

## SOCKS Protocol Version 5

### Status of this Memo

This document specifies an Internet standards track protocol for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "Internet Official Protocol Standards" (STD 1) for the standardization state and status of this protocol. Distribution of this memo is unlimited.

### Acknowledgments

This memo describes a protocol that is an evolution of the previous version of the protocol, version 4 [1]. This new protocol stems from active discussions and prototype implementations. The key contributors are: Marcus Leech: Bell-Northern Research, David Koblas: Independent Consultant, Ying-Da Lee: NEC Systems Laboratory, LaMont Jones: Hewlett-Packard Company, Ron Kuris: Unify Corporation, Matt Ganis: International Business Machines.

### 1. Introduction

The use of network firewalls, systems that effectively isolate an organizations internal network structure from an exterior network, such as the INTERNET is becoming increasingly popular. These firewall systems typically act as application-layer gateways between networks, usually offering controlled TELNET, FTP, and SMTP access. With the emergence of more sophisticated application layer protocols designed to facilitate global information discovery, there exists a need to provide a general framework for these protocols to transparently and securely traverse a firewall.

There exists, also, a need for strong authentication of such traversal in as fine-grained a manner as is practical. This requirement stems from the realization that client-server relationships emerge between the networks of various organizations, and that such relationships need to be controlled and often strongly authenticated.

The protocol described here is designed to provide a framework for client-server applications in both the TCP and UDP domains to conveniently and securely use the services of a network firewall. The protocol is conceptually a "shim-layer" between the application layer and the transport layer, and as such does not provide network-layer gateway services, such as forwarding of ICMP messages.

## 2. Existing practice

There currently exists a protocol, SOCKS Version 4, that provides for unsecured firewall traversal for TCP-based client-server applications, including TELNET, FTP and the popular information-discovery protocols such as HTTP, WAIS and GOPHER.

This new protocol extends the SOCKS Version 4 model to include UDP, and extends the framework to include provisions for generalized strong authentication schemes, and extends the addressing scheme to encompass domain-name and V6 IP addresses.

The implementation of the SOCKS protocol typically involves the recompilation or relinking of TCP-based client applications to use the appropriate encapsulation routines in the SOCKS library.

### Note:

Unless otherwise noted, the decimal numbers appearing in packet-format diagrams represent the length of the corresponding field, in octets. Where a given octet must take on a specific value, the syntax X'hh' is used to denote the value of the single octet in that field. When the word 'Variable' is used, it indicates that the corresponding field has a variable length defined either by an associated (one or two octet) length field, or by a data type field.

## 3. Procedure for TCP-based clients

When a TCP-based client wishes to establish a connection to an object that is reachable only via a firewall (such determination is left up to the implementation), it must open a TCP connection to the appropriate SOCKS port on the SOCKS server system. The SOCKS service is conventionally located on TCP port 1080. If the connection request succeeds, the client enters a negotiation for the

authentication method to be used, authenticates with the chosen method, then sends a relay request. The SOCKS server evaluates the request, and either establishes the appropriate connection or denies it.

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The client connects to the server, and sends a version identifier/method selection message:

```

+-----+-----+-----+
| VER | NMETHODS | METHODS |
+-----+-----+-----+
| 1 | 1 | 1 to 255 |
+-----+-----+-----+

```

The VER field is set to X'05' for this version of the protocol. The NMETHODS field contains the number of method identifier octets that appear in the METHODS field.

The server selects from one of the methods given in METHODS, and sends a METHOD selection message:

```

+-----+-----+
| VER | METHOD |
+-----+-----+
| 1 | 1 |
+-----+-----+

```

If the selected METHOD is X'FF', none of the methods listed by the client are acceptable, and the client MUST close the connection.

The values currently defined for METHOD are:

- o X'00' NO AUTHENTICATION REQUIRED
- o X'01' GSSAPI
- o X'02' USERNAME/PASSWORD
- o X'03' to X'7F' IANA ASSIGNED
- o X'80' to X'FE' RESERVED FOR PRIVATE METHODS
- o X'FF' NO ACCEPTABLE METHODS

The client and server then enter a method-specific sub-negotiation.

Descriptions of the method-dependent sub-negotiations appear in separate memos.

Developers of new METHOD support for this protocol should contact IANA for a METHOD number. The ASSIGNED NUMBERS document should be referred to for a current list of METHOD numbers and their corresponding protocols.

Compliant implementations MUST support GSSAPI and SHOULD support USERNAME/PASSWORD authentication methods.

#### 4. Requests

Once the method-dependent subnegotiation has completed, the client sends the request details. If the negotiated method includes encapsulation for purposes of integrity checking and/or confidentiality, these requests MUST be encapsulated in the method-dependent encapsulation.

The SOCKS request is formed as follows:

```

+-----+-----+-----+-----+-----+-----+
|VER |  CMD |  RSV |  ATYP | DST.ADDR | DST.PORT |
+-----+-----+-----+-----+-----+-----+
|  1  |   1  | X'00'|   1  | Variable |     2   |
+-----+-----+-----+-----+-----+-----+

```

Where:

- o VER protocol version: X'05'
- o CMD
  - o CONNECT X'01'
  - o BIND X'02'
  - o UDP ASSOCIATE X'03'
- o RSV RESERVED
- o ATYP address type of following address
  - o IP V4 address: X'01'
  - o DOMAINNAME: X'03'
  - o IP V6 address: X'04'
- o DST.ADDR desired destination address
- o DST.PORT desired destination port in network octet order

The SOCKS server will typically evaluate the request based on source and destination addresses, and return one or more reply messages, as appropriate for the request type.

## 5. Addressing

In an address field (DST.ADDR, BND.ADDR), the ATYP field specifies the type of address contained within the field:

- o X'01'

the address is a version-4 IP address, with a length of 4 octets

- o X'03'

the address field contains a fully-qualified domain name. The first octet of the address field contains the number of octets of name that follow, there is no terminating NUL octet.

- o X'04'

the address is a version-6 IP address, with a length of 16 octets.

## 6. Replies

The SOCKS request information is sent by the client as soon as it has established a connection to the SOCKS server, and completed the authentication negotiations. The server evaluates the request, and returns a reply formed as follows:

```

+-----+-----+-----+-----+-----+-----+
|VER | REP | RSV | ATYP | BND.ADDR | BND.PORT |
+-----+-----+-----+-----+-----+-----+
| 1 | 1 | X'00' | 1 | Variable | 2 |
+-----+-----+-----+-----+-----+-----+

```

Where:

- o VER protocol version: X'05'
- o REP Reply field:
  - o X'00' succeeded
  - o X'01' general SOCKS server failure
  - o X'02' connection not allowed by ruleset
  - o X'03' Network unreachable
  - o X'04' Host unreachable
  - o X'05' Connection refused
  - o X'06' TTL expired
  - o X'07' Command not supported
  - o X'08' Address type not supported
  - o X'09' to X'FF' unassigned
- o RSV RESERVED
- o ATYP address type of following address

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