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Patent Application Transmittal Letter

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Transmitted herewith for filing under 37 CFR 1.53(b) is a(n): Utility () Design
(X) original patent application,
() continuation-in-part application

JC825 U.S. PTO
09/703942
10/31/00

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TITLE: Method And System For Identifying And Processing Changes To A Network Topology

Enclosed are:

- The Declaration and Power of Attorney. (X) signed () unsigned or partially signed
- 26 sheets of drawings (one set) () Associate Power of Attorney
- () Form PTO-1449 () Information Disclosure Statement and Form PTO-1449
- () Priority document(s) () (Other) _____ (fee \$ _____)

CLAIMS AS FILED BY OTHER THAN A SMALL ENTITY				
(1) FOR	(2) NUMBER FILED	(3) NUMBER EXTRA	(4) RATE	(5) TOTALS
TOTAL CLAIMS	20 — 20	0	X \$18	\$ 0
INDEPENDENT CLAIMS	3 — 3	0	X \$80	\$ 0
ANY MULTIPLE DEPENDENT CLAIMS	0		\$270	\$ 0
BASIC FEE: Design (\$320.00); Utility (\$710.00)				\$ 710
TOTAL FILING FEE				\$ 710
OTHER FEES				\$
TOTAL CHARGES TO DEPOSIT ACCOUNT				\$ 710

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"Express Mail" label no. EL523338183US

Date of Deposit Oct. 31, 2000

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1 **Title**

2 Method and System for Identifying and Processing Changes to a Network Topology

3 **Field of Invention**

4 The present invention relates generally to computer networks. More particularly, it relates
5 to a method and system for identifying changes to a network topology and for acting upon the
6 network based on the changes.

7 **Background**

8 As communications networks, such as the Internet, carry more and more traffic, efficient
9 use of the bandwidth available in the network becomes more and more important. Switching
10 technology was developed in order to reduce congestion and associated competition for the
11 available bandwidth. Switching technology works by restricting traffic. Instead of broadcasting a
12 given data packet to all parts of the network, switches are used to control data flow such that the
13 data packet is sent only along those network segments necessary to deliver it to the target node.
14 The smaller volume of traffic on any given segment results in few packet collisions on that segment
15 and, thus, the smoother and faster delivery of data. A choice between alternative paths is usually
16 possible and is typically made based upon current traffic patterns.

17 The intelligent routing of data packets with resultant reduction in network congestion can
18 only be effected if the network topology is known. The topology of a network is a description of
19 the network which includes the location of and interconnections between nodes on the network.
20 The word "topology" refers to either the physical or logical layout of the network, including devices,
21 and their connections in relationship to one another. Information necessary to create the topology
22 layout can be derived from tables stored in network devices such as hubs, bridges, and switches.
23 The information in these tables is in a constant state of flux as new entries are being added and old
24 entries time out. Many times there simply is not enough information to determine where to place a
25 particular device.

26 Switches examine each data packet that they receive, read the source addresses, and log
27 those addresses into tables along with the switch ports on which the packets were received. If a
28 packet is received with a target address without an entry in the switches table, the switch receiving it

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1 broadcasts that packet to each of its ports. When the switch receives a reply, it will have identified
2 where the new node lies.

3 In a large network with multiple possible paths from the switch to the target node, this table
4 can become quite large and may require a significant amount of the switch's resources to develop
5 and maintain. As an additional complication, the physical layout of devices and their connections
6 are typically in a state of constant change. Devices are continually being removed from, added to,
7 and moved to new physical locations on the network. To be effectively managed, the topology of a
8 network must be accurately and efficiently ascertained, as well as maintained.

9 Existing mapping methods have limitations that prevent them from accurately mapping
10 topological relationships. Multiple connectivity problems are one sort of difficulty encountered by
11 existing methods. For example, connectors such as routers, switches, and bridges may be
12 interconnected devices in a network. Some existing methods assume that these devices have only a
13 single connection between them. In newer devices, however, it is common for manufacturers to
14 provide multiple connections between devices to improve network efficiency and to increase
15 capacity of links between the devices. The multiple connectivity allows the devices to maintain
16 connection in case one connection fails. Methods that do not consider multiple connectivity do not
17 present a complete and accurate topological map of the network.

18 Another limitation of existing topology methods is the use of a single reference to identify a
19 device. Existing methods use a reference interface or a reference address in a set of devices to
20 orient all other devices in the same area. These methods assumed that every working device would
21 be able to identify, or "hear," this reference and identify it with a particular port of the device. With
22 newer devices, however, it is possible that the same address or reference may be heard out of
23 multiple ports of the same device. It is also possible that the address or reference may not be heard
24 from any ports, for example, if switching technology is used.

25 Still another limitation of existing mapping systems is that they require a complete copy of
26 the topological database to be stored in memory. In larger networks, the database is so large that
27 this really is not feasible, because it requires the computer to be very large and expensive.

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1 Still another difficulty with existing systems is that they focus on the minutia without
2 considering the larger mapping considerations. Whenever an individual change in the system is
3 detected, existing methods immediately act on that change, rather than taking a broader view of the
4 change in the context of other system changes. For example, a device may be removed from the
5 network temporarily and replaced with its ports reversed. In existing systems, this swapped port
6 scenario could require hundreds or thousands of changes because the reference addresses will have
7 changed for all interconnected devices.

8 Still another disadvantage of existing methods is that they use a continuous polling paradigm.
9 These methods continuously poll network addresses throughout the day and make decisions based
10 on those continuous polling results. This creates traffic on the network that slows other processes.

11 Still another limitation of existing methods is the assumption that network parts of a
12 particular layer would be physically separated from other parts. Network layer 1 may represent the
13 physical cabling of the network, layer 2 may represent the device connectivity, and layer 3 may
14 represent a higher level of abstraction, such as the groupings of devices into regions. Existing
15 methods assume that all layer 3 region groupings are self-contained, running on the same unique
16 physical networking. However, in an internet protocol (IP) network, multiple IP domains may co-
17 exist on the same lower layer networking infrastructure. It has become common for a network to
18 employ a virtual local area network (LAN) to improve security or to simplify network maintenance,
19 for example. Using virtual LANs, a system may have any number of different IP domains sharing
20 the same physical connectivity. As a result, existing methods create confusion with respect to
21 topological mapping because networks with multiple IP addresses in different subnets for the
22 infrastructure devices cannot be properly represented because they assume the physical separation
23 of connectivity for separate IP domains. Still another limitation of existing methods is that they do
24 not allow topological loops, such as port aggregation or trunking, and switch meshing.

25 26 **Summary of Invention**

27 A method and system are disclosed for mapping the topology of a network having
28 interconnected nodes by identifying changes in the network and updating a stored network topology

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1 based on the changes. The nodal connections are represented by data tuples that store information
2 such as a host identifier, a connector interface, and a port specification for each connection. A
3 topology database stores an existing topology of a network. A topology converter accesses the
4 topology database and converts the existing topology into a list of current tuples. A connection
5 calculator calculates tuples to represent connections in the new topology. The topology converter
6 receives the new tuples, identifies changes to the topology, and updates the topology database using
7 the new tuples. The topology converter identifies duplicate tuples that appear in both the new tuples
8 and the existing tuples and marks the duplicate tuples to reflect that no change has occurred to these
9 connections. The topology converter attempts to resolve swapped port conditions and searches for
10 new singly-heard and multi-heard host link tuples in the list of existing tuples. The topology
11 converter also searches for new conflict link tuples in the existing tuples. The topology converter
12 updates the topology database with the new topology.

13 **Summary of Drawings**

14 Figure 1 is a drawing of a typical topological bus segment for representing the connectivity
15 of nodes on a network.

16 Figure 2 is a drawing of a typical topological serial segment for representing the connectivity
17 of nodes on a network.

18 Figure 3 is a drawing of a typical topological star segment for representing the connectivity
19 of nodes on a network.

20 Figure 4 is a drawing of another typical topological star segment for representing the
21 connectivity of nodes on a network.

22 Figure 5 is a drawing of the connectivity of an example network system.

23 Figure 6 is a drawing of the connectivity of another example network system.

24 Figure 7 is a block diagram of the system.

25 Figure 8 is a flow chart of the method of the system.

26 Figure 9 is a flow chart of the method used by the tuple manager.

27 Figure 10 is a flow chart of the method used by the connection calculator.

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1 Figure 11 is a flow chart of the first weeding phase of the method used by the connection
2 calculator.

3 Figures 12a-d are flow charts of an infrastructure-building phase of the method used by the
4 connection calculator.

5 Figure 13 is a flow chart of a second weeding phase of the method used by the connection
6 calculator.

7 Figure 14 is a flow chart of the noise reduction phase of the method used by the connection
8 calculator.

9 Figure 15 is a flow chart of the look-for phase of the method used by the connection
10 calculator.

11 Figures 16a-b are flow charts of the consolidation phase of the method used by the
12 connection calculator.

13 Figure 17 is a flow chart of the method used by the topology converter.

14 Figures 18a-b are flow charts of the morph topo phase of the method used by the topology
15 converter.

16 Figure 19 is a flow chart of the duplication discard phase of the method used by the
17 topology converter.

18 Figures 20a-d are flow charts of the identify different tuples phase of the method used by
19 the topology converter.

20 Detailed Description

21 The system provides an improved method for creating topological maps of communication
22 networks based. Connectivity information is retrieved from the network nodes and stored as
23 "tuples" to track specifically the desired information necessary to map the topology. These light
24 weight data structures may store the host identifier, interface index, and a port. From this tuple
25 information, the topology may be determined. A tuple may be a binary element insofar as it has two
26 parts representing the two nodes on either end of a network link or segment. A "tuco" refers to a
27 tuple component, such as half of a binary tuple.

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1 As used herein, a node is any electronic component, such as a connector or a host, or
2 combination of electronic components with their interconnections. A connector is any network
3 device other than a host, including a switching device. A switching device is one type of connector
4 and refers to any device that controls the flow of messages on a network. Switching devices
5 include, but are not limited to, any of the following devices: repeaters, hubs, routers, bridges, and
6 switches.

7 As used herein, the term "tuple" refers to any collection of assorted data. Tuples may be
8 used to track information about network topology by storing data from network nodes. In one use,
9 tuples may include a host identifier, interface information, and a port specification for each node.
10 The port specification (also described as the group/port) may include a group number and a port
11 number, or just a port number, depending upon the manufacturer's specifications. A binary tuple
12 may include this information about two nodes as a means of showing the connectivity between them,
13 whether the nodes are connected directly or indirectly through other nodes. A "conn-to-conn"
14 tuple refers to a tuple that has connectivity data about connector nodes. A "conn-to-host" tuple
15 refers to a tuple that has connectivity data about a connector node and a host node. In one use,
16 tuples may have data about more than two nodes; that is, they may be n-ary tuples, such as those
17 used with respect to shared media connections described herein.

18 A "singly-heard host" (shh) refers to a host, such as a workstation, PC, terminal, printer,
19 other device, etc., that is connected directly to a connector, such as a switching device. A singly-
20 heard host link (shhl) refers to the link, also referred to as a segment, between a connector and an
21 shh. A "multi-heard host" (mhh) refers to hosts that are heard by a connector on the same port that
22 other hosts are heard. A multi-heard host link (mhhl) refers to the link between the connector and
23 an mhh. A link generally refers to the connection between nodes. A segment is a link that may
24 include a shared media connection.

25 Figure 1 is a drawing of a typical topological bus segment 100 for representing the
26 connectivity of nodes on a network 110. In Figure 1, first and second hosts 121, 122, as well as a
27 first port 131 of a first connector 140 are interconnected via the network 110. The bus segment

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1 100 comprises the first and second hosts 121, 122 connected to the first port 131 of the first
2 connector 140.

3 Figure 2 is a drawing of a typical topological serial segment 200 for representing the
4 connectivity of nodes on the network 110. In Figure 2, the first host 121 comprises a second port
5 132 on a second connector 145 which is connected via the network 110 to the first port 131 on the
6 first connector 140. The serial segment 200 comprises the second port 132 on the second
7 connector 145 connected to the first port 131 on the first connector 140. Figure 2 is an example of
8 a connector-to-connector (“conn-to-conn”) relationship.

9 Figure 3 is a drawing of a typical topological star segment 301 for representing the
10 connectivity of nodes on the network 110. In Figure 3, the first host 121 is connected to the first
11 port 131 of the first connector 140. The star segment 301 comprises the first host 121 connected
12 to the first port 131 of the first connector 140. Figure 3 is an example of a connector-to-host
13 (“conn-to-host”) relationship.

14 Figure 4 is a drawing of another typical topological star segment 301 for representing the
15 connectivity of nodes on the network 110. In addition to the connections described with respect to
16 Figure 3, a third host 123 is connected to a third port 133 of the first connector 140 and a fourth
17 host 124 is connected to a fourth port 134 of the first connector 140. In Figure 4, the star segment
18 301 comprises the first host 121 connected to the first port 131 of the first connector 140, the third
19 host 123 connected to the third port 133 of the first connector 140, and the fourth host 124
20 connected to the fourth port 134 of the first connector 140. Thus, the star segment 301 comprises,
21 on a given connector, at least one port, wherein one and only one host is connected to that port,
22 and that host. In the more general case, the star segment 301 comprises, on a given connector, all
23 ports having one and only one host connected to each port, and those connected hosts. Since the
24 segments, or links, drawn using the topological methods of Figure 4 resemble a star, they are
25 referred to as star segments.

26 For illustrative purposes, nodes in the figures described above and in subsequent figures are
27 shown as individual electronic devices or ports on connectors. Also, in the figures the nodes are

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1 represented as terminals. However, they could also be workstations, personal computers, printers,
2 scanners, or any other electronic device that can be connected to networks 110.

3 Figure 5 is a drawing of the connectivity of an example network system. In Figure 5, first,
4 third, and fourth hosts 121, 123, 124 are connected via the network 110 to first, third, and fourth
5 ports 131, 133, 134 respectively, wherein the first, third, and fourth ports 131, 133, 134 are
6 located on the first connector 140.

7 The first, third and fourth hosts 121, 123, 124 are singly-heard hosts connected to separate
8 ports 131, 133, 134 of a common connector 140 – the first connector 140. The fifth and sixth
9 hosts 125, 126 are singly-heard hosts connected to the third and fourth connectors 142, 143. The
10 seventh and eighth hosts 127, 128 are multi-heard hosts connected to the same port 139 of the fifth
11 connector 144. The multi-heard hosts 127, 128 illustrate a shared media segment 180, also
12 referred to as a bus 180.

13 The second, third, fourth, and fifth connectors 141, 142, 143, 144 are interconnected and
14 illustrate a switch mesh 181. Each of the connectors in the switch mesh 181 is connected to each
15 other, either directly or indirectly, to create a fully meshed connection. In the mesh, traffic may be
16 dynamically routed to create an efficient flow.

17 Figure 5 also shows an example of a port aggregation 182, also referred to as trunking 182.
18 The first connector 140 is connected via the network 110 to the second connector 141 by two
19 direct links, each of which is connected to different ports on the connectors. One link is connected
20 to the sixth port 136 of the first connector 140 and to the seventh port of the second connector
21 137. The other link is connected to fifth port 135 of the first connector 140 and to the eighth port
22 138 of the second connector 141. In this example, two connectors illustrate the multiple
23 connectivity between nodes. Depending upon the device specifications, devices such as connectors
24 may be connected via any number of connectors. As explained herein, the system resolves multiple
25 connectivity problems by tracking port information for each connection.

26 Figure 6 is a drawing of the connectivity of a portion of a network having three connectors
27 171, 172, 173. A first host 151 is connected directly to the first port 161 of the first connector 171
28 and the second host 152 is connected to a sixth port 166 of the third connector 173. The second

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1 port 162 of the first connector 171 is connected directly to the third port 163 of the second, or
2 intermediate, connector 172. The fourth port 164 of the intermediate connector 172 is connected
3 directly to the fifth port 165 of the third connector 173.

4 Figure 7 shows a block diagram of the system. Figure 8 shows a flow chart of the method
5 used by the system to retrieve and update the topology of the network. A tuple manager 300, also
6 referred to as a data miner 300, gathers 902 data from network nodes and builds 904 tuples to
7 update the current topology. The topology database "topodb" 350 stores the current topology for
8 use by the system. The "neighbor data" database 310 stores new tuple data retrieved by the tuple
9 manager 300. The connection calculator 320 processes the data in the neighbor data database 310
10 to determine the new network topology. The connection calculator 320 reduces 906 the tuple data
11 and sends it to the reduced topology relationships database 330. The topology converter 340 then
12 updates 908 the topology database 350 based on the new tuples sent to the reduced topology
13 relationships database 330 by the connection calculator 320.

14 Figure 9 shows a flow chart of one operation of the tuple manager 300, as described
15 generally by the data gathering 902 and tuple building 904 steps of the method shown in Figure 8.
16 The tuple manager 300 receives 910 a signal to gather tuple data. The tuple manager 300 then
17 retrieves 912 node information of the current topology stored in the topology database 350. This
18 information tells the tuple manager 300 which devices or nodes are believed to exist in the system
19 based on the nodes that were detected during a previous query. The tuple manager 300 then
20 queries 914 the known nodes to gather the desired information. For example, the connectors may
21 maintain forwarding tables that store connectivity data used to perform the connectors' ordinary
22 functions, such as switching. Other devices may allow the system to perform queries to gather
23 information about the flow of network traffic. This data identifies the devices heard by a connector
24 and the port on which the device was heard. The tuple manager 300 gathers this data by accessing
25 forwarding tables and other information sources for the nodes to determine such information as their
26 physical address, interface information, and the port from which they "hear" other devices. Based
27 on this information, the tuple manager 300 builds 916 tuples and stores 918 them in the "neighbor
28 data" database 310. Some nodes may have incomplete information. In this case, the partial

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1 404 whether the tuple is a connector-to-host (conn-to-host) link tuple. If it is not a conn-to-host
2 link, the connection calculator 320 concludes 418 that it is a conn-to-conn link and processes 402
3 the next tuple. If the tuple is a conn-to-host link tuple, then the connection calculator 320
4 determines 406 whether the connector hears only this particular host on the port identified in the
5 tuple. If the connector hears other hosts on this port, then the tuple is classified 416 as a multi-
6 heard host link (mhh) tuple.

7 If the connector hears only the one host on the port – that is, if the host is a singly-heard
8 host – then the connection calculator 320 determines 408 whether the host is heard singly by any
9 other connectors. If no other connectors hear the host as a singly-heard host, then the tuple is
10 classified as a singly-heard host link (shhl) tuple 412 and other tuples for this host are classified 414
11 as extra host links (ehl). Another tuple for this host may be, for example, an intermediate connector
12 connected indirectly to a host. For example, Figure 6 shows three connectors 171, 172, 173 the
13 first connector is connected directly to the first host 151. This connection therefore forms an shhl
14 tuple. The intermediate connector 172 is indirectly connected to the first host 151. The tuple data
15 indicates that the intermediate connector 172 is indirectly connected to the host and hears the host
16 from a particular port. An extra host links tuple is created so that this data may be used later in
17 conjunction with other extra host links tuples from devices across the network, to verify connectivity
18 by providing hints about connections.

19 The first weeding process also attempts to identify conflicts. If other connectors hear the
20 host as a singly-heard host, then a conflict arises and the tuple is classified 410 as a singly-heard
21 conflict link (shcl) tuple to be resolved later. This conflict may arise, for example, if a host has been
22 moved within the network, in which case the forwarding table data may no longer be valid. Certain
23 connectors previously connected directly to the host may still indicate that the moved host is
24 connected. When all tuples have been processed 402 to identify singly-heard host links, the first
25 weeding phase 922 is complete.

26 Figures 12a-d show a flow chart of the infrastructure building phase 924 of the connection
27 calculator 320. The purpose of the infrastructure building phase 924 is to determine how the
28 connectors are set up in the network. The first part of the infrastructure building phase 924

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1 calculator 320 considers 472 conn2 of the tuple. If there is a tuple in conn-to-conn links for conn2
2 and conn1 474, and if there is a conn2-to-conn1 tuple in the extra host links tuples 476, and if the
3 port is the same for conn2 hearing conn1 and host1 478, then the search tuple is moved 480 into
4 the singly heard host links and other tuples containing host1 are removed 482 from the
5 singleConflictLinks.

6 Figure 14 shows a flow chart of the noise reduction phase 928. The purpose of the noise
7 reduction phase 928 is to handle those connections in which a connector is not directly connected
8 to a host or to another connector. For example, networking technology may employ shared media
9 connections between connectors, rather than dedicated media connectors. With a shared media
10 connection, the entries in the forwarding tables for connectors attached to the shared media
11 connection will include every node accessing the shared media connection and may not present a
12 useful or accurate representation of the nodal connection. For example, if the network configuration
13 in Figure 6 used a shared media connection between the first connector 171 and the intermediate
14 connector 172, then the first connector is not really connected directly to the intermediate connector
15 because other devices (not shown in Figure 6) may also use the shared media connection. These
16 other devices may include web servers, other connectors, other subnetworks, etc. Tuples will be
17 created for the connectors 171, 172 on opposing ends of the shared media. In this situation, it is
18 inefficient to maintain point-to-point binary tuples for every connection. The noise reduction phase
19 928 disproves invalid tuples created by the shared media connections.

20 For each multi-heard host links (mhhl) tuple, also referred to as multiHeardLinks (mhl)
21 tuples (sometimes referred to as the search tuple) 484, conn1 and host1 are considered 486. For
22 each extra host links tuple containing host1 488, conn2 is considered 490. If there is a tuple in
23 conn-to-conn links for conn2 and conn1 492, and if there is a conn2-to-host1 tuple in
24 extraHostLinks 494, and if the group/port for conn2 hearing conn1 and host1 is different 496, then
25 the search tuple is moved 498 to extraHostLinks.

26 Figure 15 shows a flow chart for the "look for" phase 930. The purpose of this phase is to
27 complete missing data for mhhl tuples. There may exist connections on the network that have
28 incomplete tuple data. For example, the network may simply have no traffic between certain nodes,

1 systems that track higher levels of abstraction, such as layer 3 connectivity. Also, the tuples may
 2 contain only selected information to minimize the storage space required for the topology.

3 Figures 16a-b show a flow chart of the consolidation phase 932. The purpose of this phase
 4 is to consolidate the tuples that involve shared media connections. After the noise reduction phase
 5 928, a considerable number of tuples involving shared media may remain. Rather than maintain a
 6 binary tuple for each of the connections, an n-ary tuple is created for the link using a tuco for each
 7 connector and each host connected thereto. For each mhhl tuple 518, conn1 and host1 are
 8 considered 520. If there are more conn1 group/port tuples in multiHeardLinks, and if are not any
 9 n-ary multiHeardSegments (mhs) tuples 524, then an mhs tuple is created 526. If host1 is not
 10 already in this particular mhs tuple 528, then conn2 of the tuple is considered 534. If there is a
 11 conn1-to-conn2 conn-to-connLinks tuple on the same port as conn1-to-host1 536, then all
 12 multiHeardLinks tuples for conn2-to-host1 with the same conn2 group/port as the conn1-to-conn2
 13 are added 538 to the current mhs tuple.

14 After processing each mhhl tuple 518, each singly-heard host links (shhl) tuple, also referred
 15 to as a singlyHeardLinks (shl) tuple, is considered 540. For each shhl tuple, the connector and host
 16 are considered 542. If there is no existing singlyHeardSegments (shs) tuple for the connector 544,
 17 then an shs tuple is created 546. The host tuco is then added to the shs 548.

18 Figure 17 shows a flow chart of the method used by the topology converter 340, as
 19 described generally by the topology update step 908 of the method shown in Figure 8. The
 20 topology converter 340 converts 934 the topology into tuple lists, also referred to as the “morph
 21 topo” phase 934. It then compares 936 the list from the topology currently stored in the topology
 22 database 350 with the new list generated by the connection calculator 320 and discards 936
 23 identical tuples in what is also referred to as the “discard duplicates” phase 936. It then takes
 24 action 938 on the changes in the topology as determined by the changes in the tuple lists, in what is
 25 also referred to as the “identify different tuples” phase 938.

26 Figure 18a shows a flow chart for the “morph topo” phase 934. For each node in the
 27 topology 550, the topology converter 340 determines 552 whether the node is a connector. If the
 28 node is a connector, then for each connected interface (conniface) of the connector (conn1) 554,

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1 exact match in the current tuples stored in the topodb. If an exact match is found 604, then the new
2 tuple is marked 606 as “no change” indicating that this is an identical tuple.

3 Figures 20a-d show a flow chart for the identify different tuples phase 938. The system
4 looks through each tuple in the new SinglyHeardSegments (newSHS) tuple list 608 and tries to
5 identify and fix 610 swapped ports on connectors. Swapped ports are identified by considering
6 those segment tuples in both the new topology and the existing topology that differ only by the port
7 specification in the tuco. Each tuple that is fixed as a swapped port is marked 612 as “handled.”
8 The system also looks through each tuple in the new multiHeardSegments tuple list (newMHS) 614
9 and tries to identify and fix 616 swapped ports on connectors. Each tuple that is fixed as a
10 swapped port is marked 618 as “handled.”

11 The system then processes 620 each unmarked tuple in the newSHL tuples. Four cases
12 are possible for the host of the newSHL tuples. The host of the newSHL can be found in the
13 current singlyHeardLinks (curSHL) 622, the current multiHeardLinks (curMHL) 630, the current
14 connLinks (curCL) 638, or the current UnheardOfLinks (curUOL) 642. If the host of a newSHL
15 tuple is found 622 in the current SinglyHeardLinks (curSHL) tuples, then the system determines 624
16 if there is a matching connector tuco between the newSHL tuples and the curSHL tuples. If there is
17 a matching tuco, then the system changes 626 the host connection attribute. If there is not a
18 matching tuco, then the host connection is moved 628 in the topology.

19 If the host is found in the curMHL tuples 630, then the system determines 632 whether
20 there is a matching connector tuco between the newSHL tuples and the curSHL tuples. If there is a
21 matching connector, then the segment type of connection is changed 634. If there is not a matching
22 connector, then the host connection is moved 636 in the topology. If the host is found in the curCL
23 tuples 638, then the host is moved 640 into a star segment of the connector. If it is found in the
24 curUOL 642, then the host is moved 644 into the star segment of the connector.

25 Figure 20c shows another stage of the processing undertaken during the identify different
26 tuples phase 938. For each unmarked tuple in the new multiHeardLinks tuples (newMHL) 946,
27 four cases are possible for the host of the newMHL. The host of the newMHL may be found in the
28 curSHL 648, the curMHL 656, the curCL 664, or the curUOL 668. If the host is found in the

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1 curSHL 648, then the system determines 650 whether there is a matching connector tuco between
2 the newMHL and the curMHL. If there is a matching tuco, then the segment type of connection is
3 changed 652. If there is not a matching tuco, then the host connection is moved 654 in the
4 topology.

5 If the host is found in the curMHL tuples 656, then the system determines 658 whether
6 there is a matching connector tuco in both the curMHL tuples and the newMHL tuples. If there is a
7 matching connector tuco, then the host connection attribute is changed 660. If there is not a
8 matching tuco, then the host connection is moved 662 in the topology. If the host is found in the
9 curCL tuples 664, then the host is moved into a bus segment of a connector. If the host is found in
10 the curUOL tuples 668, then the host connection is moved 670 in the topology.

11 Figure 20d shows another portion of the identify different tuples phase 938. For each
12 unmarked tuple in the newCL tuples 672, there are three possibilities for the connector. The
13 connector of the unmarked tuple in newCL can be found in the curSHL or curMHL 674, in the
14 curCL 678, or in the curUOL 682. If each connector is found in the curSHL or curMHL list 674,
15 then the system creates 676 a new point-to-point segment for the connectors. If the connectors are
16 found in the curCL 678, then the connection attributes of the connectors are changed 680. If each
17 connector is found in the curUOL tuples 682, then the host connection is moved 684 in the
18 topology.

19 Another part of the identify different tuples phase 938 is shown in blocks 686 and 688 of
20 Figure 20d. For each unmarked tuple in the newUOL tuples 686, the system checks 688 the
21 timer/configuration to determine whether the host/conn should move into the default segment from
22 its current segment.

23 An advantage of the system is that it may be schedulable. The system may map network
24 topology continuously, as done by existing systems, or it may be scheduled to run only at certain
25 intervals, as desired by the user. A further advantage of the system is that it is capable of
26 processing multiple connections between the same devices and of processing connection meshes,
27 because it tracks each nodal connection independently, without limitations on the types of
28 connections that are permitted to exist.

1 Although the present invention has been described with respect to particular embodiments
2 thereof, variations are possible. The present invention may be embodied in specific forms without
3 departing from the essential spirit or attributes thereof. It is desired that the embodiments described
4 herein be considered in all respects illustrative and not restrictive and that reference be made to the
5 appended claims for determining the scope of the invention.

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Claims

- 1
- 2 1. In a network having interconnected nodes with data tuples that represent nodal
- 3 connections, a method for mapping a network topology by identifying changes between an existing
- 4 topology and a new topology, the method comprising:
- 5 converting an existing topology into a list of existing tuples that represent existing nodal
- 6 connections;
- 7 receiving new tuples that represent new nodal connections; and
- 8 comparing the list of existing tuples with the new tuples to identify changes to the topology.
- 9 2. The method of claim 1, further comprising updating a topology database with a new
- 10 topology.
- 11 3. The method of claim 1, further comprising taking action on the changes to the
- 12 topology.
- 13 4. The method of claim 1, wherein the tuples include information about a host
- 14 identifier, a connector interface, and a port specification.
- 15 5. The method of claim 1, wherein the step of comparing comprises identifying
- 16 duplicate tuples that appear both in the list of existing tuples and in the new tuples, and maintaining a
- 17 current status of the topology for these tuples.
- 18 6. The method of claim 1, wherein the step of comparing comprises identifying a
- 19 swapped port condition on a connector.
- 20 7. The method of claim 1, wherein the step of comparing comprises searching for a
- 21 host of a new singly-heard host link tuple or a new multi-heard host link tuple in the list of existing
- 22 tuples.
- 23 8. A system for mapping a network topology by identifying changes between an
- 24 existing topology and a new topology, based on changes to data tuples that represent nodal
- 25 connections comprising:
- 26 a topology database that stores an existing topology of a network; and

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1 17. The method of claim 15, wherein the step of comparing comprises identifying
2 duplicate tuples that appear both in the list of existing tuples and in the new tuples, and maintaining a
3 current status of the topology for these tuples.

4 18. The method of claim 15, wherein the step of comparing comprises identifying a
5 swapped port condition on a connector.

6 19. The method of claim 15, wherein the step of comparing comprises searching for a
7 host of a new singly-heard host link tuple or a new multi-heard host link tuple in the list of existing
8 tuples.

9 20. The method of claim 15, wherein the step of comparing comprises searching for a
10 connector of a new conflict links tuple in the list of existing tuples.

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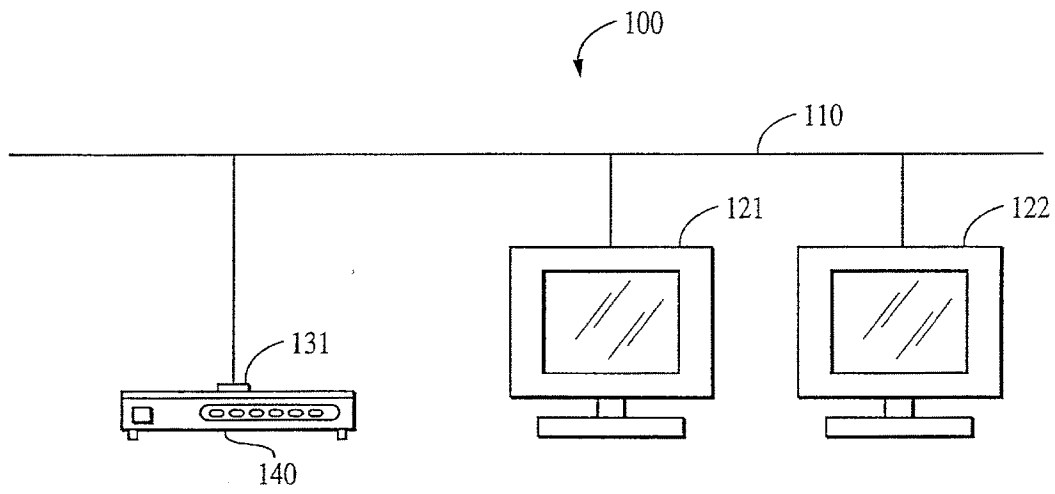


FIG. 1

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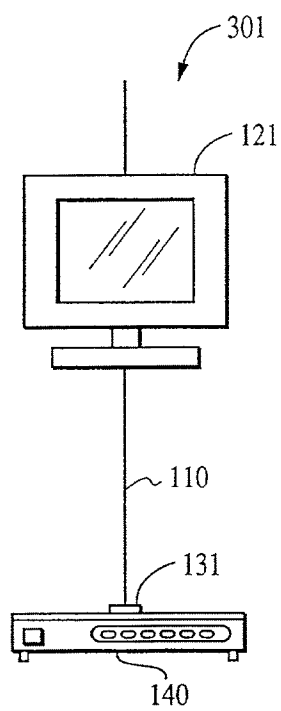
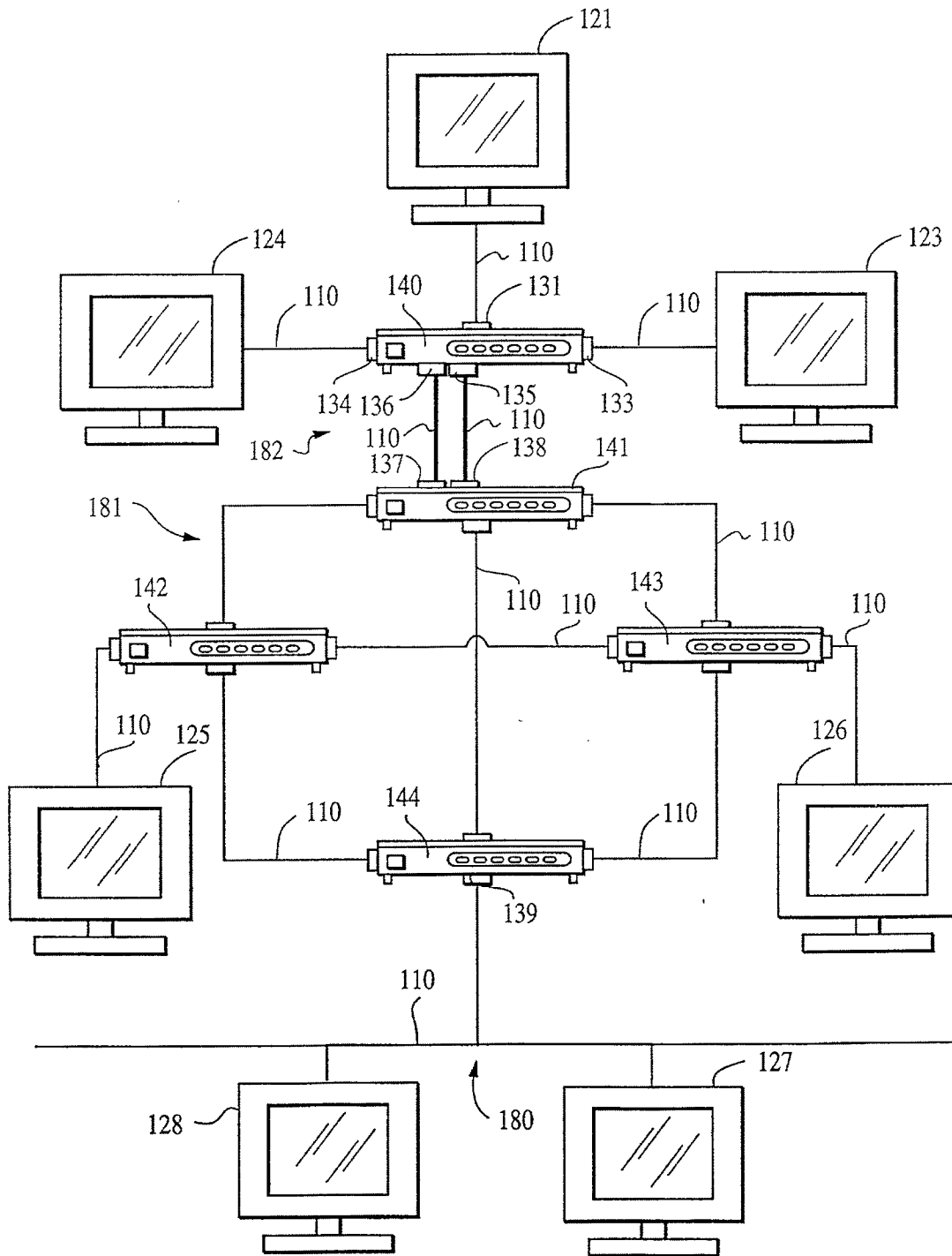


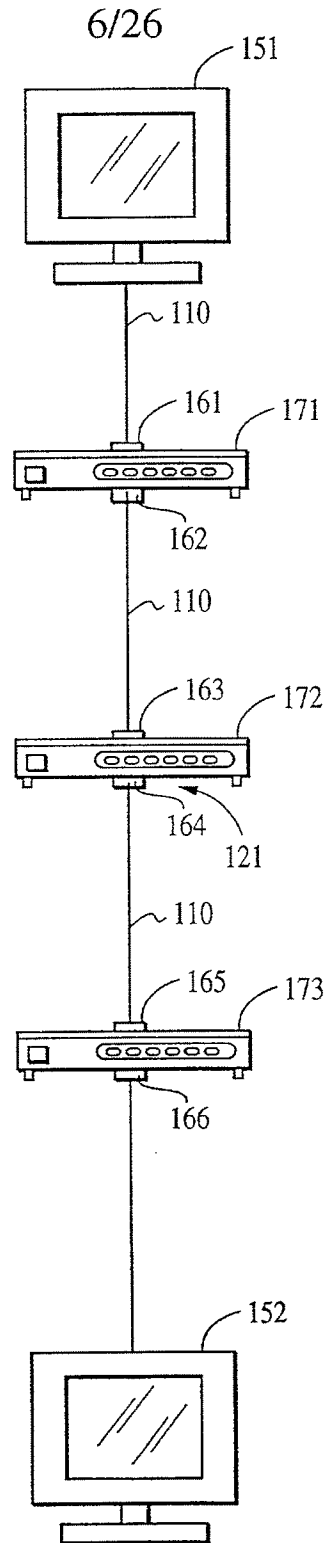
FIG. 3

FIG. 5



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FIG. 6



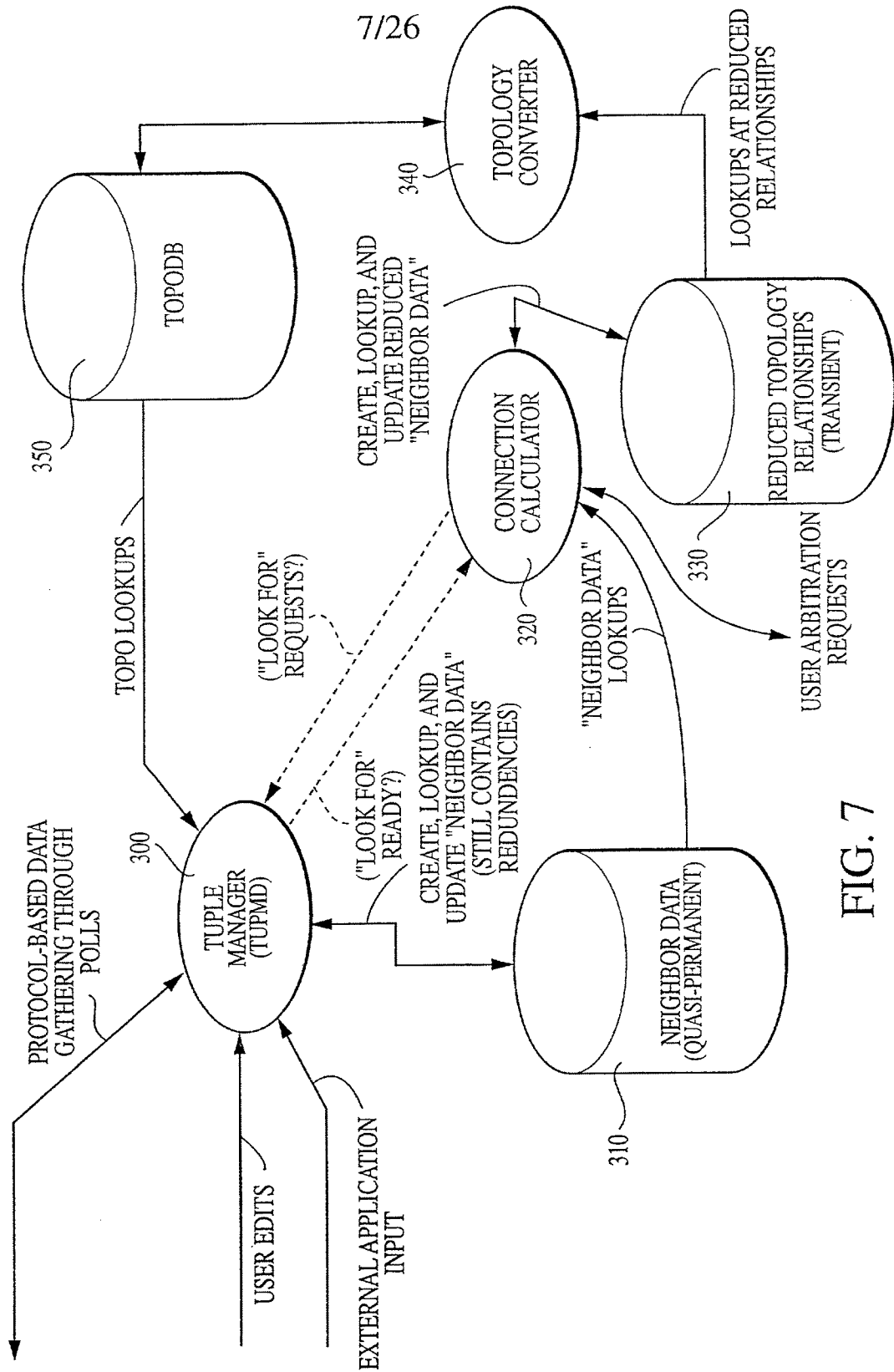


FIG. 7

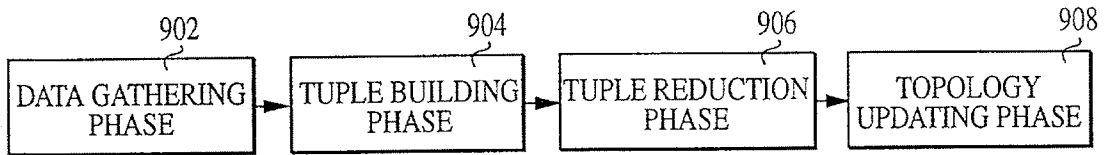


FIG. 8

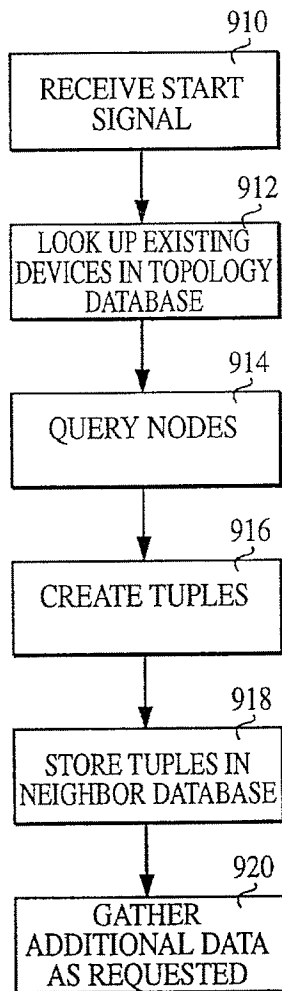


FIG. 9

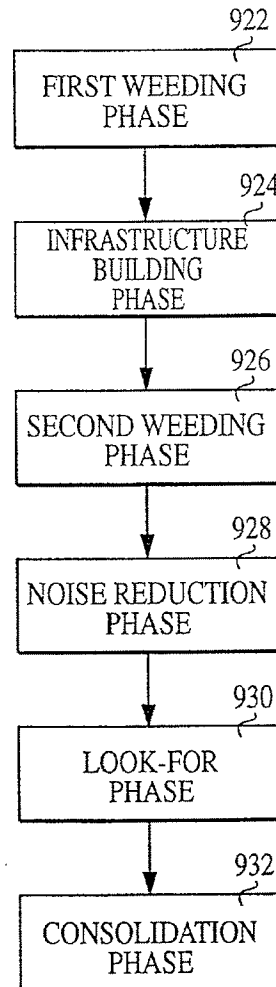


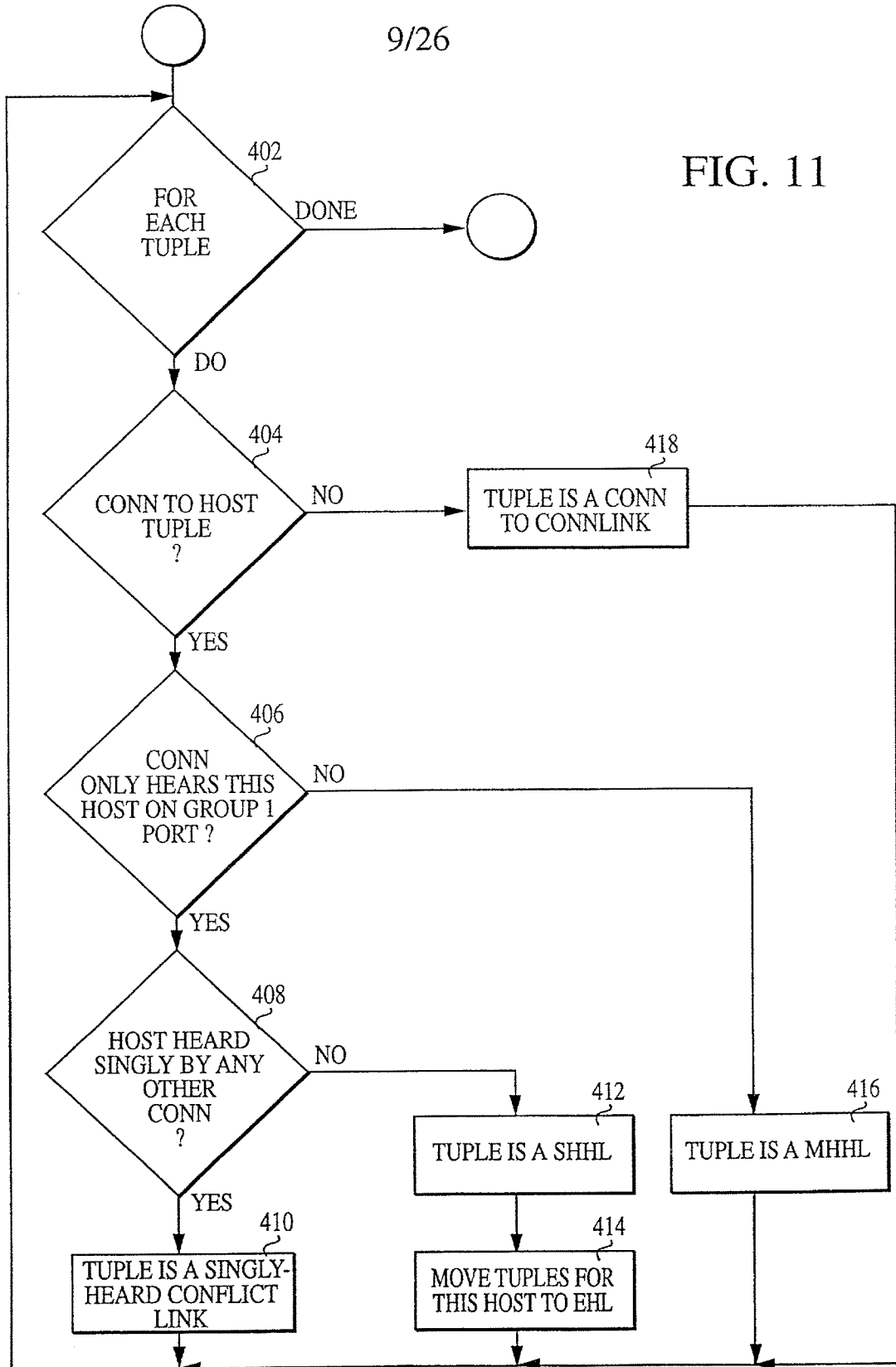
FIG. 10

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FIG. 11



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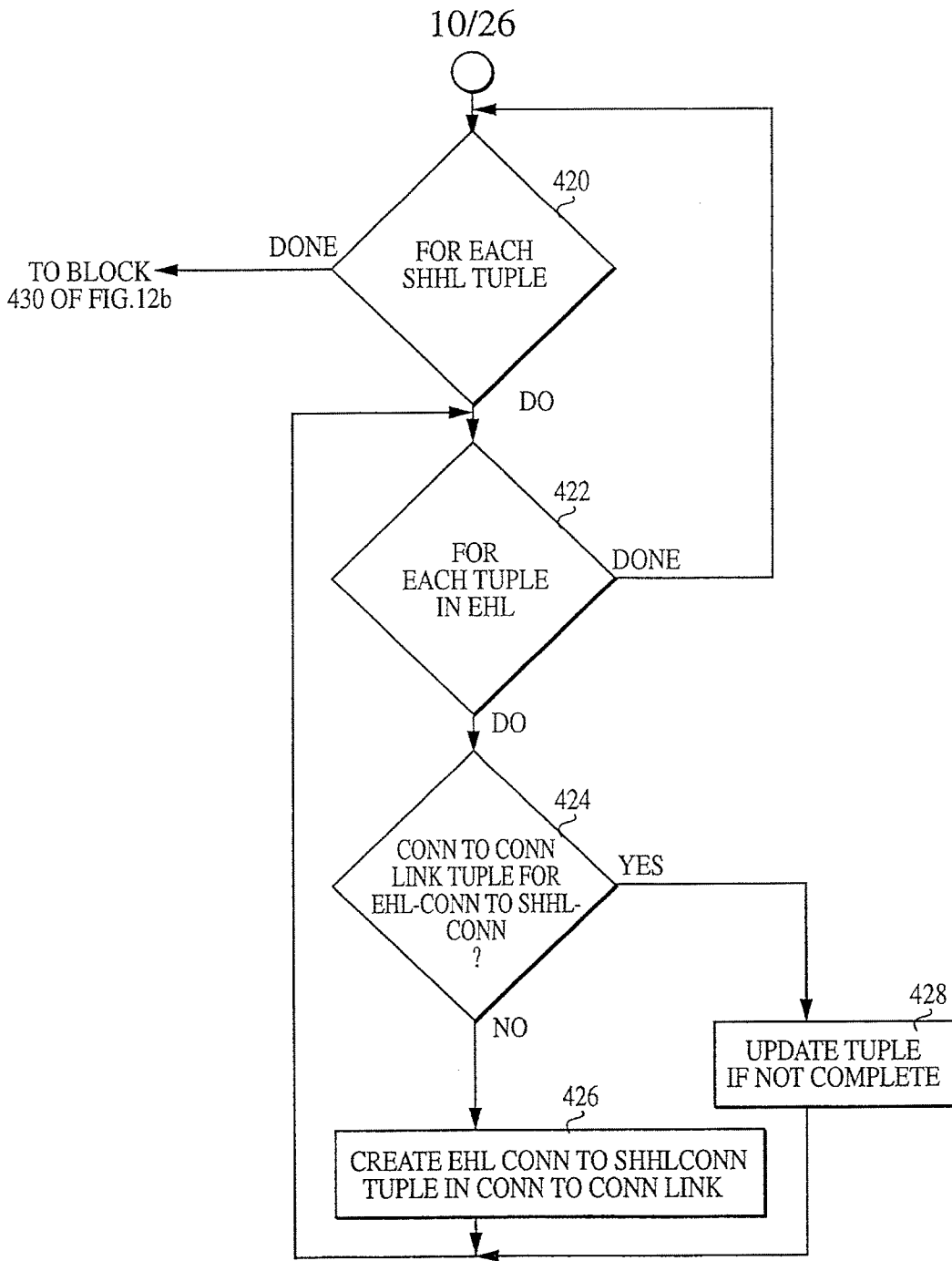
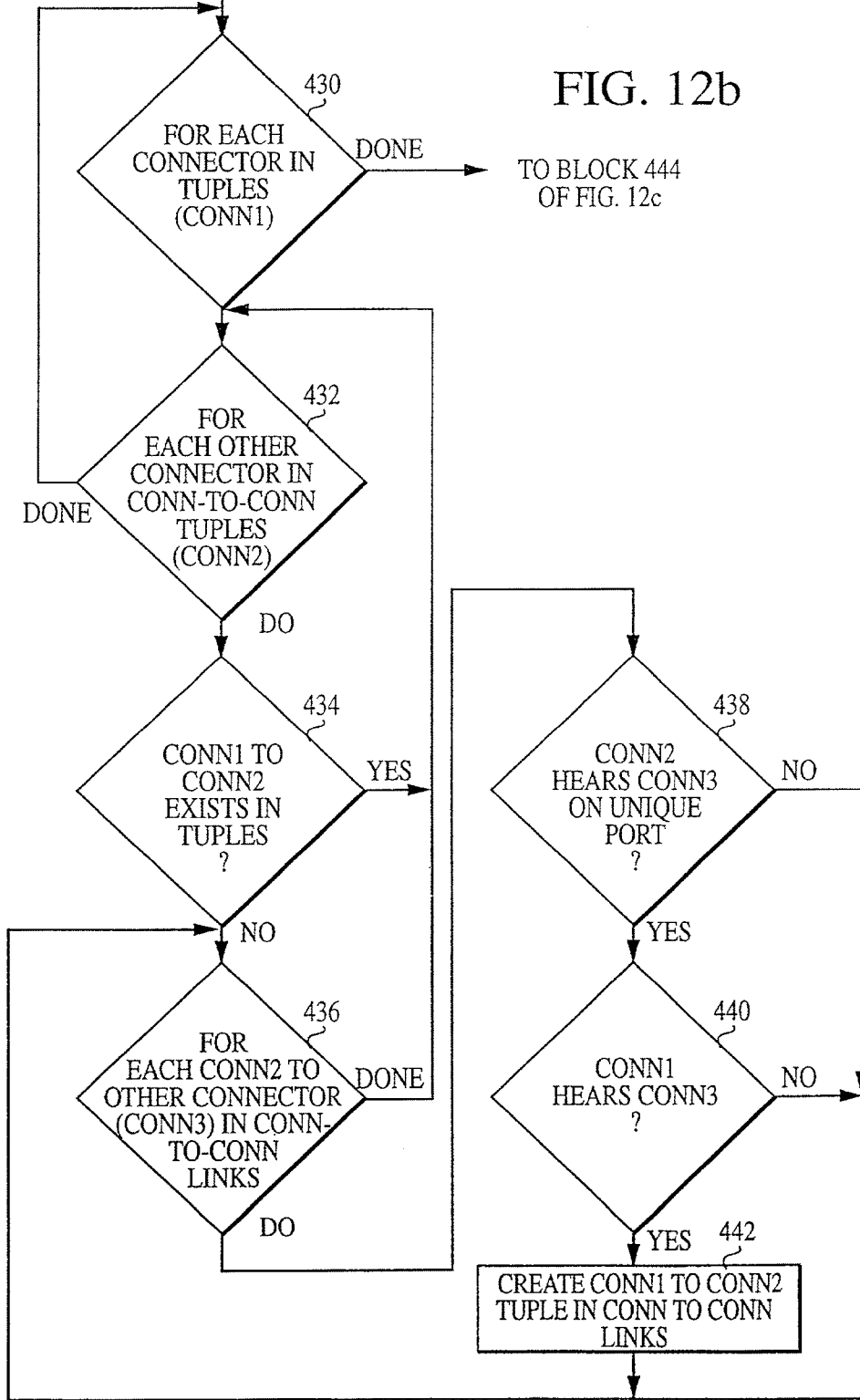


FIG. 12a

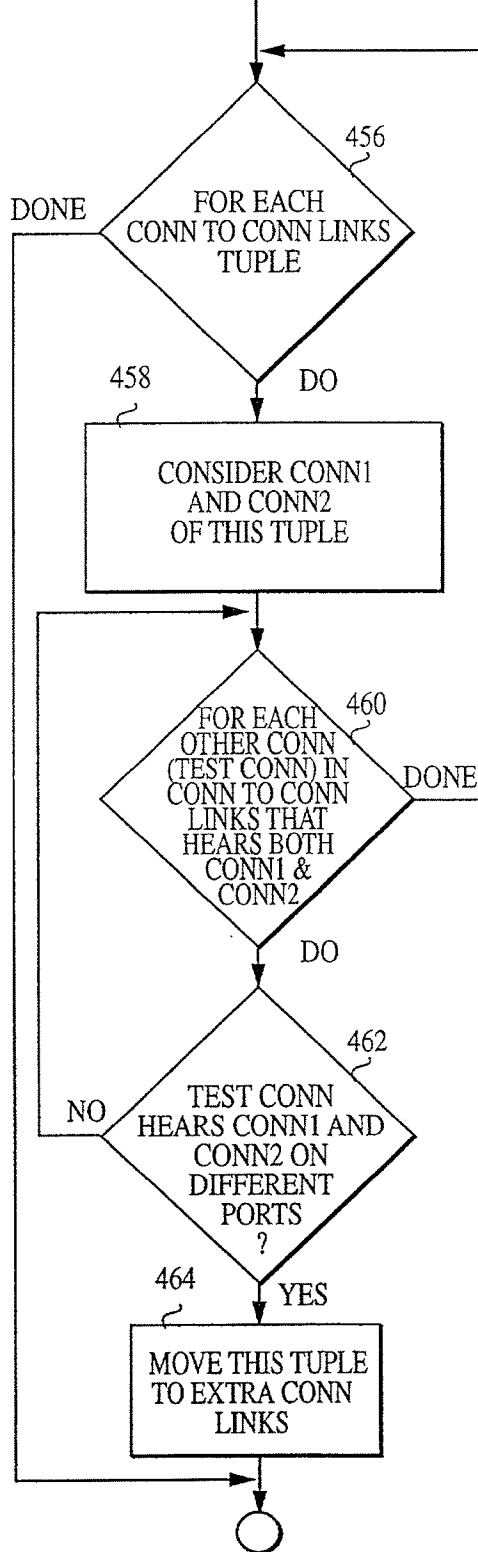
FROM BLOCK 420
OF FIG. 12a 11/26

FIG. 12b



13/26
FROM BLOCK 444 OF FIG. 12c

FIG. 12d



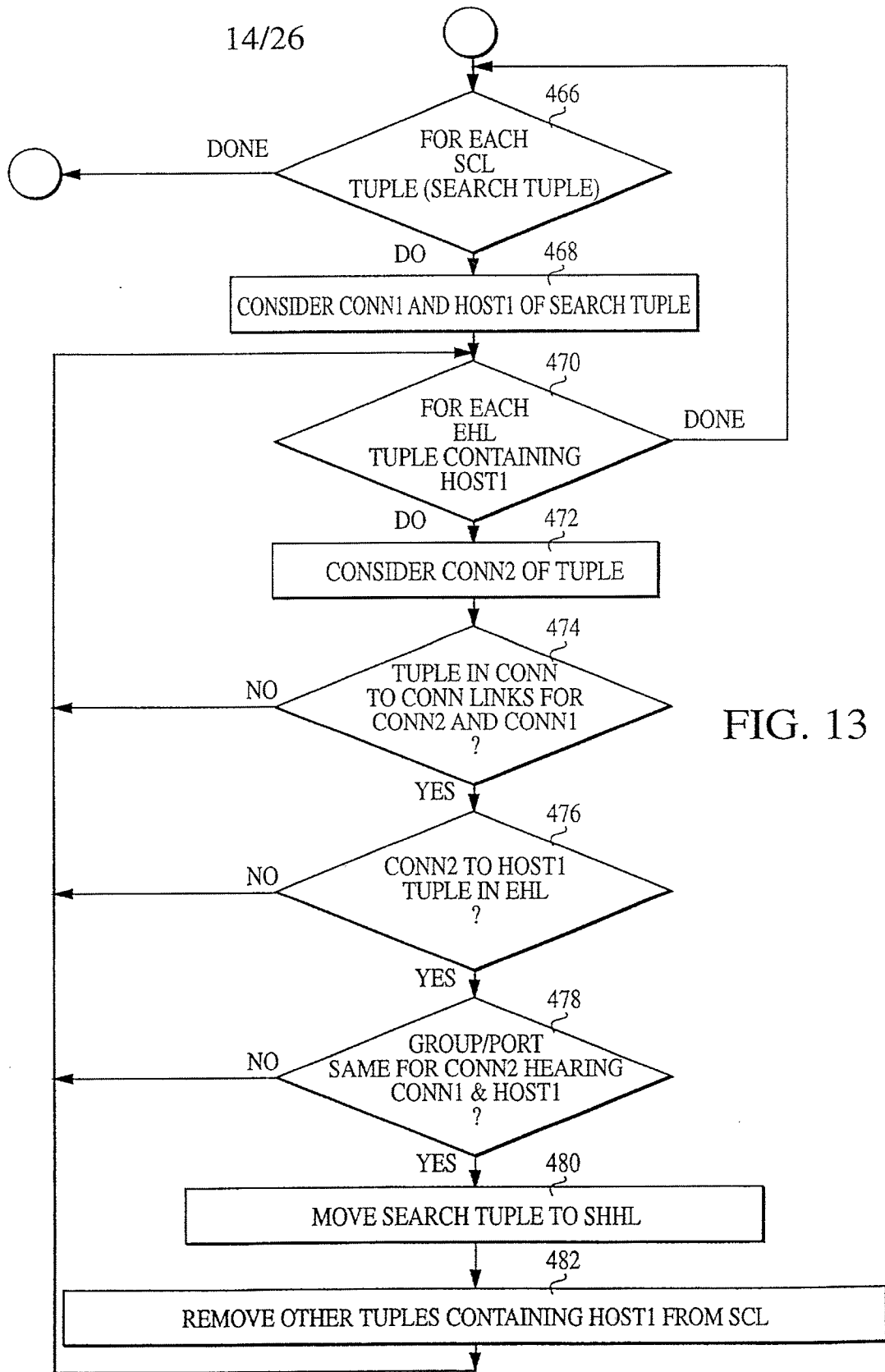


FIG. 13

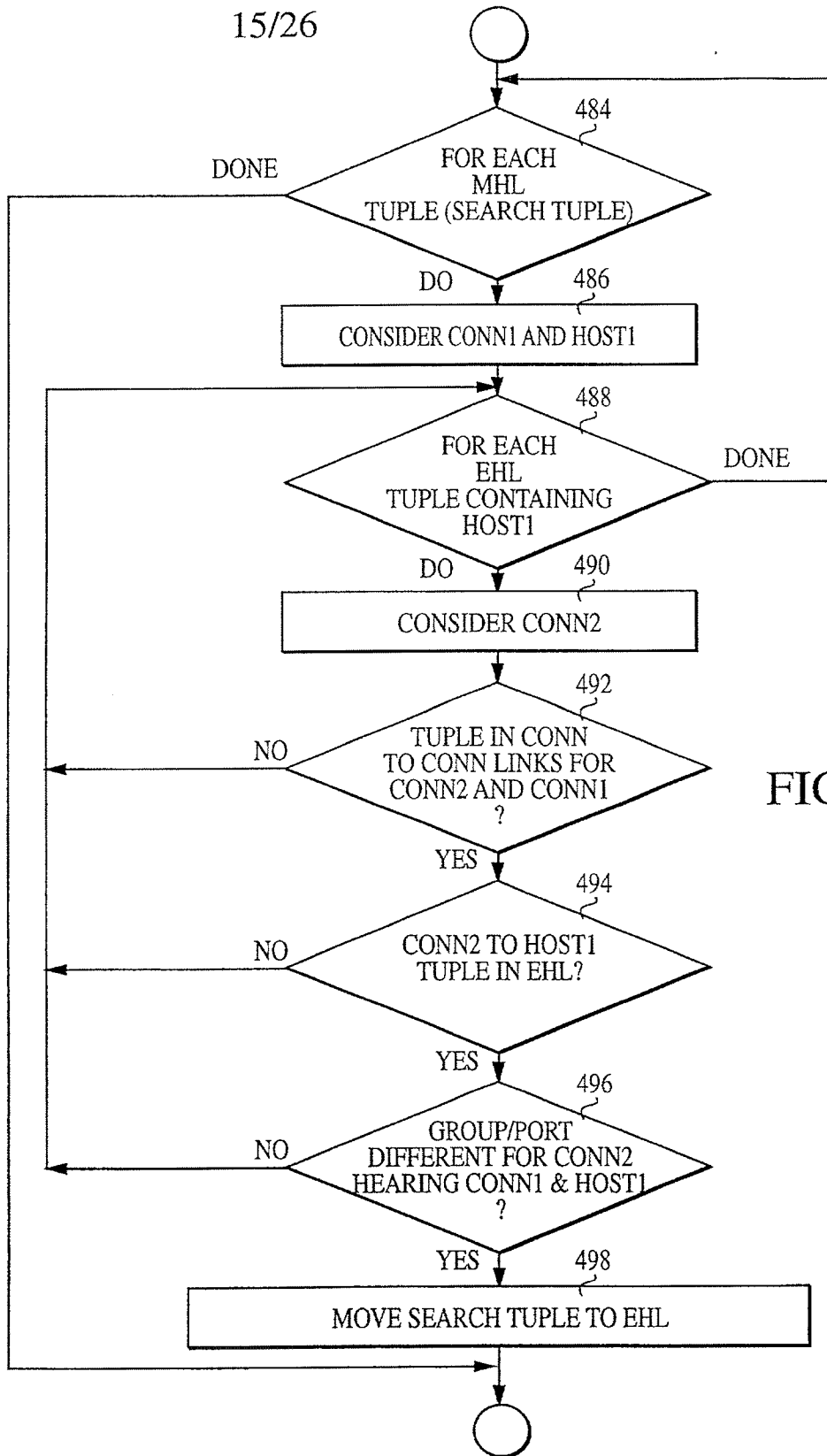
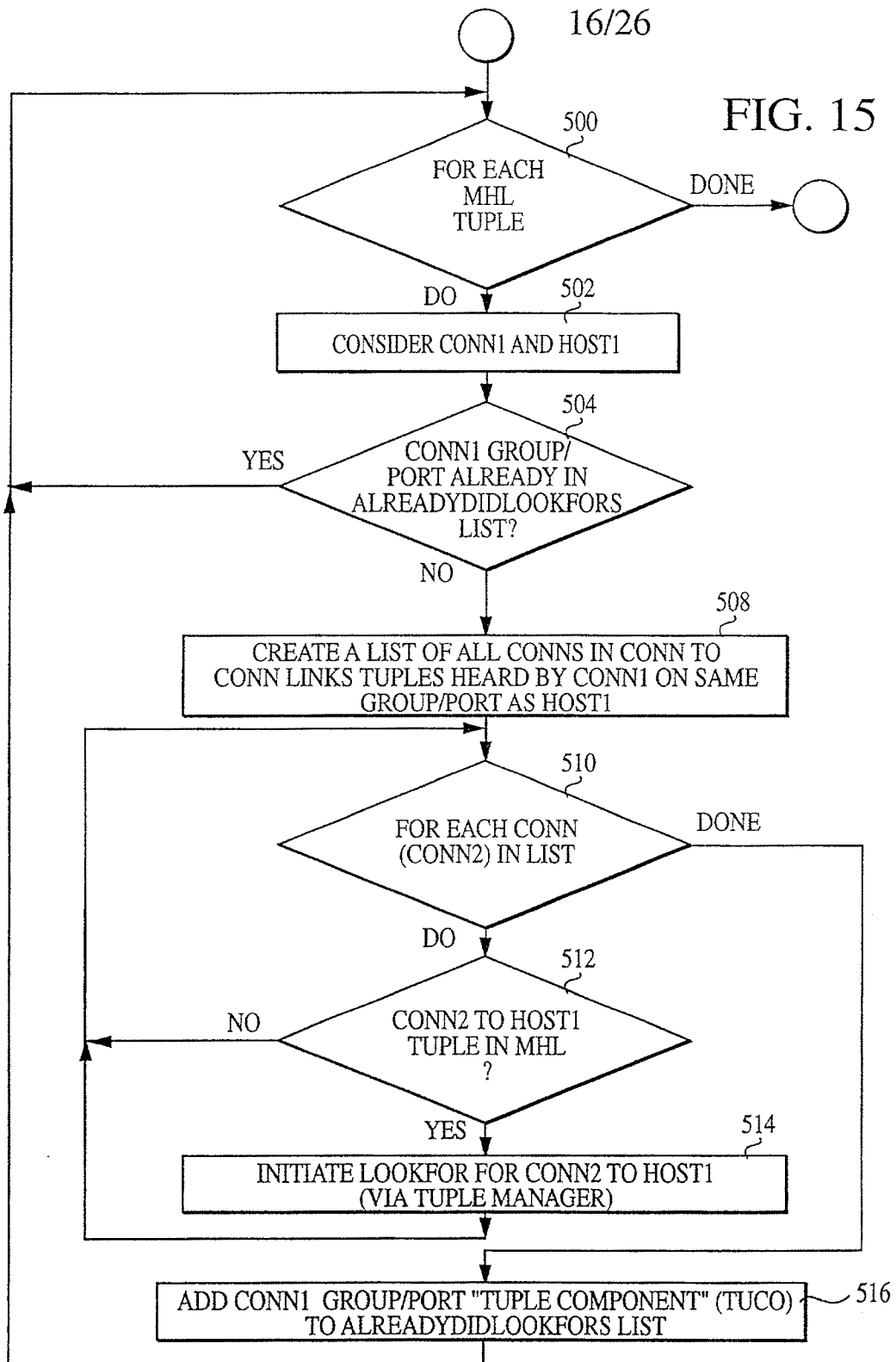


FIG. 14



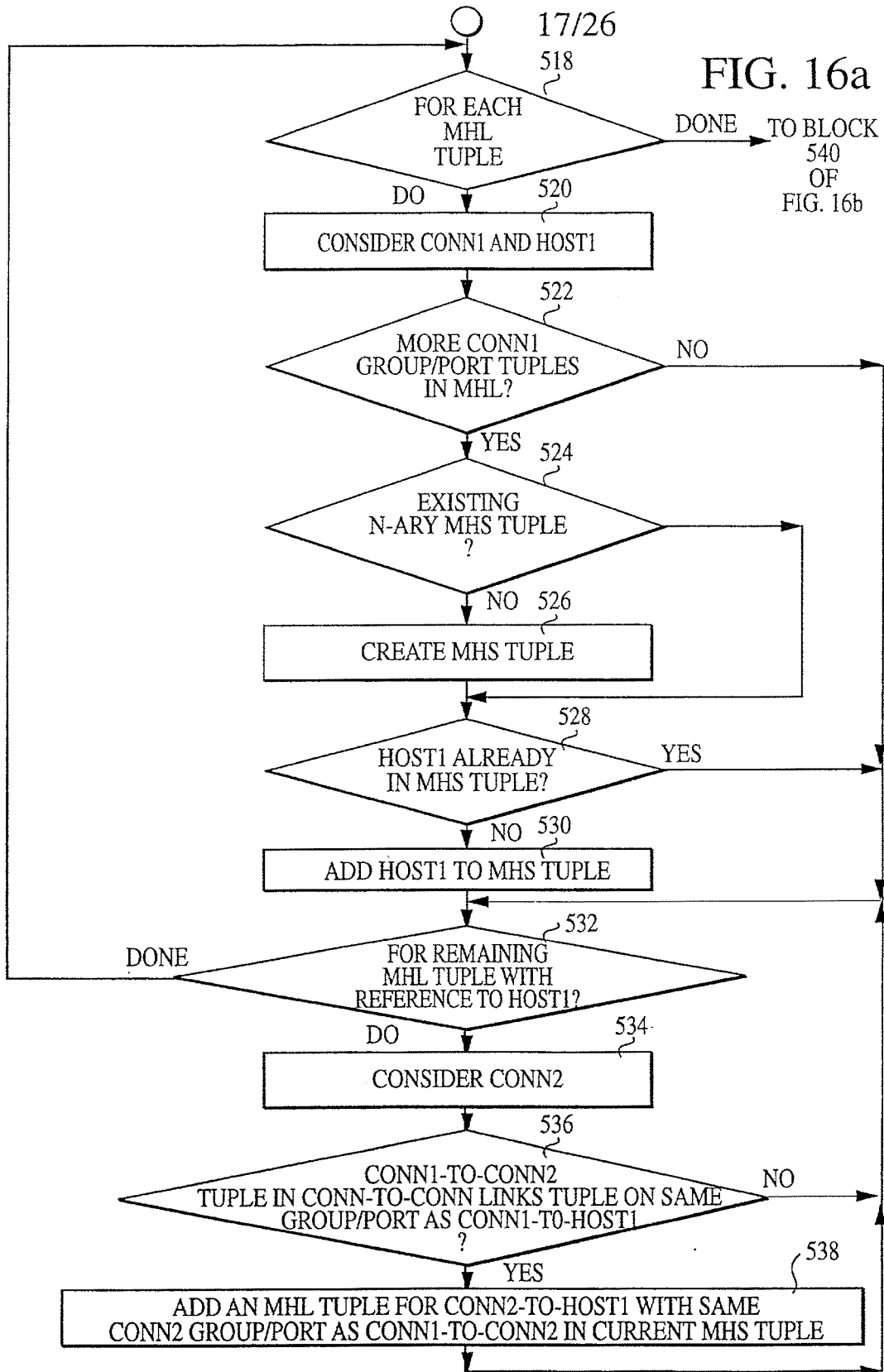
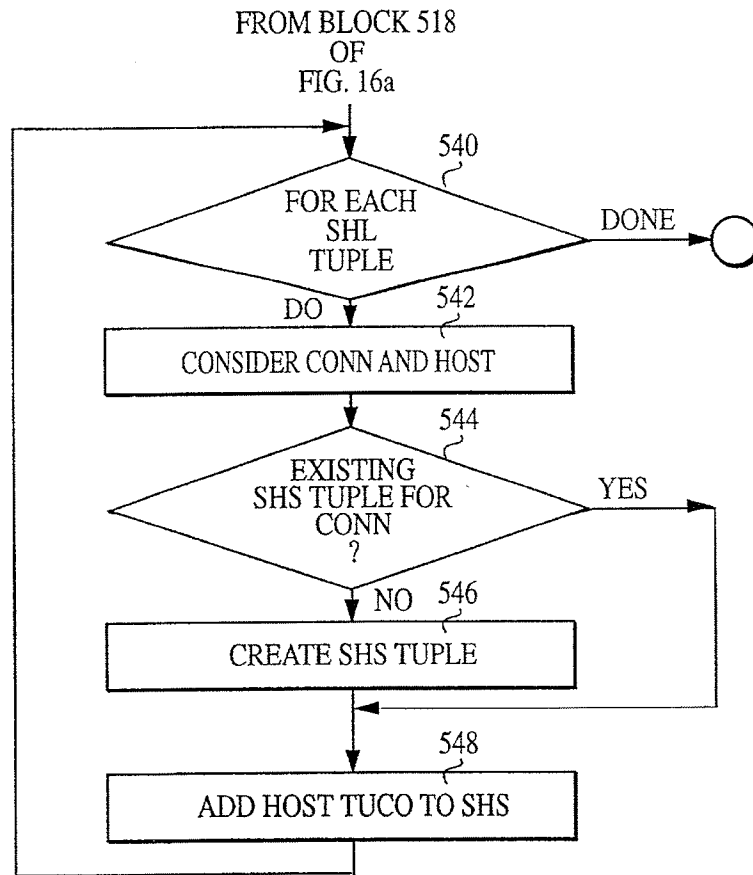
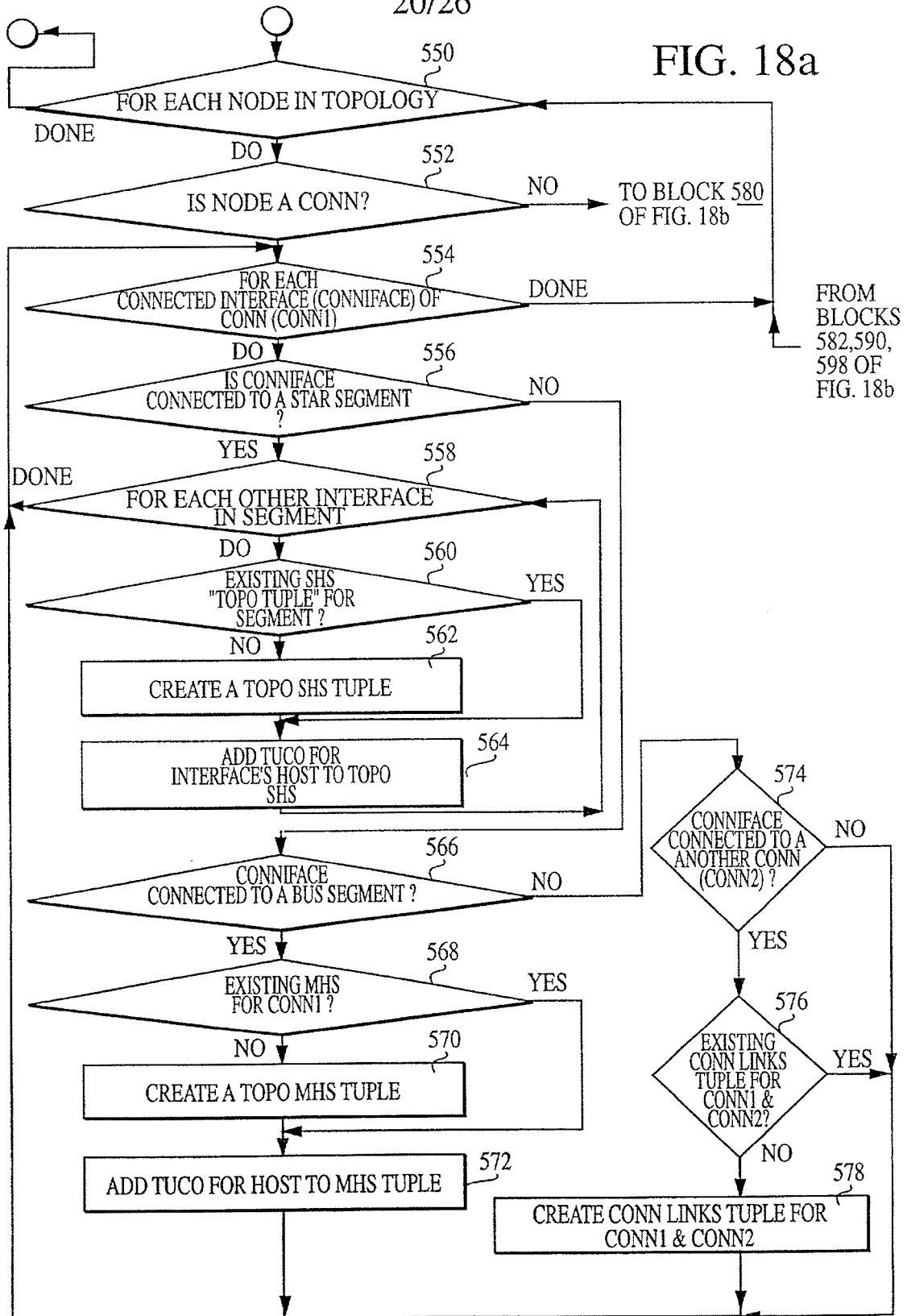


FIG. 16b



OBJECT 2160260

FIG. 18a



FROM BLOCKS 582,590, 598 OF FIG. 18b

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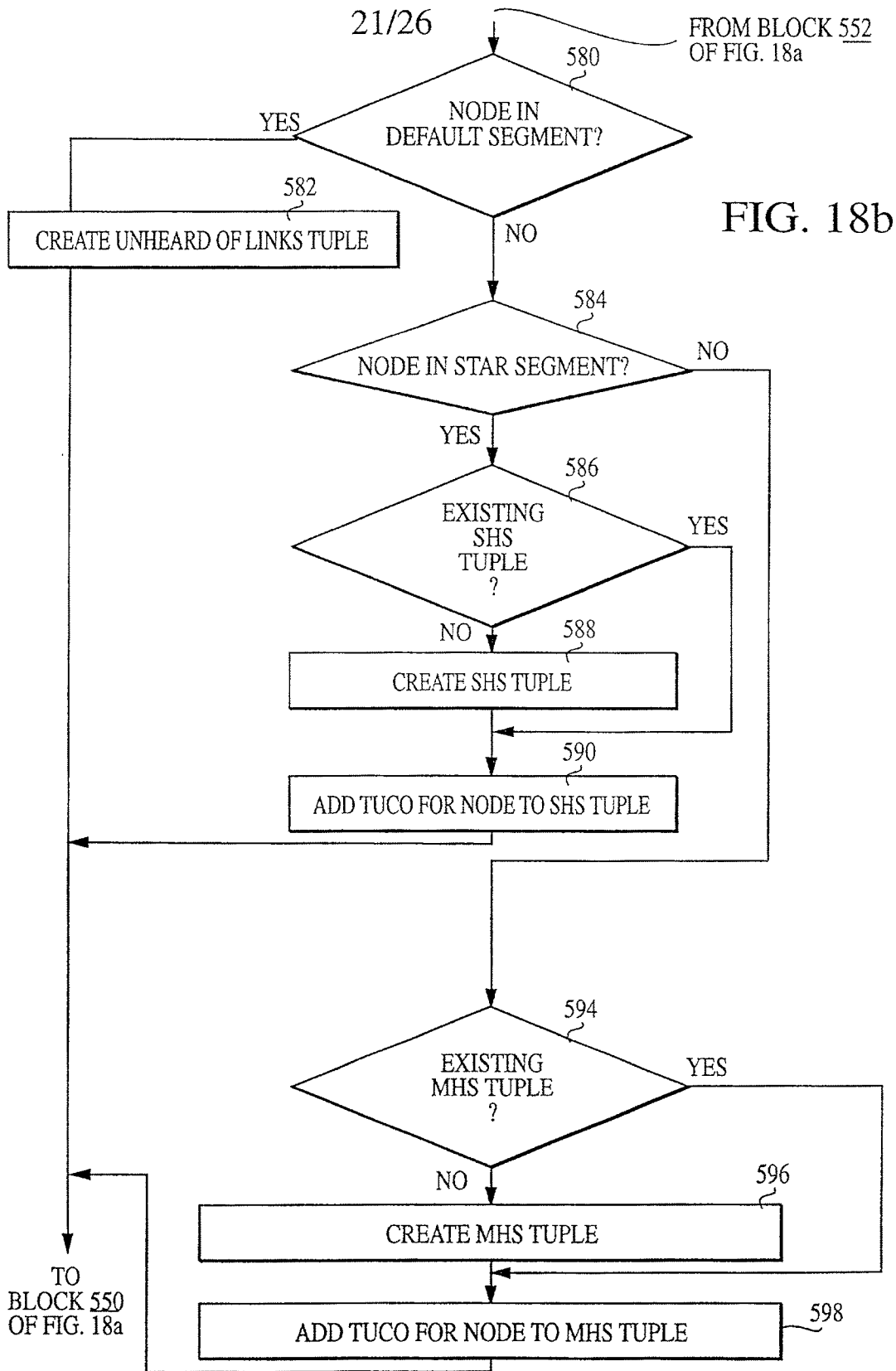
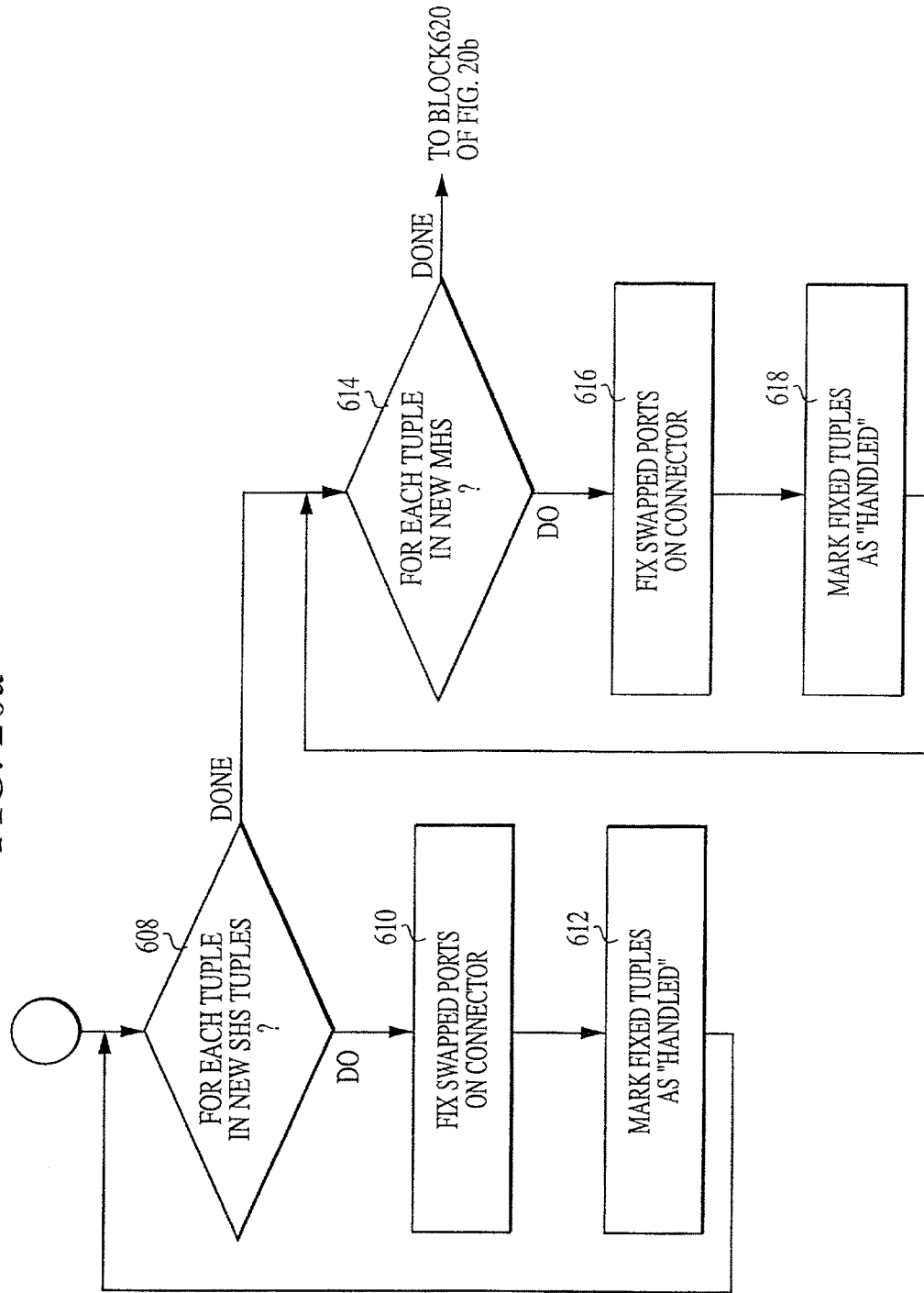


FIG. 20a



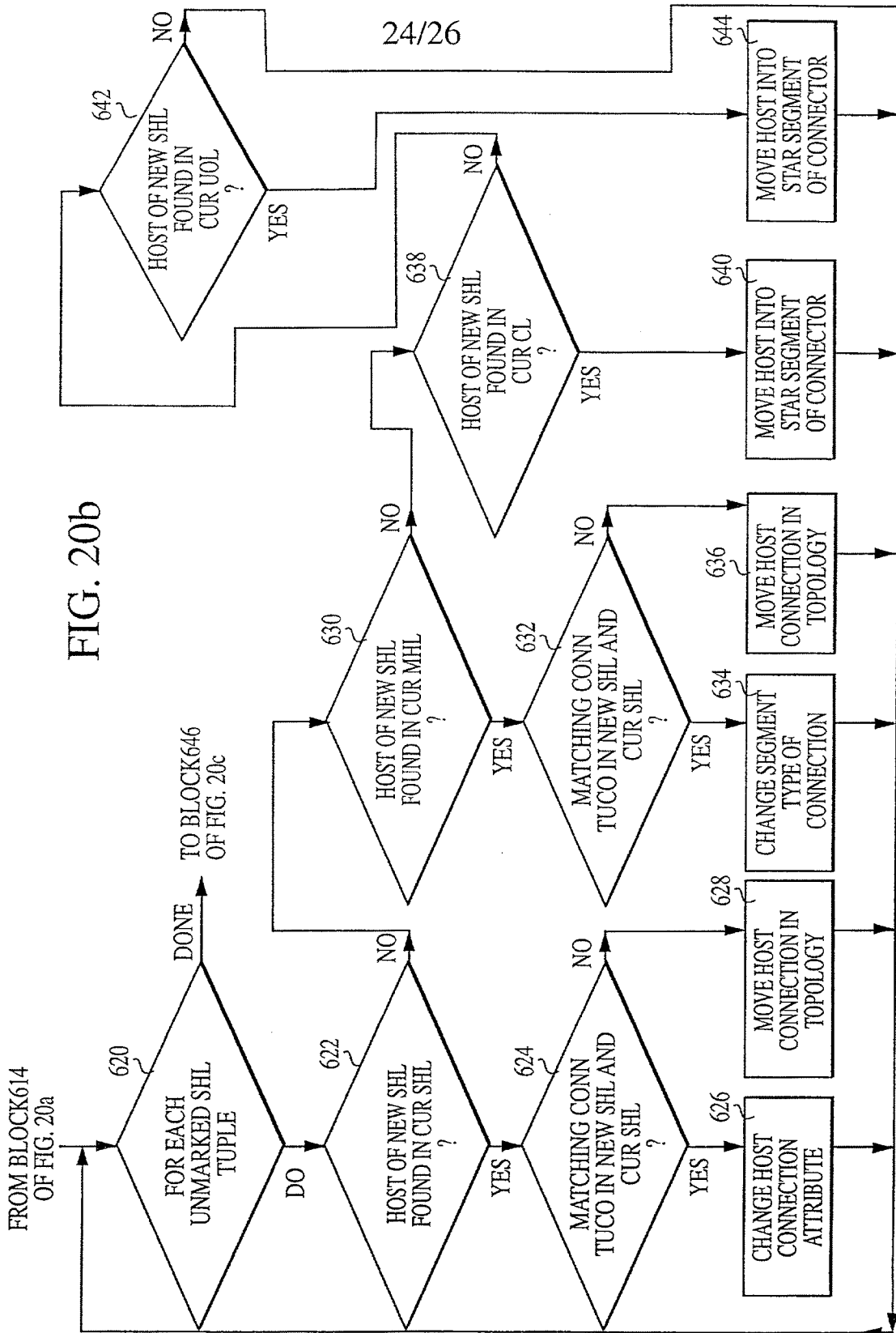


FIG. 20b

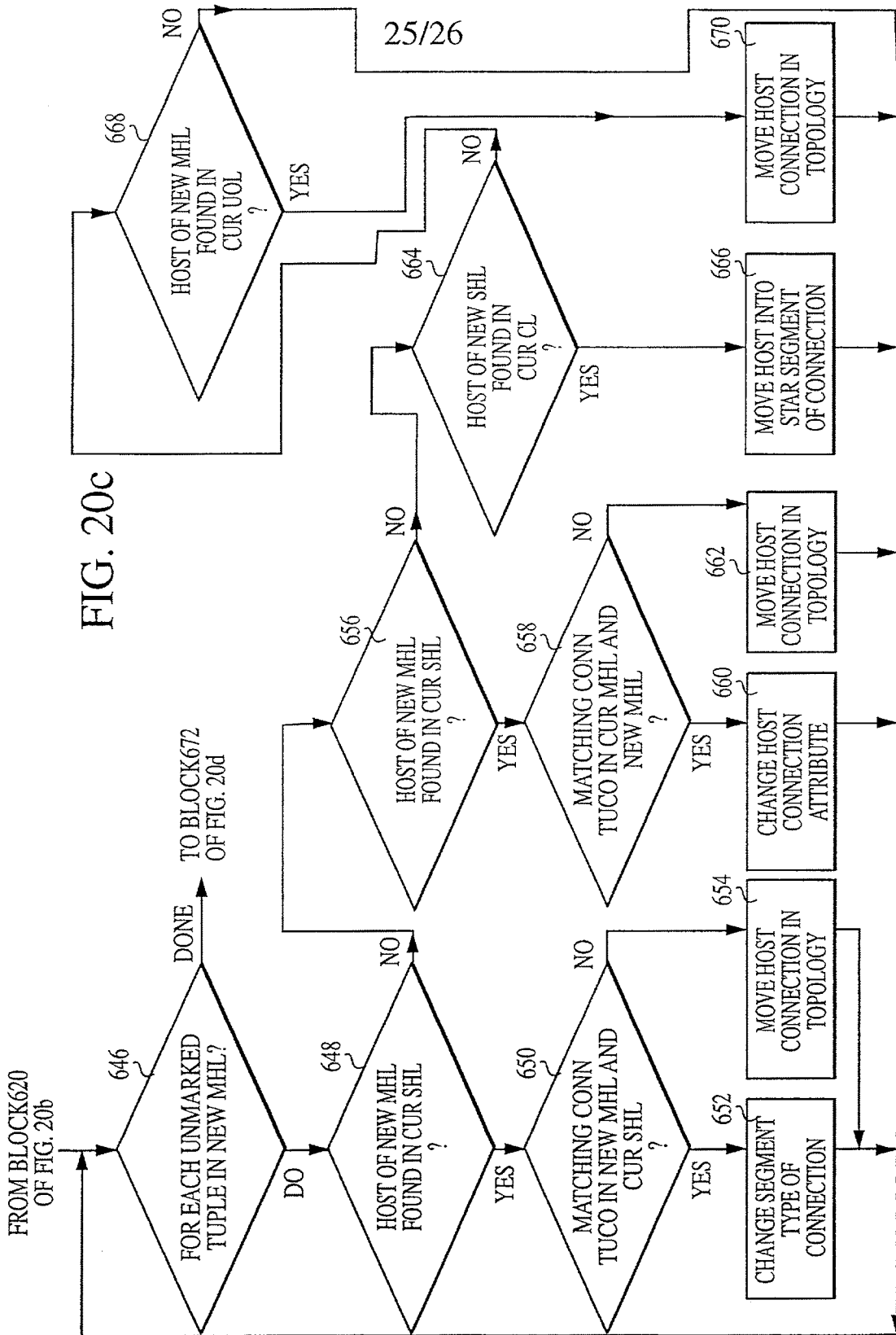
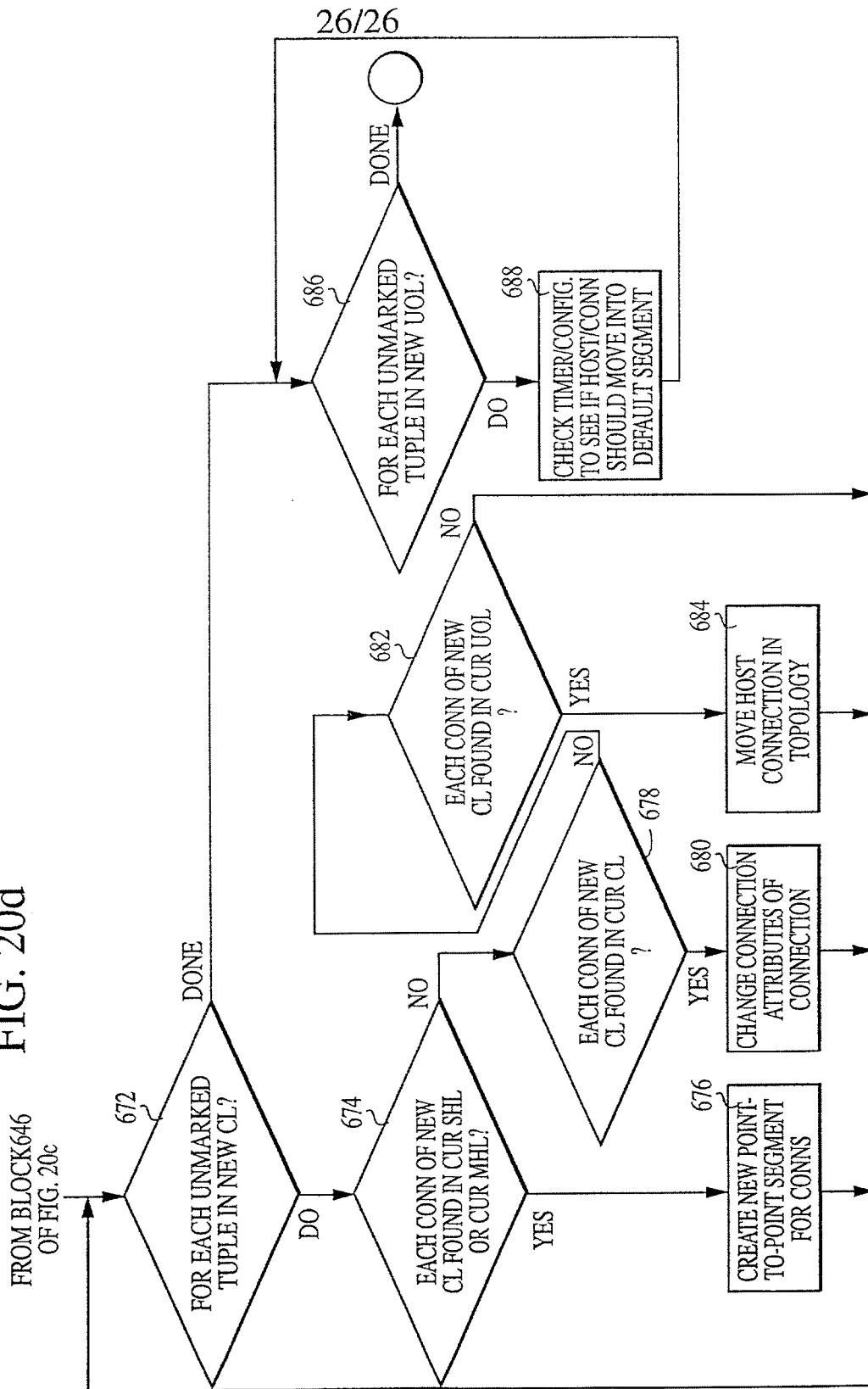


FIG. 20d



DECLARATION AND POWER OF ATTORNEY FOR PATENT APPLICATION

ATTORNEY DOCKET NO. 10008102-1

As a below named inventor, I hereby declare that:

My residence/post office address and citizenship are as stated below next to my name;

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled:

Method And System For Identifying And Processing Changes To A Network Topology

the specification of which is attached hereto unless the following box is checked:

() was filed on _____ as US Application Serial No. or PCT International Application Number _____ and was amended on _____ (if applicable).

I hereby state that I have reviewed and understood the contents of the above-identified specification, including the claims, as amended by any amendment(s) referred to above. I acknowledge the duty to disclose all information which is material to patentability as defined in 37 CFR 1.56.

Foreign Application(s) and/or Claim of Foreign Priority

I hereby claim foreign priority benefits under Title 35, United States Code Section 119 of any foreign application(s) for patent or inventor(s) certificate listed below and have also identified below any foreign application for patent or inventor(s) certificate having a filing date before that of the application on which priority is claimed:

COUNTRY	APPLICATION NUMBER	DATE FILED	PRIORITY CLAIMED UNDER 35 U.S.C. 119
N/A			YES: ___ NO: ___
			YES: ___ NO: ___

Provisional Application

I hereby claim the benefit under Title 35, United States Code Section 119(e) of any United States provisional application(s) listed below:

APPLICATION SERIAL NUMBER	FILING DATE
N/A	

U. S. Priority Claim

I hereby claim the benefit under Title 35, United States Code, Section 120 of any United States application(s) listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of Title 35, United States Code Section 112, I acknowledge the duty to disclose material information as defined in Title 37, Code of Federal Regulations, Section 1.56(a) which occurred between the filing date of the prior application and the national or PCT international filing date of this application:

APPLICATION SERIAL NUMBER	FILING DATE	STATUS (patented/pending/abandoned)
N/A		

POWER OF ATTORNEY:

As a named inventor, I hereby appoint the following attorney(s) and/or agent(s) to prosecute this application and transact all business in the Patent and Trademark Office connected therewith:

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 P.O. Box 272400
 Fort Collins, Colorado 80527-2400

Direct Telephone Calls To:

T. Grant Ritz
 (970) 898-0697

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

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Residence: 2937 Redburn Drive Ft Collins CO 80525

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Inventor's Signature:  Date: 10/31/2000

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DECLARATION AND POWER OF ATTORNEY
FOR PATENT APPLICATION (continued)

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Inventor's Signature: *Joseph R Hunt* Date: 10/31/00

Full Name of # 3 joint inventor: _____ Citizenship: _____

Residence: _____

Post Office Address: _____

Inventor's Signature _____ Date _____

Full Name of # 4 joint inventor: _____ Citizenship: _____

Residence: _____

Post Office Address: _____

Inventor's Signature _____ Date _____

Full Name of # 5 joint inventor: _____ Citizenship: _____

Residence: _____

Post Office Address: _____

Inventor's Signature _____ Date _____

Full Name of # 6 joint inventor: _____ Citizenship: _____

Residence: _____

Post Office Address: _____

Inventor's Signature _____ Date _____

Full Name of # 7 joint inventor: _____ Citizenship: _____

Residence: _____

Post Office Address: _____

Inventor's Signature _____ Date _____

Full Name of # 8 joint inventor: _____ Citizenship: _____

Residence: _____

Post Office Address: _____

Inventor's Signature _____ Date _____

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PATENT APPLICATION

ATTORNEY DOCKET NO. 10008102-1

A

IN THE U.S. PATENT AND TRADEMARK OFFICE
Patent Application Transmittal Letter

COMMISSIONER FOR PATENTS
Washington, D.C. 20231

Sir:

Transmitted herewith for filing under 37 CFR 1.53(b) is a(n): Utility () Design
 original patent application,
() continuation-in-part application

INVENTOR(S): Eric A Pulsipher et al

TITLE: Method And System For Identifying And Processing Changes To A Network Topology

Enclosed are:

- The Declaration and Power of Attorney. signed () unsigned or partially signed
- 26 sheets of drawings (one set) () Associate Power of Attorney
- () Form PTO-1449 () Information Disclosure Statement and Form PTO-1449
- () Priority document(s) () (Other) _____ (fee \$ _____)

CLAIMS AS FILED BY OTHER THAN A SMALL ENTITY				
(1) FOR	(2) NUMBER FILED	(3) NUMBER EXTRA	(4) RATE	(5) TOTALS
TOTAL CLAIMS	20 — 20	0	X \$18	\$ 0
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ANY MULTIPLE DEPENDENT CLAIMS	0		\$270	\$ 0
BASIC FEE: Design (\$320.00); Utility (\$710.00)				\$ 710
TOTAL FILING FEE				\$ 710
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TOTAL CHARGES TO DEPOSIT ACCOUNT				\$ 710

Charge \$ 710 to Deposit Account 08-2025. At any time during the pendency of this application, please charge any fees required or credit any over payment to Deposit Account 08-2025 pursuant to 37 CFR 1.25. Additionally please charge any fees to Deposit Account 08-2025 under 37 CFR 1.16, 1.17, 1.19, 1.20 and 1.21. A duplicate copy of this sheet is enclosed.

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By Laura M. Clark
Typed Name: Laura M. Clark

Respectfully submitted,

Eric A Pulsipher et al

By TGR

T. Grant Ritz

Attorney/Agent for Applicant(s)

Reg. No. 39,819

Date: Oct. 31, 2000

Telephone No.: (970) 898-0697

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Title

Method and System for Identifying and Processing Changes to a Network Topology

Field of Invention

The present invention relates generally to computer networks. More particularly, it relates to a method and system for identifying changes to a network topology and for acting upon the network based on the changes.

Background

As communications networks, such as the Internet, carry more and more traffic, efficient use of the bandwidth available in the network becomes more and more important. Switching technology was developed in order to reduce congestion and associated competition for the available bandwidth. Switching technology works by restricting traffic. Instead of broadcasting a given data packet to all parts of the network, switches are used to control data flow such that the data packet is sent only along those network segments necessary to deliver it to the target node. The smaller volume of traffic on any given segment results in few packet collisions on that segment and, thus, the smoother and faster delivery of data. A choice between alternative paths is usually possible and is typically made based upon current traffic patterns.

The intelligent routing of data packets with resultant reduction in network congestion can only be effected if the network topology is known. The topology of a network is a description of the network which includes the location of and interconnections between nodes on the network. The word "topology" refers to either the physical or logical layout of the network, including devices, and their connections in relationship to one another. Information necessary to create the topology layout can be derived from tables stored in network devices such as hubs, bridges, and switches. The information in these tables is in a constant state of flux as new entries are being added and old entries time out. Many times there simply is not enough information to determine where to place a particular device.

Switches examine each data packet that they receive, read the source addresses, and log those addresses into tables along with the switch ports on which the packets were received. If a packet is received with a target address without an entry in the switches table, the switch receiving it

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1 broadcasts that packet to each of its ports. When the switch receives a reply, it will have identified
2 where the new node lies.

3 In a large network with multiple possible paths from the switch to the target node, this table
4 can become quite large and may require a significant amount of the switch's resources to develop
5 and maintain. As an additional complication, the physical layout of devices and their connections
6 are typically in a state of constant change. Devices are continually being removed from, added to,
7 and moved to new physical locations on the network. To be effectively managed, the topology of a
8 network must be accurately and efficiently ascertained, as well as maintained.

9 Existing mapping methods have limitations that prevent them from accurately mapping
10 topological relationships. Multiple connectivity problems are one sort of difficulty encountered by
11 existing methods. For example, connectors such as routers, switches, and bridges may be
12 interconnected devices in a network. Some existing methods assume that these devices have only a
13 single connection between them. In newer devices, however, it is common for manufacturers to
14 provide multiple connections between devices to improve network efficiency and to increase
15 capacity of links between the devices. The multiple connectivity allows the devices to maintain
16 connection in case one connection fails. Methods that do not consider multiple connectivity do not
17 present a complete and accurate topological map of the network.

18 Another limitation of existing topology methods is the use of a single reference to identify a
19 device. Existing methods use a reference interface or a reference address in a set of devices to
20 orient all other devices in the same area. These methods assumed that every working device would
21 be able to identify, or "hear," this reference and identify it with a particular port of the device. With
22 newer devices, however, it is possible that the same address or reference may be heard out of
23 multiple ports of the same device. It is also possible that the address or reference may not be heard
24 from any ports, for example, if switching technology is used.

25 Still another limitation of existing mapping systems is that they require a complete copy of
26 the topological database to be stored in memory. In larger networks, the database is so large that
27 this really is not feasible, because it requires the computer to be very large and expensive.

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1 Still another difficulty with existing systems is that they focus on the minutia without
2 considering the larger mapping considerations. Whenever an individual change in the system is
3 detected, existing methods immediately act on that change, rather than taking a broader view of the
4 change in the context of other system changes. For example, a device may be removed from the
5 network temporarily and replaced with its ports reversed. In existing systems, this swapped port
6 scenario could require hundreds or thousands of changes because the reference addresses will have
7 changed for all interconnected devices.

8 Still another disadvantage of existing methods is that they use a continuous polling paradigm.
9 These methods continuously poll network addresses throughout the day and make decisions based
10 on those continuous polling results. This creates traffic on the network that slows other processes.

11 Still another limitation of existing methods is the assumption that network parts of a
12 particular layer would be physically separated from other parts. Network layer 1 may represent the
13 physical cabling of the network, layer 2 may represent the device connectivity, and layer 3 may
14 represent a higher level of abstraction, such as the groupings of devices into regions. Existing
15 methods assume that all layer 3 region groupings are self-contained, running on the same unique
16 physical networking. However, in an internet protocol (IP) network, multiple IP domains may co-
17 exist on the same lower layer networking infrastructure. It has become common for a network to
18 employ a virtual local area network (LAN) to improve security or to simplify network maintenance,
19 for example. Using virtual LANs, a system may have any number of different IP domains sharing
20 the same physical connectivity. As a result, existing methods create confusion with respect to
21 topological mapping because networks with multiple IP addresses in different subnets for the
22 infrastructure devices cannot be properly represented because they assume the physical separation
23 of connectivity for separate IP domains. Still another limitation of existing methods is that they do
24 not allow topological loops, such as port aggregation or trunking, and switch meshing.

25

26

Summary of Invention

27

28

A method and system are disclosed for mapping the topology of a network having interconnected nodes by identifying changes in the network and updating a stored network topology

1 based on the changes. The nodal connections are represented by data tuples that store information
 2 such as a host identifier, a connector interface, and a port specification for each connection. A
 3 topology database stores an existing topology of a network. A topology converter accesses the
 4 topology database and converts the existing topology into a list of current tuples. A connection
 5 calculator calculates tuples to represent connections in the new topology. The topology converter
 6 receives the new tuples, identifies changes to the topology, and updates the topology database using
 7 the new tuples. The topology converter identifies duplicate tuples that appear in both the new tuples
 8 and the existing tuples and marks the duplicate tuples to reflect that no change has occurred to these
 9 connections. The topology converter attempts to resolve swapped port conditions and searches for
 10 new singly-heard and multi-heard host link tuples in the list of existing tuples. The topology
 11 converter also searches for new conflict link tuples in the existing tuples. The topology converter
 12 updates the topology database with the new topology.

13 **Summary of Drawings**

14 Figure 1 is a drawing of a typical topological bus segment for representing the connectivity
 15 of nodes on a network.

16 Figure 2 is a drawing of a typical topological serial segment for representing the connectivity
 17 of nodes on a network.

18 Figure 3 is a drawing of a typical topological star segment for representing the connectivity
 19 of nodes on a network.

20 Figure 4 is a drawing of another typical topological star segment for representing the
 21 connectivity of nodes on a network.

22 Figure 5 is a drawing of the connectivity of an example network system.

23 Figure 6 is a drawing of the connectivity of another example network system.

24 Figure 7 is a block diagram of the system.

25 Figure 8 is a flow chart of the method of the system.

26 Figure 9 is a flow chart of the method used by the tuple manager.

27 Figure 10 is a flow chart of the method used by the connection calculator.

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1 As used herein, a node is any electronic component, such as a connector or a host, or
2 combination of electronic components with their interconnections. A connector is any network
3 device other than a host, including a switching device. A switching device is one type of connector
4 and refers to any device that controls the flow of messages on a network. Switching devices
5 include, but are not limited to, any of the following devices: repeaters, hubs, routers, bridges, and
6 switches.

7 As used herein, the term "tuple" refers to any collection of assorted data. Tuples may be
8 used to track information about network topology by storing data from network nodes. In one use,
9 tuples may include a host identifier, interface information, and a port specification for each node.
10 The port specification (also described as the group/port) may include a group number and a port
11 number, or just a port number, depending upon the manufacturer's specifications. A binary tuple
12 may include this information about two nodes as a means of showing the connectivity between them,
13 whether the nodes are connected directly or indirectly through other nodes. A "conn-to-conn"
14 tuple refers to a tuple that has connectivity data about connector nodes. A "conn-to-host" tuple
15 refers to a tuple that has connectivity data about a connector node and a host node. In one use,
16 tuples may have data about more than two nodes; that is, they may be n-ary tuples, such as those
17 used with respect to shared media connections described herein.

18 A "singly-heard host" (shh) refers to a host, such as a workstation, PC, terminal, printer,
19 other device, etc., that is connected directly to a connector, such as a switching device. A singly-
20 heard host link (shhl) refers to the link, also referred to as a segment, between a connector and an
21 shh. A "multi-heard host" (mhh) refers to hosts that are heard by a connector on the same port that
22 other hosts are heard. A multi-heard host link (mhhl) refers to the link between the connector and
23 an mhh. A link generally refers to the connection between nodes. A segment is a link that may
24 include a shared media connection.

25 Figure 1 is a drawing of a typical topological bus segment 100 for representing the
26 connectivity of nodes on a network 110. In Figure 1, first and second hosts 121, 122, as well as a
27 first port 131 of a first connector 140 are interconnected via the network 110. The bus segment

1 100 comprises the first and second hosts 121, 122 connected to the first port 131 of the first
 2 connector 140.

3 Figure 2 is a drawing of a typical topological serial segment 200 for representing the
 4 connectivity of nodes on the network 110. In Figure 2, the first host 121 comprises a second port
 5 132 on a second connector 145 which is connected via the network 110 to the first port 131 on the
 6 first connector 140. The serial segment 200 comprises the second port 132 on the second
 7 connector 145 connected to the first port 131 on the first connector 140. Figure 2 is an example of
 8 a connector-to-connector (“conn-to-conn”) relationship.

9 Figure 3 is a drawing of a typical topological star segment 301 for representing the
 10 connectivity of nodes on the network 110. In Figure 3, the first host 121 is connected to the first
 11 port 131 of the first connector 140. The star segment 301 comprises the first host 121 connected
 12 to the first port 131 of the first connector 140. Figure 3 is an example of a connector-to-host
 13 (“conn-to-host”) relationship.

14 Figure 4 is a drawing of another typical topological star segment 301 for representing the
 15 connectivity of nodes on the network 110. In addition to the connections described with respect to
 16 Figure 3, a third host 123 is connected to a third port 133 of the first connector 140 and a fourth
 17 host 124 is connected to a fourth port 134 of the first connector 140. In Figure 4, the star segment
 18 301 comprises the first host 121 connected to the first port 131 of the first connector 140, the third
 19 host 123 connected to the third port 133 of the first connector 140, and the fourth host 124
 20 connected to the fourth port 134 of the first connector 140. Thus, the star segment 301 comprises,
 21 on a given connector, at least one port, wherein one and only one host is connected to that port,
 22 and that host. In the more general case, the star segment 301 comprises, on a given connector, all
 23 ports having one and only one host connected to each port, and those connected hosts. Since the
 24 segments, or links, drawn using the topological methods of Figure 4 resemble a star, they are
 25 referred to as star segments.

26 For illustrative purposes, nodes in the figures described above and in subsequent figures are
 27 shown as individual electronic devices or ports on connectors. Also, in the figures the nodes are

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1 represented as terminals. However, they could also be workstations, personal computers, printers,
2 scanners, or any other electronic device that can be connected to networks 110.

3 Figure 5 is a drawing of the connectivity of an example network system. In Figure 5, first,
4 third, and fourth hosts 121, 123, 124 are connected via the network 110 to first, third, and fourth
5 ports 131, 133, 134 respectively, wherein the first, third, and fourth ports 131, 133, 134 are
6 located on the first connector 140.

7 The first, third and fourth hosts 121, 123, 124 are singly-heard hosts connected to separate
8 ports 131, 133, 134 of a common connector 140 – the first connector 140. The fifth and sixth
9 hosts 125, 126 are singly-heard hosts connected to the third and fourth connectors 142, 143. The
10 seventh and eighth hosts 127, 128 are multi-heard hosts connected to the same port 139 of the fifth
11 connector 144. The multi-heard hosts 127, 128 illustrate a shared media segment 180, also
12 referred to as a bus 180.

13 The second, third, fourth, and fifth connectors 141, 142, 143, 144 are interconnected and
14 illustrate a switch mesh 181. Each of the connectors in the switch mesh 181 is connected to each
15 other, either directly or indirectly, to create a fully meshed connection. In the mesh, traffic may be
16 dynamically routed to create an efficient flow.

17 Figure 5 also shows an example of a port aggregation 182, also referred to as trunking 182.
18 The first connector 140 is connected via the network 110 to the second connector 141 by two
19 direct links, each of which is connected to different ports on the connectors. One link is connected
20 to the sixth port 136 of the first connector 140 and to the seventh port of the second connector
21 137. The other link is connected to fifth port 135 of the first connector 140 and to the eighth port
22 138 of the second connector 141. In this example, two connectors illustrate the multiple
23 connectivity between nodes. Depending upon the device specifications, devices such as connectors
24 may be connected via any number of connectors. As explained herein, the system resolves multiple
25 connectivity problems by tracking port information for each connection.

26 Figure 6 is a drawing of the connectivity of a portion of a network having three connectors
27 171, 172, 173. A first host 151 is connected directly to the first port 161 of the first connector 171
28 and the second host 152 is connected to a sixth port 166 of the third connector 173. The second

1 port 162 of the first connector 171 is connected directly to the third port 163 of the second, or
2 intermediate, connector 172. The fourth port 164 of the intermediate connector 172 is connected
3 directly to the fifth port 165 of the third connector 173.

4 Figure 7 shows a block diagram of the system. Figure 8 shows a flow chart of the method
5 used by the system to retrieve and update the topology of the network. A tuple manager 300, also
6 referred to as a data miner 300, gathers 902 data from network nodes and builds 904 tuples to
7 update the current topology. The topology database "topodb" 350 stores the current topology for
8 use by the system. The "neighbor data" database 310 stores new tuple data retrieved by the tuple
9 manager 300. The connection calculator 320 processes the data in the neighbor data database 310
10 to determine the new network topology. The connection calculator 320 reduces 906 the tuple data
11 and sends it to the reduced topology relationships database 330. The topology converter 340 then
12 updates 908 the topology database 350 based on the new tuples sent to the reduced topology
13 relationships database 330 by the connection calculator 320.

14 Figure 9 shows a flow chart of one operation of the tuple manager 300, as described
15 generally by the data gathering 902 and tuple building 904 steps of the method shown in Figure 8.
16 The tuple manager 300 receives 910 a signal to gather tuple data. The tuple manager 300 then
17 retrieves 912 node information of the current topology stored in the topology database 350. This
18 information tells the tuple manager 300 which devices or nodes are believed to exist in the system
19 based on the nodes that were detected during a previous query. The tuple manager 300 then
20 queries 914 the known nodes to gather the desired information. For example, the connectors may
21 maintain forwarding tables that store connectivity data used to perform the connectors' ordinary
22 functions, such as switching. Other devices may allow the system to perform queries to gather
23 information about the flow of network traffic. This data identifies the devices heard by a connector
24 and the port on which the device was heard. The tuple manager 300 gathers this data by accessing
25 forwarding tables and other information sources for the nodes to determine such information as their
26 physical address, interface information, and the port from which they "hear" other devices. Based
27 on this information, the tuple manager 300 builds 916 tuples and stores 918 them in the "neighbor
28 data" database 310. Some nodes may have incomplete information. In this case, the partial

1 information is assembled into a tuple and may be used as a "hint" to determine its connectivity later,
2 based on other connections. The tuple manager 300 may also gather 920 additional information
3 about the network or about particular nodes as needed. For example, the connection calculator
4 320 may require additional node information and may signal the tuple manager 300 to gather that
5 information.

6 After the data is gathered and the tuples are stored in the neighbor database 310, the
7 connection calculator 320 processes the tuples to reduce them to relationships in the topology.
8 Figure 10 shows a flow chart of the process of the connection calculator 320, as shown generally in
9 the reduction step 906 of the method shown in Figure 8. The connection calculator 320 performs a
10 first weeding phase 922 to identify singly-heard hosts to distinguish them from multi-heard hosts.
11 Singly-heard hosts refer to host devices connected directly to a connector. The connection
12 calculator 320 then performs an infrastructure-building phase 924 to remove redundant connector-
13 to-connector links and to complete the details for partial tuples that are missing information. Then,
14 the connection calculator 320 performs a second weeding phase 926 to resolve conflicting reports
15 of singly-heard hosts. The connection calculator 320 then performs a noise reduction phase 928 to
16 remove redundant neighbor information for connector-to-host links. If clarification of device
17 connectivity is required, the connection calculator 320 performs a "look for" phase 930 to ask the
18 tuple manager 300 to gather additional data. The tuple data is then consolidated 932 into segment
19 and network containment relationships. The connection calculator 320 may also tag redundant
20 tuples to indicate their relevance to actual connectivity. These redundant tuples may still provide
21 hints to connectivity of other tuples. As part of the consolidation phase 932, the connection
22 calculator 320 creates new n-ary tuples (tuples having references to three or more tuocos) for shared
23 media segments.

24 Figure 11 is a flow chart of the connection calculator's first weeding process 922 for
25 distinguishing singly-heard hosts. The purpose of the first weeding process 922 is to identify the
26 direct connections between connectors and hosts; that is, those tuples having a first tuco that is a
27 connector and a second tuco that is a host. The connection calculator 320 looks through the tuple
28 list in the neighbor database 310, and for each tuple 402, the connection calculator 320 determines

1 manufactures tuples based on the list of singly-heard host link tuples identified in the first weeding
 2 phase 922. The purpose is to identify the relationship between the connectors in the extra host links
 3 tuples and the connectors directly connected to the singly-heard hosts. For each singly-heard host
 4 link 420, the connection calculator 320 processes 422 each extra host link that refers to the host.
 5 In the illustration of Figure 6, a conn-to-conn link tuple would represent the connection between the
 6 first connector 171 and the intermediate connector 172. An extra host link tuple would represent
 7 the indirect connection between the intermediate connector 172 and the first host 151. The conn-
 8 to-conn link tuple between the first connector 171 and the intermediate connector 172 is an
 9 example of an eh1Conn-to-shh1Conn tuple. If a conn-to-conn link tuple exists 424 for the extra host
 10 link connector to the singly-heard host link connector (eh1Conn-to-shh1Conn), then the connection
 11 calculator 320 updates 428 the tuple if it is incomplete. It is possible that the tuple data may be
 12 incomplete and a conn-to-conn link may not exist. In that case, a conn-to-conn tuple does not exist
 13 for the eh1Conn-to-shh1Conn, then such a tuple is created 426.

14 After processing extra host links for singly-heard host links, the connection calculator 320
 15 considers 430 each connector (referred to as conn1) in the tuples to determine the relationship
 16 between connectors. As illustrated in Figure 6, a single connector may be connected directly and
 17 indirectly to multiple other connectors. In Figure 6, the first connector 151 is connected to the
 18 intermediate connector 171 directly and also to the third connector 173 indirectly. The third
 19 connector 173 hears the first host 151 on the same part 165 that it hears the first connector 171 and
 20 the intermediate connector 172. The infrastructure building phase 924 tries to determine the
 21 relationship between other connectors heard on the same port of conn1. In a series of
 22 interconnected connectors, the connector on one end may not hear a connector on another end, but
 23 it may hear intermediate connectors, that in turn hear their own intermediate connectors. Tuples are
 24 created to represent the interconnection of conn-to-conn relationships. Based on this data, the
 25 connection calculator 320 can make inferences regarding the overall connection between
 26 connectors.

27 For every conn1, the connection calculator 320 considers 432 every other connector
 28 (conn2) to determine whether a conn1-to-conn2 tuple exists. If conn1-to-conn2 does not exist,

1 then the connection calculator 320 considers 436 every other conn-to-conn tuple containing conn2.
 2 The other connector on this tuple may be referred to as conn3. If conn2 hears conn3 on a unique
 3 port 438 and if conn1 also hears conn3 440, then the connection calculator 320 creates 442 a tuple
 4 for conn1-to-conn2 in the connector-to-connector links tuple list.

5 After processing all of the conn1 tuples, the connection calculator 320 processes 444 each
 6 conn1-to-conn2 links tuple to ensure that they have complete port data. For each incomplete tuple
 7 446, the connection calculator 320 looks 448 for a different tuple involving conn1 in the extra host
 8 links tuples on a different port. If a different tuple is found 450, then the connection calculator 320
 9 determines 452 whether conn2 also hears the host. If conn2 does hear the host, then the
 10 connection calculator 320 completes the missing port data for conn2. If conn2 does not also hear
 11 the host 452, then the connection calculator 320 continues looking 448 through different tuples
 12 involving conn1 in extra host links on different ports.

13 After attempting to complete the missing data in each of the conn-to-conn links tuples, the
 14 connection calculator 320 processes 456 each conn-to-conn links tuple. The purpose of this sub-
 15 phase is to attempt to disprove invalid conn-to-conn links. The connection calculator 320 considers
 16 458 conn1 and conn2 of each conn-to-conn links tuple. Every other connector in conn-to-conn
 17 links may be referred to as testconn. For each testconn 460, the connection calculator 320
 18 determines 462 whether the testconn hears conn1 and conn2 on different groups/ports. If testconn
 19 hears conn1 and conn2 on different ports, then the tuple is moved to extraconnlinks (ecl) 464.
 20 Otherwise, the connection calculator 320 continues processing 460 the remaining testconns.

21 Figure 13 shows a flow chart of the second weeding phase 926. The purpose of the
 22 second weeding phase 926 is to attempt to resolve conflicts involving singly-heard hosts identified in
 23 the first weeding phase 922. In the situation described herein in which more than one connector
 24 reports that a host is singly-heard, the second weeding phase 926 reviews the tuples created during
 25 the infrastructure-building phase 924 involving the connector and host in question and attempts to
 26 disprove the reported conflict. The connection calculator 320 processes 466 each
 27 singleConflictLinks (scl) tuple (sometimes referred to as the search tuple) and considers 468 conn1
 28 and host1 of the tuple. For each extra host links tuple containing host1 470, the connection

1 in which case data might not be stored in forwarding tables. In another example, a forwarding table
 2 may not have sufficient room to store all of the required information and might delete data on a
 3 FIFO basis. In the look for phase 930, the connection calculator 320 instructs the tuple manager
 4 300 to query specific nodes to retrieve the missing data. Data that was not stored in a forwarding
 5 table on the first interrogation may be present on a subsequent query. For each mhh1 tuple 500, the
 6 connection calculator 320 considers 502 conn1 and host1. If the conn1 group/port is already in an
 7 "alreadyDidLookfors" list, then a list is created 508 for all connectors in conn-to-conn links that are
 8 heard by conn1 on the same group/port as host1. For each connector (conn2) in the list 510, the
 9 connection calculator 320 determines 512 whether there is a conn2-to-host1 tuple in the mhh1
 10 tuples. If there is not such a tuple, then the connection calculator 320 initiates a look-for for conn2-
 11 to-host1 via the tuple manager 300. When each connector in the list has been processed 510, the
 12 conn1 group/port tuco is added 516 to an alreadyDidLookfors list. As an additional portion of the
 13 look for phase 930 (not shown in figures) the system may ask a user to verify or clarify information
 14 about connectivity. For example, the system may show the user the perceived connectivity or the
 15 unresolved connectivity issues and request the user to add information as appropriate.

16 The connection calculator 330 process described above collects the tuple information from
 17 the tuple manager 300, builds tuples new tuples and removes redundant or unnecessary tuples to
 18 produce the new topology. This topology may have incomplete tuples possibly resulting from
 19 extraneous information that the connection calculator 330 could not disprove. To refine the new
 20 topology, the connection calculator 330 can request the tuple manager 300 to obtain additional
 21 information about particular nodes or it may also request a user to refine the topology by adding or
 22 removing tuples. Using the process of the connection calculator 330, tuples marked as non-
 23 essential may be removed from the new topology to save space and to simplify the topology. The
 24 connection calculator 330 is not confused by multiple connectivity situations such as port
 25 aggregation 182 or switch meshing 181 as shown in Figure 5, because the tuples represent point-
 26 to-point, or neighbor-to-neighbor, connectivity showing each connection in the network. This
 27 point-to-point connectivity concept also helps enable the system to avoid difficulties that occur in

1 the topology converter 340 determines 556 whether the conniface is connected to a star segment.
 2 If it is connected to a star segment, then for every other interface in the segment 558, the topology
 3 converter 340 determines 560 whether there is an existing shs tuple, referred to as the "topo tuple"
 4 for the segment. If there is no such tuple, then the topology converter 340 creates 562 a topo shs
 5 tuple. The tuco for the interface's host-to-topo shs is then added 564 to the topo shs tuple.

6 If the connector node is not connected to a star segment 556 and is connected to a bus
 7 segment 566, the topology converter 340 determines 568 whether there is an existing mhs tuple for
 8 conn1. If there is not an existing mhs tuple for conn1, then a topo mhs tuple is created 570. A tuco
 9 is added 572 for the host to the mhs tuple.

10 If the connector node is not connected to either a star segment 556 or to a bus segment
 11 566, then the topology converter knows that it is connected to another connector (conn2). If such
 12 a connector does not already have an existing connLinks tuple for conn1 and conn2 576, then a
 13 connLinks tuple is created 578. After processing the bus segment, star segment, and conn-to-conn
 14 segment, for each conniface 554, the topology converter 340 proceeds to the next node 550.

15 Figure 18b shows a continuation of the flow chart of Figure 18a showing the steps of the
 16 method when the topology converter 340 determines that the node is not a connector 552. If the
 17 node is in the default segment, then an "unheardOfLinks" tuple is created 582 and the topology
 18 converter proceeds to the next node 550. If the node is not in the default segment 580, then the
 19 topology converter 340 determines whether the node is in a star segment 584. If the node is in a
 20 star segment, then if there is not already an shs tuple, the topology converter 340 creates 588 an shs
 21 tuple. The tuco for the node is then added 590 to the shs tuple, and the topology converter 340
 22 proceeds to the next node 550.

23 If the node is not in a star segment, then the topology converter 340 knows that it is in the
 24 bus segment. If there is not already an mhs tuple for the node, 594, then the topology converter
 25 340 creates 596 an mhs tuple. The tuco for the node is then added 598 to the mhs tuple, and the
 26 topology converter proceeds to the next node 550.

27 Figure 19 shows a flow chart for the discard duplicates phase 936 of the topology
 28 converter 340. For each tuple in the new tuples (nt) 600, the topology converter looks for 602 an

1 exact match in the current tuples stored in the topodb. If an exact match is found 604, then the new
 2 tuple is marked 606 as “no change” indicating that this is an identical tuple.

3 Figures 20a-d show a flow chart for the identify different tuples phase 938. The system
 4 looks through each tuple in the new SinglyHeardSegments (newSHS) tuple list 608 and tries to
 5 identify and fix 610 swapped ports on connectors. Swapped ports are identified by considering
 6 those segment tuples in both the new topology and the existing topology that differ only by the port
 7 specification in the tuco. Each tuple that is fixed as a swapped port is marked 612 as “handled.”
 8 The system also looks through each tuple in the new multiHeardSegments tuple list (newMHS) 614
 9 and tries to identify and fix 616 swapped ports on connectors. Each tuple that is fixed as a
 10 swapped port is marked 618 as “handled.”

11 The system then processes 620 each unmarked tuple in the newSHL tuples. Four cases
 12 are possible for the host of the newSHL tuples. The host of the newSHL can be found in the
 13 current singlyHeardLinks (curSHL) 622, the current multiHeardLinks (curMHL) 630, the current
 14 connLinks (curCL) 638, or the current UnheardOfLinks (curUOL) 642. If the host of a newSHL
 15 tuple is found 622 in the current SinglyHeardLinks (curSHL) tuples, then the system determines 624
 16 if there is a matching connector tuco between the newSHL tuples and the curSHL tuples. If there is
 17 a matching tuco, then the system changes 626 the host connection attribute. If there is not a
 18 matching tuco, then the host connection is moved 628 in the topology.

19 If the host is found in the curMHL tuples 630, then the system determines 632 whether
 20 there is a matching connector tuco between the newSHL tuples and the curSHL tuples. If there is a
 21 matching connector, then the segment type of connection is changed 634. If there is not a matching
 22 connector, then the host connection is moved 636 in the topology. If the host is found in the curCL
 23 tuples 638, then the host is moved 640 into a star segment of the connector. If it is found in the
 24 curUOL 642, then the host is moved 644 into the star segment of the connector.

25 Figure 20c shows another stage of the processing undertaken during the identify different
 26 tuples phase 938. For each unmarked tuple in the new multiHeardLinks tuples (newMHL) 946,
 27 four cases are possible for the host of the newMHL. The host of the newMHL may be found in the
 28 curSHL 648, the curMHL 656, the curCL 664, or the curUOL 668. If the host is found in the

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1 curSHL 648, then the system determines 650 whether there is a matching connector tuco between
2 the newMHL and the curMHL. If there is a matching tuco, then the segment type of connection is
3 changed 652. If there is not a matching tuco, then the host connection is moved 654 in the
4 topology.

5 If the host is found in the curMHL tuples 656, then the system determines 658 whether
6 there is a matching connector tuco in both the curMHL tuples and the newMHL tuples. If there is a
7 matching connector tuco, then the host connection attribute is changed 660. If there is not a
8 matching tuco, then the host connection is moved 662 in the topology. If the host is found in the
9 curCL tuples 664, then the host is moved into a bus segment of a connector. If the host is found in
10 the curUOL tuples 668, then the host connection is moved 670 in the topology.

11 Figure 20d shows another portion of the identify different tuples phase 938. For each
12 unmarked tuple in the newCL tuples 672, there are three possibilities for the connector. The
13 connector of the unmarked tuple in newCL can be found in the curSHL or curMHL 674, in the
14 curCL 678, or in the curUOL 682. If each connector is found in the curSHL or curMHL list 674,
15 then the system creates 676 a new point-to-point segment for the connectors. If the connectors are
16 found in the curCL 678, then the connection attributes of the connectors are changed 680. If each
17 connector is found in the curUOL tuples 682, then the host connection is moved 684 in the
18 topology.

19 Another part of the identify different tuples phase 938 is shown in blocks 686 and 688 of
20 Figure 20d. For each unmarked tuple in the newUOL tuples 686, the system checks 688 the
21 timer/configuration to determine whether the host/conn should move into the default segment from
22 its current segment.

23 An advantage of the system is that it may be schedulable. The system may map network
24 topology continuously, as done by existing systems, or it may be scheduled to run only at certain
25 intervals, as desired by the user. A further advantage of the system is that it is capable of
26 processing multiple connections between the same devices and of processing connection meshes,
27 because it tracks each nodal connection independently, without limitations on the types of
28 connections that are permitted to exist.

1 Although the present invention has been described with respect to particular embodiments
2 thereof, variations are possible. The present invention may be embodied in specific forms without
3 departing from the essential spirit or attributes thereof. It is desired that the embodiments described
4 herein be considered in all respects illustrative and not restrictive and that reference be made to the
5 appended claims for determining the scope of the invention.

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1 **Claims**

2 1. In a network having interconnected nodes with data tuples that represent nodal
3 connections, a method for mapping a network topology by identifying changes between an existing
4 topology and a new topology, the method comprising:

5 converting an existing topology into a list of existing tuples that represent existing nodal
6 connections;

7 receiving new tuples that represent new nodal connections; and

8 comparing the list of existing tuples with the new tuples to identify changes to the topology.

9 2. The method of claim 1, further comprising updating a topology database with a new
10 topology.

11 3. The method of claim 1, further comprising taking action on the changes to the
12 topology.

13 4. The method of claim 1, wherein the tuples include information about a host
14 identifier, a connector interface, and a port specification.

15 5. The method of claim 1, wherein the step of comparing comprises identifying
16 duplicate tuples that appear both in the list of existing tuples and in the new tuples, and maintaining a
17 current status of the topology for these tuples.

18 6. The method of claim 1, wherein the step of comparing comprises identifying a
19 swapped port condition on a connector.

20 7. The method of claim 1, wherein the step of comparing comprises searching for a
21 host of a new singly-heard host link tuple or a new multi-heard host link tuple in the list of existing
22 tuples.

23 8. A system for mapping a network topology by identifying changes between an
24 existing topology and a new topology, based on changes to data tuples that represent nodal
25 connections comprising:

26 a topology database that stores an existing topology of a network; and

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1 a topology converter connected to the topology database that receives new tuples that
2 represent new nodal connections; and compares the new tuples with the existing topology to identify
3 changes in the network.

4 9. The system of claim 8, wherein the topology converter converts the existing
5 topology into a list of existing tuples that represent existing nodal connections.

6 10. The system of claim 8, wherein the topology converter updates the topology
7 database with a new topology based on the new tuples.

8 11. The system of claim 8, wherein the topology converter attempts to identify swapped
9 ports on connectors.

10 12. The system of claim 8, wherein the topology converter identifies duplicate tuples
11 that appear both in the list of existing tuples and in the new tuples, and maintains a current status of
12 the topology for these tuples.

13 13. The system of claim 8, wherein the topology converter searches for a host of a new
14 singly-heard host link tuple or a new multi-heard host link tuple in the list of existing tuples.

15 14. The system of claim 8, wherein the topology converter searches for a connector of
16 a new conflict links tuple in the list of existing tuples.

17 15. A computer-readable medium having computer-executable instructions for
18 performing a method for mapping a network topology by identifying changes between an existing
19 topology and a new topology in a network having a interconnected nodes, the method comprising:

20 converting an existing topology into a list of existing tuples that represent existing nodal
21 connections;

22 receiving new tuples that represent new nodal connections;

23 comparing the list of existing tuples with the new tuples to identify changes to the topology;

24 and

25 updating a topology database with a new topology.

26 16. The method of claim 15, wherein a topology converter receives the new tuples from
27 a connection calculator that calculates connections between nodes.

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- 1 17. The method of claim 15, wherein the step of comparing comprises identifying
2 duplicate tuples that appear both in the list of existing tuples and in the new tuples, and maintaining a
3 current status of the topology for these tuples.
- 4 18. The method of claim 15, wherein the step of comparing comprises identifying a
5 swapped port condition on a connector.
- 6 19. The method of claim 15, wherein the step of comparing comprises searching for a
7 host of a new singly-heard host link tuple or a new multi-heard host link tuple in the list of existing
8 tuples.
- 9 20. The method of claim 15, wherein the step of comparing comprises searching for a
10 connector of a new conflict links tuple in the list of existing tuples.

1 **Abstract**

2 A method and system are disclosed for mapping the topology of a network having
3 interconnected nodes by identifying changes in the network and updating a stored network topology
4 based on the changes. The nodal connections are represented by data tuples that store information
5 such as a host identifier, a connector interface, and a port specification for each connection. A
6 topology database stores an existing topology of a network. A topology converter accesses the
7 topology database and converts the existing topology into a list of current tuples. A connection
8 calculator calculates tuples to represent connections in the new topology. The topology converter
9 receives the new tuples, identifies changes to the topology, and updates the topology database using
10 the new tuples. The topology converter identifies duplicate tuples that appear in both the new tuples
11 and the existing tuples and marks the duplicate tuples to reflect that no change has occurred to these
12 connections. The topology converter attempts to resolve swapped port conditions and searches for
13 new singly-heard and multi-heard host link tuples in the list of existing tuples. The topology
14 converter also searches for new conflict link tuples in the existing tuples. The topology converter
15 updates the topology database with the new topology.

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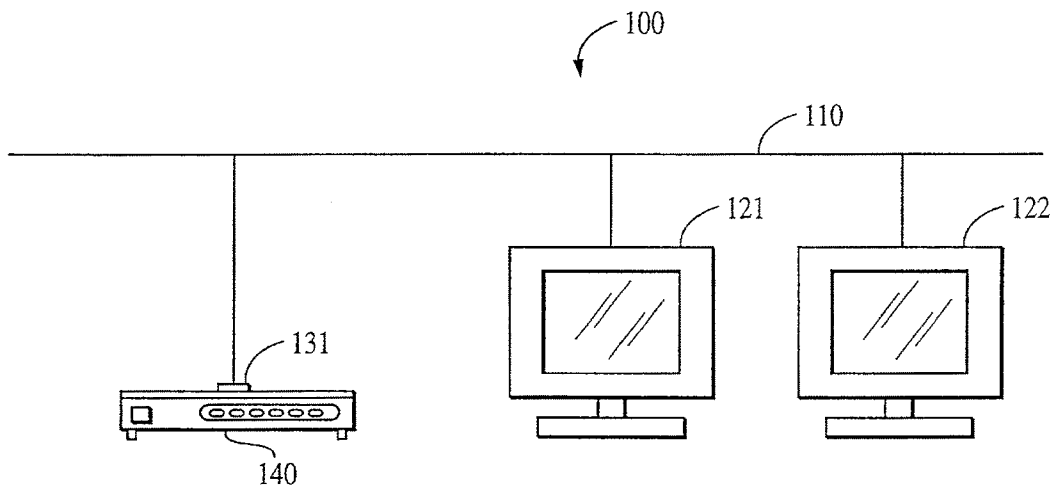


FIG. 1

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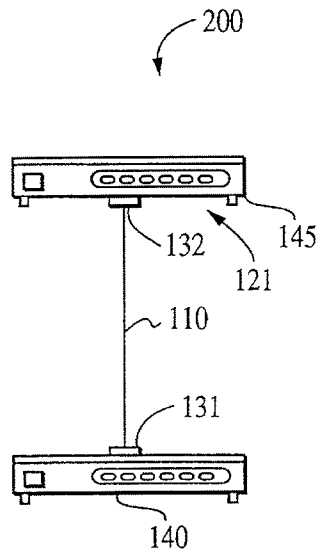


FIG. 2

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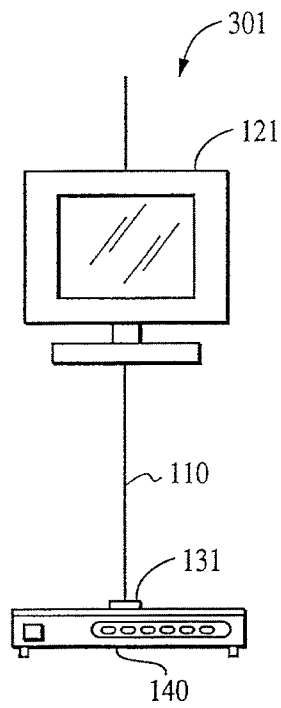


FIG. 3

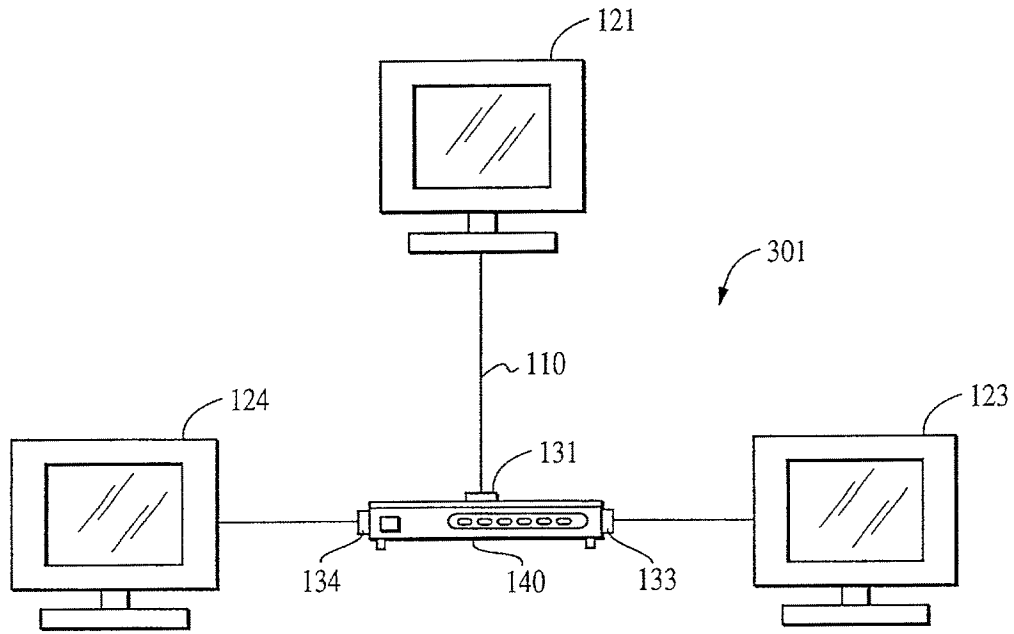
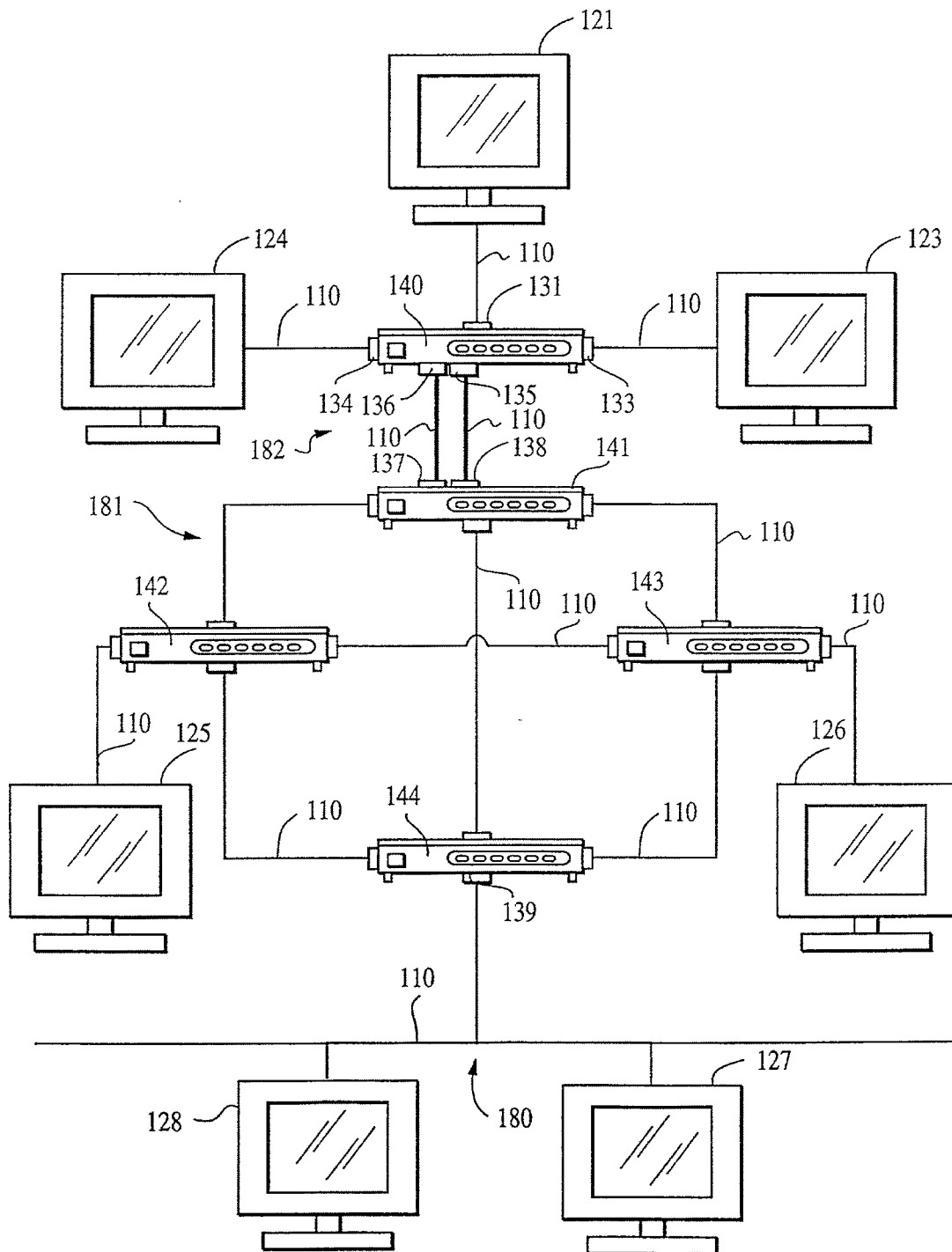


FIG. 4

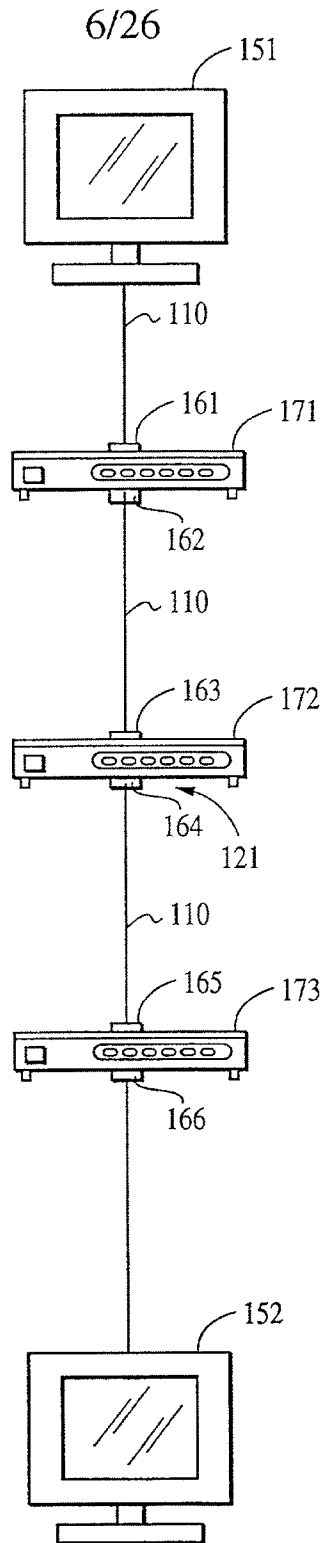
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FIG. 5



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FIG. 6



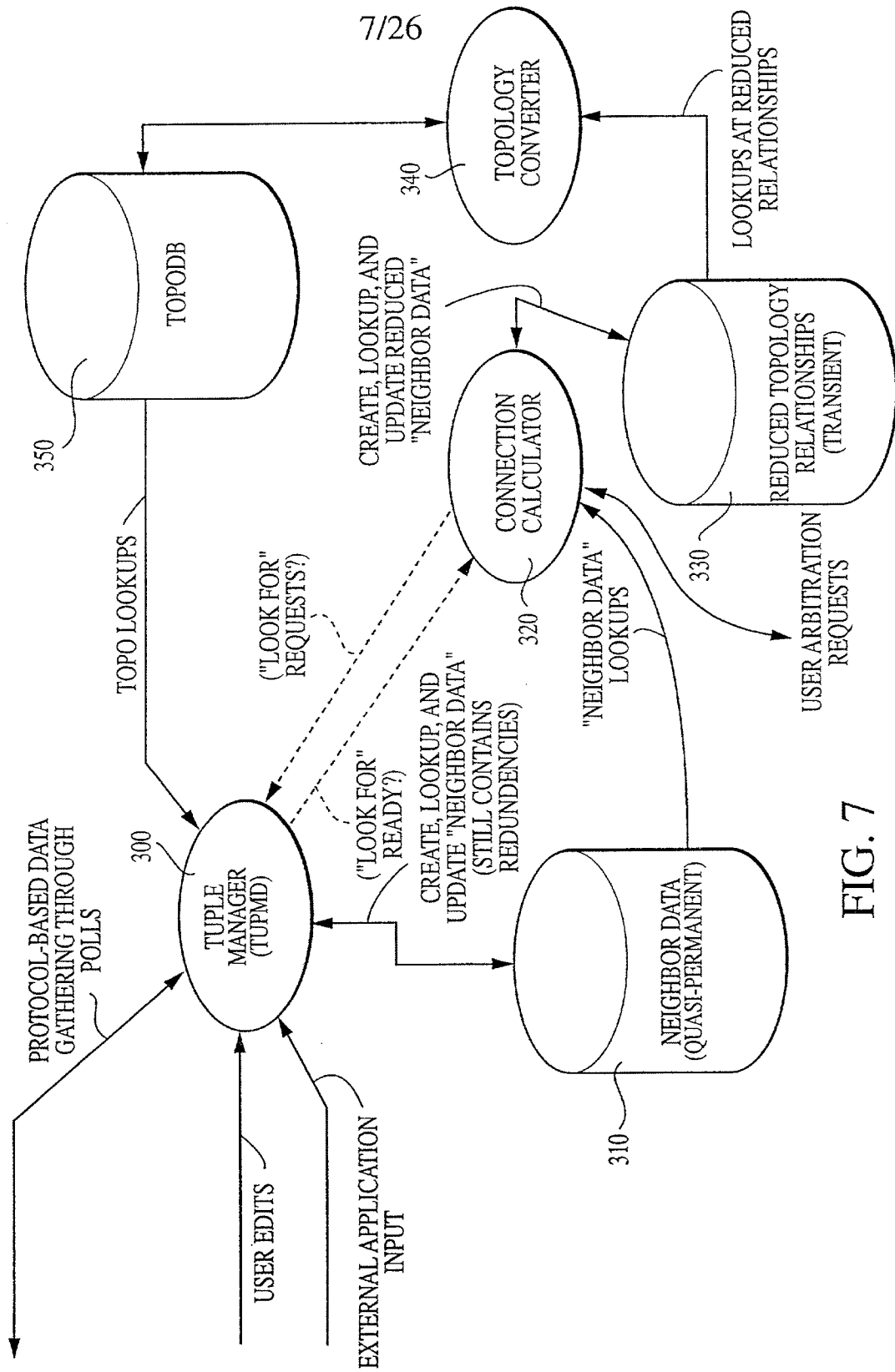
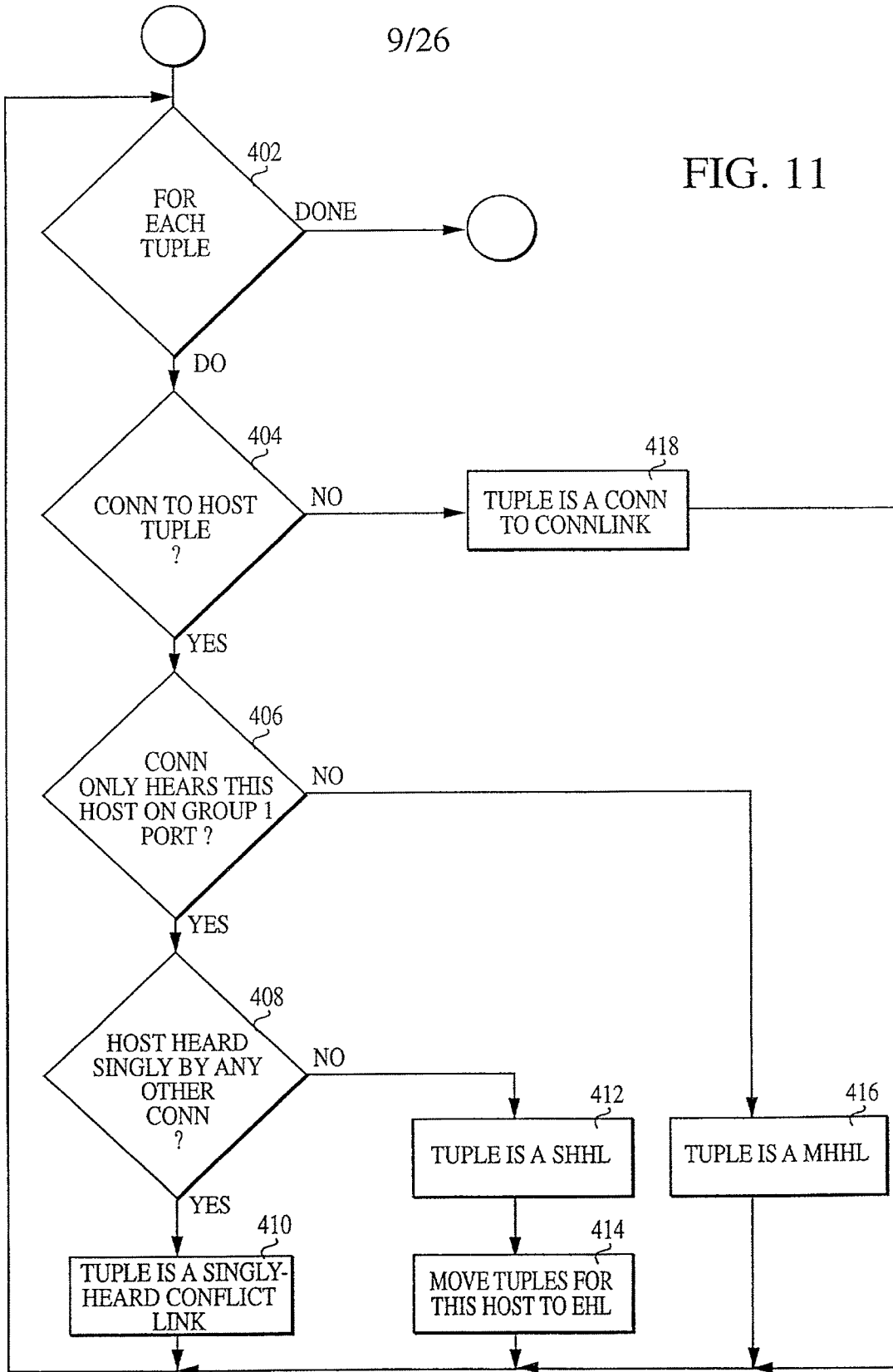


FIG. 7

9/26

FIG. 11



OFFICE OF THE ATTORNEY GENERAL

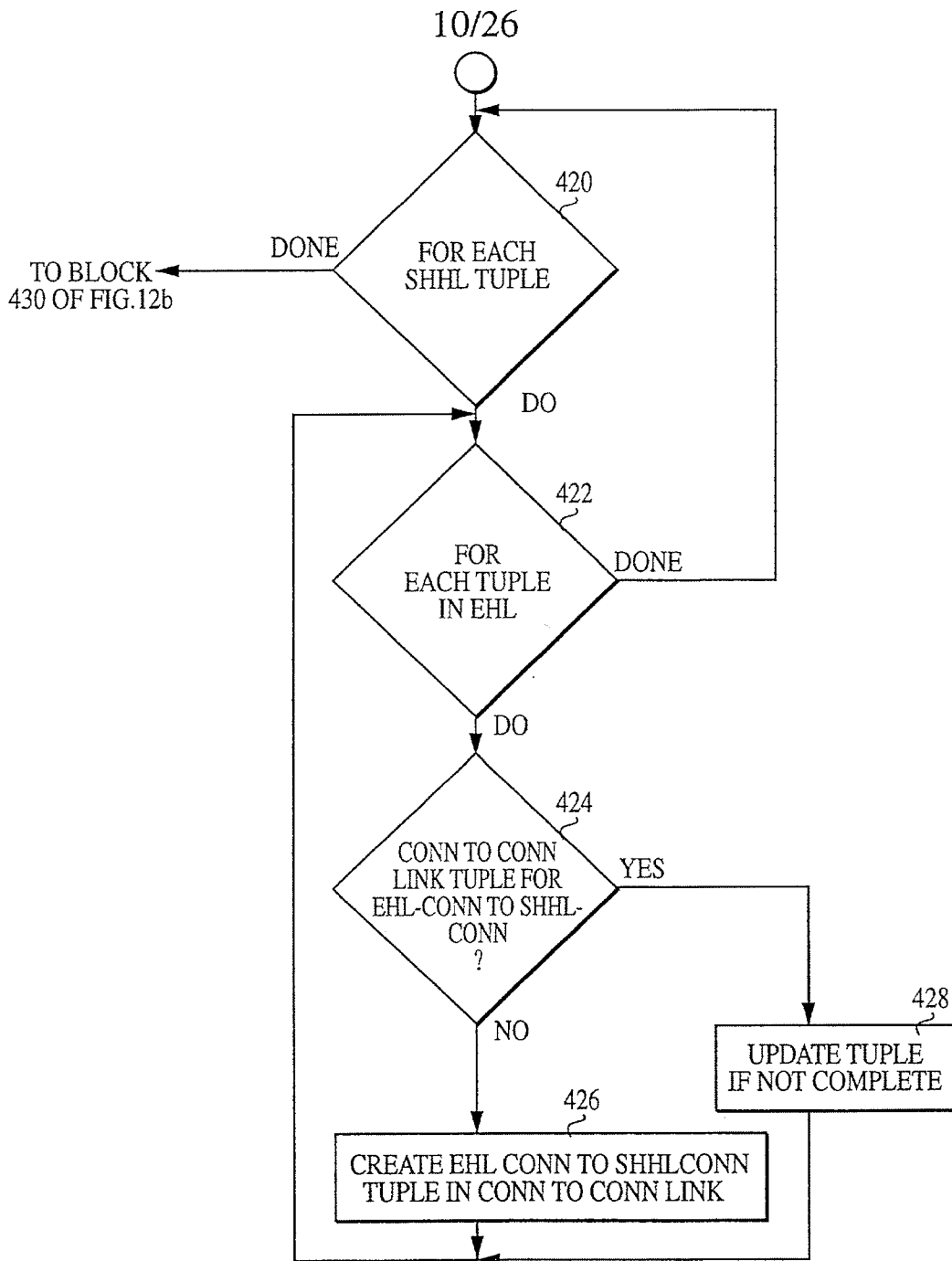


FIG. 12a

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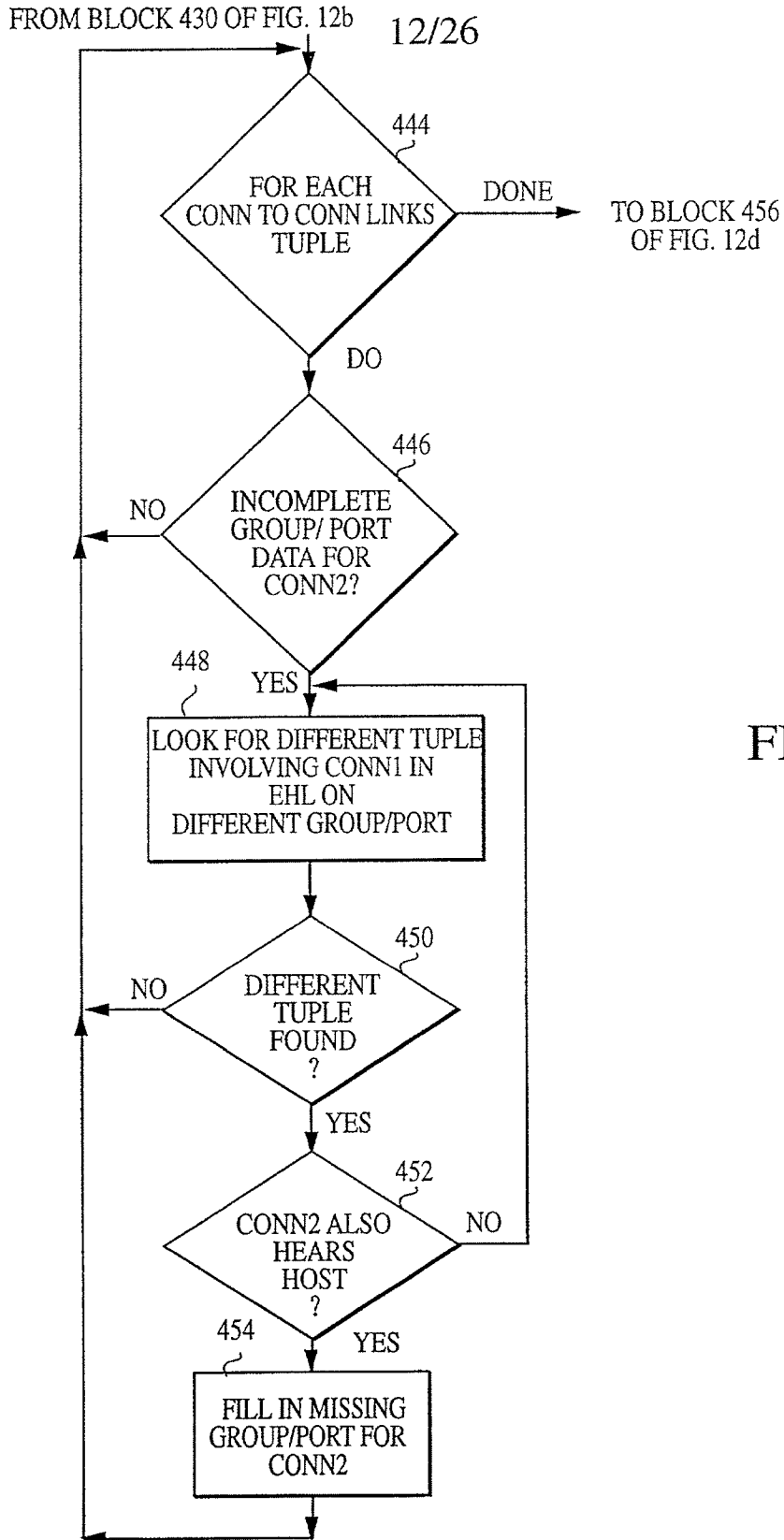


FIG. 12c

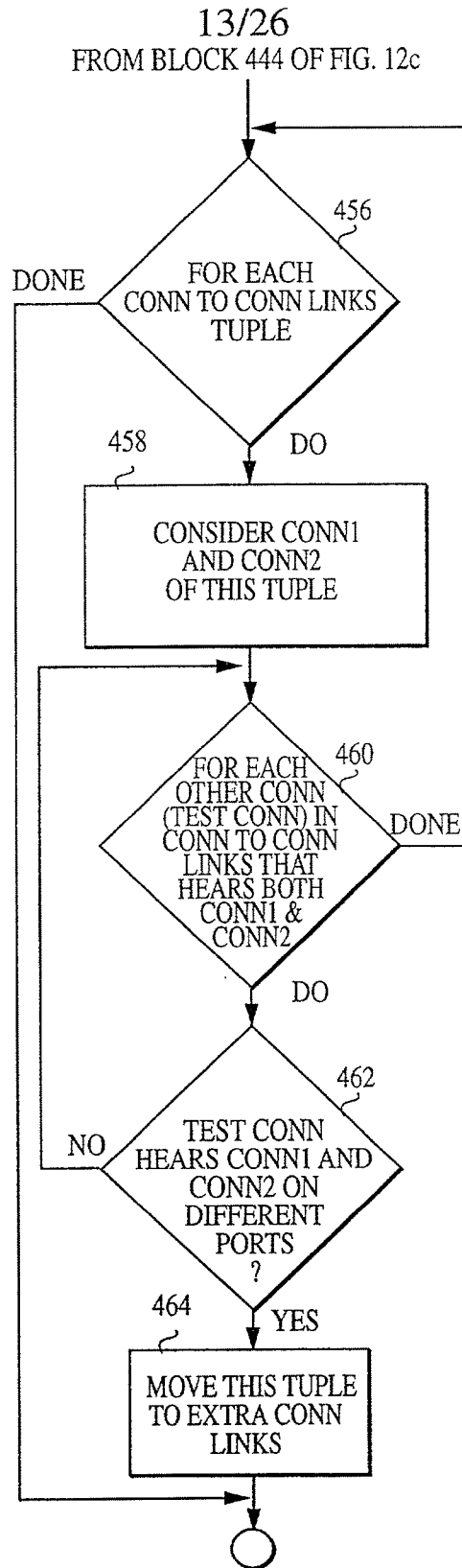


FIG. 12d

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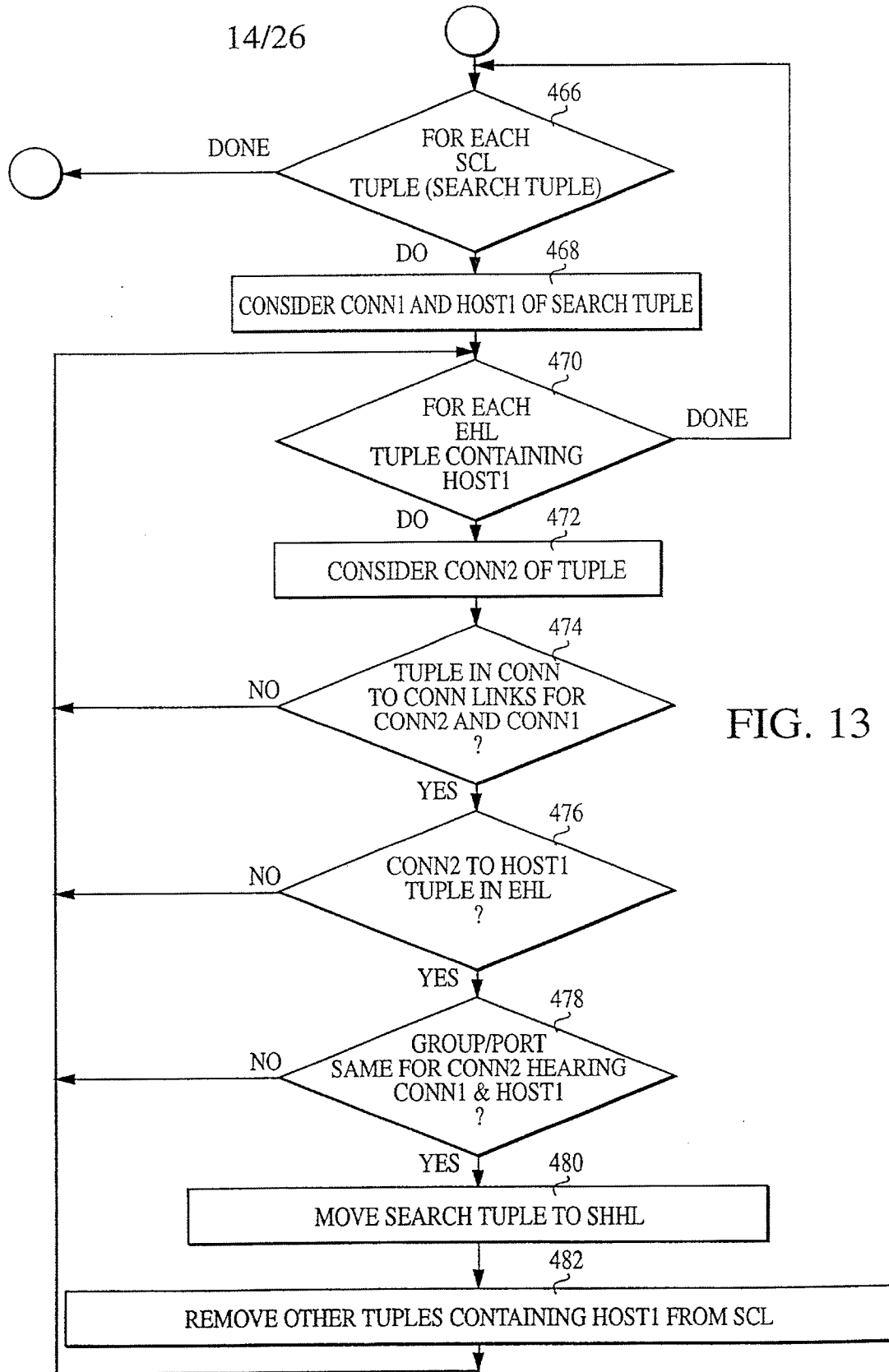


FIG. 13

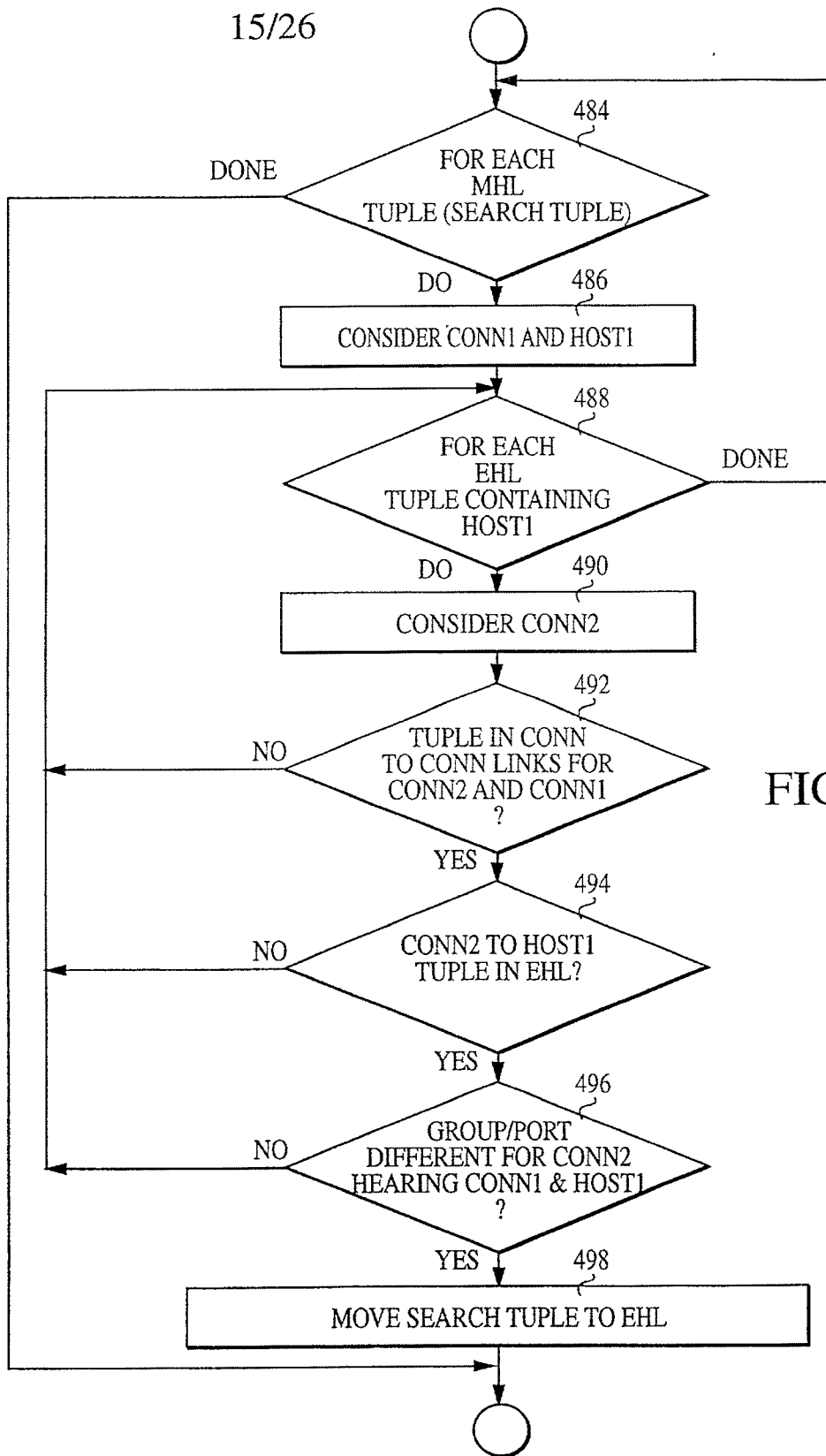
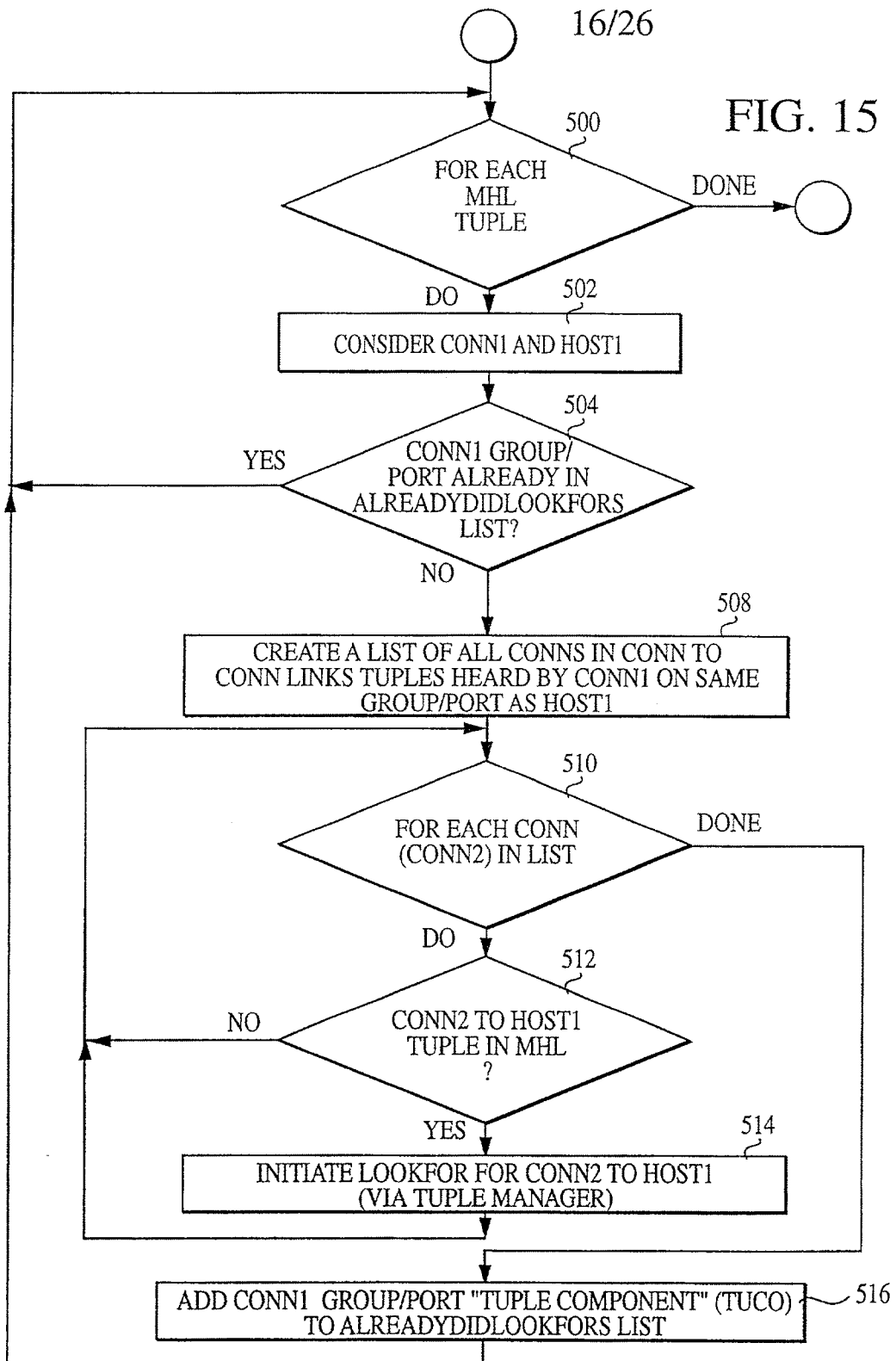


FIG. 14

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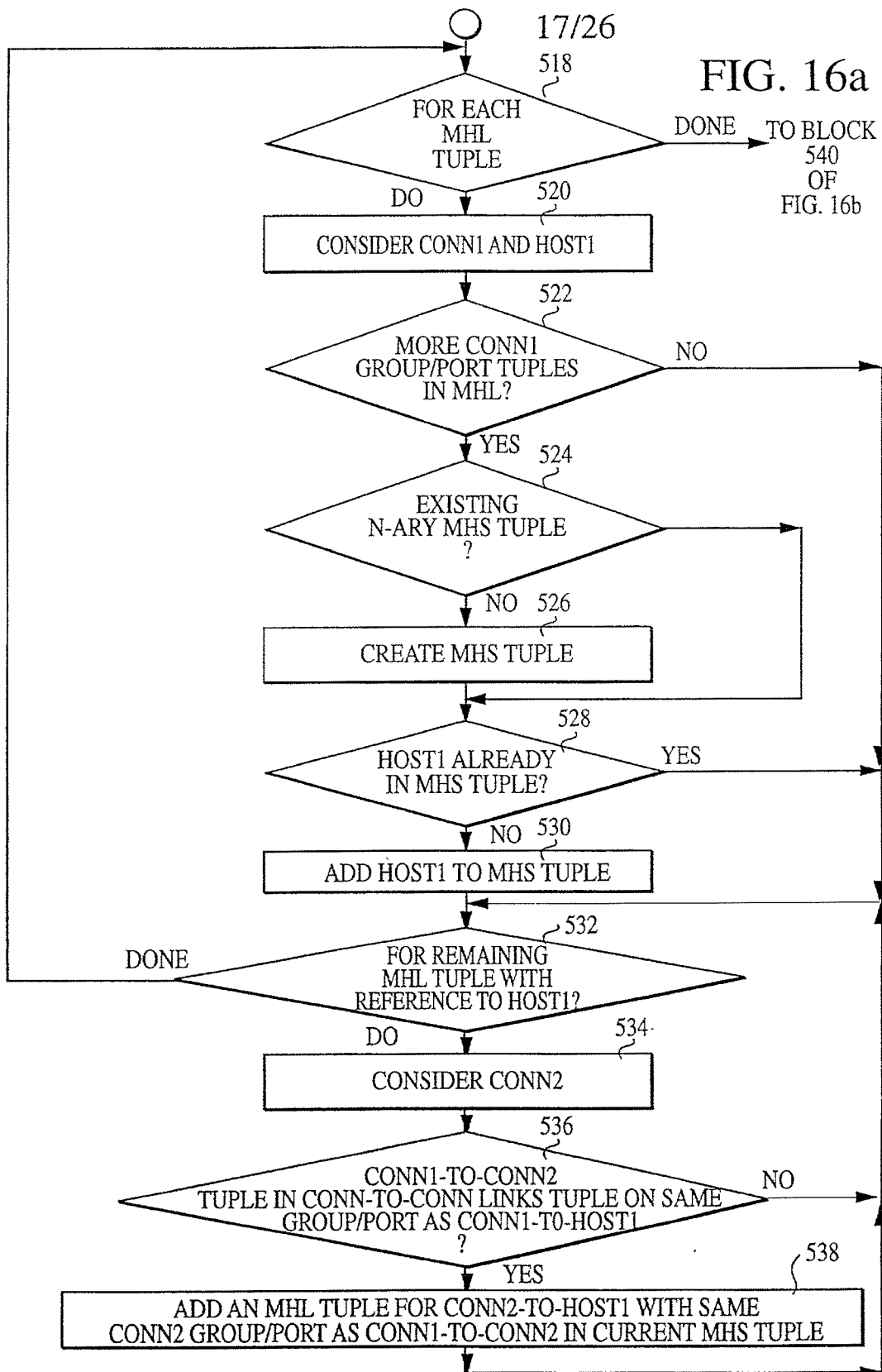
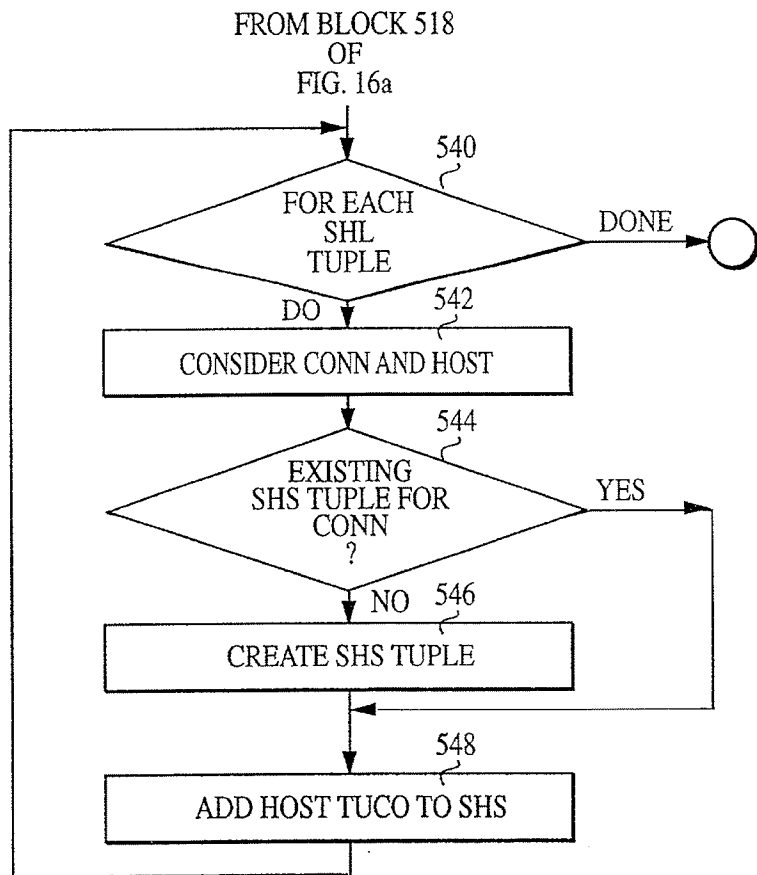
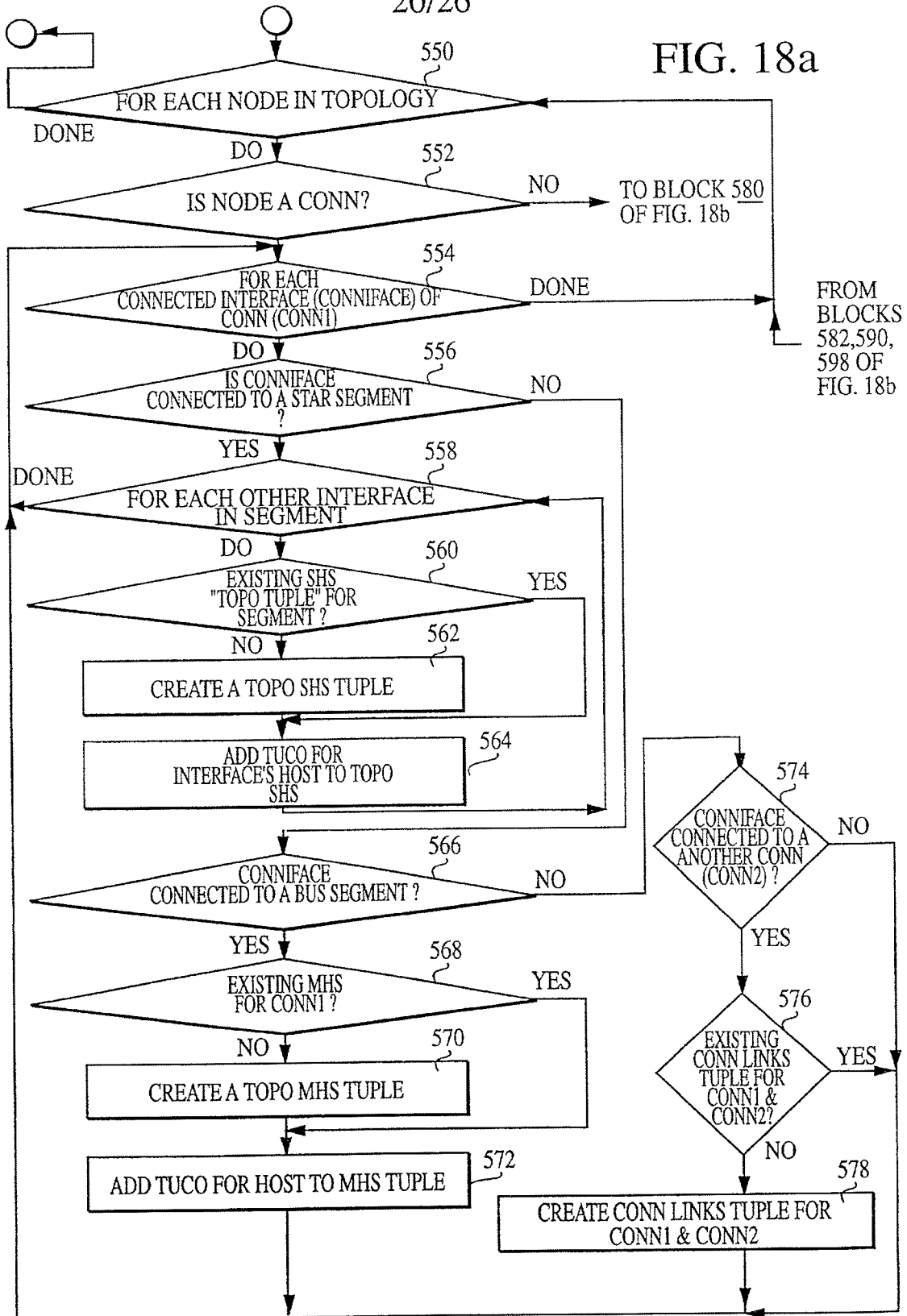


FIG. 16b



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FIG. 18a



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007607 2460060

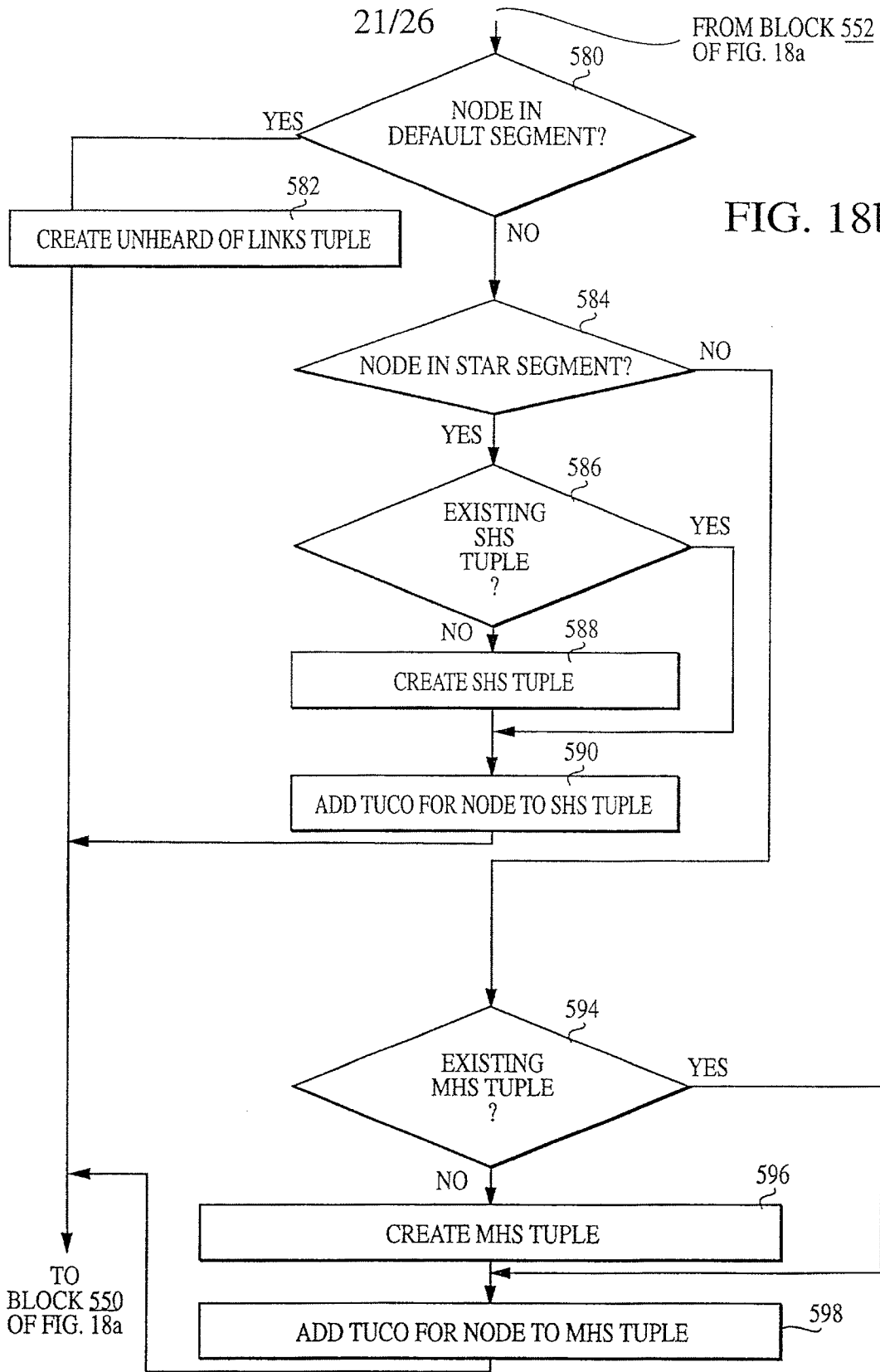
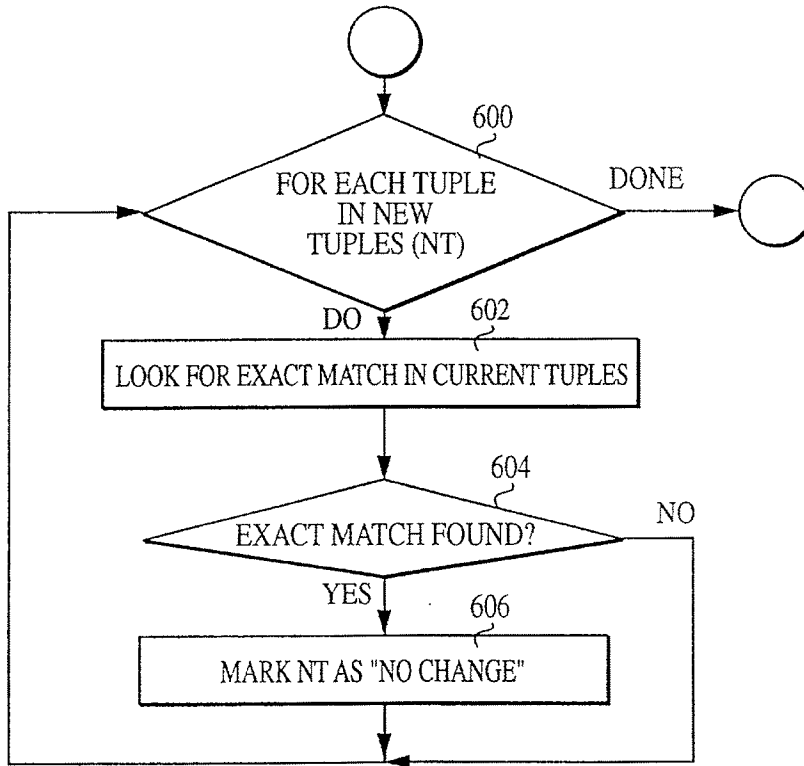
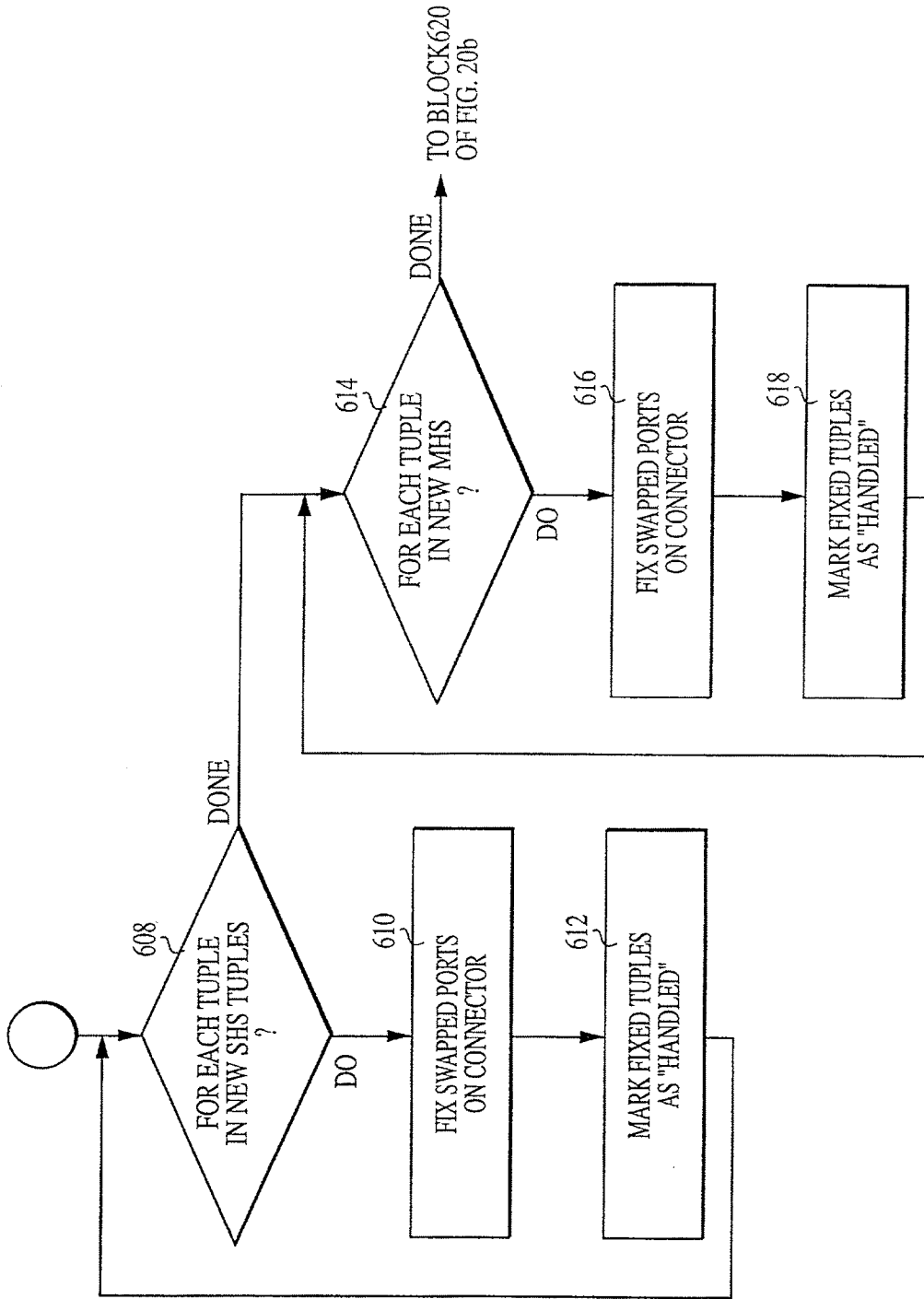


FIG. 19



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FIG. 20a



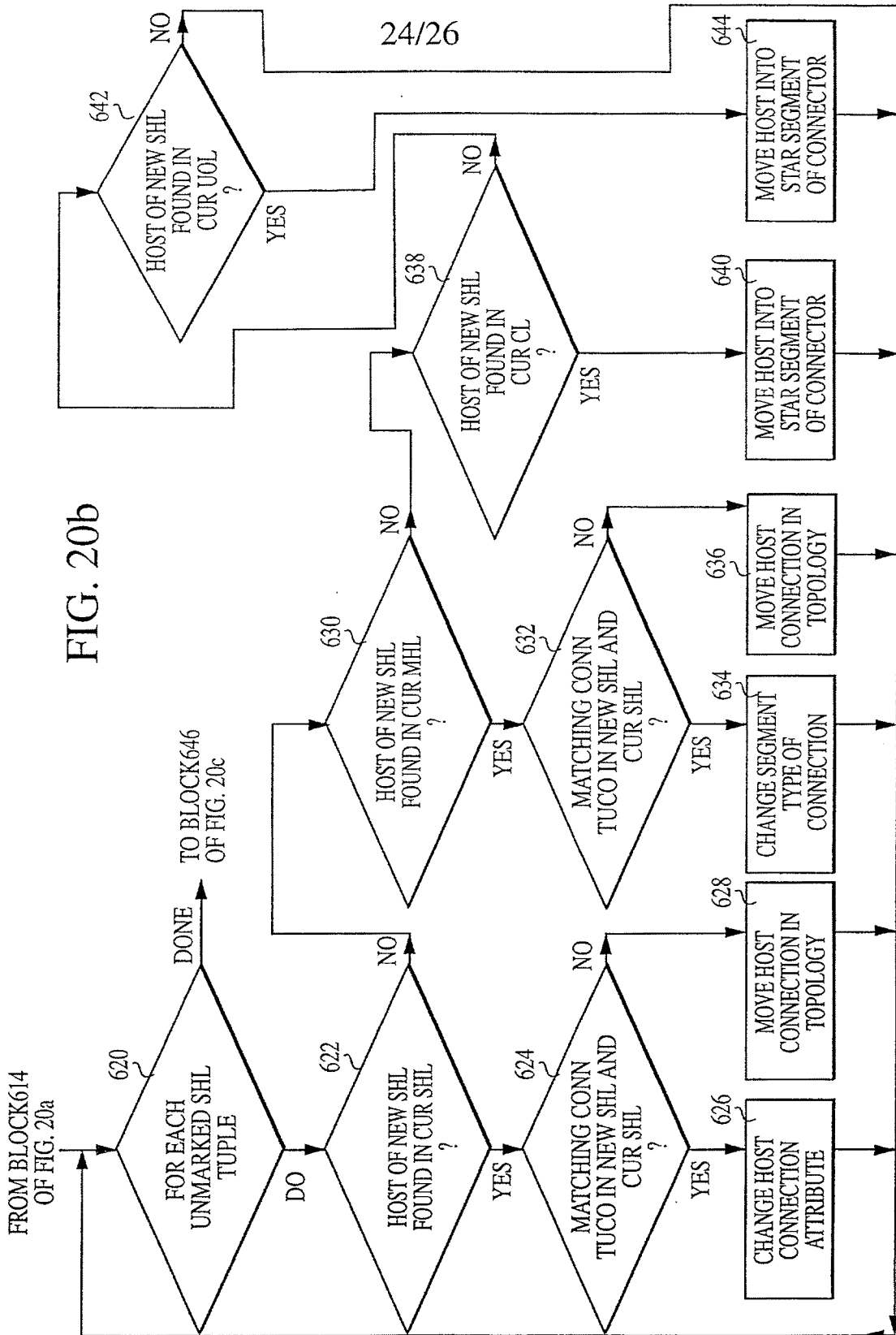


FIG. 20b

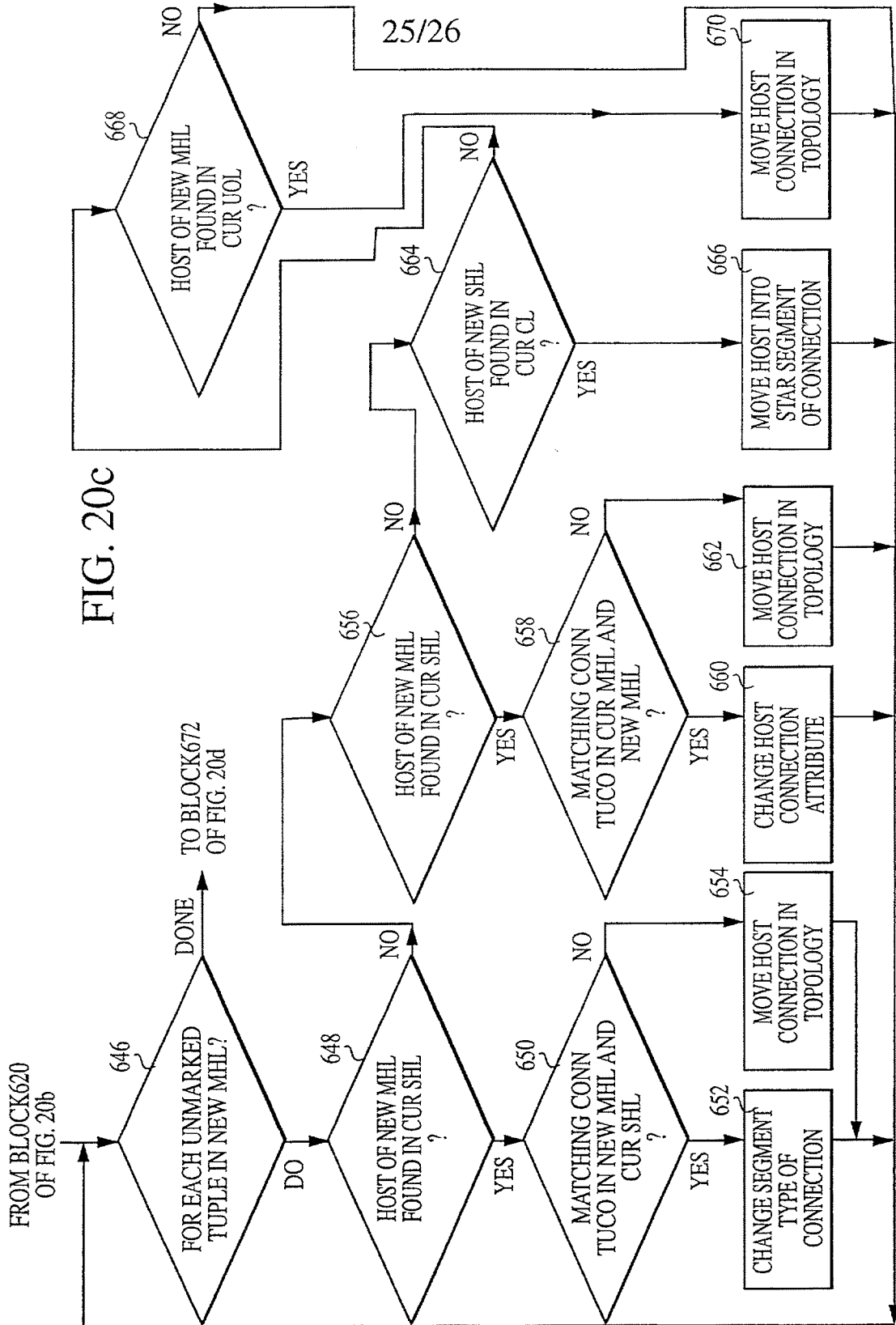
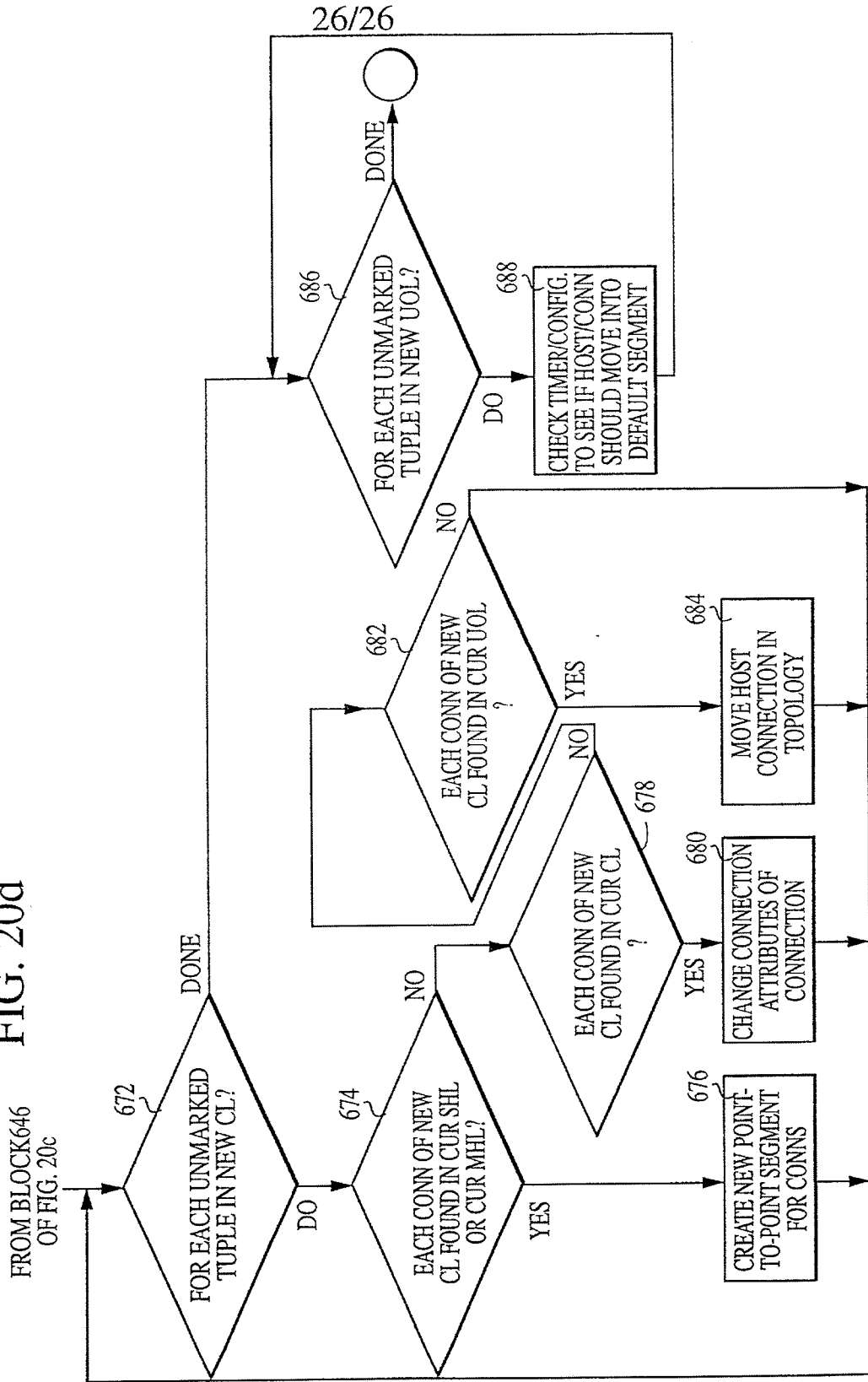


FIG. 20c

FIG. 20d



DECLARATION AND POWER OF ATTORNEY FOR PATENT APPLICATION

ATTORNEY DOCKET NO. 10008102-1

As a below named inventor, I hereby declare that:

My residence/post office address and citizenship are as stated below next to my name;

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled:

Method And System For Identifying And Processing Changes To A Network Topology

the specification of which is attached hereto unless the following box is checked:

() was filed on _____ as US Application Serial No. or PCT International Application Number _____ and was amended on _____ (if applicable).

I hereby state that I have reviewed and understood the contents of the above-identified specification, including the claims, as amended by any amendment(s) referred to above. I acknowledge the duty to disclose all information which is material to patentability as defined in 37 CFR 1.56.

Foreign Application(s) and/or Claim of Foreign Priority

I hereby claim foreign priority benefits under Title 35, United States Code Section 119 of any foreign application(s) for patent or inventor(s) certificate listed below and have also identified below any foreign application for patent or inventor(s) certificate having a filing date before that of the application on which priority is claimed:

COUNTRY	APPLICATION NUMBER	DATE FILED	PRIORITY CLAIMED UNDER 35 U.S.C. 119	
			YES: ___	NO: ___
N/A			YES: ___	NO: ___
			YES: ___	NO: ___

Provisional Application

I hereby claim the benefit under Title 35, United States Code Section 119(e) of any United States provisional application(s) listed below:

APPLICATION SERIAL NUMBER	FILING DATE
N/A	

U. S. Priority Claim

I hereby claim the benefit under Title 35, United States Code, Section 120 of any United States application(s) listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of Title 35, United States Code Section 112, I acknowledge the duty to disclose material information as defined in Title 37, Code of Federal Regulations, Section 1.56(a) which occurred between the filing date of the prior application and the national or PCT international filing date of this application:

APPLICATION SERIAL NUMBER	FILING DATE	STATUS (patented/pending/abandoned)
N/A		

POWER OF ATTORNEY:

As a named inventor, I hereby appoint the following attorney(s) and/or agent(s) to prosecute this application and transact all business in the Patent and Trademark Office connected therewith:

Customer Number 022879

Place Customer Number Bar Code Label here

Send Correspondence to:
HEWLETT-PACKARD COMPANY
 Intellectual Property Administration
 P.O. Box 272400
 Fort Collins, Colorado 80527-2400

Direct Telephone Calls To:

T. Grant Ritz
 (970) 898-0697

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

Full Name of Inventor: Eric A Pulsipher Citizenship: US

Residence: 2937 Redburn Drive Ft Collins CO 80525

Post Office Address: Same as residence

Inventor's Signature:  Date: 10/31/2000

20001031 16:00:00

PATENT NUMBER

Subclass
CLASS
ISSUE CLASSIFICATION

U.S. UTILITY Patent Application

74

O.I.P.E. PATENT DATE
 SCANNED *MH* *Am²* O.A. *CTH*

APPLICATION NO.	CONT/PRIOR	CLASS	SUBCLASS	ART UNIT	EXAMINER
09/703942		370	254	2661 2664	<i>FRAJ, M</i> <i>Schultz</i>

APPLICANTS
 Eric Pulsipher
 Joseph Hunt

TITLE
 Method and system for identifying and processing changes to a network topology

Best Available Copy

PTO-2040
12/93

ISSUING CLASSIFICATION							
ORIGINAL		CROSS REFERENCE(S)					
CLASS	SUBCLASS	CLASS	SUBCLASS (ONE SUBCLASS PER BLOCK)				
INTERNATIONAL CLASSIFICATION							

Continued on Issue Slip Inside File Jacket

<input type="checkbox"/> TERMINAL DISCLAIMER <input type="checkbox"/> The term of this patent subsequent to _____ (date) has been disclaimed. <input type="checkbox"/> The term of this patent shall not extend beyond the expiration date of U.S. Patent. No. _____ <input type="checkbox"/> The terminal _____ months of this patent have been disclaimed.	DRAWINGS Sheets Drwg. Figs. Drwg. Print Fig.			CLAIMS ALLOWED Total Claims Print Claim for O.G.	
	_____ (Assistant Examiner) (Date)			NOTICE OF ALLOWANCE MAILED _____ ISSUE FEE Amount Due Date Paid	
	_____ (Primary Examiner) (Date)			ISSUE BATCH NUMBER _____	
	_____ (Legal Instruments Examiner) (Date)				

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SEARCHED			
Class	Sub.	Date	Exmr.
370 ↓	229 216 217 221 225 254 256 257 258	6/21/04 ↓	WCS ↓

SEARCH NOTES (INCLUDING SEARCH STRATEGY)		
	Date	Exmr.
WSP# E10 Sp0	6/21/04	WCS

INTERFERENCE SEARCHED			
Class	Sub.	Date	Exmr.

(RIGHT OUTSIDE)

ISSUE SLIP STAPLE AREA (for additional cross references)

POSITION	INITIALS	ID NO.	DATE
FEE DETERMINATION			
O.I.P.E. CLASSIFIER	<i>[Handwritten]</i>		11-25-00
FORMALITY REVIEW	<i>[Handwritten]</i>	75553	2-5-01
RESPONSE FORMALITY REVIEW			

INDEX OF CLAIMS

- ✓ Rejected N Non-elected
- = Allowed I Interference
- (Through numeral) ... Canceled A Appeal
- + Restricted O Objected

Claim	Date
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HEWLETT-PACKARD COMPANY
Intellectual Property Administration
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Fort Collins, Colorado 80527-2400

11-02-00
PATENT APPLICATION
ATTORNEY SCKET NO. 10008102-1

A

IN THE U.S. PATENT AND TRADEMARK OFFICE
Patent Application Transmittal Letter

COMMISSIONER FOR PATENTS
Washington, D.C. 20231

Sir:

Transmitted herewith for filing under 37 CFR 1.53(b) is a(n): Utility () Design
 original patent application,
() continuation-in-part application

JC835 U.S. PTO
09/703942
10/31/00

INVENTOR(S): Eric A Pulsipher et al

TITLE: Method And System For Identifying And Processing Changes To A Network Topology

Enclosed are:

- The Declaration and Power of Attorney. (X) signed () unsigned or partially signed
- 26 sheets of drawings (one set) () Associate Power of Attorney
- () Form PTO-1449 () Information Disclosure Statement and Form PTO-1449
- () Priority document(s) () (Other) _____ (fee \$ _____)

CLAIMS AS FILED BY OTHER THAN A SMALL ENTITY				
(1) FOR	(2) NUMBER FILED	(3) NUMBER EXTRA	(4) RATE	(5) TOTALS
TOTAL CLAIMS	20 — 20	0	X \$18	\$ 0
INDEPENDENT CLAIMS	3 — 3	0	X \$80	\$ 0
ANY MULTIPLE DEPENDENT CLAIMS	0		\$270	\$ 0
BASIC FEE: Design (\$320.00); Utility (\$710.00)				\$ 710
TOTAL FILING FEE				\$ 710
OTHER FEES				\$
TOTAL CHARGES TO DEPOSIT ACCOUNT				\$ 710

Charge \$ 710 to Deposit Account 08-2025. At any time during the pendency of this application, please charge any fees required or credit any over payment to Deposit Account 08-2025 pursuant to 37 CFR 1.25. Additionally please charge any fees to Deposit Account 08-2025 under 37 CFR 1.16, 1.17, 1.19, 1.20 and 1.21. A duplicate copy of this sheet is enclosed.

"Express Mail" label no. EL523338183US

Date of Deposit Oct. 31, 2000

I hereby certify that this is being deposited with the United States Postal Service "Express Mail Post Office to Addressee" service under 37 CFR 1.10 on the date indicated above and is addressed to: Commissioner for Patents, Washington, D.C. 20231.

By Laura M. Clark
Typed Name: Laura M. Clark

Respectfully submitted,

Eric A Pulsipher et al

By T. Grant Ritz

T. Grant Ritz

Attorney/Agent for Applicant(s)

Reg. No. 39,819

Date: Oct. 31, 2000

Telephone No.: (970) 898-0697

HEWLETT-PACKARD COMPANY
 Intellectual Property Administration
 P. O. Box 272400
 Fort Collins, Colorado 80527-2400

11-02-00
 PATENT APPLICATION
 ATTORNEY SCKET NO. 10008102-1

A

IN THE U.S. PATENT AND TRADEMARK OFFICE
 Patent Application Transmittal Letter

COMMISSIONER FOR PATENTS
 Washington, D.C. 20231

Sir:

Transmitted herewith for filing under 37 CFR 1.53(b) is a(n): Utility () Design
 original patent application,
 continuation-in-part application

JC925 U.S. PTO
 09/703942
 10/31/00

INVENTOR(S): Eric A Pulsipher et al

TITLE: Method And System For Identifying And Processing Changes To A Network Topology

Enclosed are:

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- 26 sheets of drawings (one set) () Associate Power of Attorney
- () Form PTO-1449 () Information Disclosure Statement and Form PTO-1449
- () Priority document(s) () (Other) (fee \$ _____)

CLAIMS AS FILED BY OTHER THAN A SMALL ENTITY				
(1) FOR	(2) NUMBER FILED	(3) NUMBER EXTRA	(4) RATE	(5) TOTALS
TOTAL CLAIMS	20 — 20	0	X \$18	\$ 0
INDEPENDENT CLAIMS	3 — 3	0	X \$80	\$ 0
ANY MULTIPLE DEPENDENT CLAIMS	0		\$270	\$ 0
BASIC FEE: Design (\$320.00); Utility (\$710.00)				\$ 710
TOTAL FILING FEE				\$ 710
OTHER FEES				\$
TOTAL CHARGES TO DEPOSIT ACCOUNT				\$ 710

Charge \$ 710 to Deposit Account 08-2025. At any time during the pendency of this application, please charge any fees required or credit any over payment to Deposit Account 08-2025 pursuant to 37 CFR 1.25. Additionally please charge any fees to Deposit Account 08-2025 under 37 CFR 1.16, 1.17, 1.19, 1.20 and 1.21. A duplicate copy of this sheet is enclosed.

"Express Mail" label no. EL523338183US

Date of Deposit Oct. 31, 2000

I hereby certify that this is being deposited with the United States Postal Service "Express Mail Post Office to Addressee" service under 37 CFR 1.10 on the date indicated above and is addressed to: Commissioner for Patents, Washington, D.C. 20231.

By Laura M. Clark
 Typed Name: Laura M. Clark

Respectfully submitted,

Eric A Pulsipher et al

By T. Grant Ritz

T. Grant Ritz

Attorney/Agent for Applicant(s)

Reg. No. 39,819

Date: Oct. 31, 2000

Telephone No.: (970) 898-0697

OFFICE OF THE ATTORNEY GENERAL

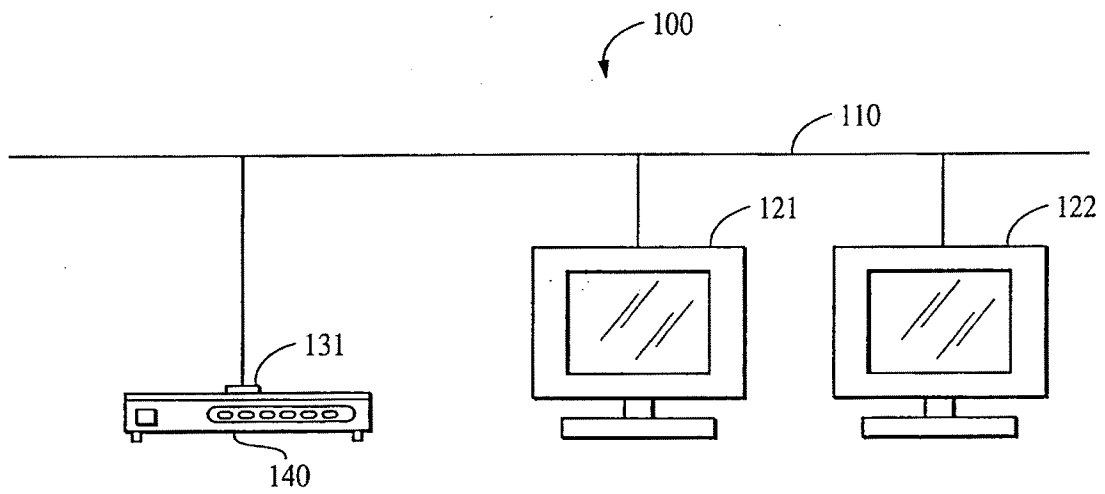


FIG. 1

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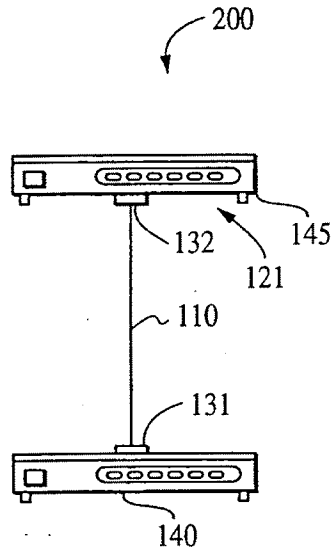


FIG. 2

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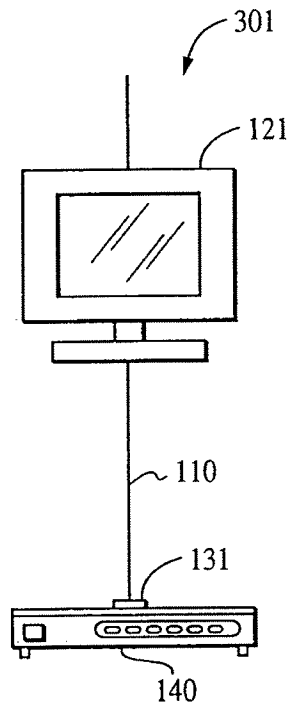


FIG. 3

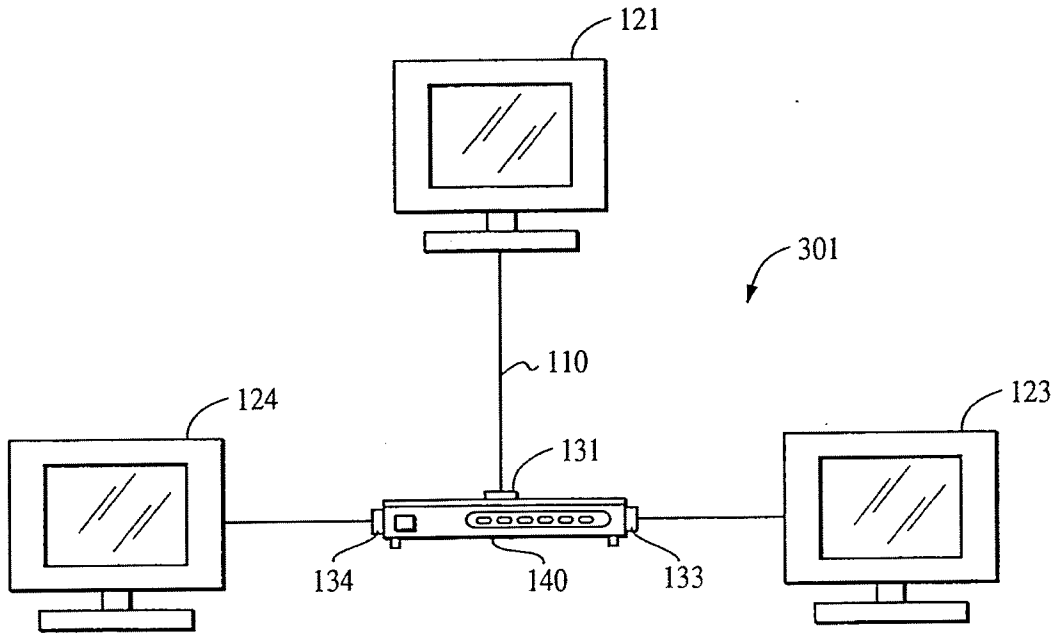


FIG. 4

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FIG. 5

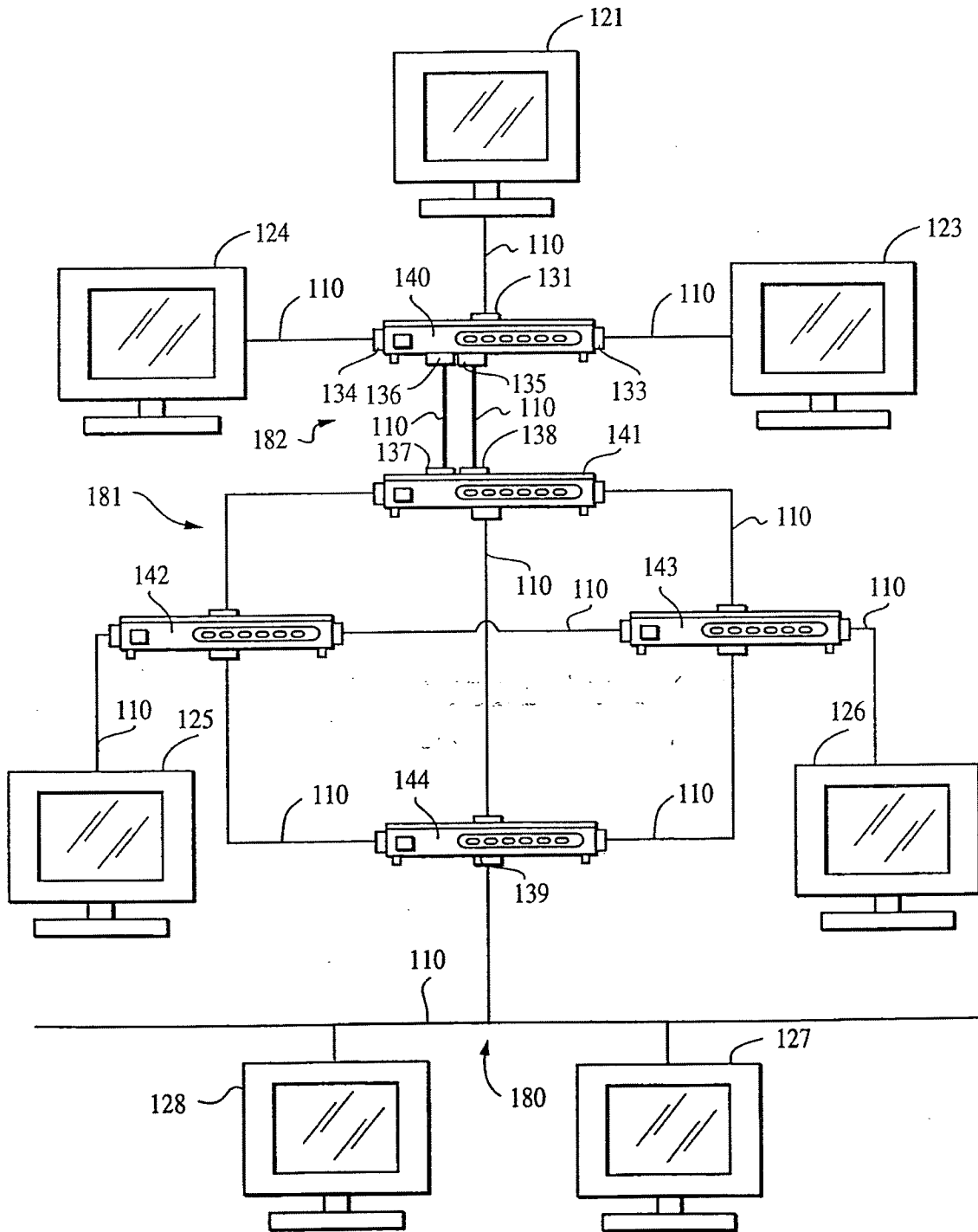
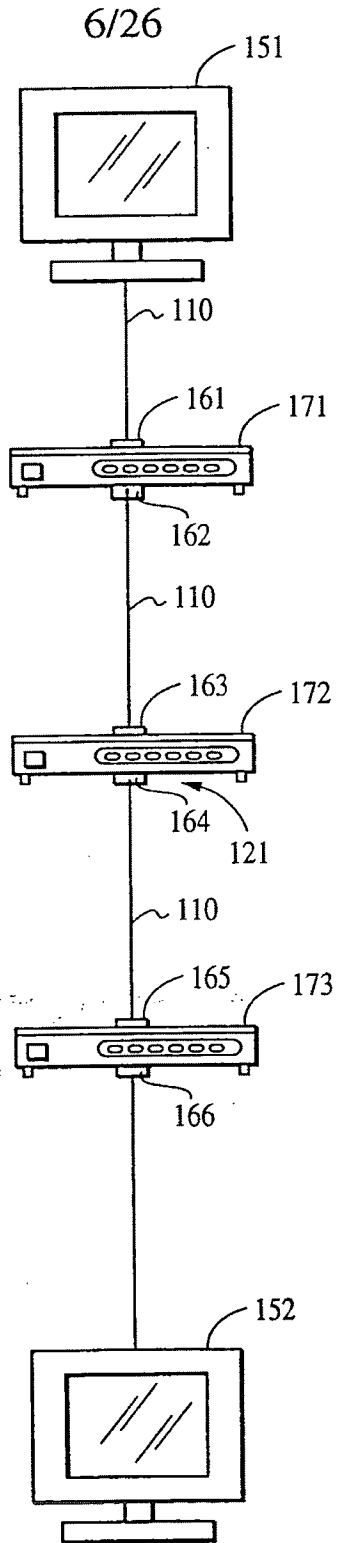
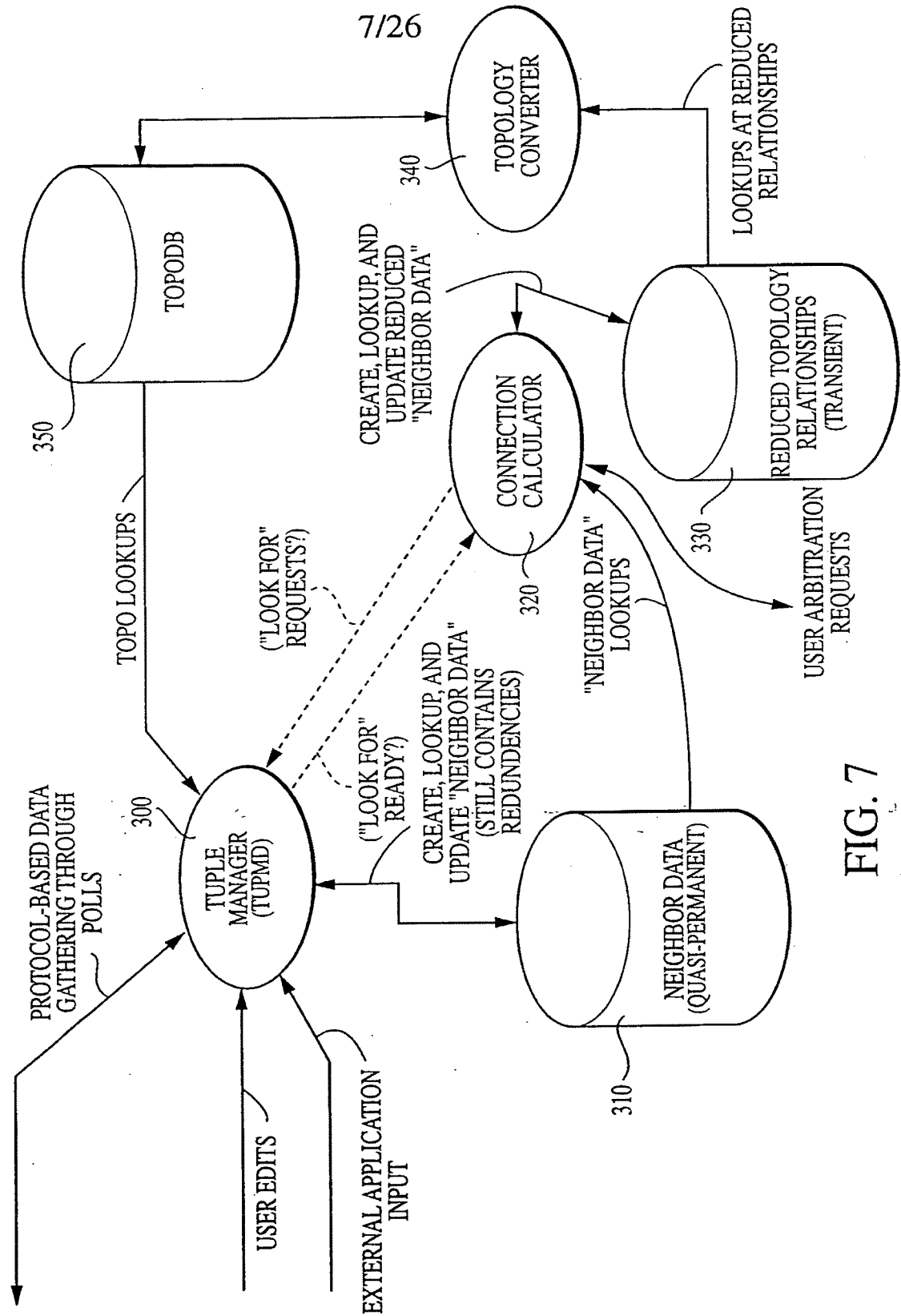


FIG. 5

FIG. 6





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FIG. 7

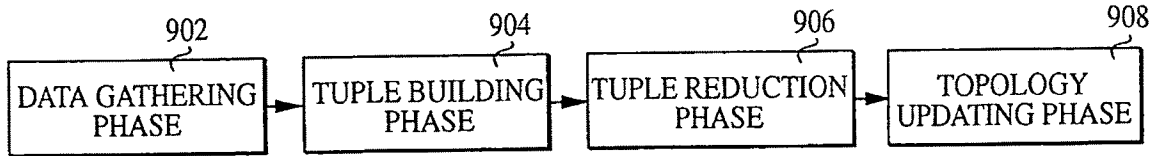


FIG. 8

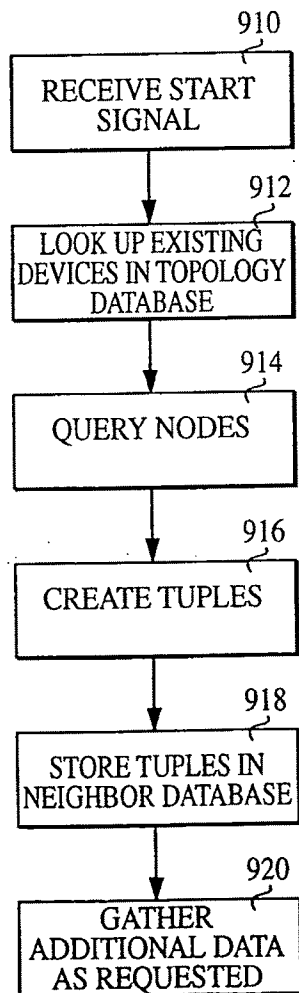


FIG. 9

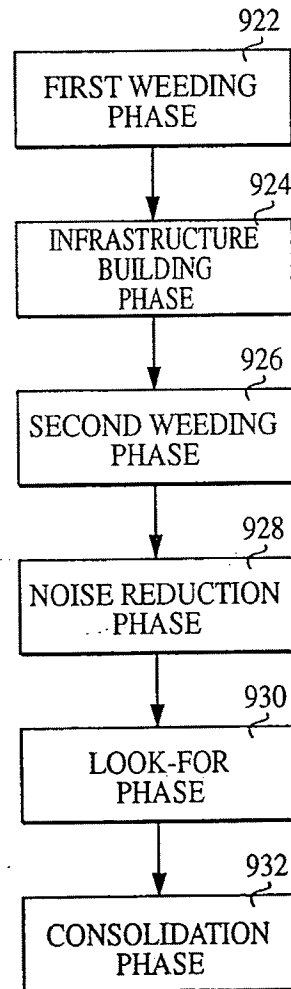


FIG. 10

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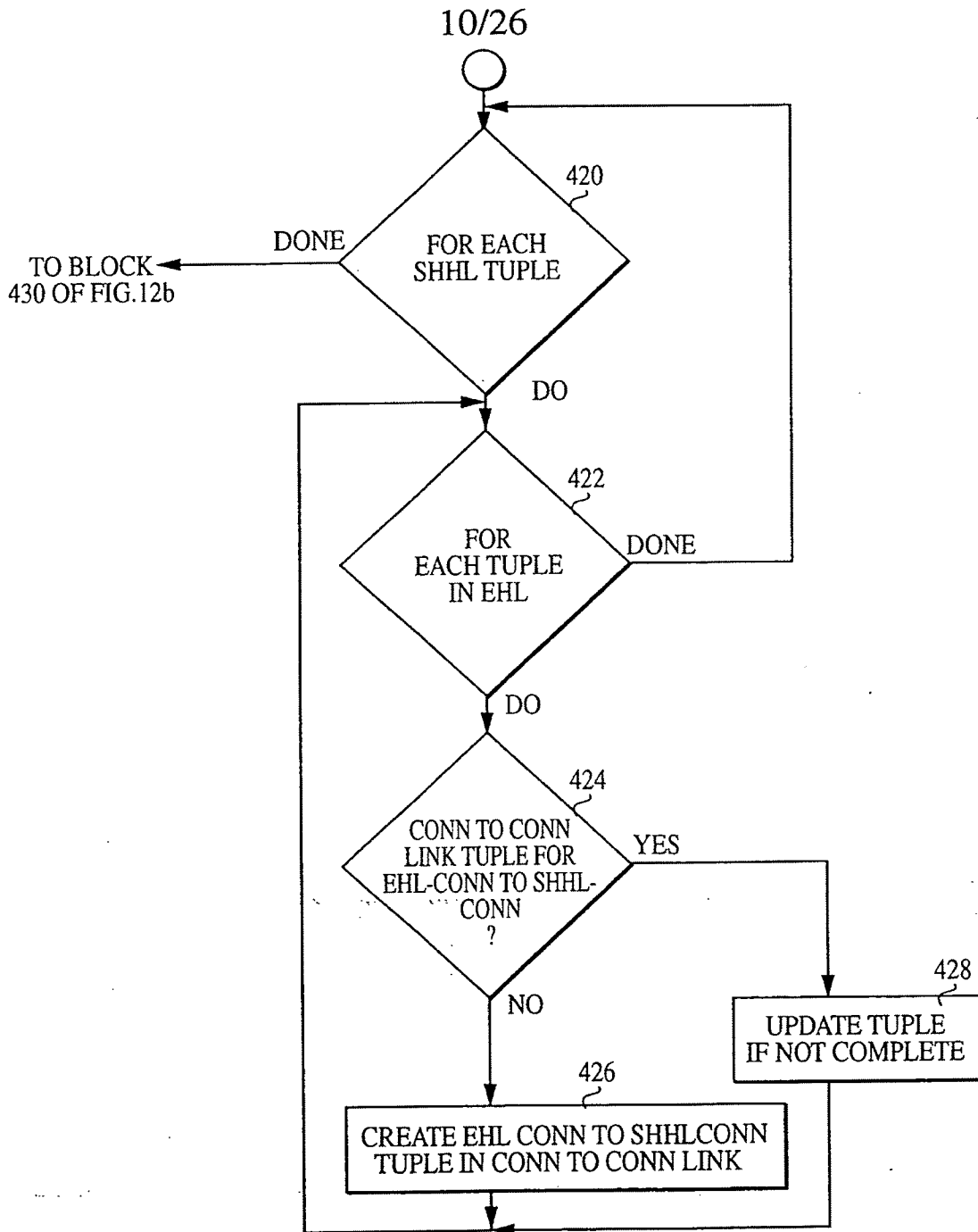
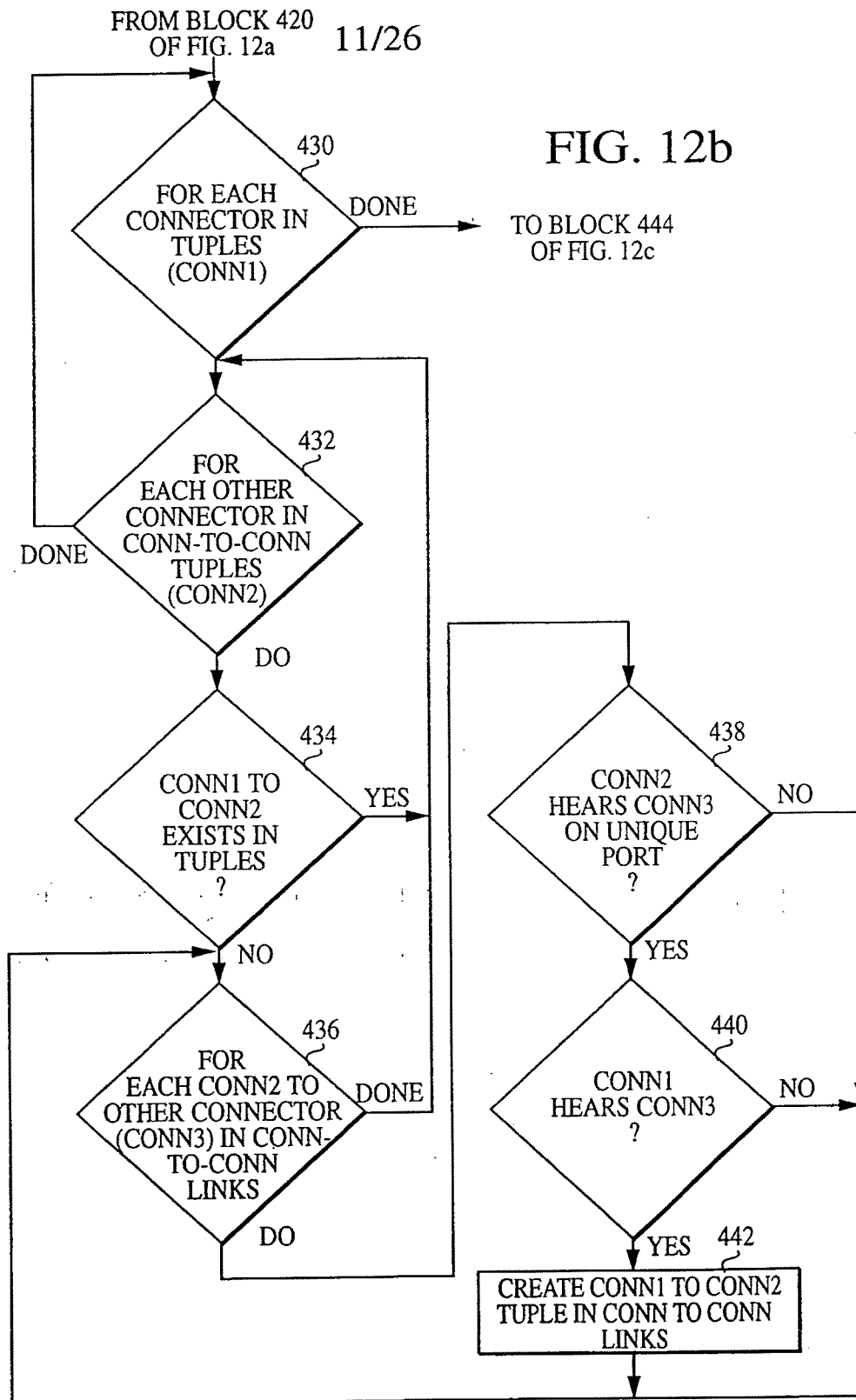


FIG. 12a

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DATE: 2/16/60

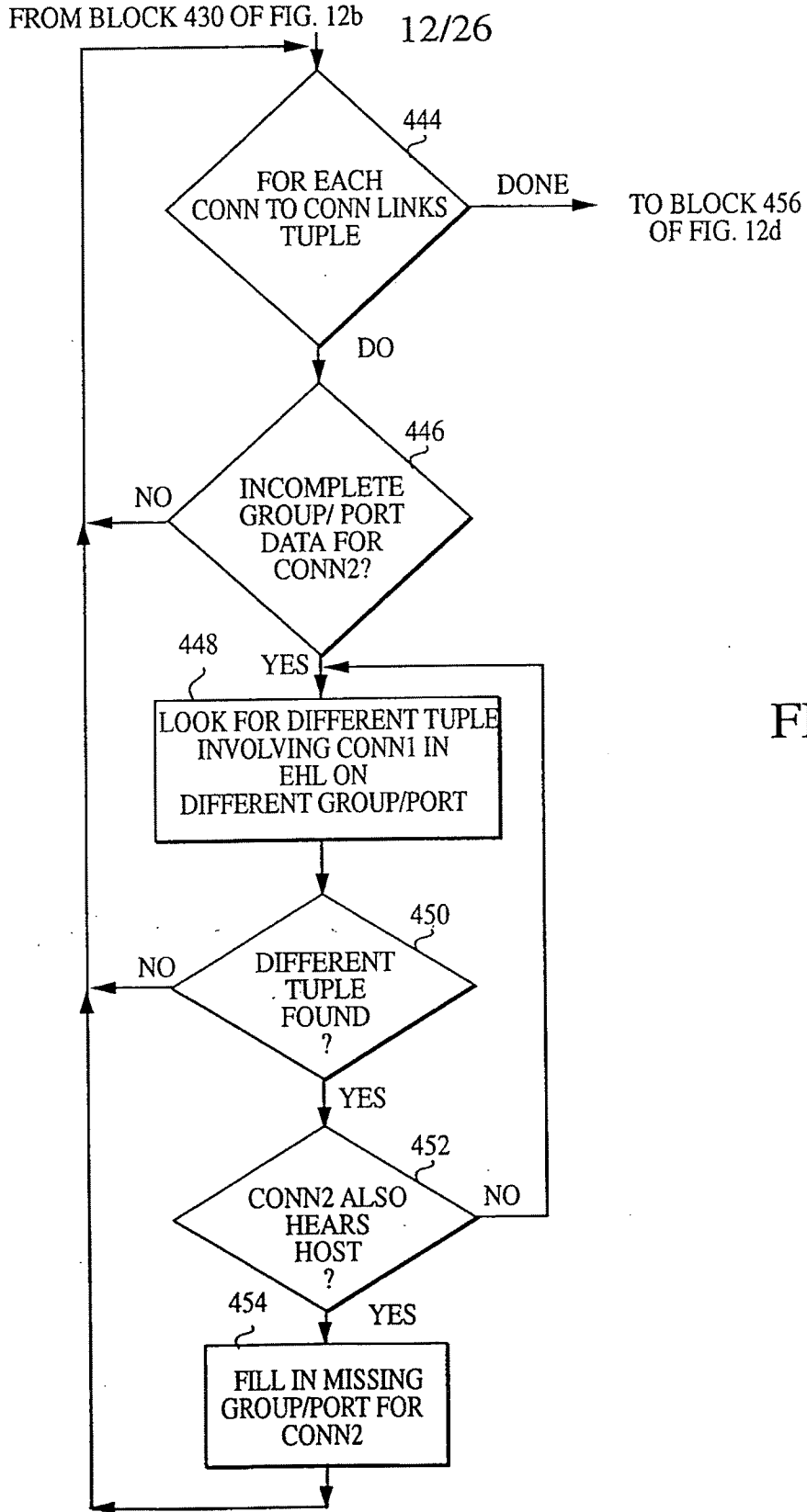
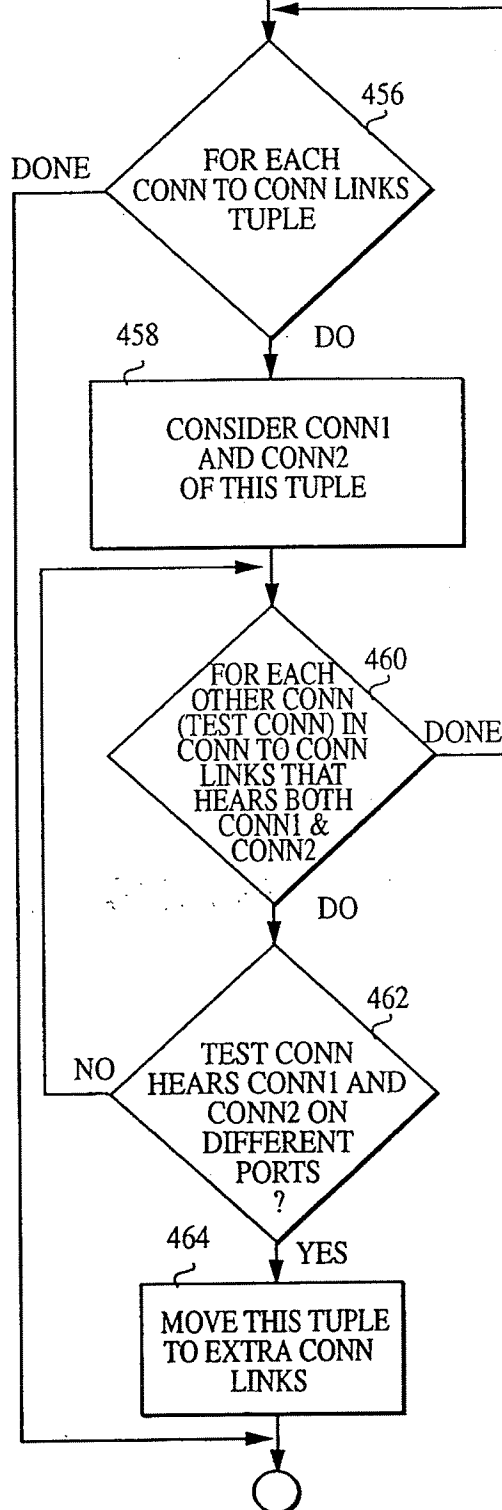


FIG. 12c

13/26
FROM BLOCK 444 OF FIG. 12c

FIG. 12d



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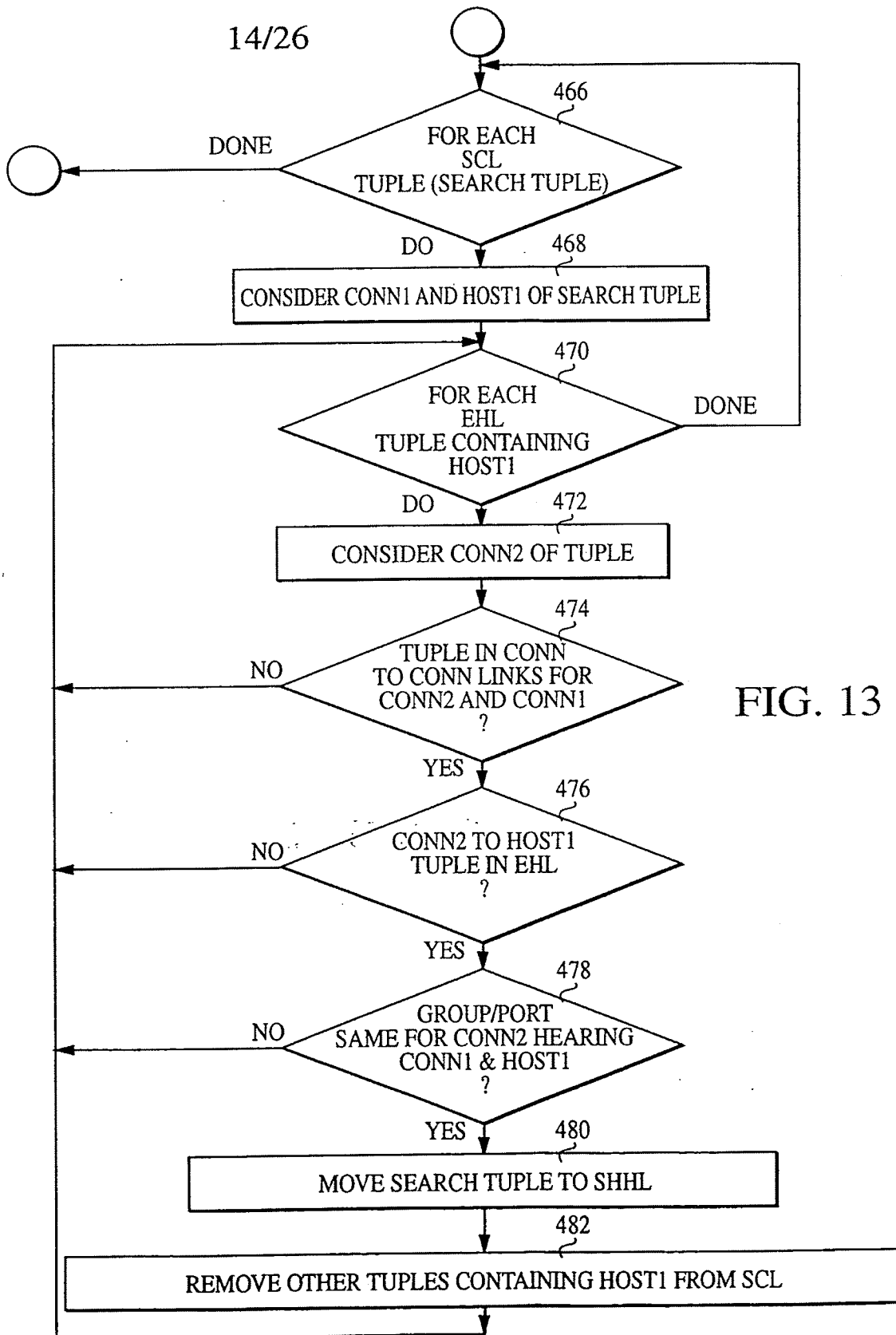


FIG. 13

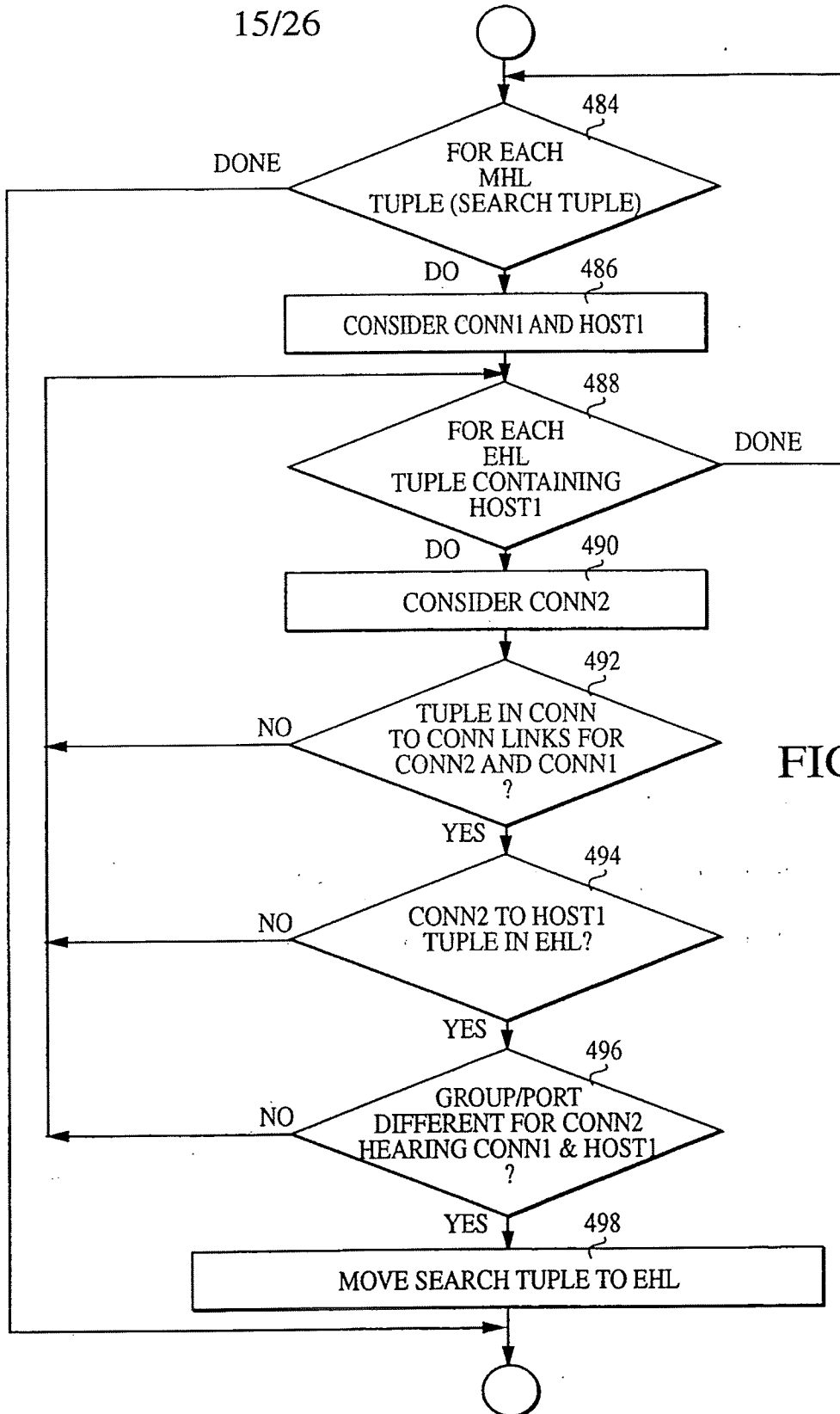
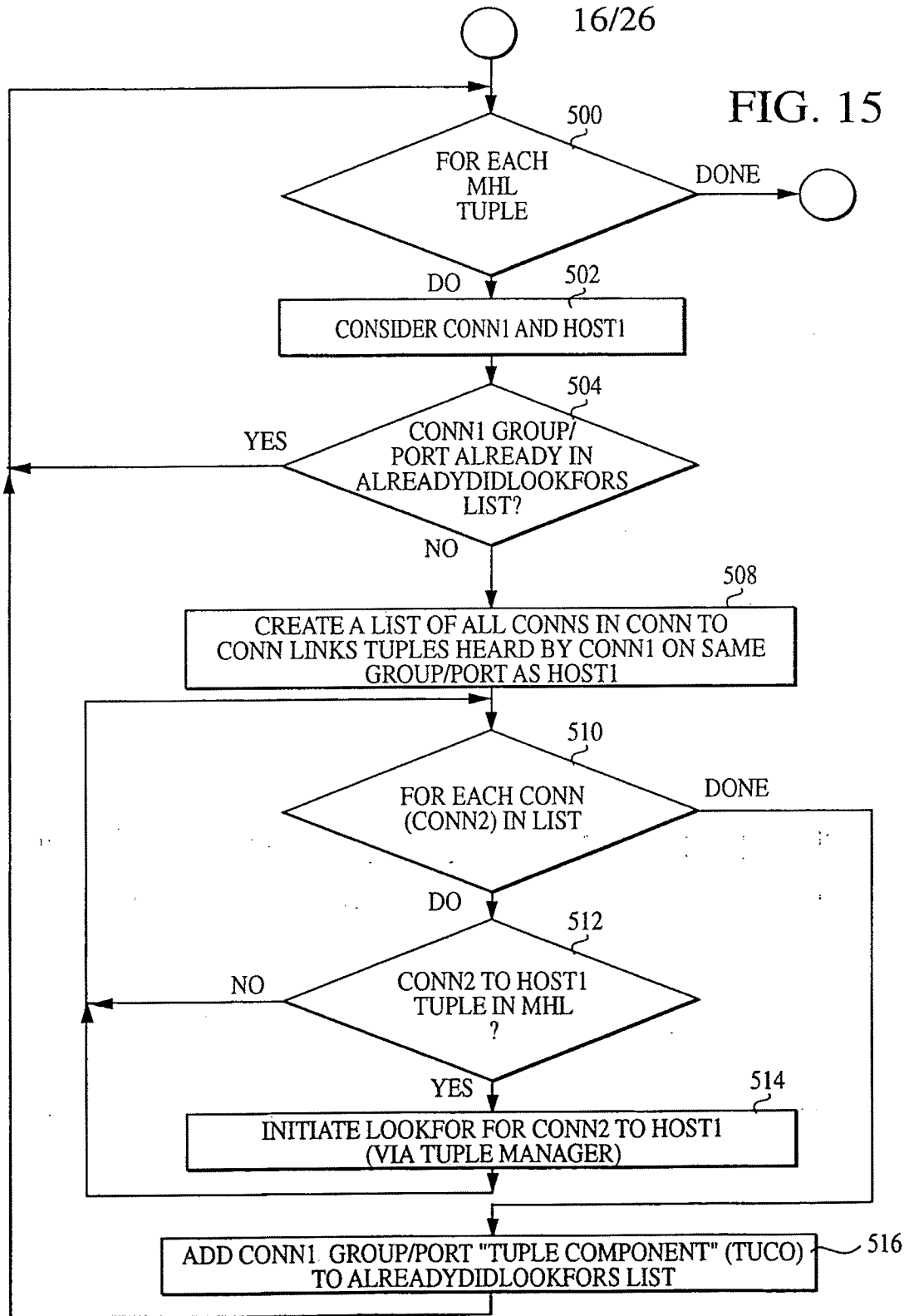


FIG. 14

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CONFIDENTIAL



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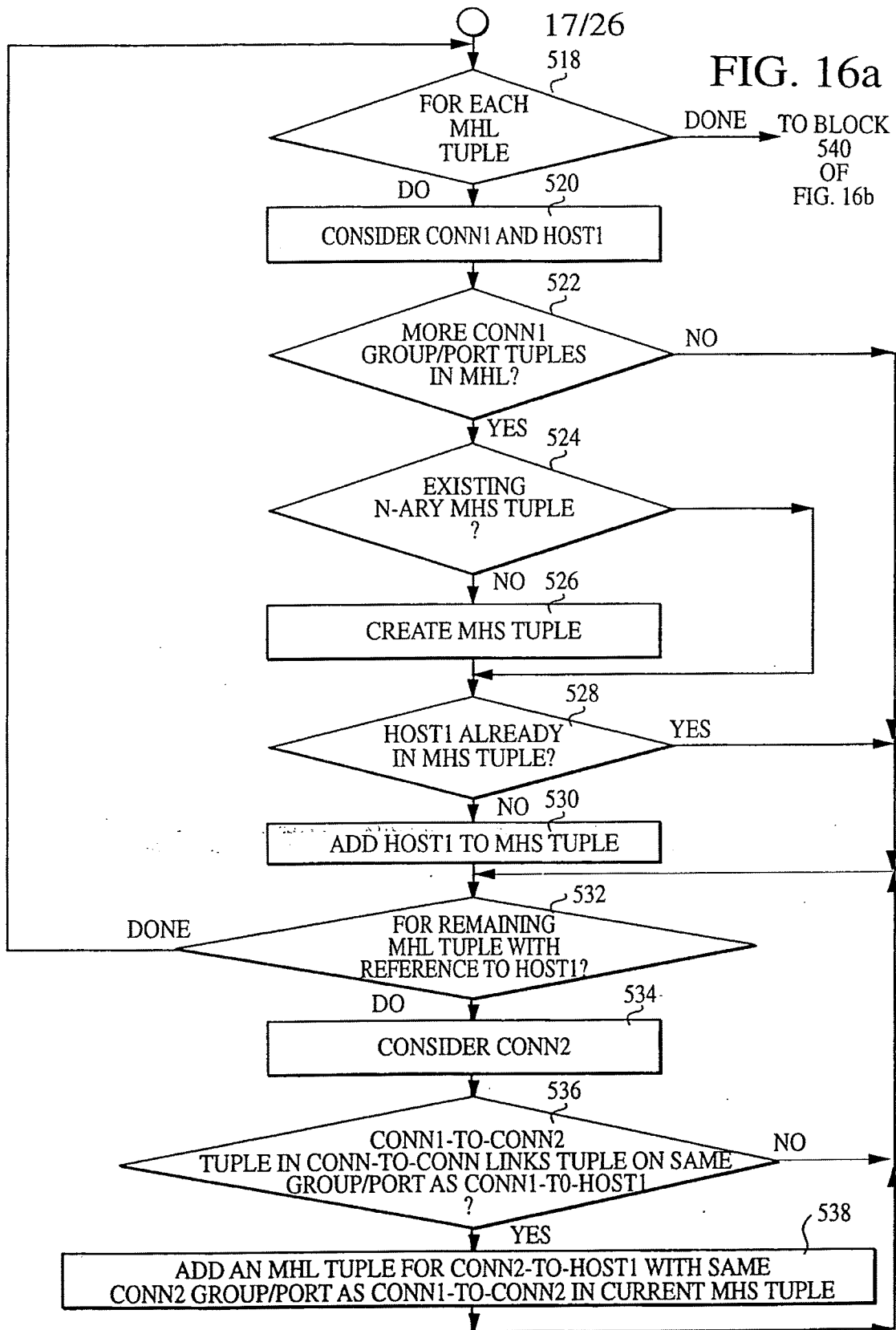
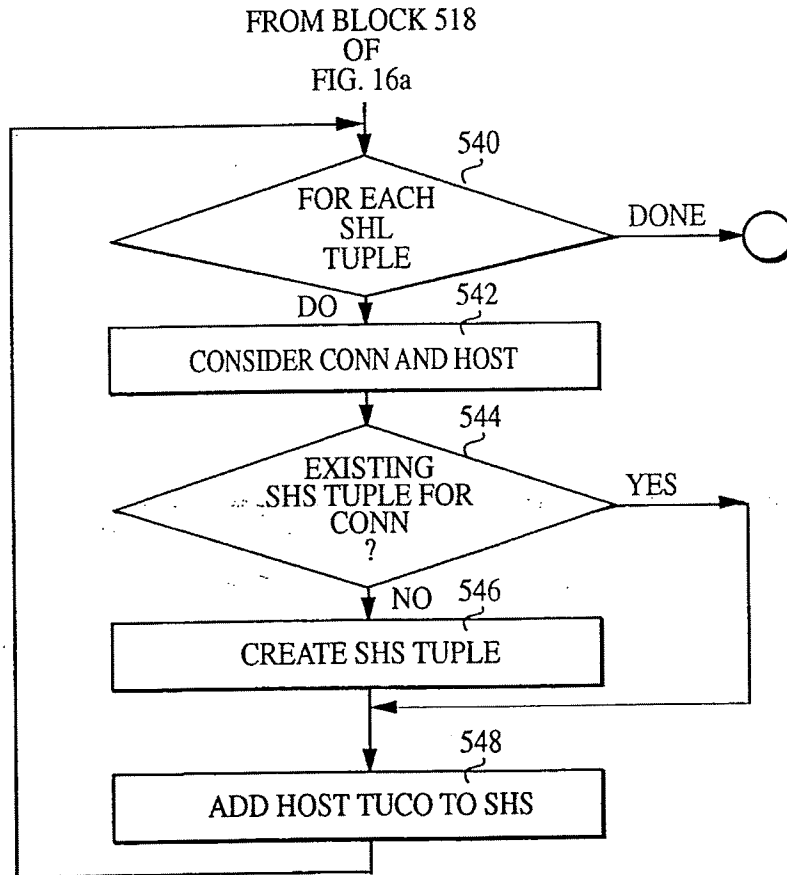
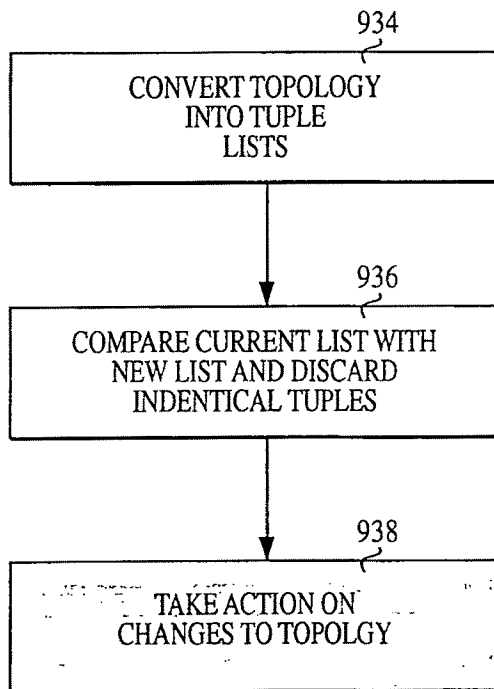


FIG. 16b



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FIG. 17



PROPERTY OF SPECTRUM

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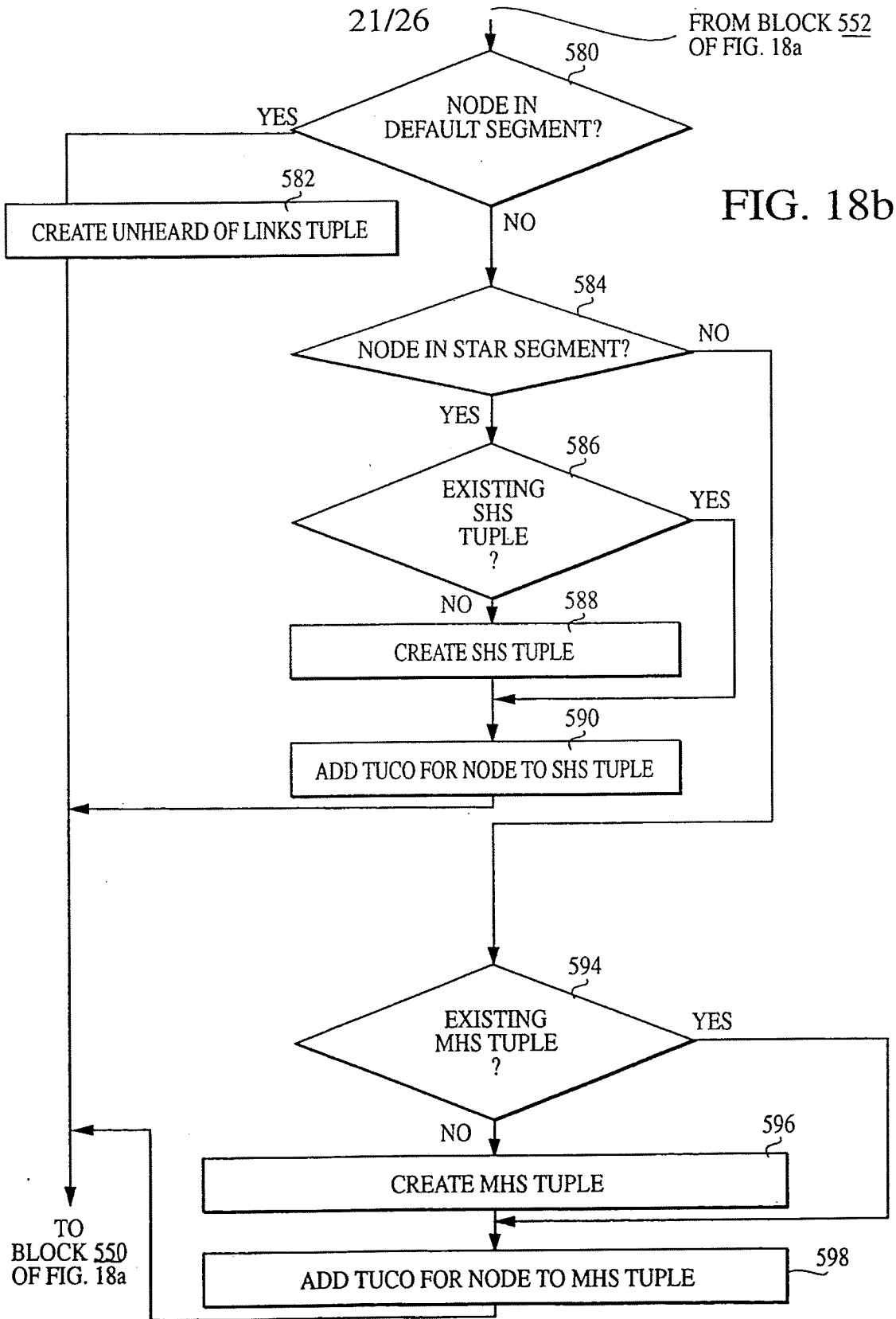
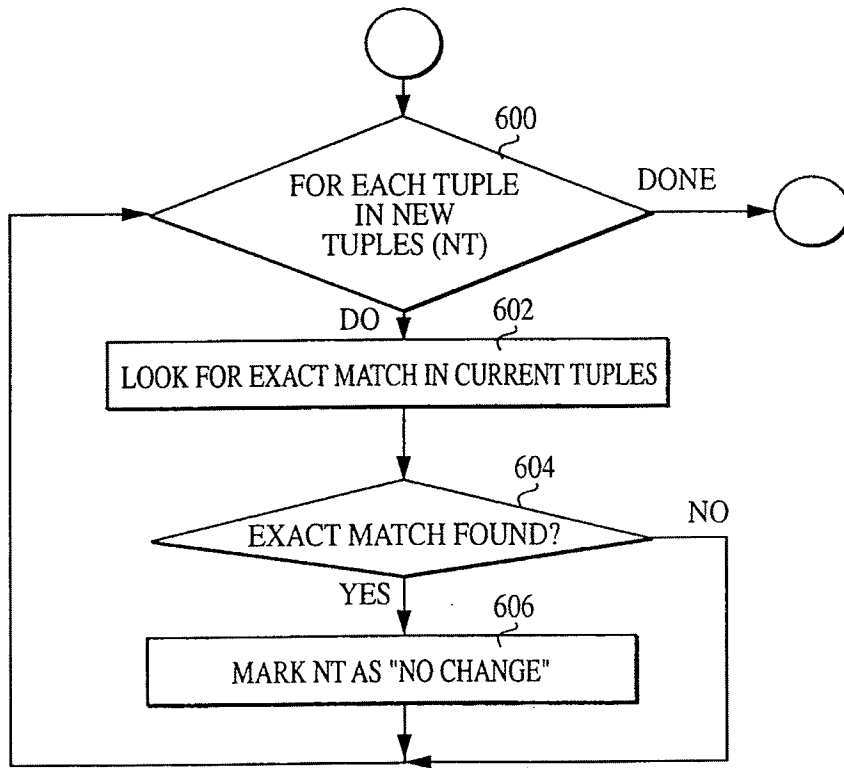
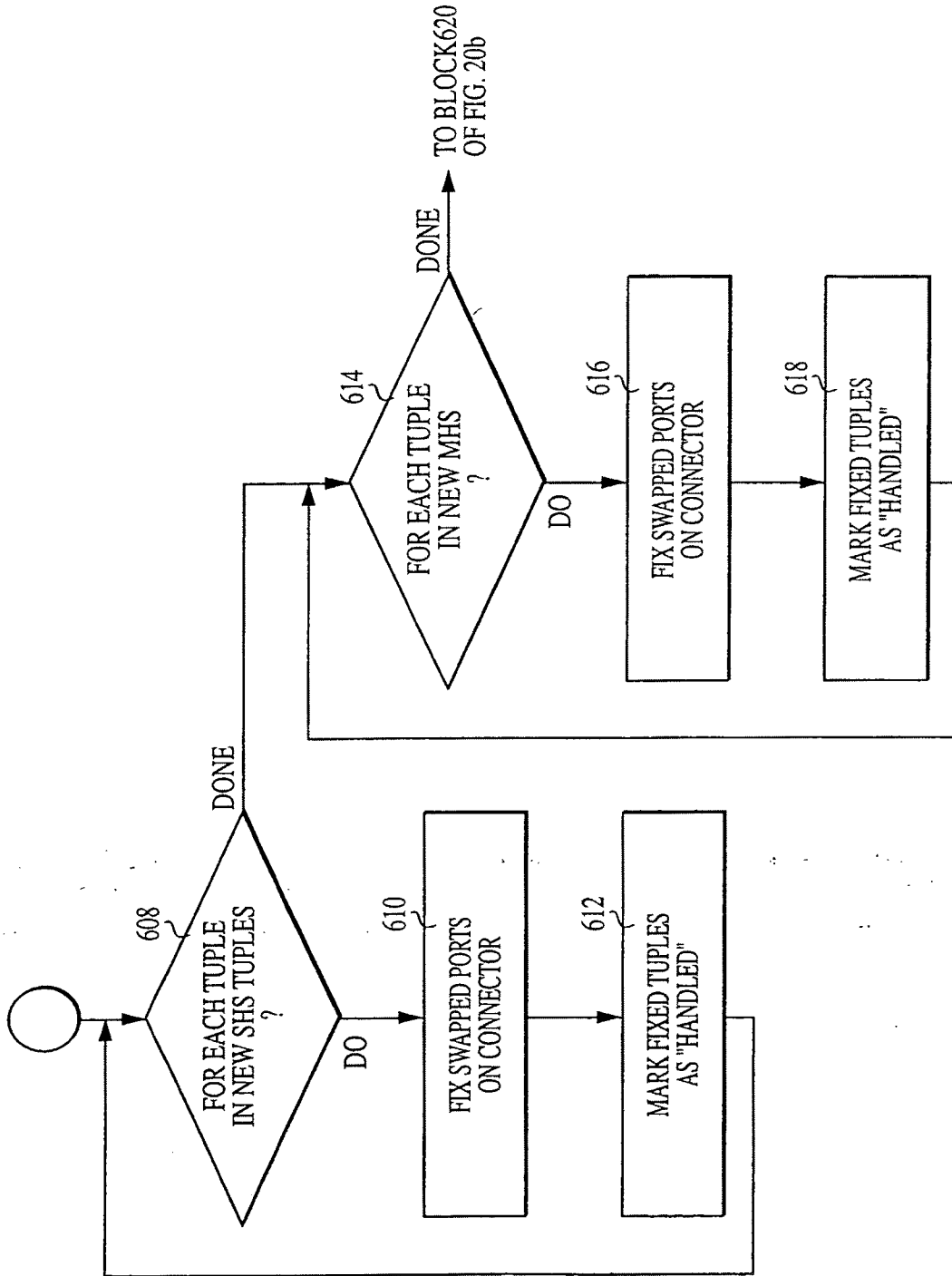


FIG. 19



OFFICE 24660460

FIG. 20a



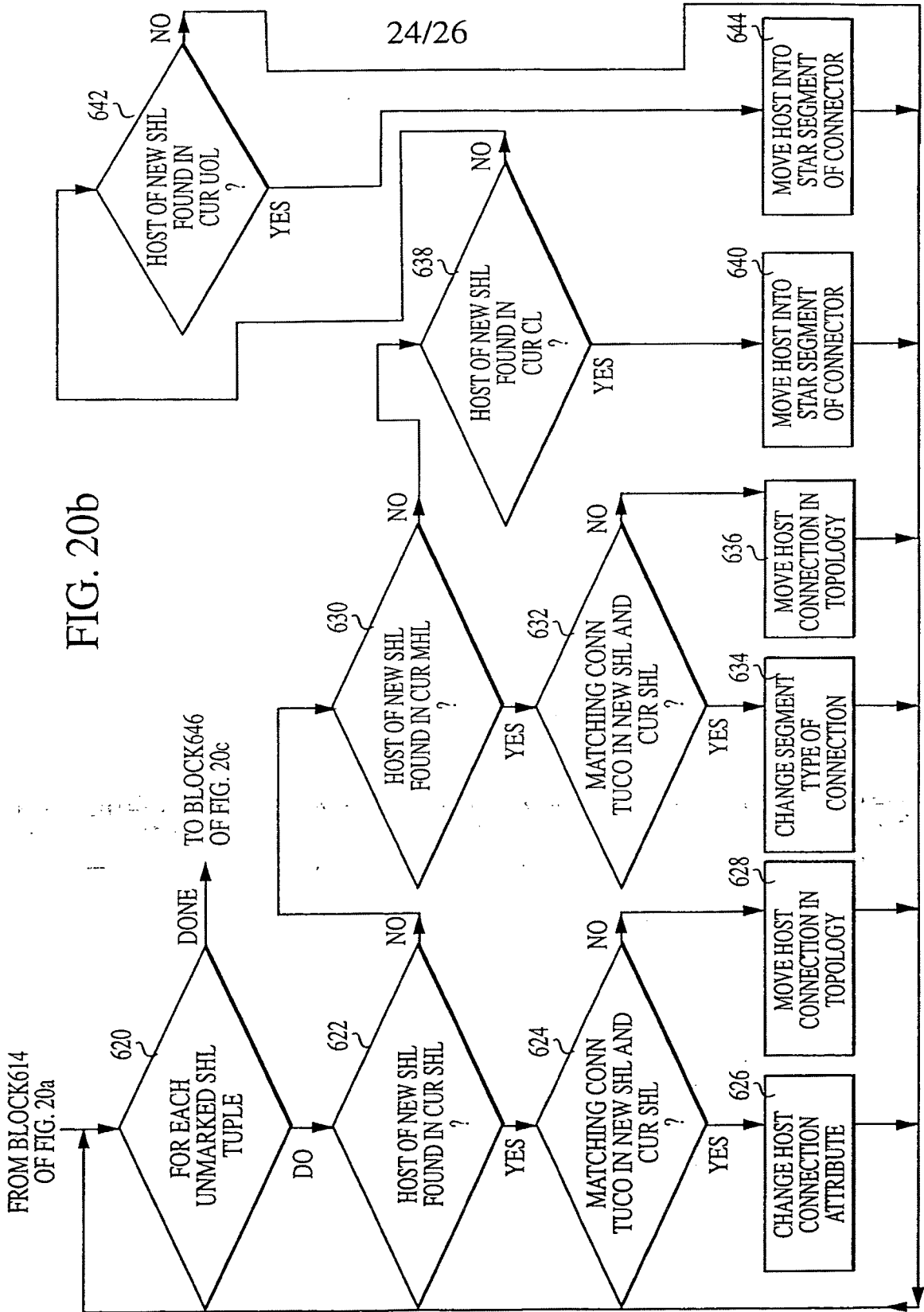


FIG. 20b

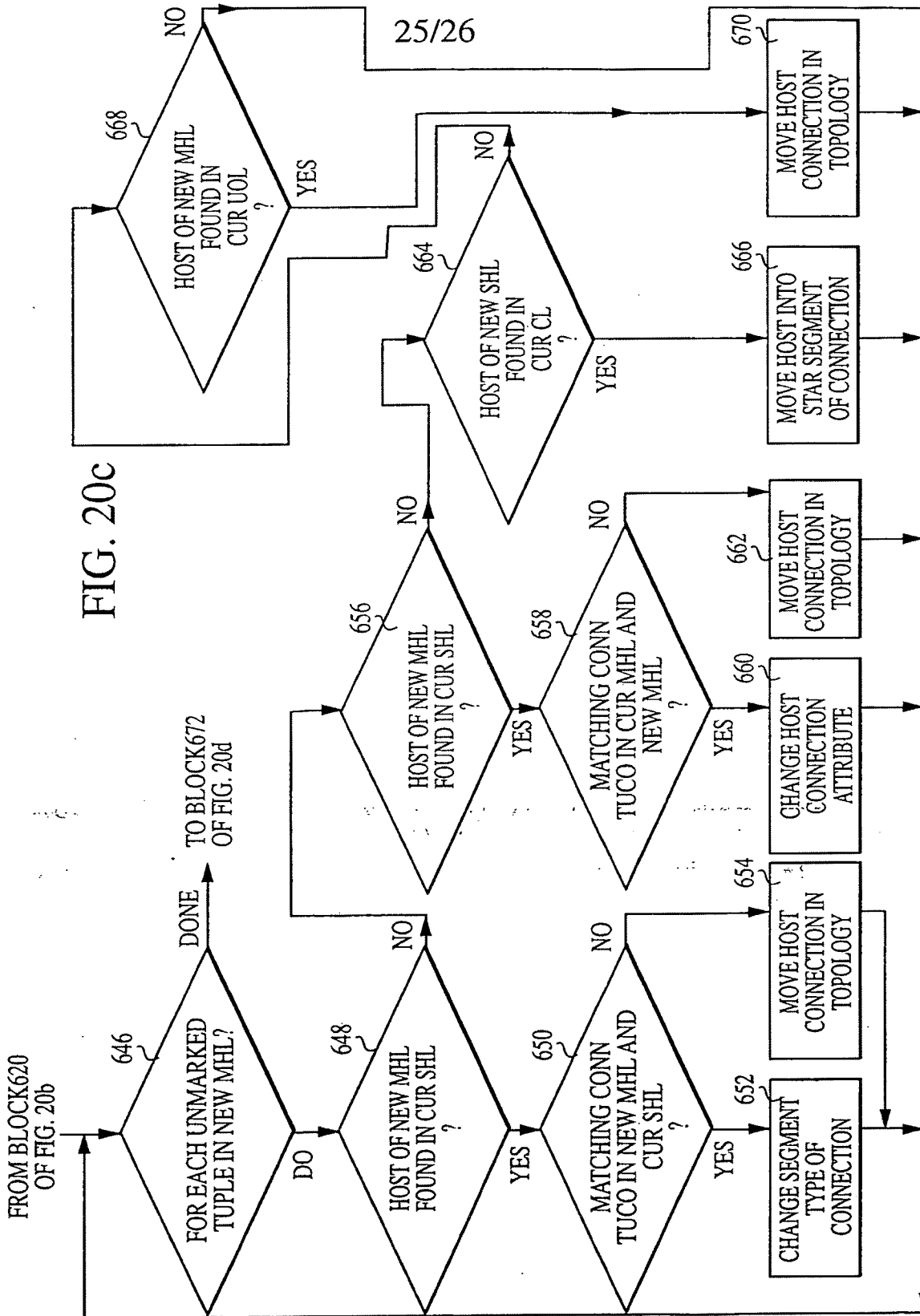
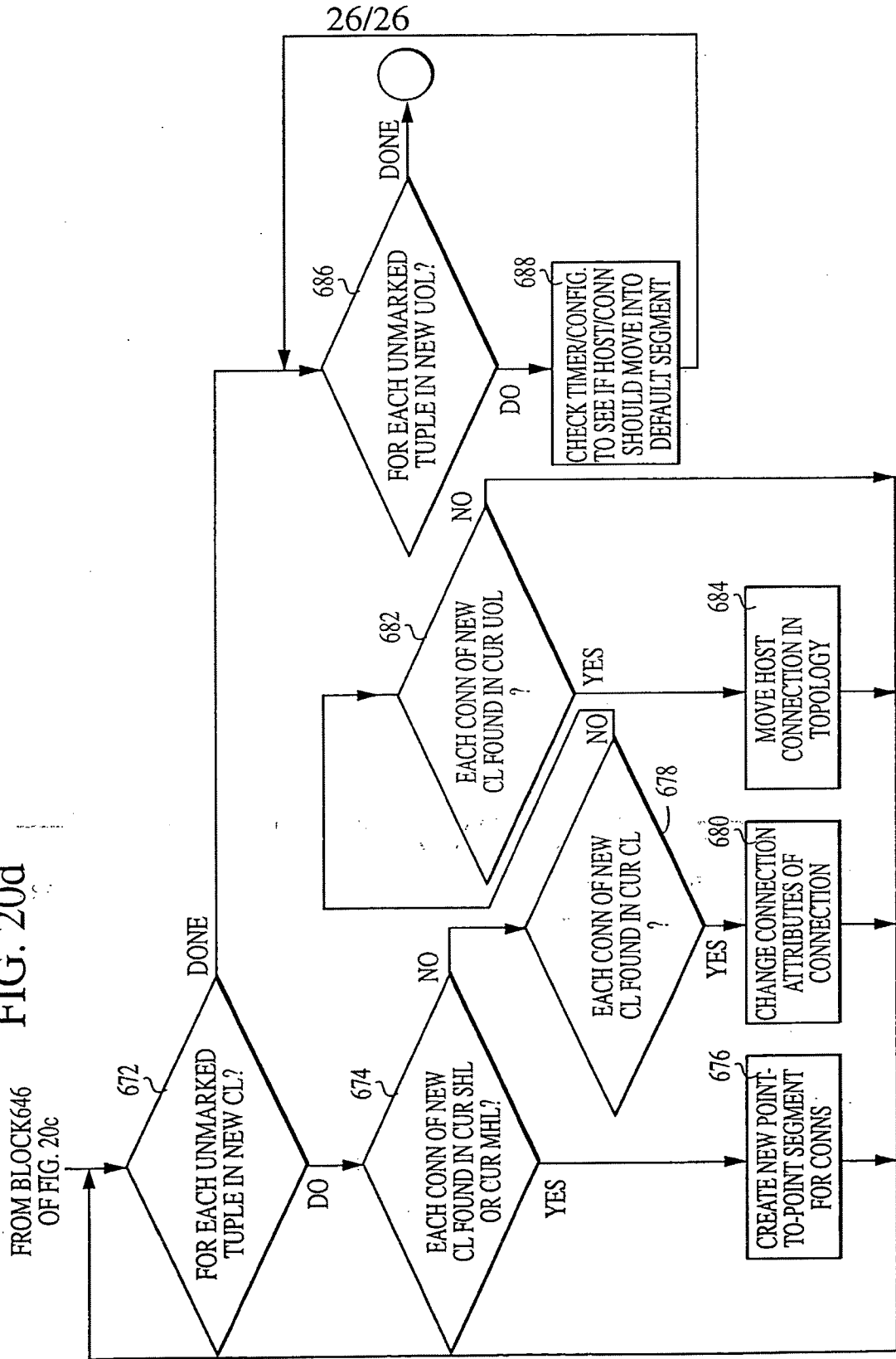


FIG. 20c

FIG. 20d



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1 Still another difficulty with existing systems is that they focus on the minutia without
2 considering the larger mapping considerations. Whenever an individual change in the system is
3 detected, existing methods immediately act on that change, rather than taking a broader view of the
4 change in the context of other system changes. For example, a device may be removed from the
5 network temporarily and replaced with its ports reversed. In existing systems, this swapped port
6 scenario could require hundreds or thousands of changes because the reference addresses will have
7 changed for all interconnected devices.

8 Still another disadvantage of existing methods is that they use a continuous polling paradigm.
9 These methods continuously poll network addresses throughout the day and make decisions based
10 on those continuous polling results. This creates traffic on the network that slows other processes.

11 Still another limitation of existing methods is the assumption that network parts of a
12 particular layer would be physically separated from other parts. Network layer 1 may represent the
13 physical cabling of the network, layer 2 may represent the device connectivity, and layer 3 may
14 represent a higher level of abstraction, such as the groupings of devices into regions. Existing
15 methods assume that all layer 3 region groupings are self-contained, running on the same unique
16 physical networking. However, in an internet protocol (IP) network, multiple IP domains may co-
17 exist on the same lower layer networking infrastructure. It has become common for a network to
18 employ a virtual local area network (LAN) to improve security or to simplify network maintenance,
19 for example. Using virtual LANs, a system may have any number of different IP domains sharing
20 the same physical connectivity. As a result, existing methods create confusion with respect to
21 topological mapping because networks with multiple IP addresses in different subnets for the
22 infrastructure devices cannot be properly represented because they assume the physical separation
23 of connectivity for separate IP domains. Still another limitation of existing methods is that they do
24 not allow topological loops, such as port aggregation or trunking, and switch meshing.

25

26

Summary of Invention

27 A method and system are disclosed for mapping the topology of a network having
28 interconnected nodes by identifying changes in the network and updating a stored network topology

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1 based on the changes. The nodal connections are represented by data tuples that store information
2 such as a host identifier, a connector interface, and a port specification for each connection. A
3 topology database stores an existing topology of a network. A topology converter accesses the
4 topology database and converts the existing topology into a list of current tuples. A connection
5 calculator calculates tuples to represent connections in the new topology. The topology converter
6 receives the new tuples, identifies changes to the topology, and updates the topology database using
7 the new tuples. The topology converter identifies duplicate tuples that appear in both the new tuples
8 and the existing tuples and marks the duplicate tuples to reflect that no change has occurred to these
9 connections. The topology converter attempts to resolve swapped port conditions and searches for
10 new singly-heard and multi-heard host link tuples in the list of existing tuples. The topology
11 converter also searches for new conflict link tuples in the existing tuples. The topology converter
12 updates the topology database with the new topology.

13 Summary of Drawings

14 Figure 1 is a drawing of a typical topological bus segment for representing the connectivity
15 of nodes on a network.

16 Figure 2 is a drawing of a typical topological serial segment for representing the connectivity
17 of nodes on a network.

18 Figure 3 is a drawing of a typical topological star segment for representing the connectivity
19 of nodes on a network.

20 Figure 4 is a drawing of another typical topological star segment for representing the
21 connectivity of nodes on a network.

22 Figure 5 is a drawing of the connectivity of an example network system.

23 Figure 6 is a drawing of the connectivity of another example network system.

24 Figure 7 is a block diagram of the system.

25 Figure 8 is a flow chart of the method of the system.

26 Figure 9 is a flow chart of the method used by the tuple manager.

27 Figure 10 is a flow chart of the method used by the connection calculator.

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1 100 comprises the first and second hosts 121, 122 connected to the first port 131 of the first
2 connector 140.

3 Figure 2 is a drawing of a typical topological serial segment 200 for representing the
4 connectivity of nodes on the network 110. In Figure 2, the first host 121 comprises a second port
5 132 on a second connector 145 which is connected via the network 110 to the first port 131 on the
6 first connector 140. The serial segment 200 comprises the second port 132 on the second
7 connector 145 connected to the first port 131 on the first connector 140. Figure 2 is an example of
8 a connector-to-connector ("conn-to-conn") relationship.

9 Figure 3 is a drawing of a typical topological star segment 301 for representing the
10 connectivity of nodes on the network 110. In Figure 3, the first host 121 is connected to the first
11 port 131 of the first connector 140. The star segment 301 comprises the first host 121 connected
12 to the first port 131 of the first connector 140. Figure 3 is an example of a connector-to-host
13 ("conn-to-host") relationship.

14 Figure 4 is a drawing of another typical topological star segment 301 for representing the
15 connectivity of nodes on the network 110. In addition to the connections described with respect to
16 Figure 3, a third host 123 is connected to a third port 133 of the first connector 140 and a fourth
17 host 124 is connected to a fourth port 134 of the first connector 140. In Figure 4, the star segment
18 301 comprises the first host 121 connected to the first port 131 of the first connector 140, the third
19 host 123 connected to the third port 133 of the first connector 140, and the fourth host 124
20 connected to the fourth port 134 of the first connector 140. Thus, the star segment 301 comprises,
21 on a given connector, at least one port, wherein one and only one host is connected to that port,
22 and that host. In the more general case, the star segment 301 comprises, on a given connector, all
23 ports having one and only one host connected to each port, and those connected hosts. Since the
24 segments, or links, drawn using the topological methods of Figure 4 resemble a star, they are
25 referred to as star segments.

26 For illustrative purposes, nodes in the figures described above and in subsequent figures are
27 shown as individual electronic devices or ports on connectors. Also, in the figures the nodes are

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1 represented as terminals. However, they could also be workstations, personal computers, printers,
2 scanners, or any other electronic device that can be connected to networks 110.

3 Figure 5 is a drawing of the connectivity of an example network system. In Figure 5, first,
4 third, and fourth hosts 121, 123, 124 are connected via the network 110 to first, third, and fourth
5 ports 131, 133, 134 respectively, wherein the first, third, and fourth ports 131, 133, 134 are
6 located on the first connector 140.

7 The first, third and fourth hosts 121, 123, 124 are singly-heard hosts connected to separate
8 ports 131, 133, 134 of a common connector 140 – the first connector 140. The fifth and sixth
9 hosts 125, 126 are singly-heard hosts connected to the third and fourth connectors 142, 143. The
10 seventh and eighth hosts 127, 128 are multi-heard hosts connected to the same port 139 of the fifth
11 connector 144. The multi-heard hosts 127, 128 illustrate a shared media segment 180, also
12 referred to as a bus 180.

13 The second, third, fourth, and fifth connectors 141, 142, 143, 144 are interconnected and
14 illustrate a switch mesh 181. Each of the connectors in the switch mesh 181 is connected to each
15 other, either directly or indirectly, to create a fully meshed connection. In the mesh, traffic may be
16 dynamically routed to create an efficient flow.

17 Figure 5 also shows an example of a port aggregation 182, also referred to as trunking 182.
18 The first connector 140 is connected via the network 110 to the second connector 141 by two
19 direct links, each of which is connected to different ports on the connectors. One link is connected
20 to the sixth port 136 of the first connector 140 and to the seventh port of the second connector
21 137. The other link is connected to fifth port 135 of the first connector 140 and to the eighth port
22 138 of the second connector 141. In this example, two connectors illustrate the multiple
23 connectivity between nodes. Depending upon the device specifications, devices such as connectors
24 may be connected via any number of connectors. As explained herein, the system resolves multiple
25 connectivity problems by tracking port information for each connection.

26 Figure 6 is a drawing of the connectivity of a portion of a network having three connectors
27 171, 172, 173. A first host 151 is connected directly to the first port 161 of the first connector 171
28 and the second host 152 is connected to a sixth port 166 of the third connector 173. The second

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1 port 162 of the first connector 171 is connected directly to the third port 163 of the second, or
2 intermediate, connector 172. The fourth port 164 of the intermediate connector 172 is connected
3 directly to the fifth port 165 of the third connector 173.

4 Figure 7 shows a block diagram of the system. Figure 8 shows a flow chart of the method
5 used by the system to retrieve and update the topology of the network. A tuple manager 300, also
6 referred to as a data miner 300, gathers 902 data from network nodes and builds 904 tuples to
7 update the current topology. The topology database "topodb" 350 stores the current topology for
8 use by the system. The "neighbor data" database 310 stores new tuple data retrieved by the tuple
9 manager 300. The connection calculator 320 processes the data in the neighbor data database 310
10 to determine the new network topology. The connection calculator 320 reduces 906 the tuple data
11 and sends it to the reduced topology relationships database 330. The topology converter 340 then
12 updates 908 the topology database 350 based on the new tuples sent to the reduced topology
13 relationships database 330 by the connection calculator 320.

14 Figure 9 shows a flow chart of one operation of the tuple manager 300, as described
15 generally by the data gathering 902 and tuple building 904 steps of the method shown in Figure 8.
16 The tuple manager 300 receives 910 a signal to gather tuple data. The tuple manager 300 then
17 retrieves 912 node information of the current topology stored in the topology database 350. This
18 information tells the tuple manager 300 which devices or nodes are believed to exist in the system
19 based on the nodes that were detected during a previous query. The tuple manager 300 then
20 queries 914 the known nodes to gather the desired information. For example, the connectors may
21 maintain forwarding tables that store connectivity data used to perform the connectors' ordinary
22 functions, such as switching. Other devices may allow the system to perform queries to gather
23 information about the flow of network traffic. This data identifies the devices heard by a connector
24 and the port on which the device was heard. The tuple manager 300 gathers this data by accessing
25 forwarding tables and other information sources for the nodes to determine such information as their
26 physical address, interface information, and the port from which they "hear" other devices. Based
27 on this information, the tuple manager 300 builds 916 tuples and stores 918 them in the "neighbor
28 data" database 310. Some nodes may have incomplete information. In this case, the partial

1 404 whether the tuple is a connector-to-host (conn-to-host) link tuple. If it is not a conn-to-host
 2 link, the connection calculator 320 concludes 418 that it is a conn-to-conn link and processes 402
 3 the next tuple. If the tuple is a conn-to-host link tuple, then the connection calculator 320
 4 determines 406 whether the connector hears only this particular host on the port identified in the
 5 tuple. If the connector hears other hosts on this port, then the tuple is classified 416 as a multi-
 6 heard host link (mhhl) tuple.

7 If the connector hears only the one host on the port – that is, if the host is a singly-heard
 8 host – then the connection calculator 320 determines 408 whether the host is heard singly by any
 9 other connectors. If no other connectors hear the host as a singly-heard host, then the tuple is
 10 classified as a singly-heard host link (shhl) tuple 412 and other tuples for this host are classified 414
 11 as extra host links (ehl). Another tuple for this host may be, for example, an intermediate connector
 12 connected indirectly to a host. For example, Figure 6 shows three connectors 171, 172, 173 the
 13 first connector is connected directly to the first host 151. This connection therefore forms an shhl
 14 tuple. The intermediate connector 172 is indirectly connected to the first host 151. The tuple data
 15 indicates that the intermediate connector 172 is indirectly connected to the host and hears the host
 16 from a particular port. An extra host links tuple is created so that this data may be used later in
 17 conjunction with other extra host links tuples from devices across the network, to verify connectivity
 18 by providing hints about connections.

19 The first weeding process also attempts to identify conflicts. If other connectors hear the
 20 host as a singly-heard host, then a conflict arises and the tuple is classified 410 as a singly-heard
 21 conflict link (shcl) tuple to be resolved later. This conflict may arise, for example, if a host has been
 22 moved within the network, in which case the forwarding table data may no longer be valid. Certain
 23 connectors previously connected directly to the host may still indicate that the moved host is
 24 connected. When all tuples have been processed 402 to identify singly-heard host links, the first
 25 weeding phase 922 is complete.

26 Figures 12a-d show a flow chart of the infrastructure building phase 924 of the connection
 27 calculator 320. The purpose of the infrastructure building phase 924 is to determine how the
 28 connectors are set up in the network. The first part of the infrastructure building phase 924

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1 then the connection calculator 320 considers 436 every other conn-to-conn tuple containing conn2.
2 The other connector on this tuple may be referred to as conn3. If conn2 hears conn3 on a unique
3 port 438 and if conn1 also hears conn3 440, then the connection calculator 320 creates 442 a tuple
4 for conn1-to-conn2 in the connector-to-connector links tuple list.

5 After processing all of the conn1 tuples, the connection calculator 320 processes 444 each
6 conn1-to-conn2 links tuple to ensure that they have complete port data. For each incomplete tuple
7 446, the connection calculator 320 looks 448 for a different tuple involving conn1 in the extra host
8 links tuples on a different port. If a different tuple is found 450, then the connection calculator 320
9 determines 452 whether conn2 also hears the host. If conn2 does hear the host, then the
10 connection calculator 320 completes the missing port data for conn2. If conn2 does not also hear
11 the host 452, then the connection calculator 320 continues looking 448 through different tuples
12 involving conn1 in extra host links on different ports.

13 After attempting to complete the missing data in each of the conn-to-conn links tuples, the
14 connection calculator 320 processes 456 each conn-to-conn links tuple. The purpose of this sub-
15 phase is to attempt to disprove invalid conn-to-conn links. The connection calculator 320 considers
16 458 conn1 and conn2 of each conn-to-conn links tuple. Every other connector in conn-to-conn
17 links may be referred to as testconn. For each testconn 460, the connection calculator 320
18 determines 462 whether the testconn hears conn1 and conn2 on different groups/ports. If testconn
19 hears conn1 and conn2 on different ports, then the tuple is moved to extraconnlinks (ecl) 464.
20 Otherwise, the connection calculator 320 continues processing 460 the remaining testconns.

21 Figure 13 shows a flow chart of the second weeding phase 926. The purpose of the
22 second weeding phase 926 is to attempt to resolve conflicts involving singly-heard hosts identified in
23 the first weeding phase 922. In the situation described herein in which more than one connector
24 reports that a host is singly-heard, the second weeding phase 926 reviews the tuples created during
25 the infrastructure-building phase 924 involving the connector and host in question and attempts to
26 disprove the reported conflict. The connection calculator 320 processes 466 each
27 singleConflictLinks (scl) tuple (sometimes referred to as the search tuple) and considers 468 conn1
28 and host1 of the tuple. For each extra host links tuple containing host1 470, the connection

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1 calculator 320 considers 472 conn2 of the tuple. If there is a tuple in conn-to-conn links for conn2
2 and conn1 474, and if there is a conn2-to-conn1 tuple in the extra host links tuples 476, and if the
3 port is the same for conn2 hearing conn1 and host1 478, then the search tuple is moved 480 into
4 the singly heard host links and other tuples containing host1 are removed 482 from the
5 singleConflictLinks.

6 Figure 14 shows a flow chart of the noise reduction phase 928. The purpose of the noise
7 reduction phase 928 is to handle those connections in which a connector is not directly connected
8 to a host or to another connector. For example, networking technology may employ shared media
9 connections between connectors, rather than dedicated media connectors. With a shared media
10 connection, the entries in the forwarding tables for connectors attached to the shared media
11 connection will include every node accessing the shared media connection and may not present a
12 useful or accurate representation of the nodal connection. For example, if the network configuration
13 in Figure 6 used a shared media connection between the first connector 171 and the intermediate
14 connector 172, then the first connector is not really connected directly to the intermediate connector
15 because other devices (not shown in Figure 6) may also use the shared media connection. These
16 other devices may include web servers, other connectors, other subnetworks, etc. Tuples will be
17 created for the connectors 171, 172 on opposing ends of the shared media. In this situation, it is
18 inefficient to maintain point-to-point binary tuples for every connection. The noise reduction phase
19 928 disproves invalid tuples created by the shared media connections.

20 For each multi-heard host links (mhhl) tuple, also referred to as multiHeardLinks (mhl)
21 tuples (sometimes referred to as the search tuple) 484, conn1 and host1 are considered 486. For
22 each extra host links tuple containing host1 488, conn2 is considered 490. If there is a tuple in
23 conn-to-conn links for conn2 and conn1 492, and if there is a conn2-to-host1 tuple in
24 extraHostLinks 494, and if the group/port for conn2 hearing conn1 and host1 is different 496, then
25 the search tuple is moved 498 to extraHostLinks.

26 Figure 15 shows a flow chart for the "look for" phase 930. The purpose of this phase is to
27 complete missing data for mhhl tuples. There may exist connections on the network that have
28 incomplete tuple data. For example, the network may simply have no traffic between certain nodes,

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1 in which case data might not be stored in forwarding tables. In another example, a forwarding table
2 may not have sufficient room to store all of the required information and might delete data on a
3 FIFO basis. In the look for phase 930, the connection calculator 320 instructs the tuple manager
4 300 to query specific nodes to retrieve the missing data. Data that was not stored in a forwarding
5 table on the first interrogation may be present on a subsequent query. For each mhl tuple 500, the
6 connection calculator 320 considers 502 conn1 and host1. If the conn1 group/port is already in an
7 "alreadyDidLookfors" list, then a list is created 508 for all connectors in conn-to-conn links that are
8 heard by conn1 on the same group/port as host1. For each connector (conn2) in the list 510, the
9 connection calculator 320 determines 512 whether there is a conn2-to-host1 tuple in the mhl
10 tuples. If there is not such a tuple, then the connection calculator 320 initiates a look-for for conn2-
11 to-host1 via the tuple manager 300. When each connector in the list has been processed 510, the
12 conn1 group/port tuco is added 516 to an alreadyDidLookfors list. As an additional portion of the
13 look for phase 930 (not shown in figures) the system may ask a user to verify or clarify information
14 about connectivity. For example, the system may show the user the perceived connectivity or the
15 unresolved connectivity issues and request the user to add information as appropriate.

16 The connection calculator 330 process described above collects the tuple information from
17 the tuple manager 300, builds tuples new tuples and removes redundant or unnecessary tuples to
18 produce the new topology. This topology may have incomplete tuples possibly resulting from
19 extraneous information that the connection calculator 330 could not disprove. To refine the new
20 topology, the connection calculator 330 can request the tuple manager 300 to obtain additional
21 information about particular nodes or it may also request a user to refine the topology by adding or
22 removing tuples. Using the process of the connection calculator 330, tuples marked as non-
23 essential may be removed from the new topology to save space and to simplify the topology. The
24 connection calculator 330 is not confused by multiple connectivity situations such as port
25 aggregation 182 or switch meshing 181 as shown in Figure 5, because the tuples represent point-
26 to-point, or neighbor-to-neighbor, connectivity showing each connection in the network. This
27 point-to-point connectivity concept also helps enable the system to avoid difficulties that occur in

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1 systems that track higher levels of abstraction, such as layer 3 connectivity. Also, the tuples may
2 contain only selected information to minimize the storage space required for the topology.

3 Figures 16a-b show a flow chart of the consolidation phase 932. The purpose of this phase
4 is to consolidate the tuples that involve shared media connections. After the noise reduction phase
5 928, a considerable number of tuples involving shared media may remain. Rather than maintain a
6 binary tuple for each of the connections, an n-ary tuple is created for the link using a tuco for each
7 connector and each host connected thereto. For each mhh1 tuple 518, conn1 and host1 are
8 considered 520. If there are more conn1 group/port tuples in multiHeardLinks, and if are not any
9 n-ary multiHeardSegments (mhs) tuples 524, then an mhs tuple is created 526. If host1 is not
10 already in this particular mhs tuple 528, then conn2 of the tuple is considered 534. If there is a
11 conn1-to-conn2 conn-to-connLinks tuple on the same port as conn1-to-host1 536, then all
12 multiHeardLinks tuples for conn2-to-host1 with the same conn2 group/port as the conn1-to-conn2
13 are added 538 to the current mhs tuple.

14 After processing each mhh1 tuple 518, each singly-heard host links (shh1) tuple, also referred
15 to as a singlyHeardLinks (shl) tuple, is considered 540. For each shh1 tuple, the connector and host
16 are considered 542. If there is no existing singlyHeardSegments (shs) tuple for the connector 544,
17 then an shs tuple is created 546. The host tuco is then added to the shs 548.

18 Figure 17 shows a flow chart of the method used by the topology converter 340, as
19 described generally by the topology.update step 908 of the method shown in Figure 8. The
20 topology converter 340 converts 934 the topology into tuple lists, also referred to as the "morph
21 topo" phase 934. It then compares 936 the list from the topology currently stored in the topology
22 database 350 with the new list generated by the connection calculator 320 and discards 936
23 identical tuples in what is also referred to as the "discard duplicates" phase 936. It then takes
24 action 938 on the changes in the topology as determined by the changes in the tuple lists, in what is
25 also referred to as the "identify different tuples" phase 938.

26 Figure 18a shows a flow chart for the "morph topo" phase 934. For each node in the
27 topology 550, the topology converter 340 determines 552 whether the node is a connector. If the
28 node is a connector, then for each connected interface (conniface) of the connector (conn1) 554,

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1 the topology converter 340 determines 556 whether the conniface is connected to a star segment.
2 If it is connected to a star segment, then for every other interface in the segment 558, the topology
3 converter 340 determines 560 whether there is an existing shs tuple, referred to as the "topo tuple"
4 for the segment. If there is no such tuple, then the topology converter 340 creates 562 a topo shs
5 tuple. The tuco for the interface's host-to-topo shs is then added 564 to the topo shs tuple.

6 If the connector node is not connected to a star segment 556 and is connected to a bus
7 segment 566, the topology converter 340 determines 568 whether there is an existing mhs tuple for
8 conn1. If there is not an existing mhs tuple for conn1, then a topo mhs tuple is created 570. A tuco
9 is added 572 for the host to the mhs tuple.

10 If the connector node is not connected to either a star segment 556 or to a bus segment
11 566, then the topology converter knows that it is connected to another connector (conn2). If such
12 a connector does not already have an existing connLinks tuple for conn1 and conn2 576, then a
13 connLinks tuple is created 578. After processing the bus segment, star segment, and conn-to-conn
14 segment, for each conniface 554, the topology converter 340 proceeds to the next node 550.

15 Figure 18b shows a continuation of the flow chart of Figure 18a showing the steps of the
16 method when the topology converter 340 determines that the node is not a connector 552. If the
17 node is in the default segment, then an "unheardOfLinks" tuple is created 582 and the topology
18 converter proceeds to the next node 550. If the node is not in the default segment 580, then the
19 topology converter 340 determines whether the node is in a star segment 584. If the node is in a
20 star segment, then if there is not already an shs tuple, the topology converter 340 creates 588 an shs
21 tuple. The tuco for the node is then added 590 to the shs tuple, and the topology converter 340
22 proceeds to the next node 550.

23 If the node is not in a star segment, then the topology converter 340 knows that it is in the
24 bus segment. If there is not already an mhs tuple for the node, 594, then the topology converter
25 340 creates 596 an mhs tuple. The tuco for the node is then added 598 to the mhs tuple, and the
26 topology converter proceeds to the next node 550.

27 Figure 19 shows a flow chart for the discard duplicates phase 936 of the topology
28 converter 340. For each tuple in the new tuples (nt) 600, the topology converter looks for 602 an

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1 curSHL 648, then the system determines 650 whether there is a matching connector tuco between
2 the newMHL and the curMHL. If there is a matching tuco, then the segment type of connection is
3 changed 652. If there is not a matching tuco, then the host connection is moved 654 in the
4 topology.

5 If the host is found in the curMHL tuples 656, then the system determines 658 whether
6 there is a matching connector tuco in both the curMHL tuples and the newMHL tuples. If there is a
7 matching connector tuco, then the host connection attribute is changed 660. If there is not a
8 matching tuco, then the host connection is moved 662 in the topology. If the host is found in the
9 curCL tuples 664, then the host is moved into a bus segment of a connector. If the host is found in
10 the curUOL tuples 668, then the host connection is moved 670 in the topology.

11 Figure 20d shows another portion of the identify different tuples phase 938. For each
12 unmarked tuple in the newCL tuples 672, there are three possibilities for the connector. The
13 connector of the unmarked tuple in newCL can be found in the curSHL or curMHL 674, in the
14 curCL 678, or in the curUOL 682. If each connector is found in the curSHL or curMHL list 674,
15 then the system creates 676 a new point-to-point segment for the connectors. If the connectors are
16 found in the curCL 678, then the connection attributes of the connectors are changed 680. If each
17 connector is found in the curUOL tuples 682, then the host connection is moved 684 in the
18 topology.

19 Another part of the identify different tuples phase 938 is shown in blocks 686 and 688 of
20 Figure 20d. For each unmarked tuple in the newUOL tuples 686, the system checks 688 the
21 timer/configuration to determine whether the host/conn should move into the default segment from
22 its current segment.

23 An advantage of the system is that it may be schedulable. The system may map network
24 topology continuously, as done by existing systems, or it may be scheduled to run only at certain
25 intervals, as desired by the user. A further advantage of the system is that it is capable of
26 processing multiple connections between the same devices and of processing connection meshes,
27 because it tracks each nodal connection independently, without limitations on the types of
28 connections that are permitted to exist.

1 Although the present invention has been described with respect to particular embodiments
2 thereof, variations are possible. The present invention may be embodied in specific forms without
3 departing from the essential spirit or attributes thereof. It is desired that the embodiments described
4 herein be considered in all respects illustrative and not restrictive and that reference be made to the
5 appended claims for determining the scope of the invention.

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Claims

- 1
- 2 1. In a network having interconnected nodes with data tuples that represent nodal
- 3 connections, a method for mapping a network topology by identifying changes between an existing
- 4 topology and a new topology, the method comprising:
- 5 converting an existing topology into a list of existing tuples that represent existing nodal
- 6 connections;
- 7 receiving new tuples that represent new nodal connections; and
- 8 comparing the list of existing tuples with the new tuples to identify changes to the topology.
- 9 2. The method of claim 1, further comprising updating a topology database with a new
- 10 topology.
- 11 3. The method of claim 1, further comprising taking action on the changes to the
- 12 topology.
- 13 4. The method of claim 1, wherein the tuples include information about a host
- 14 identifier, a connector interface, and a port specification.
- 15 5. The method of claim 1, wherein the step of comparing comprises identifying
- 16 duplicate tuples that appear both in the list of existing tuples and in the new tuples, and maintaining a
- 17 current status of the topology for these tuples.
- 18 6. The method of claim 1, wherein the step of comparing comprises identifying a
- 19 swapped port condition on a connector.
- 20 7. The method of claim 1, wherein the step of comparing comprises searching for a
- 21 host of a new singly-heard host link tuple or a new multi-heard host link tuple in the list of existing
- 22 tuples.
- 23 8. A system for mapping a network topology by identifying changes between an
- 24 existing topology and a new topology, based on changes to data tuples that represent nodal
- 25 connections comprising:
- 26 a topology database that stores an existing topology of a network; and

1 17. The method of claim 15, wherein the step of comparing comprises identifying
2 duplicate tuples that appear both in the list of existing tuples and in the new tuples, and maintaining a
3 current status of the topology for these tuples.

4 18. The method of claim 15, wherein the step of comparing comprises identifying a
5 swapped port condition on a connector.

6 19. The method of claim 15, wherein the step of comparing comprises searching for a
7 host of a new singly-heard host link tuple or a new multi-heard host link tuple in the list of existing
8 tuples.

9 20. The method of claim 15, wherein the step of comparing comprises searching for a
10 connector of a new conflict links tuple in the list of existing tuples.

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1 **Abstract**

2 A method and system are disclosed for mapping the topology of a network having
3 interconnected nodes by identifying changes in the network and updating a stored network topology
4 based on the changes. The nodal connections are represented by data tuples that store information
5 such as a host identifier, a connector interface, and a port specification for each connection. A
6 topology database stores an existing topology of a network. A topology converter accesses the
7 topology database and converts the existing topology into a list of current tuples. A connection
8 calculator calculates tuples to represent connections in the new topology. The topology converter
9 receives the new tuples, identifies changes to the topology, and updates the topology database using
10 the new tuples. The topology converter identifies duplicate tuples that appear in both the new tuples
11 and the existing tuples and marks the duplicate tuples to reflect that no change has occurred to these
12 connections. The topology converter attempts to resolve swapped port conditions and searches for
13 new singly-heard and multi-heard host link tuples in the list of existing tuples. The topology
14 converter also searches for new conflict link tuples in the existing tuples. The topology converter
15 updates the topology database with the new topology.

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DECLARATION AND POWER OF ATTORNEY FOR PATENT APPLICATION

ATTORNEY REFERENCE NO. 10008102-1

As a below named inventor, I hereby declare that:

My residence/post office address and citizenship are as stated below next to my name;

I believe I am the original, first and sole inventor (if only one name is listed below) or an original, first and joint inventor (if plural names are listed below) of the subject matter which is claimed and for which a patent is sought on the invention entitled:

Method And System For Identifying And Processing Changes To A Network Topology

the specification of which is attached hereto unless the following box is checked:

() was filed on _____ as US Application Serial No. or PCT International Application Number _____ and was amended on _____ (if applicable).

I hereby state that I have reviewed and understood the contents of the above-identified specification, including the claims, as amended by any amendment(s) referred to above. I acknowledge the duty to disclose all information which is material to patentability as defined in 37 CFR 1.56.

Foreign Application(s) and/or Claim of Foreign Priority

I hereby claim foreign priority benefits under Title 35, United States Code Section 119 of any foreign application(s) for patent or inventor(s) certificate listed below and have also identified below any foreign application for patent or inventor(s) certificate having a filing date before that of the application on which priority is claimed:

Table with 4 columns: COUNTRY, APPLICATION NUMBER, DATE FILED, PRIORITY CLAIMED UNDER 35 U.S.C. 119. Row 1: N/A, YES: NO: Row 2: YES: NO:

Provisional Application

I hereby claim the benefit under Title 35, United States Code Section 119(e) of any United States provisional application(s) listed below:

Table with 2 columns: APPLICATION SERIAL NUMBER, FILING DATE. Row 1: N/A

U. S. Priority Claim

I hereby claim the benefit under Title 35, United States Code, Section 120 of any United States application(s) listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of Title 35, United States Code Section 112, I acknowledge the duty to disclose material information as defined in Title 37, Code of Federal Regulations, Section 1.56(a) which occurred between the filing date of the prior application and the national or PCT international filing date of this application:

Table with 3 columns: APPLICATION SERIAL NUMBER, FILING DATE, STATUS (patented/pending/abandoned). Row 1: N/A

POWER OF ATTORNEY:

As a named inventor, I hereby appoint the following attorney(s) and/or agent(s) to prosecute this application and transact all business in the Patent and Trademark Office connected therewith:

Customer Number 022879

Place Customer Number Bar Code Label here

Send Correspondence to: HEWLETT-PACKARD COMPANY Intellectual Property Administration P.O. Box 272400 Fort Collins, Colorado 80527-2400

Direct Telephone Calls To: T. Grant Ritz (970) 898-0697

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine or imprisonment, or both, under Section 1001 of Title 18 of the United States Code and that such willful false statements may jeopardize the validity of the application or any patent issued thereon.

Full Name of Inventor: Eric A Pulsipher Citizenship: US

Residence: 2937 Redburn Drive Ft Collins CO 80525

Post Office Address: Same as residence

Inventor's Signature [Handwritten Signature]

Date 10/31/2000

001001-21650260

DECLARATION AND POWER OF ATTORNEY
FOR PATENT APPLICANT (continued)

ATTORNEY DOCKET NO. 10008102-1

Full Name of # 2 joint inventor: Joseph R Hunt ^{5234 Hahns Peak Dr #205} Citizenship: US

Residence: 5841 Meadow Creek Ln Loveland, CO 80538

Post Office Address: Same as Residence

Inventor's Signature: *Joseph R Hunt* Date: 10/31/00

Full Name of # 3 joint inventor: _____ Citizenship: _____

Residence: _____

Post Office Address: _____

Inventor's Signature _____ Date _____

Full Name of # 4 joint inventor: _____ Citizenship: _____

Residence: _____

Post Office Address: _____

Inventor's Signature _____ Date _____

Full Name of # 5 joint inventor: _____ Citizenship: _____

Residence: _____

Post Office Address: _____

Inventor's Signature _____ Date _____

Full Name of # 6 joint inventor: _____ Citizenship: _____

Residence: _____

Post Office Address: _____

Inventor's Signature _____ Date _____

Full Name of # 7 joint inventor: _____ Citizenship: _____

Residence: _____

Post Office Address: _____

Inventor's Signature _____ Date _____

Full Name of # 8 joint inventor: _____ Citizenship: _____

Residence: _____

Post Office Address: _____

Inventor's Signature _____ Date _____

00101-2450260



UNITED STATES PATENT AND TRADEMARK OFFICE

COMMISSIONER FOR PATENTS
 UNITED STATES PATENT AND TRADEMARK OFFICE
 WASHINGTON, D.C. 20231
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Bib Data Sheet

SERIAL NUMBER 09/703,942	FILING DATE 10/31/2000 RULE -	CLASS 370	GROUP ART UNIT 2661	ATTORNEY DOCKET NO. 10008102-1	
APPLICANTS Eric A. Pulsipher, Ft Collins, CO ; Joseph R. Hunt, Loveland, CO ;					
** CONTINUING DATA *****					
** FOREIGN APPLICATIONS *****					
IF REQUIRED, FOREIGN FILING LICENSE GRANTED ** 02/05/2001					
Foreign Priority claimed <input type="checkbox"/> yes <input checked="" type="checkbox"/> no		STATE OR COUNTRY CO	SHEETS DRAWING 26	TOTAL CLAIMS 20	INDEPENDENT CLAIMS 3
35 USC 119 (a-d) conditions met <input type="checkbox"/> yes <input checked="" type="checkbox"/> no <input type="checkbox"/> Met after Allowance					
Verified and Acknowledged		Examiner's Signature <i>wcs</i> Initials			
ADDRESS 022879					
TITLE Method and system for identifying and processing changes to a network topology					
FILING FEE RECEIVED 710	FEES: Authority has been given in Paper No. _____ to charge/credit DEPOSIT ACCOUNT No. _____ for following:		<input type="checkbox"/> All Fees <input type="checkbox"/> 1.16 Fees (Filing) <input type="checkbox"/> 1.17 Fees (Processing Ext. of time) <input type="checkbox"/> 1.18 Fees (Issue) <input type="checkbox"/> Other _____ <input type="checkbox"/> Credit		

None, wcs

PATENT APPLICATION SERIAL NO. _____

U.S. DEPARTMENT OF COMMERCE
PATENT AND TRADEMARK OFFICE
FEE RECORD SHEET

11/06/2000 KZEWIDIE 00000062 082025 09703942

01 FC:101 710.00 CH

PTO-1556

(5/87)

U.S. GPO: 1999-459-022/19144

PATENT APPLICATION FEE DETERMINATION RECORD
Effective October 1, 2000

Application or Docket Number

09/203942

CLAIMS AS FILED - PART I

	(Column 1)	(Column 2)
TOTAL CLAIMS	20	
FOR	NUMBER FILED	NUMBER EXTRA
TOTAL CHARGEABLE CLAIMS	20 minus 20 = *	0
INDEPENDENT CLAIMS	3 minus 3 = *	0
MULTIPLE DEPENDENT CLAIM PRESENT <input type="checkbox"/>		

SMALL ENTITY TYPE OR

OTHER THAN SMALL ENTITY

RATE	FEE
BASIC FEE	355.00
X\$ 9=	
X40=	
+135=	
TOTAL	

RATE	FEE
BASIC FEE	710.00
X\$18=	
X80=	
+270=	
TOTAL	

* If the difference in column 1 is less than zero, enter "0" in column 2

CLAIMS AS AMENDED - PART II

	(Column 1)	(Column 2)	(Column 3)
AMENDMENT A	CLAIMS REMAINING AFTER AMENDMENT		HIGHEST NUMBER PREVIOUSLY PAID FOR
	Total	* Minus	** =
	Independent	* Minus	*** =
FIRST PRESENTATION OF MULTIPLE DEPENDENT CLAIM <input type="checkbox"/>			

SMALL ENTITY OR

OTHER THAN SMALL ENTITY

RATE	ADDITIONAL FEE
X\$ 9=	
X40=	
+135=	
TOTAL ADDIT. FEE	

RATE	ADDITIONAL FEE
X\$18=	
X80=	
+270=	
TOTAL ADDIT. FEE	

	(Column 1)	(Column 2)	(Column 3)
AMENDMENT B	CLAIMS REMAINING AFTER AMENDMENT		HIGHEST NUMBER PREVIOUSLY PAID FOR
	Total	* Minus	** =
	Independent	* Minus	*** =
FIRST PRESENTATION OF MULTIPLE DEPENDENT CLAIM <input type="checkbox"/>			

RATE	ADDITIONAL FEE
X\$ 9=	
X40=	
+135=	
TOTAL ADDIT. FEE	

RATE	ADDITIONAL FEE
X\$18=	
X80=	
+270=	
TOTAL ADDIT. FEE	

	(Column 1)	(Column 2)	(Column 3)
AMENDMENT C	CLAIMS REMAINING AFTER AMENDMENT		HIGHEST NUMBER PREVIOUSLY PAID FOR
	Total	* Minus	** =
	Independent	* Minus	*** =
FIRST PRESENTATION OF MULTIPLE DEPENDENT CLAIM <input type="checkbox"/>			

RATE	ADDITIONAL FEE
X\$ 9=	
X40=	
+135=	
TOTAL ADDIT. FEE	

RATE	ADDITIONAL FEE
X\$18=	
X80=	
+270=	
TOTAL ADDIT. FEE	

* If the entry in column 1 is less than the entry in column 2, write "0" in column 3.
 ** If the "Highest Number Previously Paid For" IN THIS SPACE is less than 20, enter "20."
 *** If the "Highest Number Previously Paid For" IN THIS SPACE is less than 3, enter "3."
 The "Highest Number Previously Paid For" (Total or Independent) is the highest number found in the appropriate box in column 1.