

# electronics

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CHECKING SHOES FOR MISPLACED TACKS—Inspection of shoe insoles for protruding metal is speeded at E. W. Wright & Co. factory by using detector developed by United Shoe Machinery engineers (see p 144) . COVER

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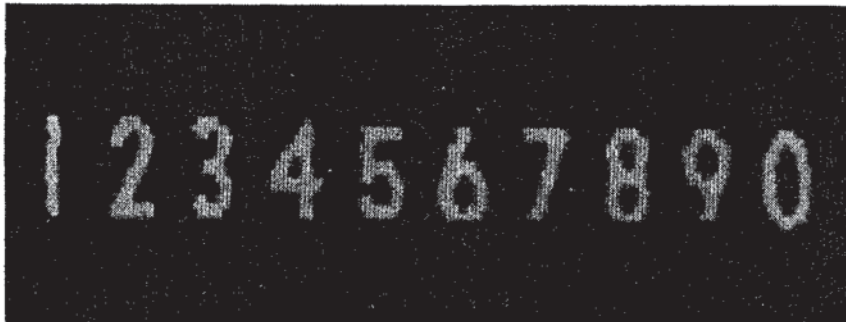
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# Character Recognition

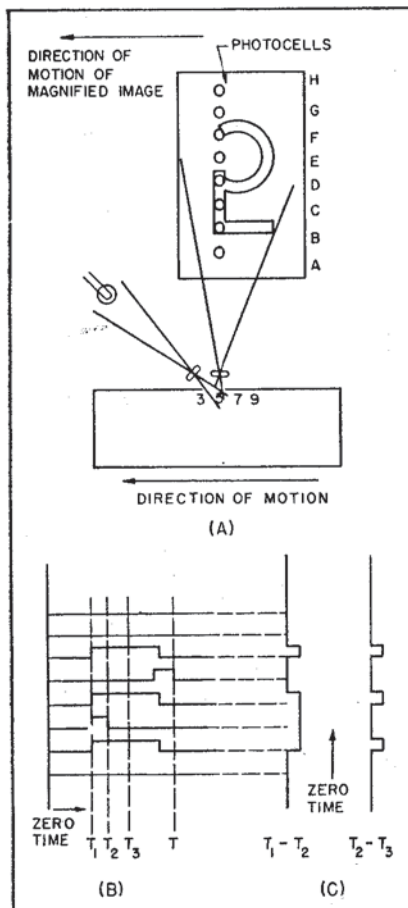
**SUMMARY** — Photoelectric scanner analyzes printed numerals and provides electrical output usable in computers and other business machines. Reader recognizes 400 characters per second. Operation is independent of type style or size of number above minimum width



Numerals printed by character reader

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**T**HIS READER recognizes Arabic numerals as printed matter passes through it.

Paper is not restricted as to color, thickness, opacity, roughness or quality. Ink need have no special qualities. The depth of the impression on the paper is of no significance, and at no time does the reader come into contact with the paper.

The character reader cannot read all type styles, but is not dependent on any one style. Minimum type width is approximately  $\frac{1}{8}$  in. for the widest digit. There is no restriction on maximum type width. The reader is not dependent on the height of the type providing the size is not larger than the aperture of the reading station. A common type height is about  $\frac{1}{8}$  inch. Most print imperfections will not impair recognition.

With checks moving into the reader at 16 per second, it recognizes 400 characters per second. The reader is able to read 1,600 characters per second, however, and could be made to operate at twice

serially and sorts by successive passes.

## Photoelectric Scanner

The scanner consists of a column of photocells whose outputs are modulated by the black portions of characters. The photocells are sequentially gated into a common buffer. Figure 1A shows how a magnified image of a number is projected on the photocells. Figure 2 shows the scanner circuits.

Figure 1B shows the photocell outputs. The effect of sequentially gating the outputs into a common buffer is shown in Fig. 1C.

Table I—Pulse-Code Combination

Total Pulses Per Scan	Long-Black Pulses Per Scan	Coded Combination
1	0	10
1	1	11
2	0	20
2	1	21\21'
2	2	22
3	0	30\30'
3	1	31

Note that the first digit of the coded combination represents the total number of black pulses per scan and the second digit represents

# for Business Machines

There are photocells above and below those required to cover the number to accommodate changes in the vertical registry. See Fig. 3. Each photocell is an input to a gate whose other input is connected to a tap of a multiple-tapped delay line. The outputs of the gates are combined in the common buffer *B*.

Each photocell is individually connected to network *B* by sending a pulse down the multiple-tapped delay line. A pulse sent down the delay line corresponds to a scan. No attempt is made to synchronize the start of a scan with the entry of a character under the column of photocells.

The uncertainty of the start of the scan with respect to the edge of the character to be scanned is called horizontal registry. The character reader overcomes this registry problem by utilizing the first scan not to recognize the character but to tell the recognition circuit to look for recognition on the second scan.

## Pulse Generator

The waveforms at the output of combining network *B* appear as in Fig. 1C. The total pulse generator produces one pulse for each black region of the number and serves as a noise filter in that an input pulse must exceed a certain width for an output pulse to be generated.

The long-black pulse generator (see Fig. 4) is also a pulse-width detector. In this case a pulse must exceed a predetermined length for a long-black pulse to be generated. For example, a pulse must equal or exceed the length shown in Fig. 1C.

The output of the total pulse generator and of the long-black pulse generator, along with the system trigger, are next sent to the encoder unit. Three pulses, each appearing once during every scan period, are derived from the tapped delay line and also sent to the encoder. These pulses, in the order of their time sequence, are the comparison pulse, the read-out pulse and the reset pulse.

Six instructions constitute the

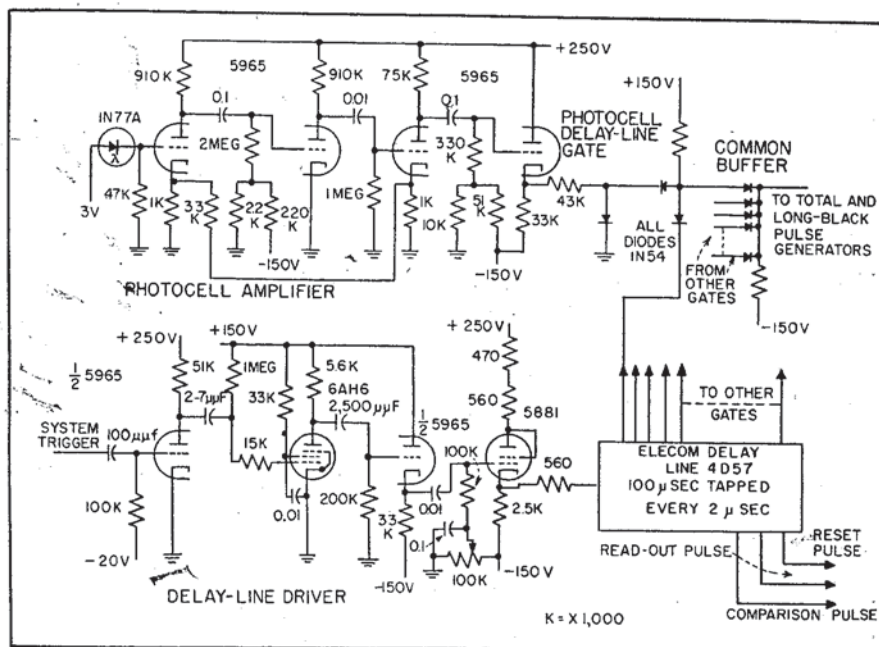


FIG. 2—Input section showing photodiode and delay line

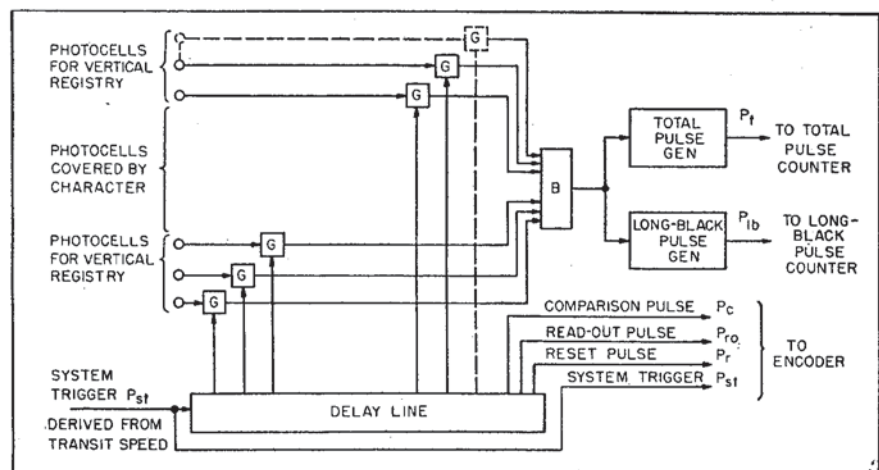


FIG. 3—Scanning station of character reader

program of the character reader.

(1) Count the total number of black pulses per scan.

(2) Count the number of long-black pulses per scan.

(3) Combine the results into discrete combinations as shown in Table I.

(4) If adjacent scans are not identical, put the most recent scan into shift-register storage and shift the register.

(5) If adjacent scans are identical, put nothing into storage and do not shift the register.

(6) When there is no character under the reading station advance

the register but put nothing into storage.

A block diagram of the encoding unit shown in Fig. 5. The system trigger  $P_{st}$  operates an electronic switch which routes the total pulses  $P_t$  alternately into total-pulse counter *A* and total-pulse counter *B* through gates  $G_a$ . The electronic switch also routes the long-black pulses  $P_{ib}$  into long-black counter *A* and long-black counter *B*, alternately, through gates  $G_b$ .

## Scan Comparison

Alternately switching input information  $P_t$  and  $P_{ib}$  into pairs of

Table II—Shift Register Storage (Numeral Three)

1 Scan	2 Total	3 Long	4 Code	5 Identity	6 Read-in	7 Shift	8 Interrogate	9 1st read-in in core row	10 2nd read-in in core row	11 3rd read-in in core row	12 4th read-in in core row
1	1	0	10	no	yes	yes	no	1	—	—	—
2	2	0	20	no	yes	yes	no	2	1	—	—
3	2	0	20	yes	no	yes	no	2	1	—	—
4	2	0	20	yes	no	no	no	2	1	—	—
5	2	0	20	yes	no	no	no	2	1	—	—
6	3	0	30 <sup>1</sup>	no	yes	no	no	2	1	—	—
7	3	0	30 <sup>1</sup>	yes	yes	yes	no	3	2	1	—
8	1	1	11	no	no	no	no	3	2	2	—
9	0	0	—	no	yes	yes	no	4	3	1	—
10	0	0	—	no	—	yes	no	5	4	2	1
11	0	0	—	no	—	yes	—	6	5	3	2
							next shift pulse	7	6	4	3
12	0	0	—	no	—	yes	yes	x	7	6	4
											5

identical counters provides for comparing adjacent scans for identity. The counters hold the input information until they are reset by reset pulse  $P_r$ , which is electronically switched through gates  $G_s$  once each scan alternately to reset both  $B$  counters and both  $A$  counters.

The comparison is accomplished in gates  $G_s$  once each scan after the counters have received their input information and at a time determined by the comparison pulse  $P_c$ . One of the six comparison gates will respond only when the input data  $P_i$  and  $P_{i+1}$  have remained constant from one scan to the next.

Scan-to-scan identity in the long-black counter allows the comparison pulse to appear at the output of buffer  $B_{di}$ . Scan-to-scan identity in the total-pulse counter allows the comparison pulse to appear in the

output of buffer stage  $B_{di}$ . A pulse appears at the output of both of the buffer stages only when scan-to-scan identity is indicated by both the total-pulse and the long-black counters. This is the necessary condition to pass the signal through gate  $G_s$  to the identity-gate generator.

After the  $A$  and  $B$  counters have been compared for identity they are read out by read-out pulse  $P_{ro}$ . If there has been no identity the read-out pulse appears at the output of gate  $G_n$ . The electronic switch then allows the read-out pulse to appear alternately at the input of the  $G_s$  gates. It enables the outputs of either the  $A$  or the  $B$  counters to appear at the outputs of the appropriate  $G_i$  gates.

The  $B_i$  buffers then show the accumulated count per scan inde-

pendently of whether the counts appear in the  $A$  counters or the  $B$  counters. The first scan of the character is detected in buffer  $B_s$  and used to generate an inhibiting signal whose duration is  $1\frac{1}{2}$  scan periods. This inhibiting signal, applied to gate  $G_s$ , prevents the recognition of an identity on the first two scans of the character. Subsequent scans of the character are prevented from generating additional inhibiting gates by the  $1\frac{1}{2}$ -character-period gate generator applied to gate  $G_i$  to disable it for the duration of the number.

### Coding

The coded combination (Table I) of total-pulse and long-black counters is formed by the  $G_m$  gates. Only one or none of these five combinations can appear during one scan. The prime combinations are obtained in the  $B_m$  buffers.

The instruction to put into storage is accomplished implicitly whenever a signal appears on one of the five code-combination lines. The storage register is shifted or advanced by the system trigger  $P_s$ , when the trigger is enabled to appear at the output of gate  $G_n$ . System trigger  $P_s$  appears and the register is enabled to advance when a permissive signal appears on either one of the two inputs to the  $B_n$  buffer. One input is connected to the identity-gate generator and allows the register to shift unless there is an identity.

Between characters, the outputs of both the total-pulse counter and the long-black counter are zero. This combination of 00 is not sensed for identity in the  $G_s$  gates; therefore no identity signal can be generated. Also, since the code combination 00 has not been

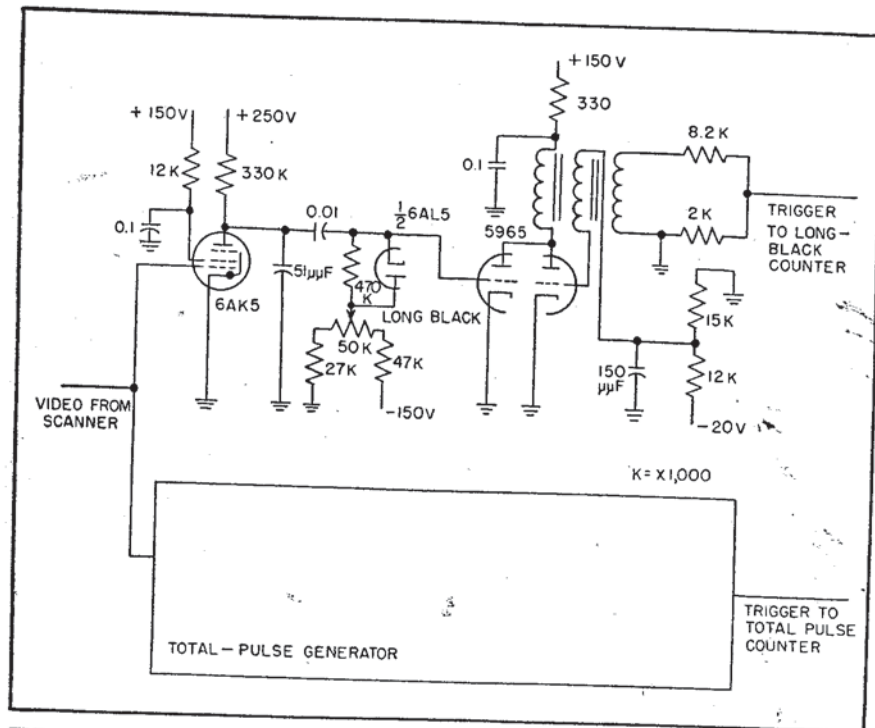


FIG. 4—Long-black pulse generator is pulse-width detector

formed in the  $G_m$  gates, there is nothing to put into storage.

Readiness for interrogation is detected in buffer  $B$ , which is connected to the seventh and last stage of storage. The first signal at the output of  $B$ , passes through gate  $G$ , and generates a gate whose duration is  $1\frac{1}{2}$  scan periods. This gate connected to  $B_n$  enables the register to advance on the next system trigger.

Subsequent signals at the output of  $B$ , during the character storage time are prevented from passing through gate  $G$ , by the  $1\frac{1}{2}$ -character-period gate generator, the latter

being energized by the system trigger that occurs after the first output signal at  $B_n$ . This particular system trigger is broadened in the matrix interrogation-pulse generator and is then used in the diode matrix.

### Shift Register Storage

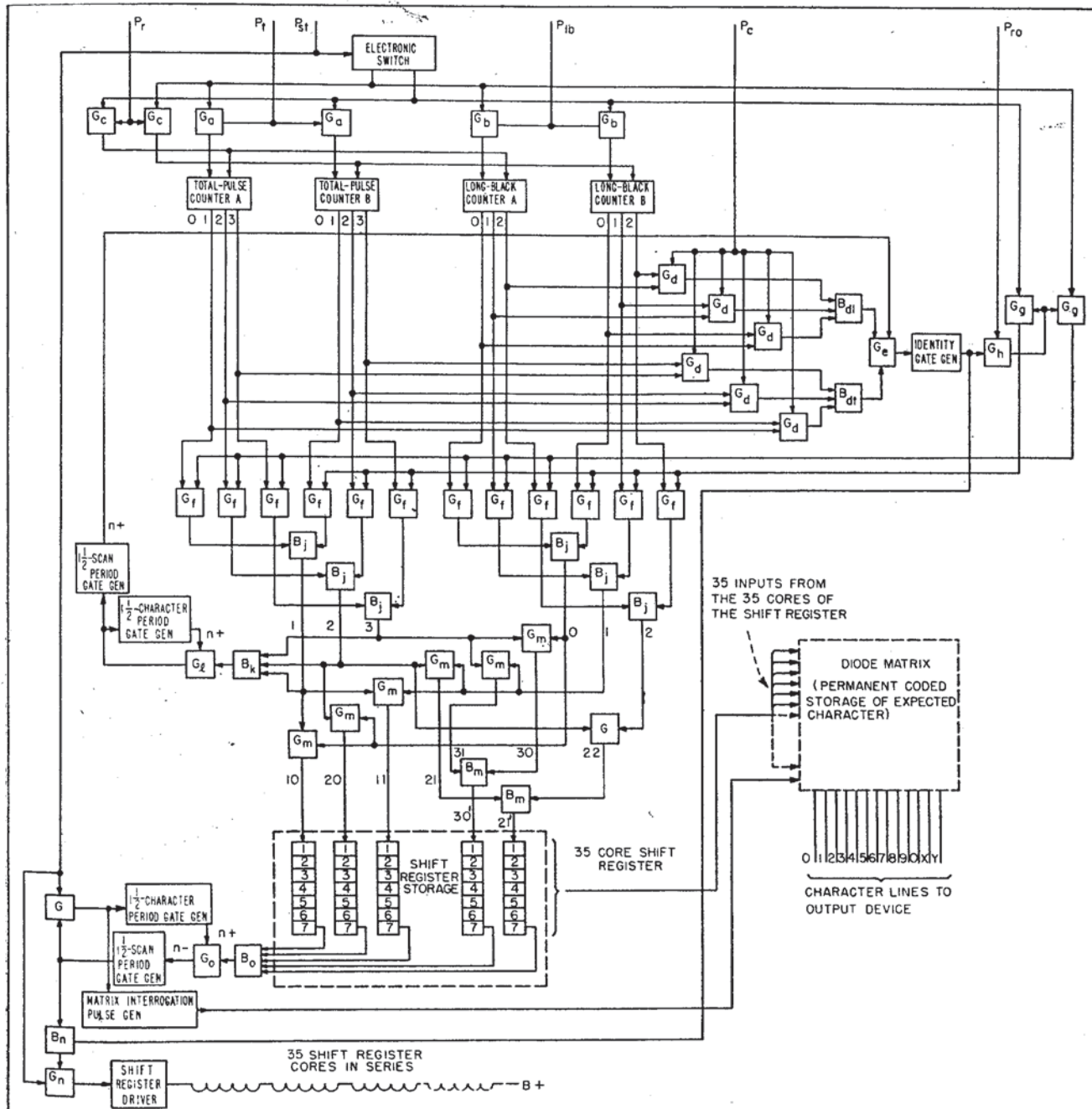
Each block in Fig. 6A represents one magnetic core. There are thirty-five cores arranged in five columns of seven each. Each of the five columns is used to store one of the five discrete code combinations.

Information is always read into row 1 and remains in row 1 until

a shift signal is applied whereupon the content of row 1 is advanced to row 2.

Access to the information in the cores is available only during the shift-pulse time. As indicated in Fig. 6A, a single shift signal suffices to advance all thirty-five cores.

Figure 6B shows the scanned character THREE. The total number of pulses  $T$  in scan one is one. There are no long pulses  $L$  (code 10). For scan two:  $T = 2, L = 0$  for code 20. Scans three, four and five also yield code 20. In scans six and seven,  $T = 3, L = 0$  for code 30. Scan eight contains only one long pulse



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