 |
| 2-4 | 1593,75 | 11,606 | 1,38 | 107 | 13 125,00 | 23,515 | 30,14 |
| 25 | 1 640,63 | 11,800 | 1,21 | 108 | 13500,00 | 23,607 | 33,67 |
| 56 | 1 1687,50 | 11,988 | 1,04 | 109 | 13875,00 | 23.694 | 37.51 |
| $\bigcirc 7$ | 1 1 734,38 | 12,170 | 0.86 | 110 | 14250,00 | 23.775 | +1,67 |
| 58 | 1781,25 | 12,347 | 0.67 | 112 | 15025,00 | 23,852 23.923 | 51,04 |
| 19 40 | 1828,13 | 12,518 | 0,49 | 113 | 15 375,00 | 23.991 | 56,29 |
| 4.0 | 1875,00 1 1 | 12,684 12,845 | 0,29 0.09 | 114 | 15750,00 | 24,054 | 51,94 |
| 4.2 | 1 1968,75 | 13,002 | -0,11 | 115 | 16 125,00 | 24,114 | 58,00 |
| 4.3 | 2015,63 | 13,154 | -0,32 | 116 | 16500,00 | 24,171 | 58,00 |
| 4.4 | 2062,50 | 13,302 | -0,54 | 117 | 16875,00 | 24,224 | 58,00 |
| 4.5 | 2109,38 | 13,446 | -0,75 | 118 | 17250,00 | 24,275 | 58,00 |
| 46 | $\begin{array}{ll}2 & 156,25 \\ 2 & 203,13\end{array}$ | 13,586 | -0,97 | 119 120 | $\begin{aligned} & 17 \\ & 18 \\ & 18\end{aligned} 0000,00$ | 24,322 24,368 | 58,00 58,00 |
| 4.7 4.8 | 2 203,13 | 13.723 | -1,20 | 121 | $18 \quad 375,00$ | 24,411 | 58,00 |
| 4.8 4.9 | $\begin{array}{ll}2 & 250,00 \\ 2 & 343,75\end{array}$ | 13,855 14,111 | $-1,43$ $-1,88$ | 122 | 18 750,00 | 24,452 | 58,00 |
| $\triangle 0$ | 2 437,50 | 14,354 | -2,34 | 123 | 19 125,00 | 24,491 | 58,00 |
| $\leq 1$ | 2531,25 | 14,585 | -2,79 | 124 | 19 500,00 | 24,528 | 58,00 |
| 4 4 4 | $\begin{aligned} & 2 \\ & 2\end{aligned} 625,00$ | 14,807 | $-3,22$ $-3,62$ | $\begin{aligned} & 125 \\ & 126 \\ & \hline \end{aligned}$ | $\begin{aligned} & 19875,00 \\ & 20 \quad 250,00 \\ & \hline \end{aligned}$ | $\begin{aligned} & 24,564 \\ & 24,597 \\ & \hline \end{aligned}$ | $\begin{aligned} & 58,00 \\ & 58.00 \end{aligned}$ |
| $\stackrel{3}{4}$ | 2 718,75 | 15,018 | -3,62 |  |  |  |  |
| $\stackrel{4}{4}$ | 2812,50 | 15,221 | -3,98 |  |  |  |  |
| 洔 | 2906,25 | 15,415 | -4,30 |  |  |  |  |
| $\pm 6$ | 3000,00 | 15,602 | -4,57 |  |  |  |  |
| $\vdots 7$ | 3093,75 | 15,783 | -4,77 |  |  |  |  |
| ¢8 | 3 3 3 187.50 | 15,956 | -4,91 |  |  |  |  |
| +19 | $\begin{aligned} & 3 \\ & 3\end{aligned} 3875,00$ | 16,124 16,287 | $-4,98$ $-4,97$ |  |  |  |  |
| til | 3 468,75 | 16.445 | -4,90 |  |  |  |  |
| 62 | 3 562,50 | 16,598 | -4,76 |  |  |  |  |
| 63 | 3 656,25 | 16,746 | -4,55 |  |  |  |  |
| 5.4 | 3 750,00 | 16.891 | -4,29 |  |  |  |  |
| 45 | 3843,75 | 17,032 | -3,99 |  |  |  |  |
| 156 | 3 937,50 | 17,169 | -3,64 |  |  |  |  |
| 157 | 4031.25 | 17,303 | -3,26 |  |  |  |  |
| 158 | $\begin{array}{ll}4 & 125,00 \\ 4 & 218,75\end{array}$ | 17.434 17.563 | $-2,86$ $-2,45$ |  |  |  |  |
| $\begin{array}{r}159 \\ \hline 70\end{array}$ | $\begin{array}{ll}4 & 218,75 \\ 4 & 312,50\end{array}$ | 17,563 17,688 | $-2,45$ $-2,04$ |  |  |  |  |

Table D.2a. -- Critical band boundaries
This table is valid for Layer I at a sampling rate of 32 kHz .
The frequencies represent the top end of each critical band.

| no | index of <br> Table F\&CB | frequency [Hz] | Bark[z] |
| :---: | :---: | :---: | :---: |
| 0 | 1 | 62,500 | 0,617 |
| 1 | 3 | 187,500 | 1,842 |
| 2 | 5 | 312,500 | 3,037 |
| 3 | 7 | 437,500 | 4,185 |
| 4 | 9 | 687,500 | 5,272 |
| 5 | 11 | 812,500 | 6,289 |
| 6 | 13 | 937,500 | 7,233 |
| 7 | 15 | 1125,000 | 8,103 |
| 8 | 18 | 1312,500 | 9,275 |
| 9 | 21 | 1500,000 | 10,301 |
| 10 | 24 | 2087,500 | 11,199 |
| 11 | 27 | 2312,000 | 11,988 |
| 12 | 32 | 2750,000 | 13,104 |
| 13 | 37 | 3250,000 | 14,027 |
| 14 | 44 | 3875,000 | 15,087 |
| 15 | 50 | 4625,000 | 16,069 |
| 16 | 55 | 6500,000 | 18,078 |
| 17 | 61 | 7700,000 | 19,095 |
| 18 | 68 | 9250,000 | 20,079 |
| 19 | 74 | 11500,000 | 21,098 |
| 20 | 89 | 15000,000 | 22,046 |
| 21 | 94 |  | 23,030 |
| 22 | 108 |  | 23,923 |

Table D.2b. .. Critical band boundaries This table is valid for Layer I at a sampling rate of $44,1 \mathrm{kHz}$.

The frequencies represent the top end of each critical band.

| no | index of <br> Table F\&CB | frequency [Hz] | Bark[z] |
| :---: | :---: | :---: | :---: |
| 0 | 1 | 86,133 | 0,850 |
| 1 | 2 | 172,266 | 1,694 |
| 2 | 3 | 258,398 | 2,525 |
| 3 | 5 | 430,664 | 4,124 |
| 4 | 6 | 516,797 | 4,882 |
| 5 | 8 | 689,063 | 6,301 |
| 6 | 9 | 775,195 | 6,959 |
| 7 | 11 | 947,461 | 8,169 |
| 8 | 13 | 119,727 | 9,244 |
| 9 | 15 | 1291,992 | 10,195 |
| 10 | 17 | 1724,258 | 11,037 |
| 11 | 20 | 1981,056 | 12,125 |
| 12 | 23 | 2325,586 | 13,042 |
| 13 | 27 | 2756,250 | 14,062 |
| 14 | 32 | 3786,914 | 15,100 |
| 15 | 37 | 45,955 |  |
| 16 | 45 | 5478,906 | 17,079 |
| 17 | 50 | 6370,234 | 17,904 |
| 18 | 55 | 7579,688 | 18,922 |
| 19 | 61 | 9302,344 | 19,963 |
| 20 | 68 | 11369,531 | 22,971 |
| 21 | 75 | 15503,906 | 22,984 |
| 22 | 81 | 19982,813 | 24,013 |
| 23 | 93 |  | 24,573 |

Table D.2c. -- Critical band boundaries
This table is valid for Layer I at a sampling rate of 48 kHz .
The frequencies represent the top end of each critical band.

| no | indexof <br> TableF\& CB | frequency[Hz] | Bark[z] |
| :---: | :---: | :---: | :---: |
| 0 | 1 | 93,750 | 0,925 |
| 1 | 2 | 187,500 | 1,842 |
| 2 | 3 | 281,250 | 2,742 |
| 3 | 4 | 375,000 | 3,618 |
| 4 | 5 | 468,750 | 4,463 |
| 5 | 6 | 562,500 | 5,272 |
| 6 | 7 | 656,250 | 6,041 |
| 7 | 9 | 843,750 | 7,457 |
| 8 | 10 | 937,500 | 8,103 |
| 9 | 12 | 1125,000 | 9,275 |
| 10 | 14 | 1312,500 | 10,301 |
| 11 | 16 | 1781,250 | 11,199 |
| 12 | 19 | 1968,750 | 12,347 |
| 13 | 21 | 2343,750 | 13,002 |
| 14 | 25 | 2718,750 | 14,111 |
| 15 | 29 | 3281,250 | 15,018 |
| 16 | 35 | 3843,750 | 17,124 |
| 17 | 41 | 4687,500 | 18,165 |
| 18 | 49 | 5437,500 | 19,028 |
| 19 | 53 | 6375,000 | 19,964 |
| 20 | 58 | 7687,500 | 21,052 |
| 21 | 65 | 9375,000 | 22,113 |
| 22 | 73 | 115375,000 | 23,072 |
| 23 | 79 | 20250,000 | 23,991 |
| 24 | 89 |  | 24,597 |
| 25 | 102 |  |  |

Table D.2d. .- Critical band boundaries
This table is valid for Layer II at a sampling rate of 32 kHz . The frequencies represent the top end of each critical band.

| no | indexof <br> TableF\&CB | frequency[Hz] | Bark[z] |
| :---: | :---: | :---: | :---: |
| 0 | 1 | 31,250 | 0,309 |
| 1 | 3 | 93,750 | 0,925 |
| 2 | 6 | 187,500 | 1,842 |
| 3 | 10 | 312,500 | 3,037 |
| 4 | 13 | 406,250 | 3,903 |
| 5 | 17 | 531,250 | 5,006 |
| 6 | 21 | 781,250 | 6,041 |
| 7 | 25 | 937,500 | 7,004 |
| 8 | 30 | 1093,750 | 8,103 |
| 9 | 35 | 1281,250 | 9,090 |
| 10 | 41 | 1468,750 | 10,139 |
| 11 | 47 | 2687,500 | 11,058 |
| 12 | 51 | 2312,000 | 11,988 |
| 13 | 56 | 2750,000 | 14,104 |
| 14 | 61 | 3250,000 | 15,087 |
| 15 | 68 | 3875,000 | 16,069 |
| 16 | 74 | 4625,000 | 17,078 |
| 17 | 79 | 5500,000 | 18,089 |
| 18 | 85 | 6500,000 | 19,095 |
| 19 | 92 | 7750,000 | 20,079 |
| 20 | 98 | 9250,000 | 21,098 |
| 21 | 103 | 11500,000 | 22,046 |
| 22 | 109 | 15000,000 | 23,030 |
| 23 | 18 |  | 23,923 |

Table D.2e. -- Critical band boundaries
This table is valid for Layer II at a sampling rate of $44,1 \mathrm{kHz}$.
The frequencies represent the top end of each critical band.

| not | indexof <br> TableF\&CB | frequency[Hz] | Bark[z] |
| :---: | :---: | :---: | :---: |
| 0 | 1 | 43,066 | 0,425 |
| 1 | 2 | 86,133 | 0,850 |
| 2 | 3 | 129,199 | 1,273 |
| 3 | 5 | 215,332 | 2,112 |
| 4 | 7 | 301,465 | 2,934 |
| 5 | 10 | 530,664 | 4,124 |
| 6 | 13 | 659,863 | 5,249 |
| 7 | 16 | 818,063 | 6,301 |
| 8 | 19 | 947,461 | 7,274 |
| 9 | 22 | 1119,727 | 8,169 |
| 10 | 26 | 1291,992 | 9,244 |
| 11 | 30 | 1507,324 | 10,195 |
| 12 | 35 | 1722,656 | 11,232 |
| 13 | 40 | 1981,055 | 12,125 |
| 14 | 46 | 2325,586 | 13,042 |
| 15 | 51 | 2756,250 | 14,062 |
| 16 | 56 | 3273,047 | 15,100 |
| 17 | 62 | 3875,977 | 16,11 |
| 18 | 69 | 4478,906 | 17,079 |
| 19 | 74 | 5340,234 | 17,904 |
| 20 | 79 | 6373,828 | 18,922 |
| 21 | 85 | 7579,688 | 19,963 |
| 22 | 92 | 9302,344 | 20,971 |
| 23 | 99 | 11369,531 | 22,074 |
| 24 | 105 | 15503,906 | 22,984 |
| 25 | 117 | 19982,813 | 24,013 |
| 26 | 130 |  | 24,573 |

Table D.2f. -. Critical band boundaries
This table is valid for Layer II at a sampling rate of 48 kHz . The frequencies represent the top end of each critical band.

| no | index of <br> Table F\&CB | frequency [Hz] | Bark[z] |
| :---: | :---: | :---: | :---: |
| 0 | 1 | 46,875 | 0,463 |
| 1 | 2 | 93,750 | 0,925 |
| 2 | 3 | 140,625 | 1,385 |
| 3 | 5 | 234,375 | 2,295 |
| 4 | 7 | 328,125 | 3,184 |
| 5 | 9 | 421,875 | 4,045 |
| 6 | 12 | 562,500 | 5,272 |
| 7 | 14 | 656,250 | 6,041 |
| 8 | 17 | 796,875 | 7,119 |
| 9 | 20 | 937,500 | 8,103 |
| 10 | 24 | 125,000 | 9,275 |
| 11 | 27 | 1265,625 | 10,057 |
| 12 | 32 | 1500,000 | 11,199 |
| 13 | 37 | 1734,375 | 12,170 |
| 14 | 42 | 1968,750 | 13,002 |
| 15 | 49 | 2343,750 | 14,111 |
| 16 | 53 | 2718,750 | 15,018 |
| 17 | 59 | 321,250 | 16,124 |
| 18 | 65 | 3843,750 | 17,032 |
| 19 | 73 | 4687,500 | 18,166 |
| 20 | 77 | 5437,500 | 19,028 |
| 21 | 82 | 6375,000 | 19,964 |
| 22 | 89 | 7687,500 | 21,052 |
| 23 | 97 | 11675,000 | 22,113 |
| 24 | 103 | 15375,0000 | 23,072 |
| 25 | 113 | 20250,000 | 23,991 |
| 26 | 126 |  | 24,597 |

## D. 2 Psychoacoustic model 2

## D.2.1 General

Psychoacoustic Model 2 is an independent psychoacoustic model that can be adjusted and adapted to any ISO/IEC 11172-3 layer. This annex presents the general Psychoacoustic Model 2, and provides sufficient information for implementation of Model 2 with Layers I and II. The Layer III psychoacoustic model is based on this implementation, with adaptations as described in the Layer III encoder.

The threshold generation process has three inputs. They are:
a) The shift length for the threshold calculation process, iblen, where 384 <iblen<640. This iblen must remain constant over any particular application of the threshold calculation process. If (as in Layer III), it is necessary to calculate thresholds for two different shift lengths, two processes, each running with a fixed shift length, will be necessary. In the case of iblen outside the range of 384 to 640 it may be necessary to calculate the psychoacoustic thresholds with a different window length as well as shift length. There are two ways to do this:

- Use a different length transform, and recalculate the startup coefficients for the model, or
- Use the same length transform, but a substantially shorter Hann window, appropriate to the data and problem at hand.

The choice of these is left to the implementation.
b) The newest iblen samples of the signal, with the samples delayed (either in the filter bank or psychoacoustic calculation) such that the window of the psychoacoustic calculation is centered in the time-window of application.
c) The sampling rate. There are sets of tables provided for the standard sampling rates. Sampling rate, like iblen, must necessarily remain constant over one implementation of the threshold calculation process.

There is one output from Psychoacoustic Model 2, a set of Signal-to-Masking Ratios, $S M R_{n}$, which are adapted to the layers as described below.

Before running the model initially, the array used to hold the preceding FFT source data window and the arrays used to hold $r$ and $f$ should be zeroed to provide a known starting point.

In Layer II, the psychoacoustic masking ratios must be calculated twice during each coder frame. The more stringent of each pair of ratios is used for bit allocation as shown in the software simulation model for Layers I and II with Psychoacoustic Model 2.

## D.2.2 Comments on notation

Throughout this threshold calculation process, three indices for data values are used. These are:
$\omega$ - indicates that the calculation is indexed by frequency in the FFT spectral line domain. An index of 1 corresponds to the DC term and an index of 513 corresponds to the spectral line at the Nyquist frequency.
$b$ - indicates that the calculation is indexed in the threshold calculation partition domain. In the case where the calculation includes a convolution or sum in the threshold calculation partition domain, $b b$ will be used as the summation variable. Partition numbering starts at 1 .
$n$ - indicates that the calculation is indexed in the coder bit (or codebook) allocation domain. An index of 1 corresponds to the lowest band in the subband filter bank.

## D.2.3 The "spreading function"

Several points in the following description refer to the "spreading function". It is calculated by the following method:

$$
\operatorname{tmpx}=1,05(j-i),
$$

Where $i$ is the Bark value of the signal being spread, $j$ is the Bark value of the band being spread into, and tmpx is a temporary variable.
$x=8$ minimum ( $\left.(t m p x-0,5)^{2}-2(t m p x-0,5), 0\right)$
Where $x$ is a temporary variable, and minimum (a,b) is a function returning the more negative of a or $b$.
$t m p y=15,811389+7,5(t m p x+0,474)-17,5\left(1,0+(t m p x+0,474)^{2}\right)^{0,5}$
where tmpy is another temporary variable.
if $(t m p y<-100)$ then $\{s p r d n g f(i, j)=0\}$ else $\left\{\operatorname{sprdngf}(i, j)=10 \frac{(\mathrm{x}+\mathrm{tmpy})}{10}\right.$ )

## D.2.4 Steps in threshold calculation

The following are the necessary steps for calculation of the $S M R_{n}$ used in the coder.
a) Reconstruct 1024 samples of the input signal.
iblen new samples are made available at every call to the threshold generator. The threshold generator must store 1024 -iblen samples, and concatenate those samples to accurately reconstru:t 1024 consecutive samples of the input signal, $s_{i}$, where $i$ represents the index, $1<=i<=1024$ of the current input stream.
b) Calculate the complex spectrum of the input signal.

First, $s_{i}$ is windowed by a 1024 point Hann window, i.e. $s w_{i}=s_{i} *\left(0,5-0,5 \cos \left(\frac{2 \pi(\mathrm{i}-0,5)}{1024}\right)\right)$. Note that in Layer III, a shorter window may be used when window switching is active, with appropriate centering of the window, per the Layer III encoder description.

Second, a standard forward FFT of $s w_{i}$ is calculated.
Third, the polar representation of the transform is calculated. $r_{\omega}$ and $f \omega$ represent the magnitude and phase components of the transformed $s w_{i}$, respectively.
c) Calculate a predicted $r$ and $f$.

A predicted magnitude, $\hat{r} \omega$, and phase, $\hat{f} \omega$ are calculated from the preceding two threshold calculation blocks' $r$ and $f$.

$$
\begin{aligned}
& \hat{r}_{\omega}=2,0 r_{\omega}(t-1)-r_{\omega}(t-2) \\
& \hat{f}_{\omega}=2,0 f \omega_{\omega}(t-1)-f \omega(t-2)
\end{aligned}
$$

where $t$ represents the current block number, $t-1$ indexes the previous block's data, and $t-2$ indexes the data from the threshold calculation block before that.
© Calculate the unpredictability measure $c_{\omega}$
$c_{\omega}$ the unpredictability measure, is:

$$
c_{\omega}=\frac{\left(\left(\mathrm{r}_{\omega} \cos \mathrm{f}_{\omega} \hat{\mathrm{r}}_{\omega} \cos \hat{f_{\omega}}\right)^{2}+\left(\mathrm{r}_{\omega} \operatorname{sinf} \operatorname{sif}^{\mathrm{f}} \omega \sin \hat{f}_{\omega}\right)^{2}\right)^{0,5}}{\mathbf{r}_{\omega}+\mathrm{abs}\left(\hat{\mathbf{r}}_{\omega}\right)}
$$

By sacrificing performance, this measure can be calculated on only a lower portion of the frequency lines. Calculations should be done from DC to at least 3 kHz and preferably to 7 kHz . An upper limit of less than $5,5 \mathrm{kHz}$ may considerably reduce performance from that obtained during the subjective testing of the audio algorithm. The $c_{\omega}$ values above this limit should be set to 0,3 . Best results will be obtained by calculating $c_{\omega}$ up to 20 kHz .

Convolve the partitioned energy and unpredictability with the spreading function.

$$
\begin{aligned}
& e c b_{b}=\sum_{b b=1}^{b m a x} e_{b b} * s_{p r d n g f\left(b v a l_{b b}, b v a l_{b}\right)}^{b t_{b}=} \sum_{b b=1}^{b \max } c_{b b} * s p r d n g f\left(b v a l_{b b}, b v a l_{b}\right)
\end{aligned}
$$

Because $c t_{b}$. is weighted by the signal energy, it must be renormalized to $c b_{b}$.

$$
c b_{b}=\frac{c t_{b}}{e c b_{b}}
$$

At the same time, due to the non-normalized nature of the spreading function, $e c b_{b}$ should be renormalized and the normalized energy en $b$, calculated.

$$
e n_{b}=e c b_{b} * r n o r m_{b}
$$

The normalization coefficient, rnormb. is:

$$
\text { rnorm }_{b}=\frac{1}{\sum_{b b=O}^{\sum_{b a x} s p r d n g f(b v a l b b, b v a l b)}}
$$

g) Convert $c b_{b}$ to $t b_{b}$, the tonality index.

$$
t b_{b}=-0,299-0,43 \log _{\mathrm{e}}\left(c b_{b}\right)
$$

Each $t b_{b}$ is limited to the range of $0<t b_{b}<1$.
h) Calculate the required SNR in each partition.
$N M T_{b}=5,5 \mathrm{~dB}$ for all $b . N M T_{b}$ is the value for noise masking tone (in dB ) for the partition. The required signal to noise ratio, $S N R_{b}$, is:

$$
S N R_{b}=\operatorname{maximum}\left(\operatorname{minval} l_{b}, t b_{b} * T M N_{b}+\left(1-t b_{b}\right) * N M T_{b}\right)
$$

Where maximum ( $a, b$ ) is a function returning the least negative of $\mathbf{a}$ or $\mathbf{b}$.
i) Calculate the power ratio.

The power ratio, $b c b$, is:

$$
b c_{b}=10 \frac{-S N R_{b}}{10}
$$

j) Calculation of actual energy threshold, $n b_{b}$.

$$
n b_{b}=e n_{b} b c_{b}
$$

k) Spread the threshold energy over FFT lines, yielding $n b_{\omega}$.

$$
n b_{\omega}=\frac{n b_{b}}{\omega_{h i g} h_{b}-\omega l o w_{b}+1}
$$

1) Include absolute thresholds, yielding the final energy threshold of audibility, thr $\omega$

$$
t h r_{\omega}=\max \left(n b_{\omega} a b s t h r_{\omega}\right)
$$

The dB values of absthr shown in tables D.4. "Absolute threshold tables" are relative to the lev:l that a sine wave of $\pm \frac{1}{2}$ lsb has in the FFT used for threshold calculation. The dB values must be converted into the energy domain after considering the FFT normalization actually used.
m) Pre-echo control

For Layer III, pre-echo control occurs at this point. The actual control is described as part of the Layer III encoder specification. This step is omitted for Layers I and II.
n) Calculate the signal-to-mask ratios, $S M R_{n}$.

Table I.5. "Layer I and II coder partition table" shows:

1. The index, $n$, of the coder partition.
2. The lower index $\omega l o w_{n}$, of the coder partition.
3. The upper index, $\omega h i g h_{n}$ of the coder partition.
4. The width index, widt $h_{n}$, where widt $h_{n}=1$ for a psychoacoustically narrow scalefactor band, and widt $h_{n}=0$ for a psychoacoustically wide scalefactor band. A psychoacoustically narrow scalefactor band is one whose width is less than approximately $\frac{1}{3}$ critical band.

The energy in the scalefactor band, epart $n$, is:

$$
\text { epart }_{n}=\sum_{\omega=\omega l o w_{n}}^{\omega h i g h_{n}}
$$

Then, if (width $h_{n}=1$ ), the noise level in the scalefactor band, $n$ part $n$ is calculated as:

$$
\text { npart }_{n}=\sum_{\omega=\omega l o w_{n}}^{\omega h i g h_{n}}{ }^{t h r_{\omega}}
$$

else,

$$
\text { npart }_{n}=\text { minimum }\left(t h r_{\text {wlow }_{n}}, \ldots, \text { thr }_{\omega h i g h_{n}}\right)^{*}\left(\omega_{n i g} h_{n}-\omega l o w_{n}+l\right)
$$

Where, in this case, minimum $(\mathrm{a}, \ldots, \mathrm{z})$ is a function returning the smallest positive argument of the arguments a...z.

The ratios to be sent to the coder, $S M R_{n}$, are calculated as:

$$
S M R_{n}=10 \log _{10}\left(\frac{\text { epart }_{n}}{\text { ppart }_{n}}\right)
$$

Table D.3a. .- Calculation partition table This table is valid at a sampling rate of 32 kHz .

| Index | $\omega 10 w$ | $\omega$ high | bval | minval | TMN |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 1 | 0,00 | 0,0 | 24,5 |
| 2 | 2 | 4 | 0,63 | 0,0 | 24,5 |
| 3 | 5 | 7 | 1.56 | 20,0 | 24,5 |
| 4 | 8 | 10 | 2,50 | 20,0 | 24,5 |
| 5 | 11 | 13 | 3,44 | 20,0 | 24,5 |
| 6 | 14 | 16 | 4,34 | 20,0 | 24,5 |
| 7 | 17 | 19 | 5,17 | 20,0 | 24,5 |
| 8 | 20 | 22 | 5,94 | 20,0 | 24,5 |
| 9 | 23 | 25 | 6,63 | 17,0 | 24,5 |
| 10 | 26 | 28 | 7,28 | 15,0 | 24,5 |
| 11 | 29 | 31 | 7,90 | 15,0 | 24,5 |
| 12 | 32 | 34 | 8,50 | 10,0 | 24,5 |
| 13 | 35 | 37 | 9,06 | 7,0 | 24,5 |
| 14 | 38 | 41 | 9,65 | 7,0 | 24,5 |
| 15 | 42 | 45 | 10,28 | 4,4 | 24,8 |
| 16 | 46 | 49 | 10,87 | 4,4 | 25,4 |
| 17 | 50 | 53 | 11,41 | 4,5 | 25,9 |
| 18 | 54 | 57 | 11,92 | 4,5 | 26,4 |
| 19 | 58 | 61 | 12,39 | 4,5 | 26,9 |
| 20 | 62 | 65 | 12,83 | 4,5 | 27,3 |
| 21 | 66 | 70 | 13,29 | 4,5 | 27,8 |
| 22 | 71 | 75 | 13,78 | 4.5 | 28,3 |
| 23 | 76 | 81 | 14,27 | 4,5 | 28,8 |
| 24 | 82 | 87 | 14,76 | 4,5 | 29.3 |
| 25 | 88 | 93 | 15,22 | 4,5 | 29,7 |
| 26 | 94 | 99 | 15,63 | 4,5 | 30,1 |
| 27 | 100 | 106 | 16,06 | 4,5 | 30,6 |
| 28 | 107 | 113 | 16,47 | 4,5 | 31,0 |
| 29 | 114 | 120 | 16,86 | 4,5 | 31,4 |
| 30 | 121 | 129 | 17,25 | 4,5 | 31,8 |
| 31 | 130 | 138 | 17,65 | 4,5 | 32,2 |
| 32 | 139 | 148 | 18,05 | 4,5 | 32,5 |
| 33 | 149 | 159 | 18,42 | 4,5 | 32,9 |
| 34 | 160 | 170 | 18,81 | 4,5 | 33,3 |
| 35 | 171 | 183 | 19,18 | 4,5 | 33,7 |
| 36 | 184 | 196 | 19.55 | 4,5 | 34,1 |
| 37 | 197 | 210 | 19,93 | 4,5 | 34,4 |
| 38 | 211 | 225 | 20,29 | 4,5 | 34,8 |
| 39 | 226 | 240 | 20,65 | 4,5 | 35,2 |
| 40 | 241 | 258 | 21,02 | 4,5 | 35,5 |
| 41 | 259 | 279 | 21,38 | 4,5 | 35,9 |
| 42 | 280 | 300 | 21,74 | 4,5 | 36,2 |
| 43 | 301 | 326 | 22,10 | 4,5 | 36,6 |
| 44 | 327 | 354 | 22,44 | 4,5 | 36,9 |
| 45 | 355 | 382 | 22,79 | 4,5 | 37,3 |
| 46 | 383 | 420 | 23,14 | 4,5 | 37,6 |
| 47 | 421 | 458 | 23,49 | 4,5 | 38.0 |
| 48 49 | 459 497 | 496 513 | 23,83 24,07 | 4,5 4,5 | 38,3 38,6 |

Table D.3b. -- Calculation partition table
This table is valid at a sampling rate of $44,1 \mathrm{kHz}$.

| Index | $\omega$ low | $\omega$ high | bval | minval | TMN |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 1 | 0,00 | 0,0 | 24,5 |
| 2 | 2 | 2 | 0,43 | 0,0 | 24.5 |
| 3 | 3 | 3 | 0,86 | 0,0 | 24,5 |
| 4 | 4 | 4 | 1,29 | 20,0 | 24,5 |
| 5 | 5 | 5 | 1,72 | 20,0 | 24,5 |
| 6 | 6 | 6 | 2.15 | 20.0 | 24,5 |
| 7 | 8 | 7 | 2,58 | 20.0 | 24,5 |
| 8 | 8 | 8 | 3.01 | 20,0 | 24,5 |
| 9 | 9 | 9 | 3,45 | 20,0 | 24,5 |
| 10 | 10 | 10 | 3.88 | 20,0 | 24,5 |
| 11 | 11 | 11 | 4.28 | 20,0 | 24.5 |
| 12 | 12 | 12 | 4,67 | 20,0 | 24,5 |
| 13 | 13 | 13 | 5,06 | 20,0 | 24,5 |
| 14 | 14 | 14 | 5,42 | 20.0 | 24.5 |
| 15 | 15 | 15 | 5.77 | 20,0 | 24.5 |
| 16 | 16 | 16 | 6,11 | 17,0 | 24,5 |
| 17 | 17 | 19 | 6,73 | 17.0 | 24.5 |
| 18 | 20 | 22 | 7.61 | 15.0 | 24,5 |
| 19 | 23 | 25 | 8,44 | 10.0 | 24,5 |
| 20 | 26 | 28 | 9.21 | 7.0 | 24,5 |
| 21 | 29 | 31 | 9,88 | 7,0 | 24,5 |
| 22 | 32 | 34 | 10,51 | 4,4 | 25.0 |
| 23 | 35 | 37 | 11,11 | 4,5 | 25,6 |
| 24 | 38 | 40 | 11.65 | 4.5 | 26,2 |
| 25 | 41 | 44 | 12,24 | 4,5 | 26,7 |
| 26 | 45 | 48 | 12,85 | 4,5 | 27,4 |
| 27 | 49 | 52 | 13,41 | 4.5 | 27,9 |
| 28 | 53 | 56 | 13,94 | 4,5 | 28,4 |
| 29 | 57 | 60 | 14,42 | 4.5 | 28,9 |
| 30 | 61 | 64 | 14,86 | 4.5 | 29.4 |
| 31 | 65 | 69 | 15,32 | 4.5 | 29.8 |
| 32 | 70 | 74 | 15.79 | 4.5 | 30,3 |
| 33 | 75 | 80 | 16,26 | 4.5 | 30,8 |
| 34 | 81 | 86 | 16,73 | 4.5 | 31,2 |
| 35 | 87 | 93 | 17,19 | 4.5 | 31.7 |
| 36 | 94 | 100 | 17.62 | 4.5 | 32,1 |
| 37 | 101 | 108 | 18,05 | 4,5 | 32,5 |
| 38 | 109 | 116 124 | 18,45 18.83 | 4.5 4.5 | 32,9 33,3 |
| 39 40 | 117 125 | 124 134 | 18,83 19.21 | 4.5 | 33,7 |
| 41 | 135 | 144 | 19.60 | 4.5 | 34,1 |
| 42 | 145 | 155 | 20.00 | 4.5 | 34,5 |
| 43 | 156 | 166 | 20,38 | 4,5 | 34,9 |
| 44 | 167 | 177 | 20.74 | 4.5 | 35,2 |
| 45 | 178 | 192 | 21,12 | 4,5 | 35,6 |
| 46 | 193 | 207 | 21,48 | 4.5 | 36,0 |
| 47 | 208 | 222 | 21.84 | 4.5 | 36.3 |
| 48 | 223 244 | 243 264 | 22,20 22,56 | 4,5 4,5 | 36,7 37,1 |
| 50 | 265 | 286 | 22.91 | 4,5 | - 37.4 |
| 51 | 287 | 314 | 23.26 | 4,5 | 37.8 |
| 52 | 315 | 342 | 23.60 | 4.5 | 38.1 |
| 53 | 343 | 371 | 23.95 | 4.5 | 38.4 |
| 54 | 372 | 401 | 24,30 | 4.5 | 38.8 |
| 55 | 402 | 431 | 24.65 | 4,5 | 39.1 |
| 56 | 432 | 469 | 25.00 | 4.5 | 39.5 |
| 57 | 470 | 513 | 25.33 | 3.5 | 39.8 |

Table D.3c. -- Calculation partition table
This table is valid at a sampling rate of 48 kHz .

| Index | $\omega 10 w$ | $\omega$ high | bval | minval | TMN |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 1 | 0,00 | 0,0 | 24,5 |
| 2 | 2 | 2 | 0,47 | 0,0 | 24.5 |
| 3 | 3 | 3 | 0,94 | 0,0 | 24,5 |
| 4 | 4 | 4 | 1,41 | 20,0 | 24.5 |
| 5 | 5 | 5 | 1,88 | 20,0 | 24.5 |
| 6 | 6 | 6 | 2,34 | 20,0 | 24,5 |
| 7 | 7 | 7 | 2,81 | 20,0 | 24,5 |
| 8 | 8 | 8 | 3,28 | 20,0 | 24,5 |
| 9 | 9 | 9 | 3,75 | 20,0 | 24,5 |
| 10 | 10 | 10 | 4,20 | 20,0 | 24,5 |
| 11 | 11 | 11 | 4,63 | 20,0 | 24.5 |
| 12 | 12 | 12 | 5,05 | 20,0 | 24,5 |
| 13 | 13 | 13 | 5,44 | 20,0 | 24,5 |
| 14 | 14 | 14 | 5.83 | 20.0 | 24,5 |
| 15 | 15 | 15 | 6,19 | 20.0 | 24,5 |
| 16 | 16 | 16 | 6,52 | 17.0 | 24,5 |
| 17 | 17 | 17 | 6.86 | 17,0 | 24,5 |
| 18 | 18 | 20 | 7,49 | 15,0 | 24,5 |
| 19 | 21 | 23 | 8,40 | 10,0 | 24,5 |
| 20 | 24 | 26 | 9.24 | 7.0 | 24.5 |
| 21 | 27 | 29 | 9.97 | 7,0 | 24,5 |
| 22 | 30 | 32 | 10,65 | 4,4 | 25,1 |
| 23 | 33 | 35 | 11,28 | 4,5 | 25,8 |
| 24 | 36 | 38 | 11,86 | 4,5 | 26,4 |
| 25 | 39 | 41 | 12,39 | 4,5 | 26,9 |
| 26 | 42 | 45 | 12.96 | 4,5 | 27,5 |
| 27 | 46 | 49 | 13.56 | 4,5 | 28,1 |
| 28 | 50 | 53 | 14,12 | 4,5 | 28,6 |
| 29 | 54 | 57 | 14.62 | 4.5 | 29.1 |
| 30 | 58 | 62 | 15,14 | 4.5 | 29.6 |
| 31 | 63 | 67 | 15.67 | 4,5 | 30,2 |
| 32 | 68 | 72 | 16,15 | 4,5 | 30,7 |
| 33 | 73 | 77 | 16,58 | 4,5 | 31.1 |
| 34 | 78 | 83 | 17,02 | 4.5 | 31,5 |
| 35 | 84 | 89 | 17.44 | 4,5 | 31,9 |
| 36 | 90 | 95 | 17.84 | 4,5 | 32,3 |
| 37 | 96 | 103 | 18,24 | 4,5 | 32,7 |
| 38 | 104 | 111 | 18,66 | 4,5 | 33,2 |
| 39 | 112 | 120 | 19,07 | 4,5 | 33,6 |
| 40 | 121 | 129 | 19.47 | 4,5 | 34,0 |
| 41 | 130 | 138 | 19.85 | 4,5 | 34,3 |
| 42 | 139 | 149 | 20.23 | 4.5 | 34,7 |
| 43 | 150 | 160 | 20,63 | 4.5 | 35,1 |
| 44 | 161 | 173 | 21.02 | 4,5 | 35,5 |
| 45 | 174 | 187 | 21,40 | 4,5 | 35,9 |
| 46 | 188 | 201 | 21.76 | 4,5 | 36,3 |
| 47 | 202 | 219 | 22,12 | 4,5 | 36,6 |
| 48 | 220 | 238 | 22,47 | 4,5 | 37,0 |
| 49 | 239 | 257 | 22,83 | 4,5 | 37,3 377 |
| 50 | 258 | 283 | 23,18 | 4,5 | 37,7 |
| 51 | 284 | 309 | 23,53 | 4.5 | 38,0 |
| 52 | 310 | 335 | 23,88 | 4,5 | 38,4 |
| 53 | 336 | 363 | 24,23 | 4,5 | 38.7 |
| 54 | 364 | 391 | 24,58 | 4,5 | 39.1 39.4 |
| 55 | 392 | 423 | 24,93 | 4,5 | 39,4 39,8 |
| 56 57 | 424 466 | 465 | 25,27 25,61 | 4.5 3.5 | 39,8 40,1 |
| 57 58 | 466 508 | 507 513 | 25,61 25,81 | 3,5 3,5 | 40,3 |

Table D.4a. -. Absolute threshold table This table is valid at a sampling rate of 32 kHz .

A value of 0 dB represents a level in the absolute threshold calculation of 96 dB below the energy of a sine wave of amplitude +-32 760 .

| index [line |  |  | absthr |
| :--- | :--- | :--- | :--- |
| lower | higher | [dB] |  |
| 1 | 1 | 58,23 |  |
| 2 | 2 | 33,44 |  |
| 3 | 3 | 24,17 |  |
| 4 | 4 | 19,20 |  |
| 5 | 5 | 16,05 |  |
| 6 | 6 | 13,87 |  |
| 7 | 7 | 12,26 |  |
| 8 | 8 | 11,01 |  |
| 9 | 9 | 10,01 |  |
| 10 | 10 | 9,20 |  |
| 11 | 11 | 8.52 |  |
| 12 | 12 | 7,94 |  |
| 13 | 13 | 7,44 |  |
| 14 | 14 | 7,00 |  |
| 15 | 15 | 6,62 |  |
| 16 | 16 | 6,28 |  |
| 17 | 17 | 5,97 |  |
| 18 | 18 | 5,70 |  |
| 19 | 19 | 5,44 |  |
| 20 | 20 | 5,21 |  |
| 21 | 21 | 5,00 |  |
| 22 | 22 | 4,80 |  |
| 23 | 23 | 4,62 |  |
| 24 | 24 | 4,45 |  |
| 25 | 25 | 4,29 |  |
| 26 | 26 | 4,14 |  |
| 27 | 27 | 4,00 |  |
| 28 | 28 | 3,86 |  |
| 29 | 29 | 3,73 |  |
| 30 | 30 | 3,61 |  |
| 31 | 31 | 3,49 |  |
| 32 | 32 | 3,37 |  |
| 33 | 33 | 3,26 |  |
| 34 | 34 | 3,15 |  |
| 35 | 35 | 3,04 |  |
| 36 | 36 | 2,93 |  |
| 37 | 37 | 2,83 |  |
| 38 | 38 | 2,73 |  |
| 39 | 39 | 2,63 |  |
| 40 | 40 | 2,53 |  |
| 41 | 41 | 2,42 |  |
| 42 | 42 | 2,32 |  |
| 43 | 43 | 2,22 |  |
| 44 | 44 | 2.12 |  |
| 45 | 45 | 2.02 |  |
| 46 | 46 | 1,92 |  |
| 47 | 47 | 1,81 |  |
|  |  |  |  |
|  |  |  |  |
|  |  |  |  |


| index [line] |  | absthr |
| :--- | :--- | :--- |
| lower | higher | [dB] |
| 48 | 48 | 1,71 |
| 49 | 50 | 1,49 |
| 51 | 52 | 1,27 |
| 53 | 54 | 1,04 |
| 55 | 56 | 0,80 |
| 57 | 57 | 0,55 |
| 59 | 60 | 0,29 |
| 61 | 62 | 0,02 |
| 63 | 64 | $-0,25$ |
| 65 | 66 | $-0,54$ |
| 67 | 68 | $-0,83$ |
| 69 | 70 | $-1,12$ |
| 71 | 72 | $-1,43$ |
| 73 | 74 | $-1,73$ |
| 75 | 76 | $-2,04$ |
| 77 | 78 | $-2,34$ |
| 79 | 80 | $-2,64$ |
| 81 | 82 | $-2,93$ |
| 83 | 84 | $-3,22$ |
| 85 | 86 | $-3,49$ |
| 87 | 88 | $-3,74$ |
| 89 | 90 | $-3,98$ |
| 91 | 92 | $-4,20$ |
| 93 | 94 | $-4,40$ |
| 95 | 96 | $-4,57$ |
| 97 | 100 | $-4,82$ |
| 101 | 104 | $-4,96$ |
| 105 | 108 | $-4,97$ |
| 109 | 112 | $-4,86$ |
| 113 | 116 | $-4,63$ |
| 117 | 120 | $-4,29$ |
| 121 | 124 | $-3,87$ |
| 125 | 128 | $-3,39$ |
| 129 | 132 | $-2,86$ |
| 133 | 136 | $-2,31$ |
| 137 | 140 | $-1,77$ |
| 141 | 144 | $-1,24$ |
| 145 | 148 | $-0,74$ |
| 149 | 152 | $-0,29$ |
| 153 | 156 | 0,12 |
| 157 | 160 | 0,48 |
| 161 | 164 | 0,79 |
| 165 | 168 | 1,06 |
| 169 | 172 | 1,29 |
| 173 | 176 | 1,49 |
| 177 | 180 | 1,66 |
| 181 | 184 | 1,81 |
|  |  |  |


| index line] |  | absthr <br> ldB] |
| :--- | :--- | :--- |
| lower | higher | [dB |
| 185 | 188 | 1,95 |
| 189 | 192 | 2,08 |
| 193 | 200 | 2,33 |
| 201 | 208 | 2,59 |
| 209 | 216 | 2,86 |
| 217 | 224 | 3,17 |
| 225 | 232 | 3,51 |
| 233 | 240 | 3,89 |
| 241 | 248 | 4,31 |
| 249 | 256 | 4,79 |
| 257 | 264 | 5,31 |
| 265 | 272 | 5,88 |
| 273 | 280 | 6,50 |
| 281 | 288 | 7,19 |
| 289 | 296 | 7,93 |
| 297 | 304 | 8,75 |
| 305 | 312 | 9,63 |
| 313 | 320 | 10,58 |
| 321 | 328 | 11,60 |
| 329 | 336 | 12,71 |
| 337 | 344 | 13,90 |
| 345 | 352 | 15,18 |
| 353 | 360 | 16,54 |
| 361 | 368 | 18,01 |
| 369 | 376 | 19,57 |
| 377 | 384 | 21,23 |
| 385 | 392 | 23,01 |
| 393 | 400 | 24,90 |
| 401 | 408 | 26,90 |
| 409 | 416 | 29,03 |
| 417 | 424 | 31,28 |
| 425 | 432 | 33,67 |
| 433 | 440 | 36,19 |
| 441 | 448 | 38,86 |
| 449 | 456 | 41,67 |
| 457 | 464 | 44,63 |
| 465 | 472 | 47,76 |
| 473 | 480 | 51,03 |
|  |  |  |
|  |  |  |

Table D.4b -- Absolute threshold table
This table is valid at a sampling rate of $44,1 \mathrm{kHz}$.
A value of 0 dB represents a level in the absolute threshold calculation of 96 dB below the energy of a sin: wave of amplitude +-32 760 .

| index [line] |  | absthr [dB] | index [line] |  | absthr [dB] | index [line] |  | at sthr [ 1 IB] |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| lower | higher |  | lower | higher |  | lower | higher |  |
| 1 | 1 | 45,05 | 48 | 48 | -0,56 | 185 | 188 | 4,98 |
| 2 | 2 | 25,87 | 49 | 50 | -0,96 | 189 | 192 | 5,35 |
| 3 | 3 | 18,70 | 51 | 52 | -1,37 | 193 | 200 | 6,15 |
| 4 | 4 | 14,85 | 53 | 54 | -1,79 | 201 | 208 | 7,07 |
| 5 | 5 | 12,41 | 55 | 56 | -2,21 | 209 | 216 | 8,10 |
| 6 | 6 | 10,72 | 57 | 58 | -2,63 | 217 | 224 | 9,25 |
| 7 | 7 | 9.47 | 59 | 60 | -3,03 | 225 | 232 | 10,54 |
| 8 | 8 | 8,50 | 61 | 62 | -3,41 | 233 | 240 | 11,97 |
| 9 | 9 | 7,73 | 63 | 64 | -3,77 | 241 | 248 | 13,56 |
| 10 | 10 | 7,10 | 65 | 66 | -4,09 | 249 | 256 | 15,3) |
| 11 | 11 | 6,56 | 67 | 68 | -4,37 | 257 | 264 | 17,23 |
| 12 | 12 | 6,11 | 69 | 70 | -4,60 | 265 | 272 | 19,33 |
| 13 | 13 | 5,72 | 71 | 72 | -4,78 | 273 | 280 | 21,64 |
| 14 | 14 | 5,37 | 73 | 74 | -4,91 | 281 | 288 | 24,15 |
| 15 | 15 | 5,07 | 75 | 76 | -4,97 | 289 | 296 | 26,83 |
| 16 | 16 | 4,79 | 77 | 78 | -4,98 | 297 | 304 | 29,8.4 |
| 17 | 17 | 4,55 | 79 | 80 | -4,92 | 305 | 312 | 33,07 |
| 18 | 18 | 4,32 | 81 | 82 | -4,81 | 313 | 320 | 36,51 |
| 19 | 19 | 4,11 | 83 | 84 | -4,65 | 321 | 328 | 40,21 |
| 20 | 20 | 3,92 | 85 | 86 | -4,43 | 329 | 336 | 44,2; |
| 21 | 21 | 3,74 | 87 | 88 | -4,17 | 337 | 344 | 48,5:3 |
| 22 | 22 | 3,57 | 89 | 90 | -3,87 | 345 | 352 | 53,21 |
| 23 | 23 | 3,40 | 91 | 92 | -3,54 | 353 | 360 | 58,17 |
| 24 | 24 | 3,25 | 93 | 94 | -3,19 | 361 | 368 | 63,4i3 |
| 25 | 25 | 3,10 | 95 | 96 | -2,82 | 369 | 376 | 69,13 |
| 26 | 26 | 2,95 | 97 | 100 | -2,06 | 377 | 384 | 69,13 |
| 27 | 27 | 2,81 | 101 | 104 | -1,33 | 385 | 392 | 69,13 |
| 28 | 28 | 2,67 | 105 | 108 | -0,64 | 393 | 400 | 69,13 |
| 29 | 29 | 2,53 | 109 | 112 | -0,04 | 401 | 408 | 69,13 |
| 30 | 30 | 2,39 | 113 | 116 | 0,47 | 409 | 416 | 69,1: |
| 31 | 31 | 2,25 | 117 | 120 | 0,89 | 417 | 424 | 69,13 |
| 32 | 32 | 2,11 | 121 | 124 | 1,23 | 425 | 432 | 69,1: |
| 33 | 33 | 1,97 | 125 | 128 | 1,51 | 433 | 440 | 69,1 |
| 34 | 34 | 1,83 | 129 | 132 | 1,74 | 441 | 448 | 69,13 |
| 35 | 35 | 1,68 | 133 | 136 | 1,93 | 449 | 456 | 69,1:1 |
| 36 | 36 | 1,53 | 137 | 140 | 2,11 | 457 | 464 | 69,13 |
| 37 | 37 | 1,38 | 141 | 144 | 2,28 |  |  |  |
| 38 | 38 | 1,23 | 145 | 148 | 2,45 |  |  |  |
| 39 | 39 | 1,07 | 149 | 152 | 2,63 |  |  |  |
| 40 | 40 | 0,90 | 153 | 156 | 2,82 |  |  |  |
| 41 | 41 | 0,74 | 157 | 160 | 3,03 |  |  |  |
| 42 | 42 | 0,56 | 161 | 164 | 3,25 |  |  |  |
| 43 | 43 | 0,39 | 165 | 168 | 3,49 |  |  |  |
| 44 | 44 | 0,21 | 169 | 172 | 3,74 |  |  |  |
| 45 | 45 | 0,02 | 173 | 176 | 4,02 |  |  |  |
| 46 | 46 | -0,17 | 177 | 180 | 4,32 |  |  |  |
| 47 | 47 | -0,36 | 181 | 184 | 4,64 |  |  |  |

Table D.4c -- Absolute threshold table This table is valid at a sampling rate of 48 kHz .

A value of 0 dB represents a level in the absolute threshold calculation of 96 dB below the energy of a sine wave of amplitude +-32 760 .

| index [line] |  | absthr |
| :--- | :--- | :--- |
| lower | higher | $[\mathrm{dB}]$ |
| 1 | 1 | 42,10 |
| 2 | 2 | 24,17 |
| 3 | 3 | 17,47 |
| 4 | 4 | 13,87 |
| 5 | 5 | 11,60 |
| 6 | 6 | 10,01 |
| 7 | 7 | 8,84 |
| 8 | 8 | 7,94 |
| 9 | 9 | 7,22 |
| 10 | 10 | 6,62 |
| 11 | 11 | 6,12 |
| 12 | 12 | 5,70 |
| 13 | 13 | 5,33 |
| 14 | 14 | 5,00 |
| 15 | 15 | 4,71 |
| 16 | 16 | 4,45 |
| 17 | 17 | 4,21 |
| 18 | 18 | 4,00 |
| 19 | 19 | 3,79 |
| 20 | 20 | 3,61 |
| 21 | 21 | 3,43 |
| 22 | 22 | 3,26 |
| 23 | 23 | 3,09 |
| 24 | 24 | 2,93 |
| 25 | 25 | 2,78 |
| 26 | 26 | 2,63 |
| 27 | 27 | 2,47 |
| 28 | 28 | 2,32 |
| 29 | 29 | 2,17 |
| 30 | 30 | 2,02 |
| 31 | 31 | 1,86 |
| 32 | 32 | 1,71 |
| 33 | 33 | 1,55 |
| 34 | 34 | 1,38 |
| 35 | 35 | 1,21 |
| 36 | 36 | 1,04 |
| 37 | 37 | 0,86 |
| 38 | 38 | 0,67 |
| 39 | 39 | 0,49 |
| 40 | 40 | 0,29 |
| 41 | 41 | 0,09 |
| 42 | 42 | $-0,11$ |
| 43 | 43 | $-0,32$ |
| 44 | 44 | $-0,54$ |
| 45 | 45 | $-0,75$ |
| 46 | 46 | $-0,97$ |
| 47 | 47 | $-1,20$ |
|  |  |  |
|  |  |  |
|  |  |  |


| index [line] |  | absthr |
| :--- | :--- | :--- |
| lower | higher | [dB] |
| 48 | 48 | $-1,43$ |
| 49 | 50 | $-1,88$ |
| 51 | 52 | $-2,34$ |
| 53 | 54 | $-2,79$ |
| 55 | 56 | $-3,22$ |
| 57 | 58 | $-3,62$ |
| 59 | 60 | $-3,98$ |
| 61 | 62 | $-4,30$ |
| 63 | 64 | $-4,57$ |
| 65 | 66 | $-4,77$ |
| 67 | 68 | $-4,91$ |
| 69 | 70 | $-4,98$ |
| 71 | 72 | $-4,97$ |
| 73 | 74 | $-4,90$ |
| 75 | 76 | $-4,76$ |
| 77 | 78 | $-4,55$ |
| 79 | 80 | $-4,29$ |
| 81 | 82 | $-3,99$ |
| 83 | 84 | $-3,64$ |
| 85 | 86 | $-3,26$ |
| 87 | 88 | $-2,86$ |
| 89 | 90 | $-2,45$ |
| 91 | 92 | $-1,04$ |
| 93 | 94 | $-1,63$ |
| 95 | 96 | $-1,24$ |
| 97 | 100 | $-0,51$ |
| 101 | 104 | 0,12 |
| 105 | 108 | 0,64 |
| 109 | 112 | 1,06 |
| 113 | 116 | 1,39 |
| 117 | 120 | 1,66 |
| 121 | 124 | 1,88 |
| 125 | 128 | 2,08 |
| 129 | 132 | 2,27 |
| 133 | 136 | 2,46 |
| 137 | 140 | 2,65 |
| 141 | 144 | 2,86 |
| 145 | 148 | 3,09 |
| 149 | 152 | 3,33 |
| 153 | 156 | 3,60 |
| 157 | 160 | 3,89 |
| 161 | 164 | 4,20 |
| 165 | 168 | 4,54 |
| 169 | 172 | 4,91 |
| 173 | 176 | 5,31 |
| 177 | 180 | 5,73 |
| 181 | 184 | 6,18 |
|  |  |  |


| index $[$ line] |  | absthr |
| :--- | :--- | :--- |
| lower | higher | [dB] |
| 185 | 188 | 6,67 |
| 189 | 192 | 7,19 |
| 193 | 200 | 8,33 |
| 201 | 208 | 9,63 |
| 209 | 216 | 11,08 |
| 217 | 224 | 12,71 |
| 225 | 232 | 14,53 |
| 233 | 240 | 16,54 |
| 241 | 248 | 18,77 |
| 249 | 256 | 21,23 |
| 257 | 264 | 23,94 |
| 265 | 272 | 26,90 |
| 273 | 280 | 30,14 |
| 281 | 288 | 33,67 |
| 289 | 296 | 37,51 |
| 297 | 304 | 41,67 |
| 305 | 312 | 46,17 |
| 313 | 320 | 51,04 |
| 321 | 328 | 56,29 |
| 329 | 332 | 61,94 |
| 333 | 340 | 68,00 |
| 341 | 348 | 68,00 |
| 349 | 356 | 68,00 |
| 357 | 364 | 68,00 |
| 365 | 372 | 68,00 |
| 373 | 380 | 68,00 |
| 381 | 388 | 68,00 |
| 389 | 396 | 68,00 |
| 397 | 404 | 68,00 |
| 405 | 412 | 68,00 |
| 413 | 420 | 68,00 |
| 421 | 428 | 68,00 |

Table D. 5 -- Layer I and Layer II coder partition table

| Index | $\omega \mathbf{l o w}_{\mathbf{n}}+\mathbf{1}$ <br> $\omega_{\text {high }}^{\mathbf{n}}$ | width $_{\mathbf{n}}$ |
| :---: | :---: | :---: |
| 0 | 1 | 0 |
| 1 | 17 | 0 |
| 2 | 33 | 0 |
| 3 | 49 | 0 |
| 4 | 65 | 0 |
| 5 | 81 | 0 |
| 6 | 97 | 0 |
| 7 | 113 | 0 |
| 8 | 129 | 0 |
| 9 | 145 | 0 |
| 10 | 161 | 0 |
| 11 | 177 | 0 |
| 12 | 193 | 1 |
| 13 | 209 | 1 |
| 14 | 225 | 1 |
| 15 | 241 | 1 |
| 16 | 257 | 1 |
| 17 | 273 | 1 |
| 18 | 289 | 1 |
| 19 | 305 | 1 |
| 20 | 321 | 1 |
| 21 | 337 | 1 |
| 22 | 353 | 1 |
| 23 | 369 | 1 |
| 24 | 385 | 1 |
| 25 | 401 | 1 |
| 26 | 417 | 1 |
| 27 | 433 | 1 |
| 28 | 449 | 1 |
| 29 | 465 | 1 |
| 30 | 481 | 1 |
| 32 | 513 |  |

## Annex E <br> (informative) <br> Bit sensitivity to errors

## E.1. General

This annex indicates the sensitivity of individual bits to random errors if application specific error protection is needed.This sensitivity is given for each bit by a value from 0 to 5 , indicating the amount of degradation resulting from one isolated error :

| 5 | catastrophic |
| :--- | :--- |
| 4 | very annoying |
| 3 | annoying |
| 2 | slightly annoying |
| 1 | audible |
| 0 | insensitive |

The values are not the results of precise measurements, rather they rely upon knowledge of the codec. They assume the error detection scheme is not in use.

Some fields in the bit stream do not have a fixed length. All bits in these fields are rated for error sensitivity, even if not in use.

For all layers, the header and error check information defined in 2.4.1.3 and 2.4.1.4 are considered to have the highest sensitivity.

## E.2. Layers I and II

| Parameters | \#bit | sensitivity |
| :--- | :--- | :--- |
| Bit allocation | all bits | 5 |
| Scalefactors select information | all bits | 5 |
| Scalefactors | 5 (msb) | 4 |
|  | 4 | 4 |
|  | 3 | 4 |
|  | 2 | 3 |
|  | 1 | 2 |
| Subband samples (*) | 0 (lsb) | 1 |
|  | $8-16$ (msb) | 3 |
|  | $5-7$ | 2 |
|  | 3,4 | 1 |
|  | (lsb) $0-2$ | 0 |

(*) according to the bit allocation

## E.3. Layer III

| Parameters | \#bit | sensitivity |
| :--- | :--- | :--- |
| scfsi | all bits | 5 |
| pari2_3_length | all bits | 4 |
| big_values | all bits | 3 |
| global_gain | all bits | 5 |
| scalefac_compress | all bits | 5 |
| window_switching_flag | 0 | 5 |
| block_type | all bits | 4 |
| mixed_block_flag | 0 | 4 |
| table_select | all bits | 5 |
| region0_count | all bits | 3 |
| region1_count | all bits | 3 |
| prellag | 0 | 2 |
| scalefac_scale | 0 | 2 |
| countltable_select | 0 | 3 |
| Subblock_gain | 2 (msb) | 4 |
|  | 1 | 3 |
|  | 0 (lsb) | 2 |
| scale_fac (**) | 3 (msb) | $3(2)$ |
|  | 2 | $3(2)$ |
|  | 1 | $2(1)$ |
| Huffmancodebits $0(* * *)$ | $0(1 \mathrm{mb})$ | $3-0$ |

(**) the scalefac length depends on scalefac_compress.
The bit sensitivity values refer to the scalefac_scale value 1 (if 0 the value is in parenthesis).
(***) If $n$ is the number of bits for Huffman coding in one block the bit sensivity decreases linearly from 3 to 0 as the bit number varies from 0 up to $n$, (from low to high frequency).

## Note:

Rearrangement of the Huffman coded values:
To get better implicit error robustness for the low frequency part of the spectrum the Huffman coded values can be transmitted not in their logical order, but in an interleaved fashion.

If max_hlen is the maximum length of a Huffman codeword over the tables which are used to code the particular block and $n$ is the number of bits used for Huffman coding of data in the block (not frame), then int(n/max_hlen) slots are filled with the first codewords, beginning from low frequencies. The remaining codewords are filled into the remaining place, again arranged from low to high frequencies.

After bit interleaving, the bit sensivity of bit $k+i^{*} \operatorname{int}\left(n / m a x \_h l e n\right)$ decreases linearly from 3 to 0 as $k$ varies from 0 up to int(n/max_hlen) -1 , where $i=0, \ldots$, max_hlen -1 , and $n$ is the number of bits for Huffman coding in one block.

This is the recommended practice for Layer III data for all channels where error robustness is important.

## Annex F <br> (informative) <br> Error concealment

An optional feature of the coded bit stream is the CRC word which provides some error detection facility to the decoder. The Hamming distance of this error detection code is $\mathrm{d}=4$, which allows for the detection of up to 3 single bit errors or for the detection of one error burst of up to 16 bit length. The amount and the position of the protected bits within one encoded audio frame generally depends on the layer, the mode, data rate, and sampling frequency.

This can be usec to control an error concealment strategy in order to avoid severe impairments of the reconstructed signal due to errors in the most sensitive information.

Some basic techniques can be used for concealment, for instance information substitution, or muting. A simple substitution technique consists, when an erroneous frame occurs, of replacing it by the previous one (if error free).

# Annex G <br> (informative) <br> Joint stereo coding 

## G.1. Intensity stereo coding Layer I, II

An optional joint stereo coding method used in Layers I and II is intensity stereo coding. Intensity stereo coding can be used to increase the audio quality and/or reduce the bitrate for stereophonic signals. The gain in bitrate is typically about 10 to $30 \mathrm{kbits} / \mathrm{s}$. It requires negligible additional decoder complexity. The increase of encoder complexity is small. The encoder and decoder delay is not affected.

Psychoacoustic results indicate that at high frequencies (above about 2 kHz ) the localization of the stereophonic image within a critical band is determined by the temporal envelope and not by the temporal fine structure of the audio signal.

The basic idea for intensity stereo coding is that for some subbands, instead of transmitting separate left an 1 right subband samples, only the sum-signal is transmitted, but with scalefactors for both the left and right channels, thus preserving the stereophonic image.

Flow diagrams of a stereo encoder and decoder, including intensity stereo mode, are shown in figure G. 1 "General stereo encoder flow-chart" and figure G. 2 "General stereo decoder flow-chart". First, an estimatio 1 is made of the required bitrate for both left and right channel. If the required bitrate exceeds the available bitrate, the required bitrate can be decreased by setting a number of subbands to intensity stereo mode. Depending on the bitrate needed, subbands

16 to 31 ,
12 to 31 ,
8 to 31, or
4 to 31
can be set to intensity stereo mode. For the quantization of such combined subbands, the higher of the bit allocations for left and right channel is used.

The left and right subband signals of the subbands in joint stereo mode are added. These new subband signals are scaled in the normal way, but the originally determined scalefactors of the left and right subband signals are transmitted according to the bitstream syntax. Quantization of common subband samples, coding of common samples, and coding of common bit allocation are performed in the same way as in independent coding.

## G.2. MS_Stereo and intensity stereo coding Layer III

In Layer III a combination of ms_stereo mode (sum/difference) and intensity stereo mode can be used.
a) MS__stereo switching

MS_stereo mode is switched on if in joint stereo mode condition

$$
\sum_{i=0}^{511}\left[\mathrm{rl}_{\mathrm{i}}^{2}-\mathrm{rr}_{\mathrm{i}}^{2}\right]<0.8 * \sum_{\mathrm{i}=0}^{511}\left[\mathrm{rl}_{\mathrm{i}}^{2}+\mathrm{rr}_{\mathrm{i}}^{2}\right]
$$

is true. The values $\mathrm{rl}_{\mathrm{i}}$ and $\mathrm{rr}_{\mathrm{i}}$ correspond to the energies of the FFT line spectrum of the left and right channel calculated within the psychoacoustic model.
b) MS_stereo processing

- MS matrix

In MS._stereo mode the values of the normalized middle/side channel $M_{i} / S_{i}$ are transmitted instead of the left/right channel values $\mathrm{L}_{\mathrm{i}} / \mathrm{R}_{\mathrm{i}}$ :

$$
M_{i}=\frac{R_{i}+L_{i}}{\sqrt{2}} \quad \text { and } \quad S_{i}=\frac{L_{i}-R_{i}}{\sqrt{2}}
$$

- Limitation of $\mathrm{S}_{\mathrm{i}}$ channel bandwith

All $\mathrm{S}_{\mathrm{i}}$ values above the highest scalefactor band are set to zero.

- Sparsing of $\mathrm{S}_{\mathrm{i}}$ channel

In every scalefactor band $s b$ all pairs of small values $\left(\mathrm{S}_{\mathrm{i}}, \mathrm{S}_{\mathrm{i}+1}\right)$ are set to zero:

$$
\begin{aligned}
& \text { if }\left(\mathrm{S}_{\mathrm{i}}{ }^{2}+\mathrm{S}_{\mathrm{i}+1^{2}}\right)<\mathrm{S}_{\mathrm{sb}}{ }^{*}\left(\mathrm{~L}_{\mathrm{i}}^{2}+\mathrm{L}_{\mathrm{i}+1^{2}}+\mathrm{R}_{\mathrm{i}}^{2}+\mathrm{R}_{\mathrm{i}+1^{2}}\right) \\
& \mathrm{S}_{\mathrm{i}}=0 ; \quad \mathrm{S}_{\mathrm{i}+1}=0 ;
\end{aligned}
$$

The following difference channel threshold coefficients apply to the scalefactor bands for block type != 2 ('long MDCT transforms'):

| sb | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| $s_{s b}$ | 0,0 | 0,0 | 0,0 | 0,0 | 0,0 | 0,10 | 0,10 | 0,10 | 0,10 | 0,10 |  |
|  |  |  |  |  |  |  | 11 | 12 | 13 | 14 | 15 |
| $s_{b}$ | 10 | 116 | 17 | 18 | 19 | 20 |  |  |  |  |  |
| $s_{s b}$ | 0,10 | 0,20 | 0,30 | 0,40 | 0,50 | 0,60 | 0,70 | 0,80 | 0,90 | 1,00 | 1,50 |

c) Intensity stereo processing

- Calculation of intensity stereo position

For each scalefactor band sb coded in intensity stereo the following steps are executed:

- is_pos ${ }_{s b}=\operatorname{NINT}\left(\frac{12}{\pi} * \arctan \left(\sqrt{\frac{\text { L_E Energy }_{s b}}{\text { R_Energy }} \mathbf{s b}}\right)\right)$
- $L_{i}=L_{i}+R_{i} \quad$ for all indices $i$ within the actual scalefactor band $s b$
- $R_{i}=0 \quad$ for all indices $i$ within the actual scalefactor band $s b$
- the intensity stereo position is_possb is transmitted instead of the scalefactor of the right channel ( 3 bits always, stereo positions $0 . .6,7=$ illegal stereo position)
where 1__Energysb/R_Energysb denote the signal energies of the left/right channel within the actual scalefactor band and $\mathrm{L}_{\mathrm{i}} / \mathrm{R}_{\mathrm{i}}$ are the transformed values.

Scalefactor bands of the right/difference channel containing only zeros after coding which do not belong to the intensity coded part should be transmitted with the scalefactor '7' to prevent intensity stereo decoding.


Figure G.1 -- General stereo encoder flow chart


Figure G. 2 -- General stereo decoder flow chart

## Annex H

(informative)

## List of patent holders

The user's attention is called to the possibility that - for some of the processes specified in this part of ISO/IEC 11172 - compliance with this International Standard may require use of an invention covered by patent rights.

By publication of this part of ISO/IEC 11172, no position is taken with respect to the validity of this claim or of any patent rights in connection therewith. However, each company listed in this annex has filed with the Information Technology Task Force (ITTF) a statement of willingness to grant a license under such rights that they hold on reasonable and nondiscriminatory terms and conditions to applicants desiring to obtain such a license.

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